# Strategy for developing efficient programs:

- 1. Design the program well
- 2. Implement the program well\*\*
- 3. Test the program well
- 4. Only after you're sure it's working, measure performance
- 5. If (and only if) performance is inadequate, find the "hot spots"
- 6. Tune the code to fix these
- 7. Repeat measure-analyse-tune cycle until performance ok

(\*\* see "Programming Pearls", "Practice of Programming", etc. etc.) Rapid development of a prototype may be the best way to discover/assess performance issues.

Hence Fred Brooks maxim - "Plan To Throw One Away".

#### Two C Fucntions which Initialialize an array

```
void test0(int x, int y, int a[x][y]) {
    fprintf(stderr, "writing to array i-j order\n");
    for (int i = 0; i < x; i++)
        for (int j = 0; j < y; j++)
            a[i][j] = i+j;
}</pre>
```

```
void test1(int x, int y, int a[x][y]) {
   fprintf(stderr, "writing to array j-i order\n");
   for (int j = 0; j < y; j++)
        for (int i = 0; i < x; i++)
        a[i][j] = i+j;
}</pre>
```

Although the loops are almost identical, the first loop runs 20x faster on a large array!

```
$ time ./cachegrind_example 0 32000 32000
allocating a 32000x32000 \text{ array} = 4096000000 \text{ bytes}
writing to array i-j order
real 0m0.893s
user 0m0.364s
sys 0m0.524s
$ time ./cachegrind example 1 32000 32000
allocating a 32000x32000 \text{ array} = 4096000000 \text{ bytes}
writing to array j-i order
real 0m15.189s
user 0m14.633s
sys 0m0.528s
```

The tool valgrind used to detect accesses to uninitialized variables at runtime also can give memory caching infomation.

The memory subsystem is beyond the scope of this course but you can see valgrind explain the performance difference between these loops.

For the first loop D1 miss rate = 24.8%

For the second loop D1 miss rate = 99.9%

Due to the C array memory layout the first loop produces much better caching performance.

Tuning caching performance is important for some application and valgrind makes this much easier.

```
$ valgrind '--tool=cachegrind' ./cachegrind example 0 10000 10000
allocating a 10000x10000 array = 400000000 bytes
writing to array i-j order
==7025==
==7025== I refs: 225,642,966
==7025== I1 misses:
                           882
==7025== LLi misses:
                         875
==7025== I1 miss rate:
                          0.00%
==7025== LLi miss rate:
                          0.00%
==7025==
==7025== D refs:
                    25,156,289 (93,484 rd + 25,062,805 wr)
==7025== D1  misses: 6,262,957 (2,406 rd + 6,260,551 wr)
==7025== LLd misses: 6,252,482 (1,982 rd + 6,250,500 wr)
==7025== D1 miss rate:
                          24.8% ( 2.5%
                                                  24.9%
==7025== LLd miss rate:
                          24.8% ( 2.1%
                                                  24.9% )
==7025==
==7025== LL refs: 6,263,839 (3,288 rd + 6,260,551 wr)
==7025== LL misses: 6,253,357 (2,857 rd + 6,250,500 wr)
==7025== LL miss rate:
                           2.4% ( 0.0%
                                                  24.9% )
```

```
$ valgrind '--tool=cachegrind' ./cachegrind example 1 10000 10000
allocating a 10000x10000 array = 400000000 bytes
writing to array j-i order
==7006==
==7006== I refs: 600,262,960
==7006== I1 misses:
                            876
==7006== LLi misses:
                          869
==7006== I1 miss rate:
                           0.00%
==7006== LLi miss rate:
                           0.00%
==7006==
==7006== D \text{ refs}: 100,056,288 (43,483 rd + 100,012,805 wr)
==7006== D1 misses:
                     100,002,957 (2,405 rd + 100,000,552 wr)
==7006== LLd misses: 6,262,481 (1,982 rd
                                           + 6,260,499 wr)
==7006== D1 miss rate:
                           99.9% ( 5.5%
                                                    99.9%)
                                           +
==7006== LLd miss rate: 6.2% ( 4.5%
                                                    6.2%
==7006==
==7006== LL refs: 100,003,833 (3,281 rd
                                           + 100,000,552 wr)
==7006== LL misses: 6,263,350 (2,851 rd
                                               6,260,499 wr)
==7006== LL miss rate:
                            0.8% (
                                                    6.2\%
                                    0.0%
                                            +
```

#### Where is execution time being spent?

Typically programs spend most of their execution time in a small part of their code.

