

# CS 225

## Data Structures

*April 9 – Graphs Intro  
Wade Fagen-Ulmschneider*

# Disjoint Sets Analysis

The **iterated log** function:

*The number of times you can take a log of a number.*

$\log^*(n) =$

$$0 \quad , n \leq 1$$

$$1 + \log^*(\log(n)) , n > 1$$

What is  $\lg^*(2^{65536})$ ?

# Disjoint Sets Analysis

In an Disjoint Sets implemented with smart **unions** and path compression on **find**:

Any sequence of **m union** and **find** operations result in the worse case running time of  $O(\underline{\hspace{2cm}})$ ,  
where **n** is the number of items in the Disjoint Sets.

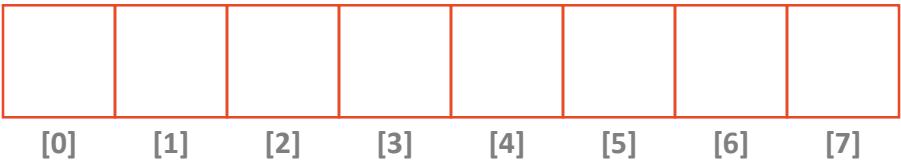
# In Review: Data Structures

## Array

- Sorted Array
- Unsorted Array
  - Stacks
  - Queues
  - Hashing
  - Heaps
    - Priority Queues
  - UpTrees
  - Disjoint Sets

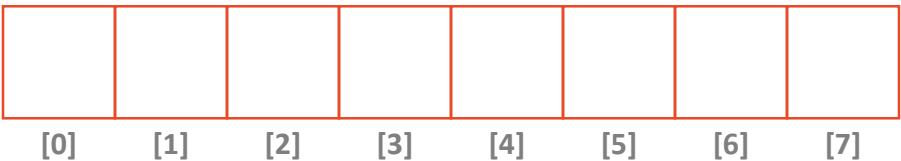
## List

- Doubly Linked List
- Skip List
- Trees
  - BTree
  - Binary Tree
  - Huffman Encoding
  - kd-Tree
  - AVL Tree

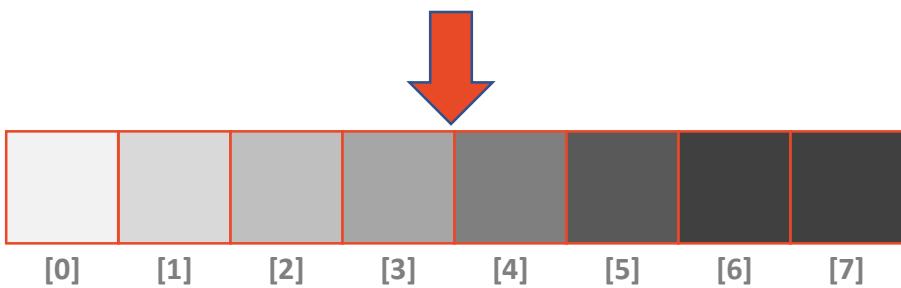


# Array

- Constant time access to any element, given an index  
 $a[k]$  is accessed in  $O(1)$  time, no matter how large the array grows
- Cache-optimized  
Many modern systems cache or pre-fetch nearby memory values due the “Principle of Locality”. Therefore, arrays often perform faster than lists in identical operations.

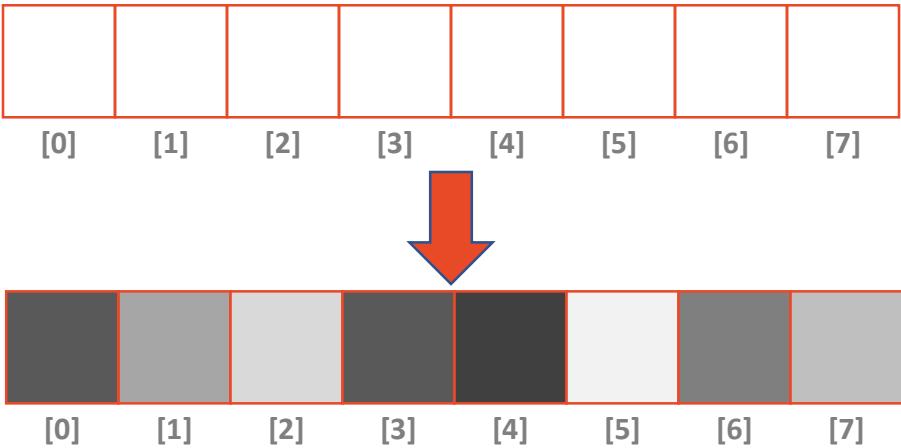


Array



Sorted Array

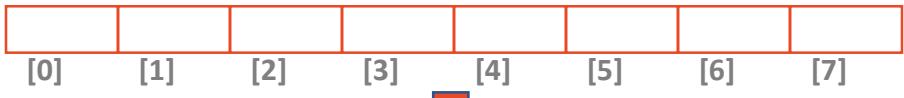
- Efficient general search structure  
**Searches on the sort property run in  $O(\lg(n))$  with Binary Search**
- Inefficient insert/remove  
**Elements must be inserted and removed at the location dictated by the sort property, resulting shifting the array in memory – an  $O(n)$  operation**



# Array

## Unsorted Array

- Constant time add/remove at the beginning/end  
**Amortized  $O(1)$  insert and remove from the front and of the array**  
**Idea: Double on resize**
- Inefficient global search structure  
**With no sort property, all searches must iterate the entire array;  $O(1)$  time**



Array

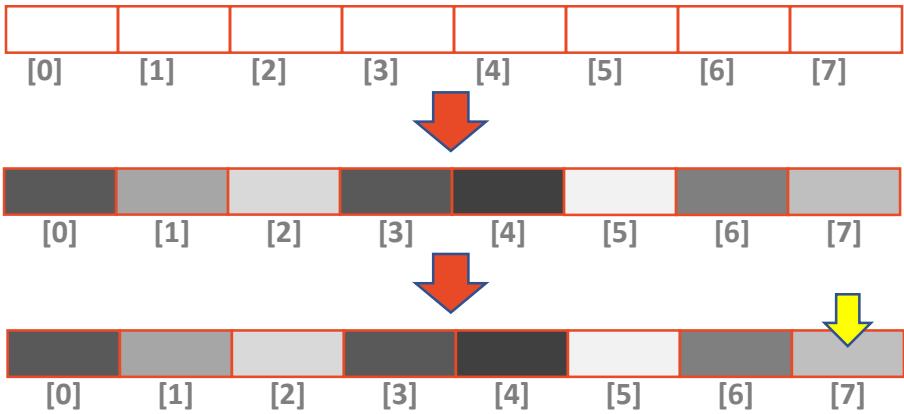


Unsorted Array



Queue (FIFO)

- First In First Out (FIFO) ordering of data  
**Maintains an arrival ordering of tasks, jobs, or data**
- All ADT operations are constant time operations  
**enqueue() and dequeue() both run in O(1) time**



Array

Unsorted Array

Stack (LIFO)

- Last In First Out (LIFO) ordering of data  
**Maintains a “most recently added” list of data**
- All ADT operations are constant time operations  
**push() and pop() both run in O(1) time**

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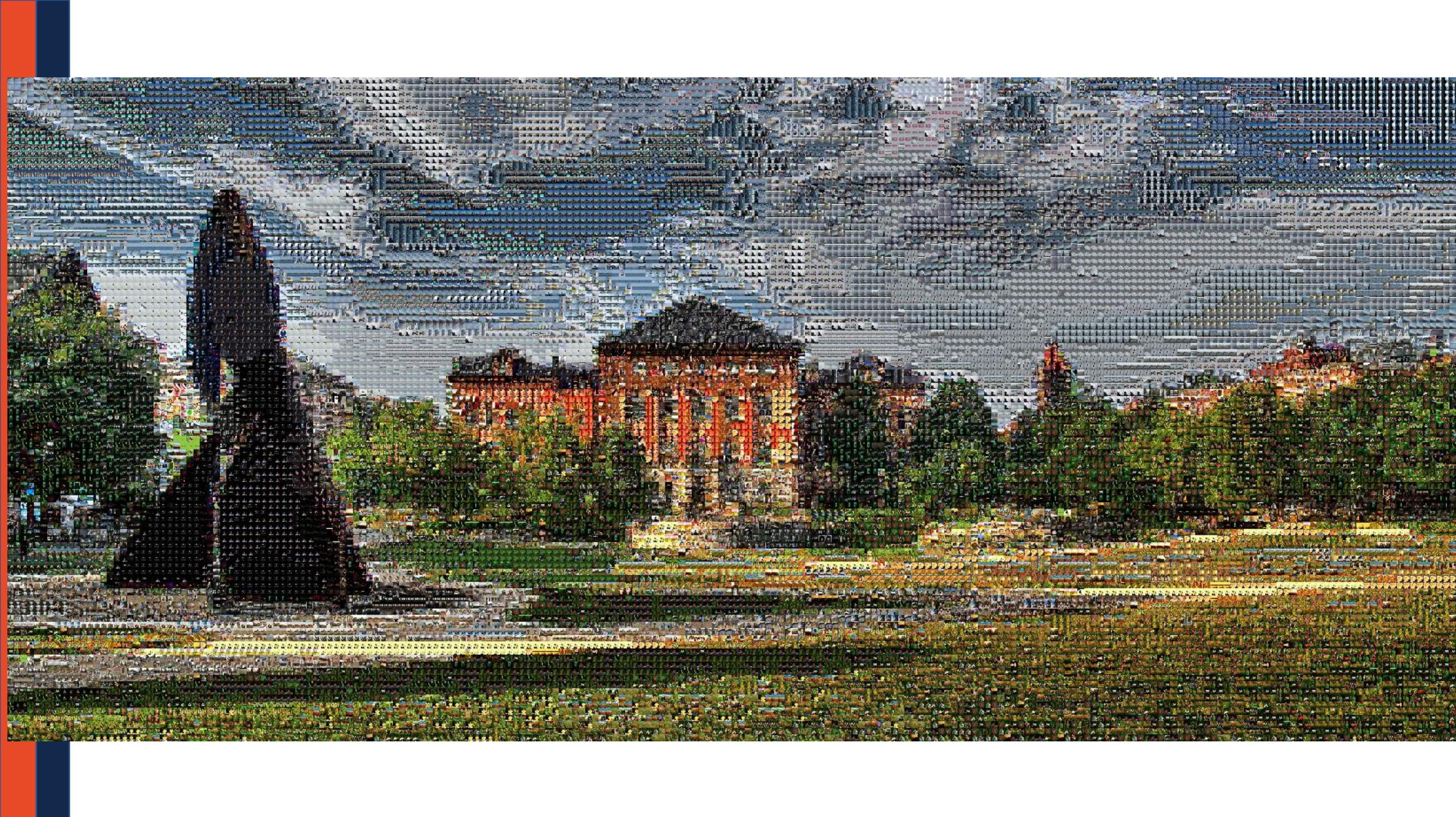
## Array

