

Lab 1: Getting Started

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1 Objective

- Prepare cross development environment for PXA270.
- Be familiar with your PXA270.

2 Prerequisite

- Finish pre-Lab and get familiar with basic Linux and gcc compiler commands.

3 Setting Up

Before we go further, please make sure that you have the Internet connected and connect your PC with development board (PXA270) correctly. Each site has four IPs. The first three IPs are assigned to your Windows (Host OS), Linux (Guest OS on Virtual Machine), and Target (PXA270), accordingly. The fourth IP is reserved. Please configure your systems as following:

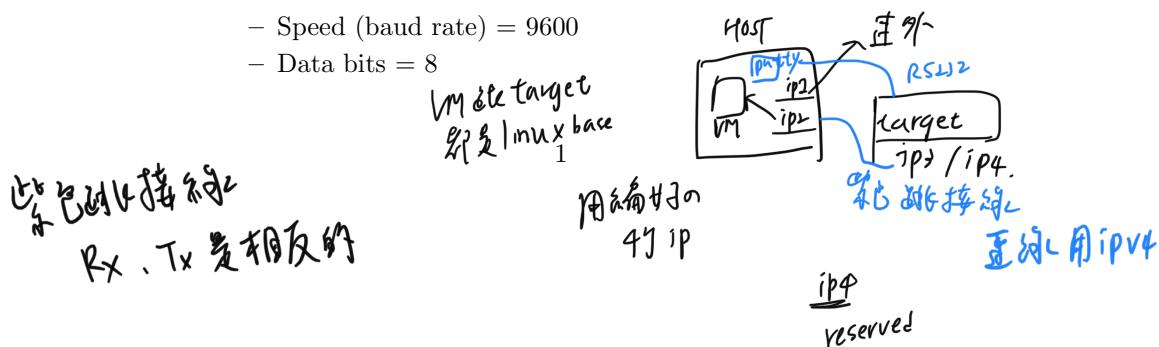
使用者自訂

3.1 Windows Configuration (ip1)

1. Configure IP as mentioned. (window IPv4)

2. Configure serial interface:

- Connect your Windows and PXA270 using RS232.
(COM/PC ↔ RS232 ↔ UART/PXA270)
- Download “Putty” from Internet and launch “Putty”.
- Choose “Connection type” as “Serial”.
- Select the “Connection→Serial” in “Category” and modify the serial control settings as following:
 - Speed (baud rate) = 9600
 - Data bits = 8



- 所有 Linux
 網路介面
 都是虛擬的
 (沒有物理的)
- Stop bit = 1
 - Parity = None
 - Flow Control = None
 - Open the connection.

3.2 Linux Configuration for VM and Target (<ip2> for VM and <ip3> for PXA)

1. Configure Network:

- Configure network temporally

pc host ifconfig: show/configure a network interface.

(左邊機器)
 (裡面)
 (terminal)

```
SHELL> sudo ifconfig eth0 <ip2> broadcast  
192.168.0.255 netmask 255.255.255.0
```

設定網路

- route: show/manipulate the IP routing table.

```
SHELL> sudo route add default gw 192.168.0.1
```

616 lab 設定 gateway

- Configure network permanently

- Edit /etc/network/interfaces

```
auto eth0  
iface eth0 inet static  
  address <ip2>  
  netmask 255.255.255.0  
  gateway 192.168.0.1
```

(這就可以出得去)
網址

立開lab 的 TCP 連接埠
設成 192.168.0.1

在這 lab 的網路是

192.168.0.X

叫這子區域

- Edit /etc/resolvconf/resolv.conf.d/base, and add the following DNS configurations to the file.

```
nameserver 140.113.6.2  
nameserver 140.113.1.1
```

- Use resolvconf to generate /etc/resolv.conf

```
SHELL> sudo resolvconf -u
```

- Restart network

重新啟動網卡

```
SHELL> sudo /etc/init.d/networking restart
```

4 Development Environment

4.1 Setting Up

1. Install package for building kernel and root filesystem:

```
SHELL> sudo apt-get install make patch libncurses32-dev
```

2. Download the following packages to your Guest OS (Linux Host), and put them under your home directory:

- ToolChain for PXA270: `arm-linux-toolchain-bin-4.0.2.tar.gz` (for building everything).
- Linux Kernel Source: `mt-linux-2.6.15.3.tar.gz` (for building kernel);
- Linux Kernel Patch for PXA270: `linux-2.6.15.3-creator-pxa270.patch` (for building kernel);
- Root Filesystem for PXA270: `rootfs.tar.gz` (for building rootfs)

Host
system

4.2 Install Toolchain

開發環境

A toolchain is the set of programming tools that are used to create a product. In the toolchain for PXA270 (`arm-linux-toolchain-bin-4.0.2.tar.gz`, a cross compiler), capable of creating executable code for a platform other than the one on which the compiler is run, is included. To compile a program running on your PXA270, you need to install the toolchain for compilation:

1. Unzip the tarball:

```
SHELL> sudo tar zxvf arm-linux-toolchain-bin-4.0.2.tar.gz  
-C /
```

2. Add the PATH to ToolChain at the end of `~/.bashrc`:

```
export PATH=$PATH:/opt/arm-unknown-linux-gnu/bin
```

3. Reload `~/.bashrc` to make the above changes effective.

```
SHELL> source ~/.bashrc
```

DEMO ⚡

4. Execute Program on PXA270

- Move or copy the executable file to <export point>:

```
SHELL> mv hello <absolute path to export point>
```

Or,

```
SHELL> cp hello <absolute path to export point>
```

- Execute program under NFS mount point on PXA270:

板子上

```
> cd /mnt  
> ./hello
```

5 Configure nfsd (Network Filesystem Daemon) on VM

- Install nfs server:

```
SHELL> sudo apt-get install nfs-kernel-server
```

- Create the export point in your home directory for nfs server:

```
SHELL> cd ~  
SHELL> mkdir <export point>
```

类似網路共享

- Configuring the nfs server on Linux.

Get the absolute path to <export point> by using 'pwd' command. Edit the `/etc/exports` file, it should contain the following item:

```
<absolute path to export point> <  
ip3 >(rw,no_root_squash,no_subtree_check)
```

- Starting the server on Linux:

重新啟動 nfs server

```
SHELL> sudo /etc/init.d/nfs-kernel-server restart
```

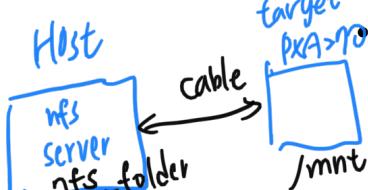
Configure nfs on PXA

- Mount to the nfs server (nfsd) on Linux:

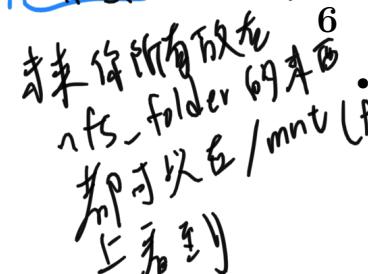
```
> portmap  
> mount <ip2>:<absolute path to export point> /mnt
```

將想要掛入的資料夾來掛上進來 ↑
VM ↑
PXA270 on /mnt ↑

export point:
nfs-folder



拿來供給 PXA 用
nfs_folder 69 有
都可以在 /mnt (PXA270 上面)
上看到



假裝 uboot 時會用 nfs 來傳 image

在 uboot " 傳達你 的 image file

在 uboot 就可執行的 網路傳輸 + 開定

Configure tftpd (Trivial File Transfer Protocol server) on VM

The tftpd is normally started by the xinetd daemon. Before that, please install necessary packages and create a root directory <tftproot> for tftpd:

```
SHELL> sudo apt-get install xinetd tftp tftpd  
SHELL> cd ~  
SHELL> mkdir <tftproot>  
SHELL> cd <tftproot>  
SHELL> pwd
```

In the above commands, apt-get is a command of APT (Advanced Packaging Tool). It helps automate the retrieval, configuration and installation of software packages. Get the absolute path, <tftproot>, of the newly created directory by using 'pwd' command. Then, edit /etc/xinetd.d/tftp to config the tftpd settings:

tftp 配置檔

```
service tftp  
{  
socket_type = dgram  
protocol = udp  
wait = yes  
user = lab616 michelle  
server = /usr/sbin/in.tftpd  
server_args = -s <tftprootpath>  
disable = no    設成 no 限制  
}
```

Then, reload xinetd:

```
SHELL> sudo /etc/init.d/xinetd reload
```

xinetd 會幫你把 tftp server 安裝好

8 Build uImage

8.1 Build Kernel

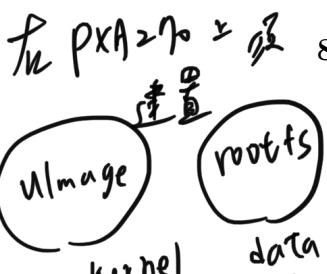
1. Unzip Kernel Source:

```
SHELL> cd ~  
SHELL> tar zxvf mt-linux-2.6.15.3.tar.gz
```

給板子用

The extraction will create a directory microtime/ in your home directory, and the following subdirectories:

- build-linux/ - tools for building kernel image and root filesystem.



Z: 是 zip 檔
X: 無解壓縮
V: 組合的压缩壓縮
f: 無解壓 file on
P: 進入

patch未補足の部分

PXA270 有自己の起動プログラム、不含む Linux 核

野生

虫卵

- `create-pxa270/` - the link to `pro/devkit/lsp/create-pxa270`.
- `linux/` - the link to `pro/devkit/lsp/create-pxa270/linux-2.6.15.3`.
- `package/` - the link to `pro/host/src/BUILD`.
- `pro/` - it contains the linux kernel source and some demo codes.

Then, we need to patch the kernel for PXA270.

修補程式人の役一=2

```
SHELL> cp linux-2.6.15.3-creator-pxa270.patch  
~/microtime/create-pxa270/  
SHELL> cd ~/microtime/create-pxa270/  
SHELL> patch -p0 < linux-2.6.15.3-creator-pxa270.patch
```

2. Build Kernel:

Create the configure file, ‘.config’, for PXA270:

```
SHELL> cd ~/microtime/linux  
SHELL> make mrproper 深度清潔  
SHELL> make creator_pxa270_defconfig
```

產生隱形的 config 檔

Build Kernel:

```
SHELL> make clean  
SHELL> make
```

It takes couple minutes to complete the compilation. Please be patient.

3. Build uImage for PXA270:

- Method 1 – Convert compressed kernel image (`zImage`) to U-Boot format (`uImage`):

After successfully compile the kernel, you will see a `zImage` files in `linux/arch/arm/boot/`. We need to convert the `zImage` to `uImage`. The `zImage` contains a compressed linux kernel and a self-decompression program at the head of the file. After system start up, the boot-loader moves the `zImage` to SDRAM, the `zImage` will execute the self-decompression program to extract the linux kernel, and start normal linux kernel boot process.

kernel 鏽彈元
含產生 zImage

```
SHELL> cd ~/microtime  
SHELL> ./build-linux/mkimage -A arm -O linux -T  
kernel -C none  
-a 0xa0008000 -e 0xa0008000 -n "2.6.15.3 kernel for  
Creator-PXA270"  
-d linux/arch/arm/boot/zImage uImage
```

- Method 2 – Convert uncompressed kernel image (`vmlinux`) to U-Boot format (`uImage`):

uImage
uboot 专用的
密档式

with bootm
command

After successfully compile the kernel, we can find a uncompressed linux kerlen vmlinux in linux/. We should compressed the vmlinux, and convert it to U-Boot readable format. After system start up, the bootloader will directly decompressed zImage from NAND FLASH to SDRAM, and start normal linux kernel boot process.

```
SHELL> cd ~/microtime  
SHELL> arm-unknown-linux-gnu-objcopy -O binary -R .note -R .comment -S linux/vmlinux linux.bin  
SHELL> gzip -9 -c linux.bin > zImage  
SHELL> ./build-linux/mkimage -A arm -O linux -T kernel -C gzip  
-a 0xa0008000 -e 0xa0008000 -n '2.6.15.3 kernel for Creator-PXA270'  
-d zImage uImage
```

Please check the size of uImage. It should be about 1.5MB. You need to repeat the above steps if you get a very small uImage.

8.2 Build Root Filesystem

The root filesystem is the filesystem that is contained on the same partition on which the `root` directory is located, and it is the filesystem on which all the other filesystems are mounted (i.e., logically attached to the system) as the system is booted up (i.e., started up). In this practice, we are going to build a root filesystem in JFFS2 format. JFFS2 (Journalling Flash File System version 2) is a log-structured filesystem for use with flash memory devices.

1. Unzip Root Filesystem Source

```
SHELL> cd ~  
SHELL> sudo tar xzvf rootfs.tar.gz  
SHELL> cd microtime/  
SHELL> sudo chown -R lab616:lab616 rootfs/
```

2. Put permanent items into your rootfs:

If you want to put file /usr/lab616/tt.bin into /usr) and you are now in the directory of microtime/rootfs/:

```
SHELL> cp /usr/lab616/tt.bin usr/
```

3. Build root filesystem

```
SHELL> cd ~/microtime  
SHELL> ./build-linux/mkfs.jffs2 -v -e 131072 --pad=0xa00000  
-r rootfs/  
-o rootfs.jffs2
```

make file system

可編輯不同格式的 fs

rootfs

9 Flash & Boot

In computing, ‘booting’ is a process that starts operating systems when the user turns on a computer system. A boot sequence is the initial set of operations that the computer performs when power is switched on. The boot loader typically loads the main operating system for the computer. In this practice, we are using ‘U-Boot’ bootloader to boot up PXA270.

9.1 Flash uImage *直到設至行前按住電鍵，就會進到 U-boot*

- Power on PXA270 and press any key before U-Boot autoboots the target.
- Configure U-Boot environment variables for tftp download:

```
UBOOT> setenv ipaddr <ip3>
UBOOT> setenv serverip <ip2> server in host
UBOOT> setenv bootargs root=/dev/mtdblock3 rw
rootfstype=jffs2 console=ttyS0,9600n8 mem=64M ip=<
<ip3>:<ip2>:<gw ip>:255.255.255.0::eth0:off
UBOOT> saveenv 192.168.1.1
```

The <gw ip> is the gateway IP. It can be skipped if PXA270 and Linux Host are located in the same subnet. You can use `printenv` to check the settings.

- Prepare uImage and rootfs.jffs2:
Build the uImage and rootfs.jffs2, then move them to the <tftpboot>.
- Download and flash uImage:

```
UBOOT> tftp a1100000 uImage PXA 故障 uImage 放在的範圍
UBOOT> protect off 100000 47ffff 把記憶體保護關掉
UBOOT> erase 100000 47ffff 把要下載的區域先做 unprotect
UBOOT> cp.b a1100000 100000 200000 用 binary copy 的方式把 uimage copy 進來
```

9.2 Flash rootfs

- Download and flash rootfs.jffs2: *若之後有修改, 重新燒, 重做這段*

```
UBOOT> tftp a1480000 rootfs.jffs2 只取 rootfs.jffs2 複製到 tftpboot
UBOOT> protect off 480000 137ffff
UBOOT> erase 480000 137ffff
UBOOT> cp.b a1480000 480000 a00000
```

9.3 Boot your PXA270

Now, you can either boot your Linux automatically or manually. For manual boot, please press ‘Enter’ before U-Boot starts the autoboot, then type:

boot from memory) 設定燒錄器的位址，從 memory 或地方先將它 boot 起來

```
UBOOT> bootm 100000
```

Otherwise, to set up a correct autoboot, you need to configure U-Boot environment variables:

```
UBOOT> setenv bootm 100000  
UBOOT> setenv bootcmd run linux  
UBOOT> saveenv
```

Then, press the 'Reset' button on PXA270 to enable the autoboot.

步驟：在板子上跑出 hello world

uImage : OS 程式

root file system : 資料 ? A check

Lab 2-1 Prepare Boot Image of Xilinx ZC702

Wen-Lin Sun

2018-9-21

In this lab, we will prepare the things that we need when booting up Linux on Xilinx ZC702. The files are first-stage boot loader(FSBL), second-stage boot loader(e.g. U-boot), device tree, kernel image and root filesystem. Make sure you have these files after you finish this lab:

- boot.bin (fsbl.elf + u-boot.elf)
- uImage
- uramdisk.image.gz
- devicetree.dtb

And put them into a directory (create it yourself).

1 Environment Setup

Before building all of these files, we should install several tools.

- Install 32-bit libraries:
Update the database before we install anything.

```
SHELL> sudo apt-get update
SHELL> sudo apt-get install lib32z1 lib32ncurses5 lib32ncurses5-dev lib32stdc++6
libbz2-1.0:i386
```

- Download & Install Xilinx Vivado Design Suite (you may need to register):
Recommend: Install it under your home directory.
Official website: <https://www.xilinx.com/support/download.html>
- Setup Environment Variables:

```
SHELL> export CROSS_COMPILE=arm-linux-gnueabihf-
SHELL> source [xilinx_directory]/Xilinx/SDK/2018.2/settings64.sh
      source /usr/xilinx/SDK/2018.3/settings64.sh
```

- Install Git:

```
SHELL> sudo apt-get install git
```

- Build & Install device tree compiler:

```

SHELL> git clone https://git.kernel.org/pub/scm/utils/dtc/dtc.git
SHELL> cd dtc
SHELL> sudo apt-get install flex bison swig python-dev, sudo apt -y install libyaml-dev pkg-config python3-dev
SHELL> make
SHELL> export PATH='pwd':$PATH

```

2 Generate FSBL

You should download the hardware definition file system.hdf from E3 first.

