

1 map concat

$(.) :: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow a \rightarrow c$
 $f . g = \lambda x \rightarrow f (g x)$

$concat :: [[a]] \rightarrow [a]$
 $concat [] = []$
 $concat (x:xs) = x ++ (concat xs)$

$map :: (a \rightarrow b) \rightarrow [a] \rightarrow [b]$
 $map f [] = []$
 $map f (x:xs) = (f x) : (map f xs)$

Claim 1.1.

$map f.concat = concat.map (map f)$

Proof. structural induction on xs
 $xs = []$

$(map f.concat) xs = (map f.concat) []$
 $= map f (concat [])$
 $= map f []$
 $= []$
 $= concat []$
 $= concat (map (map f) [])$
 $= (concat.map (map f)) []$
 $= (concat.map (map f)) xs$

$xs = (x:xs')$ structural induction on x
 $x = []$

$(map f.concat) xs = (map f.concat) (x:xs')$
 $= map f (concat ([]:xs'))$
 $= map f ([] ++ (concat xs'))$
 $= map f (concat xs')$
 $\stackrel{IH}{=} (concat.map (map f)) xs'$
 $= concat (map (map f) xs')$
 $= [] ++ (concat (map (map f) xs'))$
 $= concat ([] : (map (map f) xs'))$
 $= concat ((map f []) : (map (map f) xs'))$
 $= concat (map (map f) ([]:xs'))$
 $= concat (map (map f) xs)$
 $= (concat.map (map f)) xs$

$x = (y:ys)$

$(map f.concat) xs = (map f.concat) (x:xs')$
 $= map f (concat ((y:ys):xs'))$
 $= map f ((y:ys) ++ (concat xs'))$
 $= map f (y : (ys ++ (concat xs')))$
 $= (f y) : (map f (ys ++ (concat xs')))$

```

= (f y):(map f (concat (ys:xs')))
IH
= (f y):((concat.map (map f)) (ys:xs'))
= (f y):(concat (map (map f) (ys:xs')))
= (f y):(concat ((map f ys):(map (map f) xs'))))
= (f y):((map f ys)++(concat (map (map f) xs')))
= ((f y):(map f ys))++(concat (map (map f) xs'))
= (map f (y:ys))++(concat (map (map f) xs'))
= (map f x)++(concat (map (map f) xs'))
= concat ((map f x):(map (map f) xs'))
= concat (map (map f) (x:xs'))
= concat (map (map f) xs)
= (concat.map (map f)) xs

```

□

alternative proof

Proof. we show the equivalent claim

$\text{map } f (\text{concat } xs) = \text{concat } (\text{map } (\text{map } f) xs)$

structural induction on xs

$xs = []$

```
map f (concat xs) = map f (concat [])
                  = map f []
                  = []
                  = concat []
                  = concat (map (map f) [])
                  = concat (map (map f) xs)
```

$xs = (x:xs')$ structural induction on x

$x = []$

```
map f (concat xs) = map f (concat ([]:xs'))
                  = map f ([]++(concat xs'))
                  = map f (concat xs')
IH
                  = concat (map (map f) xs')
                  = []++(concat (map (map f) xs'))
                  = concat ([]:(map (map f) xs'))
                  = concat ((map f []):(map (map f) xs'))
                  = concat (map (map f) ([]:xs'))
                  = concat (map (map f) xs)
```

$x = (y:ys)$

```
map f (concat xs) = map f (concat ((y:ys):xs'))
                  = map f ((y:ys)++(concat xs'))
                  = map f (y:(ys++(concat xs')))
                  = (f y):(map f (ys++(concat xs')))
                  = (f y):(map f (concat (ys:xs')))
IH
                  = (f y):(concat (map (map f) (ys:xs')))
                  = (f y):(concat ((map f ys):(map (map f) xs')))
                  = (f y):((map f ys)++(concat (map (map f) xs')))
                  = ((f y):(map f ys))++(concat (map (map f) xs'))
                  = (map f (y:ys))++(concat (map (map f) xs'))
                  = (map f x)++(concat (map (map f) xs'))
                  = concat ((map f x):(map (map f) xs'))
                  = concat (map (map f) (x:xs'))
                  = concat (map (map f) xs)
```

□

2 foldl map

```
foldl :: (a -> b -> a) -> a -> [b] -> a
foldl f a []      = a
foldl f a (x:xs) = foldl f (f a x) xs
```

Claim 2.1.

```
foldl f a.map g = foldl h a
                  h b x = f b (g x)
```

