Sample problems for the 9th lecture (May 10)

1. (a) Suppose that the continuous random variable X has probability density function (pdf) $f_X(x)$. Express the probability mass function of $q_n(X)$ in terms of $f_X(\cdot)$, where

$$q_n(x) := m\Delta$$
 for $m\Delta \le x < (m+1)\Delta$

is the uniform quantization on $x \in \mathbb{R}$.

- (b) Express the entropy of $q_n(X)$ in terms of $f_X(\cdot)$.
- (c) Let $0 \le f_X(x) \le A$ be a bounded function for $x \in \mathbb{R}$. Show that the entropy of $q_n(X)$ is lower bounded by

$$\log_2 \frac{1}{\Delta A}$$
.

Hint:
$$\int_{m\Delta}^{(m+1)\Delta} f_X(x) dx \le \Delta A.$$

(d) Prove that the entropy of X is infinity.

Hint: $H(f(X)) \leq H(X)$ for any function f.

Solution.

(a)
$$\Pr[q_n(X) = m\Delta] = \int_{m\Delta}^{(m+1)\Delta} f_X(x) dx$$

(b)

$$H(q_n(X)) = \sum_{m=-\infty}^{\infty} \Pr[q_n(X) = m\Delta] \log_2 \frac{1}{\Pr[q_n(X) = m\Delta]}$$
$$= \sum_{m=-\infty}^{\infty} \left(\int_{m\Delta}^{(m+1)\Delta} f_X(x) dx \right) \log_2 \frac{1}{\left(\int_{m\Delta}^{(m+1)\Delta} f_X(x) dx \right)}$$

(c) From (b),

$$H(q_n(X)) = \sum_{m=-\infty}^{\infty} \left(\int_{m\Delta}^{(m+1)\Delta} f_X(x) dx \right) \log_2 \frac{1}{\left(\int_{m\Delta}^{(m+1)\Delta} f_X(x) dx \right)}$$

$$\geq \sum_{m=-\infty}^{\infty} \left(\int_{m\Delta}^{(m+1)\Delta} f_X(x) dx \right) \log_2 \frac{1}{\Delta A}$$

$$= \log_2 \frac{1}{\Delta A} \sum_{m=-\infty}^{\infty} \left(\int_{m\Delta}^{(m+1)\Delta} f_X(x) dx \right)$$

$$= \log_2 \frac{1}{\Delta A} \int_{-\infty}^{\infty} f_X(x) dx$$

$$= \log_2 \frac{1}{\Delta A}.$$

(d) From (c), $H(X) \ge H(q_n(X)) \ge \log_2 \frac{1}{\Delta A}$; hence,

$$H(X) \ge \lim_{\Delta \downarrow 0} \log_2 \frac{1}{\Delta A} = \infty.$$

- 2. Determine the differential entropy (in nats) of random variable X for each of the following cases.
 - (a) X is exponential with parameter λ , i.e., $f_X(x) = \lambda e^{-\lambda x}, x \geq 0$.
 - (b) X is Laplacian with parameter λ and mean zero, i.e., $f_X(x) = \frac{1}{2\lambda}e^{-\frac{|x|}{\lambda}}, x \in \mathbb{R}$.
 - (c) $X = aX_1 + bX_2$, where a and b are non-zero constants and X_1 and X_2 are independent Gaussian random variables such that $X_1 \sim \mathcal{N}(\mu_1, \sigma_1^2)$ and $X_2 \sim \mathcal{N}(\mu_2, \sigma_2^2)$.

Hint: The differential entropy of Gaussian random variable with variance σ^2 is $h(X) = \frac{1}{2} \ln{(2\pi e \sigma^2)}$.

Solution:

(a) Exponential distribution: $f_X(x) = \lambda e^{-\lambda x}$ if $x \ge 0$.

$$h(X) = -\int_0^\infty f_X(x) \ln f_X(x) dx$$
$$= -\int_0^\infty \lambda e^{-\lambda x} (\ln \lambda - \lambda x) dx$$
$$= \ln \left(\frac{e}{\lambda}\right)$$
$$= 1 - \ln(\lambda) \text{ nats}$$

where we have used the fact that $E[X] = 1/\lambda$.

