

CSC 320 - ASSIGNMENT 4 - SOLUTIONS

1. let $G = (\{R, S, T, X\}, \{a, b\}, P, S)$ where P is

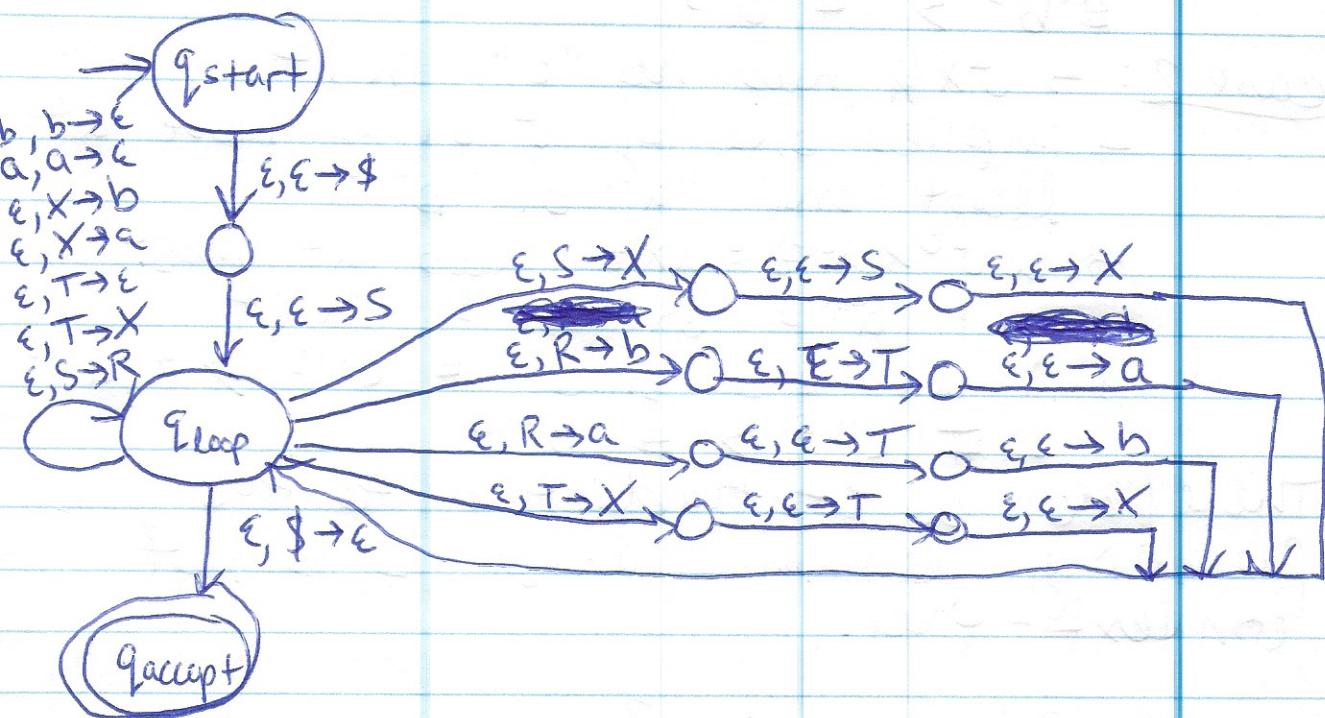
$$S \rightarrow XSX \mid R$$

$$R \rightarrow aTb \mid bTa$$

$$T \rightarrow XTX \mid X \mid \epsilon$$

$$X \rightarrow a \mid b$$

Construct PDA $D = (Q, \{a, b\}, \{a, b, \$\}, S, q_{\text{start}}, \{q_{\text{accept}}\})$ with δ as follows,



2. Let $L = \{a^n b^n c^i \mid i \leq n\}$.

- Suppose L is context-free

- then $\exists p$, the pumping length by pumping lemma
- let $s = a^p b^p c^p$, clearly $|s| = 3p \geq p$
- Then $s = uvxyz$ such that $|vxy| \leq p$, $v, y \neq \epsilon$

