

Logic Gates and Circuits

CSCB58: Computer Organization

Lecture 1: January 13th, 2023



Computer Science
UNIVERSITY OF TORONTO

A little bit about your instructor

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Department at UofT and UTSC, since 2020

■ **PhD from Carnegie Mellon University**

■ **Worked at Intel, AMD, Microsoft, and Nvidia**

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■ **My research group: embARC Lab (www.embarclab.com)**

■ **Our research: computer systems + architecture**

■ **Cross-layer spanning applications, systems, and hardware**

■ **Application-specific: E.g., machine learning, graphics, HPC, robotics**



Say hi to your awesome TAs!

Andrew



Ryan



Klein



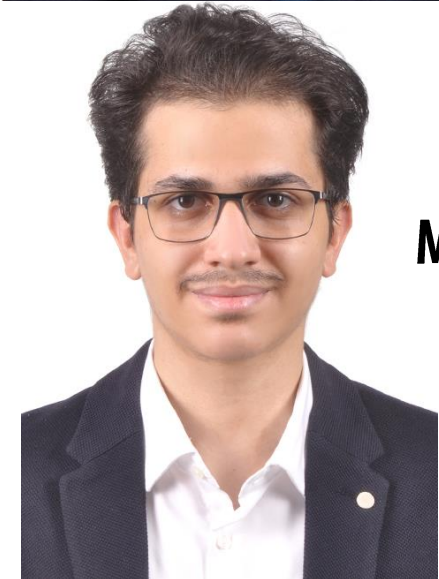
Yuanyuan



Sung



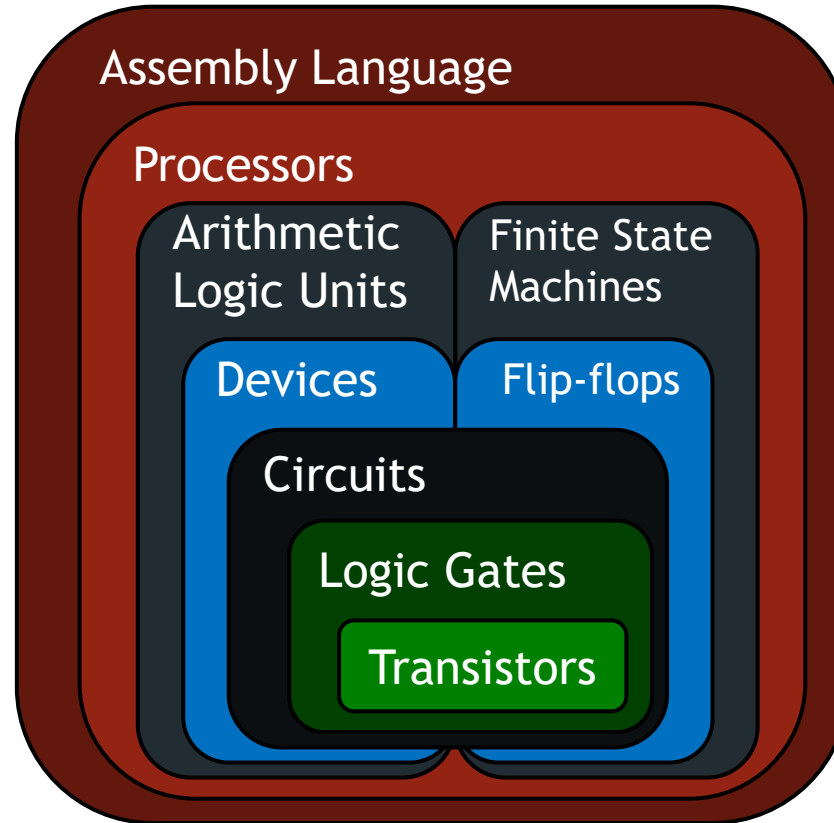
Mohannad



Today!

- **Why take B58?**
- **What's in B58**
- **How the course works?**
- **Logic gates and circuits**

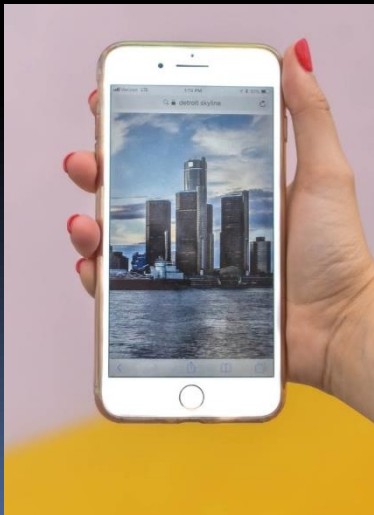
The course at a glance...



CSCB58 Is Intense

- **There is a lot to learn, practice, and do.**
- **You have tasks every week.**
- **It is historically one of the busiest courses for students.**
- **But there is good news:**
 - **Material not hard to understand, but it does require practice and effort.**
 - **All this practice makes you really good.**
 - **Course grade tend to skew high**

What is a computer?



do you need
electricity for a
computer?

CSCB58 Asks the Big Questions

- What is a computer?
- What is memory? How does a computer store information
- Why do computers work in binary?
- how does a computer actually... compute?
- How does the code we write in Python/Java/C++/etc actually translate into things happening?

Why take CSCB58?

- To better understand computers!
- See what's going on "under the hood"
- Open the black box, get rid of the mystery
- Understand the whole pipeline, from atoms to assembly
 - Everything above assembly is virtualization and abstraction
 - Everything else is an **illusion!**
- Build a cool software-based project!

CSCB58 Course Goals

- Learn to **build computers...**
 - Understand and design the underlying architecture of computer systems.
- ... and **how they work.**
 - How programs use digital structures to do computation.
- Learn **engineering.**
 - Build systems from components, understand abstractions, make tradeoffs.
- Learn **there is no magic!**
 - There is nothing you cannot understand.

