

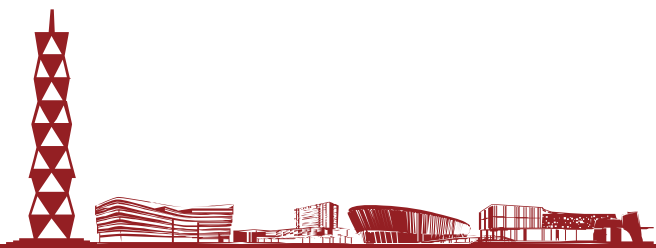


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CS271 Computer Graphics II

Lecture 6

Mesh Simplification

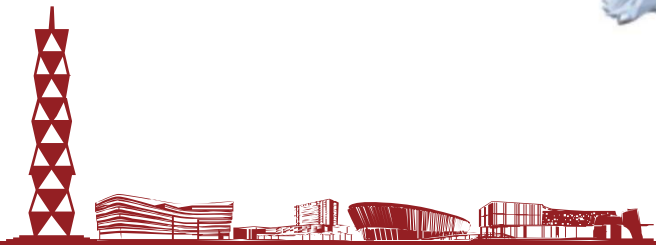
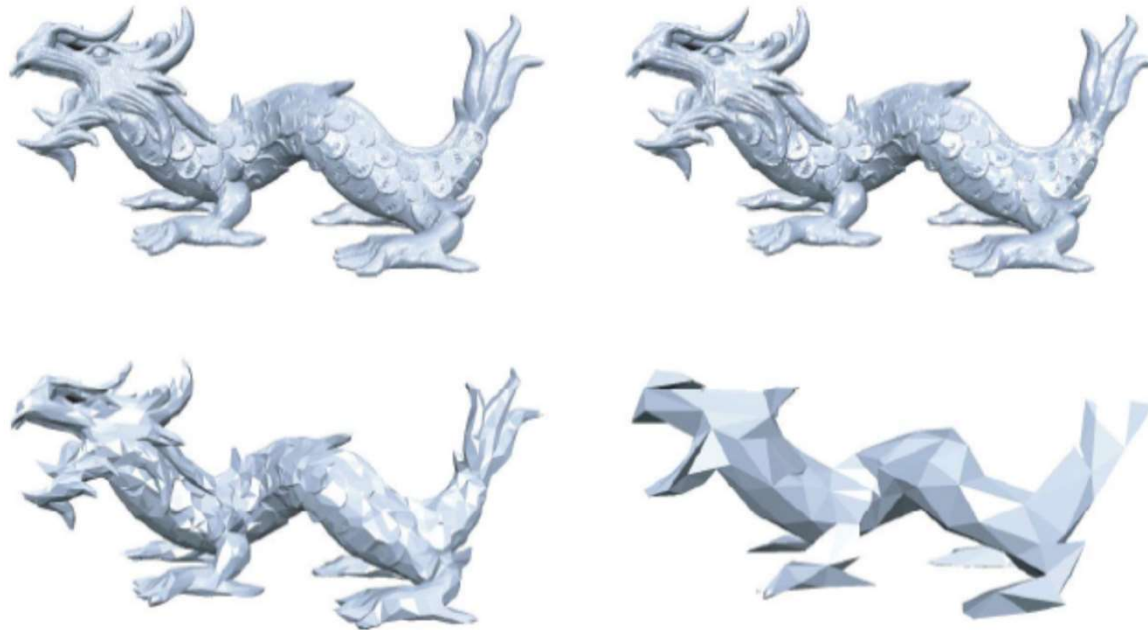


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Definition

Simplification, aka decimation, approximation, downsampling

- Transform a given polygonal mesh into another mesh with fewer faces, edges, and vertices
- The simplification or approximation procedure is usually controlled by user-defined quality criteria

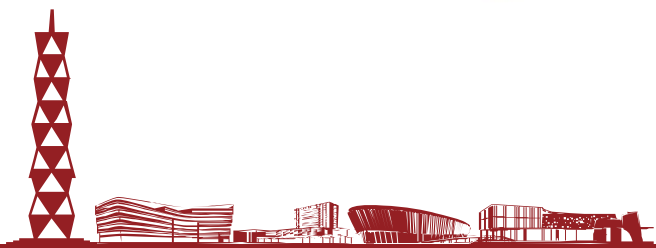
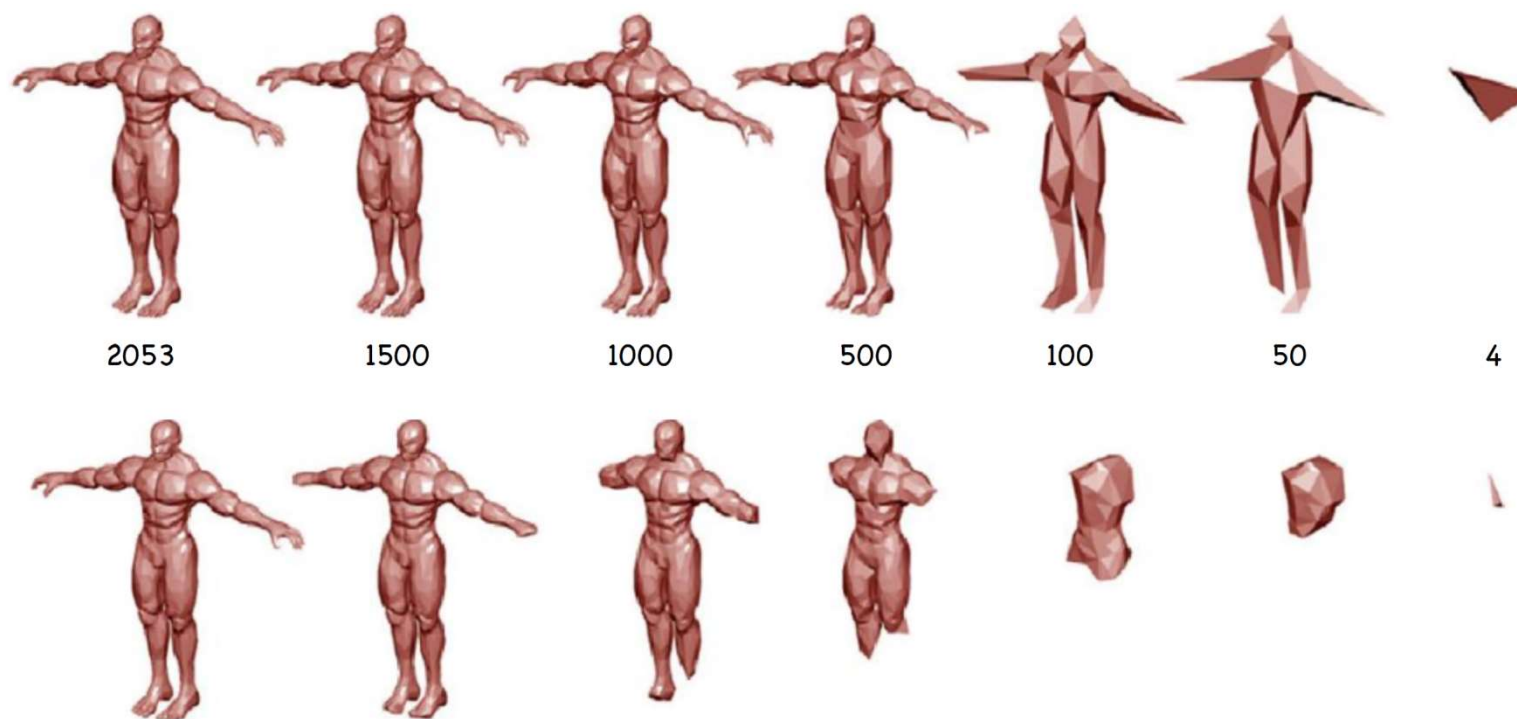




