

# STOR 614 - Linear Programming, Spring 2019

## Homework No. 4

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May 22, 2019

### Problem 1.

(a)

Phase I.

Add artificial variables  $y_1, y_2$ .

$$\begin{array}{llllll}
 \max & t = & & & -y_1 & -y_2 \\
 s.t & & 2x_1 & +x_2 & +x_3 & +y_1 & = 4 \\
 & & x_1 & +x_2 & +2x_3 & & +y_2 = 2 \\
 & & x_1, & x_2, & x_3, & y_1, & y_2 \geq 0
 \end{array}$$

Transform LP to canonical form.

$$\begin{array}{llllll}
 \max & t = & 3x_1 & +2x_2 & +3x_3 & -6 \\
 s.t & & 2x_1 & +x_2 & +x_3 & +y_1 & = 4 \\
 & & x_1 & +x_2 & +2x_3 & & +y_2 = 2 \\
 & & x_1, & x_2, & x_3, & y_1, & y_2 \geq 0
 \end{array}$$

The initial tableau is as follows

| $t$ | $x_1$ | $x_2$ | $x_3$ | $y_1$ | $y_2$ | rhs | Basic var |
|-----|-------|-------|-------|-------|-------|-----|-----------|
| 1   | -3    | -2    | -3    | 0     | 0     | -6  | $t = -6$  |
| 0   | 2     | 1     | 1     | 1     | 0     | 4   | $y_1 = 4$ |
| 0   | 1     | 1     | 2     | 0     | 1     | 2   | $y_2 = 2$ |

First iteration:  $x_1$  enters and  $y_1$  leaves.

| $t$ | $x_1$ | $x_2$ | $x_3$ | $y_1$ | $y_2$ | rhs | Basic var |
|-----|-------|-------|-------|-------|-------|-----|-----------|
| 1   | 0     | -1/2  | -3/2  | 3/2   | 0     | 0   | $t = 0$   |
| 0   | 1     | 1/2   | 1/2   | 1/2   | 0     | 2   | $x_1 = 2$ |
| 0   | 0     | 1/2   | 3/2   | -1/2  | 1     | 0   | $y_2 = 0$ |

Second iteration:  $x_2$  enters and  $y_2$  leaves.

| $t$ | $x_1$ | $x_2$ | $x_3$ | $y_1$ | $y_2$ | rhs | Basic var |
|-----|-------|-------|-------|-------|-------|-----|-----------|
| 1   | 0     | 0     | 0     | 1     | 1     | 0   | $t = 0$   |
| 0   | 1     | 0     | -1    | 1     | -1    | 2   | $x_1 = 2$ |
| 0   | 0     | 1     | 3     | -1    | 2     | 0   | $x_2 = 0$ |

The optimal  $t = 0$ . Case 2.1. Obtain a simplex tableau for the original LP.

| $z$ | $x_1$ | $x_2$ | $x_3$ | rhs | Basic var |
|-----|-------|-------|-------|-----|-----------|
| 1   | 0     | 0     | 2     | 2   | $z = 2$   |
| 0   | 1     | 0     | -1    | 2   | $x_1 = 2$ |
| 0   | 0     | 1     | 3     | 0   | $x_2 = 0$ |

First iteration:  $x_3$  enters and  $x_2$  leaves.

| $z$ | $x_1$ | $x_2$ | $x_3$ | rhs | Basic var |
|-----|-------|-------|-------|-----|-----------|
| 1   | 0     | -2/3  | 0     | 2   | $z = 2$   |
| 0   | 1     | 1/3   | 0     | 2   | $x_1 = 2$ |
| 0   | 0     | 1/3   | 1     | 0   | $x_3 = 0$ |

The original LP has an optimal solution  $(x_1, x_2, x_3) = (2, 0, 0)$  and the optimal value is  $z = 2$ .

