

P2

Ache todos $f: \mathbb{R} \rightarrow \mathbb{R}$ tais que

$$f(x+y f(x+y)) = y^2 + f(x) f(y), \quad \forall x, y \in \mathbb{R}.$$

$P(x, y)$

$P(x, 0)$: $f(x) = f(x) \cdot f(0) \Rightarrow f(x) = 0, \forall x$ {Teste e vê que não funciona.
ou $f(0) = 1$

$\Rightarrow f(0) = 1.$

$P(x, -x)$: $1 = x^2 + f(x) f(-x)$

$f(x) f(-x) = 1 - x^2 = (1+x)(1-x)$

$x \rightarrow 1$ $0 = f(1) \cdot f(-1) \Rightarrow \underline{f(1) = 0}$ ou $f(-1) = 0$

Seja $a \in \{1, -1\}$ t.q. $f(a) = 0.$

Se $f(x_0) = 0$:

$P(0, x_0)$: $f(x_0 f(x_0)) = x_0^2 + f(x_0) f(x_0)$

$1 = f(x_0) = x_0^2 \Rightarrow x_0 = \pm 1.$

$P(x, 1)$: $f(x + f(x+1)) = 1$

Se $f(x_1) = 1$.

$P(0, x_1)$: $f(x_1 f(x_1)) = x_1^2 + f(x_1) f(x_1)$

$1 = x_1^2 + 1$

$x_1 = 0$

Se $f(x_1) = 1$

$\Rightarrow x_1 = 0$

$P(x, a)$: $f(x + a f(x+a)) = 1 \Rightarrow$

$\Rightarrow x + a f(x+a) = 0$

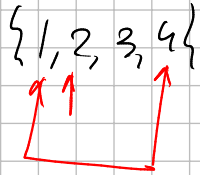
$f(x+1) = -x \Rightarrow f(x) = 1-x.$

$x = f(x-1) \Rightarrow f(x) = x+1.$

A B

P3

$p=3 \quad 2p-2=4.$

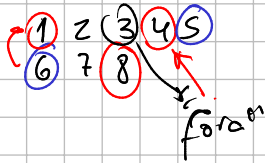


$2 \cdot 1 + 1 \equiv 0.$

Ana ganha.

A nunca perde. A tem estratégia não-perdedora: Nunca pegar p.

$p=5: \quad 2p-2=8$



$(p-1)! + 1 \equiv 0 \pmod{p}$

caso!
Ana ganha

• $n=1$. \cup .

• $n=2$. a_1, a_2 .

$(1, 1) \rightarrow (1, 2) \rightarrow (2, 1) \rightarrow \dots$

$$\sum \frac{1}{a_i} \leq \frac{1}{2} \Rightarrow \frac{1}{a_1} \leq \frac{1}{2} \Rightarrow a_1 \geq 2 \quad \underline{\text{OK!}}$$

DIMINUIR
O MAIOR

• $n=3$. a_1, a_2, a_3

$(1, 1, 1) \rightarrow (2, 2, 1) \rightarrow (3, 1, 2) \rightarrow (1, 2, 3) \rightarrow (2, 3, 1) \rightarrow \dots$

$$\sum \frac{1}{a_i} \leq \frac{1}{2} \Rightarrow \frac{1}{a_1} < \frac{1}{2} \Rightarrow a_1 > 2 \Rightarrow a_1 \geq 3 \quad \underline{\text{OK!}}$$

• $n=4$. a_1, a_2, a_3, a_4

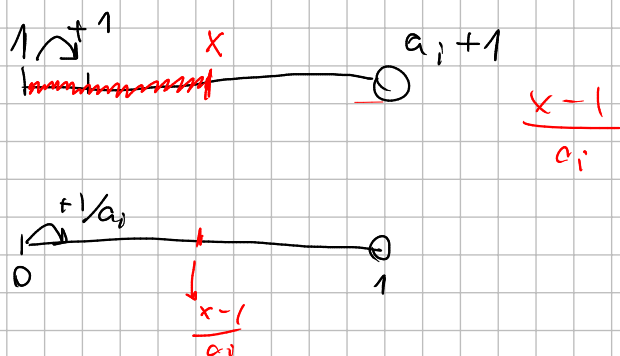
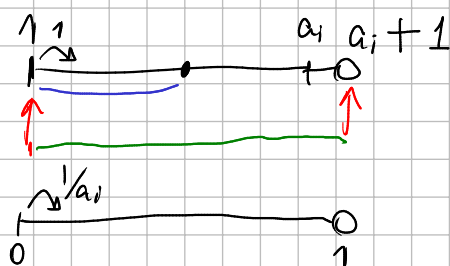
$$a_1 = 3 \quad a_2 = a_3 = a_4 = 18$$

DIMINUIR
O
MAIS PEQUENO
DO TERÇO.

1	1	1		1	2	3	\rightarrow
1	2	3	\dots	16	\rightarrow	1	2
1	2	3		16	17	\rightarrow	1
1	2	3		16	17	18	\rightarrow

1	\rightarrow	1	2	\rightarrow	1	2	\rightarrow	1	2	\rightarrow	1	2	\rightarrow	1
1		2	\rightarrow	1		2		3		4		5		6
1		2		3		4	\rightarrow	1		2		3		4
1		2		3		4		5		6	\rightarrow	1		2

DIMINUIR MAIS
PEQUENO DO TERÇO
QUE NÃO BATA
NO CHÃO.



(b_1, \dots, b_n) $b_i = \frac{0}{a_i}, \frac{1}{a_i}, \dots, 1 - \frac{1}{a_i}$.