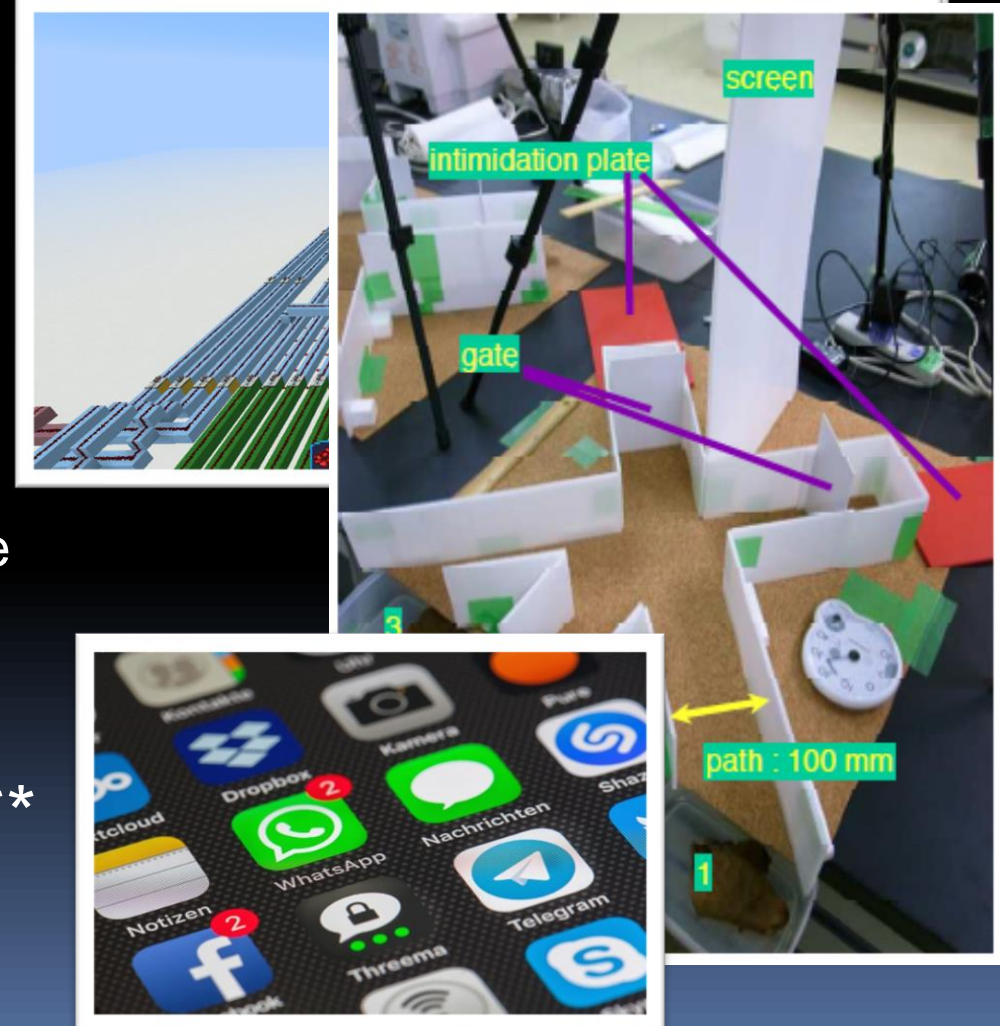
CSCB58 Course Goals

- By the end of this course, you will be able to build a computer from the atomic level upwards*
 - Build circuits from logic gates
 - Build memory/computational units from circuits
 - Build a computer from memory/computational units

(* Given an infinite amount of time, and the ability to manipulate individual atoms.)

Build a Computer!

- ...from LEGO
- ...in Minecraft
- ...from living crabs*
 - * Do not build a computer from live crabs. It is unethical.
- ...inside an iMessage to break into iPhones**
 - ** also unethical.



Some Admin



<https://villains.fandom.com/wiki/Vogons>

Administrivia

- **This course is on Quercus: q.utoronto.ca**
 - Check it regularly!
- **Get help on Piazza**
 - Use for all assignment or technical questions.
 - Read pinned posts (first posts at the top) before asking questions!
 - Private posts for questions relating to you only
- **E-mail (use official email only!)**
 - Only for emergencies and confidential personal matters
- **Office hours**
 - Dedicated lab office hours by TAs (for lecture material, labs, and assignments)
 - Instructor office hours (for lecture material and personal situations)

Course Components

- **Lectures**
 - 2 hour in person lectures every week on Friday
 - No TUT sessions
- **Quizzes** **10%**
 - 11 short weekly quizzes on Mondays (starting next week)
- **Labs** **40%**
 - 5-6 (TBD) labs starting this week (first is very easy, don't worry)
- **Project** **17%**
 - 3-week assembly project
- **Final exam** **33%**
 - Final exam - must get 30% to pass the course.

Administrivia

- **Some general principles for how we will operate:**
- **No exceptions other than AccessAbility and for medical emergencies**
- **We want consistent and fair policies for all!**
- **Course requirements: All requirements and prerequisites must be met**

Labs Marking

- **Submit each lab solution to Quercus by Monday, 6pm of following week.**
 - **Usually a Logisim file, sometimes also a PDF.**
 - **Late? You will lose marks.**
- **Each students has weekly 10-minute interview slot with your TA to demonstrate:**
 - **... your solution.**
 - **... basic understanding of the material.**
(by answering oral questions from the TA).
 - **Your mark is based on both!**

Missed Term Work

- **Life happens. We understand.**
- **We give you justification-free tokens.**
- **Can miss/do badly on 2 quizzes and 1 labs.**
 - **For any reason. We don't need to know.**
- **Works automatically, no need to tell us.**
 - **We drop lowest 2 quizzes and lowest 1 labs**
 - **You just focus on getting better.**
- **Do not waste or abuse them!**
 - **We will not give extra “tokens”, no matter what.**
 - **Best strategy: show up and do your best.**
- **No make up for late work. Late = missed.**

How to Succeed in Labs?

- **The labs are not hard, but they require work.**
- **Labs are your chance to practice.**
Do the work yourself.
 - **Submit something meaningful and on time.**
- **Don't plagiarize**
 - **You won't be able to answer TA questions...**
 - **... or answer the next quiz ...**
 - **... or the final exam.**

Lab Software

- Lab 1-3 use Logisim-evolution for hardware.
 - Lab 1 will tell you all you need to know.
 - Do not use the original Logisim or its other variations of it. Use the in the lab PDF file or the link above.
- Labs 4-5 will use MARS simulator for MIPS assembly.
 - Same idea: link will be included.

