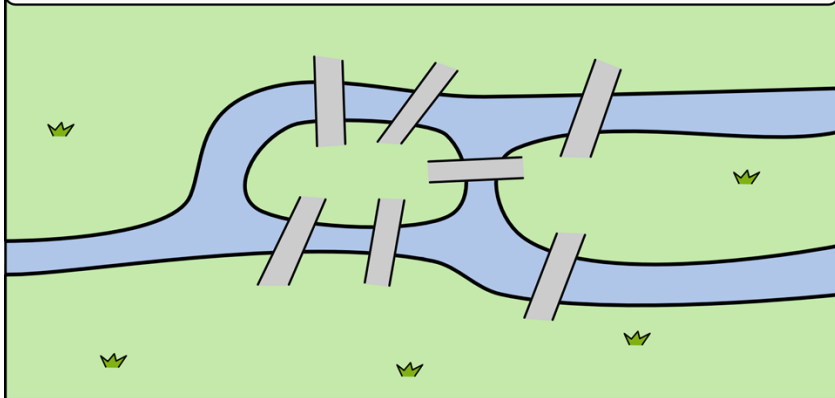


Eulerian Graphs

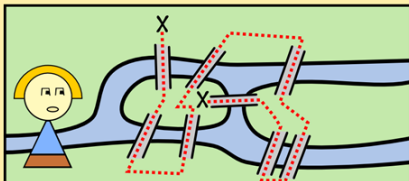
The Seven Bridges of Königsberg

Below is the city of Königsberg with four land masses and seven bridges connecting the various land masses. Can you find a walk through the city of Königsberg that crosses each bridge exactly once? You may start at any land mass you wish but may only travel between land masses by using a bridge.



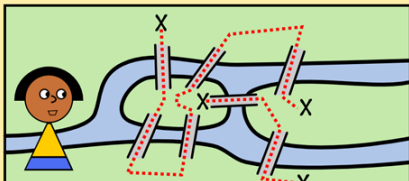
Puzzle: The Seven Bridges of Königsberg

THE ENGINEER'S SOLUTION



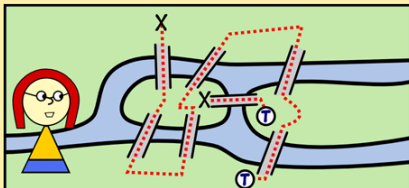
I WOULD BUILD ANOTHER BRIDGE!

THE BIOTECHNOLOGIST'S SOLUTION



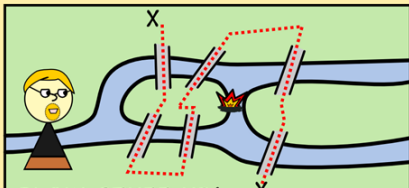
I WOULD CLONE MYSELF.

THE PHYSICIST'S SOLUTION



I WOULD TELEPORT.

THE MYTHBUSTER'S SOLUTION

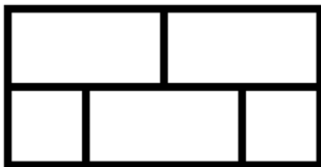


BLOW STUFF UP!

The above are jokes. Following the given rules, graph theory can show it is **impossible**.

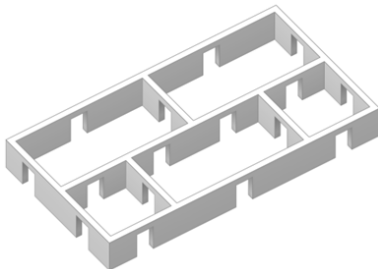
The Five Rooms Puzzle

Question: In the diagram, can you draw a path that goes through each wall exactly once? The path cannot go through any intersections! (There are 16 walls.)



Alternate formulation with doors:

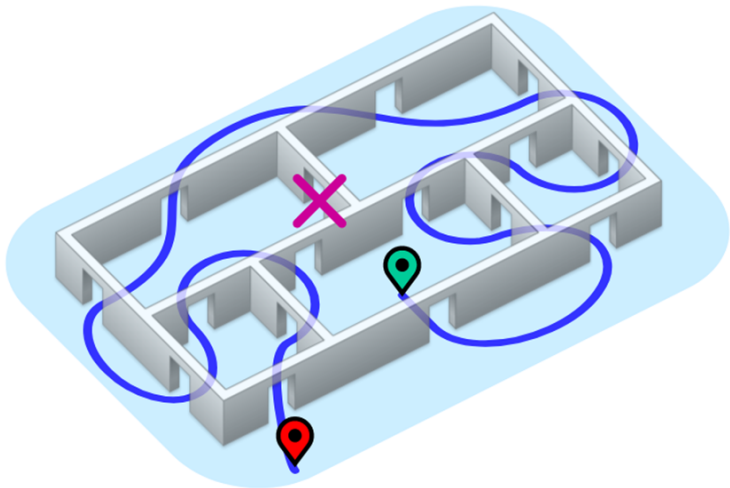
Starting in any room, can you walk through every door exactly once?



Source: Image by CMG Lee (Cmglee)

The Five Rooms Puzzle

The below shows an attempted solution, but it misses a door.

