

Problem 1

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$$f(n) = (n^2 + 1)(\log n + 1) + (n^2 \log n + 1)^2$$

$$= n^2 \log n + n^2 + \log n + 1 + n^4 (\log n)^2 + 2n^2 \log n + 1$$

$$= 3n^2 \log n + n^2 + \log n + n^4 (\log n)^2 + 2$$

$$\lim_{n \rightarrow \infty} \frac{3n^2 \log n + n^2 + \log n + n^4 (\log n)^2 + 2}{n^5}$$

$$= \lim_{n \rightarrow \infty} \frac{3 \log n}{n^3} + \frac{1}{n^3} + \frac{\log n}{n^5} + \frac{(\log n)^2}{n} + \frac{2}{n^5}$$

~~$\frac{0}{\infty} + \frac{0}{\infty} + \frac{0}{\infty} + \frac{0}{\infty} + \frac{0}{\infty}$~~ since when $n \geq 1$, $\frac{3 \log n}{n^3}$, $\frac{\log n}{n^5}$, $\frac{(\log n)^2}{n^5}$ are all less than 1

$$\text{and } \lim_{n \rightarrow \infty} \frac{1}{n^3} + \frac{2}{n^5} = 0$$

$$\text{So } \lim_{n \rightarrow \infty} \frac{f(n)}{n^5} = \frac{3 \log n}{n^3} + \frac{1}{n^3} + \frac{\log n}{n^5} + \frac{(\log n)^2}{n} + \frac{2}{n^5} < 3 + 1 + 1 + 0 = 5$$

Thus, there is an upper bound that is $\in \mathbb{R}_{>0}$

So $t=5$, $f(n) = O(n^5)$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{n^4} = \frac{3 \log n}{n^2} + \frac{1}{n^2} + \frac{\log n}{n^4} + (\log n)^2 + \frac{2}{n^4}$$

since $(\log n)^2 = \infty$

$\lim_{n \rightarrow \infty} \frac{f(n)}{n^4} = \infty$ since there is no upper bound
it should be 5 and not be 4

Problem 2

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If $f(n) = n^3$, $g(n) = 2n^3$, which ~~fulfills that~~ $f(n) = \Omega(g(n))$

$$2^{f(n)} = 2^{n^3}$$

$$2^{g(n)} = 2^{2n^3}$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \frac{n^2}{2n^2} = \frac{1}{2} > \frac{1}{2}, \text{ so } f(n) = \Omega(g(n))$$

$$\lim_{n \rightarrow \infty} \frac{2^{f(n)}}{2^{g(n)}} = \lim_{n \rightarrow \infty} 2^{f(n)-g(n)} = \lim_{n \rightarrow \infty} 2^{-n^3} = 0$$

$$\text{So } 2^{f(n)} \neq \Omega(2^{g(n)})$$

Problem 3

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since we have $f(n) = O(g(n))$

$$g(n) = O(f(n))$$

and given $f: \mathbb{N} \rightarrow \mathbb{R}_{>0}$ and $g: \mathbb{N} \rightarrow \mathbb{R}_{>0}$, we have

$$(\exists d \in \mathbb{R}_{>0}) (\exists N \in \mathbb{N}) (\forall n \geq N) \frac{f(n)}{g(n)} \leq d$$

and

$$(\exists c \in \mathbb{R}_{>0}) (\exists N \in \mathbb{N}) (\forall n \geq N) \frac{g(n)}{f(n)} \leq c$$

so we get

$$f(n) \leq d \cdot g(n)$$

$$g(n) \leq c \cdot f(n)$$

so

$$(\exists c, d \in \mathbb{R}_{>0}) (\exists N \in \mathbb{N}) (\forall n \geq N) \frac{c \cdot f(n)}{f(n)} \leq \frac{f(n)}{g(n)} \leq \frac{d \cdot g(n)}{g(n)}$$

$$\text{which is } c \leq \frac{f(n)}{g(n)} \leq d$$

so, by definition $f(n) = \Theta(g(n))$