Project

- **Multi-week assembly programming project.**
- **Worth 17%**
- **At the end of the course.**
- **Done individually.**
- **Marking based on:**
 - **One checkpoint by TAs in final week**
 - **Short project report and video.**
- **Details will follow later.**

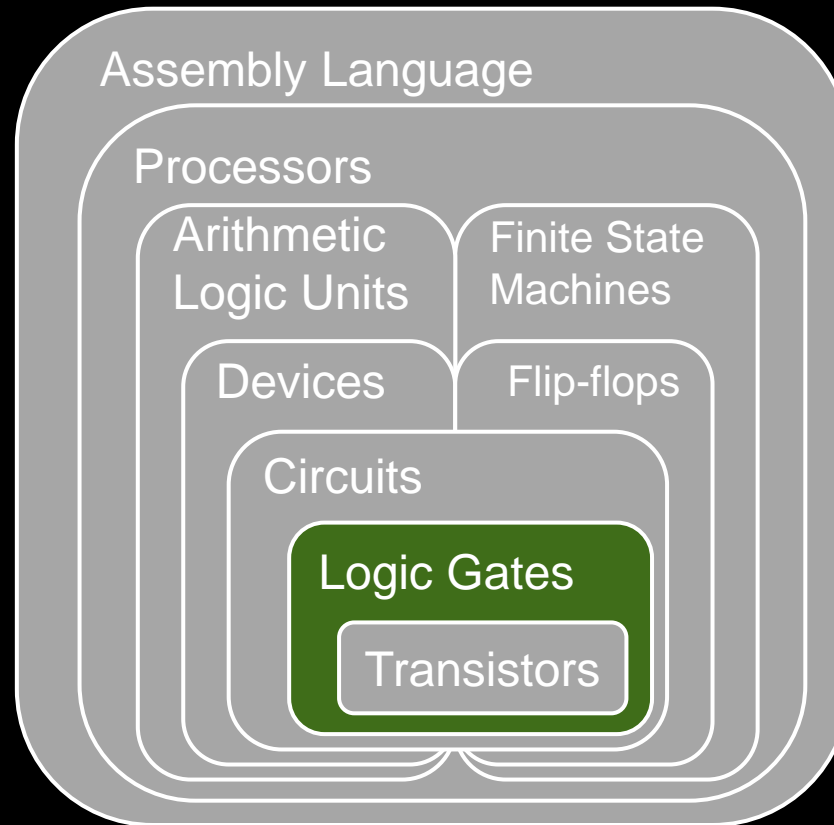
Final Exam

- **Worth 33%**
- **Closed book.**
- **In-person (as currently planned)**
- **Details will follow.**

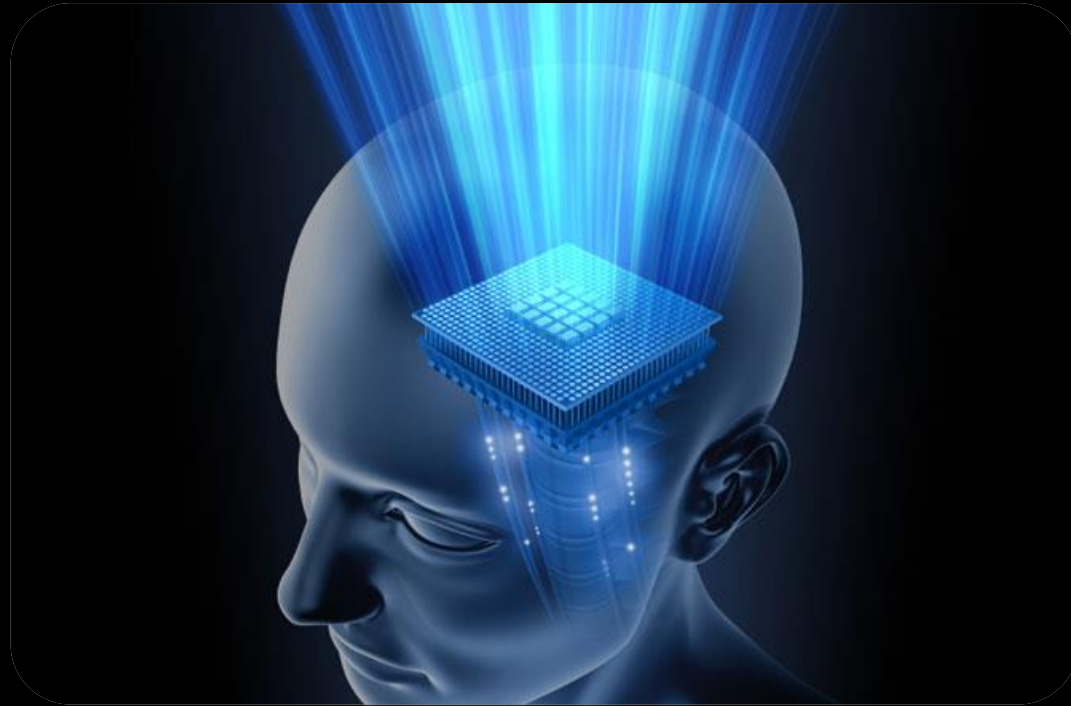
How to do well in CSCB58

- **Be Interested!**
- **Put in the effort:**
 - **Do labs!**
 - **Practice solving lecture questions on your own.**
 - **Study.**
 - **“In theory there’s no difference between practice and theory, but in practice there usually is”**
- **Interact**
 - **lectures, labs (TAs), Piazza**
 - **Ask questions in Piazza or review sessions.**

This week



You already know some important concepts...

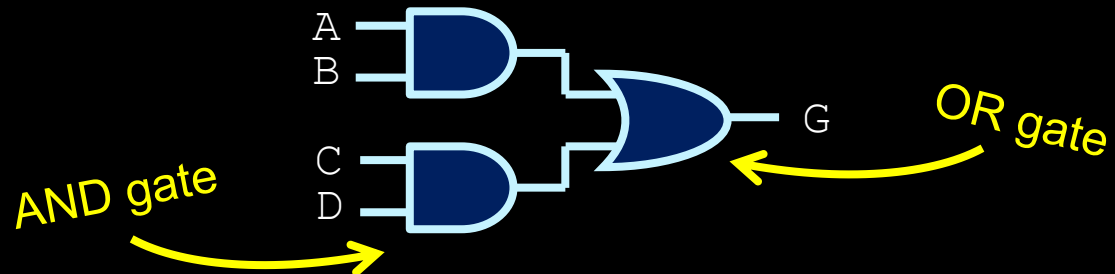


Boolean Logic from CSCA67

- Example: Create an expression that is true if both variables A and B are true, or if both C and D are true.

