

CSC 226 SUMMER 2018
ALGORITHMS AND DATA STRUCTURES II
ASSIGNMENT 4 - WRITTEN
UNIVERSITY OF VICTORIA

1. Show the longest common subsequence table, `llcs[][]`, for the following two strings:

$x = \text{"skullandbones"}$
 $y = \text{"lullabybabies"}$

What is the longest common subsequence between these strings?

2. Write a pseudocode description of the `printLCS()` algorithm, which prints the longest common subsequence of two strings x and y . Your algorithm should take as input the completed `llcs[][]` integer array of longest common subsequence lengths, and the two strings x and y . (So, you do not have the `path[][]` array - see Lecture 19, slides 100 and 101.) Your algorithm must return the specific string corresponding the length found in `llcs[n][m]` and it should run in $O(n + m)$ time, where n is the length of x and m is the length of y .
3. (a). Give the `right[]` array computed by the Boyer-Moore algorithm for the pattern **ABRACADABRA**. Assume the alphabet is $\Sigma = \{A, B, C, D, R\}$.
- (b). Give the `dfa[][]` array for the KMP algorithm for the pattern **ABRACADABRA**, and draw the transition diagram corresponding to it. Assume the alphabet is $\Sigma = \{A, B, C, D, R\}$.
4. Write an efficient method that takes a string `txt` and an integer `M` as arguments and returns the position of the first occurrence of `M` consecutive blanks in the string. If there is no such occurrence it should return the length of `txt` (`txt.length`). Give a runtime analysis of the method.
5. Modify the Rabin-Karp algorithm to search for a given pattern with the additional provision that some character in the pattern is a wildcard, that is, it can be any character. So, for some $k = 0, 1, \dots, M - 1$, where the length of the pattern is M , the k th character of the pattern can be any character. You may assume that k is given as input along with the text and pattern.