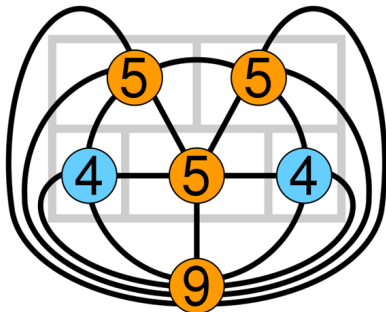
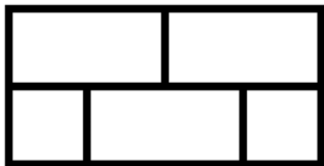


Source: Image by CMG Lee (Cmglee)

Side Note: The link above shows how it is possible if the rooms are on a torus (donut). 5/12

The Five Rooms Puzzle

We can use graph theory to show the puzzle is impossible.



- For each room (including the outer room), draw a vertex.
- Connect two vertices if there is a door between them.
- This produces a **multigraph** with some pairs of vertices having multiple edges between them. (The top right shows the degree of each vertex.)

Before proving any theorems about impossibility, we require some terminology!

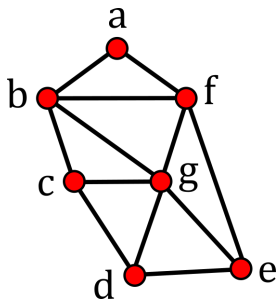
Definitions

- A walk in a (multi)-graph G is a sequence

$$W = v_0 e_1 v_1 e_2 v_2 \cdots v_{k-1} e_k v_k$$

whose terms alternate between vertices and edges (**not** necessarily distinct) such that $e_i = v_{i-1} v_i$ for $1 \leq i \leq k$. When G is a simple graph, we write $W = v_0 v_1 \cdots v_k$.

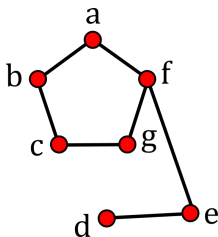
- The length of a walk is the number of edges it contains.
- A trail is a walk such that all of its edges are distinct.
- A path is a walk such that all of its vertices and edges are distinct.
- A walk is closed if the initial and terminal vertices are the same (i.e., $v_0 = v_k$).
- A walk from vertex x to vertex y is called an xy -walk.
- An Euler trail is a trail that visits every edge exactly once.
- An Euler tour is a closed Euler trail.
- A graph that has an Euler tour is called Eulerian.



Consider the graph G :

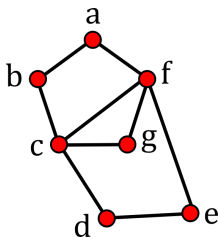
- $W = abcgbaf$ is a walk of length 6. (It is an af -walk.)
 - It is **not** a trail (edge ab is used twice) and **not** closed.
- $W = abcgb$ is a trail of length 4.
 - It is also a walk (and has distinct edges). It is **not** closed.
- $W = abgfabfa$ is a closed walk of length 7.
 - It is **not** a closed trail (edges ab and af are each repeated twice).
- $W = bcgbfab$ is a closed trail of length 6.
 - It is also a closed walk.
- $P = abcdefg$ is a path of length 6.

Example



An **Euler trail** (every edge used once) is $W = defabcgf$.

Example



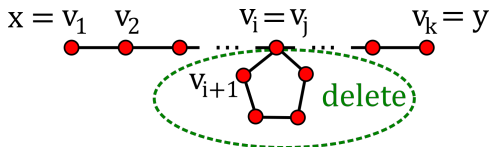
An **Euler tour** (closed Euler trail) is $W = abcdefgcfa$.

Theorem

Let G be a connected graph with $x, y \in V(G)$ and W be an xy -walk. Then W contains a path between x and y .

Proof.

- Let $W = v_1 v_2 \cdots v_k$ be an xy -walk where $v_1 = x$ and $v_k = y$.
- If W has no repeated vertices, then W is in fact a path and we are done.
- Otherwise, suppose W has a repeated vertex, say $v_i = v_j$ for some $i < j$.



- Since $v_i = v_j$, the walk W looks like

$$W = v_1 v_2 \cdots v_{i-1} v_i v_{i+1} \cdots v_j v_{j+1} \cdots v_k.$$

- Then by deleting part of W we get a walk W' of smaller length:

$$W = v_1 v_2 \cdots v_{i-1} v_i \underbrace{v_{i+1} \cdots v_{j-1} v_j}_{\text{delete}} v_{j+1} \cdots v_k \rightarrow W' = v_1 v_2 \cdots v_{i-1} v_i v_{j+1} \cdots v_k.$$

- If W' has no repeated vertices, then it is a path between x and y , otherwise, repeat the above process.

Theorem

Let G be a connected graph (or multigraph). Then G is Eulerian if and only if every vertex of G has even degree.

Proof. See Mike.

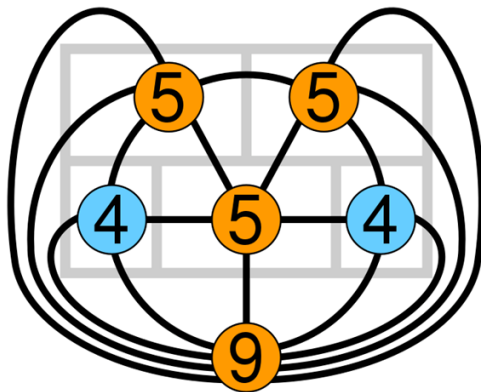
Theorem

Let G be a connected graph (or multigraph). Then G has an Euler trail if and only if G has at most two vertices of odd degree.

Proof. See Mike.

The Five Rooms Puzzle

Recall the multigraph for the five rooms puzzle:



By Euler's theorem, this graph does not have an Euler trail or Euler tour since it has four vertices of odd degree.

Therefore, the puzzle is **impossible**.