

Data Structures

Self-balancing

binary search tree

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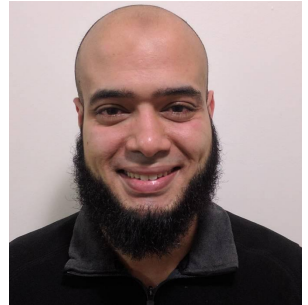
Teaching, Training and Coaching since more than a decade!

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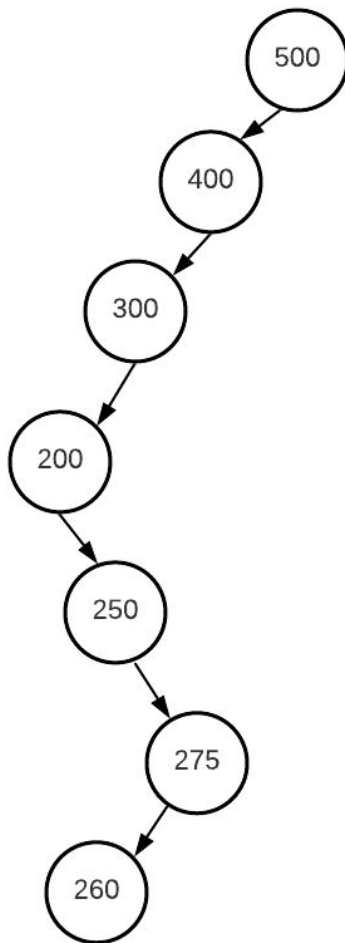
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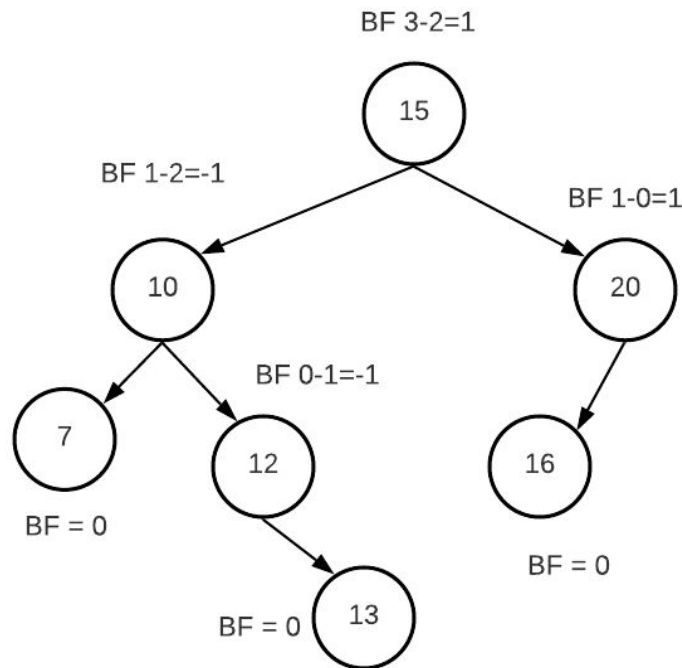
Recall: Degenerate BST

- Each node has 1 child
- From performance perspective, it is like a linkedlist, making several operations $O(n)$
- In general, trees with height much far from $\log(n)$ are not efficient!



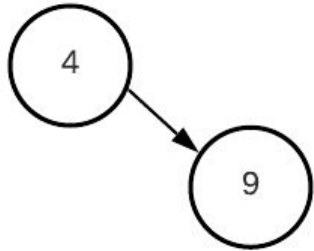
Recall: Balanced BST

- **Difference** between **heights** of left subtree and right subtree is not more than 1.
 - $\text{Height}(\text{left}) - \text{Height}(\text{right})$
 - **For visualization**, assume $\text{height}(\text{leaf}) = 1$
- Let's compute this difference for each node (**Balance factor**)
 - -7 means = right height greater with 7
 - -1, 0, 1 \Rightarrow BBST
 - $|\text{bf}| > 1 \Rightarrow$ imbalanced tree



Consider height and balance_factor()