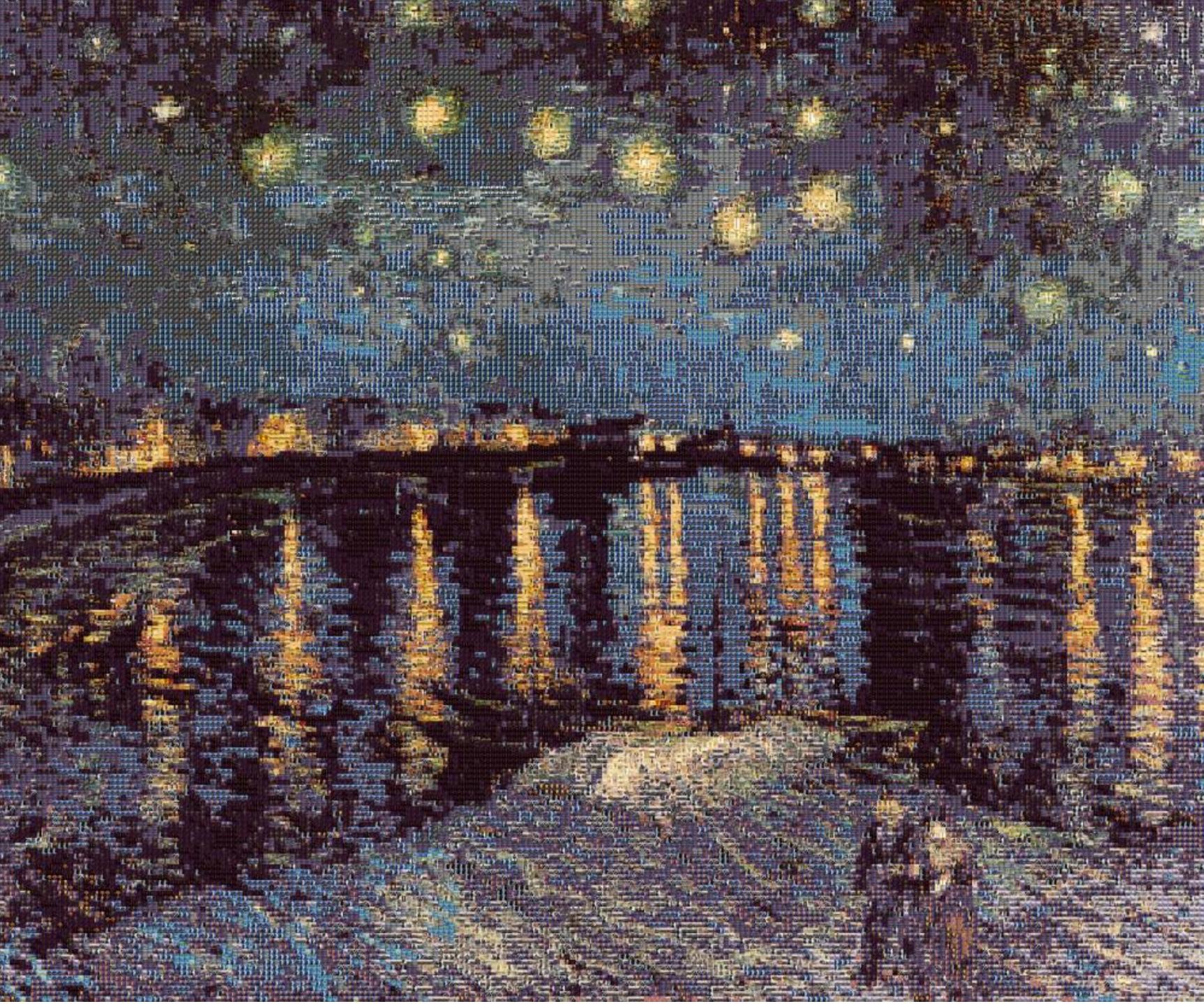
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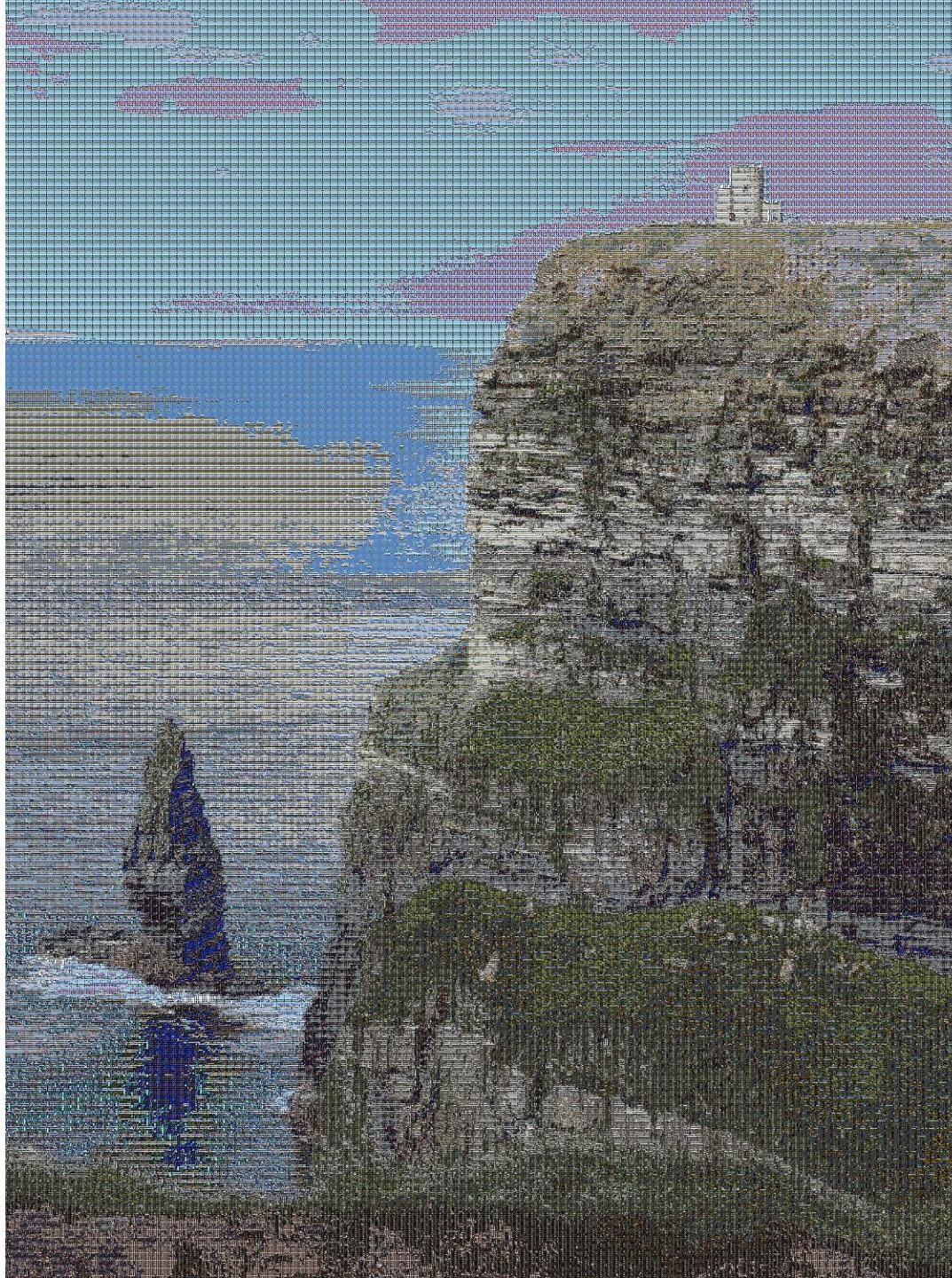
## Graphs

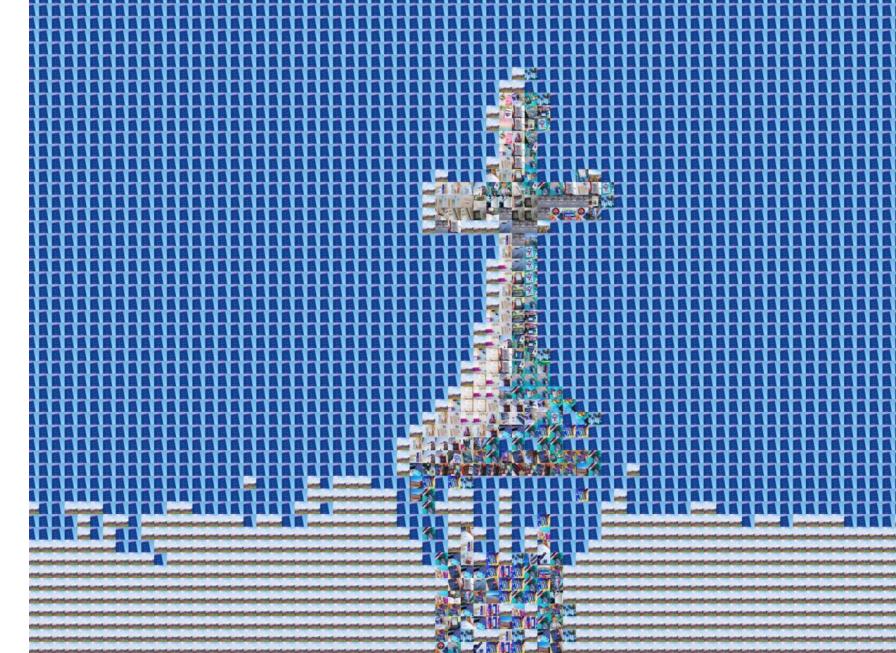
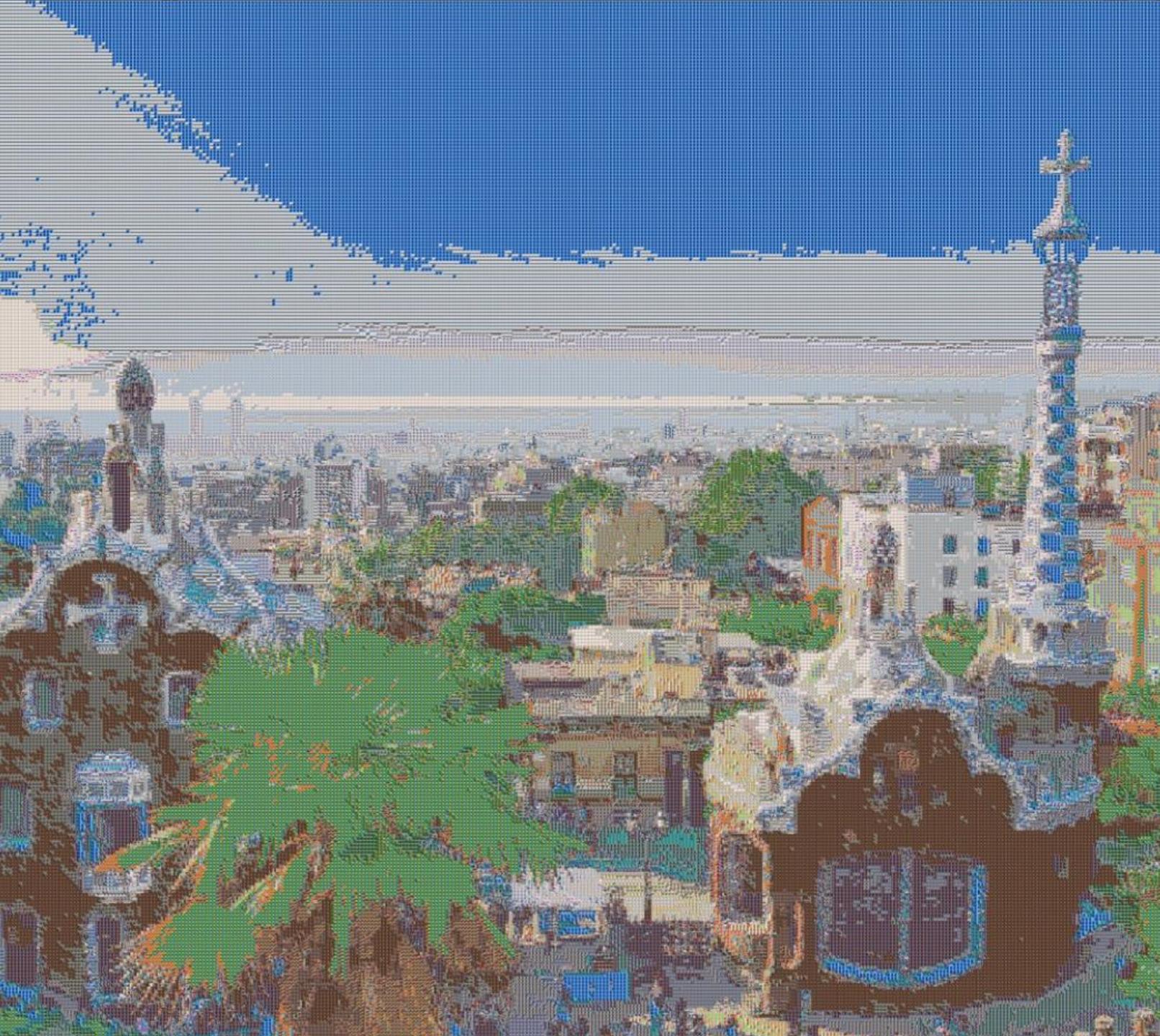