- Create a link gmake to make:

Formally, we should install gmake, but we can also use make instead here.

```

SHELL> sudo ln -s /usr/bin/make /usr/bin/gmake
SHELL> cd <hdf_dir>
SHELL> hsi
hsim% set hwdesign [open_hw_design system.hdf]
hsim% generate_app -hw $hwdesign -os standalone -proc ps7_cortexa9_0 -app zynq_fsbl
-compile -sw fsbl -dir <fsbl_dir>
hsim% exit

```

The file fsbl will be placed at <fsbl_dir> as you set in the command. Just rename it to fsbl.elf.

3 Build U-boot

In this lab, we use u-boot as our second-stage boot loader, for helping us boot from Linux. Since the tool mkimage we use when building kernel image will also be generated when building u-boot, we should make sure this step is done correctly before we continue.

- Build u-boot:

```

SHELL> git clone https://github.com/Xilinx/u-boot-xlnx.git
SHELL> cd u-boot
SHELL> make zynq_zc702_defconfig
SHELL> make

```

The file u-boot will be generated in this directory and the mkimage will be placed under the directory tools. Just rename it to u-boot.elf.

- Export mkimage tools to PATH

```

SHELL> cd tools
SHELL> export PATH='pwd':$PATH

```

4 Build uImage & Device Tree

In this section, we will build kernel image and attach it with u-boot recognizable header. Alternatively, we will build the the corresponding device tree.

- Download & Configure Linux kernel:
In this lab, we use the version 2016.4.

```
SHELL> git clone https://github.com/Xilinx/linux-xlnx.git
SHELL> git checkout tags/xilinx-v2016.4
SHELL> make ARCH=arm xilinx_zynq_defconfig
SHELL> make ARCH=arm menuconfig
```

- Build kernel image with u-boot header (uImage):

```
SHELL> make ARCH=arm UIMAGE_LOADADDR=0x8000 uImage
```

The uImage will be placed under arch/arm/boot.

- Build device tree:

```
SHELL> make ARCH=arm zynq-zc702.dtb
```

The device tree zynq-zc702.dtb will be placed under arch/arm/boot/dts. Just rename it to devicetree.dtb.

5 Modify Root Filesystem

In this lab, we will use BusyBox as our root filesystem.

- Download & Configure BusyBox:
We use the branch 1_29_stable for example.

```
SHELL> git clone git://git.busybox.net/busybox -b 1_29_stable
SHELL> make ARCH=arm menuconfig
```

Modify the configuration as you need. Make sure [Setting] -> [Build Options] -> [Build static binary (no shared libs)] is selected. And then, build it.

- Build & Generate Busybox and related files:

```
SHELL> make ARCH=arm
```

After that, use make install to gather the compiled Busybox and symlink into the directory _install.
_install.

```
SHELL> make ARCH=arm install
```

- Modify the filesystem:

- Manually add missing directories of root filesystem to _install. (/dev, /sys, /proc, /root and /etc/init.d)

```
SHELL> cd _install
SHELL> mkdir -p dev sys proc root etc/init.d
```

- Manually create initial script rcS under etc/init.d and make it executable.

```

1 #!/bin/sh
2
3 mount -t proc none /proc → 手動挂載 /proc 系統目錄
4 mount -t sysfs none /sys → 手動挂載 /sys 系統目錄
5 /sbin/mdev -s

```

- Manually create inittab under etc.

```

1 #!/bin/sh
2
3 # Init script
4 ::sysinit:/etc/init.d/rcS
5
6 # Start shell on the serial ports
7 ::respawn:/sbin/getty -L ttysP0 115200 vt100
8
9 # Execute /sbin/init when restarting the init process
10 ::restart:/sbin/init
11
12 # Umount everything before rebooting
13 ::shutdown:/bin/umount -a -r

```

- Set etc/passwd to make root be able to login without password.

```
1 root::0:0:root:/root:/bin/sh
```

- Create soft link to /init to avoid that kernel cannot find the init.

```
SHELL> ln -s /sbin/init init
```

- After changing all the directory to be owned by root, you can pack the root filesystem with cpio format. And then, compress it with gzip format. The output file will be named ramdisk.cpio.gz and placed under _install.

```

SHELL> cd ..
SHELL> sudo chown -R root:root _install
SHELL> cd _install
SHELL> find . | cpio -H newc -o | gzip -9 > ../ramdisk.cpio.gz

```

Transform the ramdisk.cpio.gz into u-boot format with mkimage.

```
SHELL> mkimage -A arm -T ramdisk -C gzip -d ramdisk.cpio.gz uramdisk.image.gz
```

Now you have the compressed root filesystem file uramdisk.image.gz under busybox.

6 Prepare Boot Image

In this section, we will generate the boot image boot.bin. Before doing this step, make sure you have fsbl.elf and u-boot.elf.

- Put them into a directory.
- Enter the directory and create a file `boot.bif` with following content (sample):

```

1 image : {
2     [bootloader]fsbl.elf
3     u-boot.elf
4 }
```

- Generate the boot image.
`bootgen` merges `fsbl.elf` and `u-boot.elf` to boot image `boot.bin`.

```
SHELL> bootgen -image boot.bif -o i boot.bin
```

PXA270 這 device tree
 ↓
 無技術
 BB 教書
 含此所有東西 LO 等寫進 drive 這裡
 ∵ kernel 就很肥

ZC702 把固件寫進一個配置文件 → device tree
 (從 driver 拉出來)

Lab 2-2 Prepare Boot Medium of Xilinx ZC702

Wen-Lin Sun

2018-9-21

In this lab, we will practice using two kinds of boot media: SD card and QSPI flash to boot up ZC702 with the image that we prepared in previous part.

1 SD Boot

Use SD card as boot medium is the most convenient way to test your boot image.

- Format the sd card with FAT format. ~~FAT32~~
- Put all the files that you have generated into the card.
 - boot.bin
 - uImage
 - uramdisk.image.gz
 - devicetree.dtb
- Check SW16 setting. (Figure 2)

2 QSPI Boot

Because there is nothing in QSPI flash now, we cannot boot from it directly now. There are several ways to write something to the QSPI flash, such as JTAG(with Xilinx SDK), Linux driver, u-boot. In this lab, we provide a way which requires the well prepared SD card in last section and write things to flash through u-boot, which is similar to what we have done in previous lab.

- Use SD boot to boot into u-boot.
After booting into u-boot, you will see the screen shows Zync> as following.

Zync>

- Load the files from SD card to Ram.

```
Zync> load mmc 0 0x6000000 ${boot_image}
Zync> load mmc 0 ${kernel_load_address} ${kernel_image}
Zync> load mmc 0 ${devicetree_load_address} ${devicetree_image}
Zync> load mmc 0 ${ramdisk_load_address} ${ramdisk_image}
```

- Init flash device on given SPI bus and chip select.

```
Zync> sf probe 0 20000000 0
```

- Write boot.bin from ram to flash.

```
Zync> sf protect unlock 0 ${boot_size}
Zync> sf erase 0 ${boot_size}
Zync> sf write 0x6000000 0 ${boot_size}
```

- Write uImage from ram to flash.

```
Zync> sf protect unlock 0x100000 ${kernel_size}
Zync> sf erase 0x100000 ${kernel_size}
Zync> sf write ${kernel_load_address} 0x100000 ${kernel_size}
```

- Write devicetree.dtb from ram to flash.

```
Zync> sf protect unlock 0x600000 ${devicetree_size}
Zync> sf erase 0x600000 ${devicetree_size}
Zync> sf write ${devicetree_load_address} 0x600000 ${devicetree_size}
```

- Write uramdisk.image.gz from ram to flash.

```
Zync> sf protect unlock 0x620000 ${ramdisk_size}
Zync> sf erase 0x620000 ${ramdisk_size}
Zync> sf write ${ramdisk_load_address} 0x620000 ${ramdisk_size}
```

- Reboot it with the QSPI boot SW16 setting (Figure 2).

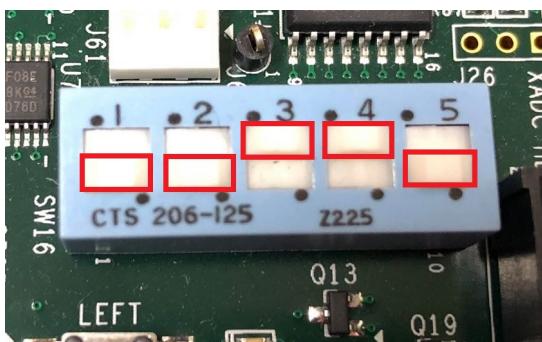


Figure 1: SD boot mode switch setting

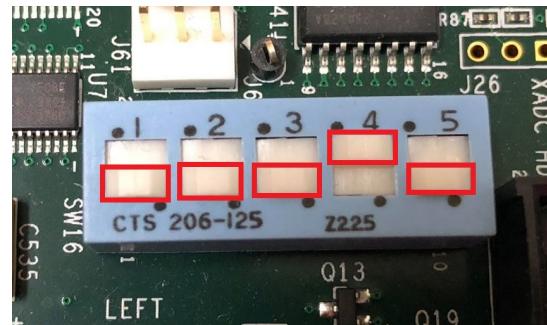


Figure 2: QSPI boot mode switch setting

Lab 3: Task

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1 Objective

- Be familiar with system calls: `fork()`, `wait()`, `waitpid()`, `nice()`, `exec()`, etc.
- Be familiar with POSIX programming: `pthread_create()`, `pthread_exit()`, etc.

2 Prerequisite

- Read man pages of the above system calls.

3 Process Control

- `fork()` creates a child process that differs from the parent process only in its PID and PPID, and in fact that resource utilizations are set to 0. The memory of child process is copied from the parent and a new process structure is assigned by the kernel. The environment, resource limits, umask, controlling terminal, current working directory, root directory, signal masks and other process resources are also duplicated from the parent in the forked child process, while file locks and pending signals are not inherited. The return value of the `fork()` function discriminates the two processes of execution. A zero is returned by the fork function in the child's process, while the parent process gets the PID (a non-zero integer) of its child.

```
1  /*
2   * process.c
3   */
4
5 #include <errno.h>      /* Errors */
6 #include <stdio.h>       /* Input/Output */
7 #include <stdlib.h>      /* General Utilities */
8 #include <sys/types.h>    /* Primitive System Data Types */
9 #include <sys/wait.h>     /* Wait for Process Termination */
10 #include <unistd.h>      /* Symbolic Constants */
11
12 pid_t childpid; /* variable to store the child's pid */
13
14 void childfunc(void)
15 {
16     int retval; /* user-provided return code */
17
18     printf("CHILD: I am the child process!\n");
19     printf("CHILD: My PID: %d\n", getpid());
```

```

20 printf("CHILD: My_parent's_PID_is : %d\n", getppid());
21 printf("CHILD: Sleeping_for_1_second ... \n");
22 sleep(1); /* sleep for 1 second */
23
24 printf("CHILD: Enter_an_exit_value_(0_to_255): ");
25 scanf("%d", &retval);
26 printf("CHILD: Goodbye!\n");
27
28 exit(retval); /* child exits with user-provided return code */
29 }
30
31 void parentfunc(void)
32 {
33     int status;      /* child's exit status */
34
35     printf("PARENT: I_am_the_parent_process!\n");
36     printf("PARENT: My_PID: %d\n", getpid());
37     printf("PARENT: My_child's_PID_is %d\n", childpid);
38     printf("PARENT: I_will_now_wait_for_my_child_to_exit.\n");
39
40     /* wait for child to exit, and store its status */
41     wait(&status);
42     printf("PARENT: Child's_exit_code_is : %d\n", WEXITSTATUS(status));
43     printf("PARENT: Goodbye!\n");
44
45     exit(0); /* parent exits */
46 }
47
48 int main(int argc, char *argv[])
49 {
50     /* now create new process */
51     childpid = fork();
52
53     if (childpid >= 0) { /* fork succeeded */
54         if (childpid == 0) { /* fork() returns 0 to the child process */
55             childfunc();
56         } else { /* fork() returns new pid to the parent process */
57             parentfunc();
58         }
59     } else { /* fork returns -1 on failure */
60         perror("fork"); /* display error message */
61         exit(0);
62     }
63
64     return 0;
65 }
```

You can check the identifiers of the parent and child process by executing `ps` command in shell. Please refer to its man page for more details.

小孩做完的結果存回 parent

- `wait()` suspends execution of the current process until one of its children terminates.
- `waitpid()` suspends execution of the current process until a child specified by pid argument has changed state. By default, `waitpid()` waits only for terminated children, but this behaviour is modifiable via the options argument. Please refer to its man page.

等
假定某
一隻 child
process

```
1  /*
2   * waitpid.c -- shows how to get child's exit status
3   */
4
5 #include <errno.h>          /* Errors */
6 #include <stdio.h>           /* Input/Output */
7 #include <stdlib.h>          /* General Utilities */
8 #include <sys/types.h>        /* Primitive System Data Types */
9 #include <sys/wait.h>         /* Wait for Process Termination */
10 #include <time.h>            /* Time functions */
11 #include <unistd.h>          /* Symbolic Constants */
12
13 pid_t childpid; /* variable to store the child's pid */
14
15 void childfunc(void)
16 {
17     int randtime;      /* random sleep time */
18     int exitstatus;    /* random exit status */
19
20     printf("CHILD: I am the child process!\n");
21     printf("CHILD: My PID: %d\n", getpid());
22
23     /* sleep */
24     srand(time(NULL));
25     randtime = rand() % 5;
26     printf("CHILD: Sleeping for %d second...\n", randtime);
27     sleep(randtime);
28
29     /* rand exit status */
30     exitstatus = rand() % 2;
31     printf("CHILD: Exit status is %d\n", exitstatus);
32
33     printf("CHILD: Goodbye!\n");
34     exit(exitstatus); /* child exits with user-provided return code */
35 }
36
37 void parentfunc(void)
38 {
39     int status;          /* child's exit status */
40     pid_t pid;
41
42     printf("PARENT: I am the parent process!\n");
43     printf("PARENT: My PID: %d\n", getpid());
44
45     printf("PARENT: I will now wait for my child to exit.\n");
```

```

47     /* wait for child to exit, and store its status */
48 do
49 {
50     pid = waitpid(childpid, &status, WNOHANG);
51     printf("PARENT: Waiting child exit ...\n");
52     sleep(1);
53 }while (pid != childpid);
54
55 if (WIFEXITED(status)) {
56     // child process exited normally.
57     printf("PARENT: Child's exit code is: %d\n",
58           WEXITSTATUS(status));
59 }else{
60     // Child process exited thus exec failed.
61     // LOG failure of exec in child process.
62     printf("PARENT: Child process executed but exited failed.\n");
63 }
64
65 printf("PARENT: Goodbye!\n");
66
67 exit(0); /* parent exits */
68 }
69
70 int main(int argc, char *argv[])
71 {
72     /* now create new process */
73     childpid = fork(); fork之後子程式與父程式做相同東西
(code相同)
74
75     if (childpid >= 0) { /* fork succeeded */
76         if (childpid == 0) { /* fork() returns 0 to the child process */
77             childfunc();
78         } else { /* fork() returns new pid to the parent process */
79             parentfunc();
80         }
81     } else { /* fork returns -1 on failure */
82         perror("fork"); /* display error message */
83         exit(0);
84     }
85
86     return 0;
87 }
88 }
```

- `nice()` adds *inc* to the nice value for the calling process. (A higher nice value means a low priority.) Only the superuser may specify a negative increment, or priority increase. The range for nice values is described in `getpriority(2)`.

```

1 #include <unistd.h>
2 ...
3 int incr = -20;
4 int ret;
5
6 ret = nice(incr); 每次 -20, 提升优先级

```

- The `exec()` family of functions initiates a new process image within a program. The initial argument for these functions is the pathname of a file which is to be executed.

```

1 /*
2  * exec.c
3  */
4
5 #include <unistd.h>
6 int main(int argc, char *argv[])
7 {
8     execl("/bin/ls", "/bin/ls", "-r", "-t", "-l", (char *) 0);
9
10    return 0;
11 }

```

使用 fork 分享 process

4 Thread

POSIX library

thread-create() in windows?

A standardized programming interface was required to take full advantages provided by threads. For UNIX systems, this interface has been specified by the IEEE POSIX 1003.1c standard (1995). Implementations adhering to this standard are referred to as **POSIX threads**, or Pthreads.

- `pthread_create()` creates a new thread, with attributes specified by *attr*, within a process. Upon successful completion, `pthread_create()` will store the ID of the created thread in the location specified by *thread*. For more information, please see its man page.
- `pthread_join()` suspends execution of the calling thread until the target thread terminates.
- `pthread_detach()` tells the underlying system that resources allocated to a particular thread can be reclaimed once it terminates. This function should be used when an exit status is not required by other threads.
- `pthread_exit()` terminates the calling thread.