- In last line of insert() function we call: `update_height()`



```
6 class AVLTree {
7 private:
8     int data { };
9     int height { };
10    AVLTree* left { };
11    AVLTree* right { };
12
13    int ch_height(AVLTree* node) { // child height
14        if (!node)
15            return -1; // -1 for null
16        return node->height; // 0 for leaf
17    }
18    void update_height() { // call in end of insert function
19        height = 1 + max(ch_height(left), ch_height(right));
20    }
21    int balance_factor() {
22        return ch_height(left) - ch_height(right);
23    }
24 }
```

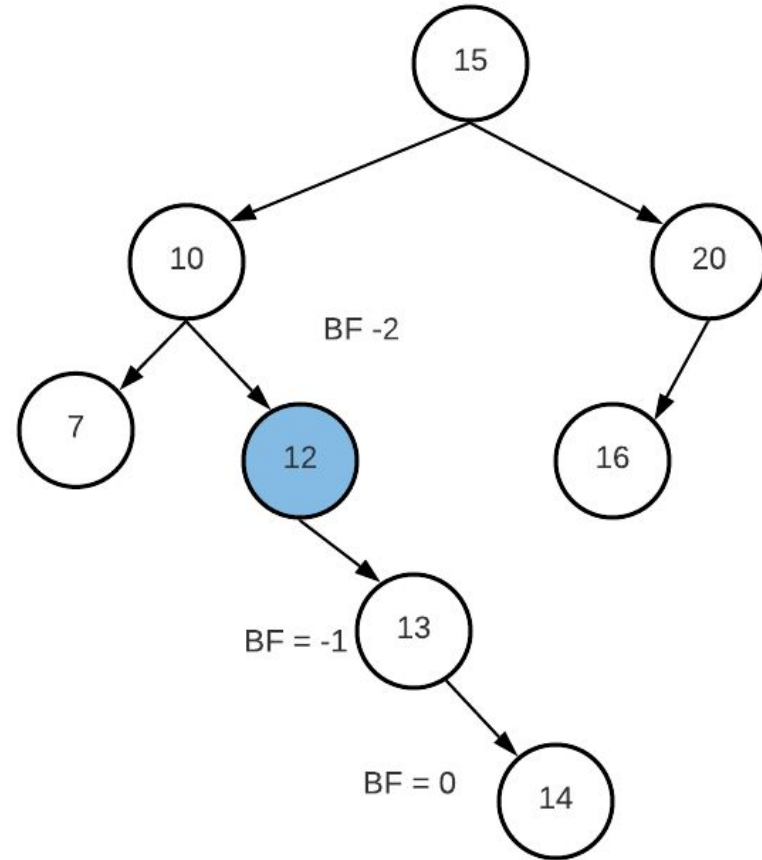
```
12 class AVLTree {
13 private:
14     struct BinaryNode {
15         int data { };
16         int height { };
17         BinaryNode* left { };
18         BinaryNode* right { };
19
20     BinaryNode(int data) :
21         data(data) {
22     }
23
24     int ch_height(BinaryNode* node) {    // child height
25         if (!node)
26             return -1;                // -1 for null
27         return node->height;          // 0 for leaf
28     }
29     int update_height() {    // call in end of insert function
30         return height = 1 + max(ch_height(left), ch_height(right));
31     }
32     int balance_factor() {
33         return ch_height(left) - ch_height(right);
34     }
35 };
36
```

Maintaining BBST

- Insertion in BST depends on input order \Rightarrow will generate imbalance trees
- **Self-balancing** BST trees follow **change and fix** approach to keep it BBST
 - E.g. Insert a new element in BST
 - Is imbalance BST?
 - Yes \Rightarrow **Fix** the tree to have again $|BF| \leq 1$
- There are several such trees that maintains tree as BBST
 - **AVL Trees**: one of oldest and simplest ways
 - Red-Black Trees, Splay Trees, Treaps
- This section is focusing on AVL Trees

