

1 WFI or not?

```
take 0 xs      = []
take n []      = []
take n (x:xs) = [x]++(take (n-1) xs)
```

```
drop 0 xs      = xs
drop n []      = []
drop n (x:xs) = drop (n-1) xs
```

Claim 1.1. The following equation holds.

$$(\text{take } n \text{ } xs)++(\text{drop } n \text{ } xs) = xs$$

Proof. by induction over n .

BC: $n=0$:

$$\begin{aligned} & (\text{take } n \text{ } xs)++(\text{drop } n \text{ } xs) \\ &= (\text{take } 0 \text{ } xs)++(\text{drop } 0 \text{ } xs) \\ &= []++xs \\ &= xs \end{aligned}$$

IH: $(\text{take } n' \text{ } xs)++(\text{drop } n' \text{ } xs) = xs$

holds for an arbitrary but fixed n' and all xs .

IS: $n=n'+1$

Case 1: $xs=[]$

$$\begin{aligned} & (\text{take } n \text{ } xs)++(\text{drop } n \text{ } xs) \\ &= (\text{take } (n'+1) \text{ } [])++(\text{drop } (n'+1) \text{ } []) \\ &= []++[] \\ &= [] \\ &= xs \end{aligned}$$

Case 2: $xs=(x:xs')$

$$\begin{aligned} & (\text{take } n \text{ } xs)++(\text{drop } n \text{ } xs) \\ &= (\text{take } (n'+1) \text{ } (x:xs'))++(\text{drop } (n'+1) \text{ } (x:xs')) \\ &= [x]++(\text{take } n' \text{ } xs')++(\text{drop } n' \text{ } xs') \\ &= [x]++xs' \\ &= xs \end{aligned}$$

□

```

zip [] ys          = []
zip xs []          = []
zip (x:xs) (y:ys) = [(x,y)] ++ (zip xs ys)

length []          = 0
length (x:xs)      = 1 + (length xs)

min x y
  | x > y  = y
  | true   = x

```

Claim 1.2. The following equation holds.

```
length (zip xs ys) = min (length xs) (length ys)
```

Proof. by structural induction over xs.

BC: xs=[]

```

length (zip xs ys)
= length (zip [] ys)
= length []
= 0
= min 0 (length ys) --since (length ys) >= 0
= min (length xs) (length ys)

```

IH: `length (zip xs' ys) = min (length xs') (length ys)`

holds for an arbitrary but fixed xs' and all ys.

IS: xs=(x:xs')

Case 1: ys=[]

```

length (zip xs ys)
= length (zip (x:xs') [])
= length []
= min (length (x:xs')) 0 --since (length (x:xs')) > 0
= min (length xs) (length ys)

```

Case 2: ys=(y:ys')

```

length (zip xs ys)
= length (zip (x:xs') (y:ys'))
= length ([(x,y)] ++ (zip xs' ys'))
= 1 + length (zip xs' ys')
= 1 + (min (length xs') (length ys'))
= min (1 + (length xs')) (1 + (length ys'))
= min (length (x:xs')) (length (y:ys'))
= min (length xs) (length ys)

```

□

But now just for the fun of it let us proof Claim 1.2 by well founded induction.

$$(xs, ys) \prec ((x : xs), (y : ys))$$

We denote the transitive closure of \prec by $<$. Obviously $<$ is a well-founded order.

Proof. by well-founded induction with the predicate

$P((xs, ys)) = (\text{length } (\text{zip } xs \text{ } ys) = \min (\text{length } xs) (\text{length } ys))$

$xs=[]$: (Note that $([], ys)$ has no successors w.r.t. $<$.)

```

length (zip xs ys)
= length (zip [] ys)
= length []
= 0
= min 0 (length ys) --since (length ys) >= 0
= min (length xs) (length ys)

```

$ys=[]$: This case can be treated analogously to $xs=[]$.

$xs=(x:xs'), ys=(y:ys')$

```

length (zip xs ys)
= length (zip (x:xs') (y:ys'))
= length [(x,y)] ++ (zip xs' ys')
= 1 + (length (zip xs' ys'))
= 1 + (min (length xs') (length ys')) {-since (xs',ys') < (xs,ys)
                                         P((xs',ys')) holds-}
= min (1 + (length xs')) (1 + (length ys'))
= min (length (x:xs')) (length (y:ys'))
= min (length xs) (length ys)

```

□

1.1 Lexicographic Order

```

a 0 m = m+1
a n 0 = a (n-1) 1
a n m = a (n-1) (a n (m-1))

alist [] ys = []++ys
alist (x:xs) [] = alist xs [x]
alist (x:xs) (y:ys) = alist xs (alist (x:xs) ys)

```

Claim 1.3. The following equation holds.

```
length (alist xs ys) = a (length xs) (length ys)
```

$$\begin{aligned}
 (xs, ys) &\prec_l (xs, (y : ys)) \\
 (xs, zs) &\prec_l ((x : xs), ys)
 \end{aligned}$$

1.2 Multiset Order

```

r [] = []
r (x:xs) = r (xs++[(x+1)]++[(x+1)])

```

Claim 1.4. The following equation holds.

```
r xs = []
```