0	\rightarrow	0	$\frac{1}{3} \rightarrow$	0	$\frac{1}{3} \rightarrow$	0	$\frac{1}{3} \rightarrow$	0
0		$\frac{1}{18} \rightarrow$	0	$\frac{1}{18} \rightarrow$	$\frac{2}{18}$	$\frac{3}{18}$	$\frac{4}{18}$	$\frac{5}{18}$
0		$\frac{1}{18}$	$\frac{2}{18}$	$\frac{3}{18} \rightarrow$	0	$\frac{1}{18}$	$\frac{2}{18}$	$\frac{3}{18}$
0		$\frac{1}{18}$	$\frac{2}{18}$	$\frac{3}{18}$	$\frac{4}{18}$	$\frac{5}{18} \rightarrow$	0	$\frac{1}{18}$

$$a_1 = 3$$

$$a_2 = a_3 = a_4 = 18.$$

\downarrow \downarrow
 $\frac{3}{18}$ $\frac{10}{18}$

\downarrow
 $\frac{12}{18}$

$$\frac{2}{3} = 1 - \frac{1}{n-2}$$

$$\begin{aligned}
 c_1 &\rightarrow c_1 + 1/a_1 \\
 c_2 &\rightarrow c_2 + 1/a_2 \\
 c_3 &\rightarrow \vdots \\
 c_n &\rightarrow \vdots \\
 &\vdots \rightarrow c_{n-1} + 1/a_{n-1} \\
 c_n &\rightarrow \underline{c_n + 1/a_n} \quad \circ
 \end{aligned}$$

$$\begin{array}{ccccccc}
 0 & \rightarrow & 0 & \rightarrow & 0 & & 0 \quad 1/3 \rightarrow 0 \quad 1/3 \rightarrow 0 \quad 1/3 \rightarrow 0 \\
 0 & & 1/18 & 2/18 & \dots & 6/18 \rightarrow 0 & 1/18 \quad 2/18 \quad 3/18 \quad 4/18 \quad 5/18 \\
 0 & & 1/18 & 2/18 & \dots & 6/18 \quad 7/18 & 8/18 \rightarrow 0 \quad 1/18 \quad 2/18 \quad 3/18 \\
 0 & & 1/18 & 2/18 & \dots & 6/18 \quad 7/18 & 8/18 \quad 9/18 \quad 10/18 \rightarrow 0 \quad 1/18
 \end{array}$$

$$c_1 \geq 1 - 0/a_1$$

Na jogada anterior,

$$c_1 \geq 1 - 1/a_1 \quad c_2 \geq 1 - 1/a_2$$

Na jogada anterior

$$c_i \geq 1 - 2/a_i, \quad \forall i \in \{1, \dots, 3\}$$

Na jogada anterior

$$c_i \geq 1 - 3/a_i, \quad \forall i \in \{1, \dots, 4\}$$

;

$$c_i \geq 1 - \frac{n-1}{a_i}, \quad \forall i \in \{1, \dots, n\}$$

$$\Rightarrow \sum c_i \geq n - (n-1) \cdot \sum \frac{1}{a_i} \geq \frac{n+1}{2}$$

$$\begin{array}{ccccccc}
 \Leftrightarrow & & \Leftrightarrow & & \Leftrightarrow & & \checkmark \\
 \left(\begin{array}{c} \leq (n-1)/a_1 \\ \leq (n-1)/a_2 \\ \vdots \\ \leq (n-1)/a_n \end{array} \right) & \Leftrightarrow & \begin{array}{c} \leq 2/a_1 \\ \vdots \\ \leq 2/a_n \end{array} & \Leftrightarrow & \begin{array}{c} \leq 1/a_1 \\ \vdots \\ \leq 1/a_n \end{array} & \Leftrightarrow & \begin{array}{c} 0/a_1 \\ \vdots \\ 2/a_1 \rightarrow 1 \\ 2/a_2 \rightarrow 1 \\ 3/a_2 \end{array}
 \end{array}$$

$$c_i \geq 1 - \frac{n-1}{a_i}$$

$$\sum c_i \geq n - (n-1) \sum \frac{1}{a_i} \geq n - \frac{n-1}{2} = \frac{n+1}{2}$$

ESTRATÉGIA INTELIGENTE
é

\Rightarrow NA RODADA X , $\sum c_i \geq \frac{n+1}{2}$.

PERDEU NA RODADA $X+N-1$

se \exists estratégia inteligente que garante $\sum c_i < \frac{n+1}{2}$. ACABOU.

$$\begin{aligned} c_1 &= (b_1 - 1) / a_1 \rightarrow \cancel{b_1 / a_1} \rightarrow 0 \text{ (} c_1 - b_1 \text{ mínimo, desde que } b_{i_1} \neq 0 \text{)} \\ c_2 &= (b_2 - 1) / a_2 \rightarrow b_2 / a_2 \\ &\vdots \\ c_n &= (b_n - 1) / a_n \rightarrow b_n / a_n \end{aligned}$$

RODADA é x-suicida

Se existam X índices i_1, \dots, i_X t.q.

$$a_{i_k} - b_{i_k} \leq X - 1.$$

1-suicida e/ A terminal.

$$c_1 = 3 \quad a_2 = a_3 = a_4 = 18$$

FALTA

$$1 \rightarrow 3 \rightarrow 2 \rightarrow 1 \rightarrow \emptyset$$

$$2 \rightarrow \emptyset$$

$$3 \rightarrow 2 \rightarrow \emptyset$$

$$4 \rightarrow 3 \rightarrow 2 \rightarrow \emptyset$$

$$AB \cap CD = P$$

$$AD \cap BC = Q$$

(A, B, C, D)