Curvature-preserved vs. Curvature-removed Criteria



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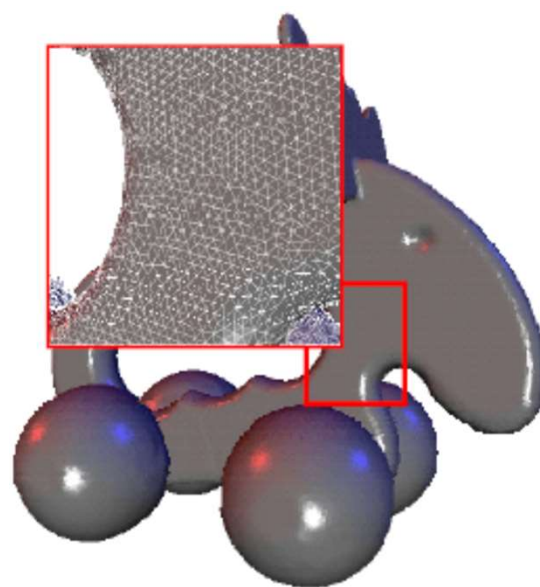


Mesh Simplification Applications

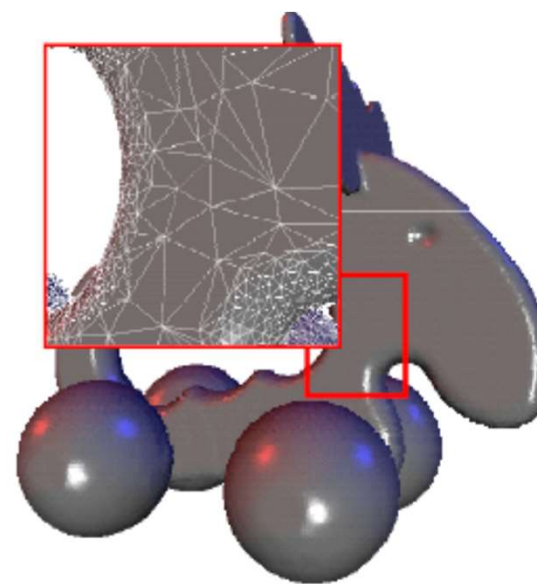


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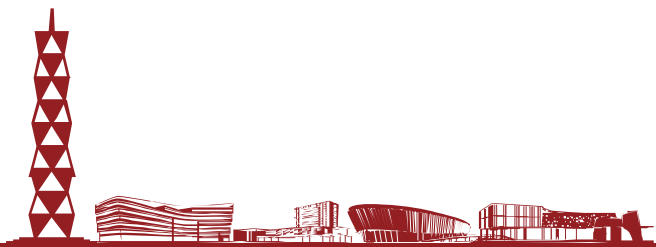
- Oversampled 3D scan data



~150k triangles



~80k triangles



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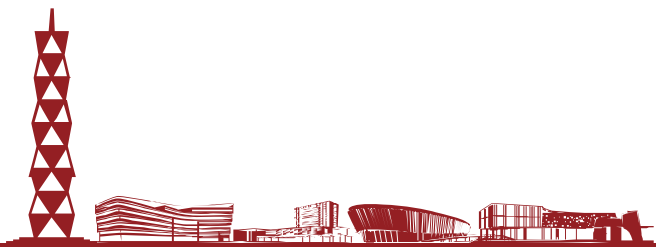
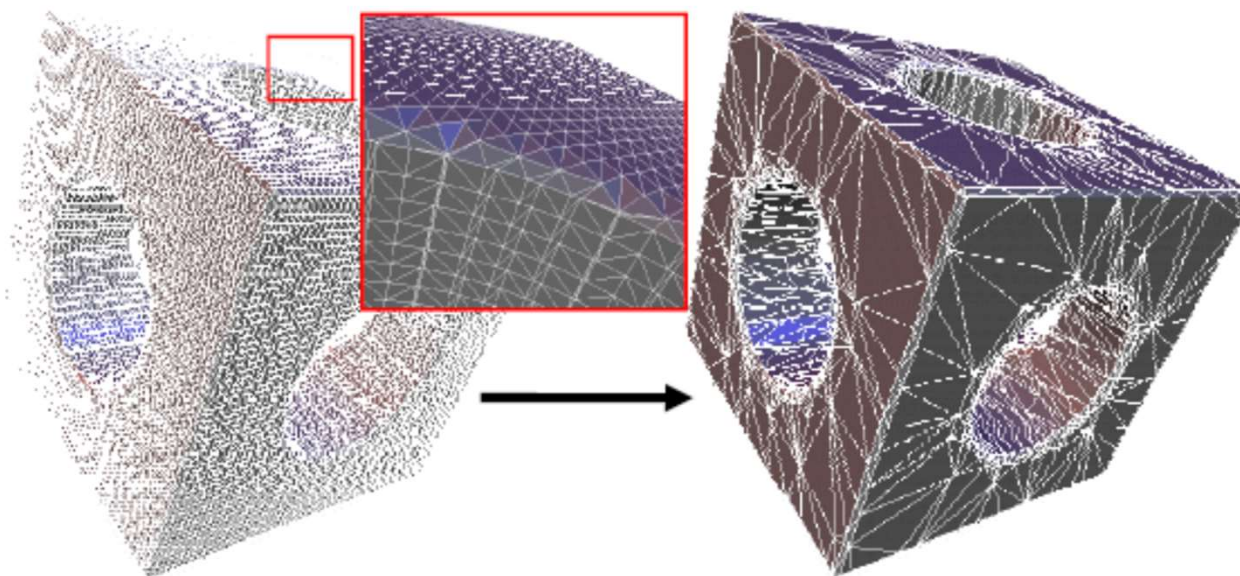


Mesh Simplification Applications



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- Over-tessellation: e.g., iso-surface extraction



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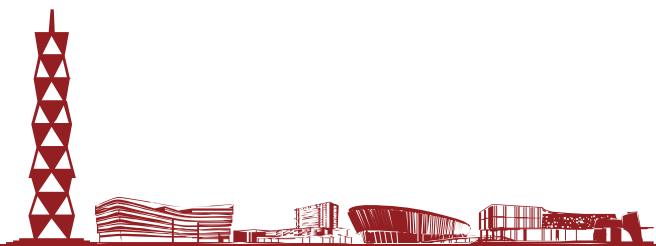
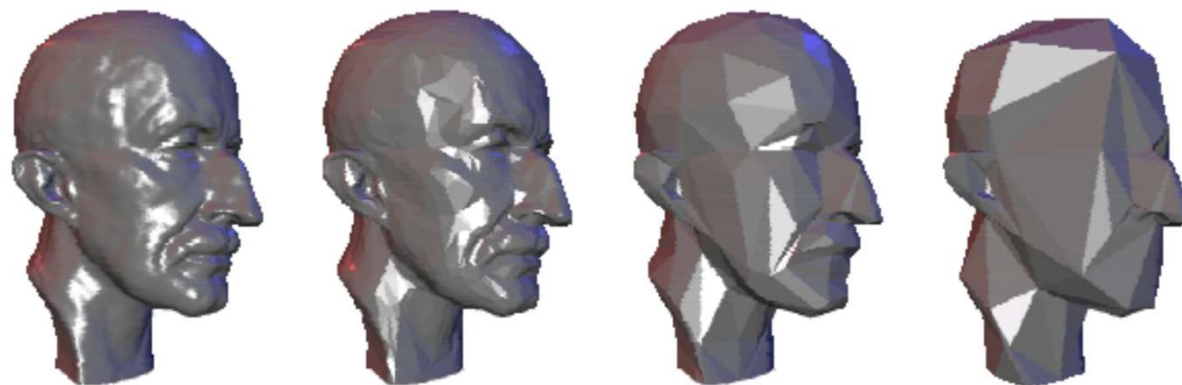
Mesh Simplification Applications



