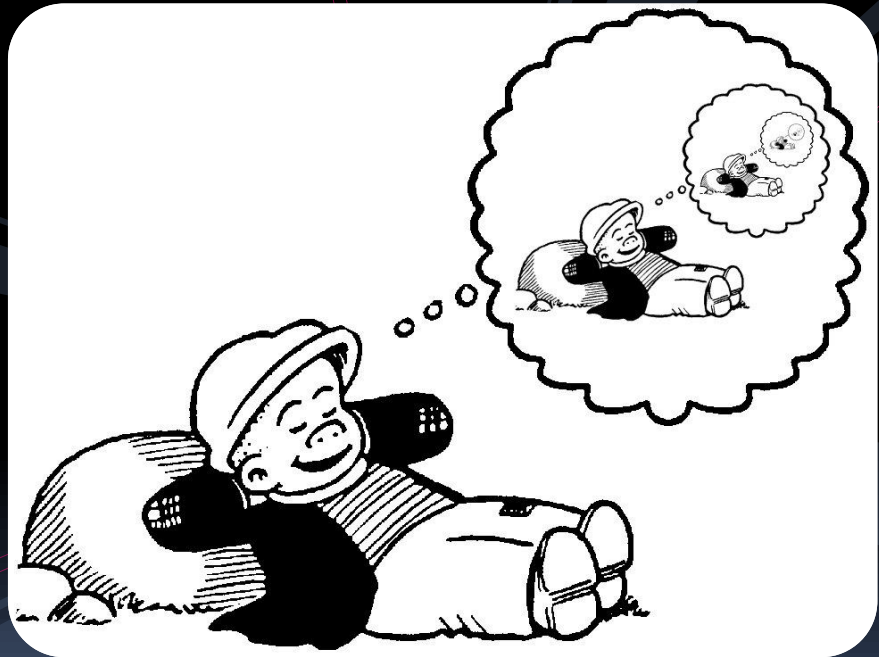


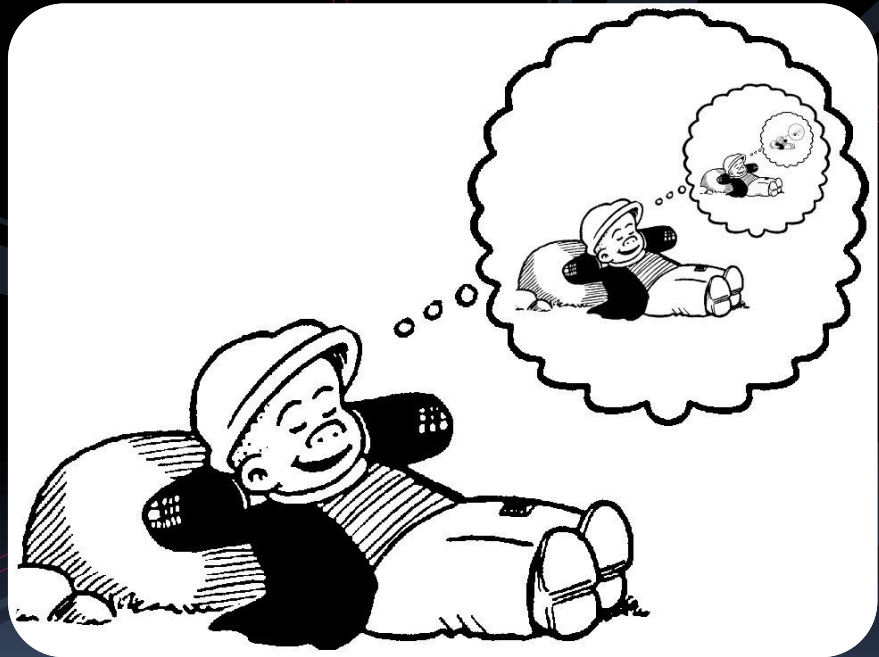
Week 10, part C:

Recursion in Assembly



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Example: factorial(int n)

Basic pseudocode for recursive factorial:

- Base Case ($n == 0$)
 - return 1
- Get factorial($n-1$)
 - Store result in "temp"
- Multiply temp by n
 - Store in result variable
- Return result



Recursive programs

- How do we handle recursive programs?
 - Still needs base case and recursive step, as with other languages.
 - Main difference: function is both caller and callee
 - So what?
 - Just make sure to preserve \$ra and saved registers
 - (As we saw before)

```
int factorial (int x) {  
    if (x==0)  
        return 1;  
    else  
        return x*fact(x-1);  
}
```



Recursive programs

- Solution: the stack!

- Before recursive call, store the register values that you use onto the stack, and restore them when you come back to that point.
- Don't forget to store \$ra, or else the program will loop forever!
 - Because it is returning to the wrong place, just after the recursive call.

```
int factorial (int x) {  
    if (x==0)  
        return 1;  
    else  
        return x*fact(x-1);  
}
```



Factorial solution

- Steps to perform:
 - ▣ Pop x off the stack.
 - ▣ Check if x is zero:
 - If $x==0$, push 1 onto the stack and return to the calling program.
 - If $x \neq 0$, push $x-1$ onto the stack and call factorial again (i.e. jump to the beginning of the code).
 - After recursive call, pop result off of stack and multiply that value by x .
 - Push result onto stack, and return to calling program.

```
int fact (int x) {  
    if (x==0)  
        return 1;  
    else  
        return x*fact(x-1);  
}
```



factorial(int n)

- Base Case ($n == 0$)
 - return 1
- Get factorial($n-1$)
 - Store result in "product"
- Multiply product by n
 - Store in "result"
- Return result

$n \rightarrow \$t0$

$n-1 \rightarrow \$t1$

$\text{fact}(n-1) \rightarrow \$t2$



factorial(int n)

- Pop n off the stack
 - Store in \$t0
- If \$t0 == 0
 - Push 1 onto stack
 - Return to caller
- If \$t0 != 0

- Base Case (n == 0)
 - return 1
- Get factorial(n-1)
 - Store result in "product"
- Multiply product by n
 - Store in "result"
- Return result

n → \$t0

n-1 → \$t1

fact(n-1) → \$t2



factorial(int n)

- Pop n off the stack
 - Store in \$t0
- If \$t0 == 0
 - Push 1 onto stack
 - Return to caller
- If \$t0 != 0
 - Calculate n-1 in \$t1
 - Save \$t0 (n) and \$ra onto stack
 - Push n-1 (\$t1) onto the stack
 - Call factorial

 - ...time passes...

- Base Case (n == 0)
 - return 1
- Get factorial(n-1)
 - Store result in "product"
- Multiply product by n
 - Store in "result"
- Return result

n → \$t0

n-1 → \$t1

fact(n-1) → \$t2

\$t0 is caller-saved!

factorial(int n)

- Pop n off the stack

- Store in \$t0

- If \$t0 == 0

- Push 1 onto stack
 - Return to caller

- If \$t0 != 0

- Calculate n-1 in \$t1
 - Save \$t0 (n) and \$ra onto stack
 - Push n-1 (\$t1) onto the stack
 - Call factorial

- ...time passes...

