

1 Judgement

Solution.1

Correct

Solution.2

Correct

Solution.3

Wrong, wait only → partly execute

Solution.4

Wrong, then → after merging and fusion

Solution.5

Correct

Solution.6

Correct

Solution.7

Correct

Solution.8

Wrong, is influenced by → is not influenced by

Solution.9

Correct

Solution.10

Wrong, linear → 1-dimension, non-linear

2 Brief Answer

Solution.1

- (1) **Jitter** might happen (interconnect between memories). To avoid it, we could add more physical memory block or stall more processes.
- (2) OS acts normally, we could level up numbers of process to improve parallelism and utilization of resources.
- (3) Processor and memory are working in a low level rate. We could improve the number of processes to execute.

Solution.2

a single page records the number of page table entry:

$$4KB/4B = 1024 = 2^{10}$$

64-bit VM needs the number of page

$$2^{64}/2^{12} = 2^{52}$$

so at least $2^{60} > 2^{52}$, 6 level page table

Solution.3

The number of page table entry is

$$2^{48}/2^{13} = 2^{35}$$

, 2^{35} inverted page table needs $2^{35}/2^{13} = 2^{19}$, 2^{19} page table entries.

Solution.4

Working Set is the processes we're gonna visit in a specific time of period(window). So, $t_1: \{1, 2, 3, 6, 7, 8, 9\}$, $t_2: \{3, 4\}$

(Referred to Link)

Solution.5

(1) The size of **page** is $2^{12}B$.

The size of **page frame** is $2^{12}B$.

The size of **VM** is $2^{32}B$.

(2) To restore page content and page table, there's gonna need $\frac{2^{10} \times 8}{2^{12}} + \frac{2^{20} \times 8}{2^{12}} = 2050$

(3) 1, that's because the biggest 10 bits of 01000000H and 01114096H is both 0000000100 and they should only go to the same Page Directory due to principles of locality.

3 Analysis

Solution.1

4 conditions to cause deadlocks:

- Mutual Exclusion
- Hold and Wait
- No preemption
- Circular wait

Banker Algorithm:

We maintain **Available** and **Need** lists.

$\text{Need}(A, B, C) = (0, 2, 5), (3, 0, 0), (2, 0, 4), (3, 1, 2), (4, 7, 1)$

$\text{Available}(A, B, C) = (3, 0, 1)$

According to that, resources could be allocated to P_1 , after executing, available is (4, 1, 3). Then resources could be allocated to P_3 , after executing, available is (4, 3, 4). Next, Then resources could be allocated

to P_2 , after executing, available is (4,6,4). At this point, both P_0 and P_4 is no longer to be satisfied. So [deadlock happens inevitably](#).

Solution.2

- (1)
- [first-fit](#) 212K - #2 417K - #5 112K - #2 426K - NULL
 - [best-fit](#) 212K - #4 417K - #2 112K - #3 426K - #5
 - [worst-fit](#) 212K - #5 417K - #2 112K - #3 426K - NULL
- (2) Best-Fit is the best. The other 2 will have an allocation error.

Solution.3

There're $100 \times 100 = 10000$ data in total. Since the capacity of a single page is 200, the number of page is 50. The elements are scheduled in rows so every 2 rows will be included in a page.

For Program A, as it executes in rows so every 2 data rows take a page. In total, [50](#) page faults are reserved.

For Program B, as it executes in columns so each column take 50 page faults, where [50,000](#) is the cumulative lies.

Solution.4

- [FIFO](#)

Page fault happens when there's no matching pages in the page table. In FIFO, there's total [14](#) page faults, so the missing rate is [30%](#). The pass process is: {1}, {1,2}, {1,2,3}, {1,2,3,4}, {5,2,3,4}, {5,6,3,4}, {5,6,2,4}, {5,6,2,1}, {3,6,2,1}, {3,7,2,1}, {3,7,6,1}, {3,7,6,2}, {1,7,6,2}, {1,3,6,2}

- [OPT](#)

There's total [8](#) page faults, so the missing rate is [60%](#). The elaborate pass process is {1}, {1,2}, {1,2,3} {1,2,3,4} {1,2,3,5} {1,2,3,6} {7,2,3,6} {1,2,3,6}

- [LRU](#).

There's total [10](#) page faults, so the missing rate is [50%](#). The elaborate pass process is {1} {1,2} {1,2,3} {1,2,3,4} {1,2,5,4} {1,2,5,6} {1,2,3,6} {1,2,3,7} {6,2,3,7} {6,2,3,1}

- [CLOCK](#)

There's total [10](#) page faults, so the missing rate is [50%](#). The elaborate pass process is {1} {1,2} {1,2,3} {1,2,3,4} {1,2,5,4} {1,2,5,6} {1,2,3,6} {1,2,3,7} {6,2,3,7} {6,2,3,1}