The above APIs are included in the header file, ‘`pthread.h`’.

```

1 /*
2  * pthread.c -- shows how to create a thread
3  */
4
5 #include <pthread.h>
6 #include <stdio.h>
7 #include <stdlib.h>

```

不使用情境，直不適用 thread

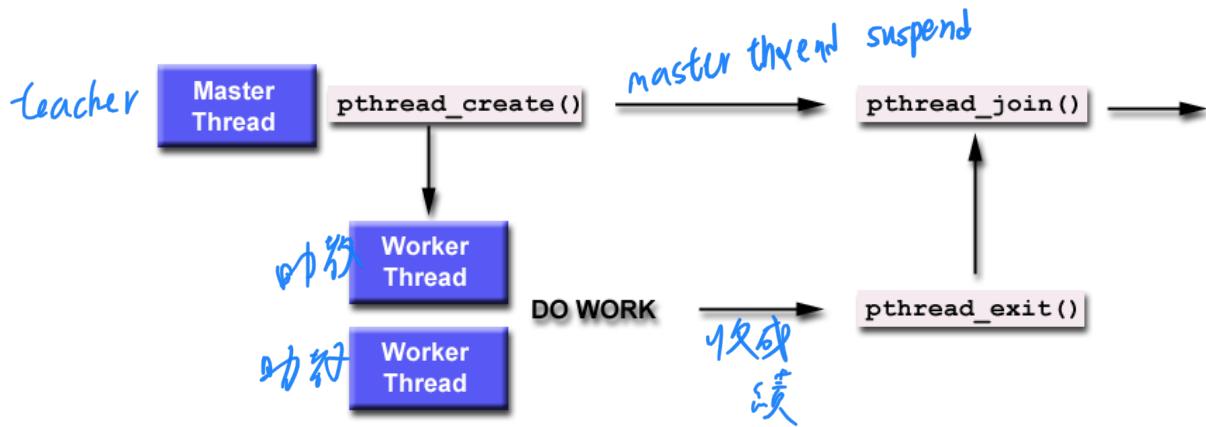


Figure 1: pthread_join()

```

8
9 #define NUM_THREADS 5
10
11 void *show(void *threadid) {
12 {
13     long tid;
14
15     tid = (long)threadid;
16     printf("Hello! I am thread %ld!\n", tid);
17
18     pthread_exit(NULL);
19 }
20
21 int main(int argc, char *argv[])
22 {
23     pthread_t threads[5];
24     int rc;
25     int t;
26
27     for(t=0;t<NUM_THREADS;t++){
28         printf("In main(): creating thread %d\n", t);
29         rc = pthread_create(&threads[t], NULL, show, (void *)t);
30         if(rc){
31             printf("ERROR: return code from pthread_create() is %d\n", rc);
32             exit(-1);
33         }
34     }
35
36     printf("Main: program completed..Exiting.\n");
37     pthread_exit(NULL);
38 }

```

Diagram annotations (handwritten in blue and red):

- Line 11: "show" is annotated with "執行 function" (execute function).
- Line 12: "}" is annotated with "回傳" (return).
- Line 15: "tid = (long)threadid;" is annotated with "Tid/pid 與 thread id function pointer" (Tid/pid and thread id function pointer).
- Line 29: "rc = pthread_create(&threads[t], NULL, show, (void *)t);" is annotated with "每執行 show 會傳 (function)" (Each execution of show will pass (function)).
- Line 30: "if (rc){" is annotated with "如果 create 失敗" (If create fails).
- Line 31: "printf("ERROR: return code from pthread_create() is %d\n", rc);" is annotated with "但又錯" (But it's wrong).
- Line 32: "exit(-1);" is annotated with "錯 -1!!" (Wrong -1!!).
- Line 33: "}" is annotated with "屬於完成 故回來的" (Belongs to completion, so it comes back).
- Line 36: "printf("Main: program completed..Exiting.\n");" is annotated with "屬主" (Belongs to main).

```

1 /*
2  * join.c -- shows how to "wait" for thread completions
3  */
4

```

```

5 #include <math.h>
6 #include <pthread.h>
7 #include <stdio.h>
8 #include <stdlib.h>
9
10 #define NUM_THREADS 4
11
12 void *BusyWork(void *t)
13 {
14     int i;
15     long tid;
16     double result = 0.0;
17
18     tid = (long)t;
19     printf("Thread %ld starting...\n", tid);
20
21     for (i = 0; i < 1000000; i++)
22     {
23         result += sin(i) * tan(i);
24     }
25
26     printf("Thread %ld done. Result = %e\n", tid, result);
27     pthread_exit((void*) t);
28 }
29
30 int main (int argc, char *argv[])
31 {
32     pthread_t thread[NUM_THREADS];
33     pthread_attr_t attr;
34     int rc;
35     long t;
36     void *status;
37
38     /* Initialize and set thread detached attribute */
39     pthread_attr_init(&attr);
40     pthread_attr_setdetachstate(&attr, PTHREAD_CREATE_JOINABLE);
41
42     for (t = 0; t < NUM_THREADS; t++) {
43         printf("Main: creating thread %d\n", t);
44         rc = pthread_create(&thread[t], &attr, BusyWork, (void *)t);
45         if (rc) {
46             printf("ERROR: return code from pthread_create() is %d\n", rc);
47             exit(-1);
48         }
49     }
50
51     /* Free attribute and wait for the other threads */
52     pthread_attr_destroy(&attr);
53     for (t = 0; t < NUM_THREADS; t++) {
54         rc = pthread_join(thread[t], &status);
55         if (rc) {

```

```

56     printf("ERROR; return code from pthread_join() is %d\n", rc);
57     exit(-1);
58 }
59     printf("Main: join with thread %ld (status: %ld)\n", t, (long)status);
60 }
61
62 printf("Main: program completed. Exiting.\n");
63 pthread_exit(NULL);
64 }
```

```

1 /*
2 * deteach.c -- shows how to detach a thread
3 */
4
5 #include <errno.h>
6 #include <pthread.h>
7 #include <stdio.h>
8 #include <stdlib.h>
9 #include <unistd.h>
10
11 void *threadfunc(void *parm)
12 {
13     printf("Inside secondary thread\n");
14     sleep(3);
15     printf("Exit secondary thread\n");
16     pthread_exit(NULL);
17 }
18
19 int main(int argc, char **argv)
20 {
21     pthread_t thread;
22     int rc = 0;
23
24     printf("Create a thread using attributes that allow detach\n");
25     rc = pthread_create(&thread, NULL, threadfunc, NULL);
26     if (rc) {
27         printf("ERROR; return code from pthread_create() is %d\n", rc);
28         exit(-1);
29     }
30
31     printf("Detach the thread after it terminates\n");
32     rc = pthread_detach(thread);
33     if (rc) {
34         printf("ERROR; return code from pthread_detach() is %d\n", rc);
35         exit(-1);
36     }
37
38     printf("Detach the thread again\n");
39     rc = pthread_detach(thread);
40     /* EINVAL: No thread could be found corresponding to that
41      * specified by the given thread ID.
```

```

42     */
43     if (rc != EINVAL) {
44         printf("Got an unexpected result ! rc=%d\n", rc);
45         exit(1);
46     }
47     printf("Second detach fails as expected.\n");
48
49     /* sleep() is not a very robust way to wait for the thread */
50     sleep(6);
51     printf("Main() completed.\n");
52     return 0;
53 }
```

The link to `libpthread.a` library should be specified to the gcc compiler when compiling a program with `pthread` calls.

In host machine, you can natively compile your code with `gcc` by:

```
SHELL> gcc -o pthread pthread.c -lpthread
```

If you want to cross compile the code for your target, please compile it with:

```
SHELL> arm UNKNOWN LINUX GNU GCC -O PTHREAD PTHREAD.C LIBRARY  

-L /OPT/ARM UNKNOWN LINUX GNU/ARM UNKNOWN LINUX GNU/LIB/  

-I /OPT/ARM UNKNOWN LINUX GNU/ARM UNKNOWN LINUX GNU/INCLUDE/  

-LPTHREAD LIBRARY
```

In case you need to compile similar codes on your host, you need to use

```
HOST SH> gcc -o pthread pthread.c -lpthread
```

Lab 4: Inter-Process Communication (Part I)

Yu-Lun Huang

2019-07-16

1 Objective

- Be familiar with inter-process communication: semaphore, mutex, etc.

2 Prerequisite

- Read man pages of `semop()`, `semctl()`, `semget()`, `pthread_mutex_init()`, `pthread_mutex_lock()`, `pthread_mutex_unlock()`, etc.

3 Semaphore

Unix allocates arrays of semaphores, rather than creating them one at a time. When you create such an array, you must supply the following information:

- An integer called the "key" which acts as the semaphore array's "name" on the system
- The number of semaphores in the array
- Ownership of semaphore array (permission bits mode)

The system call `semget(2)` is used by a program to ask the operating system if it knows of a semaphore. The program passes the semaphore's key. If the semaphore exists, then the operating system returns an integer which is used by the program as the semaphore's identifier for the duration of that program. (NOTE: The semaphore key is permanent for as long as the semaphore exists. The semaphore ID returned by the operating system is just a temporary "handle" for accessing the semaphore and may be different each time.) `semget(2)` can also be used to create semaphores that don't exist by passing additional options. `semctl(2)` is used to set the value of a semaphore, remove it from the system, etc. Think of it as the way to "manage" the semaphore array. At its simplest, `semop(2)` is used to implement P() (wait) and V() (signal). However, `semop(2)` can do multiple semaphore operations with one call. For our programs, we will only be creating arrays with one semaphore on them and doing only one operation at a time.

- Run the following commands on both your Linux host and target, and observe the status of System V IPC status.
 - `ipcs`
 - `ipcs -s`
- Here is a program, **makesem.c**. Compile it and run it with two parameters:
 - the first parameter should be a large number;
 - the second one is number '1' or '0'.

自訂
key 可以高底，也可以從外面傳進來

For example: makesem 428361733 1. 判斷何處 |

Now run ipcs again. You should see your semaphore in the list. Note that the key is the number you entered that converted to hex.

```
/* makesem.c
 *
 * This program creates a semaphore. The user should pass
 * a number to be used as the semaphore key and initial
 * value as the only command line arguments. If that
 * identifier is not taken, then a semaphore will be created.
 * If a semaphore is set so that then no semaphore will be
 * created. The semaphore is set so that anyone on the system
 * can use it.
 */

#include <errno.h>
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <sys/sem.h>

#define SEMMODE 0666 /* rw(owner)-rw(group)-rw(other) permission */

int main (int argc, char **argv)
{
    int s;
    long int key;
    int val;

    if (argc != 3)
    {
        fprintf(stderr,
                "%s: specify a key (long) and initial value (int)\n",
                argv[0]);
        exit(1);
    }

    /* get values from command line */
    if (sscanf(argv[1], "%ld", &key) != 1)
    {
        /* convert arg to long integer */
        fprintf(stderr, "%s: argument #1 must be a long integer\n",
                argv[0]);
        exit(1);
    }
    if (sscanf(argv[2], "%d", &val) != 1)
    {
        /* convert arg to long integer */
        fprintf(stderr, "%s: argument #2 must be an integer\n",
                argv[0]);
        exit(1);
    }
}
```

} *(S) get.*
/ semget() takes three parameters */*
/ Options:*
** IPC_CREAT - create a semaphore if not exists*
** IPC_EXCL - creation fails if it already exists*
** SEM_MODE - access permission*
**/*
s = semget(
key, / the unique name of the semaphore on the system */*
1, / we create an array of semaphores, but just need 1. */*
IPC_CREAT | IPC_EXCL | SEMMODE); RW-
666
/ If semget () returns -1 then it failed. However,*
** if it returns any other number >= 0 then that becomes*
** the identifier within the program for accessing the semaphore.*
**/*
if (s < 0)
{
fprintf(stderr,
"%s : creation of semaphore %ld failed : %s\n", argv[0],
key, strerror(errno));
exit(1);
}
printf("Semaphore %ld created\n", key);

/ set semaphore [s[0]] value to initial value (val) */*
if (semctl(s, 0, SETVAL, val) < 0)
{
fprintf(stderr,
"%s : Unable to initialize semaphore : %s\n",
argv[0], strerror(errno));
exit(0);
}
printf("Semaphore %ld has been initialized to %d\n", key, val);

return 0;
}

- Program **rmsem.c** removes the semaphore identified by its key from the system. Compile and run the program with your key from the last step, then run **ipcs** to see if your semaphore is still listed. Note: The easiest way to create the file is probably to copy **makesem.c** and make the modifications.

```

/* rmsem.c
 *
 * This program destroys a semaphore. The user should pass a number
 * to be used as the semaphore key.
 */

#include <errno.h>
#include <stdio.h>
#include <stdlib.h>

```

```

#include <string.h>
#include <sys/sem.h>

int main( int argc , char **argv )
{
    int s;
    long int key;

    if ( argc != 2 )
    {
        fprintf( stderr , "%s : specify a key (long)\n" , argv [ 0 ] );
        exit ( 1 );
    }

    /* get values from command line */
    if ( sscanf( argv [ 1 ] , "%ld" , &key )!=1 )
    {
        /* convert arg to long integer */
        fprintf( stderr ,
            "%s : argument #1 must be an long integer\n" , argv [ 0 ] );
        exit ( 1 );
    }

    /* find semaphore */
    s = semget( key , 1 , 0 );
    if ( s < 0 )
    {
        fprintf( stderr , "%s : failed to find semaphore %d : %s\n" ,
            argv [ 0 ] , key , strerror( errno ) );
        exit ( 1 );
    }
    printf( "Semaphore %d found\n" , key );

    /* remove semaphore */
    if ( semctl ( s , 0 , IPC_RMID , 0 ) < 0 )
    {
        fprintf ( stderr , "%s : unable to remove semaphore %d\n" ,
            argv [ 0 ] , key );
        exit ( 1 );
    }
    printf( "Semaphore %d has been removed\n" , key );

    return 0;
}

```

- The program **doodle.c** has definitions for the operations P() (wait) and V() (signal). The program first "opens" the semaphore with **semget()**, then it waits for the user to type in a small integer number (number of seconds to stay in the critical section). It then performs P() on a semaphore. Once inside the critical section, it doodles for the number of seconds you specified, then it leaves performing V() on the semaphore.

To demonstrate the operation of `doodle`, please run the following two experiments:

1. Use `makesem` to create a semaphore with initial value 1 and run three `doodle` programs to acquire the same semaphore simultaneously.
2. Use `makesem` to create a semaphore with initial value 2 and run three `doodle` programs to acquire the same semaphore simultaneously.

Observe the inter-operations between the three `doodle` programs.

```
/* doodle.c
 *
 * This program shows how P () and V () can be implemented,
 * then uses a semaphore that everyone has access to.
 */

#include <errno.h>
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <sys/sem.h>
#include <unistd.h>

#define DOODLESEMKEY 1122334455

/* P () - returns 0 if OK; -1 if there was a problem */
int P(int s)
{
    struct sembuf sop; /* the operation parameters */
    sop.sem_num = 0; /* access the 1st (and only) sem in the array */
    sop.sem_op = -1; /* wait...*/
    sop.sem_flg = 0; /* no special options needed */

    if (semop (s, &sop, 1) < 0) {
        fprintf(stderr, "P(): semop failed: %s\n", strerror(errno));
        return -1;
    } else {
        return 0;
    }
}

/* V() - returns 0 if OK; -1 if there was a problem */
int V(int s)
{
    struct sembuf sop; /* the operation parameters */
    sop.sem_num = 0; /* the 1st (and only) sem in the array */
    sop.sem_op = 1; /* signal */
    sop.sem_flg = 0; /* no special options needed */

    if (!semop(s, &sop, 1) < 0) {
        fprintf(stderr, "V(): semop failed: %s\n", strerror(errno));
        return -1;
    } else {
```

```

        return 0;
    }

int main ( int argc , char **argv )
{
    int s , secs ;
    long int key ;

    if ( argc != 2 ) {
        fprintf ( stderr , "%s : specify a key (long)\n" , argv [ 0 ] );
        exit ( 1 );
    }

    /* get values from command line */
    if ( sscanf ( argv [ 1 ] , "%ld" , &key ) != 1 )
    {
        /* convert arg to long integer */
        fprintf ( stderr ,
                  "%s : argument #1 must be an long integer\n" , argv [ 0 ] );
        exit ( 1 );
    }

    s = semget ( key , 1 , 0 ); find P

    if ( s < 0 ) {
        fprintf ( stderr ,
                  "%s : cannot find semaphore %d : %s\n" ,
                  argv [ 0 ] , key , strerror ( errno ) );
        exit ( 1 );
    }

    while ( 1 ) {
        printf ( "#secs to doodle in the critical section? (0 to exit):" );
        scanf ( "%d" , &secs );

        if ( secs == 0 )
            break;

        printf ( "Preparing to enter the critical section ..\n" );
P(s);
        printf ( "Now in the critical section! Sleep %d secs ..\n" , secs );

        while ( secs ) {
            printf ( "%d ... doodle ...\n" , secs -- );
            sleep ( 1 );
        }

        printf ( "Leaving the critical section ..\n" );
V(s);
    }
}

```

```

    }
    return 0;
}

```

- Race Condition

Race conditions arise in software when separate processes or threads of execution depend on some shared state. Operations upon shared states are critical sections that must be mutually exclusive in order to avoid harmful collision between processes or threads that share those states. Assume that two threads (T1 and T2) want to increment the value of a global integer by one. Ideally, the following sequence of operations would take place:

- Integer $i = 0$; (memory)
- T1 reads the value of i from memory into register1: 0
- T1 increments the value of i in register1: $(\text{register1 contents}) + 1 = 1$
- T1 stores the value of register1 in memory: 1
- T2 reads the value of i from memory into register2: 1
- T2 increments the value of i in register2: $(\text{register2 contents}) + 1 = 2$
- T2 stores the value of register2 in memory: 2
- Integer $i = 2$; (memory)

In the above case, the final value of i is 2, as expected. However, if the two threads run simultaneously without locking or synchronization, the outcome of the operation could be wrong. The alternative sequence of operations below demonstrates this scenario:

- race condition*
- race condition*
- critical section*
- 不希望发生*
- Integer $i = 0$; (memory)
 - T1 reads the value of i from memory into register1: 0
 - T2 reads the value of i from memory into register2: 0
 - T1 increments the value of i in register1: $(\text{register1 contents}) + 1 = 1$
 - T2 increments the value of i in register2: $(\text{register2 contents}) + 1 = 1$
 - T1 stores the value of register1 in memory: 1
 - T2 stores the value of register2 in memory: 1
 - Integer $i = 1$; (memory)

The final value of i is 1 instead of the expected result of 2. This occurs because the increment operations of the second case are not mutually exclusive.