Proof. we show the equivalent claim

```
foldl f a (map g xs) = foldl h a xs
                      h b x = f b (g x)
```

by structural induction on xs

xs=[]

```
foldl f a (map g xs) = foldl f a (map g [])
                      = foldl f a []
                      = a
                      = foldl h a []
                      = foldl h a xs
```

xs=(x:xs')

```
foldl f a (map g xs) = foldl f a (map g (x:xs'))
                      = foldl f a ([g x]++(map g xs'))
                      = foldl f (f a (g x)) (map g xs')
                       $\stackrel{IH}{=}$  foldl h (f a (g x)) xs'
                      = foldl h (h a x) xs'
                      = foldl h a (x:xs')
                      = foldl h a xs
```

□

3 bag

```

data Bag a = ListBag [(a,Integer)] deriving (Eq,Show)

add :: Eq a => (a,Integer) -> (Bag a) -> (Bag a)
add (ele,count) (ListBag bagx)
  | ele 'elem' map fst bagx = ListBag (map (inc (ele,count)) bagx)
  | otherwise               = ListBag (bagx++[(ele,count)])

inc :: Eq a => (a,Integer) -> (a,Integer) -> (a,Integer)
inc (e,c) (x,z)
  | e==x    = (e,c+z)
  | otherwise = (x,z)

unite :: Eq a => (Bag a) -> (Bag a) -> (Bag a)
unite (ListBag [])      bagy = bagy
unite (ListBag (x:xs)) bagy = unite (ListBag xs) (add x bagy)

bag :: Eq a => [a] -> (Bag a)
bag []      = ListBag []
bag (x:xs) = add (x,1) (bag xs)

```

Claim 3.1.

$\text{bag } (xs++ys) = \text{unite } (\text{bag } xs) (\text{bag } ys)$

used but unproven statements:

```

S1: map f (xs++ys) = (map f xs)++(map f ys)
S2: map f (map g xs) = map (f.g) xs
S2: (inc (e,z)).(inc (e,c)) = inc (e,z+c)
S4: (inc (f,z)).(inc (e,c)) = (inc (e,c)).(inc (f,z))

```

Lemma 3.2. *if A1: not e 'elem' map fst xs inc does not do anything*

$xs = \text{map } (\text{inc } (e,c)) \text{ } xs$

Proof. of Lemma 3.2 by structural induction on xs

A1: not e 'elem' map fst xs

xs=[]

```

xs = []
    = map (inc (e,c)) []
    = map (inc (e,c)) xs

```

xs=((f,z):xs')

```

xs = (f,z):xs'
     $\stackrel{IH}{=}$  (f,z):(map (inc (e,c)) xs')
     $\stackrel{A1}{=}$  (inc (e,c) (f,z)):(map (inc (e,c)) xs')
    = map (inc (e,c)) ((f,z),xs')
    = map (inc (e,c)) xs

```

□

Lemma 3.3. *adding same element*

$$\text{add } (e, z) \text{ (add } (e, c) \text{ (ListBag } xs)) = \text{add } (e, c+z) \text{ (ListBag } xs)$$

Proof. of Lemma 3.3

Case 1 : A1: not e 'elem' map fst xs

$$\begin{aligned} \text{add } (e, z) \text{ (add } (e, c) \text{ (ListBag } xs)) & \\ & \stackrel{A1}{=} \text{add } (e, z) \text{ (ListBag } xs++[(e, c)]) \\ & = \text{ListBag (map (inc } (e, z)) \text{ (xs++[(e, c)]))} \\ & \stackrel{S1}{=} \text{ListBag (map (inc } (e, z)) \text{ xs)++(map (inc } (e, z)) [(e, c)])} \\ & \stackrel{3.2, A1}{=} xs++(\text{map (inc } (e, z)) [(e, c)]) \\ & = xs++((\text{inc } (e, z) \text{ (e, c))} : (\text{map (inc } (e, z)) [])) \\ & = xs++((\text{inc } (e, z) \text{ (e, c))} : []) \\ & = xs++((e, z+c) : []) \\ & = xs++[(e, z+c)] \\ & \stackrel{A1}{=} \text{add } (e, c+z) \text{ (ListBag } xs) \end{aligned}$$

Case 2 : A1: e 'elem' map fst xs

$$\begin{aligned} \text{add } (e, z) \text{ (add } (e, c) \text{ (ListBag } xs)) & \\ & \stackrel{A1}{=} \text{add } (e, z) \text{ (ListBag (map (inc } (e, c)) \text{ xs))} \\ & \stackrel{A1}{=} \text{ListBag (map (inc } (e, z)) \text{ (map (inc } (e, c)) \text{ xs))} \\ & \stackrel{S2}{=} \text{ListBag (map (inc } (e, z)) \text{ (inc } (e, c)) \text{ xs)} \\ & \stackrel{S3}{=} \text{map (inc } (e, c+z)) \text{ xs} \\ & \stackrel{A1}{=} \text{add } (e, c+z) \text{ (ListBag } xs) \end{aligned}$$