(b) Laplace distribution: $f_X(x) = (1/2\lambda)e^{-\frac{|x|}{\lambda}}, x \in \mathbb{R}.$

$$h(X) = -\int_{-\infty}^{\infty} f_X(x) \left(\ln \frac{1}{2\lambda} - \frac{|x|}{\lambda} \right) dx$$
$$= \ln(2\lambda) + \frac{1}{\lambda} \int_{-\infty}^{\infty} |x| f_X(x) dx$$
$$= \ln(2\lambda) + \frac{\lambda}{\lambda}$$
$$= 1 + \ln(2\lambda) \text{ nats}$$

where we have used the fact that $E[|X|] = \lambda$.

(c) X_1 is independent of X_2 and $X_i \sim \mathcal{N}(\mu_i, \sigma_i^2)$, i = 1, 2. Hence, $X = aX_1 + bX_2$ is also Gaussian: $X \sim \mathcal{N}(a\mu_1 + b\mu_2, a^2\sigma_1^2 + b^2\sigma_2^2)$. So,

$$h(aX_1 + bX_2) = \frac{1}{2}\ln(2\pi e(a^2\sigma_1^2 + b^2\sigma_2^2))$$
 nats.

3. (a) Let X be a random variable with pdf $f_X(x)$ and support \mathcal{X} . Prove that for any positive C_1 ,

$$C_1 \cdot E[q(X)] - h(X) \ge \ln(C_2)$$

where

$$C_2 = \frac{1}{\int_{\mathcal{V}} e^{-C_1 \cdot q(X)} dx}.$$

Hint: $D(X||Y) = \int_{\mathcal{X}} f_X(x) \ln \frac{f_X(x)}{f_Y(x)} dx \ge 0$ for any two continuous random variables X and Y.

(b) When does equality hold in (a) such that

$$h(X) = E[q(X)] - \ln(C) = C_1 \cdot E[q(X)] + \ln\left(\int_{\mathcal{X}} e^{-C_1 \cdot q(x)} dx\right)?$$

- (c) Use (a) and (b) to find the random variable that maximizes the differential entropy among all variables with $E[(X-1)^2] = \sigma^2$ and support \mathbb{R} .
- (d) Use (a) and (b) to find the random variable that maximizes the differential entropy among all variables with finite support [a, b].

Solution.

(a)

$$C_1 \cdot E[q(X)] - h(X) = C_1 \cdot \int_{-\infty}^{\infty} f_X(x)q(x)dx - \int_{-\infty}^{\infty} f_X(x)\ln\frac{1}{f_X(x)}dx$$

$$= \int_{-\infty}^{\infty} f_X(x)\ln\frac{f_X(x)}{e^{-C_1 \cdot q(x)}}dx$$

$$= \int_{-\infty}^{\infty} f_X(x)\ln\frac{f_X(x)}{C_2e^{-C_1 \cdot q(x)}}dx + \ln(C_2)$$

$$= \int_{-\infty}^{\infty} f_X(x)\ln\frac{f_X(x)}{f_Y(x)}dx + \ln(C_2)$$

$$\geq \ln(C_2)$$

where $f_Y(x) = C_2 e^{-C_1 \cdot q(x)}$ is a pdf defined over $x \in \mathcal{X}$.

- (b) From the derivation in (a), equality holds iff $f_X(x) = f_Y(y)$, i.e., $f_X(x) = C_2 e^{-C_1 \cdot q(x)}$ for $x \in \mathcal{X}$.
- (c) $f_X(x) = \frac{e^{-C_1(x-1)^2}}{\int_{-\infty}^{\infty} e^{-C_1(x-1)^2} dx}$, where C_1 must satisfy

$$E[(X-1)^{2}] = \int_{-\infty}^{\infty} (x-1)^{2} f_{X}(x) dx = \sigma^{2}.$$

This implies $C_1 = \frac{1}{2\sigma^2}$ and

$$h(X) = C_1 \cdot E[(X-1)^2] + \ln\left(\int_{-\infty}^{\infty} e^{-C_1(x-1)^2} dx\right)$$
$$= \frac{1}{2} + \frac{1}{2}\ln(2\pi\sigma^2) = \frac{1}{2}\ln(2\pi e\sigma^2).$$

