

# EMPIRICAL ANALYSIS OF PERFORMANCE OF QUICKSELECT WITH (SEDFEWICK) AND (LOMUTO) PARTITION ALGORITHMS

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## INTRODUCTION

Empirical analysis of an algorithm is performed by running a program implementing the algorithm on a sample of input and analyzing the data observed.

It is an alternative to the mathematical analysis of the algorithm's efficiency. and its applicability to any algorithm is the principal strength of this approach.

In this report, we will make an Empirical analysis of the Quickselect algorithm. *Quickselect* is a selection algorithm used to find the kth smallest element in an unordered list (Levitin, 2011). And it is a widely used algorithm.

## EXPERIMENT PURPOSE

The purpose of this study is to compare the time efficiency of two approaches of an algorithm that solve the same problem but with different partition methods.

- Quickselect algorithm with Sedgewick partition.
- Quickselect algorithm with Lomuto partition.

## ALGORITHMS

Algorithms used in this experiment:

### 1. Quickselect

Fig21, page 1 (Al-Hashimi M. , 2022)

### 2. Sedgewick Partition

Fig22, Page 6 (Al-Hashimi D. M., 2022)

### 3. Lomuto Partition

*Introduction to The Design and Analysis of Algorithms* Textbook Page 159  
(Levitin, 2011)

### 4. Random Numbers Generator

Pseudocode for generating random numbers:

```
// Fill array with random numbers from 0 to 100
// Input: an integer number with the desired size
// Output: an array filled with random numbers
Algorithm Fill(size)
    a = []
    for i = 0 and i < size
        a.push(Math.floor(Math.random() * 101))
        i++
    return a
```

### 5. Find Median Algorithm

```
// Find the median value
// Input: Array
// Output: Median value
function Median(A[0..n-1])
    r = A.length - 1
    l = 0
    return  $\left\lfloor \frac{l+r-2}{2} \right\rfloor$ 
```

## TOOLS

The tools used in this experiment:

- Microsoft Excel  
*Recording data in tables, calculating the limits, and drawing graphs.*
- Visual Studio Code  
*Help in writing and editing the code*
- JavaScript Programming language
- Firefox Web browser  
*Used developer tools to run and debug the program within the browser.*

## EFFICIENCY METRIC

An algorithm's time efficiency can be measured in two ways:

1. Basic Operation Count
2. Runtime Count

Basic operation count is done by counting the number of times the algorithm's basic operation is executed. And runtime count is done by counting the time of the program implementing the algorithm.

In this study, we will use both metrics, and after collecting them we will compare between them and choose a preferred one.

## PROCEDURE OF IMPLEMENTING THE ALGORITHM

### Input Sample

We used an array as an input for both algorithms. The array had the following characteristic:

### SAMPLE SIZE

Firstly, we started with a small array size, and then we increased it until we saw a big difference between the two algorithms' efficiencies. The final size we reached was 1,000,000.

### SAMPLE RANGE

Array's values were integer values in the range from 0 to 100. The values were generated randomly, since an empirical analysis requires generating random numbers, using a random number generator in the language library, shown in Algorithm 4.

### Program Code

The program code generates different input sizes starting from 100,000 until 1,000,000 and loops into the algorithm and makes 10 trials for each input size, as the result may differ for some instances of the same size. For each trial, there's a counter that counts the number of times the basic operation has been executed in the most inner loop and calculated the algorithm's runtime. For this experiment, we chose the kth element to be sorted as always the median value in the array, using Find Median Algorithm (Algorithm 5). Moreover, we made sure to use the same instance for comparing between the two partitions by shuffling the array before using it.

## RECORDING DATA

After running the code, all the data have been recorded in tables. Then, calculated the best, average, and worst cases for each input size of the 10 trials. As shown in Tables 1, 2, 3, and 4.

