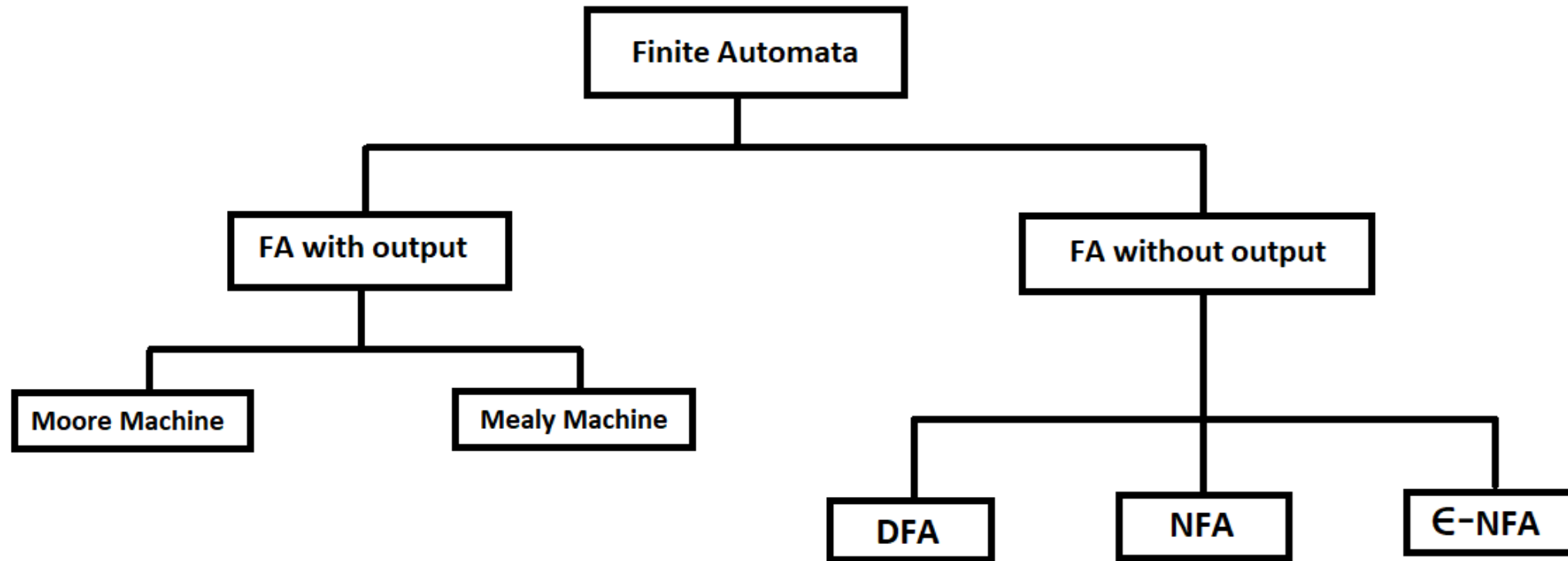




FINITE STATE MACHINE (Finite Automata)

FINITE STATE MACHINE

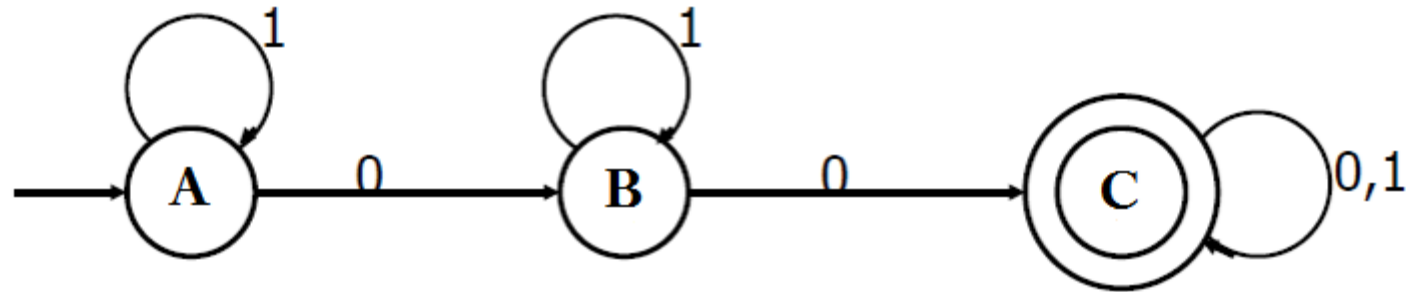


DETERMINISTIC FINITE AUTOMATA (DFA)

- it is a finite-state machine that accepts and rejects strings of symbols and only produces a unique computation (or run) of the automaton for each input string
- it is the simplest model of computation
- it has a very limited memory

DETERMINISTIC FINITE AUTOMATA (DFA)