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Multi-resolution hierarchies for

- Efficient geometry processing
- Level-of-detail (LOD) rendering



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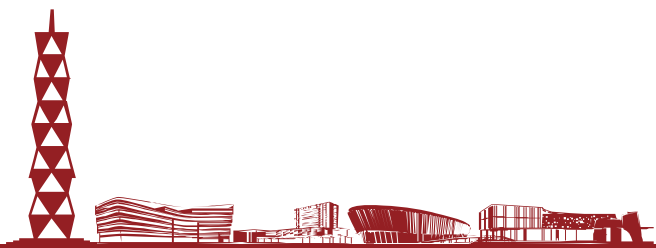


Mesh Simplification Applications



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- Adaption to hardware capabilities



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Mesh Simplification



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Adjust the complexity of a geometry data set

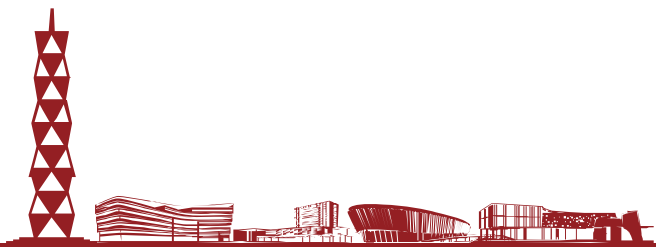
- Since many decimation schemes work iteratively, i.e., they decimate a mesh by removing one vertex at a time, they usually can be inverted
- Hierarchical method

Problem Statement

- Given: $M = (V, F)$
- Find: $M' = (V', F')$ such that
- $|V'| = n < |V|$ and $\|M - M'\|$ is minimal, or
- $\|M - M'\| < \varepsilon$ and $|V'|$ is minimal

Respect additional fairness criteria

- Normal deviation, triangle shape, scalar attributes, etc.



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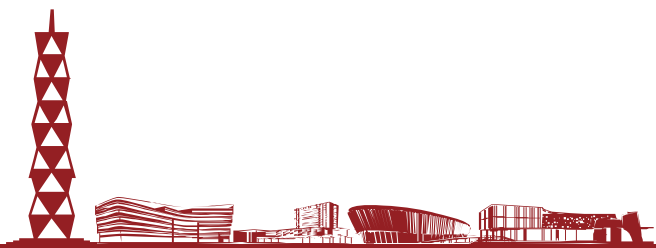
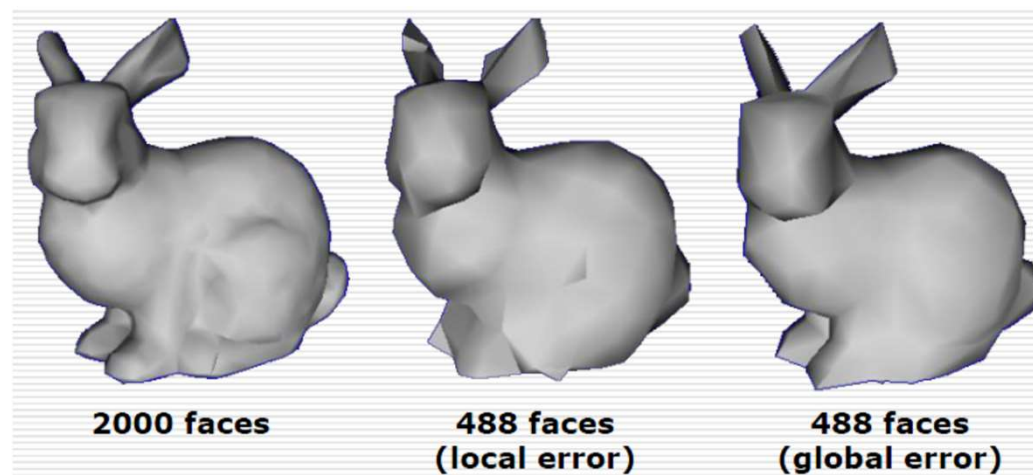
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Start with the original fine mesh

- **Simply progressively**
E.g., collapse edges, vertex clustering
- **Aim to keep original appearance**
Normal deviation, triangle shape, scalar attributes
Error control

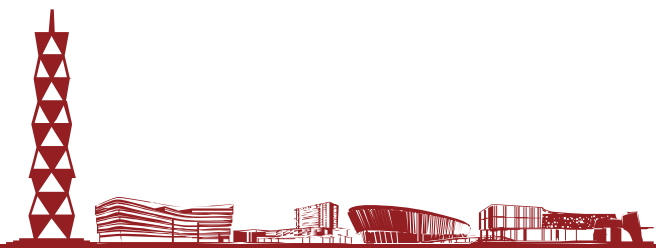


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Local Operations

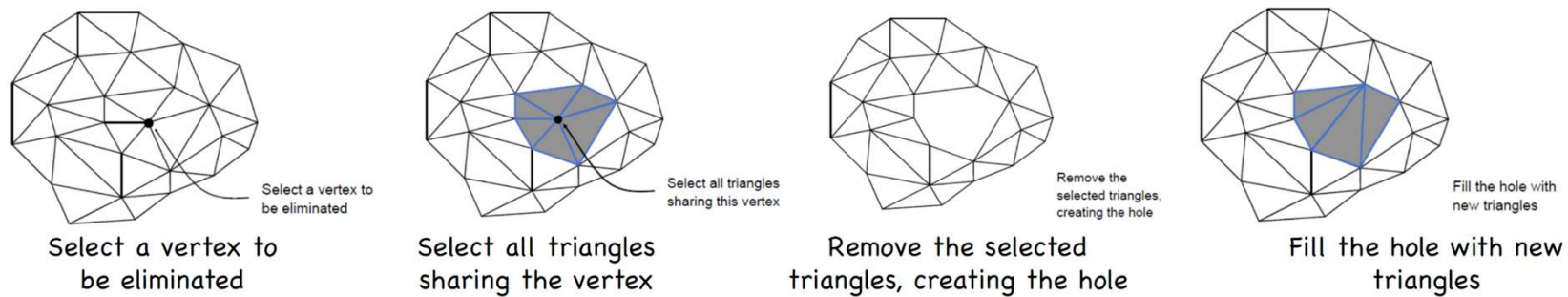
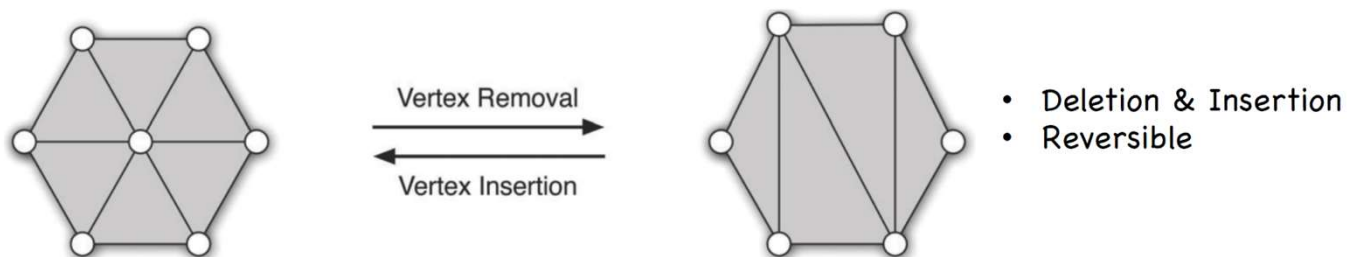


