

Basic Data Structures

C++ STL

Computational Complexity

- Complexity = Resources required
- Time
 - TLE
- Space
 - MLE
- Description
 - Too late
- Communication

Asymptotic Notation

- We often don't need precise bounds on complexity.
- For large enough n :
 - $T(n)=O(f(n))$: $T(n) \leq c_1 f(n)$ = means \in
 - $T(n)=\Omega(f(n))$: $T(n) \geq c_2 f(n)$
 - $T(n)=\Theta(f(n))$: $c_2 f(n) \leq T(n) \leq c_1 f(n)$
 - $T(n)=o(f(n))$: $\lim_{n \rightarrow \infty} \frac{T(n)}{f(n)} = 0$
 - $T(n)=\omega(f(n))$: $\lim_{n \rightarrow \infty} \frac{f(n)}{T(n)} = 0$

small- ω	big- Ω	Θ	big- O	Small- o
$>$	\geq	$=$	\leq	$<$

Growth of Function

- Factorial: $O(n!)$
- Exponential: $O(c^n)$
- Polynomial: $O(n^k)$
 - Cubic: $O(n^3)$
 - Square: $O(n^2)$
- Linearithmic: $O(n \log n)$
- Linear: $O(n)$
- Logarithmic: $O(\log n)$
- Constant: $O(1)$

STL

- Standard template library
- Very useful in competitive programming
- Composition
 - Containers
 - Iterators
 - Algorithms
 - Functors
- Will cover a small part of STL in this lecture

Useful Containers in Contest

- pair
- vector
- list
- queue
- deque
- priority_queue
- set
- map

pair

- Couple two values into a pair
- `pair<string,int> p("abc",123);`
- `make_pair(string("abc"),123);`
- `p.first` // access the first element
- `p.second` // access the second element
- Compare the field “first” first, then compare the field “second.”

vector

- An array cannot change its size but a vector can.
- vector is still fast enough in general.
- `vector<int> a;` //create a vector initially empty
- `vector<int> a(n);` //create a vector of n integers
- `vector<int> a(n,5);` // create a vector of n integers
// and the integers are all 5
- k-th element of a: `a[k]`

vector

- `a.back()` // the last element of `a`
- `a.push_back(x);` // append `x` to vector `a`
- `a.pop_back();` // remove the last element of `a`

vector as a stack

- LIFO: last-in first-out
- `vector<int> s; // make an empty stack s of integers`
- `s.empty()` // test if s is empty
- `s.size()` // number of elements in s
- `s.push_back(x)` // insert x into s
- `s.pop_back()` // remove the last inserted element
- `s.back()` // the top element in the stack
- class stack in STL is slower than vector.

queue

- FIFO: first-in first-out
- `queue<int> q;` // make an empty queue of integers
- `q.empty()` // test if q is empty
- `q.size()` // number of elements in q
- `q.push(x)` // insert x into q
- `q.pop()` // remove an element from q
- `q.front()` // the next element to be popped
- `q.back()` // the last element pushed

priority_queue

- Greatest first out
- `priority_queue<int> pq; // make an empty priority
// queue of integers`
- `pq.empty()`
- `pq.size()`
- `pq.push()`
- `pq.pop()`
- `pq.top() // get the first element`

set (Ordered Set)

- Stores **unique** elements in certain **order**
- `set<int> s; //make an empty set of integers`
 - `s.empty()`
 - `s.size()`
- `s.insert(x) // insert x into s`
- `s.erase(x) // remove x from s`
- `s.count(x) // number of copies of x in s`
- *****`s.begin()` // first element in s
- *****`s.rbegin()` // last element in s

map (Ordered Map)

- map a **key** into a unique **value**.
- Array is a kind of map. 0, ..., n-1 are keys of `int[n]`.
- map supports non-consecutive keys
- `map<string,int> cnt; // initialize an empty map
// from string to int`
- `cnt.empty()`
- `cnt.size()`
- `cnt.count(x) // number of elements having key x`

map (Ordered Map)

- `cnt["key"] = x` // map "key" into x
- `y = cnt["key"]` // get the value from "key"
- `*cnt->begin()` // the smallest key-value pair
 - `cnt->begin()->first` // key of the smallest element
 - `cnt->begin()->second` // value of the smallest element
- `*cnt->rbegin()` // the greatest key-value pair
 - `cnt->rbegin()->first`
 - `cnt->rbegin()->second`

Useful Algorithms in Contest

- `sort`
- `lower_bound`
- `upper_bound`
- `next_permutation`

sort()