(d) We can set q(x) = 1 (i.e., E[q(X)] = E[1] = 1) in (a). Hence, $f_X(x) = \frac{e^{-C_1}}{\int_a^b e^{-C_1} dx} = \frac{1}{b-a}$, where C_1 straightforwardly satisfies

$$E[1] = \int_{-\infty}^{\infty} 1 \cdot f_X(x) dx = 1.$$

As a result, C_1 can be arbitrary, and

$$h(X) = C_1 \cdot E[1] + \ln \left(\int_a^b e^{-C_1} dx \right)$$
$$= \ln(b - a).$$

- 4. Let the channel model be Y = AX + Z, where X and Y are respectively channel input and output, and A and Z are respectively channel gain and additive noise. Suppose A, X and Z are independent and they are all continuous random variables that admit pdf.
 - (a) Suppose the channel output can perfectly estimate the value of A. The channel capacity cost function is thus given by

$$C(P) = \sup_{F_X: E[t(X)] \le P} I(X; Y, A).$$

In other words, C(P) is the largest input-output mutual information attainable for any input distribution F_X that satisfies the cost constraint $E[t(X)] \leq P$. Show that C(P) can be alternatively represented as

$$C(P) = \left(\sup_{F_X: E[t(X)] \le P} h(Y|A)\right) - h(Z).$$

Hint: Translation does not change the differential entropy.

(b) Let $t(x) = x^2$, i.e., $t(\cdot)$ is a power constraint, and let Z be zero-mean Gaussian distributed with variance σ^2 . Also assume $E[A^2] = 1$. Is it possible to have

$$C(P) \stackrel{?}{=} \frac{1}{2} \ln \left(1 + \frac{P}{\sigma^2} \right).$$

Justify your answer.

Hint:
$$E[Y^2] = E[(AX + Z)^2] = E[A^2]E[X^2] + E[Z^2]$$

Solution.

(a) By chain rule for mutual information,

$$I(X;Y,A) = I(X;A) + I(X;Y|A)$$

$$= I(X;Y|A) \text{ (Because } I(X;A) = 0)$$

$$= h(Y|A) - h(Y|X,A).$$

Since h(Y|X, A) = h(AX + Z|X, A) = h(Z|X, A) as differential entropy is translation-invariant, and since h(Z|X, A) = h(Z) as Z is independent of X and A, we have the desired result.

(b) $E[X^2] \leq P$ holds iff $E[Y^2] \leq P + \sigma^2$, since

$$E[Y^2] = E[(AX + Z)^2] = E[A^2]E[X^2] + E[Z^2] = E[X^2] + \sigma^2.$$

Hence,

$$\sup_{F_X: E[X^2] \le P} h(Y|A) = \sup_{F_X: E[Y^2] \le P + \sigma^2} h(Y|A)$$

$$\le \frac{1}{2} \ln \left(2\pi e(P + \sigma^2) \right),$$

where the inequality follows from that the differential entropy is maximized when Y conditional on A is Gaussian distributed with variance $P+\sigma^2$. However, when A=a, we need X to be zero-mean Gaussian distributed with $E[X^2]=P/a$ to achieve the bound, but this requires the channel input (i.e., the transmitter) to have the perfect knowledge about the value of A and we also need $a \geq 1$ to fulfill the constraint $E[X^2] \leq P$. As the channel input is not aware of A=a,

$$C(P) < \frac{1}{2} \ln \left(1 + \frac{P}{\sigma^2} \right)$$

unless Pr[A = 1] = 1.

Note: So we can only fix X to be Gaussian distributed with $E[X^2] = P$. As such, $E[Y^2|A=a] = a^2P + \sigma^2$, and $h(Y|A=a) = \frac{1}{2}\ln\left(2\pi e(a^2P+\sigma^2)\right)$, which implies

$$h(Y|A) = E_A \left[\frac{1}{2} \ln \left(2\pi e (A^2 P + \sigma^2) \right) \right].$$

Note: When the value A is not known to the receiver but only its distribution, we shall determine the channel capacity via

$$I(X;Y) = h(Y) - h(Y|X) = h(Y) - h(AX + Z|X).$$

In such case, we cannot claim that h(AX + Z|X) = h(Z) and the determination of the exact capacity capacity formula is still an open problem.