

PERMUTATIONS AS GENOME REARRANGEMENTS

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REMARK 0.1. *RECALL: inversion*

1. BLOCK INTERCHANGES

1.1. The Cycle Graph and Its Relation to *bid*.

DEFINITION 1.1. A *Block Interchange* is an operation that interchanges two blocks of consecutive entries without rearranging said entries.

EXAMPLE 1.1. $|34|17|562| \rightarrow |562|17|34|$

DEFINITION 1.2. Let the *Block Interchange Distance* between two n -permutation p, q , denoted as $\text{bid}(p, q)$, be the smallest number of block interchanges required to transform p into q . Let $\text{bid}(p) := \text{bid}(p, \text{id})$ be the number of interchanges required to sort p .

REMARK 1.1. $\text{bid}(p, q)$ is a metric and is left-invariant, just as we have seen for $\text{btd}(p, q)$.

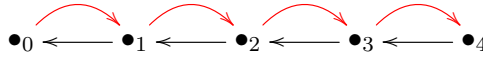
DEFINITION 1.3. Let $p = p_1 p_2 \cdots p_n$ be an n -permutation. The *Cycle Graph* $G(p)$ is a graph of coloured directed edges on the vertex set $\{0, 1, \dots, n, n+1\}$ constructed as follows:

- (1) For every i with $1 \leq i \leq n+1$: Add a black edge from p_i to p_{i-1}
- (2) For every i with $0 \leq i \leq n$: Add a red edge from i to $i+1$

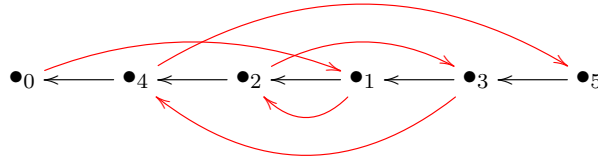
with $p_0 = 0$ and $p_{n+1} = n+1$.

REMARK 1.2. The *Cycle Graph* has $2n+2$ edges.

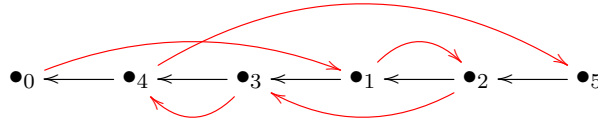
EXAMPLE 1.2. $p_1 = 123 \implies c(G(p_1)) = 4$



$p_2 = 4213 \implies c(G(p_2)) = 1$



$p_3 = 4312 \implies c(G(p_3)) = 3$



For the final permutation p_3 notice that $\text{bid}(p_3) = 1 = \frac{4+1-3}{2}$, something which will become relevant in a later theorem.

LEMMA 1.1. The cycle graph $G(p)$ has a unique decomposition into edge-disjoint directed cycles in which the colors of the edges alternate.

Proof. It is easy to see that a path of alternating colours cannot branch since by construction there is always only one edge per color directed away from every node. We begin at the 0-node and complete a cycle (this must occur since $G(p)$ is finite). Once we have done so, we remove the edges of this cycle from $G(p)$ and iterate. Notice that in removing full cycles the structure is maintained as we only remove disjointly coloured pairs of edges from every node. The cycles are then edge-disjoint by construction. \square

DEFINITION 1.4. Let $c(G(p))$ be the number of alternating directed cycles in the decomposition of $G(p)$. To avoid confusion, let $c(\Gamma(p))$ denote the number of cycles in the traditional sense of a permutation. From now on the cycles of $G(p)$ will refer to the alternating cycles.

REMARK 1.3. (1) The cycles are only edge-disjoint not vertex-disjoint.
(2) If there is an alternating path from one node to another, both nodes are in the same cycle.
(3) The identity permutation has $n + 1$ cycles. In fact, it is the only permutation with this many cycles.

Before we can find a formula for $\text{bid}(p)$ we shall examine how single block interchanges can influence the block interchange distance of a permutation in the following two lemmas.

DEFINITION 1.5. Let $p = p_1 p_2 \cdots p_n$ with $p \neq \text{id}$, or in other words, p contains at least one inversion. Select the inversion in p $(i, j \in [n] : p_i > p_j \wedge i < j)$ such that p_j is minimal and if necessary, as a second criterion that p_i is maximal.

Let $x := p_j, y := p_i$ to ease our eyes. By our selection criterion we find that $x - 1$ occurs before y in p , whilst $y + 1$ occurs after x . (In the case of $x = 1$ simply let $x - 1 = 0$)

$$p = \cdots (x - 1) | \cdots y | \cdots | x \cdots | (y + 1) \cdots$$

The Canonical Block Interchange is defined as the interchange transforming p into

$$p' = \cdots (x - 1) | x \cdots | \cdots | \cdots y | (y + 1) \cdots$$

LEMMA 1.2. 9.10. Let $p \in S_n$ with $p \neq \text{id}$. Then there exists a block interchange which increases $c(G(p))$ by two.

Proof. We shall prove that the canonical block interchange suffices, or symbolically:

$$c(G(p')) = c(G(p)) + 2.$$

Notice first that any interchange does not change any red edges in $G(p)$. The canonical interchange changes 3 black edges when x and y are adjacent, and 4 edges otherwise. We shall perform a proof by cases.