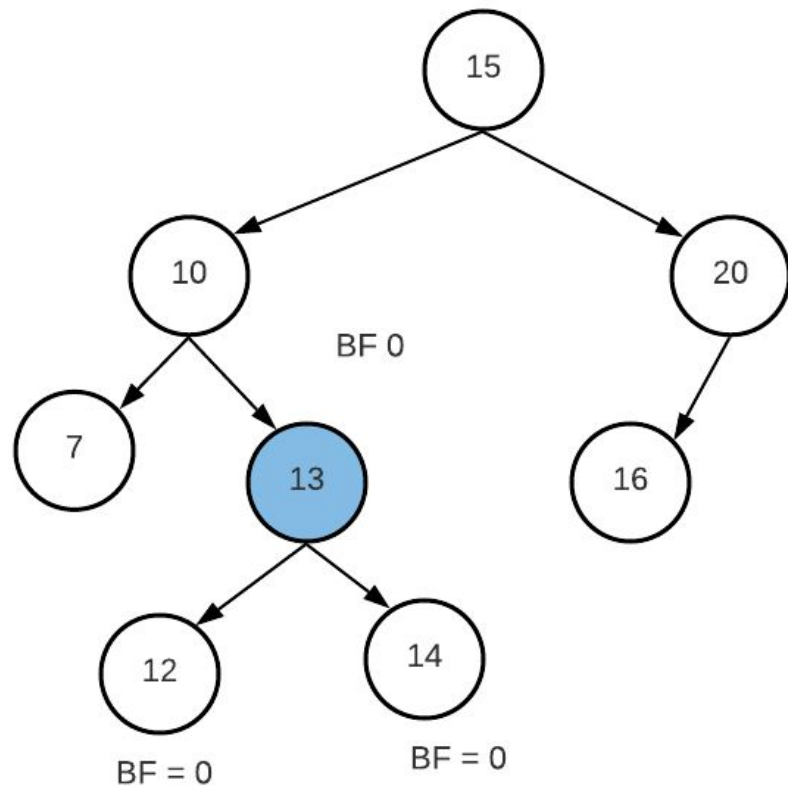
Change (insert 14)

- Let's insert 14
- Now, recompute Balance Factor
- Recursively we reach leaf nodes then go back to parents and so on
- Once you detect $|BF| > 1$, this subtree is not BST
- Think for 2 min how to restructure subtree (12) to make it balanced BST?



Fix (Tree Rotation)

- If we pushed node 13 **up** and node 12 **down** left, this subtree is fixed
 - Observe, we did not change other subtrees
 - Observe, tree remains BBST
- This kind of systematic change is called **tree rotation**
- If after insertion, in **bottom up** style we kept fixing corrupted sub-trees \Rightarrow BBST



AVL [Demo](#)

AVL Tree

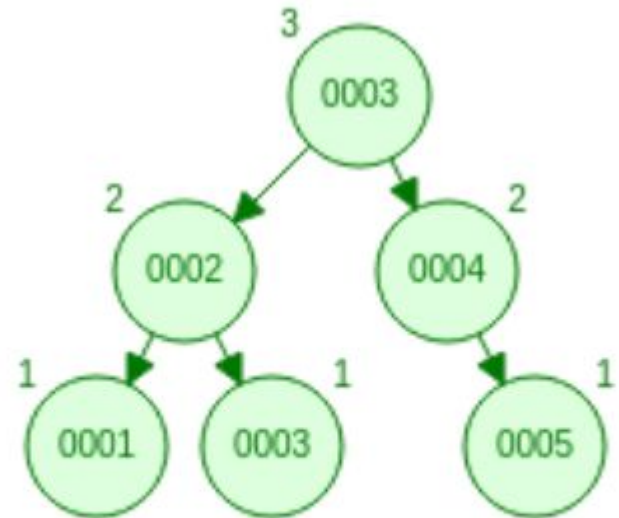
Insert

Delete

Find

Print

- Many web sites provide online demo for insertion/deletion in AVL
- Try to insert almost sorted numbers and watch the tree **restructuring**



“Acquire knowledge and impart it to the people.”

“Seek knowledge from the Cradle to the Grave.”