(b)

Convert to standard form.

$$\begin{aligned}
 \min \quad & z = 2x_1 + 3x_2 \\
 s.t \quad & 1/2x_1 + 1/4x_2 + s_1 = 4 \\
 & x_1 + 3x_2 - s_2 = 20 \\
 & x_1 + x_2 = 10 \\
 & x_1, x_2, s_1, s_2 \geq 0
 \end{aligned}$$

Add artificial variables  $y_1, y_2$ . The phase I LP is as follows.

$$\begin{array}{llllllll}
 \max & t = & & & & & -y_1 & -y_2 \\
 s.t & & 1/2x_1 & +1/4x_2 & +s_1 & & & = 4 \\
 & & x_1 & +3x_2 & & -s_2 & +y_1 & = 20 \\
 & & x_1 & +x_2 & & & & +y_2 = 10 \\
 & & x_1, & x_2, & s_1, & s_2, & y_1, & y_2 \geq 0
 \end{array}$$

Convert phase I LP to canonical form.

$$\begin{array}{llllllll}
 \max & t = & 2x_1 & +4x_2 & & -s_2 & -30 \\
 s.t & & 1/2x_1 & +1/4x_2 & +s_1 & & & = 4 \\
 & & x_1 & +3x_2 & & -s_2 & +y_1 & = 20 \\
 & & x_1 & +x_2 & & & & +y_2 = 10 \\
 & & x_1, & x_2, & s_1, & s_2, & y_1, & y_2 \geq 0
 \end{array}$$

The initial tableau of phase I LP is as follows

| $t$ | $x_1$ | $x_2$ | $s_1$ | $s_2$ | $y_1$ | $y_2$ | rhs | Basic var  |
|-----|-------|-------|-------|-------|-------|-------|-----|------------|
| 1   | -2    | -4    | 0     | 1     | 0     | 0     | -30 | $t = -30$  |
| 0   | 1/2   | 1/4   | 1     | 0     | 0     | 0     | 4   | $s_1 = 4$  |
| 0   | 1     | 3     | 0     | -1    | 1     | 0     | 20  | $y_1 = 20$ |
| 0   | 1     | 1     | 0     | 0     | 0     | 1     | 10  | $y_2 = 10$ |

First iteration:  $x_1$  enters,  $s_1$  leaves.

| $t$ | $x_1$ | $x_2$ | $s_1$ | $s_2$ | $y_1$ | $y_2$ | rhs | Basic var  |
|-----|-------|-------|-------|-------|-------|-------|-----|------------|
| 1   | 0     | -3    | 4     | 1     | 0     | 0     | -14 | $t = -14$  |
| 0   | 1     | 1/2   | 2     | 0     | 0     | 0     | 8   | $x_1 = 8$  |
| 0   | 0     | 5/2   | -2    | -1    | 1     | 0     | 12  | $y_1 = 12$ |
| 0   | 0     | 1/2   | -2    | 0     | 0     | 1     | 2   | $y_2 = 2$  |

Second iteration:  $x_2$  enters,  $y_2$  leaves.

| $t$ | $x_1$ | $x_2$ | $s_1$ | $s_2$ | $y_1$ | $y_2$ | rhs | Basic var |
|-----|-------|-------|-------|-------|-------|-------|-----|-----------|
| 1   | 0     | 0     | -8    | 1     | 0     | 6     | -2  | $t = -2$  |
| 0   | 1     | 0     | 4     | 0     | 0     | -1    | 6   | $x_1 = 6$ |
| 0   | 0     | 0     | 8     | -1    | 1     | -5    | 2   | $y_1 = 2$ |
| 0   | 0     | 1     | -4    | 0     | 0     | 2     | 4   | $x_2 = 4$ |

Third iteration:  $s_1$  enters,  $y_1$  leaves.

| $t$ | $x_1$ | $x_2$ | $s_1$ | $s_2$ | $y_1$ | $y_2$ | rhs | Basic var   |
|-----|-------|-------|-------|-------|-------|-------|-----|-------------|
| 1   | 0     | 0     | 0     | 0     | 1     | 1     | 0   | $t = 0$     |
| 0   | 1     | 0     | 0     | 1/2   | -1/2  | 3/2   | 5   | $x_1 = 5$   |
| 0   | 0     | 0     | 1     | -1/8  | 1/8   | -5/8  | 1/4 | $s_1 = 1/4$ |
| 0   | 0     | 1     | 0     | -1/2  | 1/2   | -1/2  | 5   | $x_2 = 5$   |

The optimal  $t = 0$ . Case 2.1. Obtain a simplex tableau for the original LP.

| $z$ | $x_1$ | $x_2$ | $s_1$ | $s_2$ | rhs | Basic var   |
|-----|-------|-------|-------|-------|-----|-------------|
| 1   | 0     | 0     | 0     | -1/2  | 25  | $z = 25$    |
| 0   | 1     | 0     | 0     | 1/2   | 5   | $x_1 = 5$   |
| 0   | 0     | 0     | 1     | -1/8  | 1/4 | $s_1 = 1/4$ |
| 0   | 0     | 1     | 0     | -1/2  | 5   | $x_2 = 5$   |

The current BFS is optimal. The optimal solution of the original LP is  $(x_1, x_2) = (5, 5)$  and the optimal value is  $z = 25$ .

