Week 8, part B: ALU Instructions





Source: Computer History Museum

Arithmetic instructions

Instruction	Opcode/Function	Syntax	Operation
add	100000	\$d, \$s, \$t	\$d = \$s + \$t
addu	100001	\$d, \$s, \$t	\$d = \$s + \$t
addi	001000	\$t, \$s, i	\$t = \$s + SE(i)
addiu	001001	\$t, \$s, i	\$t = \$s + SE(i)
div	011010	\$s, \$t	lo = \$s / \$t; hi = \$s % \$t
divu	011011	\$s, \$t	lo = \$s / \$t; hi = \$s % \$t
mult	011000	\$s, \$t	hi:lo = \$s * \$t
multu	011001	\$s, \$t	hi:lo = \$s * \$t
sub	100010	\$d, \$s, \$t	\$d = \$s - \$t
subu	100011	\$d, \$s, \$t	\$d = \$s - \$t

Notes: "hi" and "lo" refer to the HI and LO registers (see register slide).
"SE" = "sign extend".



R-type vs I-type arithmetic

R-Type

- add, addu
- div, divuaddiu
- mult, multu
- sub, subu

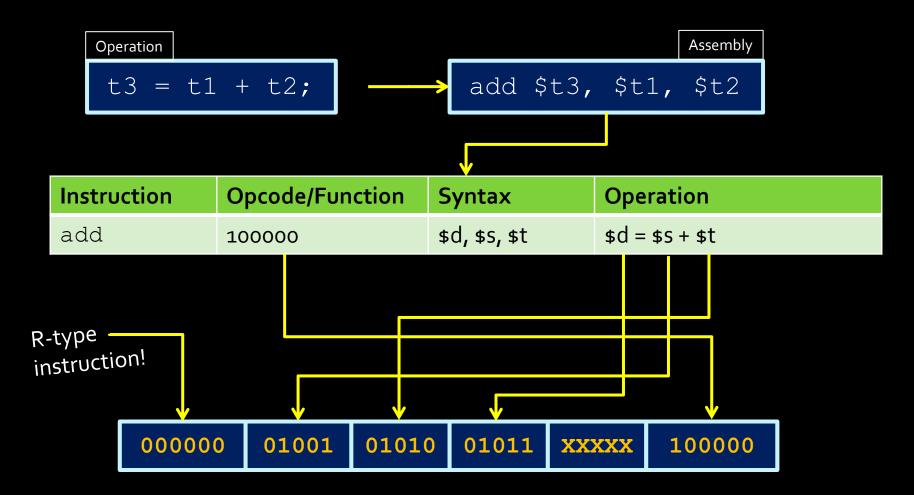
I-Type

- addi

- In general, most instructions are R-type (meaning all operands are registers) and some are I-type (meaning they use an immediate/constant value in their operation).
- Can you recognize which of the following are R-type and I-type instructions? (Hint: "i" for "immediate")



Assembly -> Machine Code



Although we specify "don't care" bits as \times values, in practice the assembler generally sets them to zero



Unsigned Instructions

- What is the difference between add or addu?
 - Both do exactly same thing! Add numbers.
- "u" stands for "unsigned"
 - Causes a "trap" (a.k.a exception) if there is overflow
 - Stops execution of current code.
 - addu ignores this overflow
- mult and multu are not the same!
 - Slight difference in operation. Use the right one!
 - Neither check for overflow.



