## **Documentation**

Since the program requirement has clearly explained the implementation for most of functionality of this program, this documentation will mainly focus on the algorithm implementation of FileServer that doesn't mentioned in the requirement.

## **General Description**

The program FileServer provide the ability to interact with the FileClient. It runs as a server which maintain the file cache. It is composed with two main functions download() and upload() which are used to be called from FileClient, and a FileCacheEntry class which denotes the file entry in the memory. FileCacheEntry also has download() and upload() method which has the ability to manipulate the file which already cached in the list.

### wait(), notify(), notifyAll(), writeback(), invalidation()

The suspend and resume of download is implemented with wait(), notify, and notifyAll() in the FileCacheEntry's download() and upload() method.

And the file consistency is implemented by the use of writeback() and invalidation() RMI call in the FileCacheEntry's download() and upload() method.

#### Solution of (5d)

If a download write request is received on server when the file state is

write\_shared, it will set the state as ownership\_changed first, and then send the RMI call of writeback to the current owner client, at last it will call wait() on itself to be suspended. After the current owner finish write, it would start running writeback() and it triggers upload() on FileServer. The FileServer will traverse the readers list and make the invalidation() RMI call to each of readers. Then it call notify() to wake up the suspend thread. When this thread got woke up, it will set the owner as itself also call notifyAll(), and the use of notifyAll() will be explained on the next section

## Solution of (5f)

If a download request is received on server when the file state is ownership\_changed, it will do nothing but call wait() at first, because the current thread cannot foresee which will be the next owner to send. After the thread which is downloading the file with write\_shared that explained in the above section call notifyAll(), the current thread with file state ownership\_changed will be resumed, and now it can make a writeback() RMI call to the current owner client(which is the client that request download on write\_shared file state), and then it call wait() again to suspend itself. After the current owner client finished uploading and call notify(), this thread will wake up, then it set the state as write\_shared, set the owner as itself, and call notifyAll(), this notifyAll() is called because there might be multiple of clients are currently suspend on ownership\_changed status, so it can wake up one of the other thread as the next to make writeback() RMI call.

### **Source Code**

```
for (FileCacheEntry fileCacheEntry : fileCacheList) {
   if (fileCacheEntry.getName().equals(filename)) {
```

```
for (FileCacheEntry fileCacheEntry : fileCacheList
```

```
olic boolean isState_notshared() {
```

```
String stateName = stateNameRetrieve();
System.out.println("Upload request by " + clientName + ".");
System.out.println("File: " + getName() + " FileSize: " + bytes.length +
FileState_AfterUpload: " + stateName);
System.out.println("Download request by " + clientName + ". Mode: " + (mode.equals("r") ? "read": "write"));

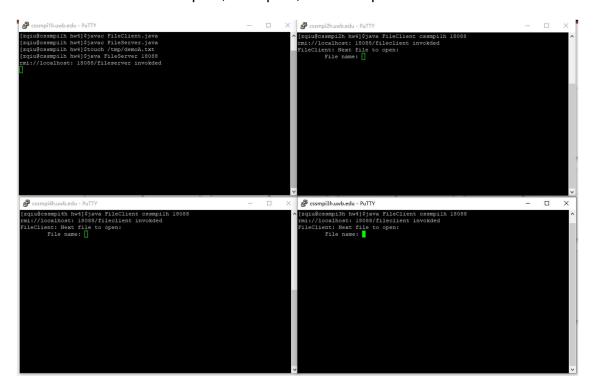
System.out.println("File: " + getName() + "FileSize: " + bytes.length + "bytes
```

```
owner = clientName;
    return new FileContents(bytes);
} else {
    return null;
}

} catch (RemoteException | NotBoundException | MalformedURLException e) {
    e.printStackTrace();
}
return null;
}
}
```

# **Test output**

- (1) Compilation: steps 1-4 (3pt)
- 1. Open 4 windows: each logging in a different cssmpiNh machine. For example, cssmpi1h, cssmpi2h, cssmpi3h, and cssmpi4h.
- 2. Compile with javac.
- 3. Create an empty file: demoA (with a Unix command "touch demoA")
- 4a. Start a server on cssmpi1h.
- 4b. Start a client on cssmpi2h, cssmpi3h, and cssmpi4h.