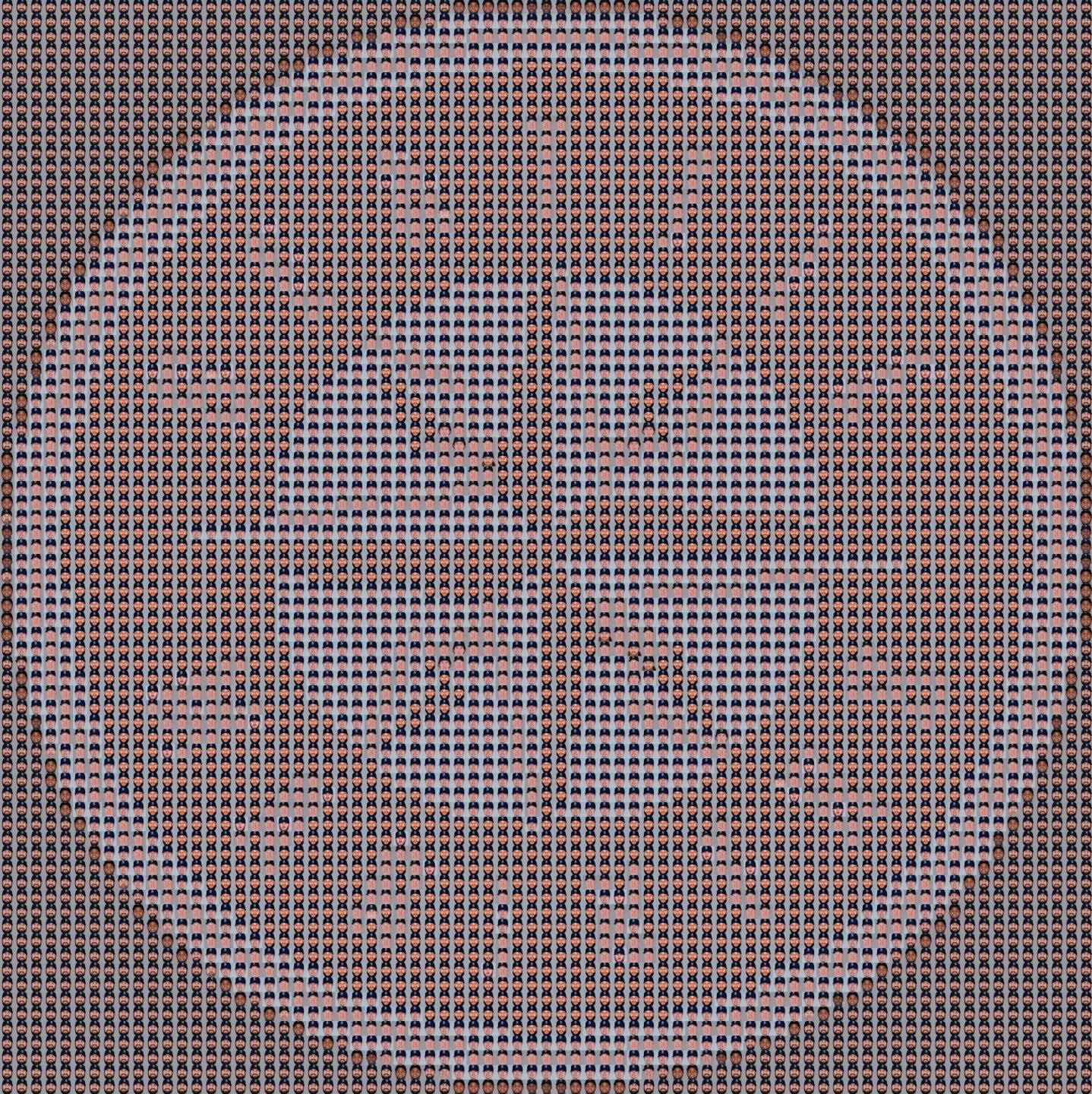


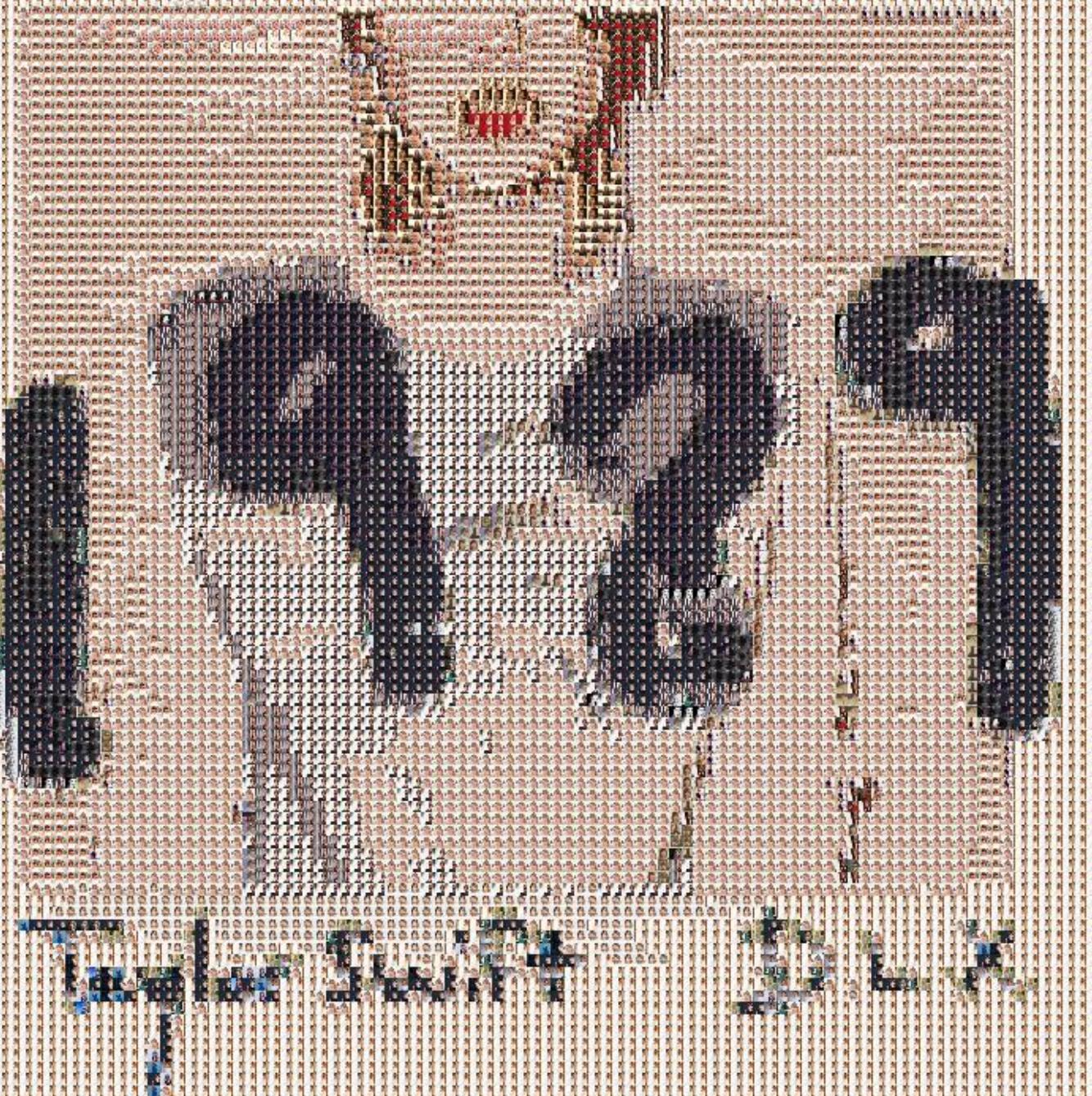






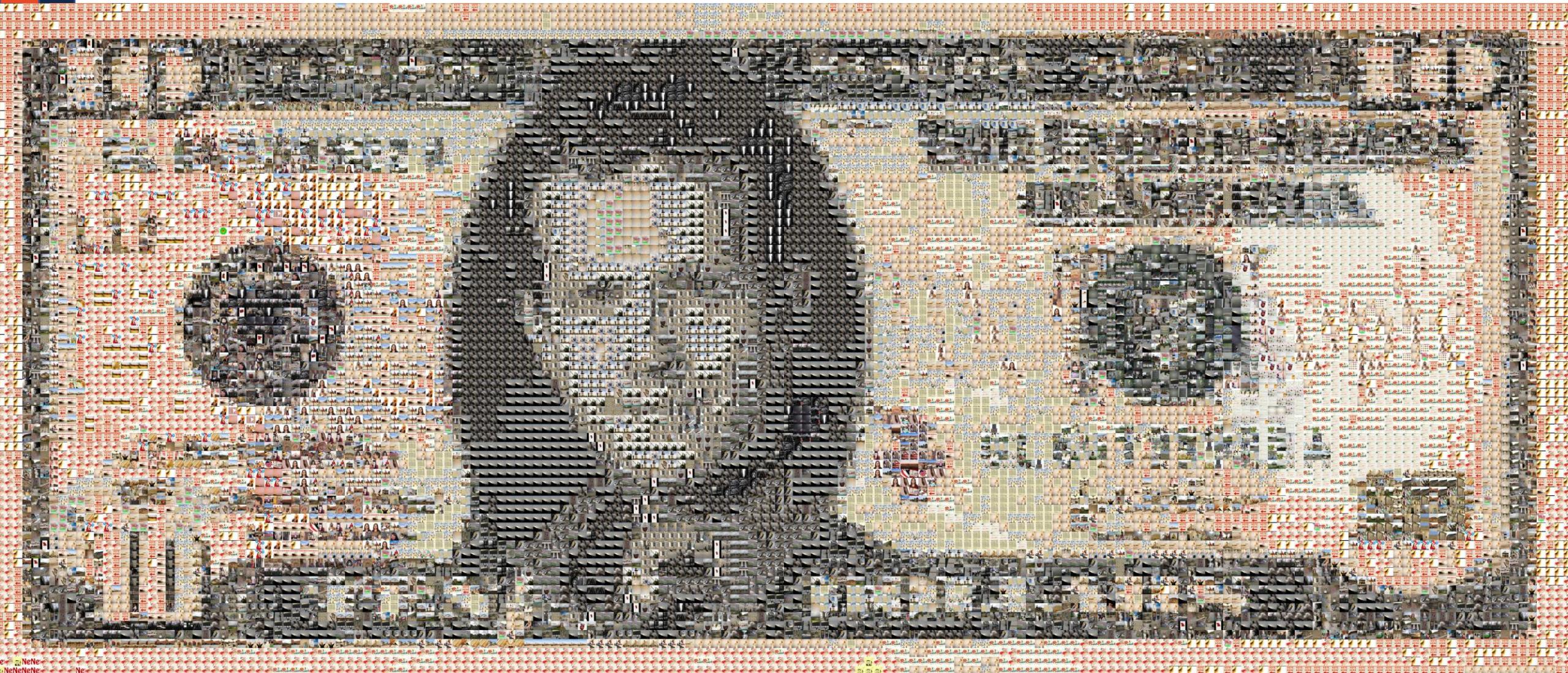


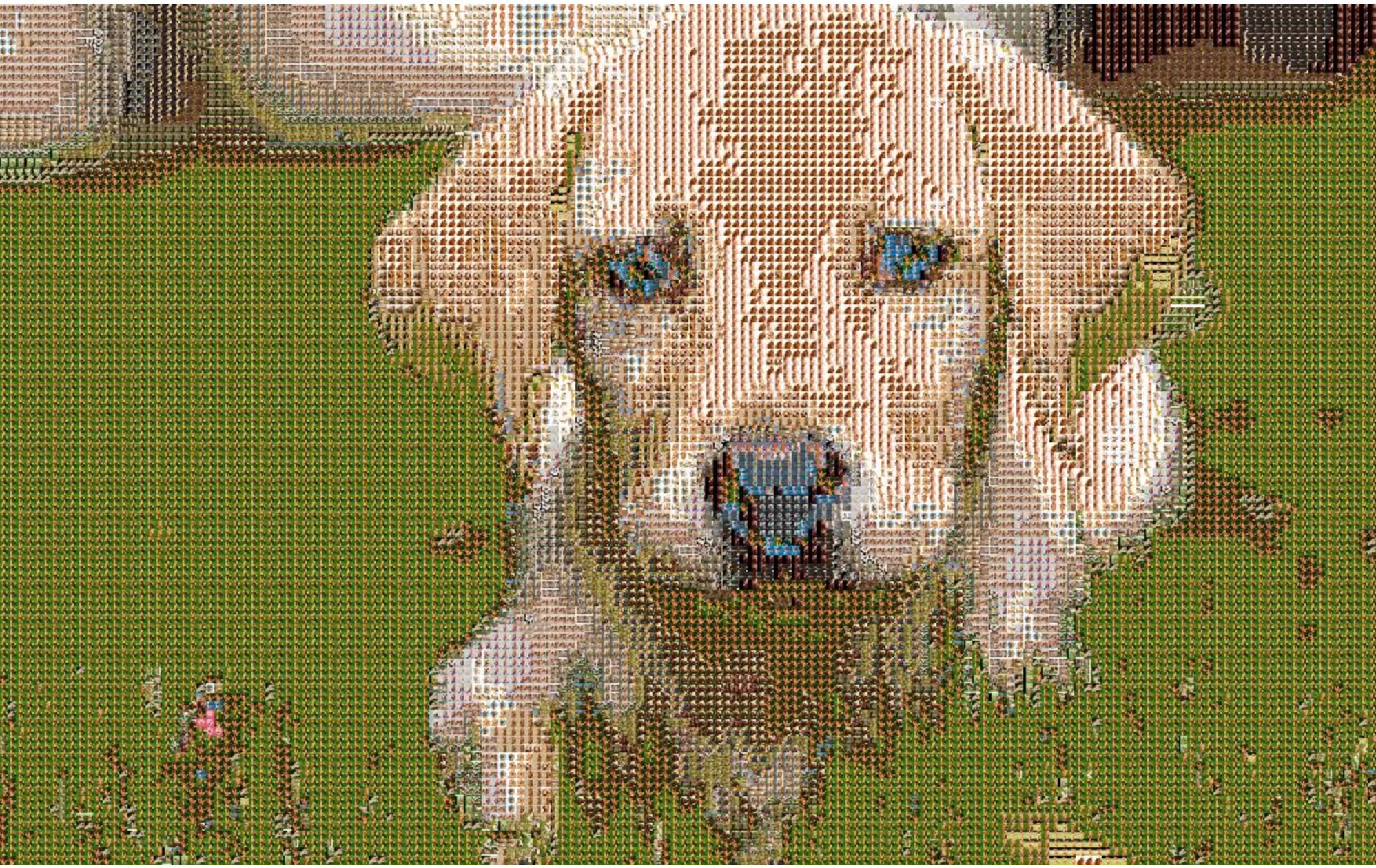