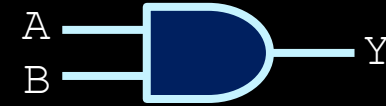
$$G = A \ \& \ B \ | \ C \ \& \ D$$

- Now create a circuit that does the same thing:

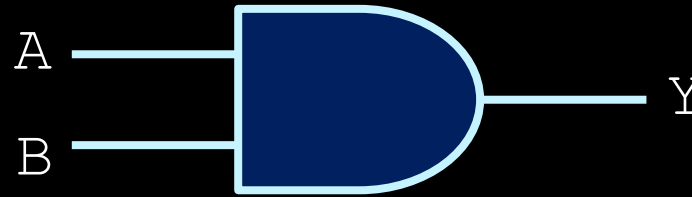


Logic Gates

- If you know how to create simple logical expressions, you already know the basics of putting logic gates together to form simple circuits.
- Just need to know which logic operations are represented by which gate!



AND Gate

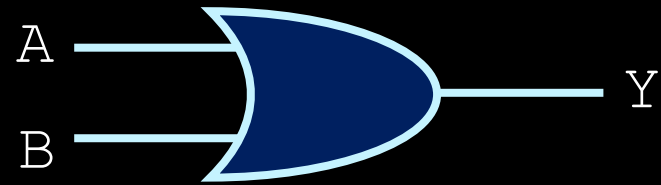


truth table
For every combination of **inputs**, what are the **outputs**?

inputs		outputs
A	B	Y
0	0	0
0	1	0
1	0	0
1	1	1

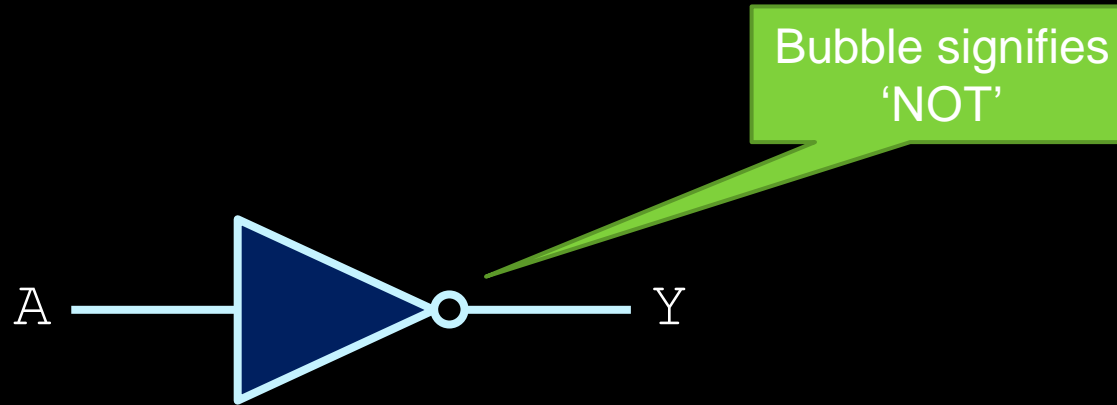
If the truth table for two circuits is the same, they are equivalent