This is often quoted as the 90/10 rule (or 80/20 rule or ...):

"90% of the execution time is spent in 10% of the code"

#### This means that

- most of the code has little impact on overall performance
- small parts of the code account for most execution time

We should clearly concentrate efforts at improving execution spped in the 10% of code which accounts for most of the execution time.

# clang -p/gprof

Given the -p flag clang instruments a C program to collect profile information

When the program executes this data is left in the file gmon.out. The program gprof analyzes this data and produces:

- number of times each function was called
- % of total execution time spent in the function
- average execution time per call to that function
- execution time for this function and its children

Arranged in order from most expensive function down. It also gives a *call graph*, a list for each function:

- which functions called this function
- which functions were called by this function

# Program for producing sorted counts of words



Program is slow on large inputs e.g.

We can instrument the program to collect profiling information and examine it with clang

```
$ clang -p -g word_frequency0.c -o word_frequency0_profile
$ head -10000 WarAndPeace.txt|word_frequency0_profile >/dev/null
$ gprof word_frequency0_profile
```

% c	cumulative	self		self	total	
time	seconds	seconds	calls	ms/call	ms/call	name
88.90	0.79	0.79	88335	0.01	0.01	get
7.88	0.86	0.07	7531	0.01	0.01	put
2.25	0.88	0.02	80805	0.00	0.00	get_word
1.13	0.89	0.01	1	10.02	823.90	read_words
0.00	0.89	0.00	2	0.00	0.00	size
0.00	0.89	0.00	1	0.00	0.00	create_map
0.00	0.89	0.00	1	0.00	0.00	keys
0.00	0.89	0.00	1	0.00	0.00	sort_words

Its clear that only the functions *get* and to a much lesser extent *put* are relevant to performance improvement.

Examine *get* and we find it traverses a linked list. So replace it with a binary tree and the program runs 200x faster on War and Peace.

```
$ clang -03 word_frequency1.c -o word_frequency1

$ time word_frequency1 <WarAndPeace.txt >/dev/null

real 0m0.277s

user 0m0.268s

sys 0m0.008s
```

Was C the best choice for our count words program?

Shell, Perl and Python are slower - but a lot less code. So faster to write, less bugs to find, easier to maintain/modify

```
$ time word_frequency1 <WarAndPeace.txt >/dev/null
       0m0.277s
real
    0m0.268s
user
SVS
      0m0.008s
$ time word frequency.sh <WarAndPeace.txt >/dev/null
real
      0m0.564s
      0m0.584s
user
sys 0m0.036s
$ time word frequency.pl <WarAndPeace.txt >/dev/null
      0m0.643s
real
    0m0.632s
user
    0m0.012s
sys
$ time word frequency.py <WarAndPeace.txt >/dev/null
real 0m1.046s
    0m0.836s
user
       0m0.024s
sys
$ wc word frequency*.*
     759 5912 word frequency1.c
 286
      19 82 word frequency.sh
 11
      38 325 word frequency.py
      43
          301 word frequency.pl
 14
```

# Performance Improvement Example - cp - read/write

Here is a cp implementation in C using low-level calls to read/write

```
while (1) {
    char c[1];
    int bytes read = read(in fd, c, 1);
    if (bytes_read < 0) {</pre>
             perror("cp: ");
             exit(1):
    if (bytes_read == 0)
        return;
    int bytes_written = write(out_fd, c, bytes_read);
    if (bytes_written <= 0) {</pre>
        perror("cp: ");
        exit(1);
```