Size	Trial 1	Trial 2	Trial 3	Trial 4	Trial 5	Trial 6	Trial 7	Trial 8	Trial 9	Trial 10	Best	Average	Worst
100000	328127	302812	542683	306074	340850	227151	328960	294025	380265	394283	227151	344523	542683
200000	593430	976699	617686	777146	624749	917600	530420	552360	640721	389398	389398	662020.9	976699
300000	1037024	816138	993992	1026354	818501	1130198	724481	1159855	875073	769394	724481	935101	1159855
400000	1423322	686483	1373128	1012208	995902	1393395	1673664	1416149	1153065	1314808	686483	1244212.4	1673664
500000	998783	1958594	1286542	1547264	1068410	1751334	1277481	1879560	1800178	1724043	998783	1529218.9	1958594
600000	1967220	2250291	1780735	1702781	1507453	1791070	1579235	2573911	1416748	1364566	1364566	1793401	2573911
700000	2301632	2414500	2671346	2041054	2168649	2736749	3142975	2144423	1998592	2328461	1998592	2394838.1	3142975
800000	2285292	2876577	2831881	2854333	1413663	3303043	2608226	2036767	2189663	2948754	1413663	2534819.9	3303043
900000	2088851	1994450	4940028	3113876	2253755	1929679	2239336	2774952	4505119	2573189	1929679	2841323.5	4940028
1000000	3358066	3912236	3623745	3493458	1873785	3304978	2677558	3412468	2843661	4471239	1873785	3297119.4	4471239

Table 1 Sedgewick Basic Operation count

Size	Trial 1	Trial 2	Trial 3	Trial 4	Trial 5	Trial 6	Trial 7	Trial 8	Trial 9	Trial 10	Best	Average	Worst
100000	15	12	13	10	11	7	8	11	13	12	7	11.2	15
200000	9	19	11	13	9	18	9	9	9	9	9	11.5	19
300000	26	17	16	19	17	20	15	20	18	12	12	18	26
400000	48	18	24	21	20	26	26	24	21	25	18	25.3	48
500000	24	33	24	27	36	46	22	36	30	42	22	32	46
600000	46	35	44	35	33	40	35	53	33	37	33	39.1	53
700000	38	78	41	32	36	43	42	35	32	43	32	42	78
800000	41	55	46	45	30	70	40	37	39	45	30	44.8	70
900000	59	44	90	41	62	48	36	38	93	43	36	55.4	93
1000000	59	63	54	53	37	74	47	53	47	58	37	54.5	74

Table 2 Sedgewick Runtime Count

Size	Trial 1	Trial 2	Trial 3	Trial 4	Trial 5	Trial 6	Trial 7	Trial 8	Trial 9	Trial 10	Best	Average	Worst
100000	487759	789758	902310	980367	589584	904434	865491	524188	675294	773424	487759	749260.9	980367
200000	1825695	2216625	1885853	2094967	2262101	1892087	2967666	2059010	2168663	2110791	1825695	2148345.8	2967666
300000	3907599	4645063	3701079	4535518	4020215	4151085	4640594	4634032	4928433	3297005	3297005	4246062.3	4928433
400000	6217023	9406307	7827253	6249519	7969872	8276252	7134592	7436505	7155862	8371841	6217023	7604502.6	9406307
500000	11462853	10782354	10410704	11410394	8880874	9535626	10833284	10170950	9806087	12977985	8880874	10627111.1	12977985
600000	11318499	13570582	17478019	13100632	12535612	15199950	14873795	15032839	14329036	13206521	11318499	14064548.5	17478019
700000	20918622	22104463	16513094	19152642	23483057	17178055	17898575	21457279	22272912	18633695	16513094	19961239.4	23483057
800000	25820752	27404584	26361537	30787289	24964896	20281246	24574535	27095645	25272237	26836685	20281246	25939940.6	30787289
900000	34427871	31446687	33077202	29097139	36318804	36690280	38240058	30710806	31540223	29682079	29097139	33123114.9	38240058
1000000	39200890	40213272	43570092	41788883	44426258	43316892	41988634	43585073	36809315	39891583	36809315	41479089.2	44426258

Table 3 Lomuto Basic Operation Count

Size	Trial 1	Trial 2	Trial 3	Trial 4	Trial 5	Trial 6	Trial 7	Trial 8	Trial 9	Trial 10	Best	Average	Worst
100000	14	23	24	24	10	20	18	11	15	18	10	17.7	24
200000	31	34	29	27	31	26	35	28	26	44	26	31.1	44
300000	69	72	44	57	47	49	47	58	52	40	40	53.5	72
400000	75	188	113	75	111	106	103	87	78	84	75	102	188
500000	152	122	110	138	90	108	124	101	110	146	90	120.1	152
600000	174	191	267	182	188	198	198	184	209	169	169	196	267
700000	277	303	199	234	366	169	193	247	233	243	169	246.4	366
800000	297	317	315	483	266	201	272	306	289	289	201	303.5	483
900000	470	369	413	310	466	444	590	326	386	465	310	423.9	590
1000000	552	488	636	661	663	619	595	544	562	502	488	582.2	663

Table 4 Lomuto Runtime Count

## GRAPH THE DATA

Using Excel tools for drawing graphs, we draw a graph for each table as shown in Figures 1, 2, 3, and 4. Drawing a graph can help us identify an algorithm's efficiency class.

