



CS/CPE 590

Huffman Coding

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Huffman Codes

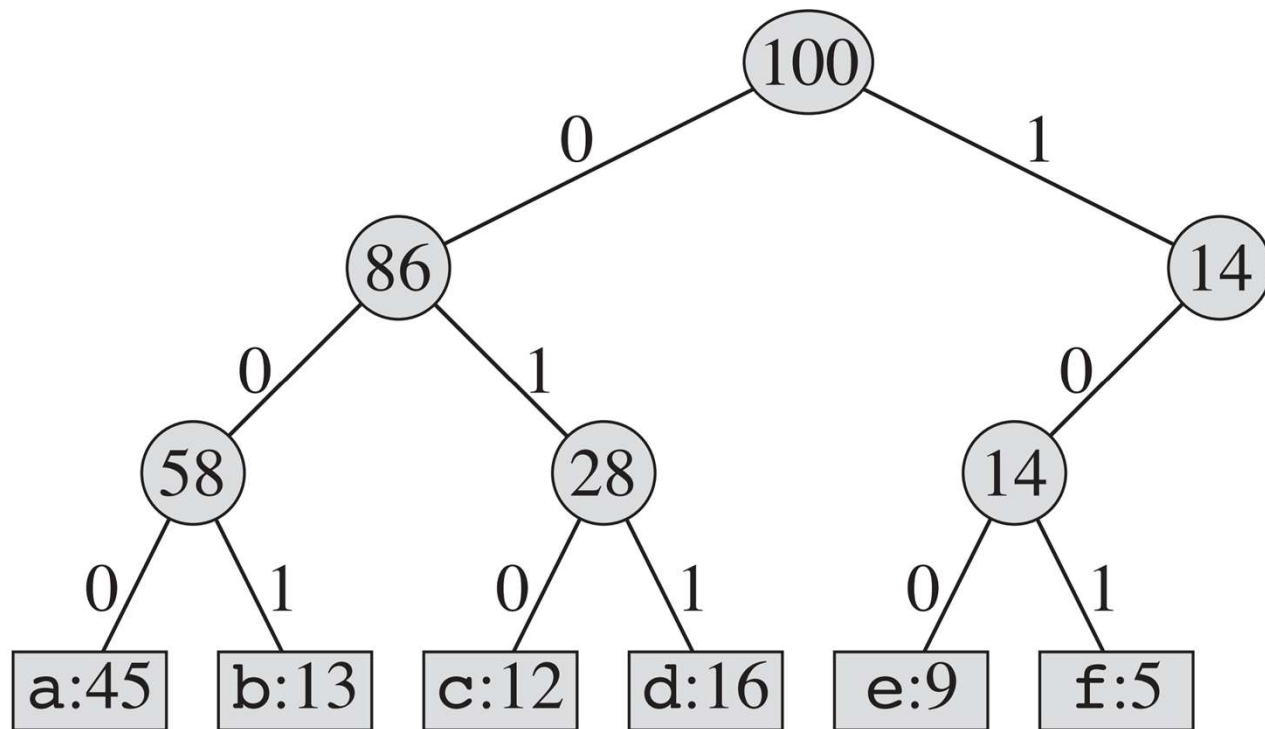
Huffman codes compress data very effectively: savings of 20% to 90% are typical, depending on the characteristics of the data being compressed. We consider the data to be a sequence of characters. Huffman's greedy algorithm uses a table giving how often each character occurs (i.e., its frequency) to build up an optimal way of representing each character as a binary string.

- Suppose we have a 100,000-character data file that we wish to store compactly.
- We observe that the characters in the file occur with the frequencies given by

	a	b	c	d	e	f
Frequency (in thousands)	45	13	12	16	9	5

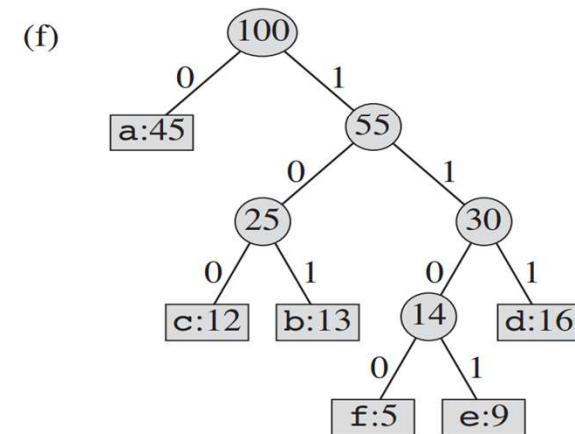
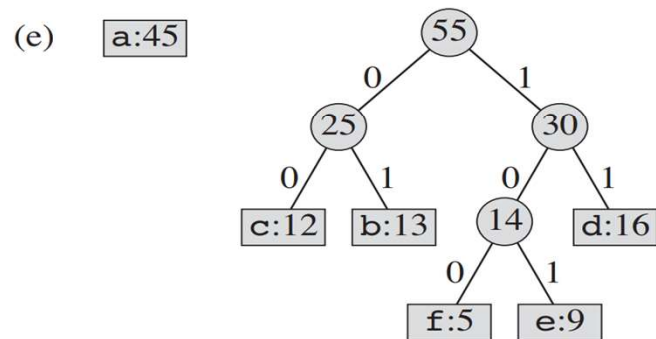
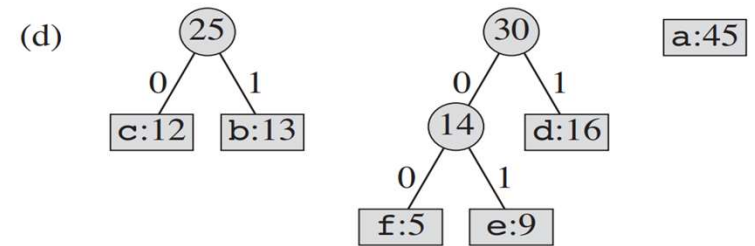
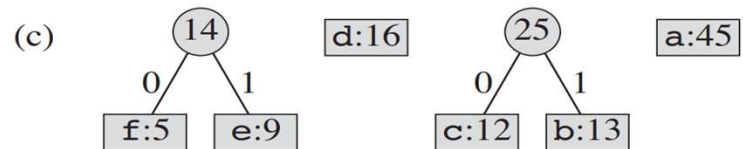
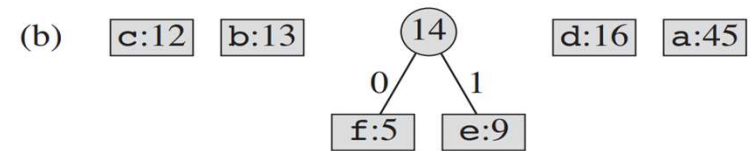
- That is, only 6 different characters appear, and the character “a” occurs 45,000 times.
- We have many options for how to represent such a file of information.
- We consider the problem of designing a binary character code