Its suprisingly slow compared to /bin/cp

```
$ time /bin/cp input_file /dev/null
real 0m0.006s
user 0m0.000s
sys 0m0.004s
$ clang cp0.c -o cp0
$ time ./cp0 input_file /dev/null
real 0m6.683s
user 0m0.932s
sys 0m5.740s
$ clang -03 cp0.c -o cp0
$ time ./cp0 input file /dev/null
real 0m6.688s
user 0m0.900s
sys 0m5.776s
```

- most execution time is spent executing system calls
- as a consequence -O3 is no help.
- 2 system calls for every byte copied a huge overhead.
- modify it to buffer its I/O
- make 2 system calls for every 8192 bytes copied should run much faster.

```
while (1) {
    char c[8192];
    int bytes read = read(in fd, c, sizeof c);
    if (bytes read < 0) {
        perror("cp: ");
        exit(1);
    if (bytes_read <= 0)</pre>
        return;
    int bytes written = write(out fd, c, bytes read);
    if (bytes_written <= 0) {
        perror("cp: ");
        exit(1);
```

```
while (1) {
   int ch = fgetc(in);
   if (ch == EOF)
        break;
   if (fputc(ch, out) == EOF) {
        perror("");
        exit(1);
   }
}
```

Using portable stdio library a byte-by-byte loop runs quite fast, because stdio buffers the I/O behind the scenes.

```
while (1) {
    if(fgets(input, sizeof input, in) == NULL) {
        break;
    }
    if (fprintf(out, "%s", input) == EOF) {
        fprintf(stderr, "cp:");
        perror("");
        exit(1);
    }
}
```

And with a little more complex code we get reasonable speed with portability:

For comparison Perl code which does a copy via an array of lines:

```
die "Usage: cp <src> <dest>\n" if @ARGV != 2;
$in = shift @ARGV;
$out = shift @ARGV;
open IN, '<', $in or
die "Cannot open $in: $!\n";
open OUT, '>', $out or die "Cannot open $out: $!\n";
print OUT <IN>;
```

And Perl code which unsets Perl's line terminator variable so a single read returns the whole file:

```
die "Usage: cp <infile> <outfile>\n" if @ARGV != 2;
$infile = shift @ARGV;
$outfile = shift @ARGV;
open IN, '<', $infile or die "Cannot open $infile: $!\n";
open OUT, '>', $outfile or die "Cannot open $outfile: $!\
undef $/;
print OUT <IN>;
```

#### Performance Improvement Example - Fibonacci

Here is a simple Perl program to calculate the n-th Fibonacci number:

```
sub fib {
    my ($n) = @_;
    return 1 if $n < 3;
    return fib($n-1) + fib($n-2);
}
printf "fib(%d) = %d\n", $_, fib($_) foreach @ARGV;</pre>
```

It becomes slow near n=35.

we can rewrite in C.

# Performance Improvement Example - Fibonacci

```
#include <stdio.h>
int fib(int n) {
    if (n < 3) return 1;
    return fib(n-1) + fib(n-2);
}
int main(int argc, char *argv[]) {
    for (int i = 1; i < argc; i++) {
        int n = atoi(argv[i]);
        printf("fib(%d) = %d\n", n, fib(n));
    }
}</pre>
```

Faster but the program's complexity doesn't change:

```
$ clang -03 -o fib0 fib0.c

$ time fib0 45

fib(45) = 1134903170

real 0m4.994s

user 0m4.976s

sys 0m0.004s
```

# Performance Improvement Example - Fibonacci

```
#!/usr/bin/perl -w
sub fib {
  my ($n) = @_;
  return 1 if $n < 3;
  $f{$n} = fib($n-1) + fib($n-2) if !defined $f{$n};
  return $f{$n};
}
printf "fib(%d) = %d\n", $_, fib($_) foreach @ARGV;</pre>
```

It is very easy to cache already computed results in a Perl hash. This changes the program's complexity from exponential to linear.

```
$ time fib1.pl 45
fib(45) = 1134903170
real     0m0.004s
user     0m0.004s
sys     0m0.000s
```

Now for Fibonanci we could also easily change the program to an iterative form which would be linear too

But memoization is a general technique which can be employed in a variety of situations to improve performnce.