

(a) Proof.

Claim: The following are loop invariants for `indexOfMax` algorithm before the while test in line (4).

i. $0 \leq i \leq n$ ii. $m = A[b]$ iii. $m = \max A[0..(i-1)]$ iv. b is the largest index at which $\max A[0..(i-1)]$ appears

$P(n)$: Loop invariants (i), (ii), (iii), (iv) hold when the loop condition is tested for the n -th time.

Base case: show $P(0)$ holds, i.e., loop invariants hold when $i \leq n$ is tested for the 0 ~~first~~ time.

At this point, $i=1, m=A[0], b=0$

since ~~$n > 0$~~ , $0 \leq i \leq n$ (n is an integer)

since $b=0$, (ii) $m = A[b] = A[0]$

since there is only one index and element in A ,

(iii) $\max A[0..(i-1)] = \max A[0] = A[0]$

and b is the largest index at which $\max A[0..(i-1)]$ appears. \square

Inductive step:

Induction hypothesis: The loop invariants hold k -th time the loop is tested. Show the loop invariants hold the $(k+1)$ -st time the loop is tested. i.e. $P(k)$

Let i_k, m_k and b_k be the values of i, m and b , respectively, after k iterations of the while loop.

IH: $0 \leq i_k \leq n, m_k = A[b_k], m_k = \max A[0..(i_k-1)]$
 b is the largest index at which $\max A[0..(i_k-1)]$ appears

Consider the $(k+1)$ st time the loop body executes. To get into the body of the loop, the while condition must be true, so we have $i_k \leq n$



Case 1: line (5) $A[i_k] \geq m_k$ is true

From line (6), we get $m_{k+1} = A[i_k]$

From line (7), we get $b_{k+1} = i_k$

From line (9), we get $i_{k+1} = i_k + 1$

Since $i_{k+1} = i_k + 1$ (~~we don't change i in this case~~)

and $0 < i_k \leq n$ (by IH) and $i_k < n$

~~(i) $0 < i_k \leq n$~~ $0 < i_k < n \Rightarrow 0 < i_k + 1 < n + 1 \Rightarrow 0 < i_{k+1} < n + 1$ (h.e.z⁺)

so (i) $0 < i_{k+1} \leq n$ holds

since $m_{k+1} = A[i_k]$ and $b_{k+1} = i_k$

(ii) $m_{k+1} = A[b_{k+1}]$ holds

Since $m_k = \max A[0..(i_k-1)]$ (by IH), $A[i_k] \geq m_k$ and $b_{k+1} = i_k$

$$\therefore m_{k+1} = A[b_{k+1}]$$

$$= A[i_k]$$

$$\geq \max A[0..(i_k-1)]$$

since $i_k = i_{k+1} - 1$

$$m_{k+1} \geq \max A[0..(i_{k+1}-1)]$$

so, m_{k+1} is the max element in $A[0..(i_{k+1}-1)]$

\therefore (iii) $m_{k+1} = \max A[0..(i_{k+1}-1)]$ holds

$$\therefore m_{k+1} = A[b_{k+1}]$$

\therefore (iv) b_{k+1} is the largest index at which $\max A[0..(i_{k+1}-1)]$ appears



Case 2: line (5) $A[i_k] \geq m_k$ is false

$$\text{so } A[i_k] < m_k$$

From line (9), we get $i_{k+1} = i_k + 1$

Since $i_k < n$

$$\text{by IH, } 0 < i_k < n$$

$$1 < i_k + 1 < n + 1$$

since i is an integer

$$0 < 1 < i_{k+1} < n + 1$$

(i) $0 < i_{k+1} \leq n$ holds

$$\text{by IH, } m_k = A[b_k]$$

$$\text{since } b_{k+1} = b_k, m_{k+1} = m_k$$

(ii) $m_{k+1} = A[b_{k+1}]$ holds

since $A[i_k] < m_k$, $m_k = \max A[0..(i_k-1)]$ (by IH)

~~$$\text{since } m_k = m_{k+1}, i_{k+1} = i_k + 1$$~~

~~$$m_{k+1} = \max A[0..(i_{k+1}-1)]$$~~

~~$$m_k = \max A[0..i_k] \Rightarrow (m_k > A[i_k])$$~~

~~$$m_k > A[i_k]$$~~

$$\max A[0..(i_k-1)] > A[i_k]$$

$$\therefore m_k = \max A[0..i_k]$$

$$\text{since } m_k = m_{k+1}, i_{k+1} = i_k + 1 \Rightarrow i_k = i_{k+1} - 1$$

(iii) $m_{k+1} = \max A[0..(i_{k+1}-1)]$ holds

$$\therefore m_{k+1} = A[b_{k+1}]$$

$$\therefore A[b_{k+1}] = \max A[0..(i_{k+1}-1)]$$

\therefore (iv) b_{k+1} is the largest index at which $\max A[0..(i_{k+1}-1)]$ appears



(6)

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Partial correctness:

Assume we have valid input and the program terminates on that ~~input~~ input.
Since the program returns b , we need to show that when the loop exits,
 b is the index of largest element ~~to~~ appear.

when the loop exits, it does so because $i < n$ evaluates to false,

This means $i \geq n$.

Loop invariant (i) says $0 < i \leq n$

So when loop exits $i = n$

Loop invariant (ii) says $m = \max A[0..(i-1)]$

so when loop exits, m is the largest element in A

Loop invariant (iv) says b is the largest index at which $\max A[0..(i-1)]$ appears,
and $m = A[b]$

so, The program returns b , which contains the correct value. \square

Termination:

Assume the input is valid, i.e. $n \in \mathbb{Z}^+$ and an array $A[0..(n-1)]$ of n integers
The only reason the program wouldn't terminate is if we had an infinite loop.

Loop exits when $i \geq n$ (when tested on line 4)

i starts as 1 (line 1)

i increase by 1 (line 9) in each iteration of the loop when if statement is false.
 n doesn't change

since n is finite and unchanging at some point $i \geq n$ becomes true.

so loop cannot be infinite

so, the implementation is correct. \square



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