The following example shows a race condition between two processes.

Please create a file `counter.txt` and set initial value 0 to the file by performing “`echo 0 > counter.txt`”. Compile the codes with and without the flag ‘`-D USE_SEM`’ and check the result in `counter.txt`. Explain the result if any difference.

```

/*
 * race.c
 */

#include <errno.h>
#include <fcntl.h>
#include <stdio.h>

```

```

#include <stdlib.h>
#include <string.h>
#include <sys/ipc.h>
#include <sys/sem.h>
#include <sys/stat.h>
#include <sys/types.h>
#include <sys/wait.h>
#include <unistd.h>

#ifndef USESEM
#define SEMMODE 0666 /* rw(owner)-rw(group)-rw(other) permission */
#define SEMKEY 1122334455

int sem; 全局

/* P () - returns 0 if OK; -1 if there was a problem */
int P (int s)
{
    struct sembuf sop; /* the operation parameters */
    sop.sem_num = 0; /* access the 1st (and only) sem in the array */
    sop.sem_op = -1; /* wait...*/
    sop.sem_flg = 0; /* no special options needed */

    if (semop (s, &sop, 1) < 0) {
        fprintf(stderr, "P(): semop failed: %s\n", strerror(errno));
        return -1;
    } else {
        return 0;
    }
}

/* V() - returns 0 if OK; -1 if there was a problem */
int V(int s)
{
    struct sembuf sop; /* the operation parameters */
    sop.sem_num = 0; /* the 1st (and only) sem in the array */
    sop.sem_op = 1; /* signal */
    sop.sem_flg = 0; /* no special options needed */

    if (semop(s, &sop, 1) < 0) {
        fprintf(stderr, "V(): semop failed: %s\n", strerror(errno));
        return -1;
    } else {
        return 0;
    }
}

#endif

/* increment value saved in file */
void Increment()

```

```

{
    int ret;
    int fd;      /* file descriptor */
    int counter;
    char buffer[100];
    int i = 10000;

    while(i)
    {
        /* open file */
        fd = open("./counter.txt", O_RDWR );
        if (fd < 0)
        {
            printf("Open_counter.txt_error.\n");
            exit(-1);
        }

#define USE_SEM
/* acquire semaphore */
P(sem);

#endif

        /***** Critical Section *****/
        /* clear */
        memset(buffer, 0, 100);

        /* read raw data from file */
        ret = read(fd, buffer, 100);
        if (ret < 0)
        {
            perror("read_counter.txt");
            exit(-1);
        }

        /* transfer string to integer & increment counter */
        counter = atoi(buffer);
        counter++;

        /* write back to counter.txt */
        lseek(fd, 0, SEEK_SET); /* reposition to the head of file */

        /* clear */
        memset(buffer, 0, 100);
        sprintf(buffer, "%d", counter);
        ret = write(fd, buffer, strlen(buffer));
        if (ret < 0)
        {
            perror("write_counter.txt");
            exit(-1);
        }
    } /* **** Critical Section *****/
}

```

```

#define USESEM
    /* release semaphore */
    V(sem);
#endif

    /* close file */
    close(fd);

    i--;
}
}

int main(int argc, char **argv)
{
    int childpid;
    int status;

#define USESEM
    /* create semaphore */
    sem = semget(SEMKEY, 1, IPC_CREAT | IPC_EXCL | SEM_MODE);
    if (sem < 0)
    {
        fprintf(stderr, "Creation of semaphore %d failed : %s\n",
                SEMKEY, strerror(errno));
        exit(-1);
    }

    /* initial semaphore value to 1 (binary semaphore) */
    if (semctl(sem, 0, SETVAL, 1) < 0)
    {
        fprintf(stderr, "Unable to initialize semaphore : %s\n",
                strerror(errno));
        exit(0);
    }

    printf("Semaphore %d has been created & initialized to 1\n",
           SEMKEY);
#endif

    /* fork process */
    if ((childpid = fork()) > 0) /* parent */
    {
        Increment();
        waitpid(childpid, &status, 0);
    }
    else if (childpid == 0) /* child */
    {
        Increment();
        exit(0);
    }
}

```

22-9
descriptor

```

    else      /* error */
    {
        perror("fork");
        exit(-1);
    }

#ifndef USE_SEM
/* remove semaphore */
if (semctl (sem, 0, IPC_RMID, 0) < 0)
{
    fprintf (stderr , "%s: unable to remove semaphore %ld\n",
            argv[0], SEM_KEY);
    exit(1);
}
printf("Semaphore %ld has been removed\n", SEM_KEY);
#endif

return 0;
}

```

4 Mutex

A mutex is abbreviated from "MUTual EXclusion", and is useful for protecting shared data structures from concurrent modifications, and implementing critical sections and monitors. A mutex has two possible states: unlocked (not owned by any thread), and locked (owned by one thread). A mutex can never be owned by two different threads simultaneously. A thread attempting to lock a mutex that is already locked by another thread is suspended until the owning thread unlocks the mutex first. Some POSIX library APIs for mutex are:

- `pthread_mutex_init()` initializes the mutex referenced by `mutex` with attributes specified by `attr`. Upon successful initialization, the state of the mutex becomes initialized and unlocked. For more information, please see its man page.
- `pthread_mutex_lock()` locks the mutex object referenced by the mutex. If the mutex is already locked, the calling thread blocks until the mutex becomes available. This operation returns with the mutex object referenced by `mutex` in the locked state with the calling thread as its owner.
- `pthread_mutex_unlock()` unlocks the mutex object referenced by the mutex.
- `pthread_mutex_destroy()` destroys the mutex object referenced by `mutex`; the mutex object becomes, ineffective, uninitialized.

```

#include <pthread.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

#define checkResults(string, val) { \
    if (val) { \
        printf("Failed with %d at %s", val, string); \
        exit(1); \
    } \
}

```

```

}

#define NUMTHREADS 3

pthread_mutex_t mutex = PTHREAD_MUTEX_INITIALIZER;
int sharedData = 0;
int sharedData2 = 0;

void *theThread(void *parm)
{
    int rc;

    printf("\tThread %lu : Entered\n", (unsigned long) pthread_self());
    /* lock mutex */ recusive lock
    rc = pthread_mutex_lock(&mutex);
    checkResults("pthread_mutex_lock()\n", rc);

    // Critical Section *****
    printf("\tThread %lu : Start_critical_section , holding lock\n",
           (unsigned long) pthread_self());

    /* Access to shared data goes here */
    ++sharedData;
    --sharedData2;

    printf("\tsharedData = %d , sharedData2 = %d\n",
           sharedData, sharedData2);

    printf("\tThread %lu : End_critical_section , release lock\n",
           (unsigned long) pthread_self());
    // ***** Critical Section *****

    /* unlock mutex */
    rc = pthread_mutex_unlock(&mutex);
    checkResults("pthread_mutex_unlock()\n", rc);

    return NULL;
}

int main(int argc, char **argv)
{
    pthread_t thread[NUMTHREADS];
    int rc = 0;
    int i;

    /* lock mutex */
    printf("Main_thread hold mutex to prevent access to shared data\n");
    rc = pthread_mutex_lock(&mutex);
    checkResults("pthread_mutex_lock()\n", rc);
}

```

```

/* create thread */
printf("Main_thread_create/start_threads\n");
for ( i = 0; i < NUMTHREADS; ++i) {
    rc = pthread_create(&thread[ i ], NULL, theThread, NULL);
    checkResults("pthread_create()\n", rc);
}

/* wait for thread creation complete */
printf("Main_thread_wait_a_bit_until_done' with the shared data\n");
sleep(3);

/* unlock mutex */
printf("Main_thread_unlock_shared_data\n");
rc = pthread_mutex_unlock(&mutex);
checkResults("pthread_mutex_lock()\n", rc);

/* wait thread complete */
printf("Main_thread_Wait_for_threads_to_complete , "
      "and_release_their_resources\n");
for ( i=0; i <NUMTHREADS; ++i) {
    rc = pthread_join(thread[ i ], NULL);
    checkResults("pthread_join()\n", rc);
}

/* destroy mutex */
printf("Main_thread_clean_up_mutex\n");
rc = pthread_mutex_destroy(&mutex);

printf("Main_thread_completed\n");

return 0;
}

```

带参数的函数（function pointer）
通过函数的参数

Lab 5: Inter-Process Communication (Part II)

Yu-Lun Huang

2019-07-16

1 Objective

- Be familiar with inter-process communication: pipe, shared memory, etc.

2 Prerequisite

- Read man pages of `pipe()`, `shmget()`, `shmctl()`, `shmat()`, `shmdt()`, etc.

3 Pipe

A pipe creates a pair of file descriptors, pointing to a pipe inode, and places them in the array pointed to by `filedes`. `filedes[0]` is for reading, `filedes[1]` is for writing. An unnamed pipe can be viewed and accessed by processes with parent-child relationship. To share a pipe among all processes, you need to create a named pipe.

The following program creates a pipe, and then `fork(2)` to create a child process. After the `fork(2)`, each process closes the descriptors that it doesn't need for the pipe (see `pipe(7)`). In the following sample code, the child process reads the file specified in `argv[1]`, and writes the file content to the parent process through the pipe. The parent process reads the data from the pipe and echoes the data on the screen.

```
1 /* pipe.c
2 *
3 * child process read the content of file
4 * and write the content to parent process through pipe
5 */
6
7 #include <fcntl.h>
8 #include <stdio.h>
9 #include <stdlib.h>
10 #include <string.h>
11 #include <sys/stat.h>
12 #include <sys/types.h>
13 #include <sys/wait.h>
14 #include <unistd.h>
15
16 int pfd[2]; /* pfd[0] is read end, pfd[1] is write end */
17
18 void ChildProcess(char *path)
19 {
20     int fd;
```

```

21 int ret;
22 char buffer[100];
23
24 /* close unused read end */
25 close(pfd[0]);
26
27 /* open file */
28 fd = open(path, O_RDONLY);
29 if (fd < 0) {
30     printf("Open %s failed.\n", path);
31     exit(EXIT_FAILURE);
32 }
33
34 /* read file and write content to pipe */
35 while (1) {
36     /* read raw data from file */
37     ret = read(fd, buffer, 100);
38
39     if (ret < 0) { /* error */
40         perror("read()");
41         exit(EXIT_FAILURE);
42     }
43     else if (ret == 0) { /* reach EOF */
44         close(fd); /* close file */
45         close(pfd[1]); /* close write end, reader see EOF */
46         exit(EXIT_SUCCESS);
47     }
48     else { /* write content to pipe */
49         write(pfd[1], buffer, ret);
50     }
51 }
52 }
53
54 void ParentProcess()
55 {
56     int ret;
57     char buffer[100];
58
59     /* close unused write end */
60     close(pfd[1]);
61
62     /* read data from pipe until reach EOF */
63     while(1) {
64         ret = read(pfd[0], buffer, 100);
65
66         if (ret > 0) { /* print data to screen */
67             printf("%.*s", ret, buffer);
68         }
69         else if (ret == 0) { /* reach EOF */
70             close(pfd[0]); /* close read end */
71             wait(NULL);

```

```

72     exit(EXIT_SUCCESS);
73 }
74 else {
75     perror("pipe_read()");
76     exit(EXIT_FAILURE);
77 }
78 }
79 }
80
81 int main(int argc, char *argv[])
82 {
83     pid_t cpid;
84
85     if (argc != 2) {
86         fprintf(stderr, "%s: -specify -a -file \n", argv[0]);
87         exit(1);
88     }
89
90     /* create pipe */
91     if (pipe(pfd) == -1) {
92         perror("pipe");
93         exit(EXIT_FAILURE);
94     }
95
96     /* fork child process */
97     cpid = fork();
98     if (cpid == -1) { /* error */
99         perror("fork");
100        exit(EXIT_FAILURE);
101    }
102
103    if (cpid == 0)
104        ChildProcess(argv[1]);
105    else
106        ParentProcess();
107
108    return 0;
109}

```

Monitoring multiple file descriptors with polling strategy may cause busy waiting, which lowers down the system performance. To effectively monitor multiple descriptors, you can use the system call `select()` to perform block waiting, rather than busy waiting. The calling process specifies the interesting file descriptors to `select()` and performs blocking waiting. When one or more descriptors become "ready" (ready for read or ready for write), `select()` informs the calling process to awake from blocking state. Then, a user can check each file descriptor to perform read/write operation. In following program, two child processes sleep a random time, then send a message to the parent process. The parent process uses `select()` to perform blocking waiting, until one of the pipes is ready to read.

```

1 /*
2  * select.c
3  */
4

```

```

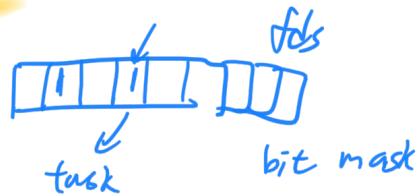
5 #include <stdio.h>
6 #include <stdlib.h>
7 #include <string.h>
8 #include <sys/time.h>
9 #include <sys/types.h>
10 #include <time.h>
11 #include <unistd.h>
12
13 #define max(a, b) ((a > b) ? a : b)
14
15 void ChildProcess(int *pfds, int sec)
16 {
17     char buffer[100];
18
19     /* close unused read end */
20     close(pfd[0]);
21
22     /* sleep a random time to wait parent process enter select() */
23     printf("Child_process_(%d)_wait_%d_secs\n", getpid(), sec);
24     sleep(sec);
25
26     /* write message to parent process */
27     memset(buffer, 0, 100);
28     sprintf(buffer, "Child_process_(%d)_sent_message_to_parent_process\n",
29             getpid());
30     write(pfd[1], buffer, strlen(buffer));
31
32     /* close write end */
33     close(pfd[1]);
34
35     exit(EXIT_SUCCESS);
36 }
37
38 int main(int argc, char *argv[])
39 {
40     int pfd1[2], pfd2[2]; /* pipe's fd */
41     int cpid1, cpid2; /* child process id */
42     fd_set rfds, arfds; ↗ why 杂乱?
43     int max_fd;
44     struct timeval tv;
45     int retval;
46     int fd_index;
47     char buffer[100];
48
49     /* random seed */
50     srand(time(NULL));
51
52     /* create pipe */
53     pipe(pfd1);
54     pipe(pfd2);
55 }
```

```

int select (int nfd, fd_set *readfds, fd_set *writefds, fd_set *exceptfds, struct timeval *timeout);
FD_SET(int fd, fd_set *fdset); FD_CLR, FD_ISSET(), FD_ZERO(fd_set *fdset);

56 /* create 2 child processes and set corresponding pipe & sleep time */
57 cpid1 = fork();
58 if (cpid1 == 0)
59     ChildProcess(pfd1, random() % 5);
60
61 cpid2 = fork();
62 if (cpid2 == 0)
63     ChildProcess(pfd2, random() % 4);
64
65 /* close unused write end */
66 close(pfd1[1]);
67 close(pfd2[1]);
68
69 /* set pfd1[0] & pfd2[0] to watch list */
70 FD_ZERO(&rfds);
71 FD_ZERO(&arfds);
72 FD_SET(pfd1[0], &arfds); → 2 pipes
73 FD_SET(pfd2[0], &arfds);
74 max_fd = max(pfd1[0], pfd2[0]) + 1;
75
76 /* Wait up to five seconds. */ 5秒間
77 tv.tv_sec = 5;
78 tv.tv_usec = 0;
79
80 while(1)
81 {
82     /* config fd_set for select */
83     memcpy(&rfds, &arfds, sizeof(rfds)); ?
84
85     /* wait until any fd response */
86     retval = select(max_fd, &rfds, NULL, NULL, &tv);
87
88     if (retval == -1) { /* error */
89         perror("select()");
90         exit(EXIT_FAILURE);
91     }
92     else if (retval) { /* # of fd got response */
93         printf("Data is available now.\n");
94     }
95     else { /* no fd response before timer expired */
96         printf("No data within five seconds.\n");
97         break;
98     }
99
100    /* check if any response */
101    for (fd_index = 0; fd_index < max_fd; fd_index++)
102    {
103        if (!FD_ISSET(fd_index, &rfds))
104            continue; /* no response */
105
106        retval = read(fd_index, buffer, 100);

```



```

107
108     if (retval > 0)      /* read data from pipe */
109         printf("%.*s", retval, buffer);
110     else if (retval < 0) /* error */
111         perror("pipe_read()");
112     else {               /* write fd closed */
113         /* close read fd */
114         close(fd_index);
115         /* remove fd from watch list */
116         FD_CLR(fd_index, &arfds); ??
117     }
118 }
119
120
121 return 0;
122 }
```

4 Shared Memory

Shared memory is a memory space that may be simultaneously accessed by multiple processes with an intent to provide inter-process communication among them or avoid redundant copies. Unlike unnamed pipes, only exist among processes with parent-child relationship, every process can access the shared memory space with the **share memory key** specified.