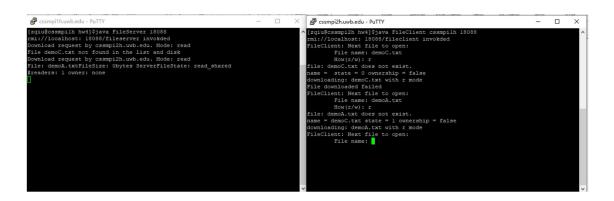
(2) File read test: step 5 (2pt)

5a. Read from demoC at cssmpi2h. (a read error must be handled at the server)

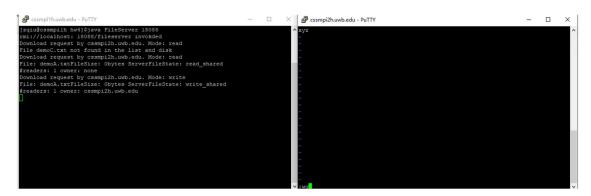


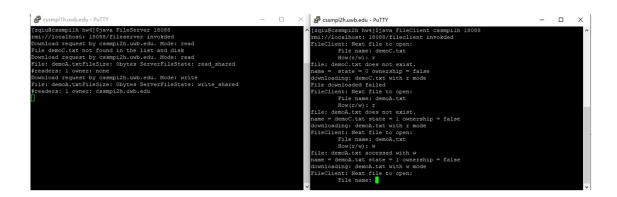
5b. Read empty from demoA at cssmpi2h. (A: clt:rs, svr:rs)



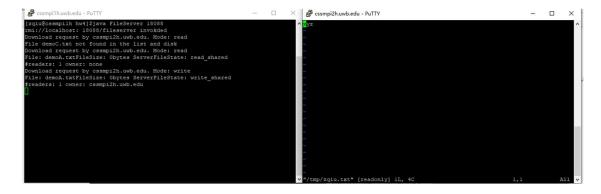


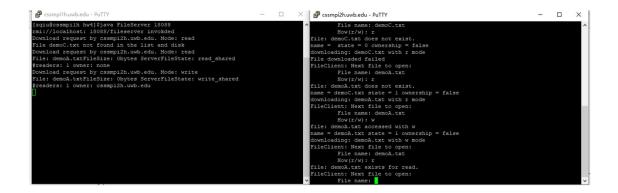
- (3) File write test: steps 6 and 7 (2pt)
- 6. Write xyz to demoA at cssmpi2h. (A: clt:wo, svr:ws)



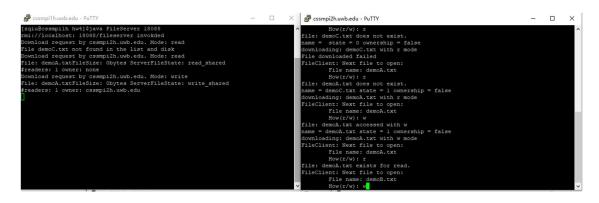


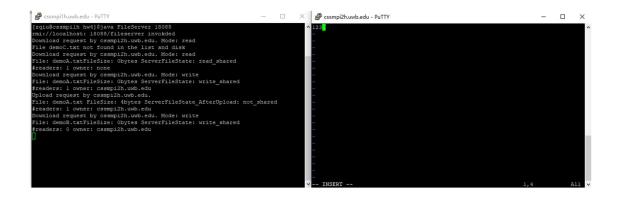
7. Read xyz from demoA at cssmpi2h. (A: clt:wo, svr:ws)

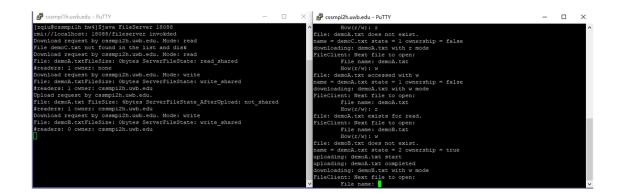




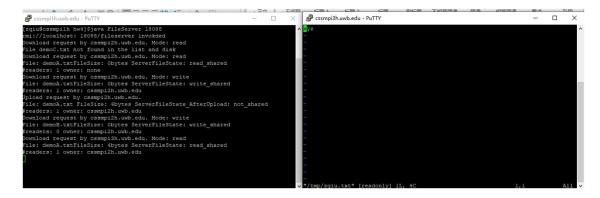
- (4) File replacement test: steps 8, 9, and 10 (3pt)
- 8. Write 123 to demoB at cssmpi2h. (A: clt:iv, svr:ns, B: clt:wo, svr:ws)