□

Lemma 3.4. *adding different element*

A1: e 'elem' map fst xs

A2: f != e

$$\text{add } (f, z) \text{ (add } (e, c) \text{ (ListBag } xs)) = \text{add } (e, c) \text{ (add } (f, z) \text{ (ListBag } xs))$$

Proof. of Lemma 3.4

A1: not e 'elem' map fst xs

Case 1 : A3: not f 'elem' map fst xs

$$\begin{aligned} \text{add } (f, z) \text{ (add } (e, c) \text{ (ListBag } xs)) & \\ & \stackrel{A1}{=} \text{add } (f, z) \text{ (ListBag (map (inc } (e, c)) \text{ xs))} \\ & \stackrel{A3}{=} \text{ListBag ((map (inc } (e, c)) \text{ xs)++[(f, z)]}) \\ & \stackrel{A3}{=} \text{ListBag ((map (inc } (e, c)) \text{ xs)++[inc } (e, c) \text{ (f, z)]})} \\ & = \text{ListBag ((map (inc } (e, c)) \text{ xs)++(map (inc } (e, c)) [(f, z)]))} \\ & \stackrel{S1}{=} \text{ListBag (map (inc } (e, c)) \text{ (xs++[(f, z)]))} \\ & \stackrel{A1, A2}{=} \text{add } (e, c) \text{ (ListBag (xs++[(f, z)]))} \\ & \stackrel{A3}{=} \text{add } (e, c) \text{ (add } (f, z) \text{ (ListBag } xs)) \end{aligned}$$

Case 2 : A3: f 'elem' map fst xs

```

add (f,z) (add (e,c) (ListBag xs))
A1 = add (f,z) (ListBag (map (inc (e,c)) xs))
A3 = ListBag (map (inc (f,z)) (map (inc (e,c)) xs))
S2 = ListBag (map (inc (f,z)).(inc (e,c)) xs)
S4 = ListBag (map (inc (e,c)).(inc (f,z)) xs)
S2 = ListBag (map (inc (e,c)) (map (inc (f,z)) xs))
A1 = add (e,c) (ListBag (map (inc (f,z)) xs))
A3 = add (e,c) (add (f,z) (ListBag xs))

```

□

Lemma 3.5. *if* A1: not e 'elem' map fst xs

```

unite (ListBag ((e,c):(xs++[(e,z)]++ys))) bagy
    = unite (ListBag (e,c+z):(xs++ys))) bagy

```

Proof. of Lemma 3.5 by structural induction on xs

A1: not e 'elem' map fst xs

xs=[]

```

unite (ListBag ((e,c):(xs++[(e,z)]++ys))) bagy
    = unite (ListBag ((e,c):([]++[(e,z)]++ys))) bagy
    = unite (ListBag ((e,c):([(e,z)]++ys))) bagy
    = unite (ListBag [(e,z)]++ys) (add (e,c) bagy)
    = unite (ListBag ys) (add (e,z) (add (e,c) bagy))
3,3 = unite (ListBag ys) (add (e,c+z) bagy)
    = unite (ListBag ((e,c+z):ys)) bagy
    = unite (ListBag ((e,c+z):(xs++ys))) bagy

```

xs=xs'++[x]

```

unite (ListBag ((e,c):(xs++[(e,z)]++ys))) bagy
    = unite (ListBag (xs++[(e,z)]++ys) (add (e,c) bagy)
    = unite (ListBag (xs'++[x]++[(e,z)]++ys) (add (e,c) bagy)
    = unite (ListBag ys) (add (e,z) (add x (add ... (add (e,c) bagy)...)))
3,4 = unite (ListBag ys) (add x (add (e,z) (add ... (add (e,c) bagy)...)))
    = unite (ListBag (xs'++[(e,z)]++[x]++ys) (add (e,c) bagy)
    = unite (ListBag ((e,c):(xs'++[(e,z)]++[x]++ys))) bagy
IH = unite (ListBag ((e,c+z):(xs'++[x]++ys))) bagy
    = unite (ListBag ((e,c+z):(xs++ys))) bagy

```

□

Lemma 3.6.

```

add (e,1) (unite (ListBag xs) bagy)
    = unite (add (e,1) (ListBag xs)) bagy

```

Proof. of Lemma 3.6 by structural induction on xs
 $xs = []$

```
add (e,1) (unite (ListBag xs) bagy)
  = add (e,1) (unite (ListBag []) bagy)
  = add (e,1) bagy
  = unite (ListBag []) (add (e,1) bagy)
  = unite (ListBag [(e,1)]) bagy
  = unite (add (e,1) (ListBag [])) bagy
  = unite (add (e,1) (ListBag xs)) bagy
```