- Pop the result of factorial(n-1) from stack, store in \$t2
 - Restore \$ra and \$t0 from stack
 - Multiply factorial(n-1) (stored in \$t2) and n (in \$t0)
 - Push result onto stack
 - Return to calling program

- Base Case (n == 0)
 - return 1
- Get factorial(n-1)
 - Store result in "product"
- Multiply product by n
 - Store in "result"
- Return result

n → \$t0

n-1 → \$t1

fact(n-1) → \$t2

\$t0 is caller-saved!

Translated recursive program (part 1)

```
main:      add $t3, $zero, 6      #
           addi $sp, $sp, -4      #
           sw $t3, 0($sp)         #
           jal factorial          #
           ...

factorial: lw $t0, 0($sp)         # pop n off the
           addi $sp, $sp, 4       # stack
           bne $t0, $zero, rec    # if n is zero?

base:      addi $t0, $zero, 1     # base case:
           addi $sp, $sp, -4     # push 1 to the
           sw $t0, 0($sp)        # stack
           jr $ra                # return to caller

rec:       ...                   # recursive case
```

Translated recursive program (part 2)

```
# (... continuing from part 1)
rec:      addi $t1, $t0, -1          # calculate n-1
          addi $sp, $sp, -4         # save n on the
          sw $t0, 0($sp)           # stack
          addi $sp, $sp, -4         # save $ra on the
          sw $ra, 0($sp)           # stack
          addi $sp, $sp, -4         # push n-1 on the
          sw $t1, 0($sp)           # stack

          jal factorial             # call factorial(n-1)

          # returning from recursive call...
          lw $t2, 0($sp)           # pop result of
          addi $sp, $sp, 4          # factorial(n-1)
          lw $ra, 0($sp)           # restore saved $ra
          addi $sp, $sp, 4          # from the stack
          lw $t0, 0($sp)           # restore saved n
          addi $sp, $sp, 4          # from the stack
```



Translated recursive program (part 2)

```
# (... continuing from part 2)

    mult $t0, $t2           # compute result
    mflo $t0                #  n * factorial(n-1)
                             #  assume no overflow

    addi $sp, $sp, -4       # push result onto
    sw $t0, 0($sp)          #  stack
    jr $ra                  # return to caller
```

The code is available on Quercus.

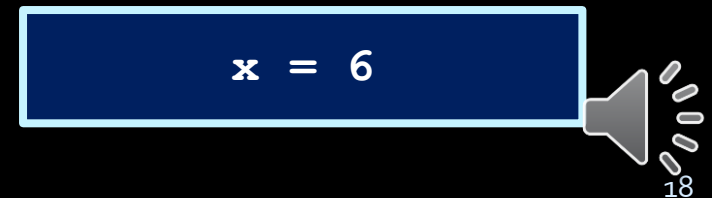
Highly recommended: download and execute it on MARS.



Factorial Stack View

Now running: `main`

Initial call to `factorial` with `x=6`



Factorial Stack View

Now running: `fact(6)`

`pop x=6`

Factorial Stack View

Now running: `fact(6)`

save x on the stack

saved x = 6



Factorial Stack View

Now running: `fact(6)`

save `$ra` in preparation for recursion

`$ra` for `fact(6)`

saved `x = 6`

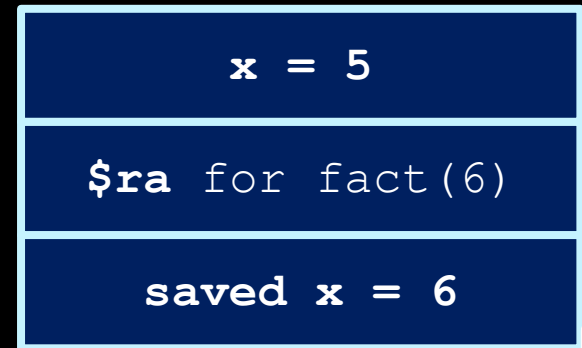
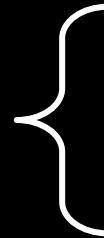


Factorial Stack View

Now running: `fact(6)`

push argument for factorial(5)

factorial(6)

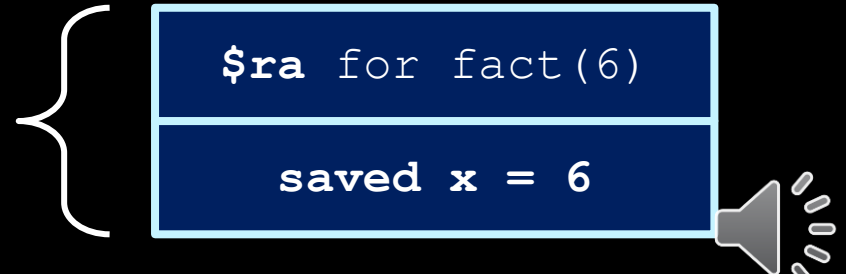


Factorial Stack View

Now running: `fact(5)`

`pop x = 5`

factorial(6)