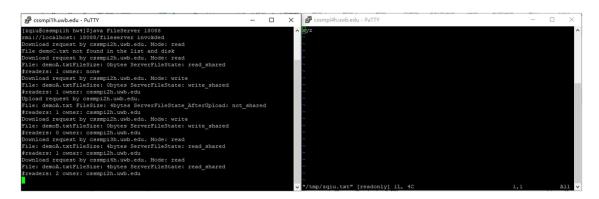




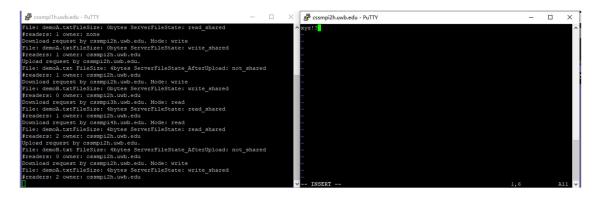
## 9. Read xyz demoA at cssmpi3h. (A: clt:rs, svr:rs)



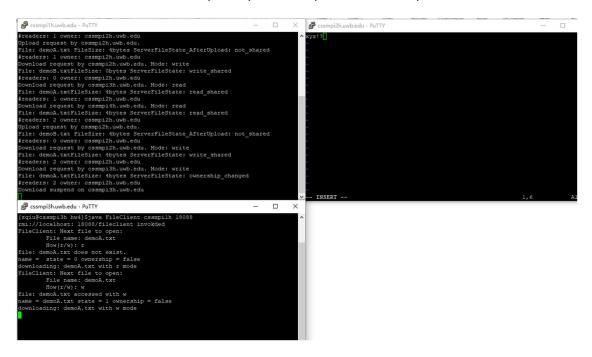
# 10. Read xyz demoA at cssmpi4h. (A: clt:rs, svr:rs)



- (5) File writeback test: steps 11, 13, and 15 (3pt)
- 11. Write xyz?! to demoA at cssmpi2h. (A: clt:wo, svr:ws, B: clt:iv, srv:ns)

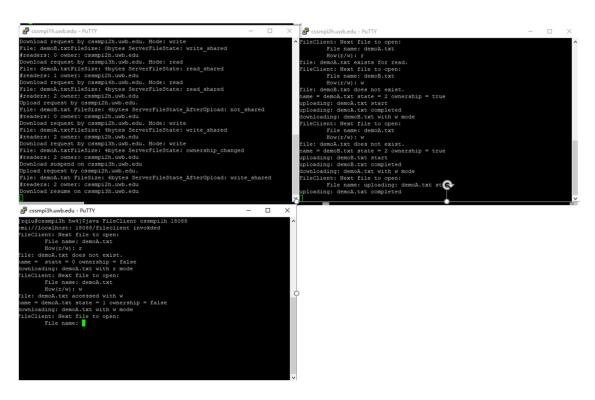


- 12. Keep emacs open at cssmpi2h. (A: clt:wo, svr:ws, B: clt:iv, srv:ns)
- 13. Write to demoA at cssmpi3h. (A: clt:suspended, svr:oc)

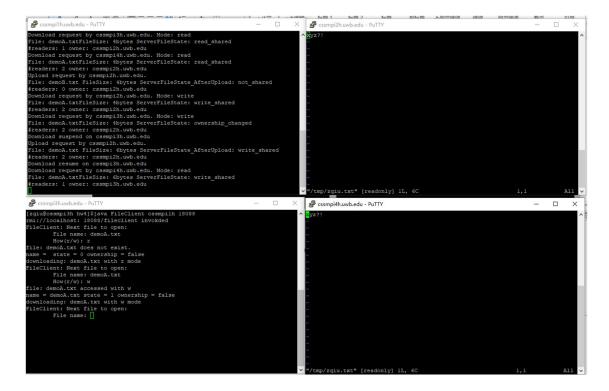


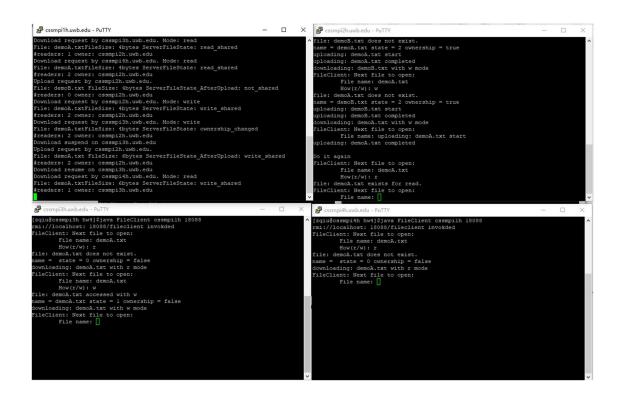
- 14. close emacs at cssmpi2h. (A: clt:rs, svr:ws)
- 15. Write xyz?!abc to demoA at cssmpi3h (A: clt:wo, srv;ws)