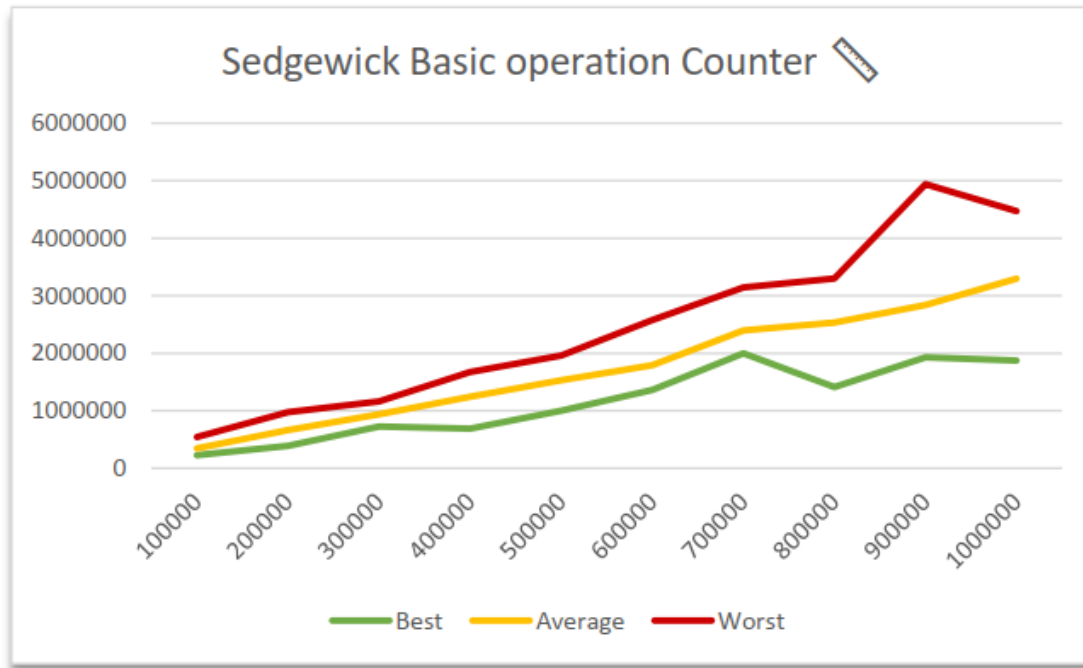


Figure 1 Sedgewick Basic operation Counter

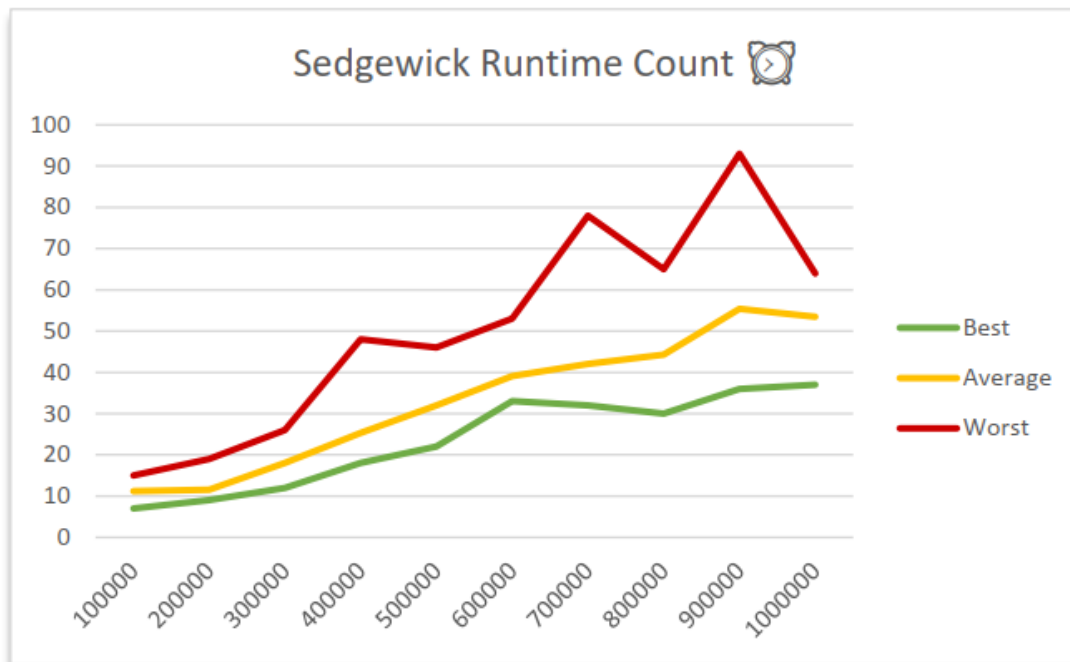


Figure 2 Sedgewick Runtime Count



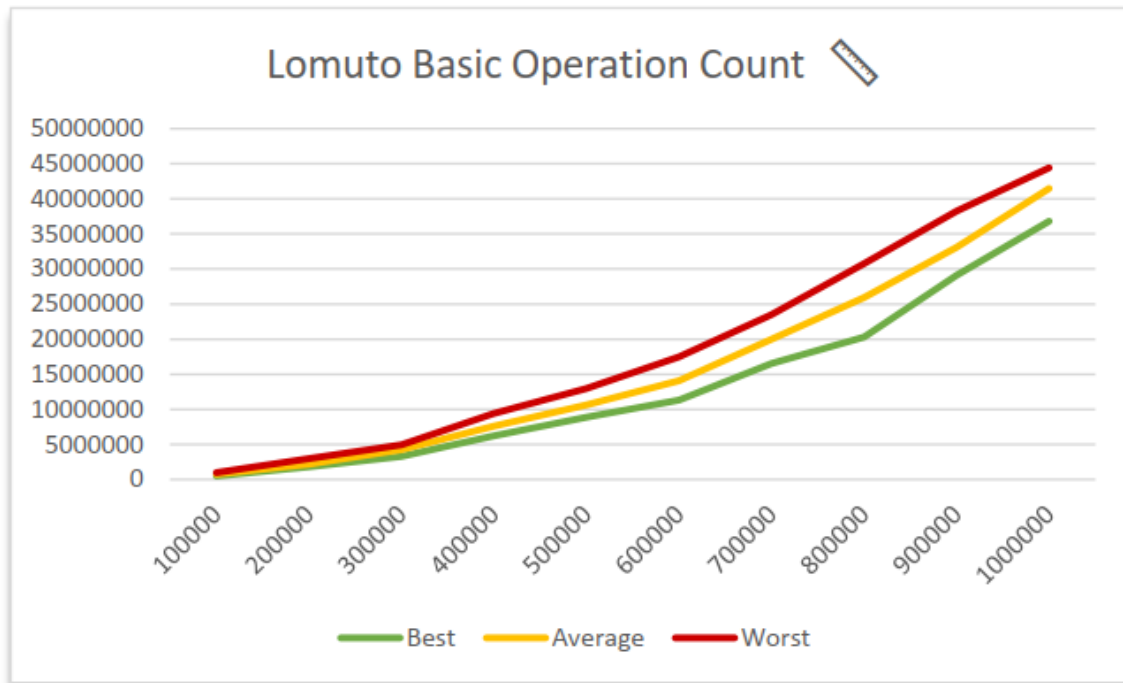


Figure 3 Lomuto Basic Operation Count

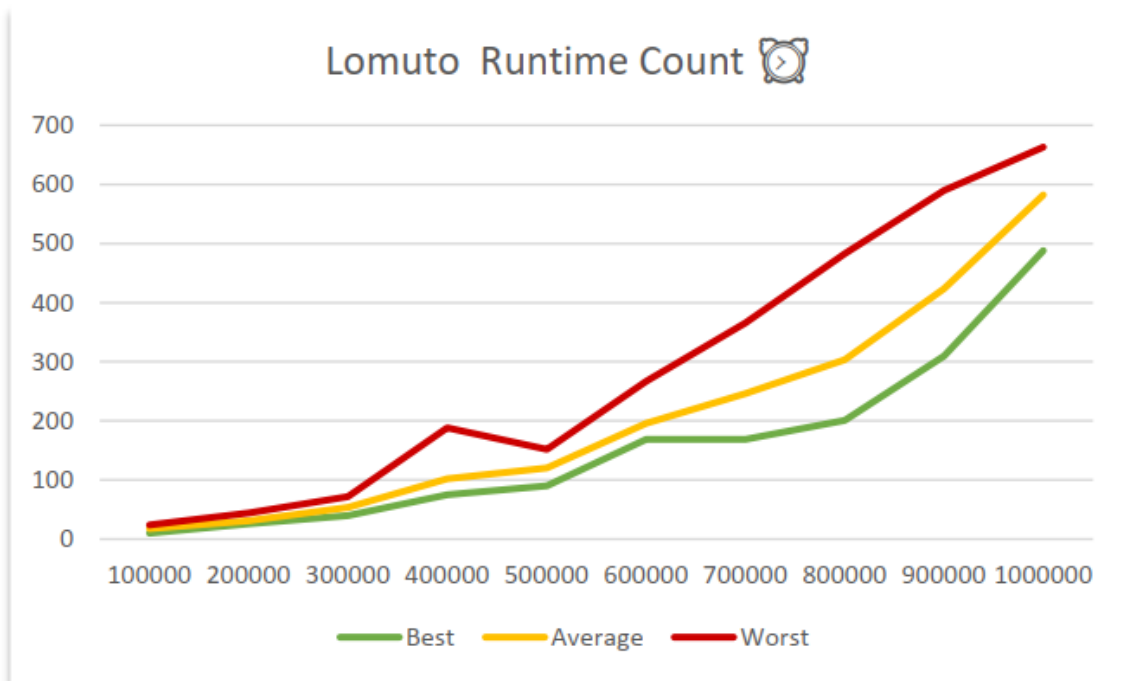


Figure 4 Lomuto Runtime Count

## LIMITS

By using Excel tools in calculating limits as  $n$  gets large  $\lim_{n \rightarrow \infty} \left( \frac{t(n)}{g(n)} \right)$

$$\text{Sedgewick Basic operation Counter} = \lim_{n \rightarrow \infty} \left( \frac{t(n)}{n} \right) = +ve \text{ Constant}$$

$$\text{Sedgewick Runtime} = \lim_{n \rightarrow \infty} \left( \frac{t(n)}{n} \right) = +ve \text{ Constant}$$

$$\text{Lomuto Basic operation Counter} = \lim_{n \rightarrow \infty} \left( \frac{t(n)}{n \log n} \right) = +ve \text{ Constant}$$

$$\text{Lomuto Runtime} = \lim_{n \rightarrow \infty} \left( \frac{t(n)}{n \log n} \right) = +ve \text{ Constant}$$

## ANALYZING DATA

In this experiment, we used two methods for investigating the order of growth, and each method has its unique strength and weakness.

