03. Process



Backgrounds

- a program in execution
- owns the entire computer

Process Store

Concepts

• how process stored in memory

o consists of

- text section / code section
- PC
- stack: contains temporary data (function args...)
- data section: contains constant, variables
- heap: allocate memory dynamically

max stack heap data text

0

Process state

- o new
- running
- waiting
- ready
- terminate

PCB

- a data structure
- consists
 - process state
 - process number: pid
 - program counter
 - registers:
 - memory limits
 - list of open files

process state
process number
program counter
registers
memory limits
list of open files

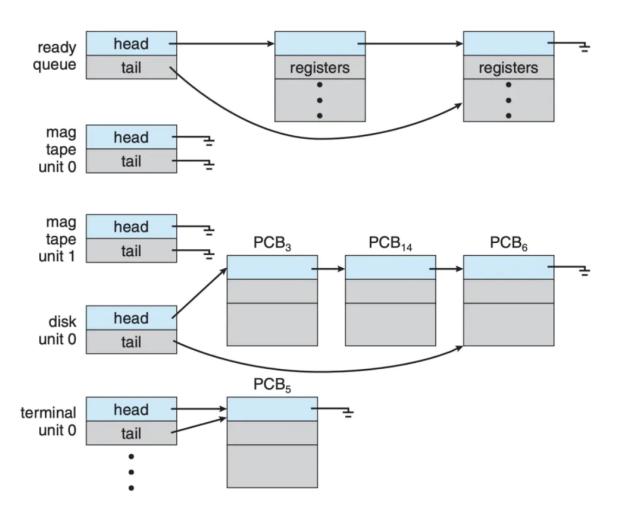
Process in Linux

• task_struct in linux/sched.h>

```
pid t pid; /* process identifier */
long state; /* state of the process */
unsigned int time slice /* scheduling information */
struct task_struct *parent; /* this process's parent */
struct list_head children; /* this process's children */
struct files_struct *files; /* list of open files */
struct mm struct *mm; /* address space of this pro */
```

Process Scheduling

 $Queue \; \leftarrow \; Link \; Method$



- job queue
 - when a process get into a system, it is added to job queue
- ready queue
 - contains all ready processes waiting to be executed
- device queue
 - shared device has a queue to store processes
- job queue and ready queue are implemented by a link
 - o 2 pointers → PCB
 - head
 - tail

CPU Scheduling

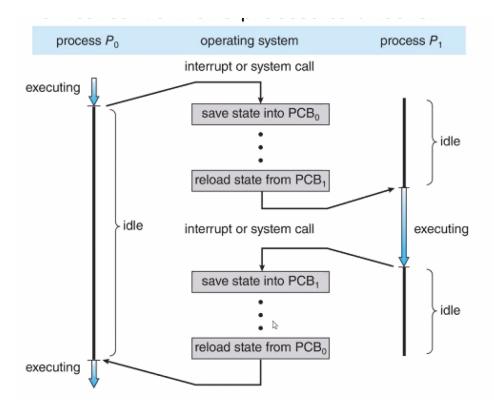
- process enter a system
 - o add to ready queue
 - state: ready
- process enter real CPU
 - execute
- process need to visit devices when executing
 - add to device queues
 - state: waiting

Scheduling Program

- process scheduling is done by scheduling program
- generally, in a system need to execute a lot of processes
 - processes can't be allocated all to real memory
 - stored in disk → virtual memory
- 长期调度程序:会把在磁盘上的进程加载到内存,把最近不会被执行的进程从内存中挪出。
 - 。 I/O密集型进程
 - 。 CPU密集型进程。
- 短期调度程序(CPU调度程序):从就绪队列中选择进程分配CPU。

Process Switch (上下文切换)

- save old state
- load new state



Process Management

Process Set-up

- fork()
 - parent process generate children processes, forming a tree of processes
 - resources in child process
 - from OS (no share)
 - from parent (share)
 - execution
 - concurrent
 - parent waits until children terminate
 - return pid
- pid = 0 child

 $\bullet \ \ \text{pid} > 0 \ \ \text{parent}$

Process Execution

- exec()
 - load new .exe file into child PCB to execute
 - parent enters waiting queue

Process Terminates

- exit()
 - child calls it and ask OS to delete itself.
 - return parent
 - if parent is waiting, all resources will be released by OS
 - wait() contains process table
- others
 - only can be terminated by parent
 - child resources overflow
 - no need for child's tasks
 - parent is terminating
- 僵尸进程
 - child terminates
 - parent not waiting
- 孤儿进程
 - child terminates
 - parent terminates without waiting

Status

User Status

Kernel Status

Files, IPC (Interprocess Cooperation)

- Aim
 - information sharing
 - computation speedup
 - modularity
 - convenience
- 2 modules
 - shared memory
 - message passing

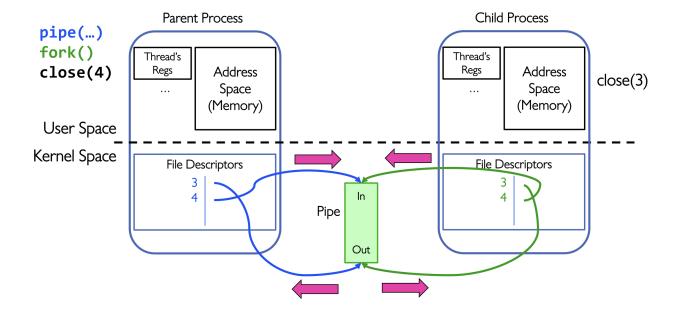
Shared Memory

process

Message Passing

• common in distributed system

Pipe



Socket

- consists:
 - o IP
 - port
 - each process needs a port
- RPC (Remote-Process-Communication)