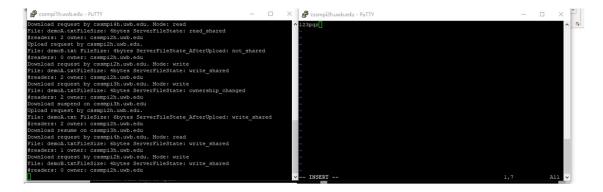


- (6) Session semantics read test: steps 16 and 17 (2pt)
- 16. Read xyz?! from demoA at cssmpi2h.(A: clt:rs, svr:ws)
- 17. Read xyz?! from demoA at cssmpi4h.(A: clt:rs, svr:ws)



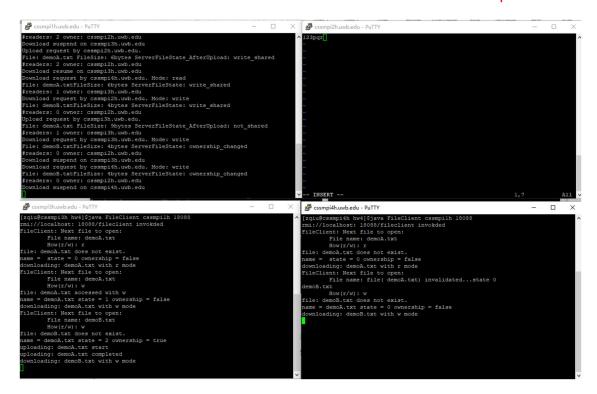


- (7) Multiple write test: steps 18, 20, and 21 (3pt)
- 18. Write 123pqr to demoB at cssmpi2h.(A: clt:iv, svr:ws, B: clt:wo, svr:ws)



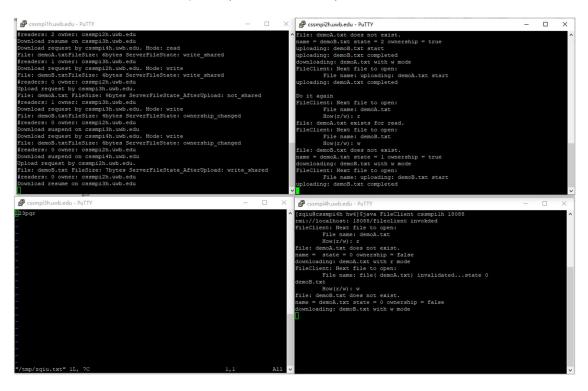
- 19. Keep emacs open at cssmpi2h.
- 20. Write 123pqr456 to demoB at cssmpi3h.(A: clt:iv, svr:rs, B: clt:suspended, svr:oc)
- 21. Write 123pqr456abc to demoB at cssmpi4h. (A: clt:iv, srv:ns, B: clt:suspended, svr:oc)

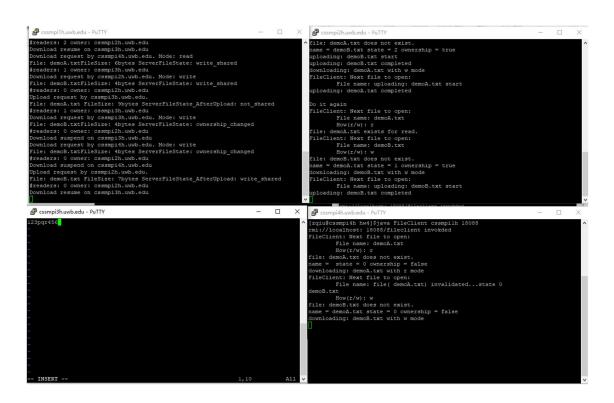
The screenshot of write is shown on 22&23 since the download is suspend



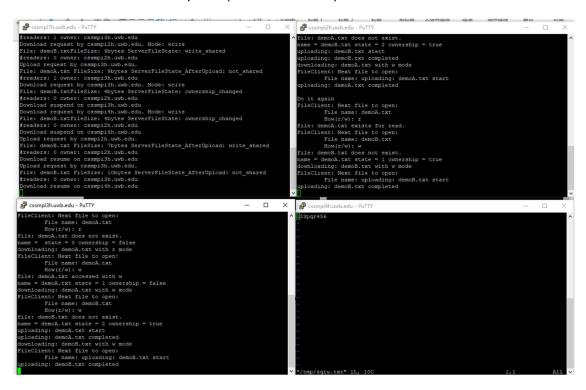
# (8) Quiet termination: steps 22, 23, 24, 25, and 26 (2pts)

