

# Transshipment Problem

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# 01

## Background

Mission, Goal, and Data



# Background



## Mission

To meet customer demand and reduce expected long-run average cost by replenishment (order-up-to  $S$ ) policy and transshipment policy.

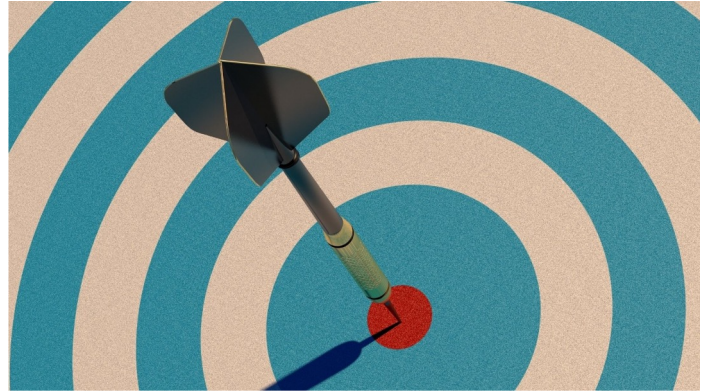
Given an order-up-to  $S$  policy, we show how to determine an optimal transshipment policy, using an all-in-one linear programming and a benders decomposition model.



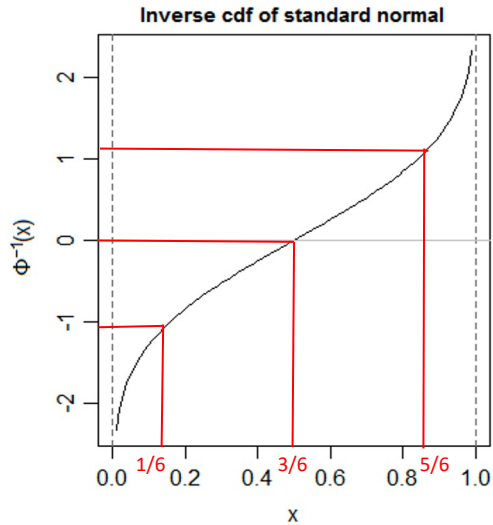
# Background

## Goal

Solve and report **objective** and **first stage decisions**:  
**order-up-to quantities**



# Background: Data



## Demand (Independent & Normally Distributed)

Location	1	2	3	4	5	6	7
Mean	100	200	150	170	180	170	170
Std Dev	20	50	30	50	40	30	50

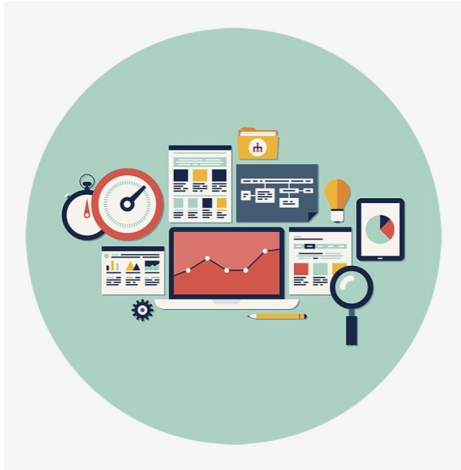
## Discretize the Normal distribution

- Low - Medium - High
- Normal density  $\rightarrow$  3 intervals with equal probability
- $[1/6, 3/6, 5/6]$  quantile

Equal probability for each of the  $3^7$  scenarios.