OR Gate



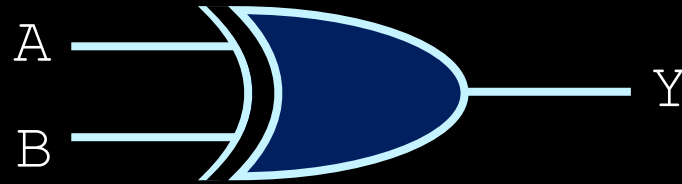
A	B	Y
0	0	0
0	1	1
1	0	1
1	1	1

NOT Gates



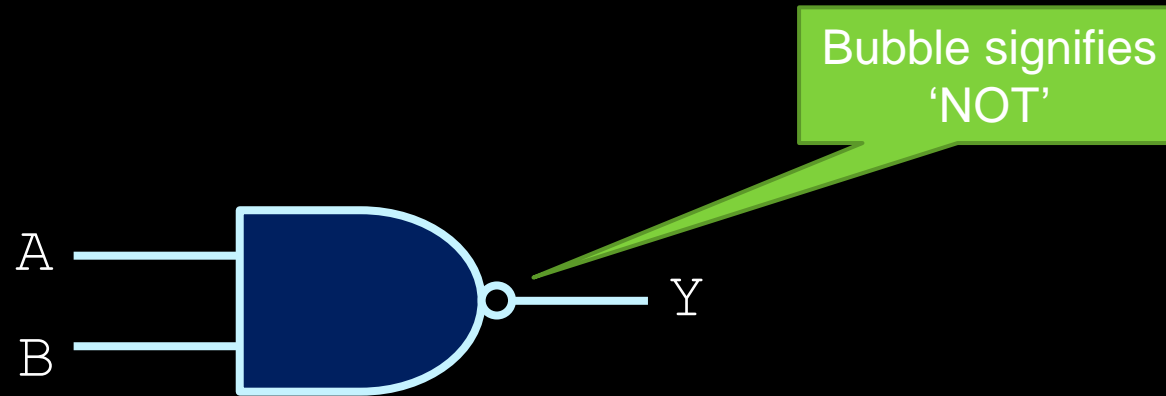
A	Y
0	1
1	0

XOR Gates



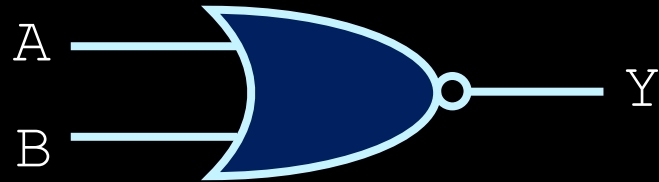
A	B	Y
0	0	0
0	1	1
1	0	1
1	1	0

NAND Gates



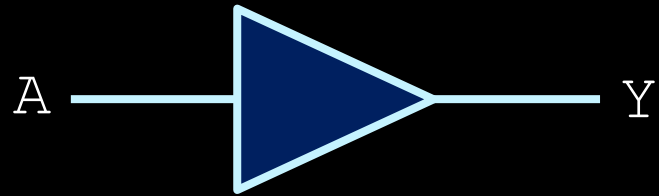
A	B	Y
0	0	1
0	1	1
1	0	1
1	1	0

NOR Gates



A	B	Y
0	0	1
0	1	0
1	0	0
1	1	0

Buffer



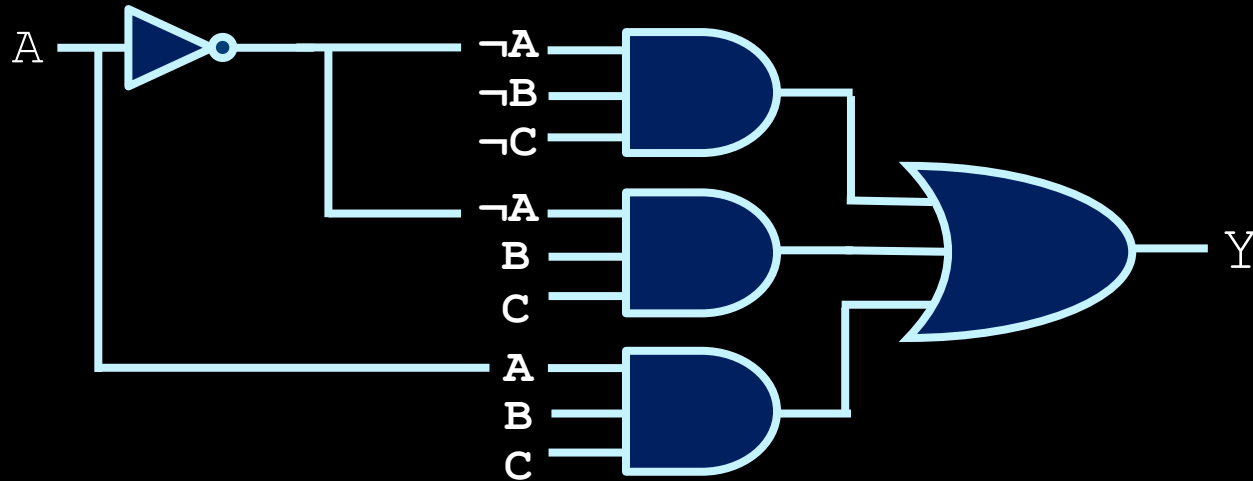
This is not as silly as you might think now, as we'll see later...

A	Y
0	0
1	1

Circuit Design?

- Creating circuit logic can be similar to creating Boolean logic in Python, C or Java:

```
Y = (!A and !B and !C) or  
    (!A and B and C) or  
    (A and B and C)
```

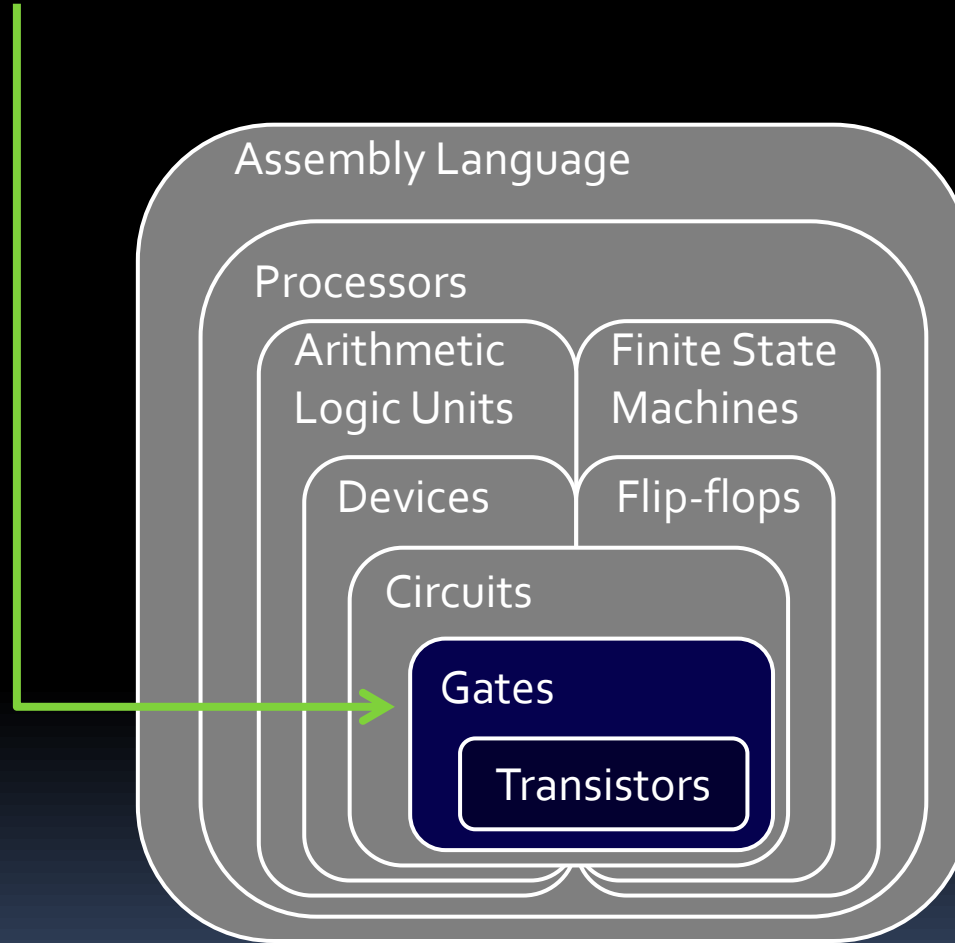


The real challenge in circuit design...

- Given a truth table or description...
- ...find a circuit that implements it.
- Many ways of tackling the problem

A	B	C	Y
0	0	0	1
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

You are here



Making logic with gates

- Logic gates create an output value, based on one or more input values.
- These correspond to Boolean logic that we've seen before in CSCAo8/A48/A67:

AND



A	B	Y
0	0	0
0	1	0
1	0	0
1	1	1

OR



A	B	Y
0	0	0
0	1	1
1	0	1
1	1	1

NOT



A	Y
0	1
1	0

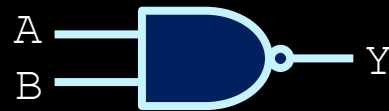
Other gates

XOR



A	B	Y
0	0	0
0	1	1
1	0	1
1	1	0

NAND



A	B	Y
0	0	1
0	1	1
1	0	1
1	1	0

NOR



A	B	Y
0	0	1
0	1	0
1	0	0
1	1	0

Aside: notation

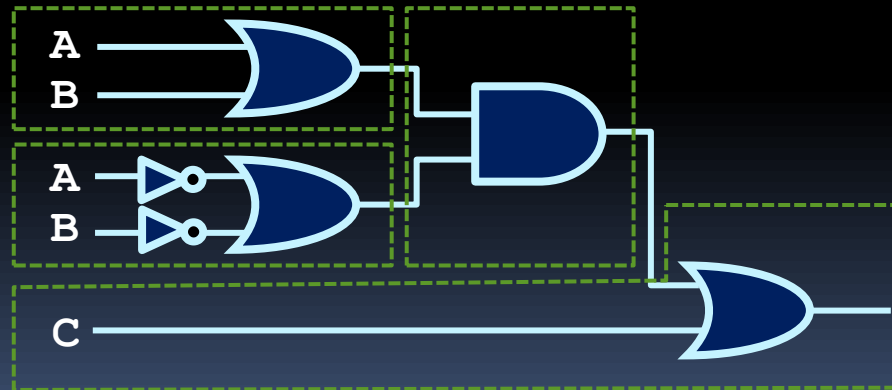
- While we're talking about notation...
 - **AND** operations are denoted in these expressions by the multiplication symbol.
 - e.g. $A \cdot B \cdot C$ or ABC or $A * B * C \approx A \wedge B \wedge C$
 - **OR** operations are denoted by the addition symbol.
 - e.g. $A + B + C \approx A \vee B \vee C$
 - **NOT** is denoted by multiple symbols.
 - e.g. $!A$ or $\sim A$ or \bar{A} or A' or $\neg A$
 - **XOR** occurs rarely in circuit expressions.
 - e.g. $A \oplus B$

