# STOR 614 - Linear Programming, Spring 2019 Homework No. 2

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#### Problem 1.

Suppose that P contains a line  $Q = \{x + \lambda d \mid \lambda \in \mathbb{R}\}$ , where  $x \in \mathbb{R}, d \in \mathbb{R}, d \neq 0$ . Because the set  $\{a_1, \ldots, a_m\}$  contains n linearly independent vectors, there exists  $a_k$ , such that  $a_k^T d \neq 0$ . Therefore,

$$\left\{a_k^T q \mid q \in Q\right\} = \left\{a_k^T x + \lambda a_k^T d \mid \lambda \in \mathbb{R}\right\} = \mathbb{R},$$

which contradicts with

$$a_k^T q \ge b_k$$
, for all  $q \in Q$ 

#### Problem 2.

Basic feasible solutions:

$$\begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1/3 \\ 2/3 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 1/3 \\ 2/3 \end{bmatrix}, \begin{bmatrix} 1/2 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 0 \end{bmatrix}.$$

Degenerate BFS:

$$\begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1/3 \\ 2/3 \end{bmatrix}$$

because it has 3 active constraints:

$$x_1 + x_2 \le 1$$
$$4x_1 + x_2 \le 2$$
$$2x_1 + x_2 \le 4/3.$$

#### Problem 3.

$$\begin{array}{lll} \min & z = 3x_1 + x_2 \\ \mathrm{s.t} & x_1 & -\beta_1 & = 3 \\ & x_1 + x_2 & +\beta_2 = 4, \\ & 2x_1 - x_2 & = 3, \\ & x_1, \ x_2, \ \beta_1, \ \beta_2 \geq 0 \end{array}$$

### Problem 4. Problem 4(a)

True.

Let  $A \in \mathbb{R}^{m \times n}$ . Because A has full row rank,  $m \leq n$ . Let  $x_{B1}$  and  $x_{B2}$  be the bases of x. Let  $x_{N1}$  and  $x_{N2}$  be the corresponding collections of nonbasic variables, i.e.

$$x_{N1} = \{x_1, \dots, x_n\} \setminus x_{B1}$$
$$x_{N2} = \{x_1, \dots, x_n\} \setminus x_{B2}$$

Let  $n_0$  be the number of zeros in x. Then,

$$x_{B1} \neq x_{B2}$$
  
 $\Rightarrow x_{N1} \neq x_{N2}$   
 $\Rightarrow |x_{N1} \cup x_{N2}| > n - m \quad (|\cdot| \text{ denotes cardinality})$   
 $\Rightarrow n_0 > n - m$ 

Therefore, x is degenerate.

#### Problem 4(b)

False.

Counter example:  $A = \begin{bmatrix} 1 & 0 \end{bmatrix}$ ,  $b = \begin{bmatrix} 0 \end{bmatrix}$ .  $x = \begin{bmatrix} 0 & 0 \end{bmatrix}^T$  is a degenerate basic solution but has only one basis,  $\begin{bmatrix} 1 \end{bmatrix}$ .

# Problem 5. Problem 5(a)

**Proposition 1.** Let  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^n$ ,  $b \neq 0$ . If

- 1. Ax = b has a solution  $x^*$ ,  $x_i^* > 0$ , for i = 1, ..., n
- 2. The columns of A are not linearly independent

then there exists a solution of Ax = b, x', such that

- 1.  $x_i' \ge 0$ , for i = 1, ..., n,
- 2. The columns of A corresponding to nonzero entries of x' are linearly independent.

Proof. Assume the induction hypothesis that the proposition holds for all  $n < N(N \ge 3)$ . Consider n = N. Let d be a nontrivial solution of Ax = 0. Then,  $\exists \lambda \in \mathbb{R}$ , such that  $y = (x^* + \lambda d)$  has at least one zero entry and no negative entries. Let  $A' \in \mathbb{R}^{m \times N'}$  be the matrix containing the columns of A corresponding to nonzero entries of y (N' < N). If the columns of A' are linearly independent, then x' = y. Otherwise, by induction hypothesis, there exists a solution of A'x = b, y', such that  $y'_i \ge 0$ , for  $i = 1, \ldots, N'$ , and the columns of A' corresponding to nonzero entries of y' are linearly independent. Then, x' can be obtained by replacing the nonzero entries of y by corresponding entries of y'.

Proof for problem 5(a).

If b = 0, then 0 is a degenerate basic feasible solution of P, so  $b \neq 0$ .

Suppose that x is not a basic feasible solution, then the columns of A corresponding to nonzero entries of x are not linearly independent. By proposition 1, there exists  $x' \in \mathbb{R}^n$ , such that x' is a feasible solution, x' has more than (n-m) zeros, and the columns of A corresponding to nonzero entries of x' are linearly independent. Therefore, x' is a degenerate basic feasible solution, which contradicts with that all basic feasible solutions are nondegenerate.

## **Problem 5(b)** A counter example:

$$A = \begin{bmatrix} 2 & 0 \end{bmatrix}$$
$$b = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$
$$x = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$

x has exactly 1 positive component but is not a basic feasible solution.