- Two cases, uxy contains a "c" or it doesn't.
- case 1: $\rightarrow uxy$ has a "c" in it
 - \rightarrow implies uxy has no "a"'s
 - \rightarrow one of u or y have at least one "b" or "c"
 - $\rightarrow uv^2xy^2z$ has at least $p+1$ "b"'s or $p+1$ "c"'s in it
 - $\rightarrow \therefore uv^2xy^2z \notin L$ since ~~# "c"'s > # "a"'s~~ and ~~# "b"'s > # "a"'s~~
- case 2: $\rightarrow uxy$ has no "c" in it
 - \rightarrow implies one of u or y have at least one "a" or "b"
 - $\rightarrow uv^0xy^0z = uxz$ has at most $p-1$ "a"'s or "b"'s in it
 - $\rightarrow \therefore uxz \notin L$ since $\# "c"'s > \# "a"'s$ and $\# "c"'s > \# "b"'s$
- Therefore a contradiction to the pumping lemma and L is not context-free.

3. Let $L = \{0^p \mid p \text{ is prime}\}$.

- Suppose L is context-free, then by the pumping lemma

- $\exists P$ the pumping length
- Let $p \geq n+2$ be a prime number
- Let $s = 0^p$, so $|s| \geq n$
- Let $s = uvxyz$ such that $|uxy| \leq n$, $v, y \neq \epsilon$
- Let $|vxy| = m \geq 1$, then $|uxz| = p-m$

- Now, consider $uv^{p-m}xy^{p-m}z$.
- $|uv^{p-m}xy^{p-m}z| = |uxz| + (p-m)|v| + (p-m)|y|$
 $= (p-m) + (p-m)(|v| + |y|)$
 $= (p-m) + (p-m)m$
 $= (m+1)(p-m)$
- Thus, $uv^{p-m}xy^{p-m}z = O^{(m+1)(p-m)}$ where
 $m \geq 1 \Rightarrow m+1 \geq 2$ and $p \geq n+2$, and $m \leq n$
- Implies that $p-m \geq 2$, thus $(m+1)(p-m)$ is not prime.
- Therefore $uv^{p-m}xy^{p-m}z \notin L$ and L is not context-free.

4. Let $L_1 = \{a^n b^{2n} c^m \mid n, m \geq 0\}$ and $L_2 = \{a^n b^m c^{2m} \mid n, m \geq 0\}$

(a) Let $G_1 = (\{S_1\}, A_1, \{a, b, c\}, P_1, S_1)$ and $G_2 = (\{S_2\}, A_2, \{a, b, c\}, P_2, S_2)$ where P_1 and P_2 are as follows:

$$\begin{aligned} P_1: \quad S_1 &\rightarrow aS_1, bbaA_1, |A_1| \epsilon \\ &A_1 \rightarrow cA_1, |A_1| \epsilon \end{aligned}$$

$$\begin{aligned} P_2: \quad S_2 &\rightarrow A_2 b S_2 c c | A_2 | \epsilon \\ &A_2 \rightarrow aA_2 | A_2 | \epsilon \end{aligned}$$

Therefore, $L_1 + L_2$ are context-free as G_1 and G_2 are context-free grammars that accept L_1 and L_2 , respectively.

$$(b) L_1 \cap L_2 = \{a^n b^{2n} c^{4n} \mid n \geq 0\}.$$

- Assume $L_1 \cap L_2$ is context-free then
- $\exists p$ such that it satisfies the pumping lemma.

- let $s = a^p b^{2p} c^{4p}$, where $|s| = 7p \geq p$
- Let $s = uvxyz$ such that $|vxy| \leq p$ and $u, y \neq \epsilon$.
- Two cases, either vxy contains a c or it doesn't

- case 1: \rightarrow Suppose vxy contains ~~no~~ c
 - \rightarrow then v or y has at least one a or b
 - \rightarrow if an a, then uv^2xy^2z has at least $p+1$ a's and at most $2p$ a's
 - \rightarrow But $4(p+1) = 4p+4 > 4p$ so ratio of a's + c's is off
 - \rightarrow if a b, then uv^2xy^2z has at least $2p+1$ b's
 - \rightarrow But $2(2p+1) = 4p+2 > 4p$ so ratio of b's + c's is off.
 - \rightarrow Thus, $uv^2xy^2z \notin L_1 \cap L_2$
- case 2: \rightarrow Suppose vxy contains a c
 - \rightarrow Then, vxy has no a's since $|vxy| \leq p$.
 - \rightarrow v or y has at least one b or c.
 - \rightarrow If a b, then uv^2xy^2z has at least $2p+1$ b's
 - \rightarrow But $2p \leq 2p+1$ and ratio of a's + b's is off
 - \rightarrow If a c, then uv^2xy^2z has at least $4p+1$ c's, by $4p \leq 4p+1$ and the ratio of a's + c's is off.
 - \rightarrow Thus, $uv^2xy^2z \notin L_1 \cap L_2$
- Therefore, by the pumping lemma, $L_1 \cap L_2$ is not context-free.