Making Boolean expressions

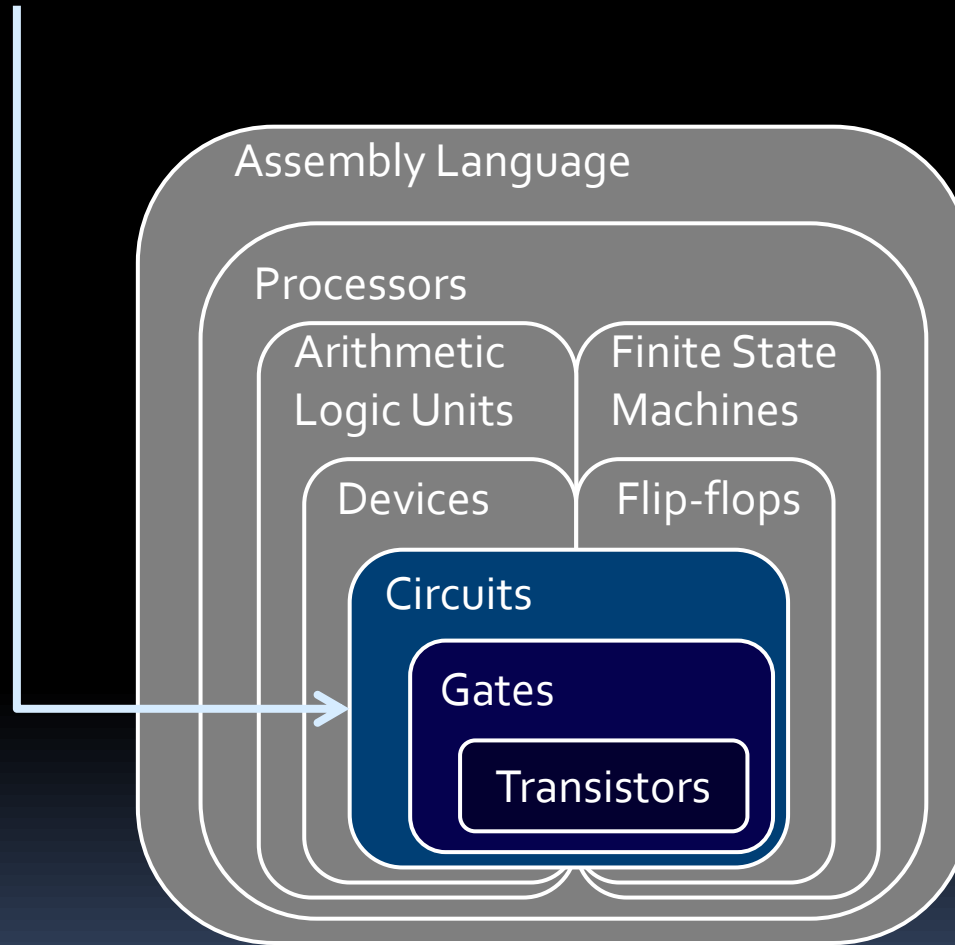
- So how would you represent boolean expressions using logic gates?

$$Y = (A \text{ or } B) \text{ and } (\text{not } A \text{ or not } B) \text{ or } C$$

- Like so:

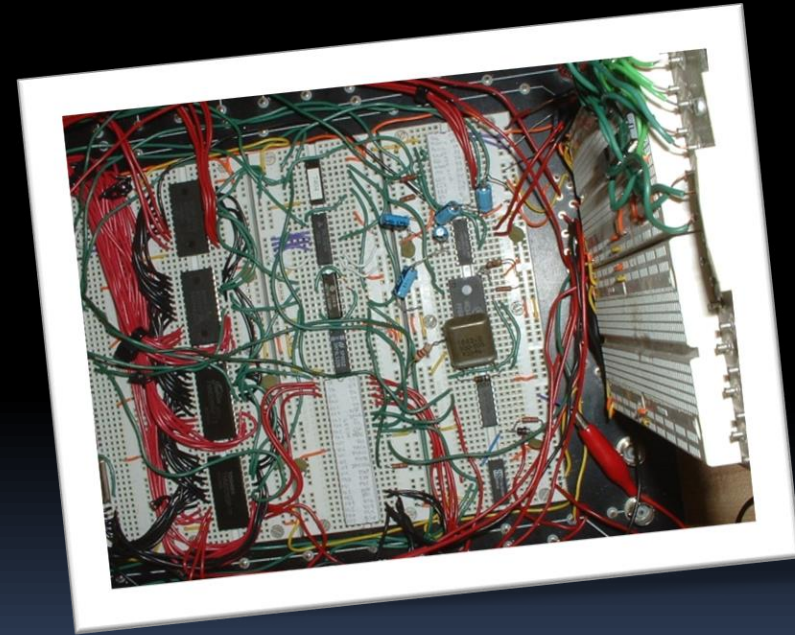


Now you are here



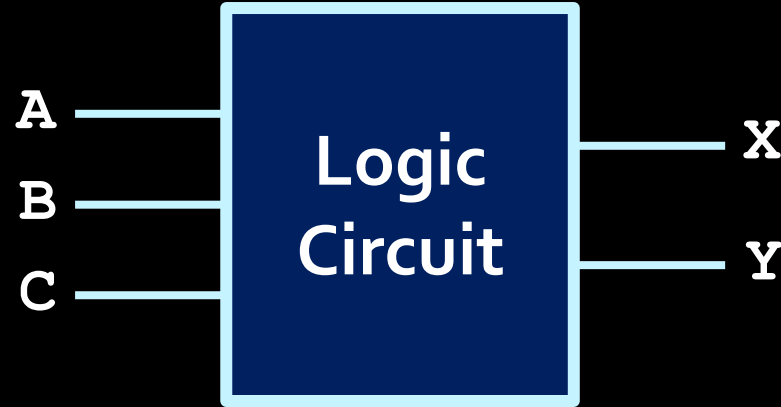
Creating complex circuits

- What do we do in the case of more complex circuits, with several inputs and more than one output?
 - If you're lucky, a truth table is provided to express the circuit.
 - Usually the behaviour of the circuit is expressed in words, and the first step involves creating a truth table that represents the described behaviour.