The followings are two processes communicating via shared memory: **shm_server.c** and **shm_client.c**. The two programs here illustrate the passing of a simple piece of memory (a string) between the processes if running simultaneously:

```

1 /*
2  * shm_server.c -- creates the string and shared memory.
3  */
4
5 #include <stdio.h>
6 #include <stdlib.h>
7 #include <sys/ipc.h>
8 #include <sys/shm.h>
9 #include <sys/types.h>
10 #include <unistd.h>
11
12 #define SHMSZ    27
13
14 int main(int argc, char *argv[])
15 {
16     char c;
17     int shmid;
18     key_t key;
19     char *shm, *s;
20     int retval;
21
22     /* We'll name our shared memory segment "5678" */
23     key = 5678;
24 }
```

```

25  /* Create the segment */
26  if ((shmid = shmget(key, SHMSZ, IPC_CREAT | 0666)) < 0) {
27      perror("shmget");
28      exit(1);
29  }
30
31  /* Now we attach the segment to our data space */
32  if ((shm = shmat(shmid, NULL, 0)) == (char *) -1) {
33      perror("shmat");
34      return the addr of the attached shm segment : ok 1
35  }           ok :-|
36  printf("Server-create-and-attach-the-share-memory.\n");
37
38  /* Now put some things into the memory for the other process to read */
39  s = shm;
40
41  printf("Server-write-a-z-to-share-memory.\n");
42  for (c = 'a'; c <= 'z'; c++)
43      *s++ = c;
44  *s = '\0';
45
46  /*
47   * Finally, we wait until the other process changes the first
48   * character of our memory to '*', indicating that it has read
49   * what we put there.
50  */
51  printf("Waiting-other-process-read-the-share-memory....\n");
52  while (*shm != '*')
53      sleep(1);
54  printf("Server-read-* from the share memory.\n");
55
56  /* Detach the share memory segment */
57  shmdt(shm); Success failed -1          ok 使用 shm 成功 失败
58
59  /* Destroy the share memory segment */
60  printf("Server-destroy-the-share-memory.\n");
61  retval = shmctl(shmid, IPC_RMID, NULL);
62  if (retval < 0)
63  {
64      fprintf(stderr, "Server-remove-share-memory-failed\n");
65      exit(1);
66  }
67
68  return 0;
69 }

```

```

1  /*
2  * shm_client.c -- attaches itself to the created shared memory
3  *                   and uses the string (printf).
4  */
5

```

```

6 #include <stdio.h>
7 #include <stdlib.h>
8 #include <sys/ipc.h>
9 #include <sys/shm.h>
10 #include <sys/types.h>
11
12 #define SHMSZ      27
13
14 int main(int argc, char *argv[])
15 {
16     int shmid;
17     key_t key;
18     char *shm, *s;
19
20     /* We need to get the segment named "5678", created by the server */
21     key = 5678;           same as server's
22
23     /* Locate the segment */
24     if ((shmid = shmget(key, SHMSZ, 0666)) < 0) {
25         perror("shmget");
26         exit(1);
27     }           server write rig
28
29     /* Now we attach the segment to our data space */
30     if ((shm = shmat(shmid, NULL, 0)) == (char *) -1) {
31         perror("shmat");
32         exit(1);
33     }
34     printf("Client attached the share memory created by server.\n");
35
36     /* Now read what the server put in the memory */
37     printf("Client read characters from share memory...\n");
38     for (s = shm; *s != '\0'; s++)
39         putchar(*s);
40     putchar('\n');
41
42     /*
43      * Finally, change the first character of the segment to '*',
44      * indicating we have read the segment.
45      */
46     printf("Client wrote * to the share memory.\n");
47     *shm = '*';
48
49     /* Detach the share memory segment */
50     printf("Client detached the share memory.\n");
51     shmdt(shm);
52
53     return 0;
54 }
```

修改板子，廠商含 12 位 I/O 在 source code
(microtime)

Lab 6-1: I/O Control for PXA270

Yu-Lun Huang

2019-07-12

There is an 8-bit LED lamps on the motherboard of PXA270, numbered from D9(1) to D16(8). We use the `creator-pxa270-lcd.ko` module to control the LED lamps, it is also used to drive the LCD, 7-segment LED, KeyPAD, and DIP Switch.

1 Compile Modules

Rebuild your kernel and rootfile system to support `creator-pxa270-lcd.ko` module:

- Get the source code of `creator-pxa270-lcd.ko`:

Download Creator_PXA270_LCD_Device_Driver.src.tar.gz from E3 website and decompressed it to your kernel source.

tar board 1325P? tar z

A: 解開

Z

- Configure kernel source:

SHELL> cd ~
SHELL> tar xzvf Creator_PXA270_LCD_Device_Driver.src.tar.gz
" " Z
) 儲存一些東西解到資料夾 最後面有 Z f 到檔案那層

(device driver character device, 開機程序

開等

In the window of "Linux Kernel Configuration", select "Load an Alternate Configuration from File" and load the configuration file `arch/arm/configs/creator_pxa270_defconfig`.

Select "Device Drivers" → "Character devices" and mark "Creator-pxa270 LCD" as [M].

Save and exit kernel configuration.

default configne 進來選

- Make Image:

Compile Linux kernel and `creator-pxa270-lcd.ko` module.

SHELL> make clean
SHELL> make

N - 不要選
Y - Yes 選
M - mod of
有掛到的

上傳

The `creator-pxa270-lcd.ko` module will be placed at `microtime/linux/drivers/char/`.

mount module

- patch
- Make new root filesystem:

Copy `creator-pxa270-lcd.ko` module into root filesystem.

SHELL> cp ~/microtime/linux/drivers/char/creator-pxa270-lcd.ko
~/microtime/rootfs/lib/modules/2.6.15.3/kernel/drivers/char/

.exe

image

data

Then, rebuild and flash root filesystem.

重編 重建

rootfs 的 file 拷貝

1

linux
K src

rootfs

2 Load Modules

Type the following command to load the creator-pxa270-lcd.ko on PXA270.

把它掛載上去
insert module
insmod .ko

3 PXA I/O: Control LED

- LED programming guide

vmmmod
modprobe ?

Header file: 開源模擬の i 不可用 < >

```
1 #include "asm-arm/arch-pxa/lib/creator_pxa270_lcd.h"
```

Commands:

```
1 LED_IOCTL_SET // set the specified LED (D9 - D16)
2 LED_IOCTL_CLEAR // clear the specified LED (D9 - D16)
```

Values:

1	LED_ALL_ON	0xFF
2	LED_ALL_OFF	0x00
3	LED_D9_INDEX	1
4	LED_D10_INDEX	2
5	LED_D11_INDEX	3
6	LED_D12_INDEX	4
7	LED_D13_INDEX	5
8	LED_D14_INDEX	6
9	LED_D15_INDEX	7
10	LED_D16_INDEX	8

1 7 與 D9 到 D15

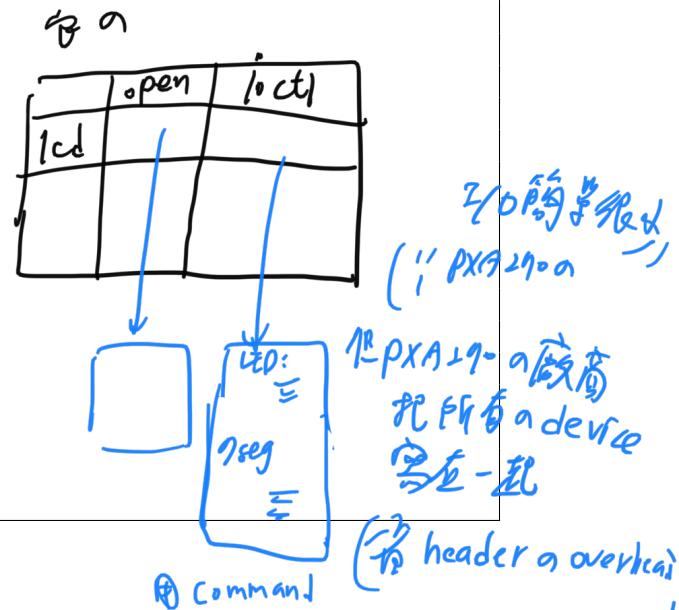
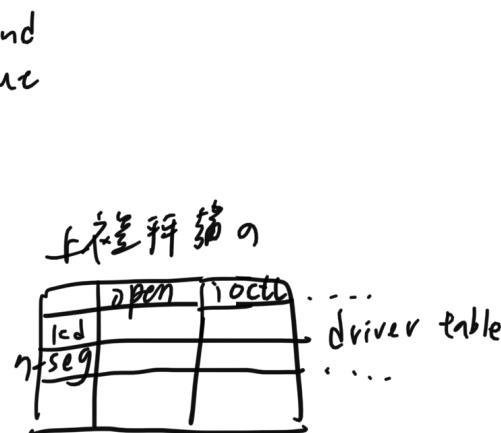
Sample code:

```
1 /*
2  * led.c -- the sample code for controlling LEDs on Creator.
3  */
4
5 #include <stdio.h>
6 #include <stdlib.h>
7 #include <sys/fcntl.h>
8 #include <sys/ioctl.h>
9 #include <unistd.h>
10 #include "asm-arm/arch-pxa/lib/creator_pxa270_lcd.h"
11
12 int main(int argc, char *argv[])
13 {
14     int fd; /* file descriptor for /dev/lcd */
15     int retval;
16
17     unsigned short data;
```

装置表 device table
↑ (instant)

ls /dev/loop* 看有没有

```
19 /* Open device /dev/lcd */
20 if( (fd = open( "/dev/lcd" , O_RDWR)) < 0)
21 {
22     printf( "Open /dev/lcd failed .\n" );
23     exit( -1);
24 }
25
26 /* Turn on all LED lamps */
27 data = LED_ALL_ON;
28 ioctl(fd , LED_IOCTL_SET , &data );
29 printf("Turn_on_all_LED_lamps\n");
30 sleep( 3);
31
32 /* Turn off all LED lamps */
33 data = LED_ALL_OFF;
34 ioctl(fd , LED_IOCTL_SET , &data );
35 printf("Turn_off_all_LED_lamps\n");
36 sleep( 3);
37
38 /* Turn on D9 */
39 data = LED_D9_INDEX;
40 ioctl(fd , LED_IOCTL_BIT_SET , &data );
41 printf("Turn_on_D9\n");
42 sleep( 3);
43
44 /* Turn off D9 */
45 data = LED_D9_INDEX;
46 ioctl(fd , LED_IOCTL_BIT_CLEAR , &data );
47 printf("Turn_off_D9\n");
48 sleep( 3);
49
50 /* Close fd */
51 close(fd );
52
53 return 0;
54 }
```



Add the header search path when compile `led.c`.

```
SHELL> arm-unknown-linux-gnu-gcc -o led led.c  
-L /opt/arm-unknown-linux-gnu/arm-unknown-lin  
-I /opt/arm-unknown-linux-gnu/arm-unknown-lin
```

-I /home/lab616/microtime/linux/include/ -I /mnt

建立新呼叫 module

4 PXA I/O: 7-Segment

- Header files: `asm-arm/arch-pxa/lib/creator_pxa270_lcd.h`
 - Function: `ioctl(fd, command, data)`
 - Command:

- _7SEG_IOCTL_ON: turn on 7 segment LED (no data is needed)
- _7SEG_IOCTL_OFF: turn off 7 segment LED (no data is needed)
- _7SEG_IOCTL_SET: set 7 segment LED (_7seg_info_t)

不
用
看
自
己
就
可
以

- Data:

```

1 typedef struct _7Seg_Info{
2     unsigned char Mode; // _7SEG_MODE_PATTERN or _7SEG_MODE_HEX_VALUE
3     unsigned char Which; // D5 ~ D8
4     unsigned long Value; // pattern or hex
5 } _7seg_info_t;

```

- The setting of _7Seg_Info:

- flags used in Mode:

1 #define _7SEG_MODE_PATTERN	0	show 自定义的图方便
2 #define _7SEG_MODE_HEX_VALUE	1	show 数字方便

- flags used in Which:

1 #define _7SEG_D5_INDEX	8 // Segment	D5 (1)
2 #define _7SEG_D6_INDEX	4 // Segment	D6 (2)
3 #define _7SEG_D7_INDEX	2 // Segment	D7 (3)
4 #define _7SEG_D8_INDEX	1 // Segment	D8 (4)
5 #define _7SEG_ALL		
6 (_7SEG_D5_INDEX _7SEG_D6_INDEX _7SEG_D7_INDEX _7SEG_D8_INDEX)		

The following is the sample code of 7 segment display control.

```

1 _7seg_info_t    data;
2 int             fd, ret, i;
3 if ((fd = open("/dev/lcd", O_RDWR)) < 0)  return (-1);
4
5 ioctl(fd, _7SEG_IOCTL_ON, NULL);
6 data.Mode = _7SEG_MODE_HEX_VALUE;          开放 OnF
7 data.Which = _7SEG_ALL;                    HEX_VALUE      D5 D6 D7 D8
8 data.Value = 0x2004;                      2 0 + 4
9
10 ioctl(fd, _7SEG_IOCTL_SET, &data);
11 sleep (3);
12 data.Mode = _7SEG_MODE_PATTERN;
13 data.Which = _7SEG_D5_INDEX | _7SEG_D8_INDEX;
14 data.Value = 0x0d7f /* change to 5008 */;
15 ioctl(fd, _7SEG_IOCTL_SET, &data);
16 ioctl(fd, _7SEG_IOCTL_OFF, NULL);
17 close(fd);

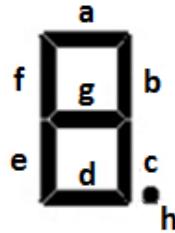
```

5 PXA I/O: Keypad

- Keypad I/O

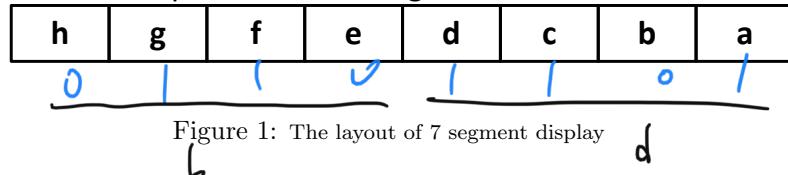
7-segment

The bit value 1: on
The bit value 0: off



不亮
亮

8 bits to represents each Segment



- Header files: asm-arm/arch-pxa/lib/creator_pxa270_lcd.h
- Function: ioctl(fd, command, data)
 - * Command:
 - KEY_IOCTL_GET_CHAR: unsigned short, get its ASCII value.
 - KEY_IOCTL_WAIT_CHAR: wait until get a character.
 - KEY_IOCTL_CHECK_EMPTY
 - KEY_IOCTL_CLEAR
 - KEY_IOCTL_CANCEL_WAIT_CHAR

這樣子

- * Definition

```

1 #define VK_S2 1      /* ASCII = '1' */
2 #define VK_S3 2      /* ASCII = '2' */
3 #define VK_S4 3      /* ASCII = '3' */
4 #define VK_S5 10     /* ASCII = 'A' */
5 #define VK_S6 4
6 #define VK_S7 5
7 #define VK_S8 6
8 #define VK_S9 11
9 #define VK_S10 7
10 #define VK_S11 8
11 #define VK_S12 9
12 #define VK_S13 12
13 #define VK_S14 14    /* ASCII = '*' */
14 #define VK_S15 0
15 #define VK_S16 15    /* ASCII = '#' */
16 #define VK_S17 13

```

- * Sample

```

1 unsigned short key;
2 int fd, ret;

```

```

3   if ((fd = open("/dev/lcd", O_RDWR)) < 0) return (-1);
4
5   ioctl(fd, KEY_IOCTL_CLEAR, &key); 前面可能有殘存的buffer
6   while(1) {
7       ret = ioctl(fd, KEY_IOCTL_CHECK_EMTPY, &key)
8       if (ret < 0) {
9           sleep(1);
10      continue;
11
12     }
13     ret = ioctl(fd, KEY_IOCTL_GET_CHAR, &key)
14     if (key & 0xff) == '#' break; /* terminate */
15 } low bit留下來
16 close(fd);

```

→最後接 input a 字母

6 PXA I/O: LCD Control 2進I/O

- Header files: asm-arm/arch-pxa/lib/creator_pxa270_lcd.h
- Function: ioctl(fd, command, data)
 - Device Name: /dev/lcd
 - Data Structure

```

1 /* Data structure for writing char to LCD screen */
2 typedef struct lcd_write_info {
3     unsigned char Msg[512];          /* the array for saving input */
4     unsigned short Count;           /* the number of input char */
5     int CursorX, CursorY;         /* X, Y axis of cursor */
6 } lcd_write_info_t;
7
8 /* Data structure for writing a picture to LCD screen */
9 typedef struct lcd_full_image_info {
10    unsigned short data[0x800];      /* the array for saving picture */
11 } lcd_full_image_info_t;