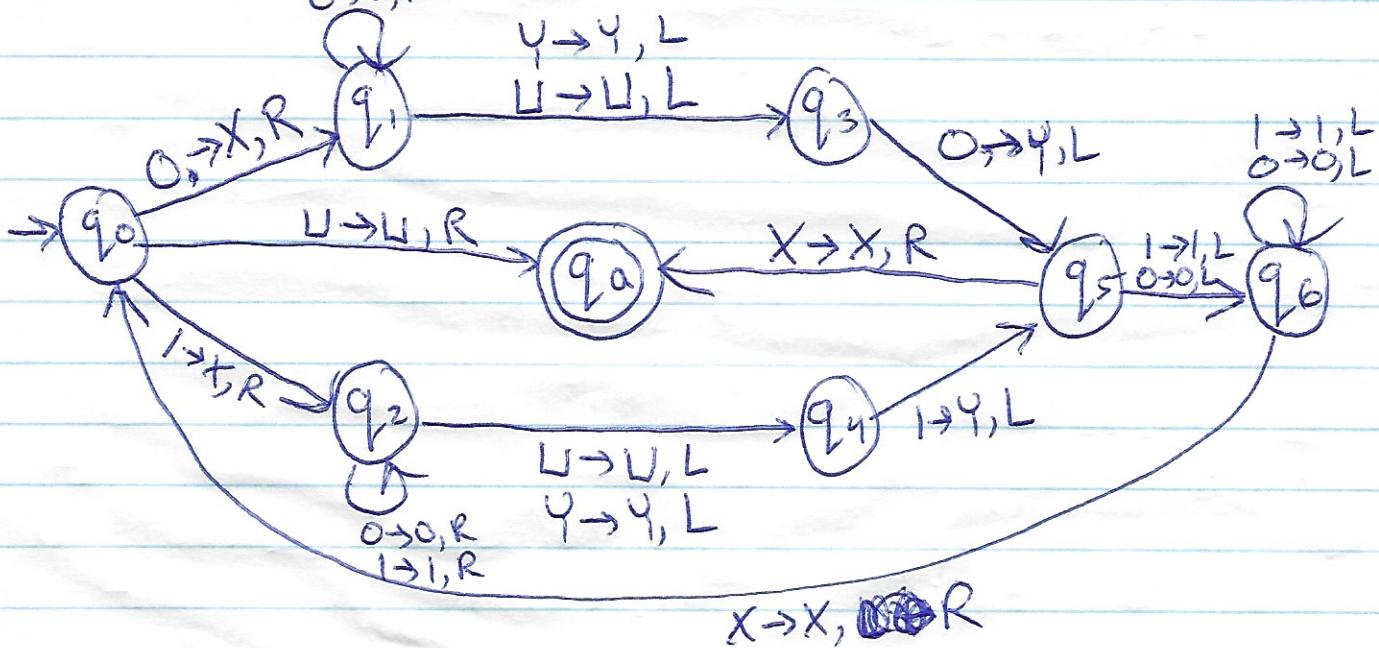
(3)

5. Let $L = \{ww^R \mid w \in \{0,1\}^*\}$. We will design a Turing Machine $M = (Q, \{0,1\}, \Gamma, \delta, q_0, q_a, q_r)$ with the following high-level description.

High-level:

1. Read the first symbol, 0 or 1, and mark it.
2. Move to the other end of the tape and read the last symbol
3. if it's the same mark it and move back to the beginning
4. if it's not, reject
5. repeat until all symbols are marked and accept.

Formal: let $Q = \{q_0, q_1, q_2, q_3, q_4, q_5, q_6, q_a, q_r\}$, $\Gamma = \{a, b, X, Y, U, L\}$ and δ as follows



→ All other transitions go to q_r