Circuit example

- The circuit on the right has three inputs (A, B and C) and two outputs (X and Y).



- What logic is needed to set X high when all three inputs are high?
- What logic is needed to set Y high when the number of high inputs is odd?

Combinational circuits

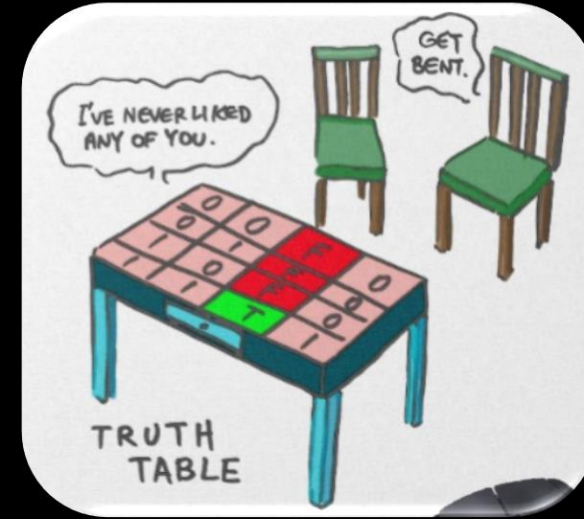
- Small problems can be solved easily.



- Larger problems require a more systematic approach.
 - ▣ Example: Given three inputs A, B, and C, make output Y high in the case where all of the inputs are low, or when A and B are low and C is high, or when A and C are low but B is high, or when A is low and B and C are high.

Creating complex logic

- How do we approach problems like these (and circuit problems in general)?
- **Basic steps:**
 1. Create truth tables.
 2. Express as Boolean expression.
 3. Convert to gates.
- The key to an efficient design?
 - Spending extra time on Step #2.



Example truth table

- Consider the following example:
 - *"Given three inputs A , B , and C , make output Y high wherever any of the inputs are low, except when all three are low or when A and C are high."*
 - This leads to the truth table on the right.
 - Is there a **more compact way** to describe this function?

A	B	C	Y
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0

Warm-Up Exercise

For each of the following logic expressions, what are the A, B, C values that make the expression evaluate to 1?

$A'B'C'$

ABC

$A'BC$

ABC'

Answers to Warm-Up Exercise

For each of the following logic expressions, what are the A, B, C values that make the expression evaluate to 1?

$A'B'C'$

- A=0, B=0, C=0, and only this!

ABC

- 111 and only this!

$A'BC$


- 011 and only this!

ABC'

- 110 and only this!

Maxterms, informally


- Assume a standard truth table format.
 - Sort rows as if input ABC is binary number.
- Maxterms tell us **which rows have low output**.
 - These rows are referred to as **maxterms**.
- Express circuit behaviour by listing those rows.
 - In this example, we only need maxterms M_0 M_5 M_7

Row index	A	B	C	Y		Maxterm	Y
0	0	0	0	0		M_0	0
1	0	0	1	1		M_1	1
2	0	1	0	1		M_2	1
3	0	1	1	1		M_3	1
4	1	0	0	1		M_4	1
5	1	0	1	0		M_5	0
6	1	1	0	1		M_6	1
7	1	1	1	0		M_7	0

Minterms, informally

- A more popular alternative:
list which input rows cause **high output**.
 - These rows are referred to as **minterms**.
 - In this case we have the minterms m_1 m_2 m_3 m_4 m_6

Row index	A	B	C	Y
0	0	0	0	0
1	0	0	1	1
2	0	1	0	1
3	0	1	1	1
4	1	0	0	1
5	1	0	1	0
6	1	1	0	1
7	1	1	1	0



Minterm	Y
m_0	0
m_1	1
m_2	1
m_3	1
m_4	1
m_5	0
m_6	1
m_7	0

Minterms and maxterms

- A more formal description:
 - **minterm** = an AND expression with every input present in true or complemented form.
 - **maxterm** = an OR expression with every input present in true or complemented form.
- For example, given four inputs (A, B, C, D):
 - Valid minterms:
 - $A \cdot B \cdot \bar{C} \cdot D$, $A \cdot \bar{B} \cdot C \cdot \bar{D}$, $A \cdot B \cdot C \cdot D$
 - Valid maxterms:
 - $A+B+\bar{C}+D$, $A+\bar{B}+C+\bar{D}$, $A+B+C+D$
 - Neither minterm nor maxterm:
 - $A \cdot B+C \cdot D$, $A \cdot B \cdot D$, $A+B$

Naming

- Given n inputs, there are 2^n minterms and maxterms possible (same as rows in a truth table).
 - **minterms** are labeled as m_x
maxterms are labeled as M_x
 - The x subscript indicates the row in the truth table.
 - x starts at 0 and ends with $n-1$.
 - Minterms are about when output is 1:
 - From m_0 for $\bar{A} \cdot \bar{B} \cdot \bar{C}$ to m_7 for $A \cdot B \cdot C$
 - Maxterms are about when output is 0:
 - $M_0 (A+B+C)$ to $M_7 (\bar{A}+\bar{B}+\bar{C})$
- } $n=3$

m_0 vs M_0

- Minterm m_0 is \bar{A} and \bar{B} and \bar{C}
 - $m_0 = 1$ if and only if: $A = B = C = 0$ (row 0)
- Maxterm M_0 is A or B or C
 - $M_0 = 0$ if and only if: $A = B = C = 0$ (row 0)
- **Minterms** tell us when **the output is 1**
- **Maxterms** tell us when **the input is 0**

Examples

For example, given four inputs (A, B, C, D):

- Valid minterms:
 - $A \cdot B \cdot \bar{C} \cdot D$ m_{13} (because $1101 = 13$)
 - $A \cdot \bar{B} \cdot C \cdot \bar{D}$ m_{10} (because $1010 = 10$)
 - $A \cdot B \cdot C \cdot D$ m_{15} (because $1111 = 15$)
- Valid maxterms:
 - $A+B+\bar{C}+D$ M_2 (because $0010 = 2$)
 - $A+\bar{B}+C+\bar{D}$ M_5 (because $0101 = 5$)
 - $A+B+C+D$ M_0 (because $0000 = 0$)
- Neither minterm nor maxterm:
 - $A \cdot B+C \cdot D$ mixes AND and OR
 - $A \cdot B \cdot D$
 - $A+B$ missing some of the inputs

Quick Exercises

- Given 4 inputs $\underline{A}, \underline{B}, C$ and D write:

- $m_9 \Rightarrow$

- $m_{15} \Rightarrow$

- $m_{16} \Rightarrow$

- $M_2 \Rightarrow$

- Which minterm is this?