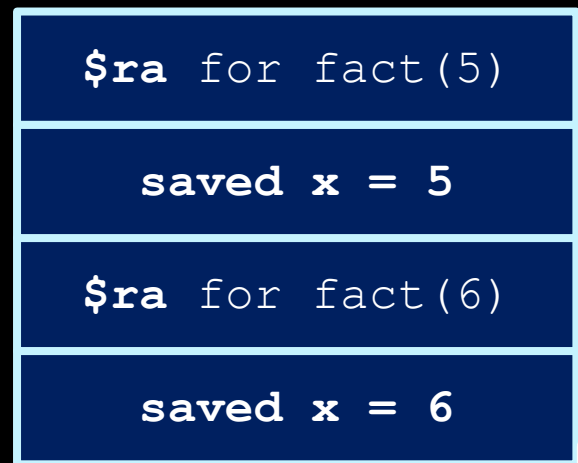


Factorial Stack View

Now running: `fact(5)`

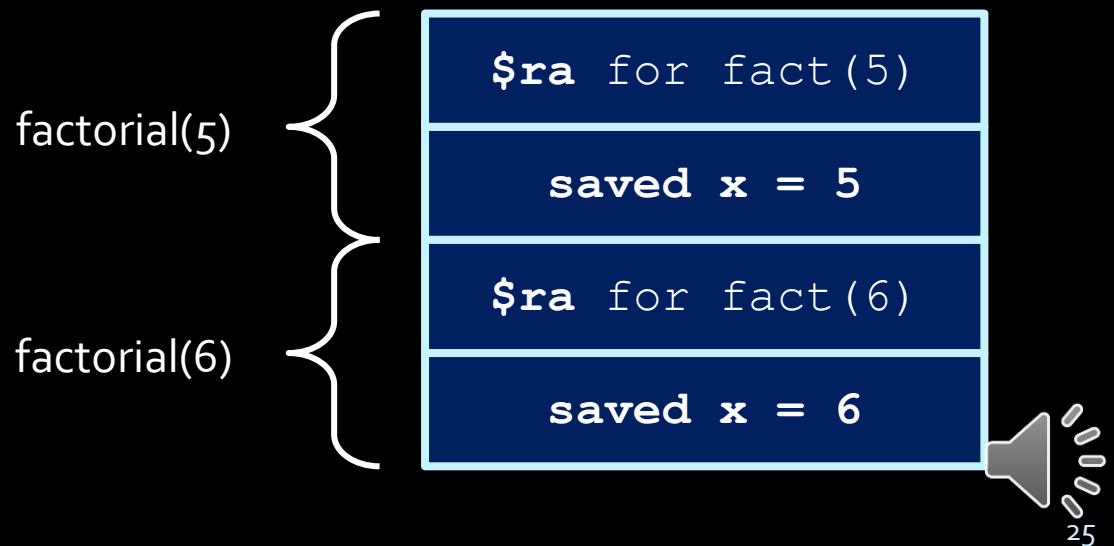
save `$ra` in preparation for recursion

factorial(6)