```

- Command

```

1 /* Clear LCD data and move cursor back to the Upper-left corner */
2 #define LCD_IOCTL_CLEAR    LCD_IO( 0x0 )
3
4 /* Write char to LCD */ write
5 #define LCD_IOCTL_WRITE    LCD_IOW( 0x01, lcd_write_info_t )
6
7 /* Turn On or Off cursor */
8 #define LCD_IOCTL_CUR_ON   LCD_IO( 0x02 )
9 #define LCD_IOCTL_CUR_OFF  LCD_IO( 0x03 )
10
11 /* Get and Set the position (X, Y) of cursor */
12 #define LCD_IOCTL_CUR_GET  LCD_IOR( 0x04, lcd_write_info_t )
13 #define LCD_IOCTL_CUR_SET  LCD_IOW( 0x05, lcd_write_info_t )

```

寫程目前值

```

14
15 /* Write a picture to LCD */
16 #define LCD_IOCTL_DRAW_FULL_IMAGE LCD_IOW(0x06, lcd_full_image_info_t)

```

The following is the sample code of LCD control.

```

1 /*
2 * lcd.c -- The sample code to print "Hello World" on LCD screen.
3 */
4
5 #include <stdio.h>
6 #include <sys/fcntl.h>
7 #include <sys/ioctl.h>
8 #include <unistd.h>
9 #include "asm-arm/arch-pxa/lib/creator-pxa270_lcd.h"
10
11 int main()
12 {
13     int fd;
14     lcd_write_info_t display; /* struct for saving LCD data */
15
16     /* Open device /dev/lcd */
17     if ((fd = open("/dev/lcd", O_RDWR) < 0))
18     {
19         printf("open-/dev/lcd-error\n");
20         return (-1);
21     }
22
23     /* Clear LCD */
24     ioctl(fd, LCD_IOCTL_CLEAR, NULL);
25
26     /* Save output string to display data structure */
27     display.Count = sprintf((char *) display.Msg, "Hello-World\n");
28     /* Print out "Hello World" to LCD */
29     ioctl(fd, LCD_IOCTL_WRITE, &display);
30
31     /* Get the cursor position */
32     ioctl(fd, LCD_IOCTL_CUR_GET, &display);
33     printf("The cursor position is at (x,y) = (%d,%d)\n",
34           display.CursorX, display.CursorY);
35
36     close(fd);
37     return 0;
38 }

```

取回後會放在 CursorX
CursorY

7 PXA I/O: Audio Control

- Header files: asm-arm/arch-pxa/lib/creator-pxa270_lcd.h
- Function: ioctl(fd, command);

– Device Name: /dev/lcd

– Command

```
1 IOCTL_RECORD_START  
2 IOCTL_RECORD_STOP  
3 IOCTL_PLAY_START  
4 IOCTL_PLAY_STOP
```

The following is the sample code of Audio control.

```
1 #include <sys/ioctl.h>  
2 #include <stdio.h>  
3 #include <stdlib.h>  
4 #include "creator_s3c4510_codec.h" //must write the absolutely path  
5  
6 #define RECORDING_SIZE 8*8000 /* no of bytes for 4 seconds at 8000 per second */  
7 #define SIZE_16 4*8000 /* no of int for 4 seconds */  
8 main() {  
9     int codec_fd;  
10    int TotalReadSize, nRead, count, i;  
11    char AudioBuffer[RECORDING_SIZE];  
12    FILE *audio_fd;  
13  
14    codec_fd = open( "/dev/codec" , O_RDWR); /* Open the codec device driver */  
15    if (codec_fd < 0) {  
16        printf ( "Open-/dev/codec-error\n" );  
17        return (-1);  
18    }  
19  
20    if( ioctl (codec_fd , IOCTL_RECORD_START) < 0) { /*Start recording */  
21        printf ( "Audio_recording_start_error\n" );  
22        close(codec_fd);  
23        return (-1);  
24    }  
25    printf("Say_something_to_the_microphone\n");  
26    sleep(4); /* record 4 seconds of data */  
27  
28    if( ioctl (codec_fd , IOCTL_RECORD_STOP) < 0) { /* Stop recording */  
29        printf ( "Audio_recording_stop_error\n" );  
30        close(codec_fd);  
31        return (-1);  
32    }  
33    printf("Recording_stopped\n");  
34  
35    sleep(3); /* sleep for 3 seconds */  
36  
37    if( ioctl (codec_fd , IOCTL_PLAY_START) < 0) { /*Start playback */  
38        printf ( "Audio_playback_start_error\n" );  
39        close(codec_fd);  
40        return (-1);  
41    }  
42    printf("Start_playing_recorded_data_repeatedly\n");
```

插入代码前的 menu config 的操作

PXA 有固定地址的 buffer]
播放器 play 与 音频往那位置地址

```

43 | sleep(12); /* Sleep for 12 seconds to allow repeating twice */
44 |
45 | if( ioctl (codec_fd , IOCTL_PLAY_STOP) < 0) { /*Stop playback */
46 |     printf ( "Audio_playback_stop_error\n");
47 |     close(codec_fd );
48 |     return (-1);
49 |
50 | printf("Playback_stopped\n");
51 |
52 | /* Begin reading the data*/
53 | TotalReadSize = 0;
54 | count =RECORDING_SIZE;
55 | do {
56 |     if (count + TotalReadSize > RECORDING_SIZE)
57 |         count = RECORDING_SIZE- TotalReadSize ;
58 |
59 |     nRead = read(codec_fd , AudioBuffer+TotalReadSize , count);
60 |     if (nRead > 0 )
61 |         TotalReadSize += nRead;
62 |     else if (nRead == 0) /* EOF */
63 |         break;
64 |     else {
65 |         printf("Reading_audio_data_failed!\n");
66 |         close(codec_fd );
67 |         exit(1);
68 |     }
69 | } while (TotalReadSize < RECORDING_SIZE);
70 |
71 | /* write the audio data to a file */
72 | audio_fd = fopen("myaudio.txt" , "w");
73 | for( i = 0; i < SIZE_16; i = i+2) {
74 |     fprintf(audio_fd ,"%d\n" , (int) AudioBuffer[ i ]); /* convert into 16-bit 2's complement */
75 |
76 |
77 | fclose( audio_fd );
78 | close( codec_fd );
79 |

```

Lab 6-2: Socket Programming

Yu-Lun Huang

2019-07-16

1 Objective

- Understand basic socket programming and client-server paradigms.

2 Prerequisite

- Read man pages of `socket()`, `read()`, `write()`, etc.

3 Client-Server Model

The client–server programming model partitions a task between the service providers (servers) and service requesters (clients). Generally, clients initiate communication sessions by sending requests to a server which await incoming requests.

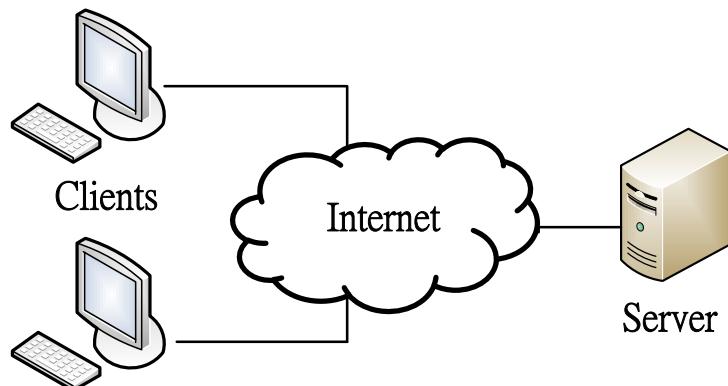


Figure 1: Client-Server Model

If clients and servers communicate over the network, then socket-based programming skills may be used to exchange messages in between. Figure 1 shows a schematic client-server interaction model.

4 Socket Fundamentals

Socket is an interface between an application process and transport layer. With the socket interface, the application process can send and receive messages to or from another application process running on a remote host. Similar to accessing an I/O device, a socket is accessible via a descriptor, as illustrated in Figure 2.

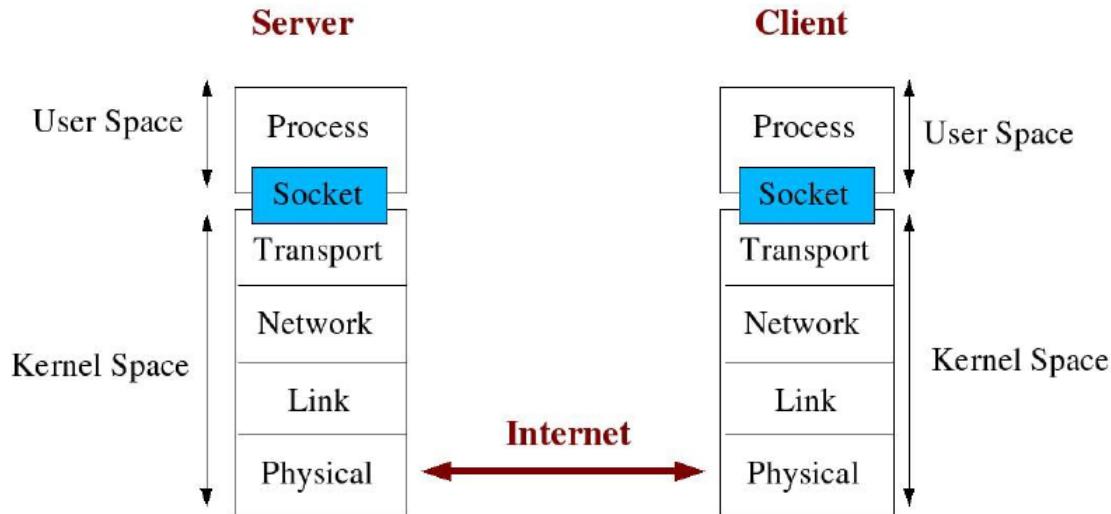


Figure 2: Socket Description.

There are several types of sockets, including internet socket, unix socket, X.25 socket, etc. Internet sockets can be further classified as STREAM socket and DATAGRAM socket, which is used for TCP protocol and UDP protocol. The STREAM socket, relying on TCP to provide reliable two-way communication, is used for connection-oriented protocols, as shown in Figure 3.

There are several important system calls for creating sockets for client/server programs:

- `socket()` creates an endpoint for communication and returns a descriptor.
- `connect()` initiates a connection on a socket. Normally used in a client program for sending requests.
- `bind()` binds a name to a socket. Normally, `listen()` is invoked after binding to socket. Used in a server program.
- `listen()` waits for connections on a socket. Used in a server program to wait for incoming requests.
- `accept()` accepts a connection on a socket. Used in a server program.

In a server program, you generally create a socket (using `socket()`, bind to it (`bind()`) and listen `listen()` for incoming requests. Once a request is initiated by a client, `accept()` is invoked to accept the connection.

To simplify the sample programs, the common APIs for creating different types of sockets are listed in 'sockop.c':

- `passivesock()` creates a socket and binds to the corresponding socket descriptor to passively accept the incoming connection requests.
- `connectsock()` creates a socket and initiates a connection to the remote host (specified in `struct sockaddr_in sin`).

Note that, we also create a new header file (`sockop.h` to include all header files required for socket programming. In `sockop.h`, `sys/socket.h` is necessary for those socket related APIs. In addition, the function prototypes of `passivesock()` and `connectsock()` are also defined in `sockop.h`.

```

1  /*
2   * sockop.h
3   */

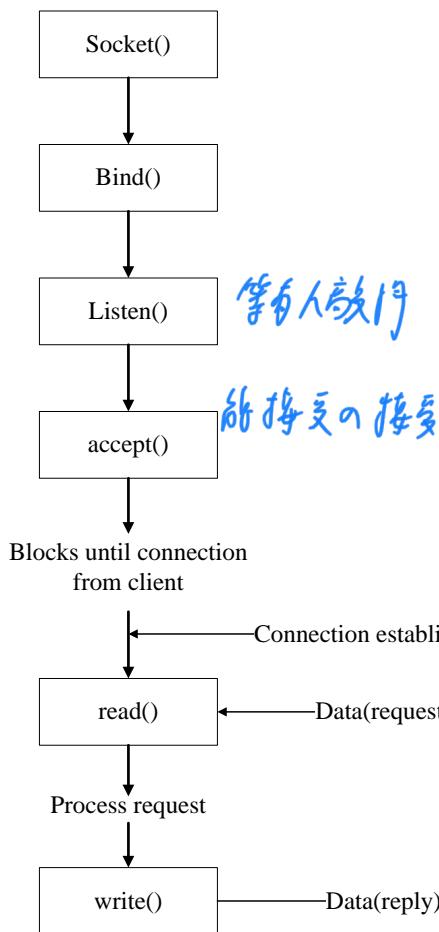
```

Socket.



for all

Server



Client

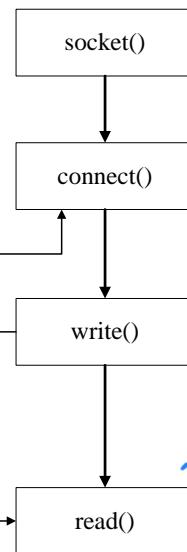


Figure 3: The design flow of Connection-Oriented Protocol.

```

4
5 #ifndef _SOCKOP_H_
6 #define _SOCKOP_H_
7
8 #include <arpa/inet.h>
9 #include <errno.h>
10 #include <netdb.h>
11 #include <netinet/in.h>
12 #include <stdio.h>
13 #include <stdlib.h>
14 #include <string.h>
15 #include <sys/socket.h>
16 #include <sys/types.h>
17 #include <sys/wait.h>
18
  
```

```

19 #define erexit(format,arg...) exit(printf(format,##arg))
20
21 /* Create server */
22 int passivesock(const char *service, const char *transport, int qlen);
23
24 /* Connect to server */
25 int connectsock(const char *host, const char *service, const char *transport);
26
27 #endif /* _SOCKOP_H_ */

```

```

1 /*
2  * sockop.c
3  */
4
5 #include "sockop.h"
6
7 /*
8  * passivesock - allocate & bind a server socket using TCP or UDP
9  *
10 * Arguments:
11 *   service - service associated with the desired port
12 *   transport - transport protocol to use ("tcp" or "udp")
13 *   qlen      - maximum server request queue length
14 */
15 int passivesock(const char *service, const char *transport, int qlen)
16 {
17     struct servent *pse; /* pointer to service information entry */
18     struct sockaddr_in sin; /* an Internet endpoint address */
19     int s, type; /* socket descriptor and socket type */
20
21     memset(&sin, 0, sizeof(sin));
22     sin.sin_family = AF_INET;
23     sin.sin_addr.s_addr = INADDR_ANY; server 諸君が来店可能接受: 不指定IP
24
25     /* Map service name to port number */
26     if ((pse = getservbyname(service, transport)))
27         sin.sin_port = htons((unsigned short)pse->s_port);
28     else if ((sin.sin_port = htons((unsigned short)atoi(service))) == 0)
29         erexit("Can't find \"%s\"-service-entry\n", service);
30
31     /* Use protocol to choose a socket type */
32     if (strcmp(transport, "udp") == 0)
33         type = SOCK_DGRAM;
34     else
35         type = SOCK_STREAM;
36
37     /* Allocate a socket */
38     s = socket(PF_INET, type, 0); 先去开一个socket, 然后看有没有空插位
39     if (s < 0)
40         erexit("Can't create socket: %s\n", strerror(errno));
41

```

插頭插到那插座

```
42 /* Bind the socket */
43 if (bind(s, (struct sockaddr *)&sin, sizeof(sin)) < 0)
44     erexit("Can't bind to port %s: %s\n",
45            service, strerror(errno));
46
47 /* Set the maximum number of waiting connection */
48 if (type == SOCK_STREAM && listen(s, qlen) < 0)
49     erexit("Can't listen on port %s: %s\n",
50            service, strerror(errno));
51
52 return s;
53 }
54
55 /*
56 * connectsock - allocate & connect a socket using TCP or UDP
57 *
58 * Arguments:
59 *   host      - name of host to which connection is desired
60 *   service    - service associated with the desired port
61 *   transport  - name of transport protocol to use ("tcp" or "udp")
62 */
63 int connectsock(const char *host, const char *service, const char *transport)
64 {
65     struct hostent *phe; /* pointer to host information entry */
66     struct servent *pse; /* pointer to service information entry */
67     struct sockaddr_in sin; /* an Internet endpoint address */
68     int s, type; /* socket descriptor and socket type */
69
70     memset(&sin, 0, sizeof(sin));
71     sin.sin_family = AF_INET;
72
73     /* Map service name to port number */
74     if ((pse = getservbyname(service, transport)))
75         sin.sin_port = pse->s_port;
76     else if ((sin.sin_port = htons((unsigned short)atoi(service))) == 0)
77         erexit("Can't get \"%s\" service entry\n", service);
78
79     /* Map host name to IP address, allowing for dotted decimal */
80     if ((phe = gethostbyname(host))) // 記住192.168.1.100, 這個IP address
81         memcpy(&sin.sin_addr, phe->h_addr, phe->h_length);
82     else if ((sin.sin_addr.s_addr = inet_addr(host)) == INADDR_NONE)
83         erexit("Can't get \"%s\" host entry\n", host);
84
85     /* Use protocol to choose a socket type */
86     if (strcmp(transport, "udp") == 0)
87         type = SOCK_DGRAM; // udp 無出沒回應, 快
88     else
89         type = SOCK_STREAM; // tcp 一定得有回應
90
91     /* Allocate a socket */
92     s = socket(PF_INET, type, 0); // 可放上面, 把屬性設好才插得上去
```

