

1.

We assume that t is the time we use before the enhancement, and t_1 is the time the part which will be enhanced uses.

So we can know the rest part time is $t - t_1$.

The enhancement improves some mode of execution by a factor of 3. Thus we can know that the part could use $\frac{t}{3}$ as the processing time.

And then we can get a equation: $\frac{t_1}{3} = t - t_1$

So we can get the $t_1 = \frac{3}{4}t$, and we can solve the a&b questions.

(a) Before time: t

After time: $\frac{t}{2}$ We have the speedup $2\times$.

(b) 75%.

2.

(a) There are no penalty to any other instructions, so we can use Amdahl's law directly.

$$\text{So we get } speedup_{overall} = \frac{1}{0.7 + \frac{0.3}{3}} = 1.25\times$$

(b) We have data cache processing slowed down, we use the Amdahl's law in the special forms.

$$\text{So we get } speedup_{overall} = \frac{1}{\frac{0.1}{0.5} + 0.6 + \frac{0.3}{3}} \approx 1.11\times$$

(c) The percentage of execution time spent on floating-point operations is

$$\frac{0.1}{0.9} \approx 11\%$$

The percentage of execution time spent on data cache accessed is $\frac{0.2}{0.9} \approx 22\%$

3.

$$(a) \text{ speedup} = \frac{1}{0.5 + \frac{0.5}{2}} \approx 1.33 \times$$

$$(b) \text{ speedup} = \frac{1}{0.2 + \frac{0.8}{2}} \approx 1.67 \times$$

$$(c) \text{ speedup} = \frac{1}{0.6 \times \frac{0.5}{2} + 0.6 \times 0.5 + 0.4} \approx 1.18 \times$$

$$(d) \text{ speedup} = \frac{1}{0.6 + 0.4 \times 0.2 + 0.4 \times \frac{0.8}{2}} \approx 1.19 \times$$

4.

$$(a) \text{ speedup} = \frac{1}{0.2 + \frac{0.8}{N}} = \frac{N}{0.2N + 0.8}$$

$$(b) \text{ speedup} = \frac{1}{0.2 + 0.08 + \frac{0.8}{8}} \approx 2.63 \times$$

$$(c) \text{ speedup} = \frac{1}{0.2 + \frac{0.8}{8} + 0.03} \approx 3.03$$

$$(d) \text{ speedup} = \frac{1}{0.2 + \frac{0.8}{N} + 0.01 \times \log_2 N}$$

$$(e) \text{ speedup} = \frac{1}{1 - P\% + \frac{P\%}{N} + 0.01 \times \log_2 N}$$

And we can find that when $N = P \ln 2$, We can have the highest speedup.