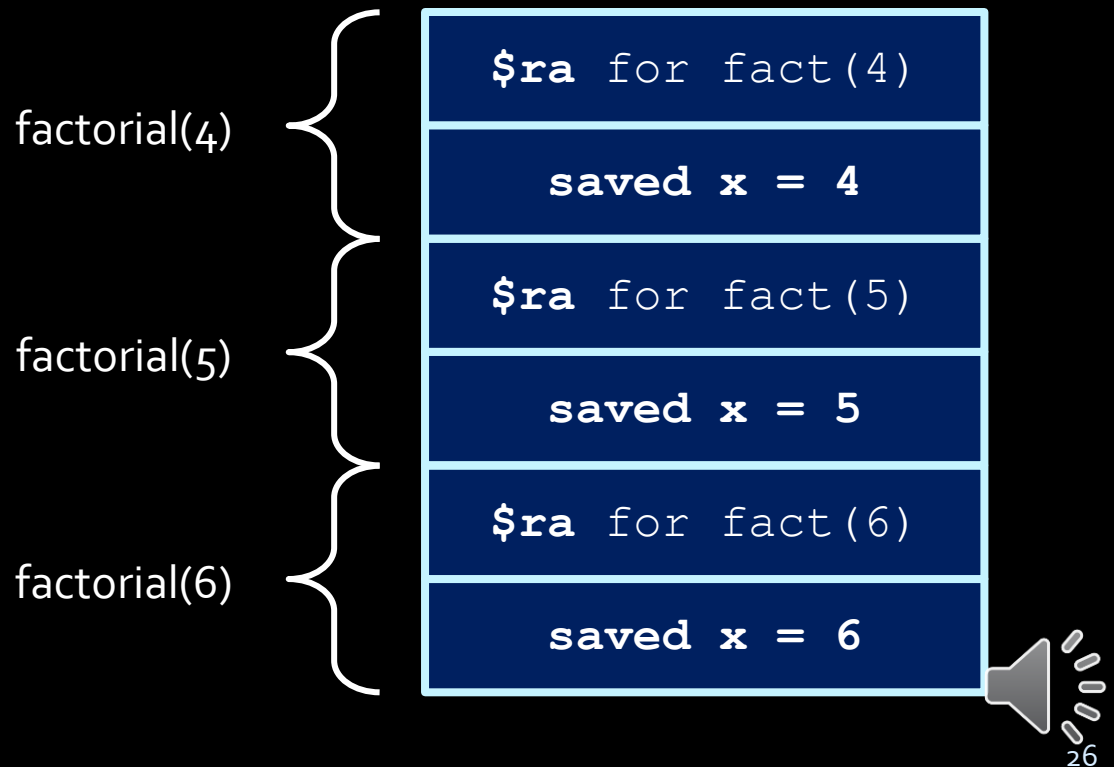
Factorial Stack View

Now running: `fact(5)`



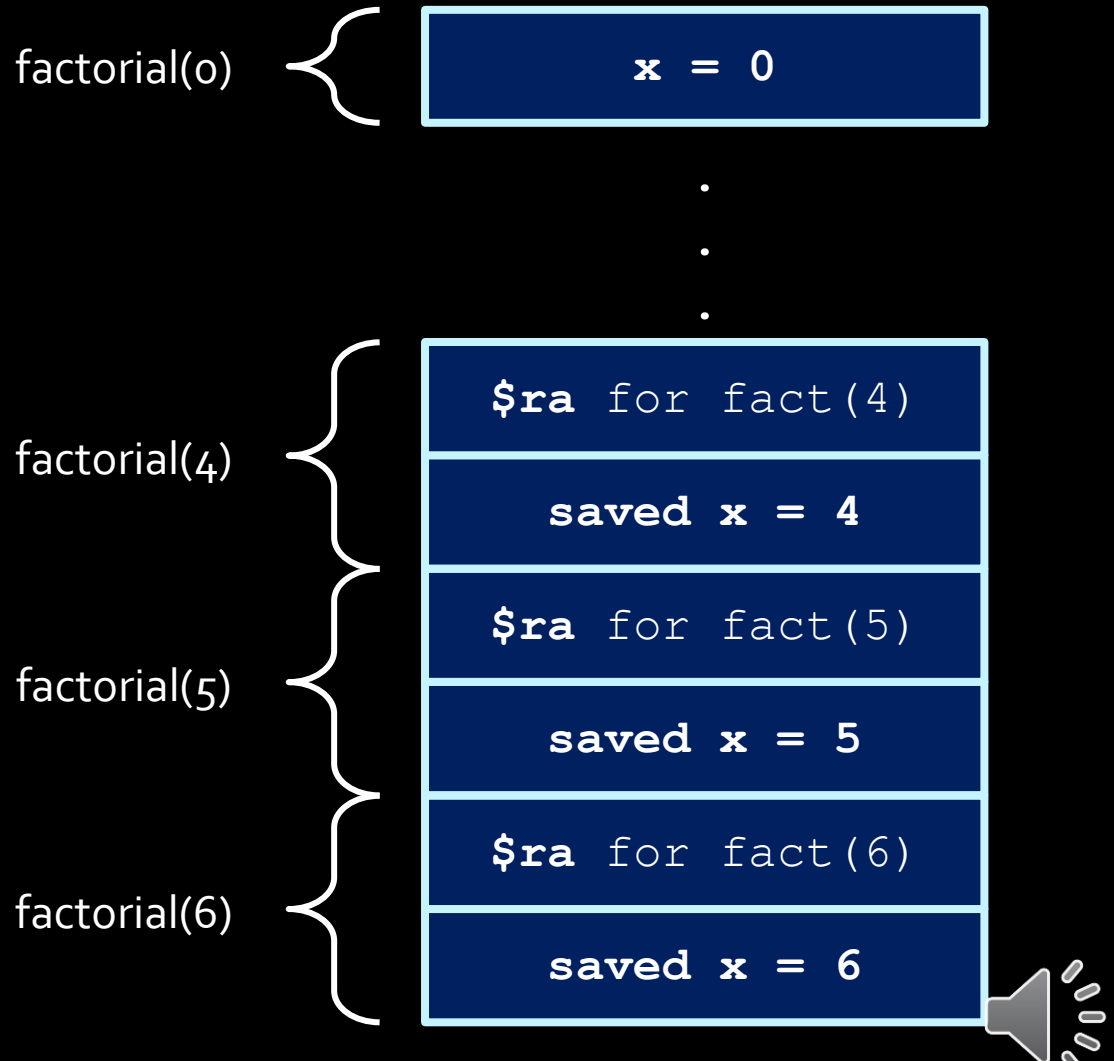
Factorial Stack View

Now running: `fact(4)`



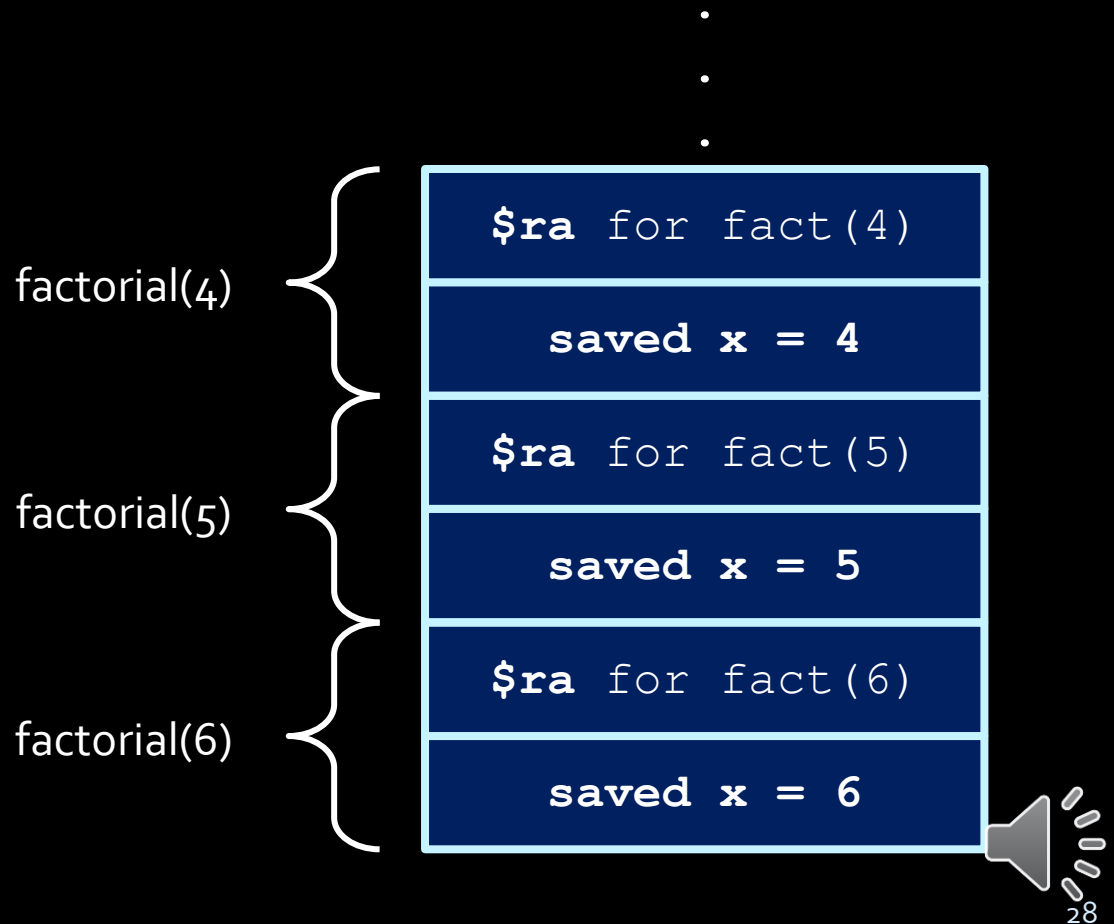
Factorial Stack View

Now running: `fact(0)`



Factorial Stack View

