第二十讲:分布式系统

第 4 节: LegoOS

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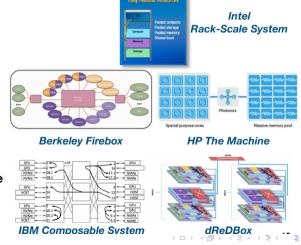
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Paper: LegoOS

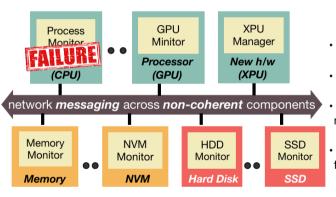
Slides: LegoOS

分布式 I/O 计算逐渐浮现在大型数据中心和移动终端领域。

- Network is faster
 - InfiniBand (200Gbps, 600ns)
 - Optical Fabric (400Gbps, 100ns)
- More processing power at device
 - SmartNIC, SmartSSD, PIM
- · Network interface closer to device
 - · Omni-Path, Innova-2



出现了新型的 splitkernel 架构

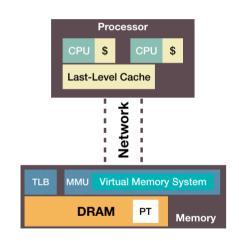


- Split OS functions into monitors
- Run each monitor at h/w device
- Network messaging across non-coherent components
- Distributed resource mgmt and failure handling

LegoOS: 把 CPU, Memory, Storage 分布在不同机器上,通过高速网络 RDMA 形成一个虚拟的大机器 OS 抽象:

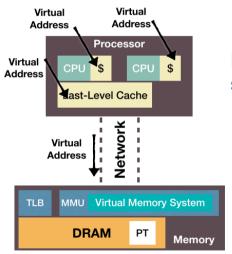
- virtual Nodes (vNodes) -> hardware devices
- Similar semantics to virtual machines
- Unique vID, vIP, storage mount point
- Can run on multiple processor, memory, and storage components

Separate Processor and Memory



Separate and move virtual memory system to memory component

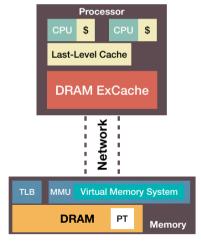
Separate Processor and Memory



Processor components only see virtual memory addresses
All levels of cache are virtual cache.

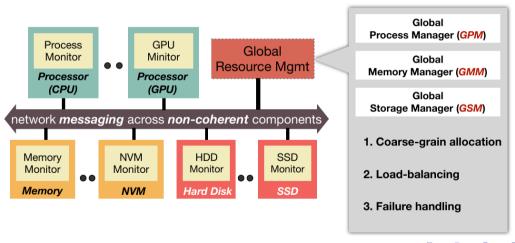
Memory components manage virtual and physical memory

Add Extended Cache at Processor



- · Add small DRAM/HBM at processor
- Use it as Extended Cache, or ExCache
 - Software and hardware co-managed
 - Inclusive
 - · Virtual cache

Distributed Resource Management



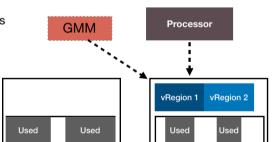
Distributed Memory Management

0 User Virtual Address Space max

vRegion 1 vRegion 2 vRegion 3

fix-sized, coarse-grain virtual region (vRegion) (e.g., 1GB)

- GMM assigns vRegions to mem components
- On virtual mem alloc syscalls (e.g., mmap)
- Make decisions based on global loads
- Owner of a vRegion
 - Fine-grained virtual memory allocation
 - On-demand physical memory allocation
 - Handle memory accesses



mmap 1.5GB

write 1GB

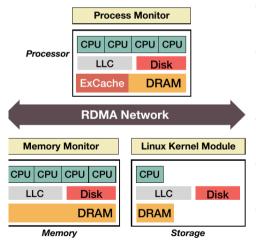
(Physical Memory)

Memory (M1)

(Physical Memory)

Memory (M2)

Implementation



- Status
 - 206K SLOC, runs on x86-64, 113 common Linux syscalls
- Processor
 - Reserve DRAM as ExCache (4KB page as cache line)
 - · h/w only on hit path, s/w managed miss path
- Memory
 - Limit number of cores, kernel-space only
- Storage/Global Resource Monitors
 - · Implemented as kernel modules on Linux
- Network
 - RDMA RPC stack based on LITE [SOSP'17]