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Local Simplification Operator

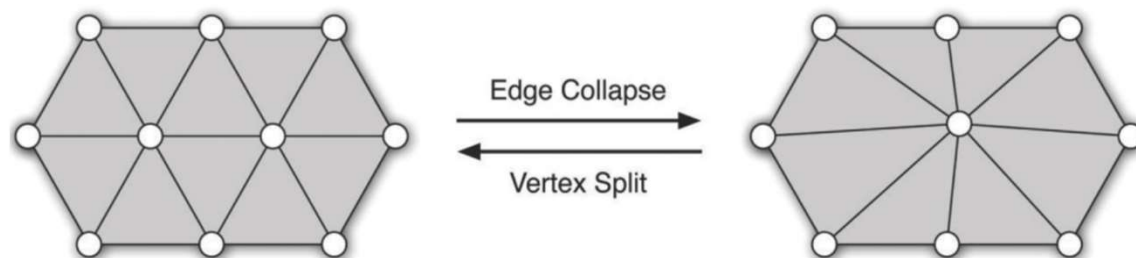
- Vertex Removal



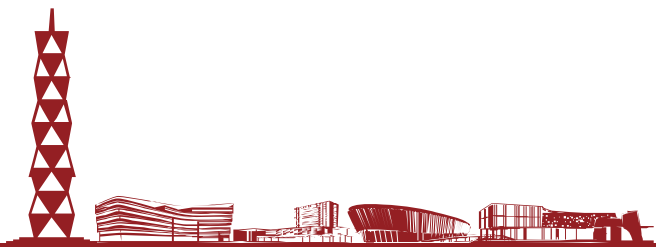


Local Simplification Operator

- Edge Collapse



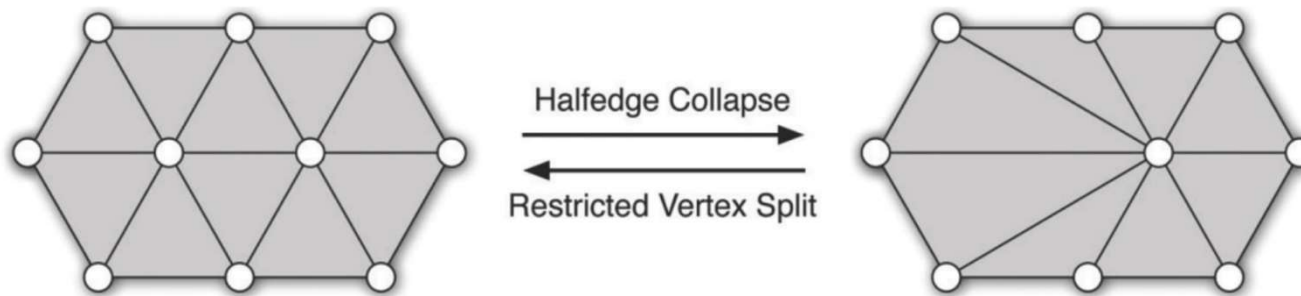
- Merge two adjacent vertices
- Simple to implement
- Well-suited for implementing geomorphing





Local Simplification Operator

- Half-edge Collapse

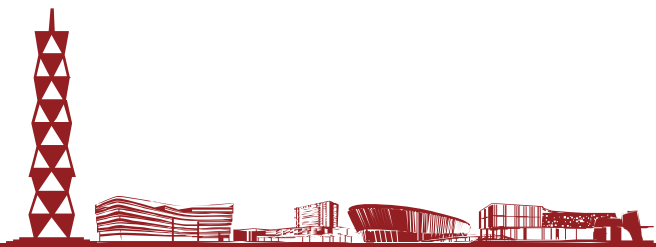


- Collapse edge into one end point

Special case of vertex removal

Special case of edge collapse

- After collapse: $n(E) - 3$, $n(V) - 1$, $n(F) - 2$
- According to Euler Formula: unchanged
- Half-edge collapsing would not change the genus of a mesh
- Should determine whether collapse is ok (may introduce non-manifold structure)

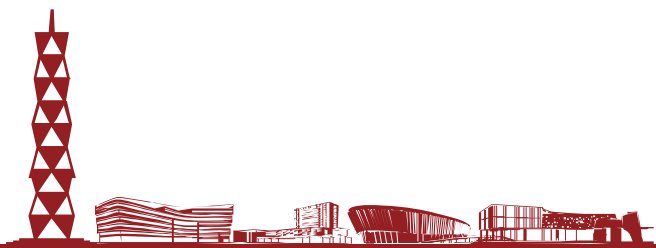
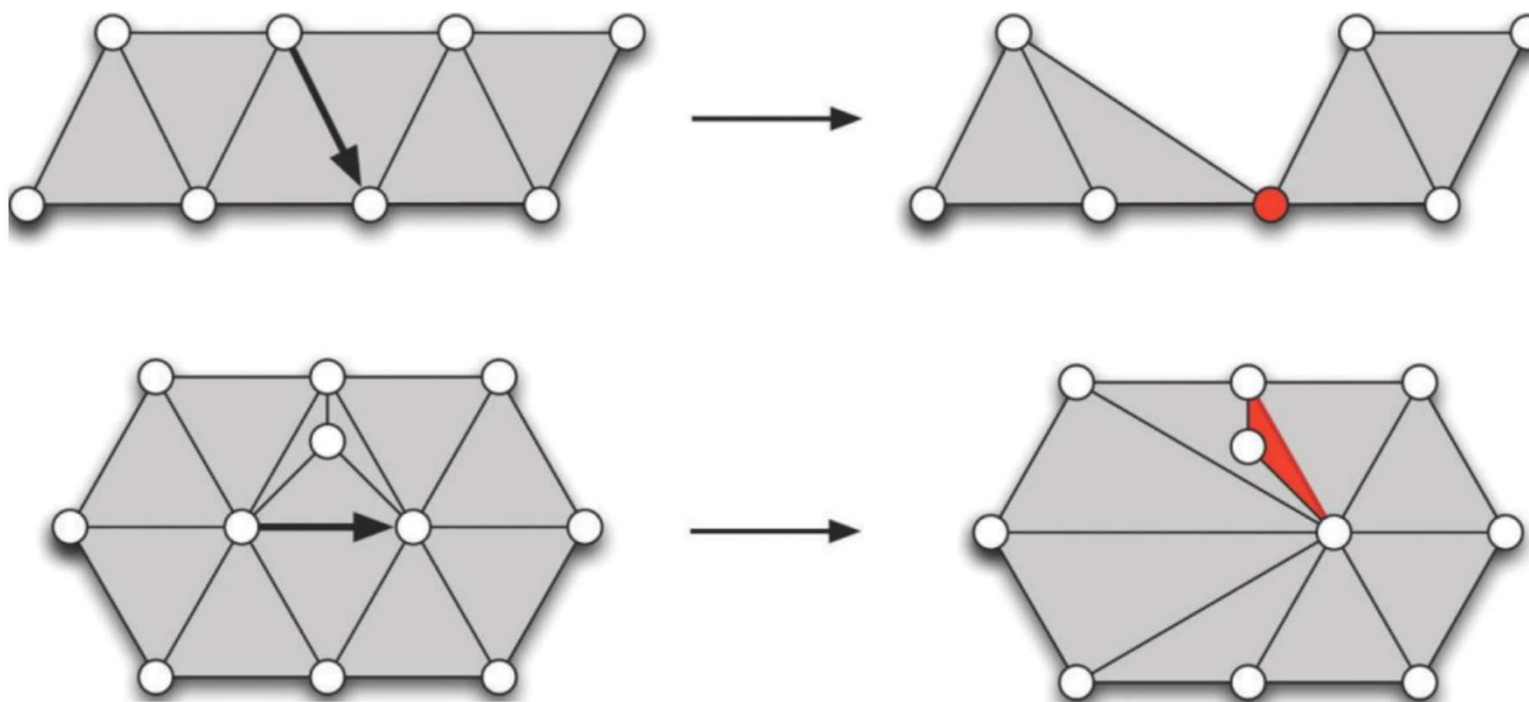