“When you're asked about kd-trees in an interview and Wade comes to mind:”





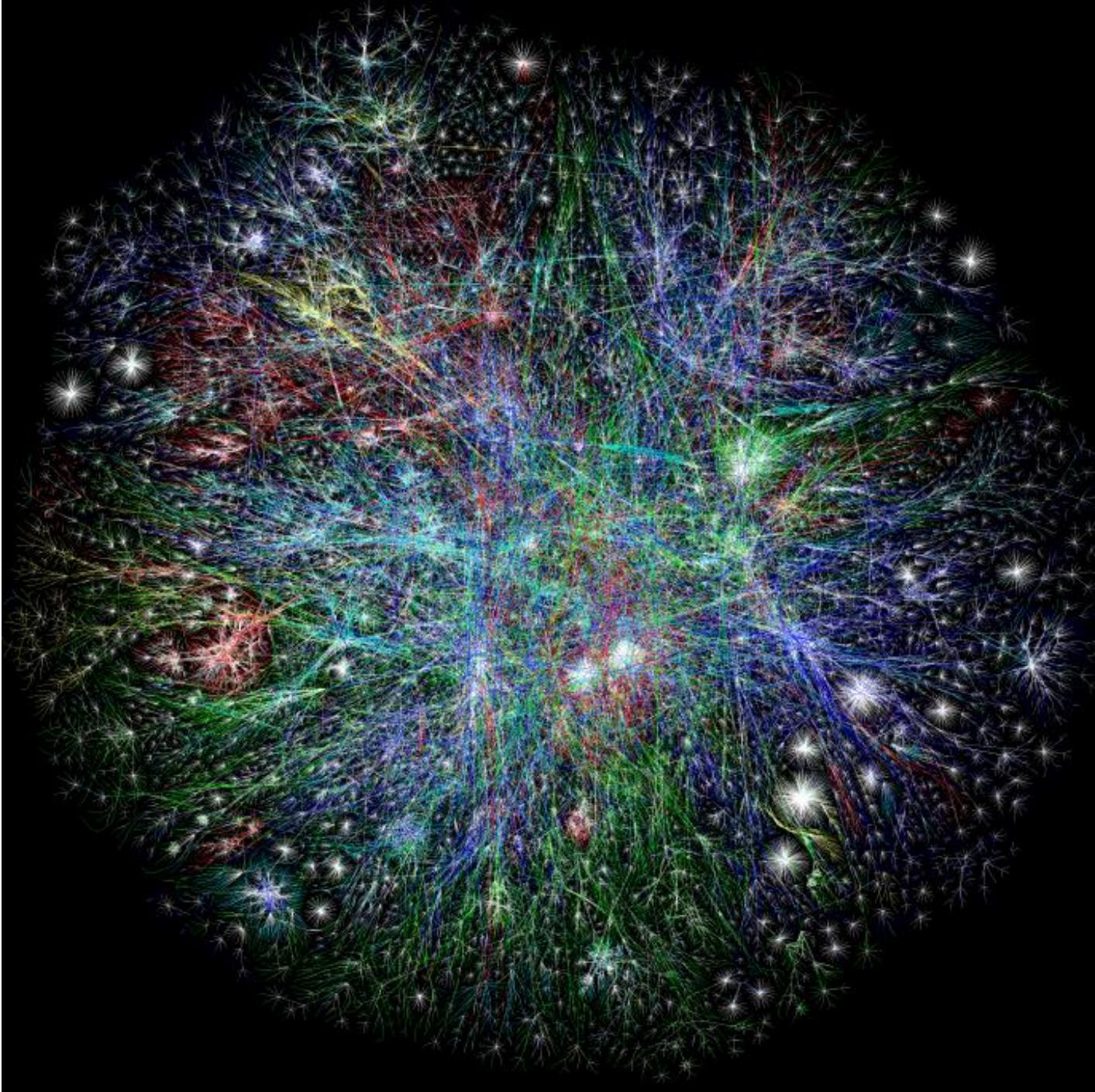


Aw, Snap!