(1) x, y adjacent: The cycle decomposition of $G(p)$ will contain an alternating path through the following nodes in respective order:

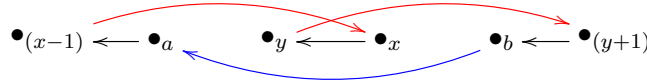
$$a \rightarrow (x - 1) \rightarrow x \rightarrow y \rightarrow (y + 1) \rightarrow b$$

where a is the entry in p directly to the right of $(x - 1)$ and b the one to the left of $(y + 1)$.

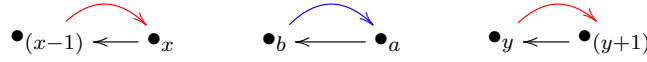
$$p = \cdots (x - 1) | a \cdots y | \cdots | x \cdots b | (y + 1) \cdots$$

Since a path exists, we find that the entries are part of the same cycle C . After performing the canonical interchange, C is split into 3 cycles. In the black edges we now have $x \rightarrow (x-1)$, $(y+1) \rightarrow y$, and $a \rightarrow b$ while the red have been untouched as seen in the following illustration:

p



$\rightarrow p'$



Notice the arrow in blue is not necessarily a singular edge, however it is a well defined alternating path.

- (2) x, y not adjacent: We must introduce another notation in this case. Let a, b, c, d be such that they are adjacent to $(x-1), y, x, (y+1)$ respectively pair-wise. Since they are neighbours we will find the alternating path

$$P_a = a \rightarrow (x-1) \rightarrow x \rightarrow c$$

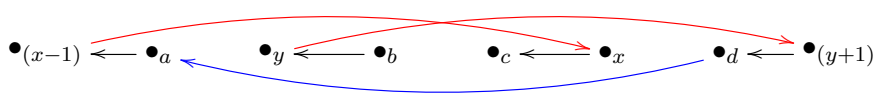
and

$$P_b = b \rightarrow y \rightarrow y+1 \rightarrow d$$

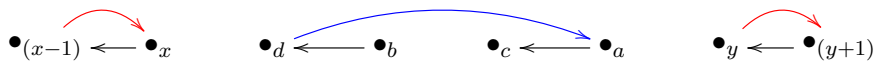
We must make a further division of cases at this point.

- (a) P_a, P_b are of one cycle: Denote the cycle as C_{ab} . In this case there exists an alternating path A_{da} from d to a with the final edge being grey. See the following illustration:

p



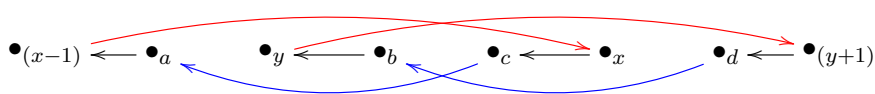
$\rightarrow p'$

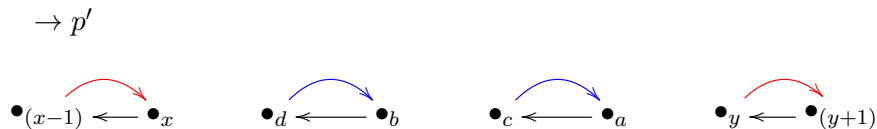


where the blue edge is our path A_{da} , the red edges are kept static, and once more we have split one cycle into three, all else remaining the same.

- (b) P_a, P_b are of distinct cycles: Denote these cycles C_a and C_b respectively. We find that each cycle shall split into two, we shall cement this only with a short illustration as the argument is analogue.

p





Where the blue edges are our paths C_a and C_b exempted $(x-1) \rightarrow x$ and $y \rightarrow (y+1)$.

Having exhausted all cases, we conclude the proof. □

LEMMA 1.3. 9.11 A block interchange cannot increase $c(G(p))$ by more than two.

Proof. Proof Sketch Only □

THEOREM 1.1. 9.9 Let $p \in S_n$. Then $\text{bid}(p) = \frac{n+1-c(G(p))}{2}$

Proof. By the previous two lemmata we have for any $p \in S_n$ that

$$\frac{n+1-c(G(p))}{2} \leq \text{bid}(p) \leq \frac{n+1-c(G(p))}{2}$$

since we require this amount of block interchanges to attain $n+1$ cycles in p , thus culminating our desired equality. □

REMARK 1.4. We have discovered that we need at most $\lfloor \frac{n}{2} \rfloor$ block interchanges to sort an n -permutation. Furthermore $\text{bid}(p) \in \mathbb{N}$ implies that $(n+1)$ and $c(G(p))$ must either both be odd or even.

1.2. Average Number of Block Interchanges Required to Sort p .

DEFINITION 1.6. Define the Hultman Number as

$$S_H(n, k) := |\{\pi \in S_n : c(G(\pi)) = k\}|$$

EXAMPLE 1.3. $S_H(3, 4) = 1$ and $S_H(3, 2) = 5$. Draw sample graphs quickly? Maybe at least two. Continue example after 9.14 et al.

THEOREM 1.2. 9.14 and Corollary? Proof Sketch?

REMARK 1.5. The previous theorem illustrates the connection between $c(G(p))$ and $c(\Gamma(p))$. It allows us to translate ?????????? Skip all the other stuff, useless?

THEOREM 1.3. 9.20? 9.21, 9.22? just write theorems and quickly talk about at most.

THEOREM 1.4. 9.23. The average number of block interchanges needed to sort an n -permutation is

$$\frac{1}{2} \left(n - \frac{1}{\left\lfloor \frac{n+2}{2} \right\rfloor} - \sum_{i=2}^n \frac{1}{i} \right)$$