Topologically Illegal (half-)edge Collapses



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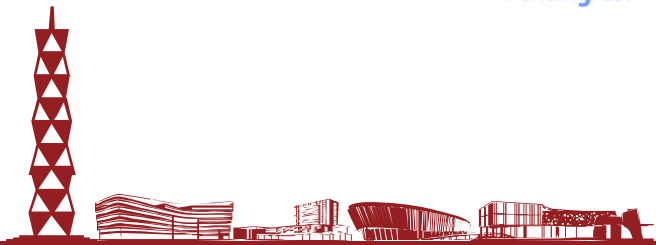


Simplification via Edge Collapse

One popular scheme: iteratively collapse edges

Greedy algorithm from a general overview:

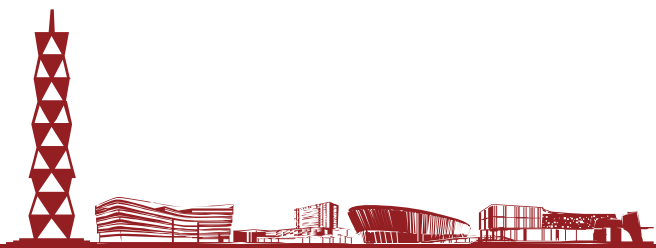
- Assign each edge a cost
- Collapse edge with least cost
- Repeat until target number of elements is reached
- Particularly effective cost function: **Quadric Error Metric**





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Quadric Error Metric

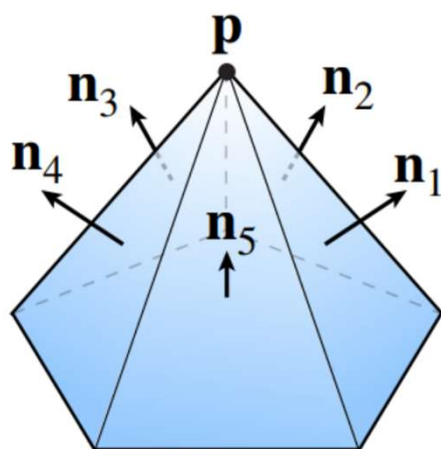
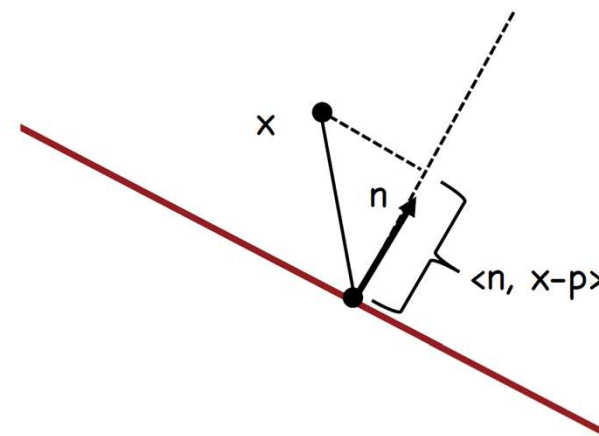


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Quadric Error Metric (QEM)

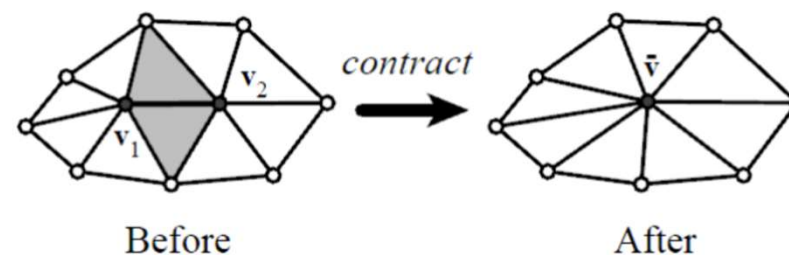
Approximate distance to a collection of triangles

- Q: Distance to plane w/ normal n passing through point p ?
- A: $\text{dist}(x) = \langle n, x \rangle - \langle n, p \rangle = \langle n, x - p \rangle$
- Quadric error is then sum of squared point-to-plane distances



$$Q(x) := \sum_{i=1}^k \langle n_i, x - p \rangle^2$$

$$Q^e = Q_1^v + Q_2^v$$





Quadric Error – Homogeneous Coordinates

Suppose in coordinates we have

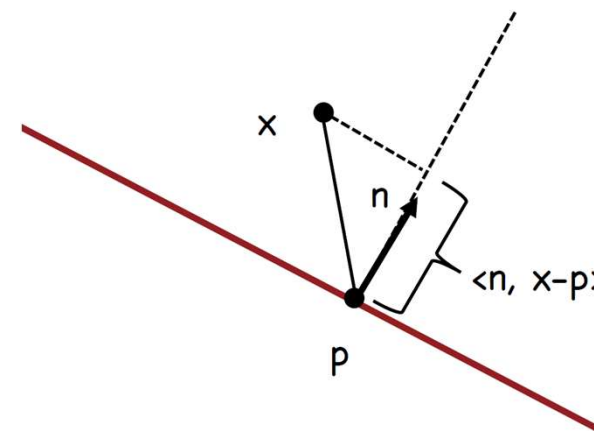
- A query point $\mathbf{x} = (x, y, z)$
- A normal $\mathbf{n} = (a, b, c)$
- An offset $\mathbf{d} := \langle \mathbf{n}, \mathbf{p} \rangle$