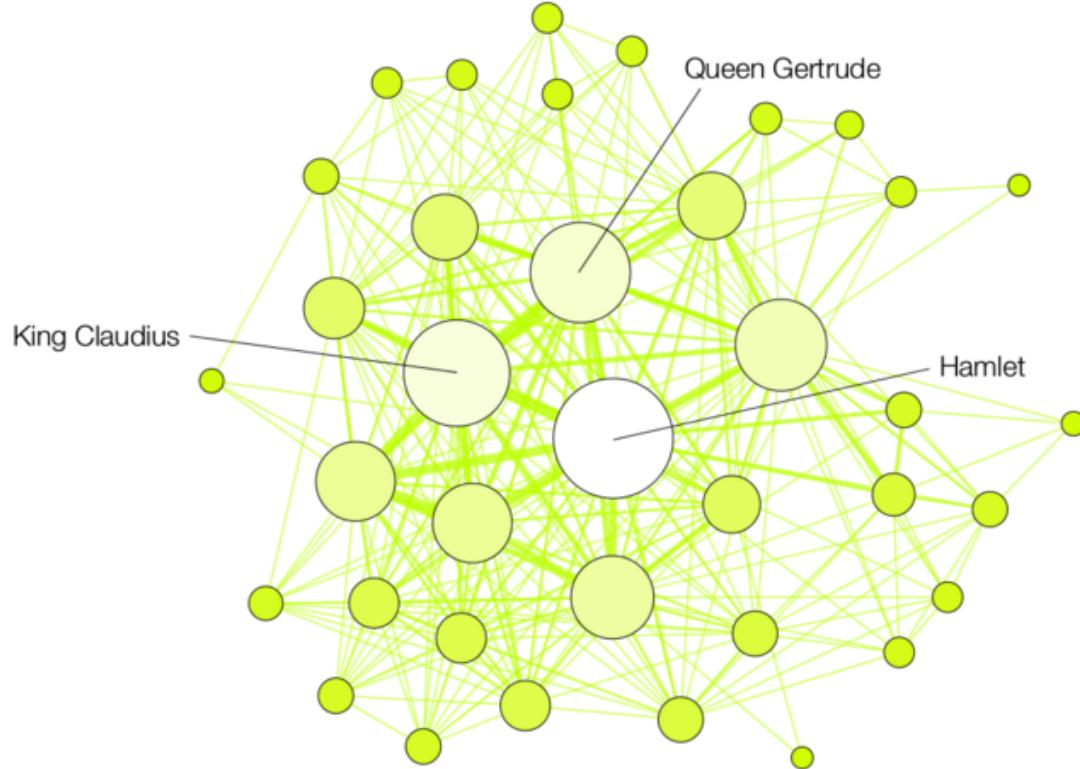
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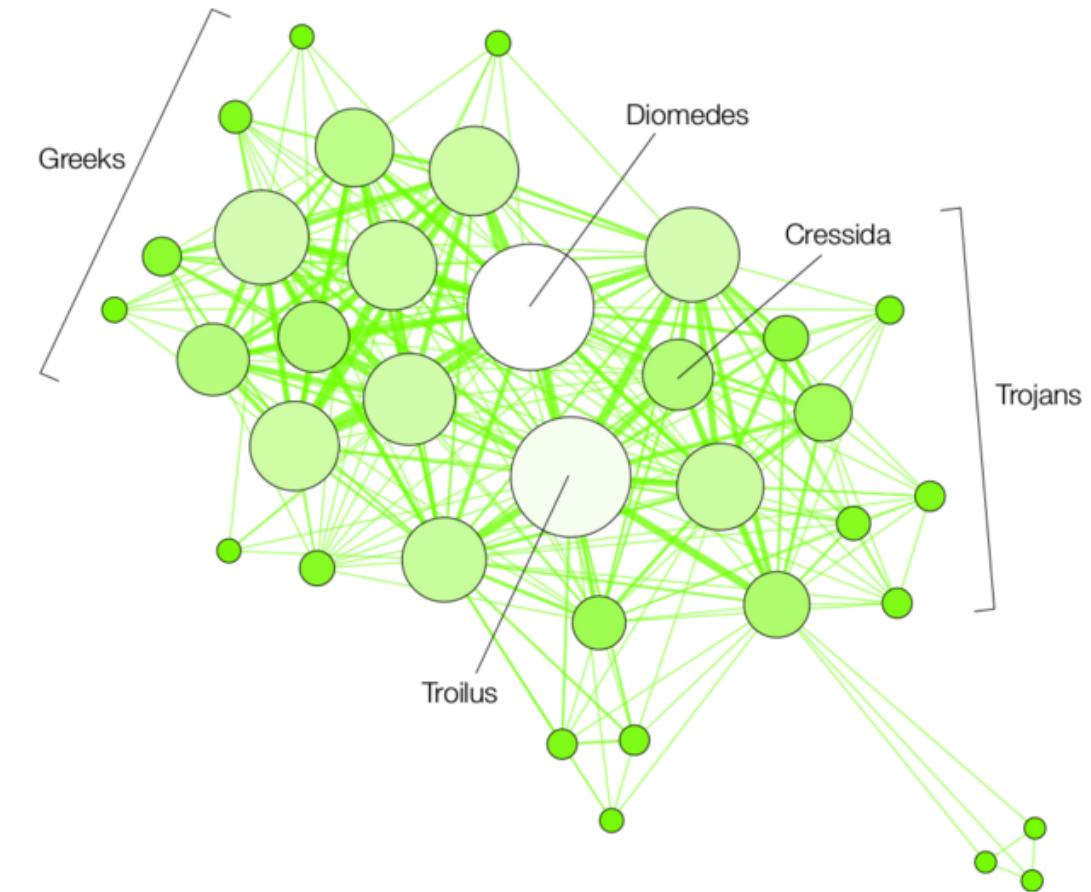
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**The Internet, 2003**  
*The OPTE Project (2003)*  
Map of the entire internet; nodes  
are routers; edges are connections.



## HAMLET

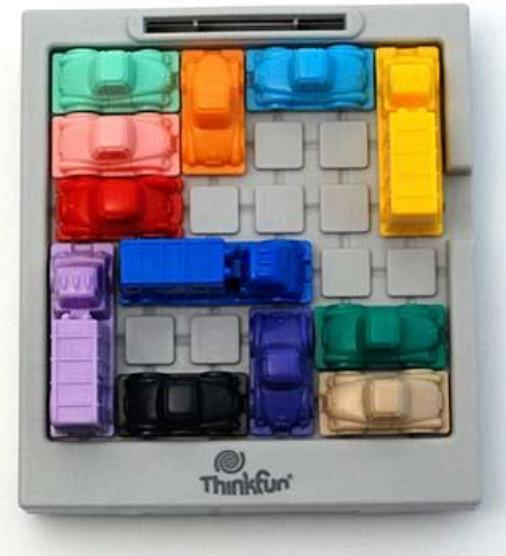
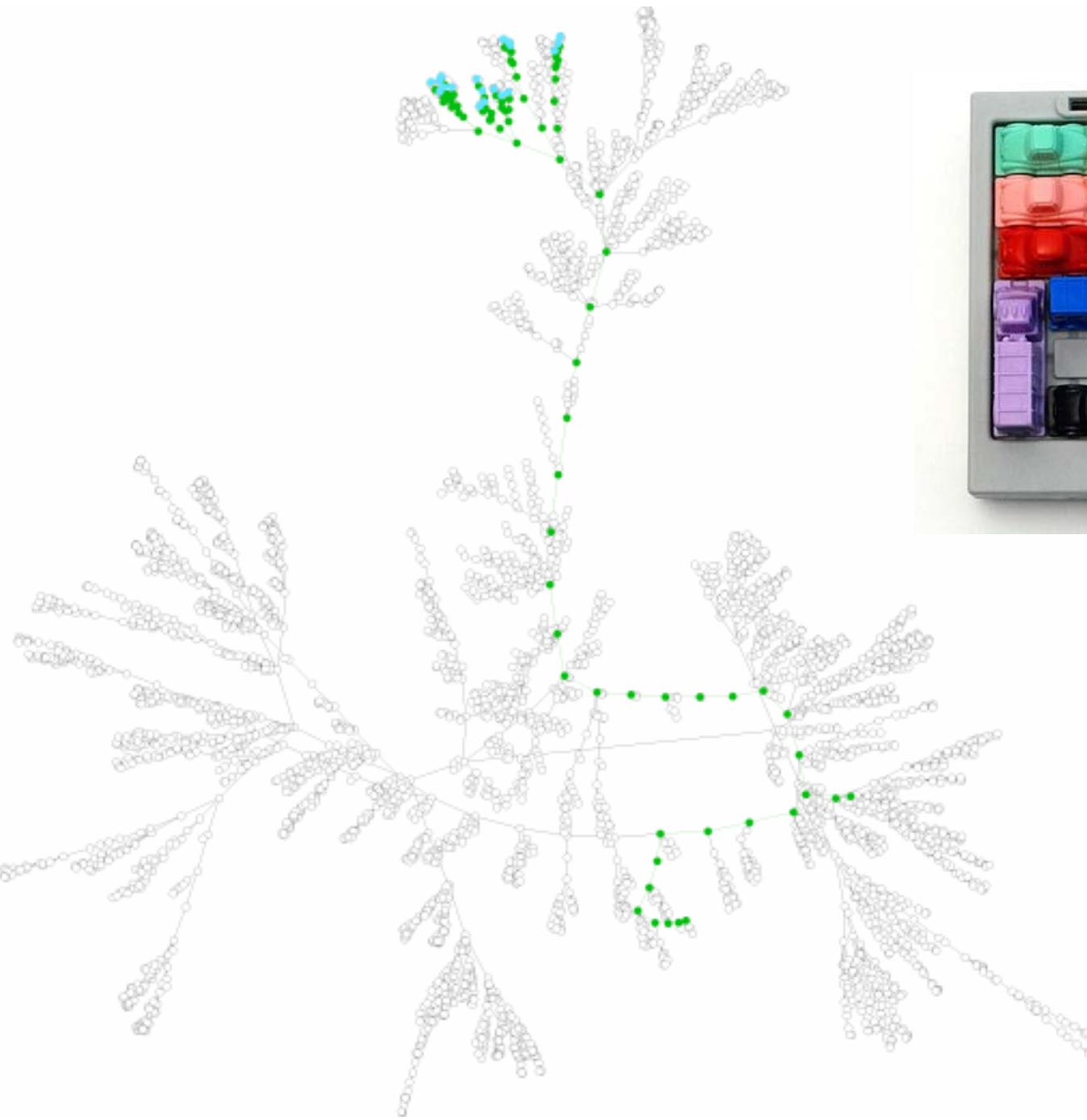


## TROILOUS AND CRESSIDA

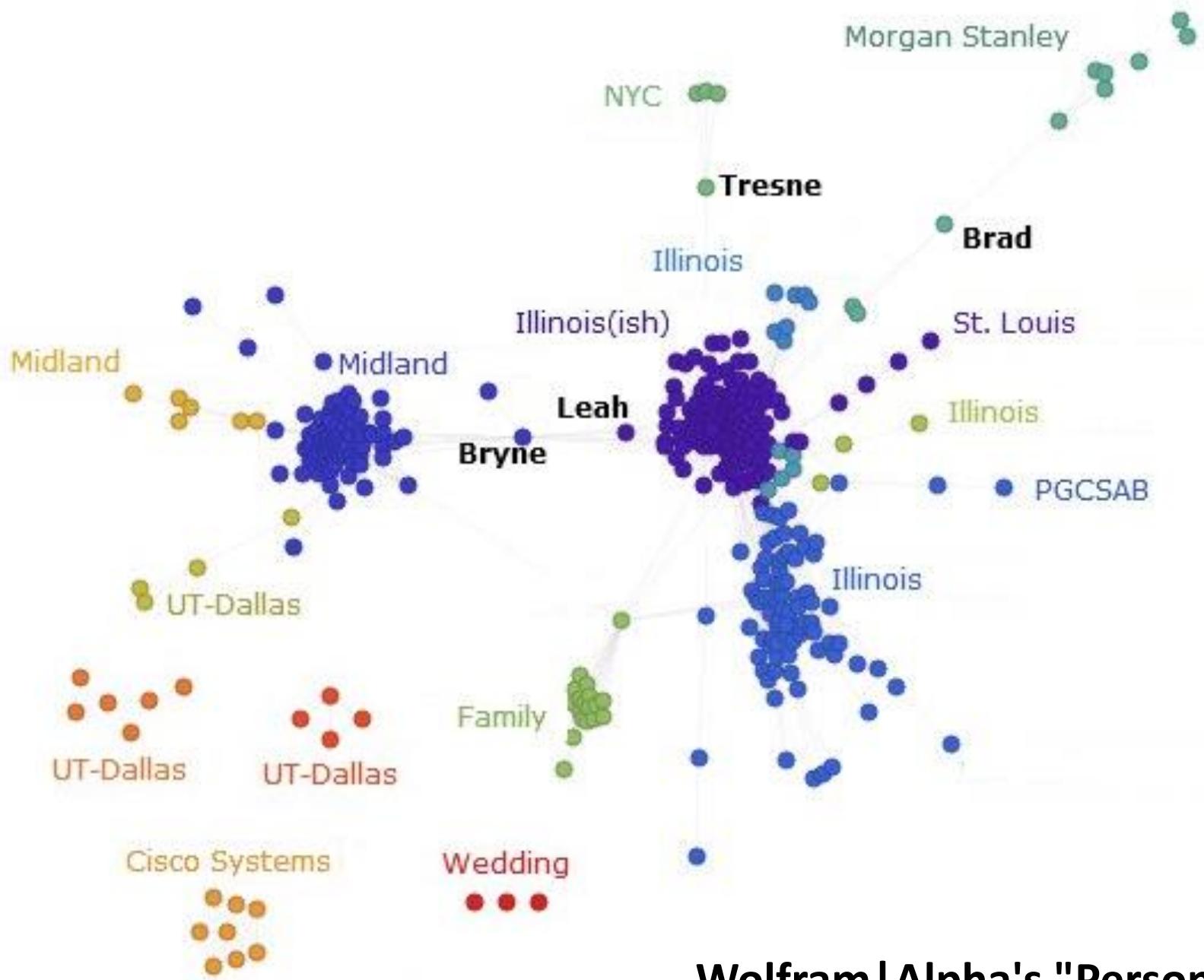
**Who's the real main character in Shakespearean tragedies?**

