

**YONEDA EMBEDDING**

**CATEGORY THEORY STUDY GROUP**

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# YONEDA LEMMA

**Lemma (Yoneda).** *If  $K : D \rightarrow \mathbf{Set}$  is a functor from  $D$  and  $r$  an object in  $D$  (for  $D$  a category with small hom-sets), there is a bijection*

$$y : \text{Nat}(D(r, -), K) \cong Kr \quad (4)$$

*which sends each natural transformation  $\alpha : D(r, -) \rightarrow K$  to  $\alpha_r 1_r$ , the image of the identity  $r \rightarrow r$ .*

The proof is indicated by the following commutative diagram:

$$\begin{array}{ccc}
 D(r, r) & \xrightarrow{\alpha_r} & K(r) \\
 \downarrow f_* = D(r, f) & & \downarrow K(f) \\
 D(r, r) & & K(r)
 \end{array}
 \qquad
 \begin{array}{c}
 r \\
 \downarrow f \\
 r
 \end{array}
 \quad (5)$$

# NATURAL ISOMORPHISM

The Yoneda map  $y$  of (4) is natural in  $K$  and  $r$ . To state this fact formally, we must consider  $K$  as an object in the functor category  $\mathbf{Set}^D$ , regard both domain and codomain of the map  $y$  as functors of the pair  $\langle K, r \rangle$ , and consider this pair as an object in the category  $\mathbf{Set}^D \times D$ . The codomain for  $y$  is then the evaluation functor  $E$ , which maps each pair  $\langle K, r \rangle$  to the value  $Kr$  of the functor  $K$  at the object  $r$ ; the domain is the functor  $N$  which maps the object  $\langle K, r \rangle$  to the set  $\text{Nat}(D(r, -), K)$  of all natural transformations and which maps a pair of arrows  $F: K \rightarrow K'$ ,  $f: r \rightarrow r'$  to  $\text{Nat}(D(f, -), F)$ . With these observations we may at once prove an addendum to the *Yoneda Lemma*:

**Lemma.** *The bijection of (4) is a natural isomorphism  $y: N \xrightarrow{\sim} E$  between*

# PROOF IN HASKELL

```
isoLeft :: (Functor f) => (forall x. (a -> x) -> f x) -> f a
```

```
isoLeft nt = nt id
```

```
isoRight :: (Functor f) => f a -> (forall x. (a -> x) -> f x)
```

```
isoRight fa g = fmap g fa
```

# (CO)YONEDA EMBEDDING

The object function  $r \mapsto D(r, -)$  and the arrow function

$$(f: s \rightarrow r) \mapsto D(f, -): D(r, -) \rightarrow D(s, -)$$

for  $f$  an arrow of  $D$  together define a full and faithful functor

$$Y: D^{\text{op}} \rightarrow \mathbf{Set}^D \quad (6)$$

called the *Yoneda functor*. Its dual is another such functor

$$Y': D \rightarrow \mathbf{Set}^{D^{\text{op}}} \quad (7)$$

(also faithful) which sends  $f: s \rightarrow r$  to the natural transformation

# IN HASKELL

```
arrowMap :: (b -> a) -> (forall x. (a -> x) -> (b -> x))  
arrowMap f g = g . f
```

# CHALLENGE 1

Express the co-Yoneda embedding in Haskell.

## SOLUTION

```
arrowMap :: (a -> b) -> (forall x. (x -> a) -> (x -> b))  
arrowMap f g = f . g
```

## CHALLENGE 2

Show that the bijection we established between  $\text{fromY}$  and  $\text{btoa}$  is an isomorphism (the two mappings are the inverse of each other).



# CHALLENGE 3

Work out the Yoneda embedding for a monoid. What functor corresponds to the monoid's single object? What natural transformations correspond to monoid morphisms?

# CHALLENGE 4

What is the application of the covariant Yoneda embedding to preorders? (Question suggested by Gershon Bazerman.)

## SOLUTION

$a \leq b$  if and only if  $\forall x, b \leq x \Rightarrow a \leq x$

# CHALLENGE 5

Yoneda embedding can be used to embed an arbitrary functor category  $[C, D]$  in the functor category  $[[C, D], \text{Set}]$ . Figure out how it works on morphisms (which in this case are natural transformations).

**THANK YOU!**