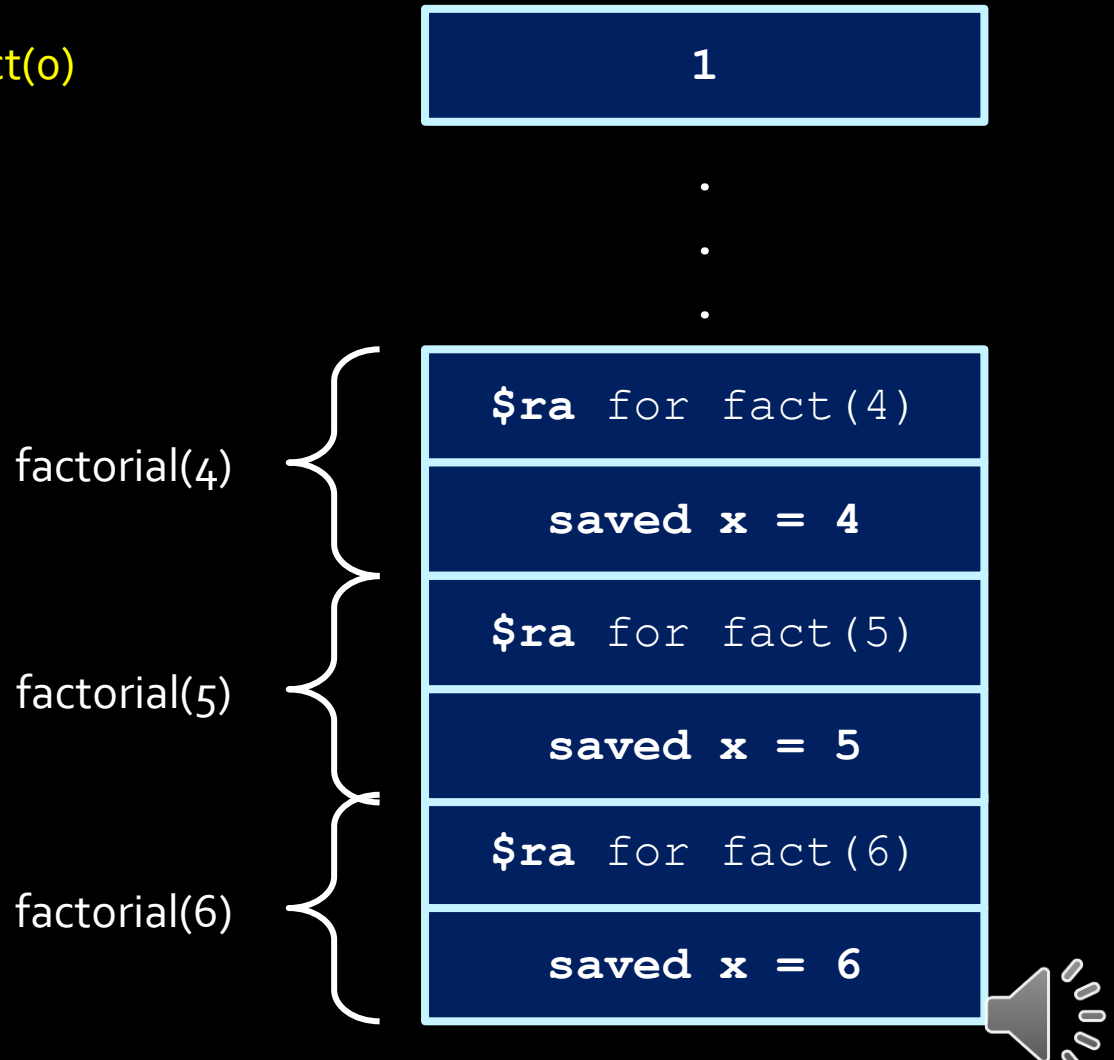
Now running: `fact(o)`



Factorial Stack View

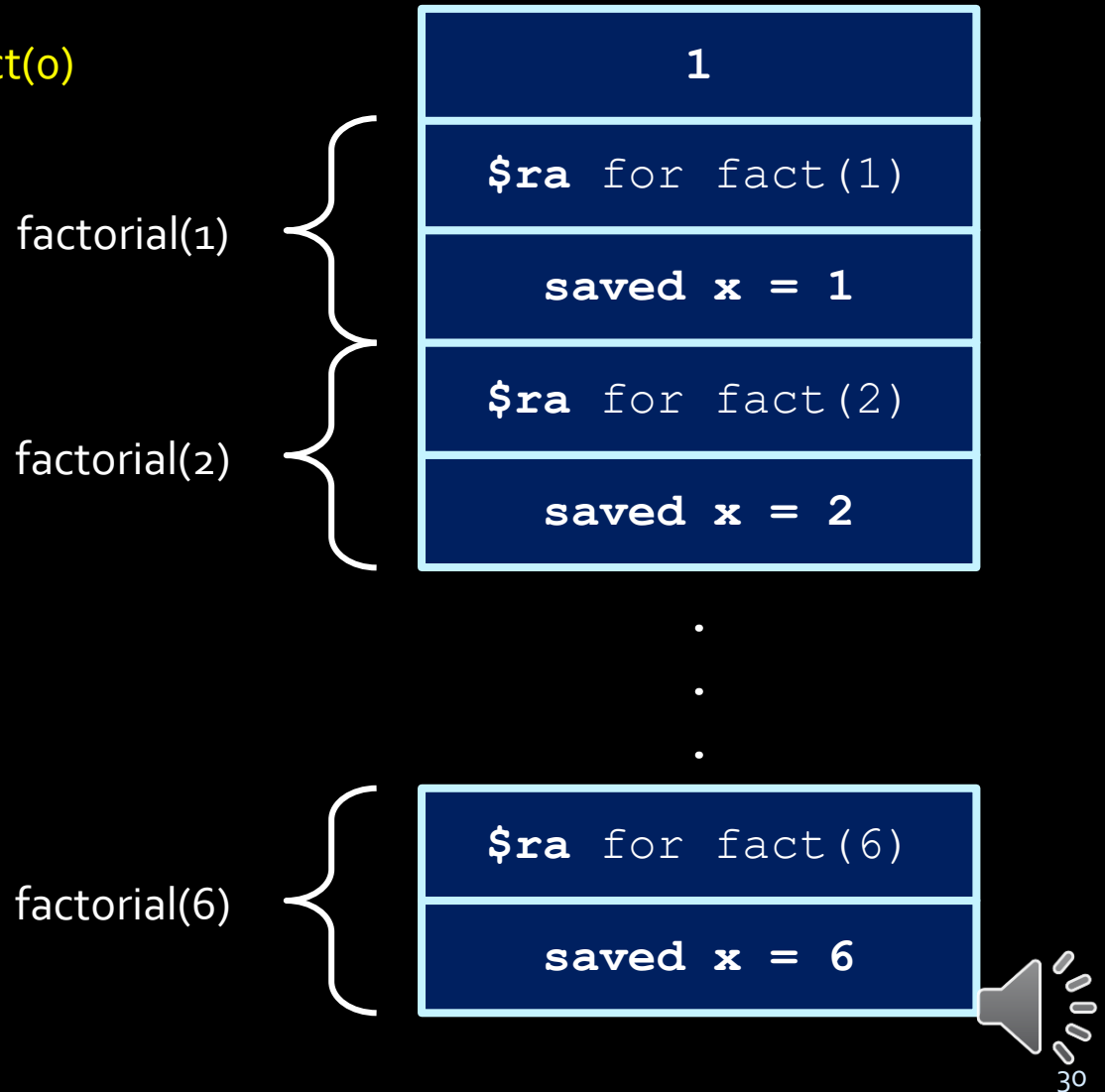
Now running: `fact(o)`

return value of `fact(o)`



Factorial Stack View

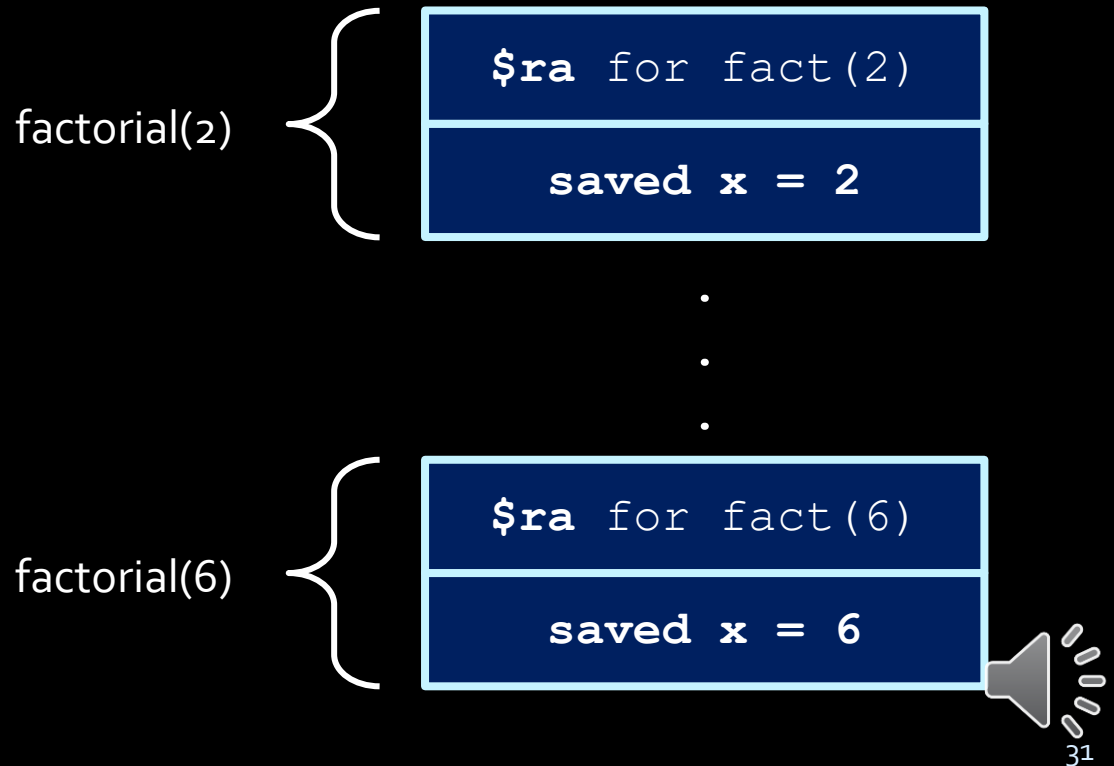
fact(o) is now returning
return value of fact(o)



Factorial Stack View