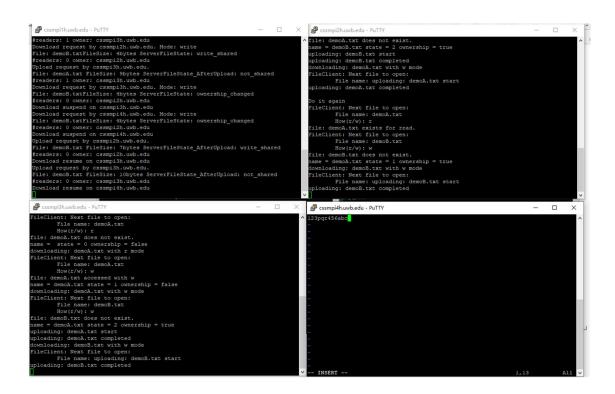
# 22. Close emacs at cssmpi2h. (A: clt:iv, svr:ws)



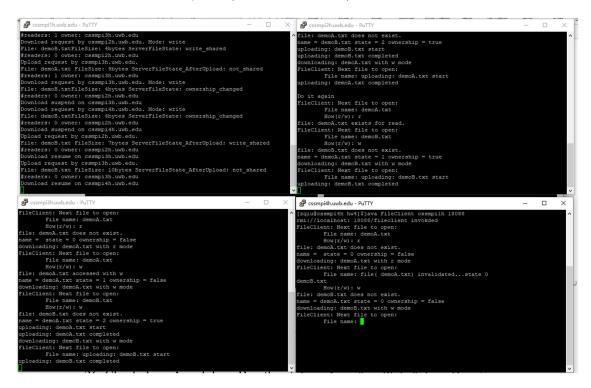


## 23. Close emacs at cssmpi3h. (A: clt:iv, svr:ws)

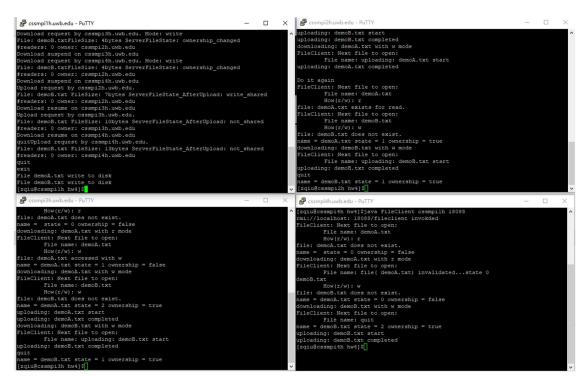




# 24. Close emacs at cssmpi4h. (A: clt:wo, svr:ws)



# 25. quit cssmpi1h, cssmpi2h, cssmpi3h, and cssmpi4h.



(FileServer only recognizes "exit" as termination demand during the test, it has gotten fixed in the code)

26. Check demoA and demoB with cat demoA == "xyz?!abc" demoB =="123pqr456abc"



# **Discussion**

In the discussion, I would like to discuss about two optional algorithm that I came up with during the design of this program. They might don't really have an improvement, but they are various design of implementation.

### **Functional improvements**

 Ask writeback() right after each download for write() rather than request writeback() on busy write request (a download request for write when state is write\_shared/ ownership\_changed)

The goal of this design is to request the modified file update as soon as possible to avoid the case that the client crash may lost the file modification on cache. But it also has a weakness of this design, if the client frequently changes the same in one session login, it needs to upload on each modification, so it might cause high traffic issue. Moreover, a change also needs to be implemented on ownership algorithm, because in the current design the ownership of this file would be lost after upload and that would cause multiple downloads on each write request. My solution addressed to this issue is to have an ownershipInvalidation() RMI call on client side, and the server will call it when there's another client request the download for write.

### **Performance improvement**

Use inQueue index instead of wait() & notify() for suspending

The goal of this design is to let the thread resume followed the sequence of first-in first-out.

This algorithm can be implemented by having a currentSeq variable which is used to record the largest current queue index, and a readyInQueueSeq variable which is used to notify which suspend thread in queue is now ready to be resumed, when a thread starts FileCacheEntry's download(), the currentSeq will be increased 1 if it realize it will be suspend at this time, and then it record the currentSeq as an its own sequence number in the local function variable ownSeq, after that it runs into a dead loop which only get break when (ownSeq == readyInQueueSeq). The readyInQueueSeq will be increased 1 on the execution of upload(). With this design, the threads which suspend on ownership\_changed status would not be random resumed but follow the queue sequence.