The first one is by drawing a graph, and the second is by using limits as  $n$  increases to  $\infty$ .

Analyzing the graphs, all four graphs show no variability, which means that we do not have to investigate the best, average, and worst cases.

Graphs 1 and 2 follow the same behavior, they show a straight line, and that leads to efficiency class  $\Theta(n)$ .

Graphs 3 and 4 show a convex shape and their efficiency class is  $\Theta(n \log n)$ .

By using limits to investigate the order of growth, all the limits' results confirm the order of growth we obtained from the graphs. When  $\lim_{n \rightarrow \infty} \left( \frac{t(n)}{g(n)} \right) = c$  that implies that  $t(n)$  has the same order of growth of  $g(n)$  which means that  $t(n) \in \theta(g(n))$

## PREFERRED METRIC CHOICE

The preferred metric we chose is the basic operation counter.

The physical runtime depends on a specific machine, as we mentioned earlier, Quickselect is a widely used algorithm, and if we used it on a supercomputer the results we obtained will significantly change, unlike the basic operation counter will remain the same.

## DISCUSSION OF RESULTS

### Comparing to Theory

The mathematical analysis of Quickselect with Lomuto and Sedgewick partitions:

$$C_{Best, Avg} \in \theta(n), C_{Worst} \in \theta(n^2)$$

And by comparing the mathematical with the empirical analysis we can indicate that the empirical analysis is not accurate.

### Reasons for Incorrect Analysis

The reason for not getting the right answers is because of the input values range.

In the experiment, we limited our input range by only values starting from 0 to 100. And that was not as random input as we expected.

Another reason is, the number of trials was not enough to investigate the algorithm's best, average, and worst cases.

## CONCLUSION

In conclusion, the empirical analysis does not always give an accurate analysis of an algorithm's performance. The weakness of empirical analysis lies in depending on a particle sample of instances, and the computer used in the experiment.

## REFERENCES

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