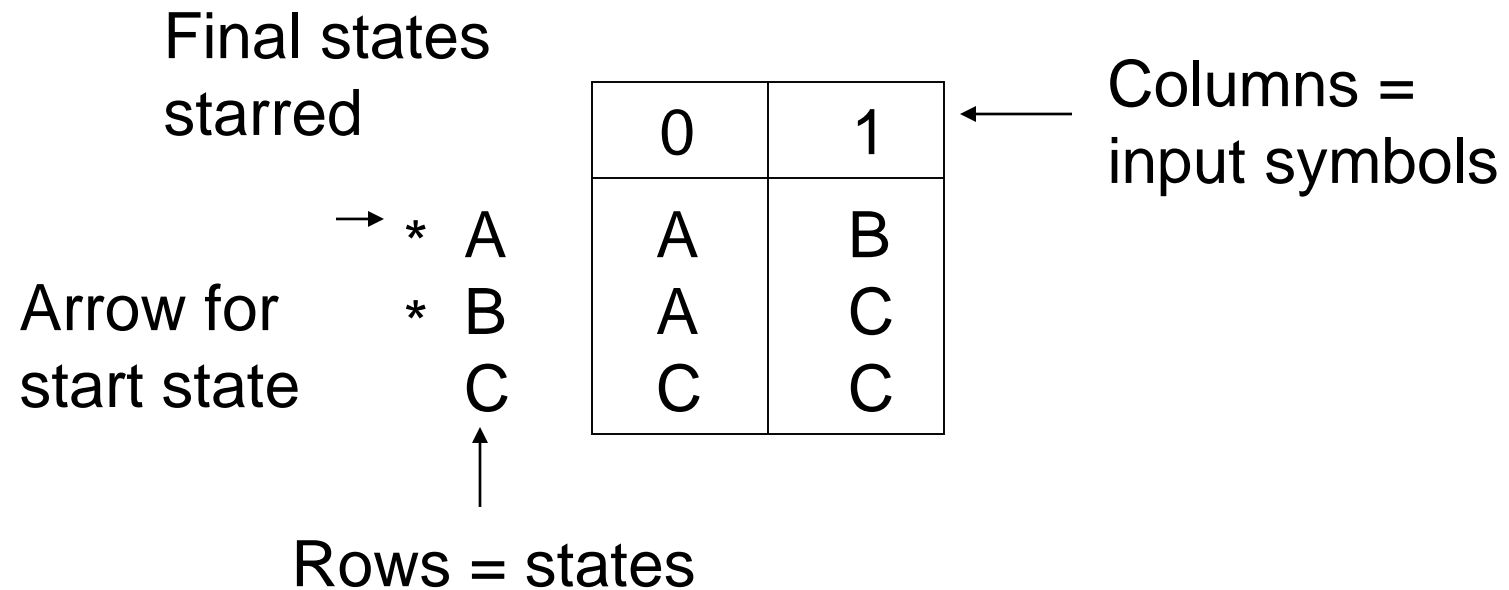
- DFAs are easiest to present pictorially: Such a graph is called a state transition diagram



FORMAL DEFINITION

- A deterministic finite automaton M is a 5-tuple, $(Q, \Sigma, \delta, q_0, F)$, consisting of:
 - Q = set of all states
 - Σ = input symbols
 - δ = transition function ($Q \times \Sigma \rightarrow Q$)
 - q_0 = start state / initial state
 - F = set of final states

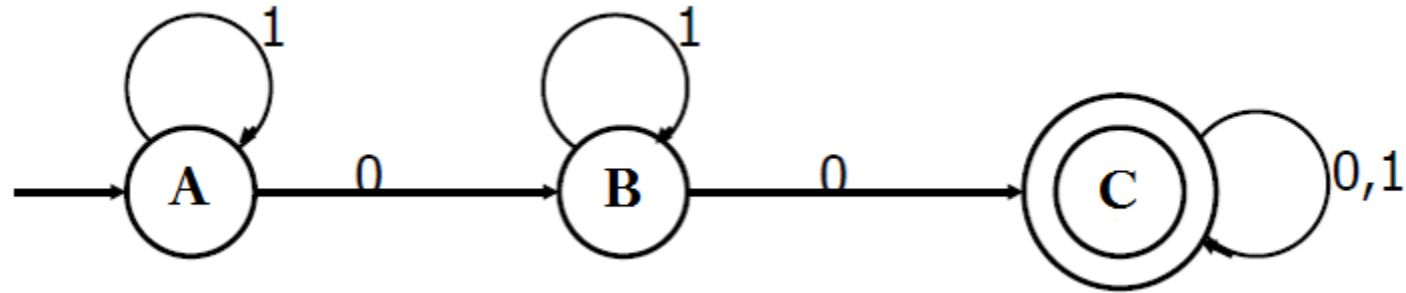
ALTERNATIVE REPRESENTATION: TRANSITION TABLE



EXAMPLE 1

- From the given DFA diagram, describe a DFA accepting the language by:

- a.) determining the 5-tuple
- b.) building a DFA transition table



5-tuple

- $Q = \{A, B, C\}$
- $\Sigma = \{0, 1\}$
- $\delta = \{A \times 0 \rightarrow B; A \times 1 \rightarrow A; B \times 0 \rightarrow C; B \times 1 \rightarrow B; C \times 0 \rightarrow C; C \times 1 \rightarrow C\}$
- $q_0 = A$
- $F = C$

Transition Table

	0	1
→ A	B	A
B	C	B
* C	C	C



EXAMPLE 2

- For each of the following languages, describe a DFA accepting the language by drawing a DFA diagram
 - L1 = set of all strings over $\{0,1\}$ that starts with 0
 - L2 = set of all strings over $\{0,1\}$ of length 2