In **homogeneous coordinates**, let

- $\mathbf{u} := (x, y, z, 1)$
- $\mathbf{v} := (a, b, c, d)$
- Signed distance to plane is then just $\langle \mathbf{u}, \mathbf{v} \rangle = ax + by + cz + d$
- Squared distance is $\langle \mathbf{u}, \mathbf{v} \rangle^2 = \mathbf{u}^T (\mathbf{v} \mathbf{v}^T) \mathbf{u} =: \mathbf{u}^T \mathbf{K} \mathbf{u}$
- Matrix $\mathbf{K} = \mathbf{v} \mathbf{v}^T$ encodes squared distance to plane

Key idea: sum of matrices \mathbf{K} distance to union of planes

$$\mathbf{u}^T \mathbf{K}_1 \mathbf{u} + \mathbf{u}^T \mathbf{K}_2 \mathbf{u} = \mathbf{u}^T (\mathbf{K}_1 + \mathbf{K}_2) \mathbf{u}$$



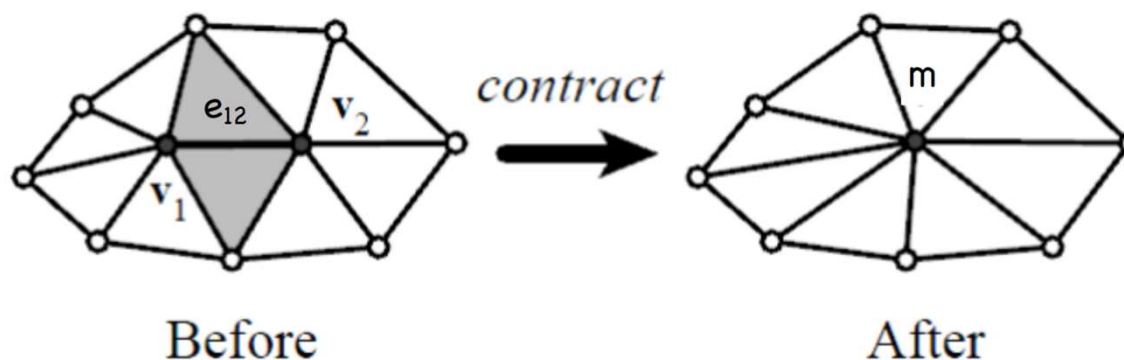
$$\mathbf{K} = \begin{bmatrix} a^2 & ab & ac & ad \\ ab & b^2 & bc & bd \\ ac & bc & c^2 & cd \\ ad & bd & cd & d^2 \end{bmatrix}$$





Quadric Error of Edge Collapse

- How much does it cost to collapse an edge e_{12} ?
Idea: compute midpoint m , measure error $Q(m) = m^T(K_1 + K_2)m$
- Error becomes “score” for e_{12} , determining priority



- Better idea: find point x that minimize error!
- But how to minimize quadric error?



Revisit: Minimizing a Quadratic Function

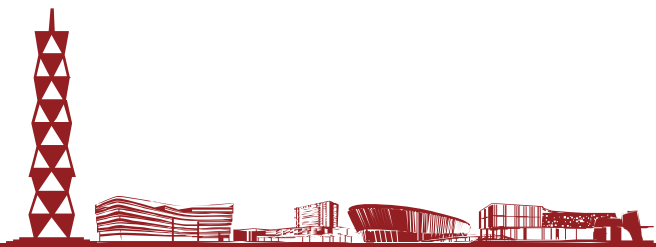
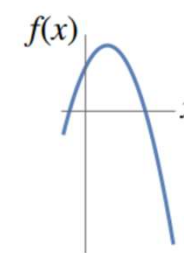
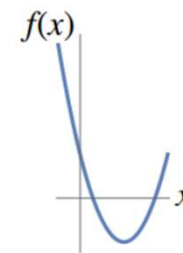


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Suppose you have a function $f(x) = ax^2 + bx + c$

- Q: What does the graph of this function look like?
- Q: How do we find the minimum?
- A: Find where the function looks “flat” if we zoom in really close
- i.e., find point x where 1st derivative vanishes:

$$\begin{aligned}f'(x) &= 0 \\2ax + b &= 0 \\x &= -b/2a\end{aligned}$$





Minimizing Quadratic Polynomial

Not much harder to minimize a quadratic polynomial in n variables

- Can always write in terms of a symmetric matrix A
- E.g., in 2D: $f(x, y) = ax^2 + bxy + cy^2 + dx + ey + g$

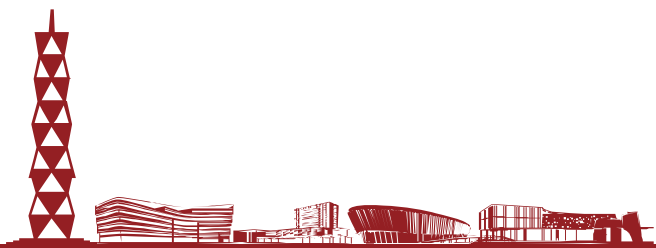
$$\mathbf{x} = \begin{bmatrix} x \\ y \end{bmatrix} \quad A = \begin{bmatrix} a & b/2 \\ b/2 & c \end{bmatrix} \quad \mathbf{u} = \begin{bmatrix} d \\ e \end{bmatrix}$$

$$f(\mathbf{x}, y) = \mathbf{x}^T A \mathbf{x} + \mathbf{u}^T \mathbf{x} + g$$

(will have the same form for any n)

- Q: How do we find a critical point (min/max/saddle)?
- A: Set derivative to zero!

$$2A\mathbf{x} + \mathbf{u} = 0$$
$$\mathbf{x} = -1/2 A^{-1}\mathbf{u}$$





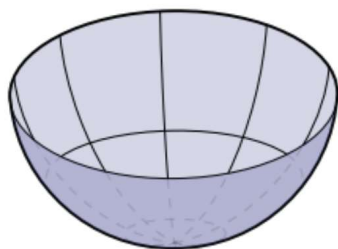
Positive Definite Quadratic Form

Just like our 1D parabola, critical point is not always a min!

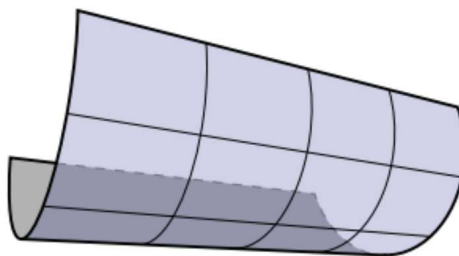
- Q: In 2D, 3D, nD, when do we get a minimum?
- A: When matrix A is positive-definite:

$$\mathbf{X}^T \mathbf{A} \mathbf{X} > 0 \quad \forall \mathbf{X}$$

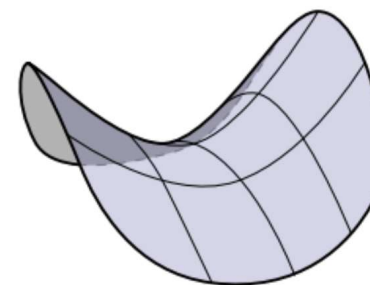
- **1D:** Must have $xax = ax^2 > 0$, i.e., a is positive!
- **2D:** Graph of function looks like a “bowl”:



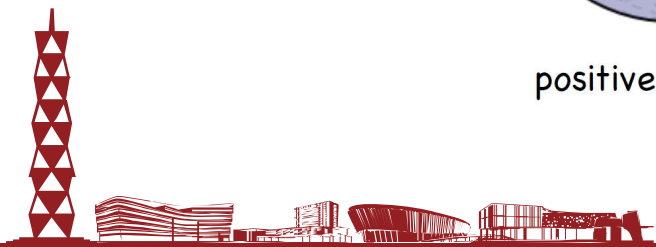
positive definite



positive semidefinite



indefinite





Minimizing Quadric Error



Find “best” point for edge collapse by minimizing quadratic form

$$\min_{\mathbf{u} \in \mathbb{R}^4} \mathbf{u}^T \mathbf{K} \mathbf{u}$$

- Already know 4th (homogeneous) coordinate for a point is 1
- So, break up our quadratic function into two pieces:

$$\begin{bmatrix} \mathbf{x}^T & 1 \end{bmatrix} \begin{bmatrix} B & \mathbf{w} \\ \mathbf{w}^T & d^2 \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ 1 \end{bmatrix} \\ = \mathbf{x}^T B \mathbf{x} + 2\mathbf{w}^T \mathbf{x} + d^2$$

- Now we have a quadratic polynomial in the unknown position $\mathbf{x} \in \mathbb{R}^3$
- Can minimize as before:

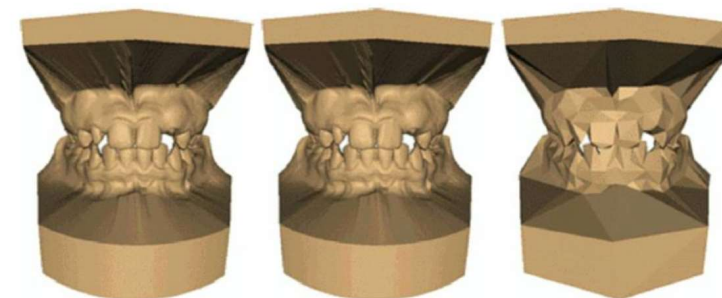
$$2B\mathbf{x} + 2\mathbf{w} = 0 \quad \Longleftrightarrow \quad \mathbf{x} = -B^{-1}\mathbf{w}$$





QEM Simplification: Final Algorithm

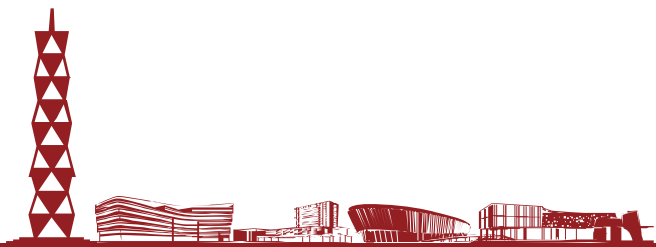
- Input: a mesh
- Output: a simplified mesh
- Initialization:
 - Compute \mathbf{K} for each triangle (squared distance to plane)
 - Set \mathbf{K}_i at each vertex to sum of \mathbf{K} s from incident triangles
 - For each edge \mathbf{e}_{ij} :
 - Set $\mathbf{K}_{ij} = \mathbf{K}_i + \mathbf{K}_j$
 - Find point x minimizing error, set cost to $\mathbf{K}_{ij}(x)$
- Until we reach target number of triangles:
 - Collapse edge \mathbf{e}_{ij} with smallest cost to optimal point x
 - Set quadric at new vertex to \mathbf{K}_{ij}
 - Update cost of edges touching new vertex



Full
Resolution

60,000
triangles

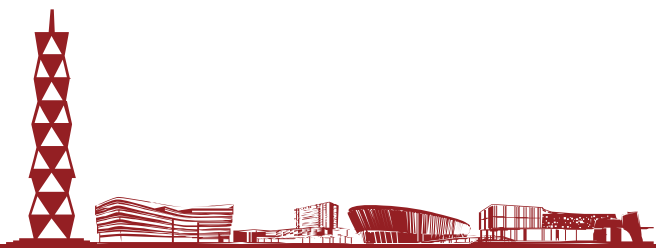
1000
triangles





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Variational Shape Approximation



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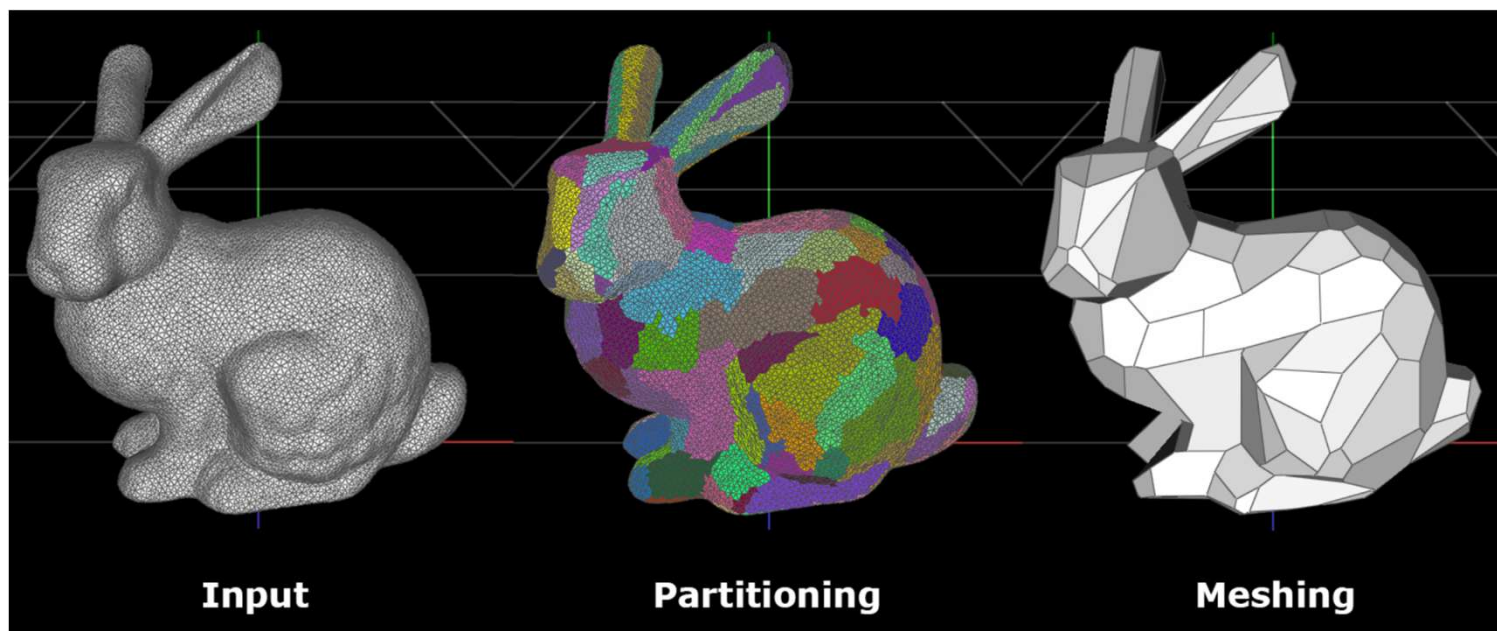


Variational Shape Approximation (VSA)



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VSA is highly sensitive to features and symmetries and produces anisotropic meshes of high approximation quality