Now running: `fact(1)`

pop `fact(0)` and saved registers

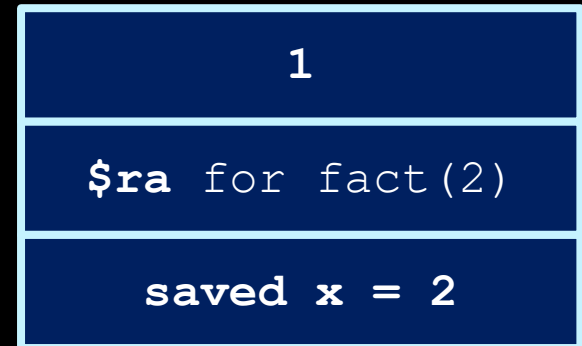


Factorial Stack View

Now running: `fact(1)`

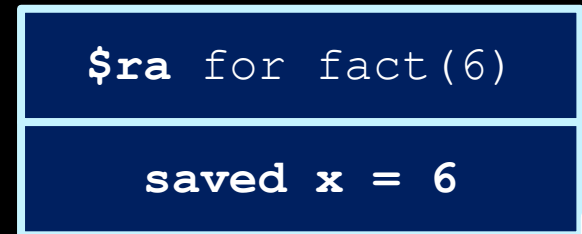
push `1 * fact(0)` and return

factorial(2)



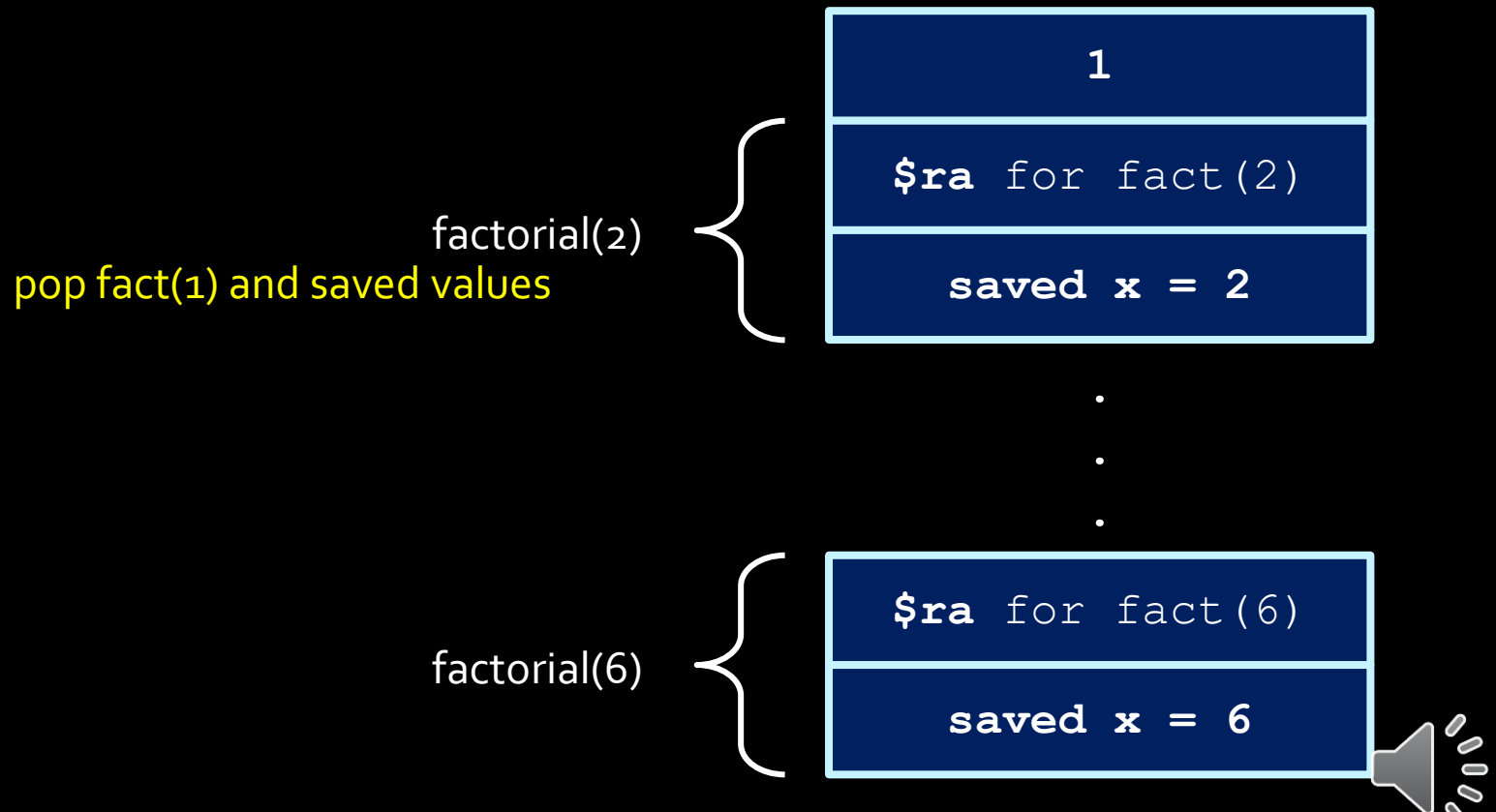
⋮

factorial(6)