使用1000以下的port

8080 → http

```

93     if (s < 0)
94         erexit("Can't create socket: %s\n", strerror(errno));
95
96     /* Connect the socket */
97     if (connect(s, (struct sockaddr *)&sin, sizeof(sin)) < 0)
98         erexit("Can't connect to %s.%s: %s\n", host,
99                service, strerror(errno));
100
101    return s;
102 }
```

The following program shows the client side sample program (echoc.c). *client主子c檔*

```

1 /*
2  * echoc.c
3  */
4
5 #include <stdio.h>
6 #include <stdlib.h>
7 #include <string.h>
8 #include <unistd.h>
9
10 #include "sockop.h"
11
12 #define BUFSIZE 1024
13
14 int main(int argc, char *argv[])
15 {
16     int connfd; /* socket descriptor */
17     int n;
18     char buf[BUFSIZE];
19
20     if (argc != 4)
21         erexit("Usage: %s host_address host_port message\n", argv[0]);
22
23     /* create socket and connect to server */
24     connfd = connectsock(argv[1], argv[2], "tcp");
25
26     /* write message to server */ depend on什麼
27     if ((n = write(connfd, argv[3], strlen(argv[3]))) == -1)
28         erexit("Error: write()\n");
29
30     /* read message from the server and print */
31     memset(buf, 0, BUFSIZE);
32     if ((n = read(connfd, buf, BUFSIZE)) == -1)
33         erexit("Error: read()\n");
34     printf("%s\n", buf);
35
36     /* close client socket */ 不close會一直卡在那
37     close(connfd); 沒close就會新config就卡 up, down
38
39     return 0;
```

40 }

The server side sample program (echod.c)

passive server

```
/*
 * echod.c
 */
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <unistd.h>
#include "sockop.h"

#define BUFSIZE 1024

int main(int argc, char *argv[])
{
    int sockfd, connfd; /* socket descriptor */
    struct sockaddr_in addr_cln;
    socklen_t sLen = sizeof(addr_cln);
    int n;
    char snd[BUFSIZE], rcv[BUFSIZE];

    if (argc != 2)
        erexit("Usage: %s -port\n", argv[0]);

    /* create socket and bind socket to port */
    sockfd = passivesock(argv[1], "tcp", 10);

    while(1) { /* waiting for connection */
        connfd = accept(sockfd, (struct sockaddr *) &addr_cln, &sLen);
        if (connfd == -1)
            erexit("Error: -accept()\n");

        /* read message from client */
        if ((n = read(connfd, rcv, BUFSIZE)) == -1)
            erexit("Error: -read()\n");

        /* write message back to the client */
        n = sprintf(snd, "Server: %.*s", n, rcv);
        if ((n = write(connfd, snd, n)) == -1)
            erexit("Error: -write()\n");

        /* close client connection */
        close(connfd);
    }

    /* close server socket */
}
```

在第一次在data中傳輸時
OS會把資料傳回來暫存
之後的OS就會掉封包

先去給別人?

多國路卡依測到-
對已進來
OS發送 file server 有符合
會把對應的 server
發回 accept

```
49     close ( sockfd );
50
51     return  0;
52 }
```

The above 'echod.c' is the simplest server code that accepts one client request at a time. To concurrently accept multiple client requests, you can either use the '`select()`' system call or create multiple threads/processes to handle the incoming requests.

5 Reference Books

You can obtain more information about socket programming from the following books:

1. Douglas E. Comer, "Internetworking with TCP/IP Vol.1: Principles, Protocols, and Architecture (4th Edition)," Prentice Hall, 2000.
2. Douglas E. Comer and David L. Stevens, "Internetworking with TCP/IP Vol. II: ANSI C Version: Design, Implementation, and Internals (3rd Edition)," Prentice Hall, 1998.
3. Douglas E. Comer, David L. Stevens and Michael Evangelista, "Internetworking with TCP/IP, Vol. III: Client-Server Programming and Applications, Linux/Posix Sockets Version," Prentice Hall, 2000.

Lab 7 Linux Driver

Wen-Lin Sun

2018-10-19

In last lab, we used PXA270's LCD driver module to control the I/O on it. This time, similarly, we will use a driver to control the I/O, but the driver and the device are created by yourself on ZC702. There are several kinds of devices in Linux, such as block devices, character devices, pipe devices and socket devices. In this lab, we will create our own character devices step by step.

1 Environment Setup

Before we start to build our own devices, make sure you have done the environment setup in Lab2. The way we build up driver modules is just the same as the way we build up the boot image. So now, go to the directory where you build up the boot image (maybe `linux-xlnx`), create a directory to place your driver sources, and check the environment variables (`PATH` and `CROSS_COMPILE`) before you start this lab.

2 A Basic Driver

A simple driver can be implemented in two lines.

```
1 #include <linux/module.h>
2 MODULE_LICENSE("GPL");
```

And you can compile the source with the `Makefile` as following.

```
1 CC=arm-linux-gnueabihf-gcc
2
3 obj-m := mydev.o
4 系统内建的值
5 all:          目录下
6     make -C .. M=$(PWD) modules    所有的 object 都编译 module
7 clean:        可以直接，每不需要东西都删掉
8     make -C .. M=$(PWD) clean
```

Now, you can just build the driver with the `make` tool.

```
SHELL> make ARCH=arm
```

In this example, the source file is `mydev.c`, and the output file is `mydev.ko`.

3 Initial/Remove Module

We have created a driver module without any functionality so far. In this section, we are going to add two functions to define the actions that the driver will do on its initialization

and remove.

```
1 #include <linux/module.h>
2 #include <linux/init.h>
3 #include <linux/kernel.h>
4 MODULE_LICENSE("GPL");
5
6 static int my_init(void) {
7     printk("call_init\n");
8
9     return 0;
10}
11
12 static void my_exit(void) {
13     printk("call_exit\n");
14}
15 module_init(my_init);
16 module_exit(my_exit);
```

my_init is the function which will be called when we use insmod to load the driver. Relatively, my_exit is called when we use rmmod to unload the driver. To declare the relationship between the function and the action, the functions module_init and module_exit should be used. So now, after compiling the source, you can test the driver on your board. Try to load and unload it, and check if the behavior is just the same as you defined. By the way, the function printk is the way we use to print out the message in kernel space, instead of printf.

xx.ko

xx.init()

xx.cleanup()

insmod xx.ko & my

rmmod xx.ko

4 Define the Operations

To define the reactions of the actions which may be done on our device, such as, read, write, open, we should use the library linux/fs.h, which defines the basic operations use on file. (Everything is seen as file in Linux, including the devices.) The definition of the file operations in the linux/fs.h are showed as following.

```
1 struct file_operations {
2     struct module *owner;
3     loff_t (*llseek) (struct file *, loff_t, int);
4     ssize_t (*read) (struct file *, char __user *, size_t, loff_t *);
5     ssize_t (*write) (struct file *, const char __user *, size_t,
6                      loff_t *); ↳ I/O subsystem uniform API
7     ssize_t (*read_iter) (struct kiocb *, struct iov_iter *);
8     ssize_t (*write_iter) (struct kiocb *, struct iov_iter *);
9     int (*iterate) (struct file *, struct dir_context *);
10    unsigned int (*poll) (struct file *, struct poll_table_struct *);
11    long (*unlocked_ioctl) (struct file *, unsigned int, unsigned long
12                           );
13    long (*compat_ioctl) (struct file *, unsigned int, unsigned long);
14    int (*mmap) (struct file *, struct vm_area_struct *);
15    int (*open) (struct inode *, struct file *);
```

```

16     int (*fsync) (struct file *, loff_t, loff_t, int datasync);
17     int (*aio_fsync) (struct kiocb *, int datasync);
18     int (*fasync) (int, struct file *, int);
19     int (*lock) (struct file *, int, struct file_lock *);
20     ssize_t (*sendpage) (struct file *, struct page *, int, size_t,
21                         loff_t *, int);
22     unsigned long (*get_unmapped_area)(struct file *, unsigned long,
23                                      unsigned long, unsigned long, unsigned long);
24     int (*check_flags)(int);
25     int (*flock) (struct file *, int, struct file_lock *);
26     ssize_t (*splice_write)(struct pipe_inode_info *, struct file *,
27                            loff_t *, size_t, unsigned int);
28     ssize_t (*splice_read)(struct file *, loff_t *, struct
29                           pipe_inode_info *, size_t, unsigned int);
30     int (*setlease)(struct file *, long, struct file_lock **, void **)
31     ;
32     long (*fallocate)(struct file *file, int mode, loff_t offset,
33                        loff_t len);
34     void (*show_fdinfo)(struct seq_file *m, struct file *f);
35 #ifndef CONFIG_MMU
36     unsigned (*mmap_capabilities)(struct file *);
37 #endif
38     ssize_t (*copy_file_range)(struct file *, loff_t, struct file *,
39                               loff_t, size_t, unsigned int);
40     int (*clone_file_range)(struct file *, loff_t, struct file *,
41                            loff_t,
42                            u64);
43     ssize_t (*dedupe_file_range)(struct file *, u64, u64, struct file
44                                 *,
45                                 u64);
46 }

```

We don't have to define all the operations. In this example, we define read, write and open.

```

1 #include <linux/init.h>
2 #include <linux/module.h>
3 #include <linux/kernel.h>
4 #include <linux/fs.h>
5 MODULE_LICENSE("GPL");
6
7 // File Operations
8 static ssize_t my_read(struct file *fp, char *buf, size_t count, loff_t *
9                      fpos) {
10
11     printk("call\read\n");
12
13     return count;
14 }
15 static ssize_t my_write(struct file *fp, const char *buf, size_t count,
16                       loff_t *fpos) {

```

```

17     printk("call_write\n");
18
19     return count;
20 }
21
22 static int my_open(struct inode *inode, struct file *fp) {
23
24     printk("call_open\n");
25
26     return 0;
27 } 非uniform API
28
29 struct file_operations my_fops = {
30     read: my_read,
31     write: my_write,
32     open: my_open
33 };
34
35 #define MAJOR_NUM 244 → 用通常沒人用的 244, 245
36 #define DEVICE_NAME "my_dev"
37
38 static int my_init(void) {
39     printk("call_init\n");
40     if(register_chrdev(MAJOR_NUM, DEVICE_NAME, &my_fops) < 0) {
41         printk("Can't get major %d\n", MAJOR_NUM);
42         return (-EBUSY);
43     }
44
45     printk("My device is started and the major is %d\n", MAJOR_NUM);
46     return 0;
47 }
48
49 static void my_exit(void) {
50     unregister_chrdev(MAJOR_NUM, DEVICE_NAME);
51     printk("call_exit\n");
52 }
53
54
55 module_init(my_init);
56 module_exit(my_exit);

```

To declare the relationship between, for example, my_read and the read operation, we initialize a structure file_operations called my_fops. So now, we can register a character device with my_fops to the system when the driver is initialized. And don't forget to unregister the character device when the driver is removed.

To see your device under /dev, use should load your module first and get the major number. And then you can create the device node with the major number (in this example is 244).

SHELL> mknod /dev/mydev c 244 0 *→ 如果註冊後在 /dev 下還是沒有東西，可用以下*

The c in the command means character device.

So now, you can design and test your own device and its driver!

文件 " "

屬於 a device : 4 在 console 畫出東西

Lab 8: Signal and Timer

Yu-Lun Huang

2011-07-20

1 Objective

- Be familiar with signal, timer and process reaper.

2 Prerequisite

- Read man pages of `kill()`, `sigaction()`, `setitimer()`, etc.

3 Signal

send
sig
註冊/
安裝

Signals notify tasks of events that occurred during the execution of other tasks or Interrupt Service Routines (ISRs). It diverts the notified task from its normal execution path and triggers the associated signal handler. Signal handler is a function run in "asynchronous mode", which gets called when the task receives that signal, no matter which code the task is executed on. This is just like the relationship between hardware interrupts and ISRs, which are also asynchronous to the execution of OS and do not occur at any predetermined point of time. The difference between a signal and a normal interrupt is that signals are software interrupts, which are generated by other task or OS to trap normal execution of task. Another difference is that task can only receive a signal when it is running in user mode. If the receiving process is running in kernel mode, the execution of the signal will start only after the process returns to user mode. When the signal handler completes, the normal execution resumes. The task continues execution from wherever it happened to be before the signal was received.

Signals have been used for approximately 30 years without any major modifications. Although we now have advanced synchronization tools and many IPC mechanisms, signals play a vital role in Linux for handling exceptions and interrupts. For example, when child process exits its execution, it sends a `SIGCHLD` signal to parent process and becomes a zombie process. When the parent process catches this signal, it performs `wait()` or `waitpid()` to reclaim the resources used by child process in the signal handler. This design prevents the parent process from blocking in `wait()` or `waitpid()` function calls. (We detail this implementation in later section.)

3.1 Sending Signals

Linux supports various types of signal, you should specify the desired type when sending signal to other tasks. Each signal in Linux is identified by an integer value, which is called signal number or vector number. The first 31 signals are standard signals, some of which date back to 1970s UNIX from Bell Labs. The POSIX (Portable Operating Systems and Interface for UNIX) standard introduced a new class of signals designated as real-time signals, with numbers ranging from 32 to 63. Usually, all signal numbers as well as the associated symbolic name are defined in `/usr/include/signal.h`, or you can use the command '`kill -1`' to see a list of signals supported by your Linux system.

man kill
sigaction

ps: to see currently process running
ipcs: to see shm, message queue

Signals can be generated from within the shell using the kill command. (Man the command "kill" to obtain detailed information.) The name of kill may seem strange, but actually most signals serve the purpose of terminating processes, so this is not really that unusual.

For example, the following command sends the SIGUSR1 signal to process 3423.

SHELL> kill -USR1 3423
① terminate SIGKILL
X的信号
且SIGKILL

Another method is to use the kill system call within a C program to send a signal to a process. This call takes a process ID and a signal number as parameters.

```
int kill(pid_t pid, int sig_no);
```

15 9

One common use of this mechanism is to end another process by sending it a SIGTERM or SIGKILL signal. Another common use is to send a command to a running program. Two "userdefined" signals are reserved for this purpose: SIGUSR1 and SIGUSR2. The SIGHUP signal is sometimes used for this purpose as well, commonly to wake up an idling program or cause a program to reread its configuration files.

3.2 Catching Signals

When a process receives a signal, it may do one of several things, depending on the signal's disposition. For each signal, there is a default disposition, which determines what happens to the task if the program does not specify any signal handler. If a signal handler is used, the currently executing task is paused, the signal handler is executed; and when the signal handler returns, the task resumes.

For most signal types, a program can specify some other behavior – either to ignore the signal or to call a special signal handler function to respond to the signal. We can use sigaction system call to register the signal handler or set a signal disposition.

The syntax of sigaction is:

```
int sigaction (int signum, const struct sigaction *act, struct sigaction *oldact)
```

The first argument, signum, is a specified signal. The next two parameters are pointers to sigaction structures; the second argument, is used to set the new disposition of the signal signum; and the third argument is used to store the previous disposition, usually NULL.

The sigaction structure is defined as:

```
struct sigaction {  
    void (*sa_handler)(int);  
    void (*sa_sigaction)(int, siginfo_t *, void *);  
    sigset_t sa_mask;  
    int sa_flags;  
}
```

The members of the sigaction structure are described as follows.

- **sa_handler:**

sa_handler is a pointer pointed to a user-defined signal handler or default signal handler:

- SIG_DFL, which specifies the default disposition for the signal.
- SIG_IGN, which specifies that the signal should be ignored.
- A pointer to a signal-handler function. This function receives the signal number as its only argument

- **sa_mask:**

sa_mask gives a mask of signals which should be blocked during execution of the signal handler. The **sigset_t** structure consists an array of unsigned long, each bit inside controls the block state of a particular signal type. In the first unsigned long in **sigset_t**, the least significant bit (0) is unused since no signal has the number 0, the other 31 bits represents the 31 standard UNIX signals. The bits in the second unsigned long are the real-time signal numbers from 32 to 64.

- **sa_flags:**

sa_flags specifies the action of signal. Sets of flags are available for controlling the signal in a different manner. More than one flag can be used by OR operation:

- **SA_NOCLDWAIT:**
If signum is SIGCHLD, do not transform child processes into zombies when they terminate.
- **SA_RESETHAND:**
SA_RESETHAND restores the default action of the signal after the user-defined signal handler has been executed.
- **SA_NODEFER:**
The signal which triggered the handler will be blocked. Set the **SA_NODEFER** allows signal can be received from within its own signal handler.
- **SA_SIGINFO:**
When **SA_SIGINFO** is set, the **sa_sigaction** should be used, instead of specifying the signal handler in **sa_handler**.

For more flags information, please refer the **sigaction** manpages.

- **sa_sigaction:**

If the **SA_SIGINFO** flag is set in **sa_flags**, **sa_sigaction** should be used. **sa_sigaction** is a pointer to a function that takes three arguments, for example:

```
void my_handler (int signo, siginfo_t *info, void *context);
```

Here, **signo** is the signal number, and **info** is a pointer to the structure of type **siginfo_t**, which specifies the signal-related information; and **context** is a pointer to an object of type **ucontext_t**, which refers to the receiving process context that was interrupted with the delivered signal. For the detail structure of **siginfo_t** and **ucontext_t**, please refer the manpage of **sigaction** and **getcontext**.

