

# Important thing for search

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## 1 Introduction

Just list the thing I think it is important but hard to remember. If want to use it, you need to fully understand all things in lectures 16–19.

## 2 pseudo-code

This section is the pseudo code for the template search algorithms.

Listing 1: pseudo code for search algorithm template

```
1 0.: initialize
2   agenda = [ [start] ]           # a list of paths or partial solutions
3   [extended_list = {}];          # optional: closed list of nodes
4   while agenda is not empty:
5       1. path = agenda.pop(0)     # remove first element from agenda
6       2. if is-path-to-goal(path, goal)
7           return path
8       3. otherwise extend path [if the ending node is not in extended_list]
9           [add the ending node to extended_list]
10          for each connected node (child node of the ending node of path)
11              make a new path that extends the connected node
12          4. reject all new paths from step 3 that contain cycles
13          5. add new paths from step 4 to agenda and reorganise agenda
14              # algorithms differ here!
15 6. return failure
```

Change part 5 can create new algorithms.

### 3 Comparison of different search algorithms

	Complete	Optimal	Time complexity	Space complexity	agenda reorganisation
DFS	No	No	$O(b^d)$	$O(bd)$	paths added to the front agenda; or stack agenda
BFS	Yes	No in general	$O(b^d)$	$O(b^d)$	paths added to the back of agenda; or queue agenda
NDS	No	No	$O(b^d)$	$O(b^d)$	paths randomly inserted to the agenda
UCS	Yes	Yes	$O(b^{C^*/\epsilon})$	$O(b^{C^*/\epsilon})$	paths added agenda then sort the agenda based on path costs (optional: trim down agenda by only retaining the most promising ones, i.e. if there are multiple paths ending with the same node, keep the one with the optimal cost); or PriorityQueue agenda

For UCS,  $C^*$  is the optimal cost,  $\epsilon$  is the lower bound of action costs,

So the optimal depth  $d = C^*/\epsilon$