$xs = (x:xs')$

```
add (e,1) (unite (ListBag xs) bagy)
  = add (e,1) (unite (ListBag (x:xs')) bagy)
  = add (e,1) (unite (ListBag xs') (add x bagy))
IH = unite (add (e,1) (ListBag xs')) (add x bagy)
```

Case 1 : A1: not e 'elem' map fst xs

```
IH = unite (add (e,1) (ListBag xs')) (add x bagy)
A1 = unite (ListBag (xs'++[(e,1)])) (add x bagy)
  = unite (ListBag (x:(xs'++[(e,1)]))) bagy
  = unite (ListBag ((x:xs')++[(e,1)])) bagy
  = unite (ListBag (xs++[(e,1)])) bagy
A1 = unite (add (e,1) (ListBag xs)) bagy
```

Case 2.1 : A1: e 'elem' map fst xs and A2: $x = (f,c)$ $f \neq e$
 A1 and A2 imply A3: e 'elem' map fst xs'

```
IH = unite (add (e,1) (ListBag xs')) (add x bagy)
A3 = unite (ListBag (map (inc (e,1)) xs')) (add x bagy)
  = unite (ListBag (x:(map (inc (e,1)) xs'))) bagy
A2 = unite (ListBag ((inc (e,1) x):(map (inc (e,1)) xs'))) bagy
  = unite (ListBag (map (inc (e,1)) (x:xs'))) bagy
A1 = unite (add (e,1) (ListBag xs)) bagy
```

Case 2.2 : A1: e 'elem' map fst xs and A2: $x = (e,c)$
 multiset, A1 and A2 imply A3: not e 'elem' map fst xs'

```
IH = unite (add (e,1) (ListBag xs')) (add x bagy)
A3 = unite (ListBag (xs'++[(e,1)])) (add x bagy)
  = unite (ListBag (x:(xs'++[(e,1)]))) bagy
  = unite (ListBag ((e,c):(xs'++[(e,1)]))) bagy
L3,5 = unite (ListBag ((e,c+1):xs')) bagy
  = unite (ListBag ((inc (e,1) (e,c)):xs')) bagy
L3,2 = unite (ListBag ((inc (e,1) x):(map (inc (e,1)) xs'))) bagy
  = unite (ListBag (map (inc (e,1)) (x:xs'))) bagy
  = unite (ListBag (map (inc (e,1)) xs)) bagy
A1 = unite (add (e,1) (ListBag xs)) bagy
```


□

Lemma 3.7. *Associativity of unite*

```
unite (ListBag xs) (unite (ListBag ys) (ListBag zs))
  = unite (unite (ListBag xs) (ListBag ys)) (ListBag zs)
```

Proof. by structural induction on xs

xs=[]

```
unite (ListBag xs) (unite (ListBag ys) (ListBag zs))
  = unite (ListBag []) (unite (ListBag ys) (ListBag zs))
  = unite (ListBag ys) (ListBag zs)
  = unite (unite (ListBag []) (ListBag ys)) (ListBag zs)
  = unite (unite (ListBag xs) (ListBag ys)) (ListBag zs)
```

xs=(x:xs')

```
unite (ListBag xs) (unite (ListBag ys) (ListBag zs))
  = unite (ListBag (x:xs')) (unite (ListBag ys) (ListBag zs))
  = unite (ListBag xs') (add x (unite (ListBag ys) (ListBag zs)))
L3.6 = unite (ListBag xs') (unite (add x (ListBag ys)) (ListBag zs))
IH = unite (unite (ListBag xs') (add x (ListBag ys))) (ListBag zs)
  = unite (unite (ListBag (x:xs')) (ListBag ys)) (ListBag zs)
  = unite (unite (ListBag xs) (ListBag ys)) (ListBag zs)
```

□

Proof. of Claim 3.1 by structural induction on xs

xs=[]

```
bag (xs++ys) = bag ([]++ys)
  = bag ys
  = unite (ListBag []) (bag ys)
  = unite (bag []) (bag ys)
  = unite (bag xs) (bag ys)
```

xs=(x:xs')

```
bag (xs++ys) = bag ((x:xs')++ys)
  = bag (x:(xs'++ys))
  = add (x,1) (bag (xs'++ys))
IH = add (x,1) (unite (bag xs') (bag ys))
  = add (x,1) (unite (bag xs') (bag ys))
  = unite (ListBag []) (add (x,1) (unite (bag xs') (bag ys)))
  = unite (ListBag [(x,1)]) (unite (bag xs') (bag ys))
Ass = unite (unite (ListBag [(x,1)]) (bag xs')) (bag ys)
  = unite (unite (ListBag []) (add (x,1) (bag xs'))) (bag ys)
  = unite (add (x,1) (bag xs')) (bag ys)
  = unite (bag (x:xs')) (bag ys)
  = unite (bag xs) (bag ys)
```

□