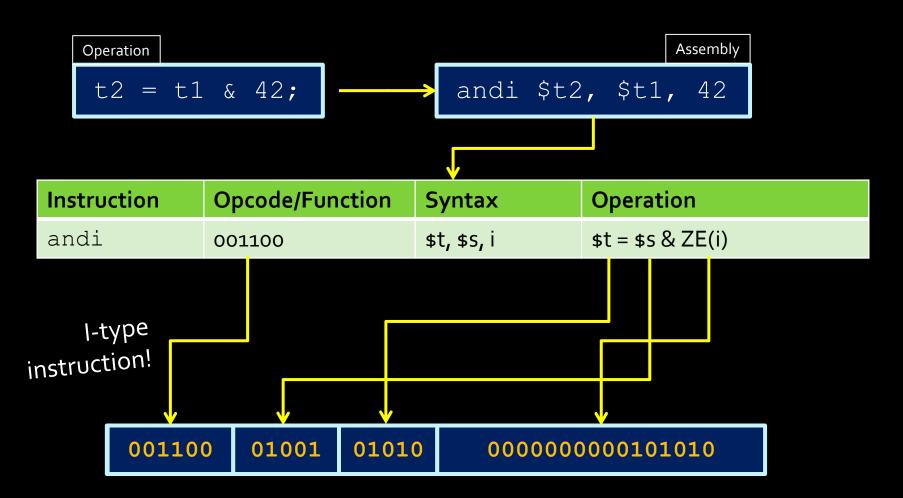
Logical instructions

Instruction	Opcode/Function	Syntax	Operation
and	100100	\$d, \$s, \$t	\$d = \$s & \$t
andi	001100	\$t, \$s, i	\$t = \$s & ZE(i)
nor	100111	\$d, \$s, \$t	\$d = ~(\$s \$t)
or	100101	\$d, \$s, \$t	\$d = \$s \$t
ori	001101	\$t, \$s, i	\$t = \$s ZE(i)
xor	100110	\$d, \$s, \$t	\$d = \$s ^ \$t
xori	001110	\$t, \$s, i	\$t = \$s ^ ZE(i)

Note: ZE = zero extend (pad upper bits with 0 value).



Assembly → Machine Code II





Shift instructions

Instruction	Opcode/Function	Syntax	Operation
sll	000000	\$d, \$t, a	\$d = \$t << a
sllv	000100	\$d, \$t, \$s	\$d = \$t << \$s
sra	000011	\$d, \$t, a	\$d = \$t >> a
srav	000111	\$d, \$t, \$s	\$d = \$t >> \$s
srl	000010	\$d, \$t, a	\$d = \$t >>> a
srlv	000110	\$d, \$t, \$s	\$d = \$t >>> \$s

- Order is \$d, \$t, \$s or \$d, \$t, a (not \$d, \$s, \$t as before!)
- srl = "shift right logical"
- sra = "shift right arithmetic".
- The "v" denotes a variable number of bits, specified by \$s.
- a is shift amount, and is stored in shamt when encoding the R-type machine code instructions.

Data movement instructions

Instruction	Opcode/Function	Syntax	Operation
mfhi	010000	\$d	\$d = hi
mflo	010010	\$d	\$d = lo
mthi	010001	\$ S	hi = \$s
mtlo	010011	\$ S	lo = \$s

 These are instructions for operating on the HI and LO registers described earlier (for multiplication and division)



lui – load upper immediate

Instruction	Opcode/Function	Syntax	Operation
lui	001111	\$ t, i	\$t = i << 16

- Load 16-bit immediate into upper half of the register.
- The lower 16 bits of the register are set to zero.

iiiiiiiiiiiii000000000000000000





ALU instructions in RISC

R type

- add, div, mult, sub
- addu, divu, multu, subu
- or, and, nor, xor

I type

- addi
- addiu
- andi, ori, xori

- Most ALU instructions are R-type instructions.
 - The six-digit codes in the tables are therefore the function codes (opcodes and opcodes)
 - Except the few I-type instructions (addi, andi, ori, etc.)



ALU instructions in RISC

- Not all R-type instructions have an I-type equivalent.
 - We have addi but not subi
 - We have ori but not nori
 - div but not divi
- RISC principle: an operation doesn't need an instruction if it can be performed through multiple existing operations.
 - □ addi \$t0, -1 → "subi" \$t0, 1
 - □ addi + div → "divi"



Pseudoinstructions

- Pseudo instructions look like assembly instructions...
- ...but don't have a dedicated machine code instruction.
- Provided by the assembler
 - Mapping ASM to machine code is more like a many-to-one mapping...
- If a temporary register is needed, use sat



Pseudoinstructions

- Move data from \$t4 to \$t5?
 - \blacksquare move \$t5,\$t4 \rightarrow

add \$t5,\$t4,\$zero

- Multiply and store in \$s3?
 - □ mul \$s1, \$t4, \$t5 →

mult \$t4,\$t5
mflo \$s1

- Load a 32-bit immediate?
 - $^{\circ}$ li \$s0,0x1234ABCD \rightarrow

lui \$s0,\$s0,0x1234
ori \$s0,\$s0,0xABCD



Time to write our first assembly program





Making an assembly program

- Assembly language programs typically have structure similar to simple Python or C programs:
 - They set aside registers to store data.
 - They have sections of instructions that manipulate this data.
- It is always good to decide at the beginning which registers will be used for what purpose!
 - More on this later ©



