

# Pre-processing: Constraint propagation

- Node consistency: reduce the domain for unary constraints
- Arc consistency: every element in the domain satisfies binary constraints
- $x_i \in X_i : \exists x_j \in X_j$  such that  $(x_i, x_j)$  satisfies binary consistency.



# Arc-consistency: AC-3 algorithm

## Can arc-consistency guarantee global consistency?

- Why re-add  $(X_k, X_i)$ ?
- Why not  $(X_i, X_k)$ ?
- Time complexity?
- $\mathcal{O}(dn^2 \cdot d^2)$
- Remove-Inconsistent-Values
  - $\mathcal{O}(d^2)$

**function** AC-3(*csp*) **returns** the CSP, possibly with reduced domains  
**inputs:** *csp*, a binary CSP with variables  $\{X_1, X_2, \dots, X_n\}$   
**local variables:** *queue*, a queue of arcs, initially all the arcs in *csp*

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while queue is not empty do  
   $(X_i, X_j) \leftarrow \text{REMOVE-FIRST}(\textit{queue})$   
  if REMOVE-INCONSISTENT-VALUES( $X_i, X_j$ ) then  
    for each  $X_k$  in NEIGHBORS[ $X_i$ ] do  
      add  $(X_k, X_i)$  to queue  
return true
```

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**function** REMOVE-INCONSISTENT-VALUES( $X_i, X_j$ ) **returns** true iff succeeds  
*removed*  $\leftarrow$  false  
**for each**  $x$  **in** DOMAIN[ $X_i$ ] **do**  
 **if** no value  $y$  in DOMAIN[ $X_j$ ] allows  $(x, y)$  to satisfy the constraint  $X_i \leftrightarrow X_j$   
 **then** delete  $x$  from DOMAIN[ $X_i$ ]; *removed*  $\leftarrow$  true  
**return removed**