*Martin Grandjean (2016)*

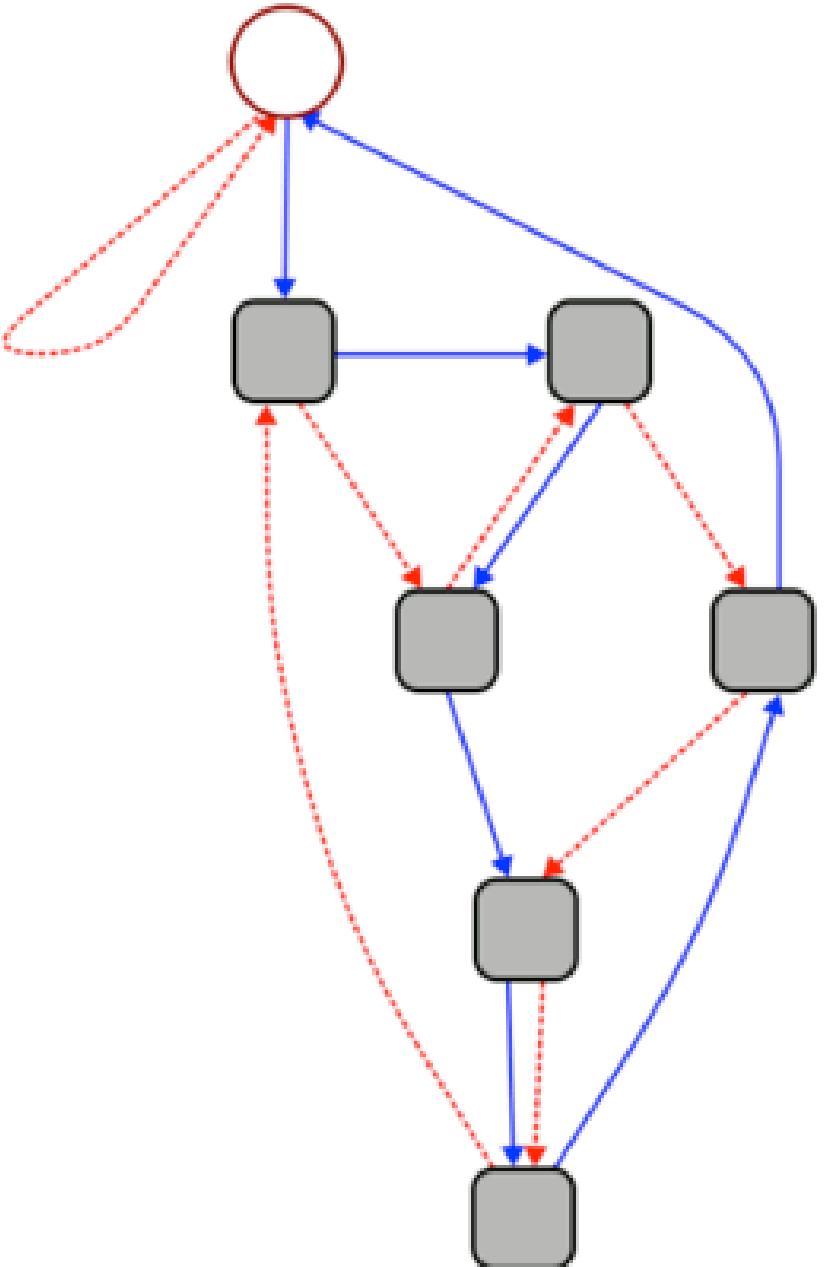
<https://www.pbs.org/newshour/arts/whos-the-real-main-character-in-shakespearean-tragedies-heres-what-the-data-say>



**“Rush Hour” Solution**  
*Unknown Source*  
*Presented by Cinda Heeren, 2016*



**Wolfram|Alpha's "Personal Analytics" for Facebook**  
Generated: April 2013 using Wade Fagen-Ulmschneider's Profile Data



This graph can be used to quickly calculate whether a given number is divisible by 7.

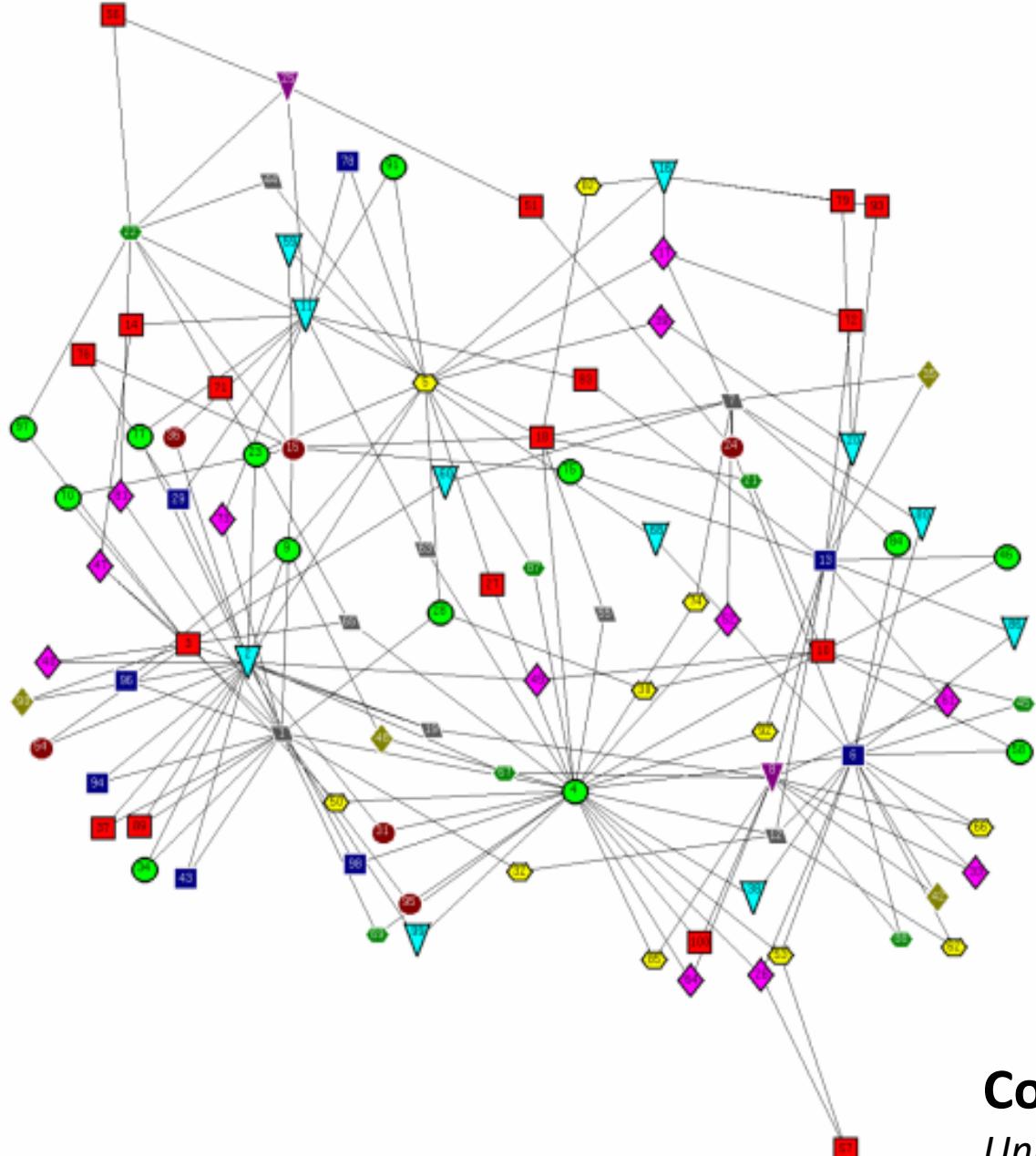
1. Start at the circle node at the top.
2. For each digit **d** in the given number, follow **d blue (solid) edges** in succession. As you move from one digit to the next, follow **1 red (dashed) edge**.
3. If you end up back at the circle node, your number is divisible by 7.

