

◇ **Best before:** term test 1.

1. Do exercise 4 on page 71 of the course notes (about number of occurrences of an item in an array).
2. This question concerns programs that find the largest integer power of 2 which does not exceed m/n . I.e., we want to compute $2^{\lfloor \log_2(m/n) \rfloor}$.

- (a) Prove this iterative program to be correct with respect to its precondition/postcondition pair — prove both partial correctness and termination.

▷ Precondition: $m, n \in \mathbb{N}, 1 \leq n \leq m$.

▷ Postcondition: Return x such that $x = 2^t$ for some $t \in \mathbb{N}$ and $nx \leq m < 2nx$.

POW2ITER(m, n)

```

1   $x = 1; \quad y = 2 * n$ 
2  while  $y \leq m$ :
3       $y = 2 * y; \quad x = 2 * x$ 
4  return  $x$ 
```

- (b) Prove this recursive program to be correct with respect to its precondition/postcondition pair.

▷ Precondition: $m, n \in \mathbb{N}, 1 \leq n \leq m$.

▷ Postcondition: Return x such that $x = 2^t$ for some $t \in \mathbb{N}$ and $nx \leq m < 2nx$.

POW2REC(m, n)

```

1  if  $m < 2 * n$ :
2       $result = 1$ 
3  else:
4       $result = \text{POW2REC}(m, n * 2) * 2$ 
5  return  $result$ 
```

3. (a) Here is a recursive version of the program on page 59 of the course notes — the specification is the same.

▷ Precondition: $x, y \in \mathbb{N}$.

▷ Postcondition: True.

PAGE59(x, y)

```

1  if  $x == 0$  and  $y == 0$ :
2      return
3  elif  $x \neq 0$ :
4      PAGE59( $x - 1, y$ )
5  else:
6      PAGE59(16,  $y - 1$ )
```

Prove that PAGE59 is correct with respect to its specification.

- (b) Do exercises 8 and 9 on pages 72-73 of the course notes (about proving program termination).
 - (c) Similar to part (a), write recursive versions of the programs from part (b), and provide proofs of correctness for them.
Advice: For some programs, you may find it convenient to modify the precondition slightly.
4. (a) The following recursive program performs *long division* in base 2 and returns both quotient and remainder. Prove that it is correct with respect to its precondition/postcondition pair.

\triangleright Precondition: $x, y \in \mathbb{N}, y > 0$.
 \triangleright Postcondition: Return $\langle x \text{ div } y, x \text{ mod } y \rangle$.
 $D(x, y)$

```

1  if  $x == 0$ :
2      return  $\langle 0, 0 \rangle$ 
3  else:
4       $\langle q, r \rangle = D(x \text{ div } 2, y)$ 
5      if  $2 * r + (x \text{ mod } 2) < y$ :
6          return  $\langle 2 * q, 2 * r + (x \text{ mod } 2) \rangle$ 
7      else:
8          return  $\langle 2 * q + 1, 2 * r + (x \text{ mod } 2) - y \rangle$ 

```

- (b) Write an iterative version of the program $D(x, y)$ from part (a), then prove it to be correct with respect to the same precondition/postcondition pair.

Advice: You may find it helpful to use the POW2ITER program from question 2.

5. Consider the problem of finding the largest index of a largest element in a slice of a list.

- (a) In the following recursive program, A is a list of integers and all other variables are integers. Prove it to be correct with respect to its precondition/postcondition pair.

\triangleright Precondition: $0 \leq f < l \leq \text{len}(A)$.
 \triangleright Postcondition: Return an integer u such that
 \triangleright (i) $f \leq u < l$,
 \triangleright (ii) $A[u]$ is the largest value in $A[f : l]$, and
 \triangleright (iii) $A[u]$ is greater than every integer in $A[u + 1 : l]$.
 $\text{MAX}(A, f, l)$

```

1  if  $f + 1 == l$ :
2      return  $f$ 
3  else:
4       $m = (f + l) \text{ div } 2$ 
5       $x = \text{MAX}(A, f, m)$ 
6       $y = \text{MAX}(A, m, l)$ 
7      if  $A[x] > A[y]$ :
8          return  $x$ 
9      else:
10         return  $y$ 

```

- (b) Write an iterative version of the program $\text{MAX}(A, f, l)$ from part (a), then prove it to be correct with respect to the same precondition/postcondition pair.

6. For this question we consider numbers represented in base 3 and whether they are divisible by 4.

Let $\Sigma_3 = \{0, 1, 2\}$. For a string $x \in \Sigma_3^*$ (we call x a *ternary* string), we define $V(x)$ to be the number represented in base 3 by x . More formally, if $x = x[0]x[1] \cdots x[n-1]$, where $n > 0$ and each $x[i] \in \Sigma_3$,

then $V(x) = \sum_{i=0}^{n-1} (x[n-1-i] \cdot 3^i)$.¹ For example, $V(021101) = 199$.

¹Note that a reference to an element of x can mean either a symbol from Σ_3 or the numeric value of that symbol. E.g., if x is the string 10021, then $x[3]$ is both the symbol 2 and the number 2. You should be able to tell which by its context.

- (a) Prove that the following program is correct with respect to its precondition/postcondition pair.

▷ Precondition: $x \in \Sigma_3^*$ and $x \neq \epsilon$ (i.e., x is a nonempty ternary string).
 ▷ Postcondition: Return **True** if $V(x) \bmod 4 = 0$; Otherwise return **False**.
 BASE3DIVBY4(x)

```

1     $t = x[0]; \quad i = 1$ 
2    while  $i < \text{len}(x)$ :
3        if  $t == 0$ :
4             $t = x[i]$ 
5        else:
6             $t = 4 - t + x[i]$ 
7            if  $t \geq 4$ :  $t = t - 4$ 
8         $i = i + 1$ 
9    return ( $t == 0$ )
```

- (b) Write a recursive version of the program BASE3DIVBY4(x) from part (a), then prove it to be correct with respect to the same precondition/postcondition pair.

7. *Extra — a programming question!*

The following definition may look familiar.

We say that a nonempty list t of distinct integers has the *bifurcate property* iff this condition holds.

Bifurcate property: For any nonempty slice $t[p : q]$, there is an integer k such that

- $p < k \leq q$,
- every integer in $t[p + 1 : k]$ is less than $t[p]$,
- every integer in $t[k : q]$ is greater than $t[p]$.

Write a program that takes a nonempty list t of distinct integers as input and returns True or False depending on whether t has the bifurcate property, then prove that it is correct.

- Make your program as efficient as possible.
 There are $O(n^2)$ slices in a list of length n . It takes $O(n)$ time to check a slice for the property. So it seems the total complexity of this program is $O(n^3)$. Can you do better?
- You should consider breaking the program into smaller pieces — *modular design!*
 It could simplify your proof.
- Try using only recursion, only iteration or perhaps a combination of both.