- $\overline{A} \cdot B \cdot \overline{C} \cdot \overline{D}$

- Which maxterm is this?

- $A+B+C+\overline{D}$

What is This For?

- Recap:
 - Minterms and maxterms are a shorthand to refer to **rows of the truth table**.
 - **minterms** describe rows where output is **high**.
 - **maxterms** describe rows where output is **low**.
- Use minterms and maxterms to go from truth table to logic expression
 - Define output by **OR-ing minterms** or **AND-ing maxterms**.
 - Don't mix them both

Using minterms and maxterms

- What are minterms used for?
 - A single minterm indicates a set of inputs that will make the output go high.
 - Example: m_2
 - Output **only goes high in third row** of truth table.

Row index	A	B	C	D	m_2
0	0	0	0	0	0
1	0	0	0	1	0
2	0	0	1	0	1
3	0	0	1	1	0
4	0	1	0	0	0
5	0	1	0	1	0
6	0	1	1	0	0
7	0	1	1	1	0
8	1	0	0	0	0
9	1	0	0	1	0
10	1	0	1	0	0
11	1	0	1	1	0
12	1	1	0	0	0
13	1	1	0	1	0
14	1	1	1	0	0
15	1	1	1	1	0

Using minterms and maxterms

- What happens when you **OR** two **minterms**?
- Result is output that goes high in both minterm cases.
- **For m_2+m_8 , both third and ninth rows** of truth table result in high output.

	A	B	C	D	m_2	m_8	m_2+m_8
0	0	0	0	0	0	0	
1	0	0	0	1	0	0	
2	0	0	1	0	1	0	
3	0	0	1	1	0	0	
4	0	1	0	0	0	0	
5	0	1	0	1	0	0	
6	0	1	1	0	0	0	
7	0	1	1	1	0	0	
8	1	0	0	0	0	1	
9	1	0	0	1	0	0	
10	1	0	1	0	0	0	
11	1	0	1	1	0	0	
12	1	1	0	0	0	0	
13	1	1	0	1	0	0	
14	1	1	1	0	0	0	
15	1	1	1	1	0	0	

Creating Boolean expressions

- Two canonical forms of Boolean expressions:
- **Sum-of-Minterms** (SOM):
 - Since each minterm corresponds to a single high output in the truth table, the combined high outputs are a **union** of these minterm expressions.
 - Also known as: Sum-of-Products.
- **Product-of-Maxterms** (POM):
 - Since each maxterm only produces a single low output in the truth table, the combined low outputs are an **intersection** of these maxterm expressions.
 - Also known as Product-of-Sums.

$$Y = m_2 + m_6 + m_7 + m_{10} \quad (\text{SOM})$$

	A	B	C	D	m_2	m_6	m_7	m_{10}	Y
0	0	0	0	0					
1	0	0	0	1					
2	0	0	1	0					
3	0	0	1	1					
4	0	1	0	0					
5	0	1	0	1					
6	0	1	1	0					
7	0	1	1	1					
8	1	0	0	0					
9	1	0	0	1					
10	1	0	1	0					
11	1	0	1	1					
12	1	1	0	0					
13	1	1	0	1					
14	1	1	1	0					
15	1	1	1	1					

$$Y = m_2 + m_6 + m_7 + m_{10} \quad (\text{SOM})$$

	A	B	C	D	m_2	m_6	m_7	m_{10}	Y
0	0	0	0	0	0	0	0	0	0
1	0	0	0	1	0	0	0	0	0
2	0	0	1	0	1	0	0	0	1
3	0	0	1	1	0	0	0	0	0
4	0	1	0	0	0	0	0	0	0
5	0	1	0	1	0	0	0	0	0
6	0	1	1	0	0	1	0	0	1
7	0	1	1	1	0	0	1	0	1
8	1	0	0	0	0	0	0	0	0
9	1	0	0	1	0	0	0	0	0
10	1	0	1	0	0	0	0	1	1
11	1	0	1	1	0	0	0	0	0
12	1	1	0	0	0	0	0	0	0
13	1	1	0	1	0	0	0	0	0
14	1	1	1	0	0	0	0	0	0
15	1	1	1	1	0	0	0	0	0

$$Y = M_3 \cdot M_5 \cdot M_7 \cdot M_{10} \cdot M_{14} \quad (\text{POM})$$

A	B	C	D	M ₃	M ₅	M ₇	M ₁₀	M ₁₄	Y
0	0	0	0						
0	0	0	1						
0	0	1	0						
0	0	1	1						
0	1	0	0						
0	1	0	1						
0	1	1	0						
0	1	1	1						
1	0	0	0						
1	0	0	1						
1	0	1	0						
1	0	1	1						
1	1	0	0						
1	1	0	1						
1	1	1	0						
1	1	1	1						

$$Y = M_3 \cdot M_5 \cdot M_7 \cdot M_{10} \cdot M_{14} \quad (\text{POM})$$

[illegible]

Using Sum-of-Minterms

- Sum-of-Minterms is a way of expressing which inputs cause the output to go high. Product-of-Maxterms is a way of expression which inputs cause the output to go low.
 - Assumes that the truth table columns list the inputs according to some logical or natural order.
- Minterm and maxterm expressions are used for efficiency reasons:
 - More compact than displaying entire truth tables.
 - Sum-of-minterms (SOM) are useful when very few input combinations that produce high output.
 - Product-of-maxterms (POM) useful when expressing truth tables that have very few low output cases.

Converting SOM to gates

- Once you have a Sum-of-Minterms expression, it is easy to convert this to the equivalent combination of gates:

$$m_0 + m_1 + m_2 + m_3 =$$

$$\overline{A} \cdot \overline{B} \cdot \overline{C} + \overline{A} \cdot \overline{B} \cdot C + \overline{A} \cdot B \cdot \overline{C} + \overline{A} \cdot B \cdot C =$$

