Data Structures AVL Insertion

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Insertion

- To insert, we insert normally
- But, we need to fix any corruption (|BF| > 1) immediately
- To do so, in end of function, we update height then call balance which will check if |BF| > 1 or not
- We need also to update left & right as they might be changed
- We will follow the last homework code (rewritten BST with extra node struct)

Insertion: before

- This is our insertion (before change)
- Observe: it assumes tree nodes never rebalanced (changed)
- We need to balance
 And
 Make code flexible

```
void insert node(int target, BinaryNode* node) {
    if (target < node->data) {
        if (!node->left)
            node->left = new BinaryNode(target);
        else
            insert node(target, node->left);
    } else if (target > node->data) {
        if (!node->right)
            node->right = new BinaryNode(target);
        else
            insert node(target, node->right);
```

Insertion: after

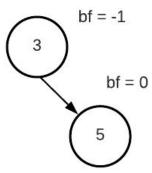
```
BinaryNode* insert node(int target, BinaryNode* node) {
    if (target < node->data) {
        if (!node->left)
            node->left = new BinaryNode(target);
        else // change left. update left as it might be balanced
            node->left = insert node(target, node->left);
    } else if (target > node->data) {
        if (!node->right)
            node->right = new BinaryNode(target);
        else
            node->right = insert node(target, node->right);
    node->update height();
    return balance(node);
```

Insertion: caller

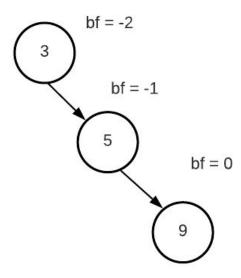
• The root can be changed from the caller, so consider that

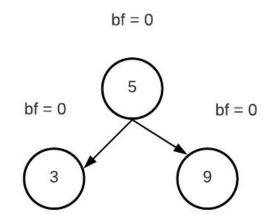
```
void insert_value(int target) {
   if (!root)
      root = new BinaryNode(target);
   else
      root = insert_node(target, root);
}
```

- Let's insert the first 2 values
- No problems so far

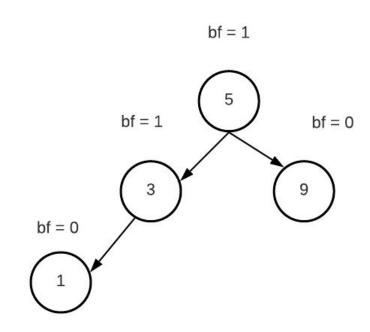


- With 9, node(3) is imbalance
- Right-Right case
- Do Left rotation at 3

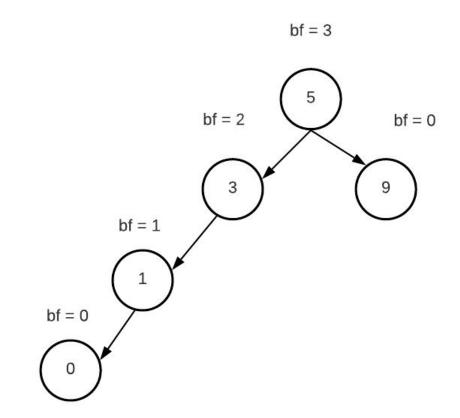


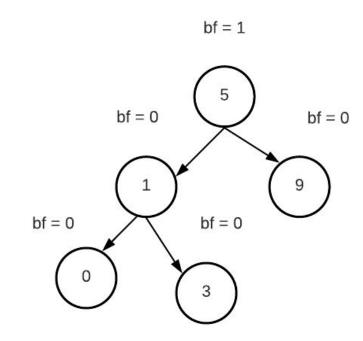


Still good with 1

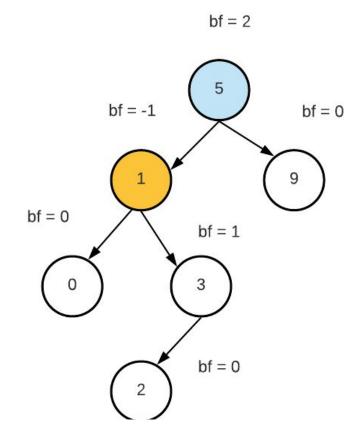


- With 0, node(3) is imbalance
- Left-Left case
- Do right rotation at 3

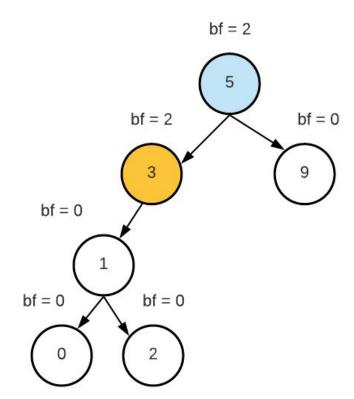


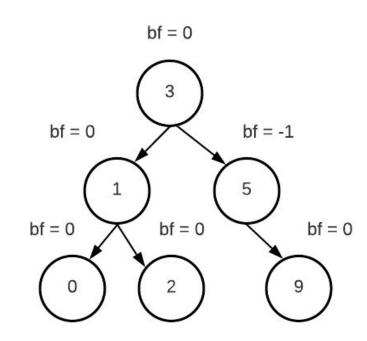


- With 2, node(5) is imbalance
- Left-Right case
 - o Bf: 2 -1
- First perform left_rotation(1)
 - o This pushes 1 down and 3 up
 - A=0, B=2, C=Null
 - So B's parent will change from 3 to 1
- Then perform right_rotation(5)