3703

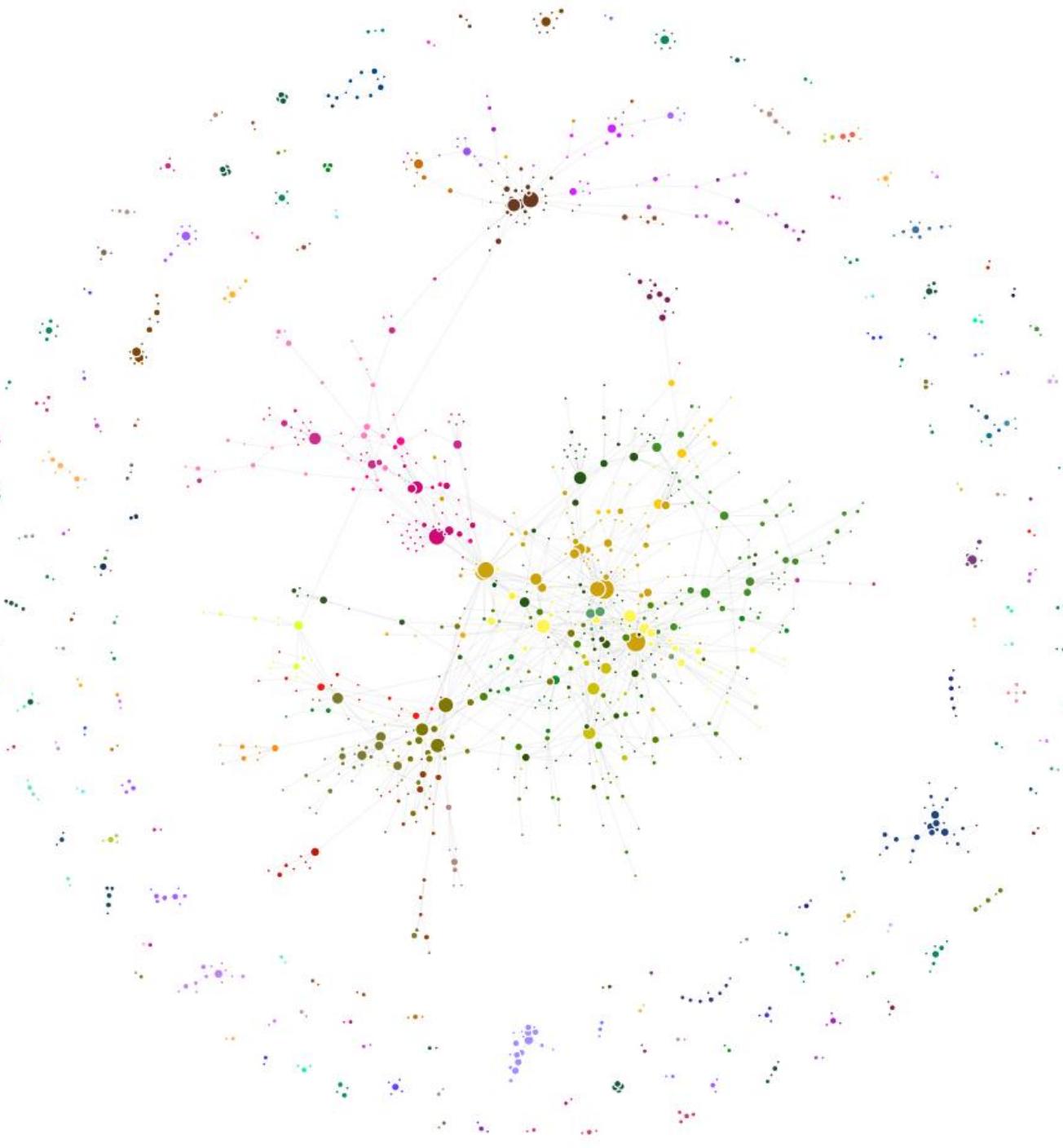
**“Rule of 7”**

*Unknown Source*

*Presented by Cinda Heeren, 2016*



**Conflict-Free Final Exam Scheduling Graph**  
*Unknown Source*  
Presented by Cinda Heeren, 2016

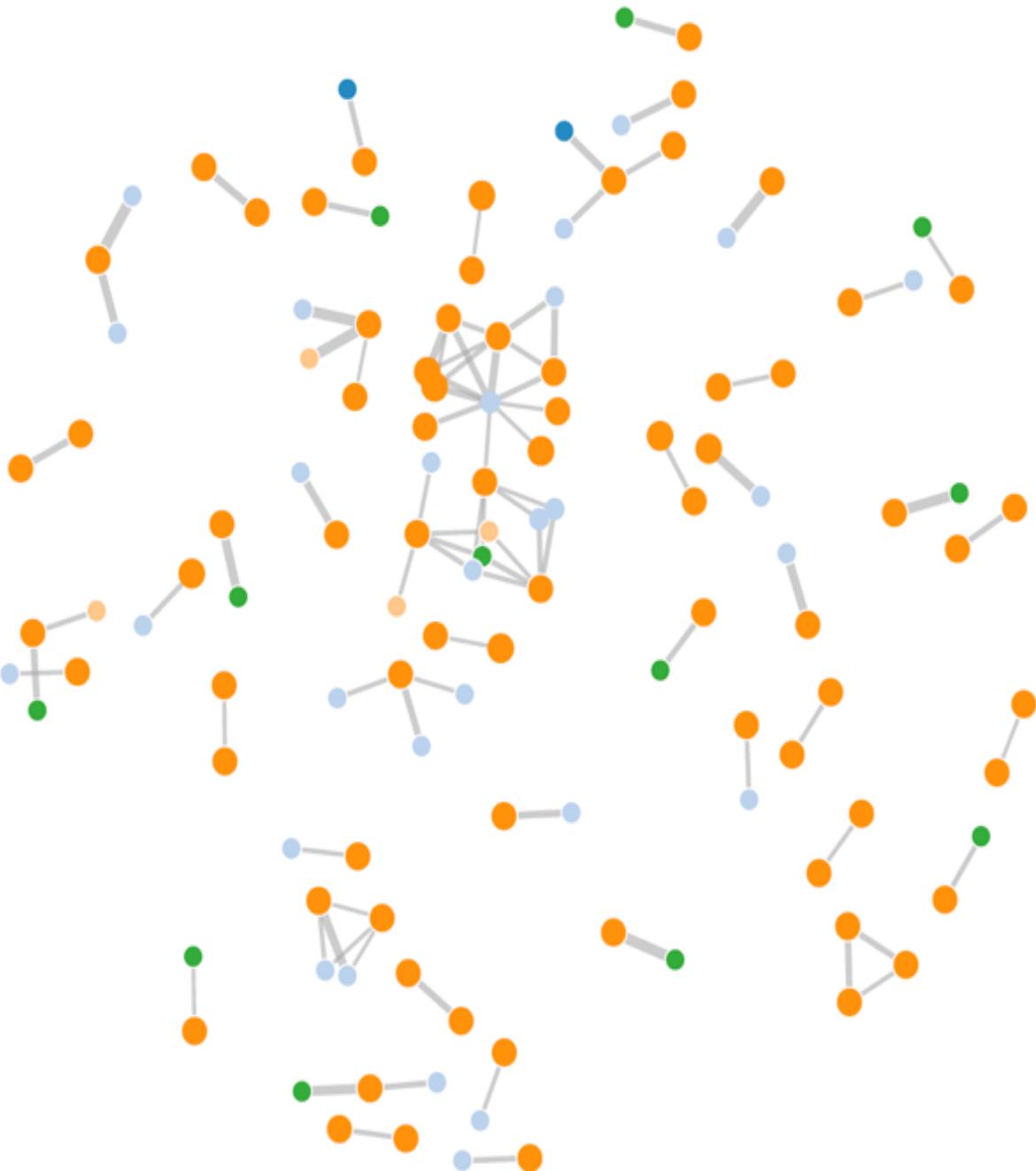


## Class Hierarchy At University of Illinois Urbana-Champaign

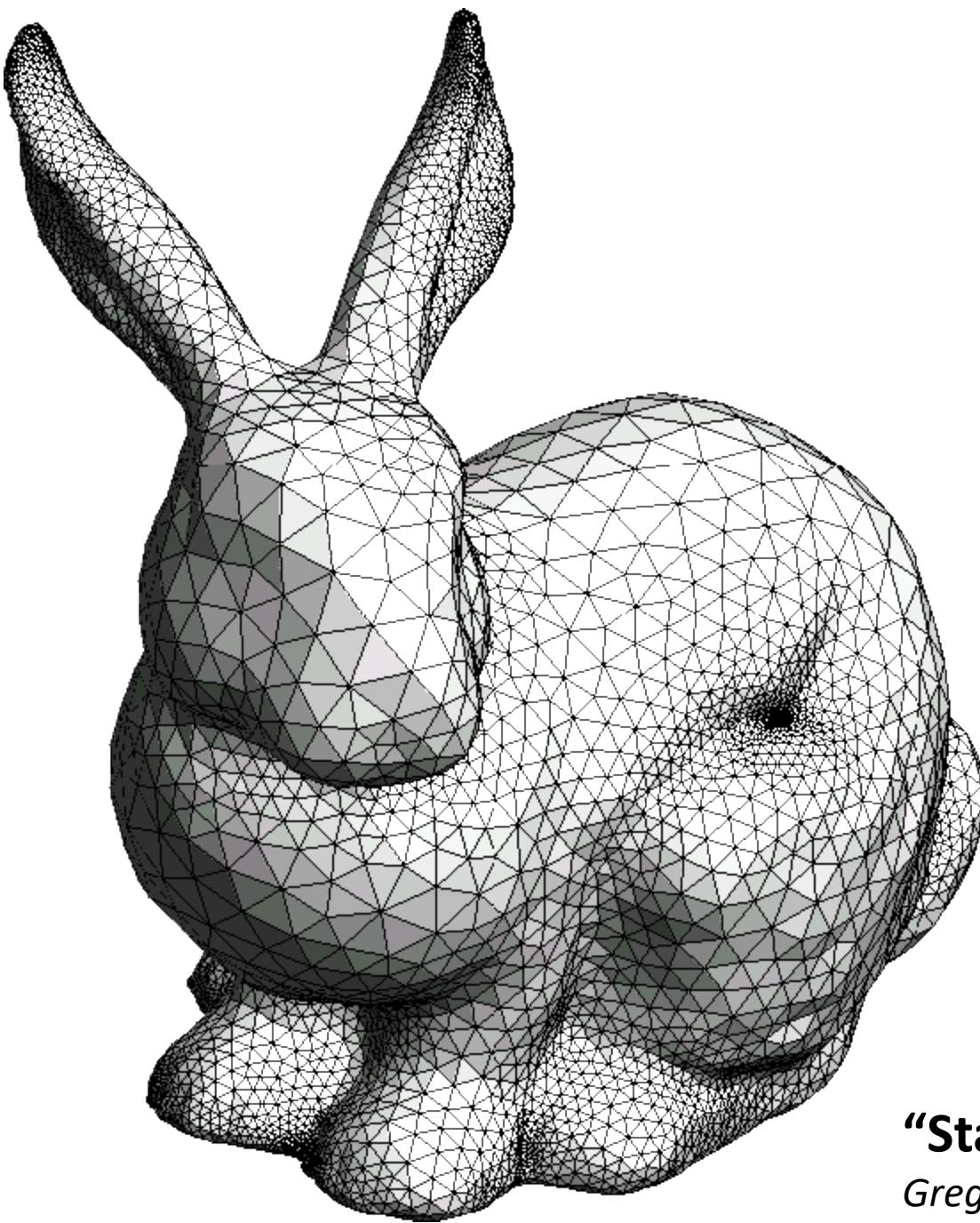
*A. Mori, W. Fagen-Ulmschneider, C. Heeren*

Graph of every course at UIUC; nodes are courses, edges are prerequisites

[http://waf.cs.illinois.edu/discovery/class\\_hierarchy\\_at\\_ilinois/](http://waf.cs.illinois.edu/discovery/class_hierarchy_at_ilinois/)

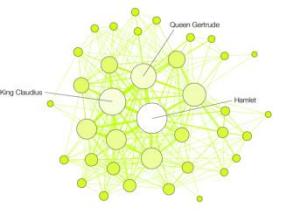
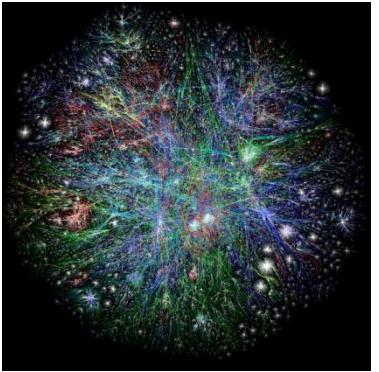
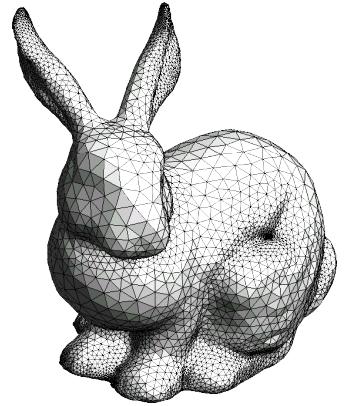