# Background: Data



## Cost (Uniform)

Symbol	Unit	Meaning	Value
h	dollar/unit	Per unit holding cost when there is excess inventory at one location	1.0
c	dollar/unit	Per unit transshipment cost to move inventory from one location to satisfy demand at another location	0.5
p	dollar/unit	Per unit shortage cost when demand cannot be met at one location	4.0

## Assumption

- Transshipment cost, holding cost and shortage cost are uniform across all locations



# 02

## All-In-One LP

Two-stage stochastic, All-in-one model





# All-In-One LP

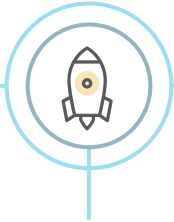
**1. Parameters**



**3. Objective Function**



**2. Decision Variables**



**4. Constraints**



# Two-stage Stochastic Problem

**Stage 1: Find optimal order-up-to quantity**

$$\begin{aligned} \min \quad & c^\top x + E[h(x, \tilde{\omega})] \\ \text{s.t.} \quad & Ax \leq b \end{aligned}$$

Decision Variable:  $x$

**Stage 2: Find transshipment quantity to minimize COST**

$$\begin{aligned} h(x, \omega) = \min \quad & g^\top y \\ \text{s.t.} \quad & Wy = r(\omega) - Tx \\ & y \geq 0 \end{aligned}$$

Decision Variable:  $y$   
Random Variable:  $\omega$



# All - In - One Model

## 1. Parameters

1.  $\mathbf{d_{3 \times 7}}$ : Demand for 7 locations with low, medium, high scenarios
2.  $\mathbf{h}$ : Per unit holding cost when there is excess inventory for all locations
3.  $\mathbf{p}$ : Per unit transshipment cost from one location to satisfy demand at another location
4.  $\mathbf{c}$ : Per unit shortage cost when demand cannot be met for all locations

# All - In - One Model

## 2. Decision Variables

1.  $\mathbf{s}_{[1 \times 7]}$ : Order-up-to quantity for each of the 7 locations
  2.  $\mathbf{e}_{[3^7 \times 7]}$ : Ending inventory held at retailer  $i$  for scenario  $s$
  3.  $\mathbf{f}_{[3^7 \times 7]}$ : Stock at retailer  $i$  used to satisfy demand at retailer  $i$  for scenario  $s$
  4.  $\mathbf{q}_{[3^7 \times 7]}$ : Inventory at retailer  $i$  increased through replenishment for scenario  $s$
  5.  $\mathbf{r}_{[3^7 \times 7]}$ : Amount of shortage met after replenishment at retailer  $i$  for scenario  $s$
  6.  $\mathbf{t}_{[3^7 \times 7 \times 7]}$ : Stock at retailer  $i$  used to meet demand at retailer  $j$  for scenario  $s$
-

# All - In - One Model

## 3. Objective

$$\min \sum_s Prob_s \times \left( \sum_i \underbrace{h \times e_{si}}_{\text{Holding Cost}} + \sum_{i \neq j} \underbrace{c \times t_{sij}}_{\text{Shipment Cost}} + \sum_i \underbrace{p \times r_{si}}_{\text{Shortage Cost}} \right)$$

Based on the equal probability of each scenarios

$$E[h(S, \tilde{D})] = \sum_s \left( \sum_i h \times e_{si} + \sum_{i \neq j} c \times t_{sij} + \sum_i p \times r_{si} \right) / M$$

$$M = 3^7$$

# All - In - One Model

## 4. Constraint

$$\text{s.t. } f_i + \sum_{j \neq i} t_{sij} + e_{si} = s_i, \forall i, s$$

$$f_i + \sum_{j \neq i} t_{sji} + r_{si} = d_i, \forall i, s$$

$$\sum_i r_{si} + \sum_i q_{si} = \sum_i d_{si}, \forall s$$

$$e_{si} + q_{si} = s_i, \forall i, s$$

$$e_{si}, f_{si}, q_{si}, r_{si}, s_i, t_{sij} \geq 0, \forall i, j, s$$

# 03

## Benders Decomposition

Master problem, Sub-problem, Build Cut



# Benders Decomposition

## Solve Master Problem

Check for stopping  
criteria

1

## Solve Subproblem

Obtain the dual  
extreme points

2

## Build Benders Cut

Add cut to master;  
repeat

3





# Benders Decomposition

## 1. Objective

Master problem:

$$\min_S \eta$$

Sub-problem:

$$\min h(S, \tilde{D})$$

$$E[h(S, \tilde{D})] = \sum_i h \times e_i + \sum_{i \neq j} c \times t_{ij} + \sum_i p \times r_i$$

The parameters and variable used in this model are the same as the all-in-one model.