Compute result = $a^2 + 2b + 10$

- Set up values in registers
 - $a \rightarrow $to, b \rightarrow $t1$
- temp = 10
- temp = temp + b
- temp = temp + b (again!)
- result = a*a

result = result + temp

```
addi $t0, $zero, 7
addi $t1, $zero, 9
addi $t6, $zero, 10
add $t6, $t6, $t1
add $t6, $t6, $t1
mult $t0, $t0
mflo $t4
add $t4, $t4, $t6
```

Formatting Assembly Code

- Start file with .text
 - (we'll see other options later)
- Follow this with:
- .globl main
 - Makes the main label visible to the OS
- main:
 - Tells OS which line of code should run first.
- Write instructions, up to 3 columns per line
 - label: <instr> <params> # comments
 - Labels and comments as needed
- At the end of the program, tell the OS to finish:

```
li $v0, 10 syscall
```



```
# Compute the following result: r = a^2 + 2b + 10
.text
.globl main
main: addi $t0, $zero, 7 # set a=7 for testing
      addi $t1, $zero, 9 # set b=9 for testing
# $t0 will be a, $t1 will be b, $t5 will be r
# $t6 will be temp
      addi $t6, $zero, 10 # add 10 to r
      add $t6, $t6, $t1  # then add b
      add $t6, $t6, $t1 # then add b again
      mult $t0, $t0  # multiply a * a
      mflo $t4
                    # move the low result of a^2
                          # into the register for r
      add $t4, $t4, $t6
                          # add the temporary value
                          \# (2b + 10) to the result
      addi $v0, $zero, 10 # end program
      syscall
```

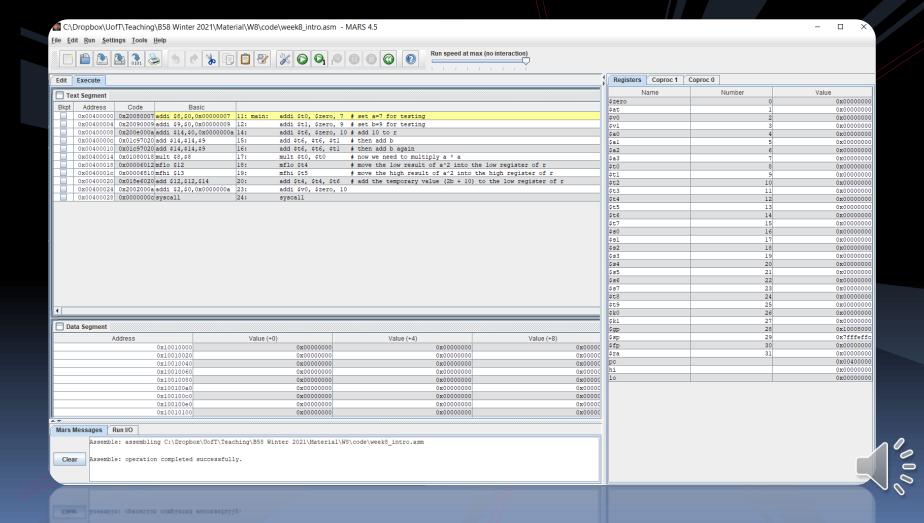
How can we run this?

- We don't have a MIPS CPU handy.
- We'll use a simulator instead.
- A program that simulates the operation of the MIPS CPU on your own computer,



Simulating MIPS

The MARS Simulator



MARS Simulator

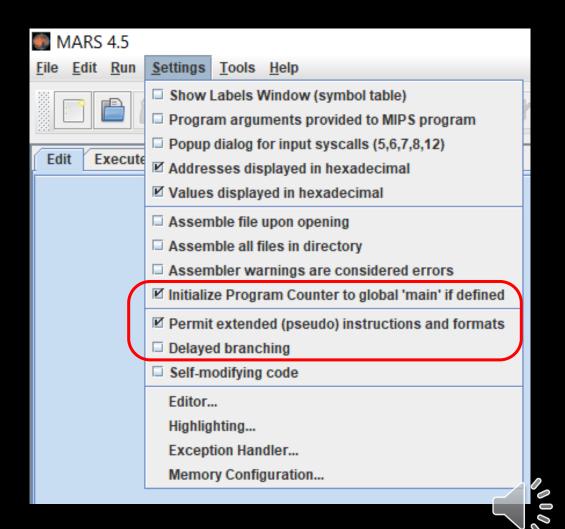
- MARS Simulator official site:
 - http://courses.missouristate.edu/kenvollmar/mars/
 - As with Logisim, you will need Java.
- Official version sometimes freezes during debugging.
- Download <u>alternative</u> MARS jar file from Quercus module on assembly