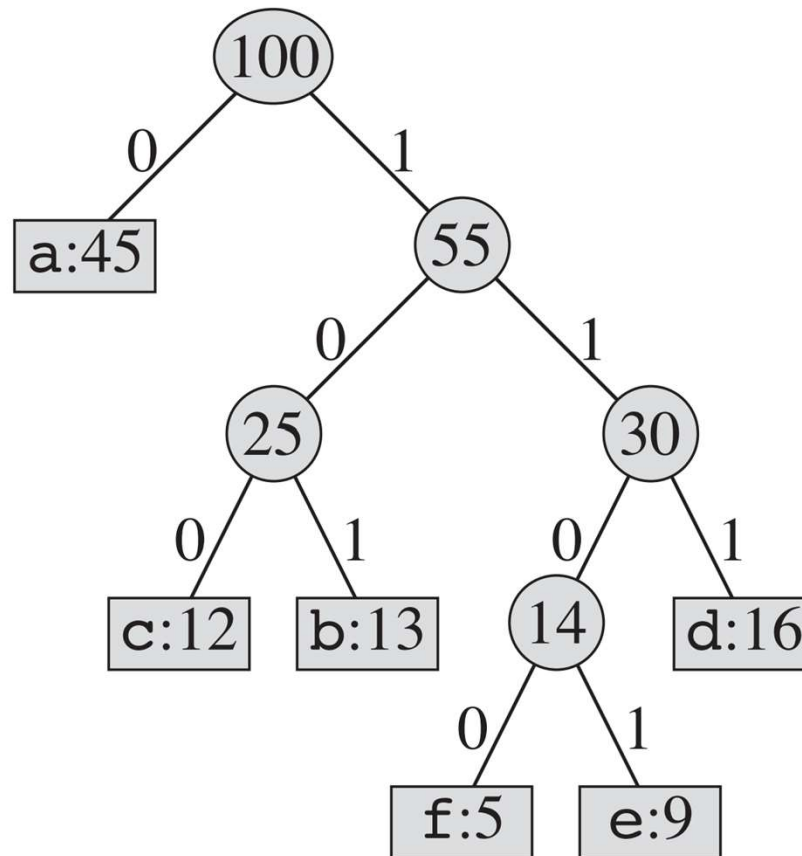
Fixed-length Coding



Variable-length Coding

(a) f:5 e:9 c:12 b:13 d:16 a:45





String 001011101 parses uniquely as 0 0 101 1101, which decodes to “aabe”

Algorithm

Input: a collection C of objects containing a character and its frequency.

Output: the root of a Huffman tree

HUFFMAN(C)

```
1   $n = |C|$ 
2   $Q = C$ 
3  for  $i = 1$  to  $n - 1$ 
4      allocate a new node  $z$ 
5       $z.left = x = \text{EXTRACT-MIN}(Q)$ 
6       $z.right = y = \text{EXTRACT-MIN}(Q)$ 
7       $z.freq = x.freq + y.freq$ 
8      INSERT( $Q, z$ )
9  return  $\text{EXTRACT-MIN}(Q)$     // return the root of the tree
```

Running time: ($O(n \lg n)$)

Compression Savings

- 100,000 characters each taking 1 byte = 8 bits
-- 800,000 bits (without any compression)
- With fixed-length codes, 3 bits for each character
-- 300,000 bits (62.5% savings)
- Variable-length codes (25.33% savings w.r.t fixed length codes & 72% savings w.r.t when stored as characters)

$$(45 \cdot 1 + 13 \cdot 3 + 12 \cdot 3 + 16 \cdot 3 + 9 \cdot 4 + 5 \cdot 4) \cdot 1,000 = 224,000 \text{ bits}$$

Contents of this presentation are based on
Book Chapter- 16, **Introduction to Algorithms** by *Cormen, Leiserson,
Rivest, & Stein*



THANK YOU

Stevens Institute of Technology