# Benders Decomposition

## 2. Constraints

### Master problem

Add one more constraint (cut) in each iteration k

$$\eta \geq \alpha_k + \beta_k \times S^k$$

$$\alpha = \sum_i d_i * (\underbrace{\pi_{Mi}} + \underbrace{\pi_{Ri}})$$

$$\beta = \underbrace{\pi_B} + \underbrace{\pi_E}$$

$\pi$  : The dual variables of the constraints  
In the sub-problems

### Sub-problem

Constraint B

$$f_i + \sum_{j \neq i} t_{ij} + e_i = s_i, \forall i$$

Constraint M

$$f_i + \sum_{j \neq i} t_{ji} + r_i = d_i, \forall i$$

Constraint R

$$\sum_i r_i + \sum_i q_i = \sum_i d_i$$

Constraint E

$$e_i + q_i = s_i, \forall i$$
$$e_i, f_i, q_i, r_i, t_{ij} \geq 0, \forall i, j$$

# Benders Decomposition

## 3. Algorithm

In the  $k$ -th iteration:

- Solve the master problem to get a **lower bound**  $LB_k$  and a **first stage decision**  $S_k$ .
- Use the new first stage decision to solve the sub-problem and get **the coefficients**  $(\alpha_k, \beta_k)$  of the new cut, and **upper bound**  $UB_k$ .
- Add a **new cut (constraint)**  $\eta \geq \alpha_k + \beta_k * S_k$  to the master problem.
- If  $UB_k - LB_k \leq \text{Stopping Criteria}$ , then stop; else, go on to the next iteration.

**Stopping Criteria: Upper Bound-Lower Bound  $\leq 1e-6$**

# 04

## Results Comparison

All-in-one v.s. Bender's, All-in-one LP v.s. IP



# Results Comparison

## All-In-One v.s. Benders

Model	Time (second)	Objective Value	Value of s
All-In-One	6.598277	144.5956842831869	Location 1: 107.2557 Location 2: 213.3020 Location 3: 158.4650 Location 4: 183.3020 Location 5: 190.8835 Location 6: 180.8835 Location 7: 183.3020
Benders	99.032718 ( 52 iterations)	144.59568428320085	Location 1: 99.9999 Location 2: 219.3484 Location 3: 179.0226 Location 4: 169.9999 Location 5: 180.0000 Location 6: 199.0226 Location 7: 169.9999



# Results Comparison w/wo integer restriction

Model	Time (second)	Objective Value	Value of s
All-In-One	6.598277	144.5956842831869	Location 1: 107.2557 Location 2: 213.3020 Location 3: 158.4650 Location 4: 183.3020 Location 5: 190.8835 Location 6: 180.8835 Location 7: 183.3020
All-In-One (Int) * Decision variables are integers; Demands are rounded to integers	8.608392	143.93209876543145	Location 1: 103 Location 2: 210 Location 3: 159 Location 4: 188 Location 5: 190 Location 6: 179 Location 7: 188



# Plan

1. 4-level scenarios: Very Low - Low - High - Very High

[1/8, 3/8, 5/8, 7/8] quantile

Model	Time (second)	Objective Value	Value of s
All-In-One (4-level of scenarios)	\	150.56341192779723	Location 1: 106.3728 Location 2: 215.9320 Location 3: 159.5592 Location 4: 185.9320 Location 5: 192.7456 Location 6: 179.5592 Location 7: 185.9320

2. Compare the second stage decisions of Bender's method and all-in-one model.

3. Holding cost, transshipment cost, shortage cost could vary at different locations.



# Code



**Link:**

<https://github.com/YuanLuo/Transshipment.git>





# THANKS!

Team 2

