

Arc-consistency: AC-3 algorithm

Can arc-consistency guarantee global consistency?

- Why re-add (X_k, X_i) ?
- Why not (X_i, X_k) ?
- Time complexity?
- $\mathcal{O}(dn^2 \cdot d^2)$
- Remove-Inconsistent-Values
 - $\mathcal{O}(d^2)$

function AC-3(*csp*) **returns** the CSP, possibly with reduced domains
inputs: *csp*, a binary CSP with variables $\{X_1, X_2, \dots, X_n\}$
local variables: *queue*, a queue of arcs, initially all the arcs in *csp*

```
while queue is not empty do  
   $(X_i, X_j) \leftarrow \text{REMOVE-FIRST}(\textit{queue})$   
  if REMOVE-INCONSISTENT-VALUES( $X_i, X_j$ ) then  
    for each  $X_k$  in NEIGHBORS[ $X_i$ ] do  
      add  $(X_k, X_i)$  to queue  
return true
```

function REMOVE-INCONSISTENT-VALUES(X_i, X_j) **returns** true iff succeeds
removed \leftarrow false
for each x **in** DOMAIN[X_i] **do**
 if no value y in DOMAIN[X_j] allows (x, y) to satisfy the constraint $X_i \leftrightarrow X_j$
 then delete x from DOMAIN[X_i]; *removed* \leftarrow true
return removed

CSP search algorithms

- Depth-first search (depth = n)
- Initial state: $\{\}$
- Successor function: assign **any variable**: $d \cdot (n - l)$
- Goal test: all variable are assigned.
- Inefficient: $d^n \cdot n!$
- Improvement: assign variable in a **fixed order (backtracking!)**