**MP Collaborations in CS 225**  
*Unknown Source*  
*Presented by Cinda Heeren, 2016*

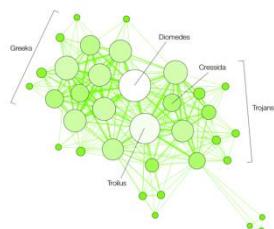


**“Stanford Bunny”**  
*Greg Turk and Mark Levoy (1994)*

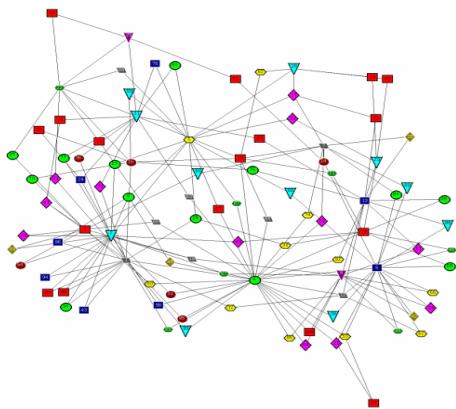
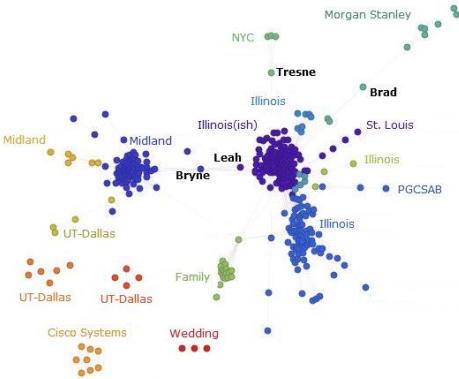
# Graphs



HAMLET

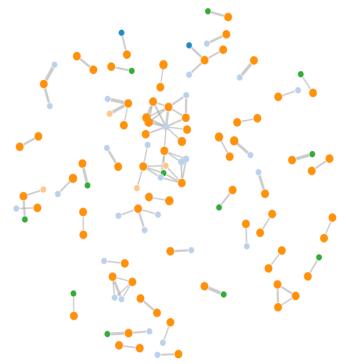
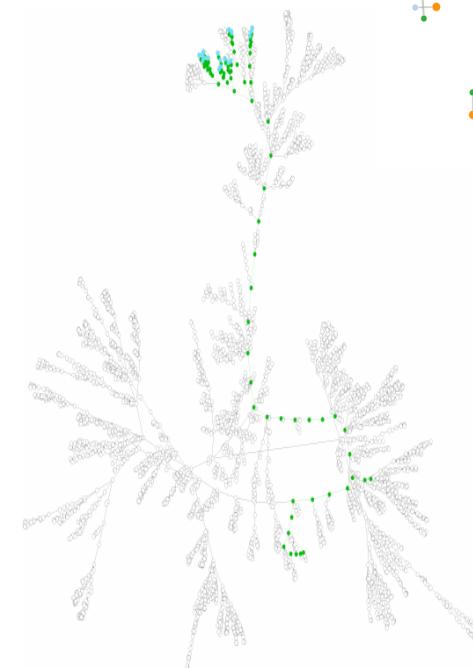
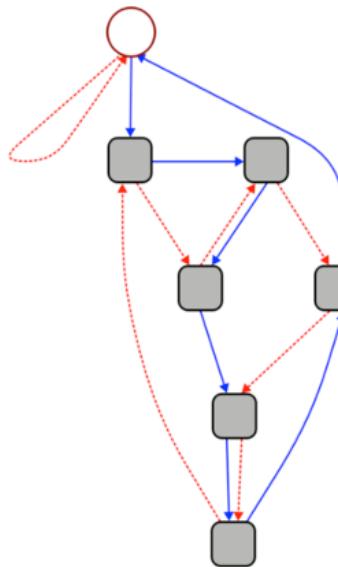


TROILUS AND CRESSIDA



To study all of these structures:

1. A common vocabulary
2. Graph implementations
3. Graph traversals
4. Graph algorithms

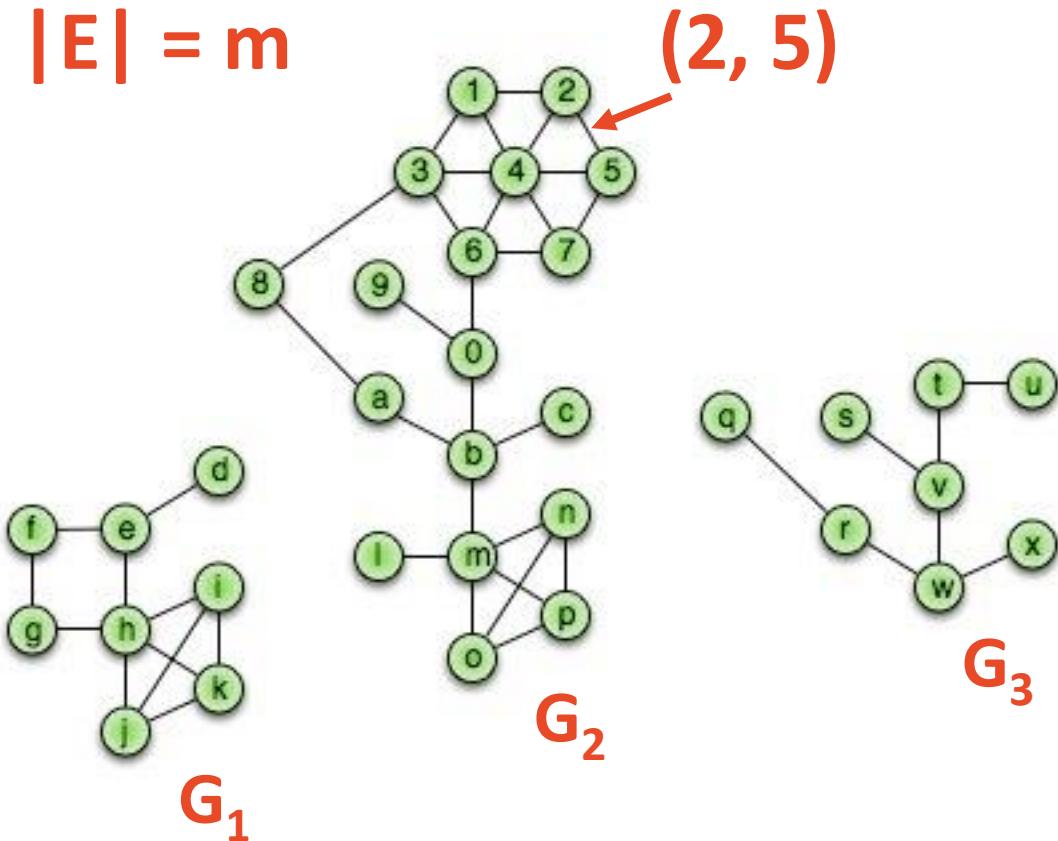


# Graph Vocabulary

$G = (V, E)$

$|V| = n$

$|E| = m$



Incident Edges:  
 $I(v) = \{ (x, v) \text{ in } E \}$

Degree( $v$ ):  $|I|$

Adjacent Vertices:  
 $A(v) = \{ x : (x, v) \text{ in } E \}$

Path( $G_2$ ): Sequence of vertices connected by edges

Cycle( $G_1$ ): Path with a common begin and end vertex.

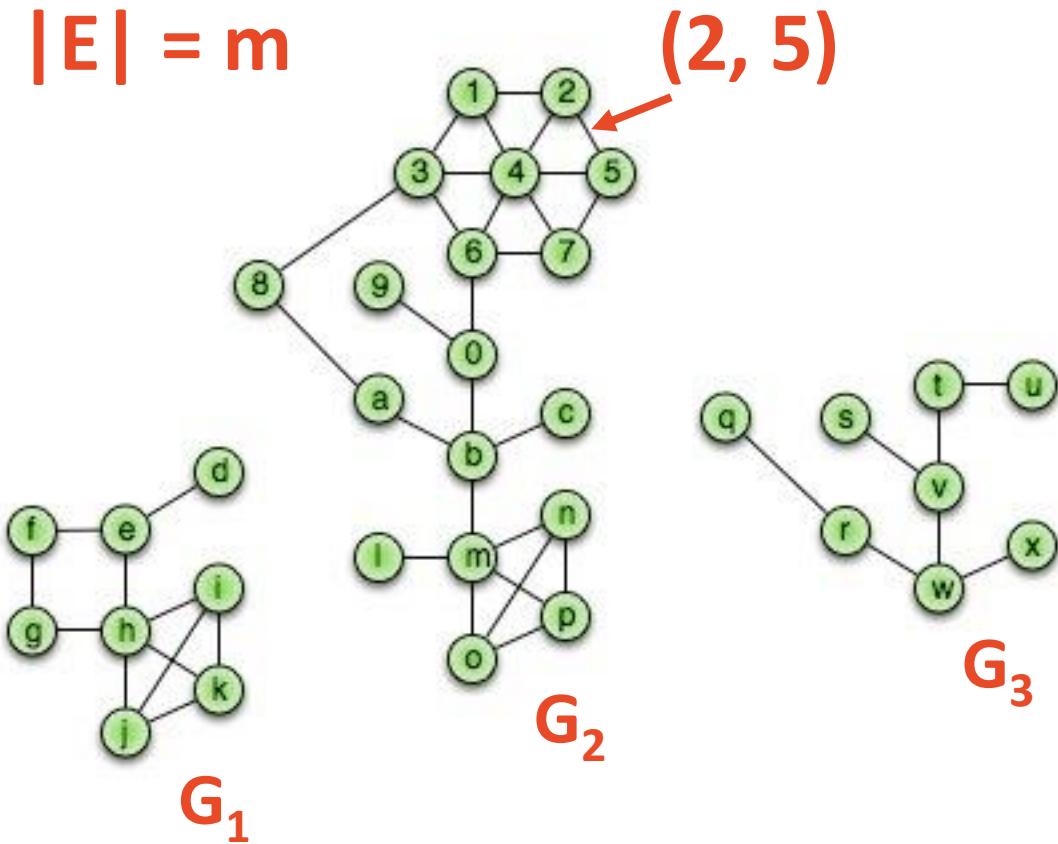
Simple Graph( $G$ ): A graph with no self loops or multi-edges.

# Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$



**Subgraph( $G$ ):**

$G' = (V', E')$ :

$V' \subseteq V, E' \subseteq E$ , and  
 $(u, v) \in E \rightarrow u \in V', v \in V'$

**Complete subgraph( $G$ )**

**Connected subgraph( $G$ )**

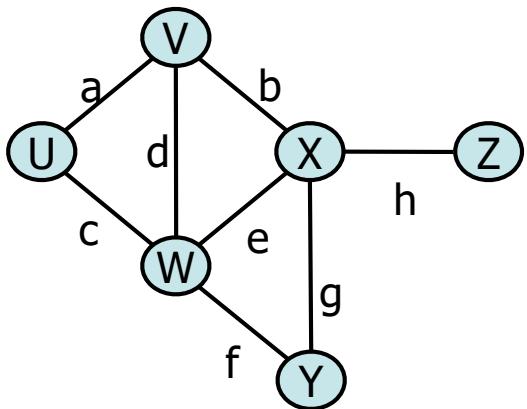
**Connected component( $G$ )**

**Acyclic subgraph( $G$ )**

**Spanning tree( $G$ )**

Running times are often reported by  $n$ , the number of vertices, but often depend on  $m$ , the number of edges.

How many edges? **Minimum edges:**  
Not Connected:



Connected\*:

**Maximum edges:**  
Simple:

Not simple:

$$\sum_{v \in V} \deg(v) =$$