## Problem 2.

(a)

This problem is feasible if and only if there exists  $k$  such that  $b/a_k \geq 0$ .

*Proof.* If there exists  $k$  such that  $b/a_k \geq 0$ , then  $\{x_i = 0 \ (i \neq k), x_k = b/a_k\}$  is a feasible solution.

If there doesn't exist  $k$  such that  $b/a_k \geq 0$ , then for all  $a_i \neq 0$ ,  $b/a_i < 0$ . Suppose  $b > 0$ , then all  $a_i \leq 0$ . Then  $\sum_{i=1}^n a_i x_i \leq 0$ , thus this problem is not feasible. If  $b < 0$ , then  $\sum_{i=1}^n a_i x_i \geq 0$ , and this problem is not feasible.  $\square$

(b)

Let  $K = \{i \mid b/a_i \geq 0\}$ . The set of BFS's is  $\{x_i = 0 \ (i \neq k), x_k = b/a_k \mid k \in K\}$ . Thus the set of objective function values at BFS's is  $\{c_k b/a_k \mid k \in K\}$ .

Find

$$p = \min_{i \in K} c_i b/a_i$$

The optimal solution is  $x_i = 0 \ (i \neq p), x_p = b/a_p$  and the optimal value is  $c_p b/a_p$ .

### Problem 3.

Let  $z$  be the BFS at a certain step. In each step, replace a basis of  $z$  that is not a basis of  $y$  by a basis of  $y$  that is not a basis of  $z$ . In this way, the number of shared bases of  $y$  and  $z$  increases by 1 in each iteration. Thus we can go from  $x$  to  $y$  in a finite number of steps.

### Problem 4.

Proposition 1. When formulating the phase I LP, if a non-artificial variable  $x$  satisfies:

$x$  has coefficient 1 in an equation and zero coefficients in all the other equations. (condition 1)

then we do not need to add an artificial variable in the equation where  $x$  has coefficient 1.

When an artificial variable  $y$  becomes nonbasic, a non-artificial variable becomes basic and satisfies condition 1. By proposition 1, we can formulate a new phase I LP without  $y$ . The new phase I LP is equivalent to the LP obtained by eliminating the column of  $y$  from the old phase I LP.

### Problem 5.

Use Matlab to solve the problems.

(a) The problem is unbounded.

(b) The optimal solution is  $(0, 5, 0, 1, 0, 4)$  and the optimal value is  $-74$ .

Matlab code for revised simplex method:

```
1 %% LP (a)
2 A=[ 0 1 2 1 1 0 -5;
```

```

3      0 2 1 0 -2 1 0;
4      1 1 -2 0 1 0 3];
5  c = -[2;3;-4;3;1;-4;6];
6  b = [2;1;3];
7  B = [4 6 1];
8
9  %% LP (b)
10 A = [1 0 0 -1 1 1;
11      1 1 0 -1 3 0;
12      1 1 1 -3 0 1];
13 c = -[14;-19;0;21;52;0];
14 b = [3;4;6];
15 B = [1 2 6];
16
17 %% revised simplex method
18 k=0;
19 while(1)
20     N = setdiff(1:length(c), B);
21     k=k+1;
22     k
23     AB = A(:,B);
24     AB_inv = inv(AB);
25
26     reduced_cost = -c.' + c(B).'*AB_inv*A
27     coef_x = AB_inv*A
28     rhs_z=c(B).'*AB_inv*b
29     rhs_x=AB_inv*b
30
31     t = find(reduced_cost(N)<0);
32     if length(t) == 0
33         fprintf('found optimal value\n')
34         break
35     end
36     pivot_column = N(t(1));
37     u = AB_inv*A(:,pivot_column);
38     t = find(u>0);

```

```
39     if length(t) == 0
40         fprintf('unbounded\n')
41         break
42     end
43     [i,min_i] = min(rhs_x(t) ./ u(t));
44     pivot_row = t(min_i);
45     B(pivot_row) = pivot_column;
46 end
47 z=c(B) .*AB_inv*b;
```