MARS Settings

Make sure:

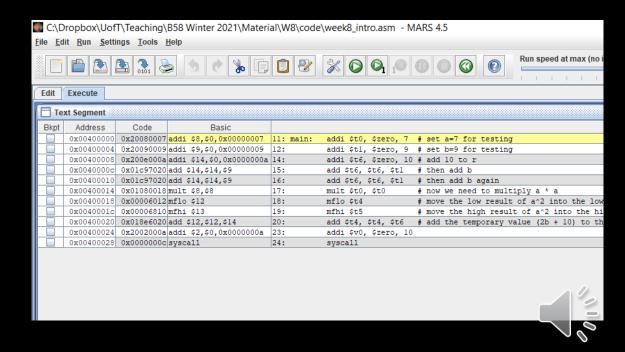
- delayed branching is turned off
- permit extended instruction is on
- Initialize program counter to global 'main' is on



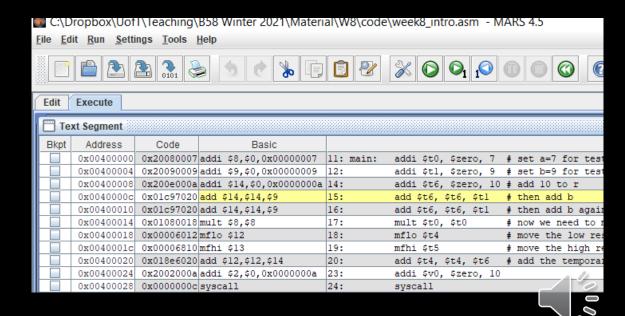
- MARS works like a simple IDE
 - Write assembly program in code editor
 - Save it to an .asm or .s file (doesn't matter)

```
Edit Execute
week8_intro.asm
          mips1.asm
   # Compute the following result
   \# r = a^2 + 2b + 10
   # Register assignment:
      $t0 = a, $t1 = b
      $t6 = temp
  .globl main
  main:
          addi $t0, $zero, 7 # set a=7 for testing
          addi $t1, $zero, 9 # set b=9 for testing
  # compute:
          addi St6, Szero, 10 # add 10 to r
                            # then add b again
```

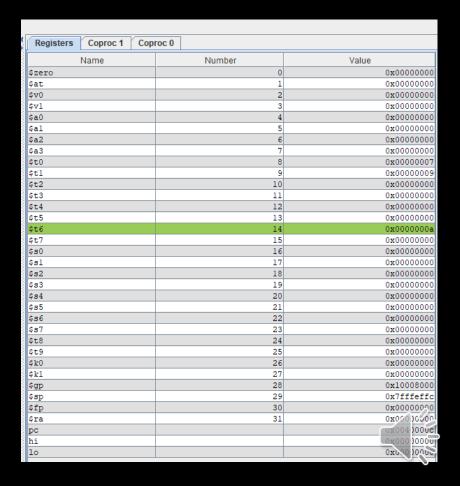
- MARS works like a simple IDE
 - Assemble the program (F₃ or Run → Assemble)
 - Mars will switch to execute view



- MARS works like a simple IDE
 - Run the whole program (F₅ or Run → Go)
 - Or execute line by line (F7 or Run -> Step)



- MARS works like a simple IDE
 - Check the register window to see what is going on.



Get MARS

- You need it for labs.
- And for practice.
- Code from lectures is available on OneDrive.
- Try to execute it on MARS.
- See how it works.
- Play with it yourself.
- Learn!





r = (2a + 5) * (7b)

```
.text
.globl main
# $t0 = a, $t1 = b, $t4 = r
# $t7 = left side, $t8 = right side
main: addi $t0, $zero, 7 # load up some values to test
      addi $t1, $zero, 9
# calculate left side
calc_left: add $t7, $t0, $t0 # ls <- 2a
            addi $t7, $t7, 5  # ls <- ls + 5
# calculate right side
calc_right: addi $t8, $zero, 7 # rs <- 7</pre>
            mult $t8, $t1  # multiply 7 * b
            mflo $t8 # put result back into rs
# multiply left * right and put result into r
mulitply: mult $t7, $t8
            mflo $t4
```

Implement c = max(a,b)?

- Most code does not simply execute linearly from start to finish.
- For example, how would we implement:

Move to next part!