- Sorting elements into ascending order
- `sort(a,a+100); // sort int a[100]`
- `sort(v.begin(), v.end()); // sort vector v`
- The elements must be comparable
 - int and string are comparable
 - `public bool operator<(const T& rhs) const`
- In most cases, `sort()` is fast enough.
- Time complexity: $O(n \log n)$

lower_bound

- `lower_bound(v.begin(), v.end(), x)`
 - Finding the first position that is not $< x$ in vector `v`
 - Returns an iterator
 - Works only if all elements are not $< x$ after that position.
For example, `v` is sorted in ascending order.
- Implementation: variant binary search
- Time complexity: $O(\log n)$

lower_bound

1	2	2	3	3	3	4	4	4	4	5	5	5	5	5
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---



lower_bound of 4

1	2	2	3	3	3	3	3	3	5	5	5	5	5	5
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---



lower_bound of 4

1	2	1	3	1	3	1	5	4	5	6	7	6	4	5
<	<	<	<	<	<	<	≥	≥	≥	≥	≥	≥	≥	≥

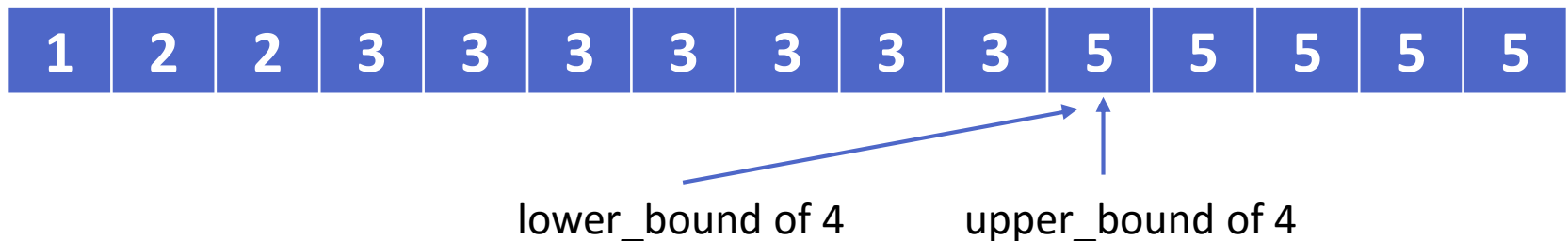
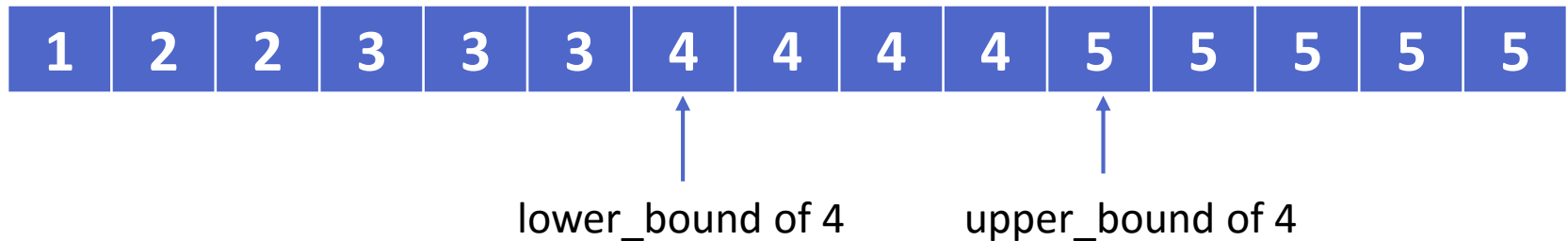


lower_bound of 4

upper_bound

- `upper_bound(v.begin(), v.end(), x)`
 - Finding the first position that is not $\leq x$ in vector `v`
 - Returns an iterator
 - Works only if all elements are not $\leq x$ after that position.
For example, `v` is sorted in ascending order.
- Implementation: variant binary search
- Time complexity: $O(\log n)$

lower_bound and upper_bound



Might be too slow in
some application!

next_permutation

- Permutation: reordering of a sequence
- Lexicographic order of two permutation p and q
 - $p < q$ if there exists i such that $p[i] < q[i]$ and $p[j] = q[j]$ for every $j < i$.
- Next permutation of p : the minimum of all permutation greater than p .
- `next_permutation`:
 - If the input has next permutation, then modify the input into its next permutation and return `true`.
 - Otherwise, return `false`.

Next Permutation