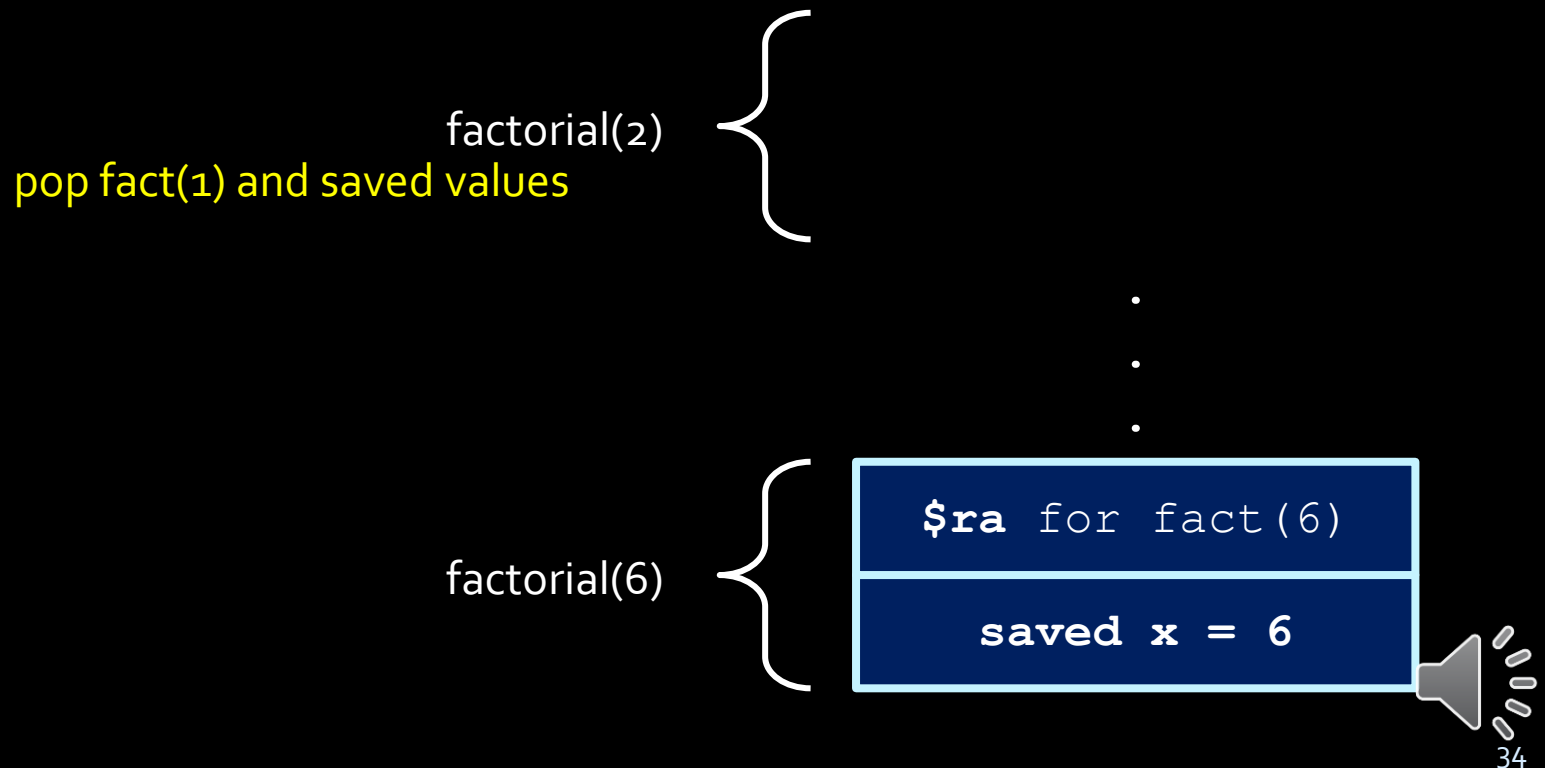
Factorial Stack View

Now running: `fact(1)`



Factorial Stack View

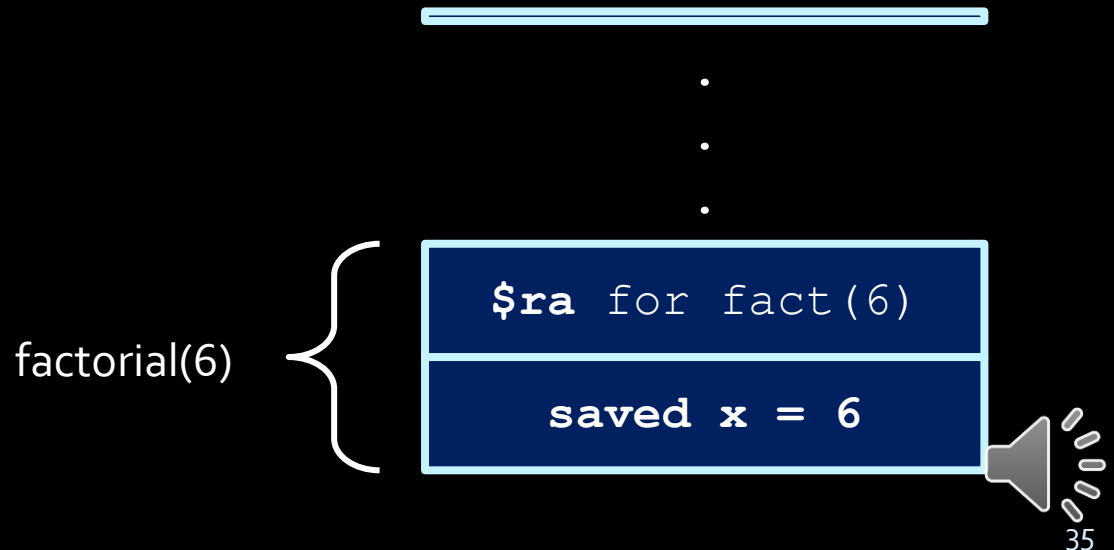
Now running: `fact(1)`



Factorial Stack View

Now running: `fact(2)`

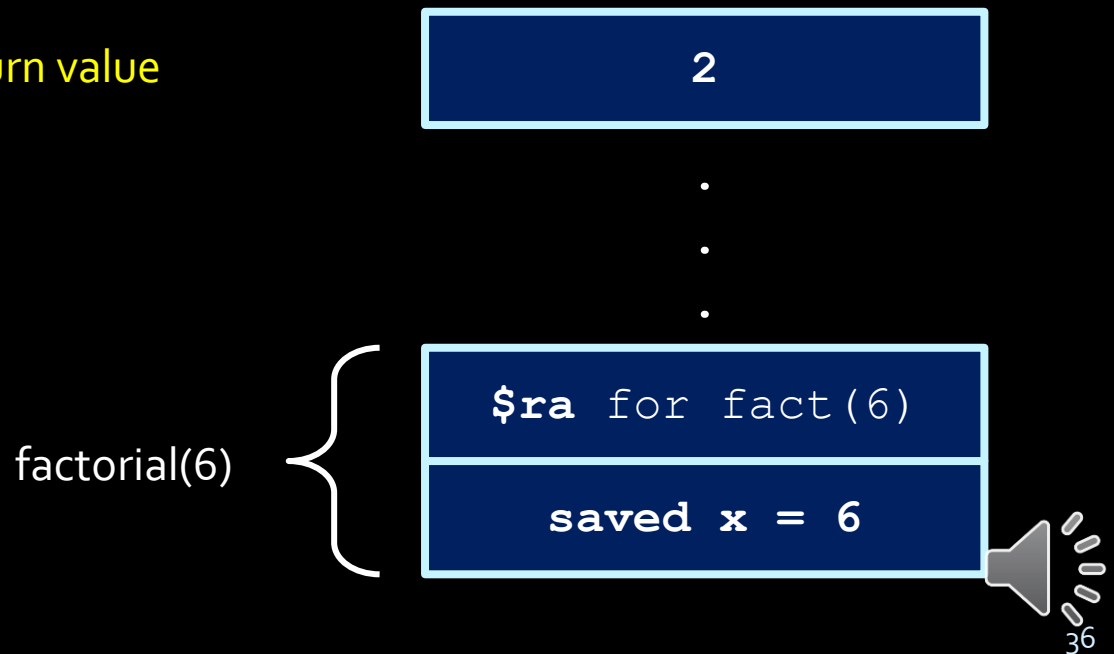
push $2 * \text{fact}(1)$ as return value



Factorial Stack View

Now running: `fact(2)`

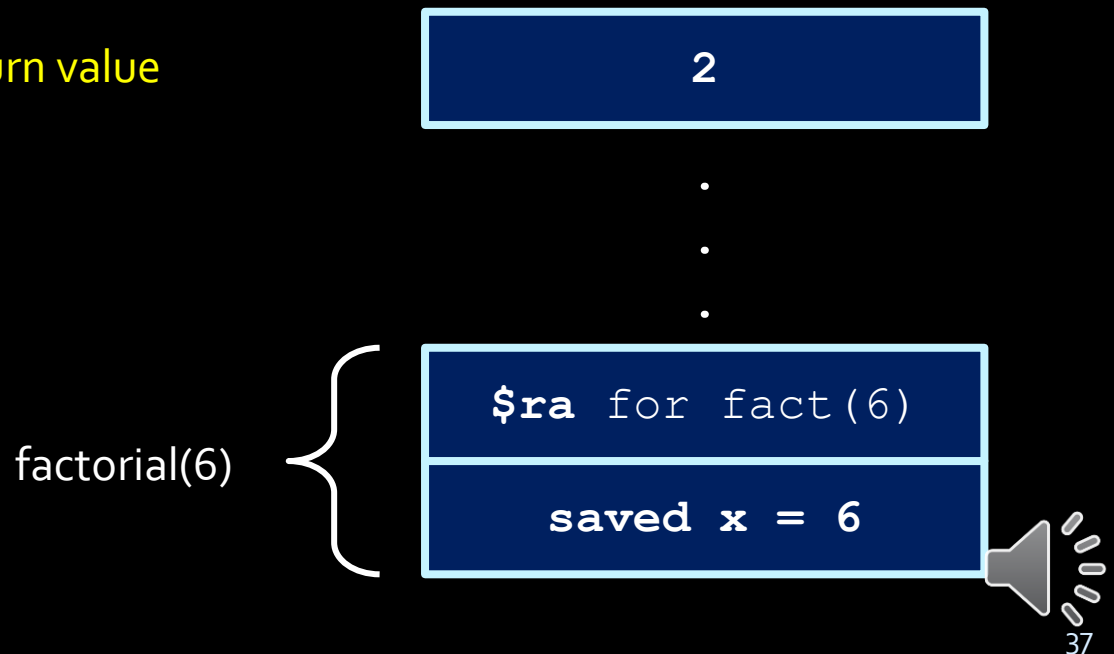
push $2 * \text{fact}(1)$ as return value



Factorial Stack View

Now running: `fact(2)`

push $2 * \text{fact}(1)$ as return value

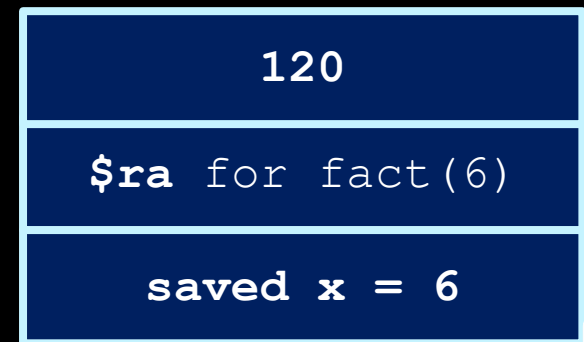


Factorial Stack View

Now running: `fact(5)`

return value of `fact(5)`

`factorial(6)`



Factorial Stack View

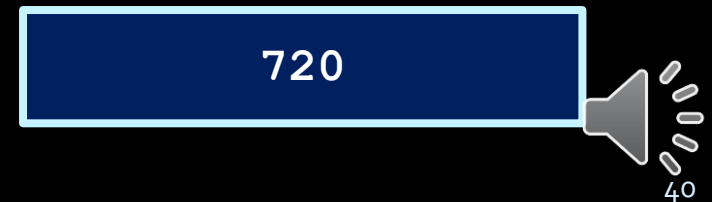
Now running: `fact(6)`

push return value $6 * \text{fact}(5)$

Factorial Stack View

Now running: `fact(6)`

push return value $6 * \text{fact}(5)$



Factorial Stack View

Now running: `main`

`main` pops result from stack (720)



Reflections on Recursion

- **Assembly does not understand recursion.**
- Assembly programs are just a linear sequence of assembly instructions.
- Recursion comes from:
 - Jumping to the beginning of the function over and over again and....
 - **Using the stack sensibly** to store and retrieve remembered values from the stack
- C function signatures help the compiler implement the recursion (what to push/pop)



The Stack is Finite


- You can recurse too much!
 - Maximum stack size **limits** the number of recursive calls that you can make.
 - Also depends on how much the function is using the stack.
- Exceeding the limit is called a **stack overflow**.
 - Older systems – overwrite variables, crashes...
 - On later systems, you might get an exception.
- Modern systems can grow the stack dynamically



Local Variables

- Sometimes you just need **local variables**
 - You ran out of registers.
 - Or you want a local array or local struct.
 - You are compiling C code and the programmer is using many local variables.

→ In next part!



```
int func(int a, int b) {  
    int local_array[256];  
    ...  
}
```