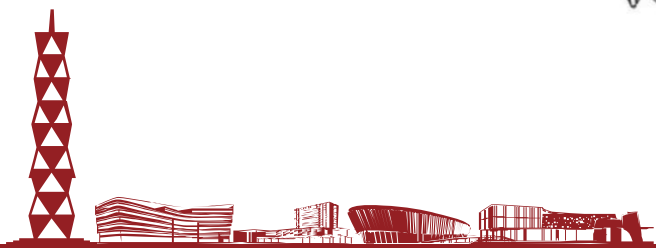
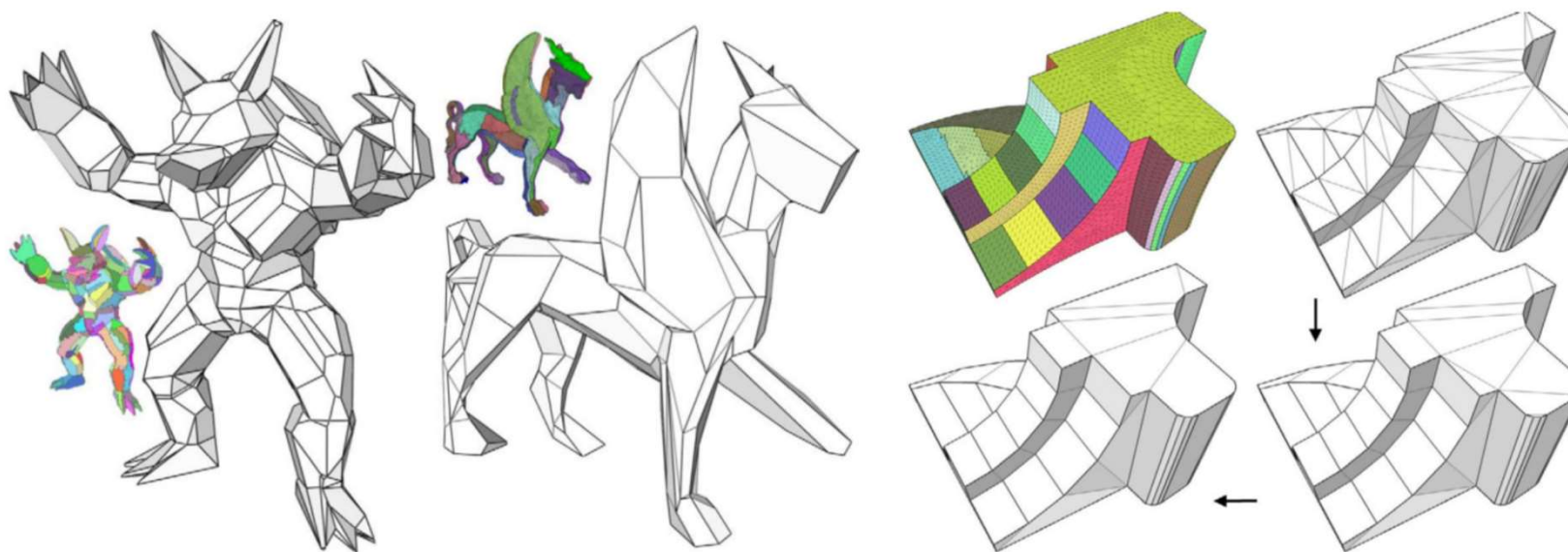


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Variational Shape Approximation (VSA)

- The input shape is approximated by a set of proxies
- A plane in space through the point x_i with normal direction n_i





Region Representation

- \mathbf{M} : a triangle mesh
- $\mathbf{R} = \{\mathbf{R}_1, \dots, \mathbf{R}_k\}$: a partition of \mathbf{M} into k regions $R_1 \cup \dots \cup R_k = M$
- Proxies: $\mathbf{P} = \{\mathbf{P}_1, \dots, \mathbf{P}_k\}$, $\mathbf{P}_i = (\mathbf{x}_i, \mathbf{n}_i)$

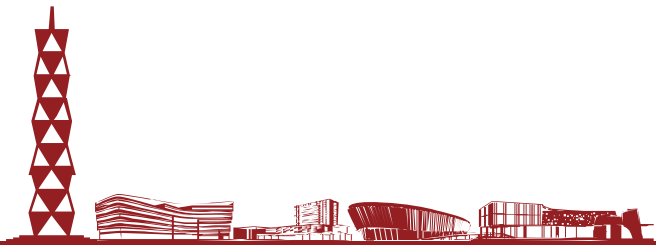
Distance metrics between \mathbf{R}_i and \mathbf{P}_i

- The squared orthogonal distance of x from the plane \mathbf{P}_i

$$L^2(R_i, P_i) = \int_{x \in R_i} (n_i^T x - n_i x_i)^2 dA$$

- A measure of the normal field:

$$L^{2,1}(R_i, P_i) = \int_{x \in R_i} \|n(x) - n_i\|^2 dA$$





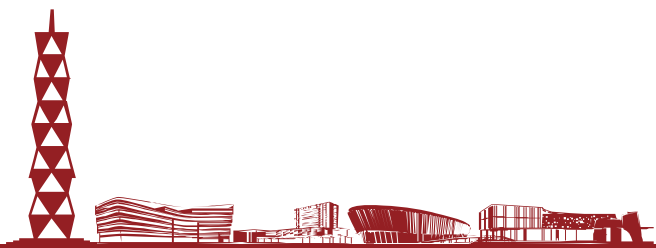
Goal of VSA



Given a number k and an error metric $E(L_2 \text{ or } L_{2,1})$, find a set $\mathbf{R} = \{\mathbf{R}_1, \dots, \mathbf{R}_k\}$ of regions and a set $\mathbf{P} = \{\mathbf{P}_1, \dots, \mathbf{P}_k\}$ of proxies such that the global distortion

$$E(R, P) = \sum_{i=1}^k E(R_i, P_i)$$

is minimized

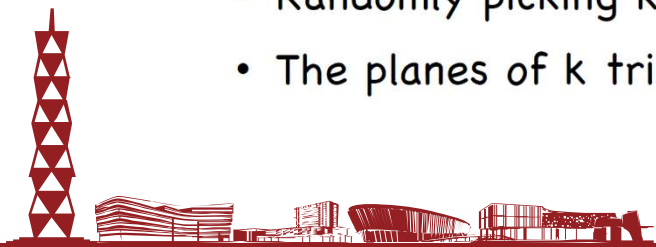




Lloyd's Clustering Algorithm



- The algorithm iteratively alternates between a **geometry partitioning** phase and a **proxy fitting** phase
- **Geometry partitioning phase**
 - A set of regions that best fit a **given** set of proxies
 - Modifies the set R of regions to achieve a lower approximation error while keeping the proxies P fixed
- **Proxy fitting phase**
 - The partitioning is kept **fixed**, and the proxies are adjusted to minimize approximation error
 - L^2 metric: the best proxy is the least-squares fitting plane
 - $L^{2,1}$ metric: the proxy normal n_i is just the area-weighted average of the triangle normals
- **Initialization**
 - Randomly picking k triangles as R
 - The planes of k triangles are used to initialize P





More Paper

- Liu Y J, Xu C X, Fan D, et al. **Efficient construction and simplification of Delaunay meshes**[J]. ACM Transactions on Graphics (TOG), 2015, 34(6): 1-13.
- Yi R, Liu Y J, He Y. **Delaunay mesh simplification with differential evolution**[J]. ACM Transactions on Graphics (TOG), 2018, 37(6): 1-12.
- Liang Y, He F, Zeng X. **3D mesh simplification with feature preservation based on whale optimization algorithm and differential evolution**[J]. Integrated Computer-Aided Engineering, 2020, 27(4): 417-435.