The following example uses **SA_SIGINFO** and **sa_sigaction** to extract information from a signal. Use the **kill** command to send a **SIGUSR1** signal to **catch** program.

```

1  /*
2   * sig_catch.c
3   */
4
5 #include <signal.h>
6 #include <stdio.h>
7 #include <string.h>
8 #include <sys/types.h>
9 #include <unistd.h>
10
11 void handler (int signo, siginfo_t *info, void *context)
12 {
13     /* show the process ID sent signal */
```

usage :

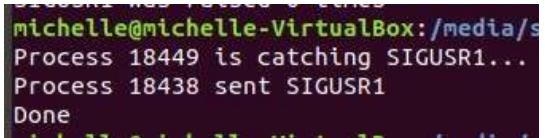
3

terminal ① : ./sig-catch

terminal ② : kill -10 process pid (or SIGUSR1)

```

14     printf ("Process -(%d) -sent -SIGUSR1.n" , info->si_pid );
15 }
16
17 int main (int argc , char *argv [])
18 {
19     struct sigaction my_action;
20
21     /* register handler to SIGUSR1 */
22     memset(&my_action , 0 , sizeof (struct sigaction)); 设置信号处理函数
23     my_action.sa_flags = SA_SIGINFO; 设置信号属性
24     my_action.sa_sigaction = handler; 设置信号处理函数名
25
26     sigaction(SIGUSR1 , &my_action , NULL); 注册信号
27
28     printf ("Process -(%d) -is -catching -SIGUSR1... \n" , getpid ());
29     sleep(10);
30     printf ("Done.\n");
31
32     return 0;
33 }
```



Remember that some OSs implement `sa_handler` and `sa_sigaction` as a union, do not assign values to these two arguments at same time.

3.3 Atomic Access in Signal Handler

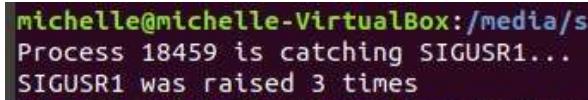
Assigning a value to a global variable can be dangerous because the assignment may actually be carried out in two or more machine instructions, and a second signal may occur between them, leaving the variable in a corrupted state. If you use a global variable to flag a signal from a signal handler function, it should be of the special type `sig_atomic_t`. In Linux, `sig_atomic_t` is an ordinary integer data type, but which one it is, and how many bits it contains, may vary from machine to machine. In practice, you can also use `sig_atomic_t` as a pointer. If you want to write a program which is portable to any standard UNIX system, though, use `sig_atomic_t` for these global variables. Reading and writing `sig_atomic_t` is guaranteed to happen in a single instruction, so there's no way for a handler to run "in the middle" of an access.

The following program listing uses a signal-handler function to count the number of times that the program receives SIGUSR1, one of the signals reserved for application use.

```

1 /*
2  * sig_count.c
3  */
4
5 #include <signal.h>
6 #include <stdio.h>
7 #include <string.h>
8 #include <sys/types.h>
9 #include <time.h> to use the timespec structure, nanosleep
10 #include <unistd.h>
11
12 sig_atomic_t sigusr1_count = 0;
13
14 void handler (int signal_number)
```

*usage: 使用 -g -f 运行
kill -10 process pid kill 3 次
输出：SIGUSR1 was raised 3 times*



```

15 {
16     ++sigusr1_count; /* add one, protected atomic action */
17 }
18
19 int main ()
20 {
21     struct sigaction sa;
22     struct timespec req; it allows the user to store time in seconds & nanoseconds
23     int retval;
24
25     /* set the sleep time to 10 sec */
26     memset(&req, 0, sizeof(struct timespec));
27     req.tv_sec = 10; Store time in seconds
28     req.tv_nsec = 0; Store time in nanoseconds
29
30     /* register handler to SIGUSR1 */
31     memset(&sa, 0, sizeof(sa));
32     sa.sa_handler = handler;
33
34     sigaction(SIGUSR1, &sa, NULL);
35
36     printf("Process -(%d)- is -catching -SIGUSR1-...\n", getpid());
37
38     /* sleep 10 sec */
39     do{ high resolution (of time) sleep
40         retval = nanosleep(&req, &req); ① On & off sleep for the requested interval.
41     } while(retval); return 0
42
43     printf("SIGUSR1-was-raised-%d-times\n", sigusr1_count);
44
45     return 0;
46 }

```

② interrupted by signal handler or encounters an error:

return -1

4 Timer

The timer schedules an event according to a predefined time value in the future. When the timer expired, it delivers a signal to the calling process. Timers are used everywhere in Unix-like systems, from basic delay function implementation to network transmission and performance monitoring. Linux provides two basic POSIX standard timer functions, `alarm` and `setitimer`. The `alarm` system call arranges an alarm clock for delivering a `SIGALRM` signal after the specified time elapsed. The `setitimer` system call is a generalization of the `alarm` call, it provides three different types of timer for counting the elapsed time. The syntax of `setitimer` is shown as following:

```
int setitimer(int which, const struct itimerval *new_val, struct itimerval *old_val);
```

The first argument `which` is the timer code, specifying which timer to set.

- If the timer code is `ITIMER_REAL`, the process is sent a `SIGALRM` signal after the specified wall-clock time has elapsed.
- If the timer code is `ITIMER_VIRTUAL`, the process is sent a `SIGVTALRM` signal after the process has

soft timer

ITIMER_REAL : t₁+t₂+t₃+t₄+t₅

task
t₁

process 在運作的時候才有
執行 (system call)

ITIMER_VIRTUAL : t₁+t₃+t₅ 這就是
I/O 等待時間

I/O t₂

ITIMER_PROF : t₁+t₃+t₄+t₅

I/O & 與 CPU

用不用 soft-time pfi claim a 5 sec ts
 不用, 只有 ITIMER_REAL a 5 sec 是
 在 5 secs, 其他都會 ≥ 5 secs

executed for the specified time. Time in which the process is not executing (that is, when the kernel or another process is running) is not counted.

- If the timer code is **ITIMER_PROF**, the process is sent a **SIGPROF** signal when the specified time has elapsed either during the process's own execution or the execution of a system call on behalf of the process.

process 身子 + 系統幫忙執行的時間 (system call)

The second argument is a pointer to a **itimerval** structure specifying the new settings for that timer. The third argument, if not null, is a pointer to another **itimerval** structure which receives the old timer settings. The struct **itimerval** variable has two fields:

- it_value** is a **struct timeval** field that contains the time until the timer next expires and a signal is sent. If this is 0, the timer is disabled.
- it_interval** is another **struct timeval** field containing the value to which the timer will be reset after it expires. If this is 0, the timer will be disabled after it expires. If this is nonzero, the timer is set to expire repeatedly after this interval.

And the struct **timeval** has the following two members:

- tv_sec** represents the number of whole seconds of elapsed time.
- tv_usec** is the rest of the elapsed time (a fraction of a second), represented as the number of microseconds. It is always less than one million.

The following program illustrates the use of **setitimer** to track the execution time of a program. A timer is configured to be expired every 250 milliseconds and send a **SIGVTALRM** signal.

```

1  /*
2   * timer.c
3   */
4
5 #include <signal.h>
6 #include <stdio.h>
7 #include <string.h>
8 #include <sys/time.h>
9
10 void timer_handler (int signum)
11 {
12     static int count = 0;
13
14     printf ("timer_expired %d times\n", ++count);
15 }
16
17 int main (int argc, char **argv)
18 {
19     struct sigaction sa;
20     struct itimerval timer;
21
22     /* Install timer_handler as the signal handler for SIGVTALRM */
23     memset (&sa, 0, sizeof (sa));
24     sa.sa_handler = &timer_handler;
25     sigaction (SIGVTALRM, &sa, NULL);
26
27     /* Configure the timer to expire after 250 msec */

```

```

michelle@michelle-VirtualBox: ~
timer_expired 1 times
timer_expired 2 times
timer_expired 3 times
timer_expired 4 times
timer_expired 5 times
timer_expired 6 times
timer_expired 7 times
timer_expired 8 times
timer_expired 9 times
timer_expired 10 times
timer_expired 11 times
timer_expired 12 times
timer_expired 13 times
timer_expired 14 times
timer_expired 15 times
timer_expired 16 times

```

它會將 **it_value** 設到 0 時，產生一個錯誤，
 並將 **it_value** 的值設為 **it_interval** 的值，然後重新開始計時，如此往復

SIGVTALRM \rightarrow 睡眠
 16 ITIMER_VIRTUAL \rightarrow 200ms
 第一次启动时 200ms

```

28 16 timer.it_value.tv_sec = 0;
29 16 timer.it_value.tv_usec = 250000;
30 16 /* Reset the timer back to 250 msec after expired */
31 16 timer.it_interval.tv_sec = 0;
32 16 timer.it_interval.tv_usec = 250000;
33
34
35 /* Start a virtual timer */
36 setitimer(ITIMER_VIRTUAL, &timer, NULL); 放射
37
38 /* Do busy work */
39 while (1); 用
40
41 return 0; Sgaction
42 } (. . . . .)
    
```

- why RA sleep 做?
- ① embedded device dependent
(有的系统不叫 sleep \rightarrow VS)
 - ② 放射 y 值的子函数固定
可以用 sleep, 不建议使用
睡眠请求 \rightarrow signal

OS辐射导致好
ITIMER_VIRTUAL
交叉
SIGVTALRM

The following program demonstrates the difference among three different timers. Write a simple “while(1) loop” program first and execute several instances of this program as the number of cores in your machine. For example, if you have a dual core machine, then run two “while(1) loop” programs simultaneously. After occupy all computing resources, execute the `timer_diff.c` program and observe the results of three counters.

```

1 /*
2 * timer_diff.c
3 */
4
5 #include <fcntl.h>
6 #include <signal.h>
7 #include <stdio.h>
8 #include <stdlib.h>
9 #include <string.h>
10 #include <sys/stat.h>
11 #include <sys/time.h>
12 #include <sys/types.h>
13 #include <unistd.h>
14
15 /* counter */
16 int SIGALRM_count = 0;
17 int SIGVTALRM_count = 0;
18 int SIGPROF_count = 0;
19
20 /* handler of SIGALRM */
21 void SIGALRM_handler (int signum)
22 {
23     SIGALRM_count++;
24 }
25
26 /* handler of SIGVTALRM */
27 void SIGVTALRM_handler (int signum)
28 {
29     SIGVTALRM_count++;
30 }
31
    
```

```

32 /* handler of SIGPROF */
33 void SIGPROF_handler (int signum)
34 {
35     SIGPROF_count++;
36 }
37
38 void IO_WORKS();
39
40 int main (int argc, char **argv)
41 {
42     struct sigaction SA_SIGNALRM, SA_SIGVTALRM, SA_SIGPROF;
43     struct itimerval timer;
44
45     /* Install SIGNALRM_handler as the signal handler for SIGNALRM */
46     memset (&SA_SIGNALRM, 0, sizeof (SA_SIGNALRM));
47     SA_SIGNALRM.sa_handler = &SIGALRM_handler;
48     sigaction (SIGALRM, &SA_SIGNALRM, NULL);
49
50     /* Install SIGVTALRM_handler as the signal handler for SIGVTALRM */
51     memset (&SA_SIGVTALRM, 0, sizeof (SA_SIGVTALRM));
52     SA_SIGVTALRM.sa_handler = &SIGVTALRM_handler;
53     sigaction (SIGVTALRM, &SA_SIGVTALRM, NULL);
54
55     /* Install SIGPROF_handler as the signal handler for SIGPROF */
56     memset (&SA_SIGPROF, 0, sizeof (SA_SIGPROF));
57     SA_SIGPROF.sa_handler = &SIGPROF_handler;
58     sigaction (SIGPROF, &SA_SIGPROF, NULL);
59
60     /* Configure the timer to expire after 100 msec */
61     timer.it_value.tv_sec = 0;
62     timer.it_value.tv_usec = 100000;
63
64     /* Reset the timer back to 100 msec after expired */
65     timer.it_interval.tv_sec = 0;
66     timer.it_interval.tv_usec = 100000;
67
68     /* Start timer */
69     setitimer(ITIMER_REAL, &timer, NULL);
70     setitimer(ITIMER_VIRTUAL, &timer, NULL);
71     setitimer(ITIMER_PROF, &timer, NULL);
72
73     /* Do some I/O operations */
74     IO_WORKS();
75
76     printf ("SIGALRM_count===%d\n", SIGALRM_count);
77     printf ("SIGVTALRM_count===%d\n", SIGVTALRM_count);
78     printf ("SIGPROF_count===%d\n", SIGPROF_count);
79
80     return 0;
81 }
82

```

```

michelle@michelle-VirtualBox:~/Documents$ ./sigcount
SIGALRM_count = 148
SIGVTALRM_count = 2
SIGPROF_count = 22
michelle@michelle-VirtualBox:~/Documents$
```

```

83 void IO_WORKS()
84 {
85     int fd, ret;
86     char buffer[100];
87     int i;
88
89     /* Open/Read/Close file 300000 times */
90     for (i = 0; i < 300000; i++) {
91         if ((fd = open("/etc/init.d/networking", O_RDONLY)) < 0) {
92             perror("Open /etc/init.d/networking");
93             exit(EXIT_FAILURE);
94         }
95
96         do {
97             ret = read(fd, buffer, 100);
98         } while (ret);
99
100        close(fd);
101    }
102 }
```

W.R.T.

5 Process Reaper

When a child process completes its execution, rather than reclaims all memory resources associated with it, the OS puts this process into a zombie state and allows the parent process read the exit status of child process. In Lab 3, we let the parent process executes the `wait` and `waitpid` system call to ‘reap’ (remove) the zombie process. But executing `wait` and `waitpid` blocks the normal execution of parent process. Even configures the `waitpid` with `WNOHANG` option, the parent should perform periodic checking on status of child processes.

One solution is that we can allow parent process explicitly ignore `SIGCHLD` by setting its handler to `SIG_IGN` (rather than simply ignoring the signal by default) or has the `SA_NOCLDWAIT` flag set, all child exit status information will be discarded and no zombie processes will be left. But sometimes, we need to perform specific operation when child process end its execution, e.g. check the exit status of child processes. In such case, we register the signal handler with `SIGCHLD` to reap the child process manually. This signal handler is also known as the **process reaper**.

The following program illustrates the basic operation of **process reaper**. This program uses another system call `sigaction` to register the `reaper` function for `SIGCHLD`. (Please refer the manpage to see the difference between `sigaction` and `sigsetSIG`.) When the parent process traps into the `reaper` function, it performs a nonblocking `waitpid` call to reclaim zombie processes. The `reaper` function checks the return value of `waitpid` to determine any zombie process to reap. If the return value is equal to zero, which indicates no zombie exists, we exit the handler to resume normal execution.

```

1 /*
2  * reaper.c
3  * Demonstrate the work of process reaper
4  */
5
6 #include <signal.h>
7 #include <stdio.h>
8 #include <stdlib.h>
9 #include <sys/types.h>
```

```

10 #include <sys/wait.h>
11 #include <time.h>
12 #include <unistd.h>
13
14 #define FORKCHILD 5
15
16 volatile sig_atomic_t reaper_count = 0;
17
18 /* signal handler for SIGCHLD */
19 void Reaper(int sig)
20 {
21     pid_t pid;
22     int status;
23
24     while((pid = waitpid(-1, &status, WNOHANG)) > 0)
25     {
26         printf("Child %d is terminated.\n", pid);
27         reaper_count++;
28     }
29 }
30
31 void ChildProcess()
32 {
33     int rand;
34
35     /* rand a sleep time */
36     srand(time(NULL));
37     rand = random() % 10 + 2;
38
39     printf("Child %d sleep %d sec.\n", getpid(), rand);
40     sleep(rand);
41     printf("Child %d exit.\n", getpid());
42
43     exit(0);
44 }
45
46 int main(int argc, char *argv[])
47 {
48     int cpid;
49     int i;
50
51     /* regist signal handler */
52     signal(SIGCHLD, Reaper);
53
54
55     /* fork child processes */
56     for (i = 0; i < FORKCHILD; i++)
57     {
58         if ((cpid = fork()) > 0) /* parent */
59             printf("Parent fork child process %d.\n", cpid);
60         else /* child */

```

```

mitchette@mitchette-VirtualBox:~/mcu
Parent fork child process 18771.
Child 18771 sleep 6 secs.
Parent fork child process 18772.
Child 18772 sleep 5 secs.
Parent fork child process 18773.
Child 18773 sleep 10 secs.
Parent fork child process 18774.
Child 18774 sleep 10 secs.
Parent fork child process 18775.
Child 18775 sleep 4 secs.
Child 18771 exit.
Child 18771 is terminated.
Child 18772 exit.
Child 18772 is terminated.
Child 18775 exit.
Child 18775 is terminated.
Child 18773 exit.
Child 18773 is terminated.
Child 18774 exit.
Child 18774 is terminated.

```

```
61     ChildProcess ();
62     sleep (1);
63 }
64
65 /* wait all child exit */
66 while(reaper_count != FORKCHILD)
67     sleep (1);
68
69 return 0;
70
71 }
```