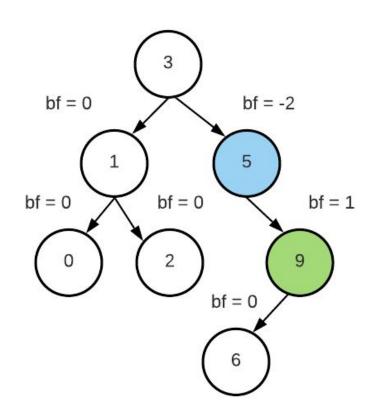


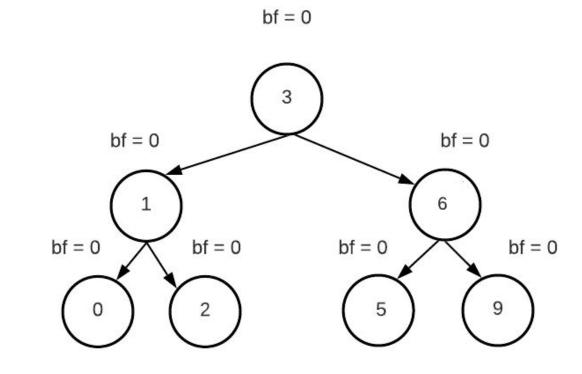
- Now 3 is left-left case
 - Observe: bf(3) = 2
 - Don't let that confuse u
- Remaining perform right_rotation(5)
 - o This pushes 5 down and 3 up
 - A=1, B=null, C=9



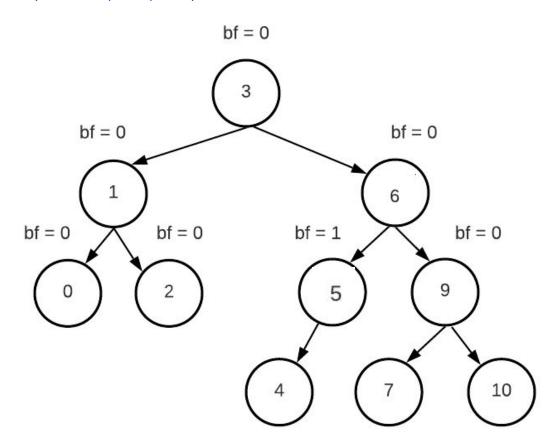


- After inserting 6, we have right-left case
 Bf: -2, 1
- Right-rotation(9) to convert to right-right
- Then Left-Rotation(5)
- As no children A/B/C this is visually direct

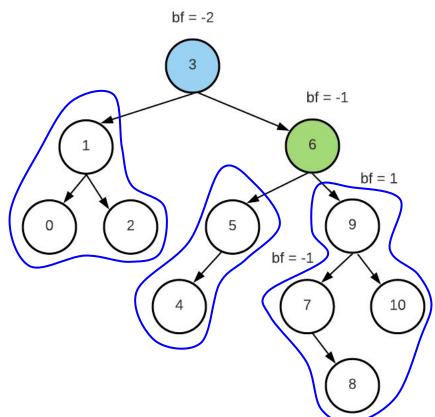


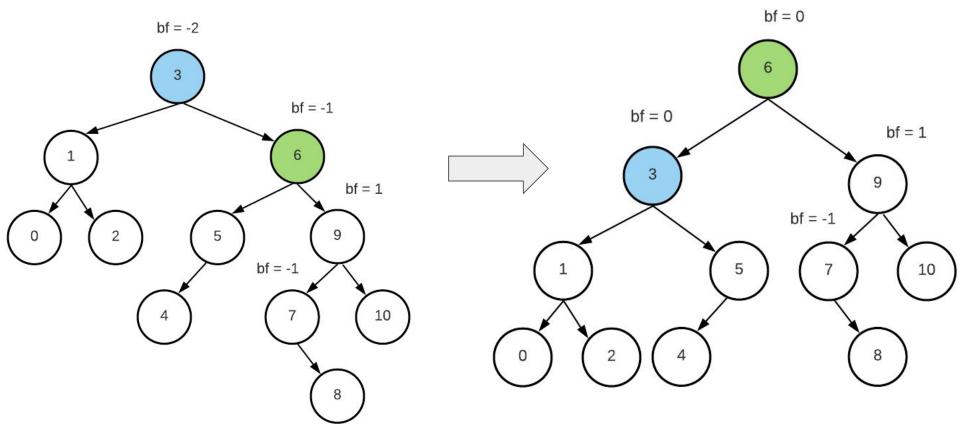


 Inserting next 3 values cause no problem



- After inserting 8, we have right-right case
 - o Bf: -2, -1
- Left-rotation(3)
 - o Gets 3 down and 6 up
 - \circ A = subtree(1)
 - \circ B = subtree(5)
 - It will change from left of 6 to right of 3
 - \circ C = subtree(9)





"Acquire knowledge and impart it to the people."

"Seek knowledge from the Cradle to the Grave."