PERMUTATIONS AS GENOME REARRANGEMENTS

SIMON GRÜNING

Remark 0.1. RECALL: inversion

1. Block Interchanges

1.1. The Cycle Graph and Its Relation to bid.

Definition 1.1. A Block Interchange is an operation that interchanges two blocks of consecutive entries without rearranging said entries.

Example 1.1. $|34|17|562| \rightarrow |562|17|34|$

DEFINITION 1.2. Let the Block Interchange Distance between two n-permutation p, q, denoted as bid(p,q), be the smallest number of block interchanges required to transform p into q. Let bid(p) := bid(p, id) be the number of interchanges required to sort p.

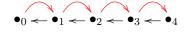
Remark 1.1. bid(p,q) is a metric and is left-invariant, just as we have seen for btd(p,q).

DEFINITION 1.3. Let $p = p_1 p_2 \cdots p_n$ be an n-permutation. The Cycle Graph G(p) is a graph of coloured directed edges on the vertex set $\{0, 1, \dots, n, n+1\}$ constructed as follows:

- (1) For every i with $1 \le i \le n+1$: Add a black edge from p_i to p_{i-1}
- (2) For every i with $0 \le i \le n$: Add a red edge from i to i+1 with $p_0 = 0$ and $p_{n+1} = n+1$.

Remark 1.2. The Cycle Graph has 2n + 2 edges.

Example 1.2. p = 123



p = 4312



draw cycle graph count cycles find minimal transpositions to id and show how it works

Lemma 1.1. The cycle graph G(p) has a unique decomposition into edge-disjoint directed cycles in which the colors of the edges alternate.

Proof. It is easy to see that a path of alternating colours cannot branch since by construction there is always only one edge per color directed away from every node. We begin at the 0 node and complete a cycle (this must occur since G(p) is finite). Once we have done so, we remove the edges of this cycle from G(p) and iterate. Notice that in removing full cycles the structure is maintained as we only remove disjointly coloured pairs of edges from every node. The cycles are then edge-disjoint by construction.

DEFINITION 1.4. Let c(G(p)) be the number of alternating directed cycles in the decomposition of G(p). To avoid confusion, let $c(\Gamma(p))$ denote the number of cycles in the traditional sense of a permutation. From now on the cycles of G(p) will refer to the alternating cycles.

Remark 1.3. The cycles are only edge-disjoint not vertex-disjoint.

Remark 1.4. If there is an alternating path from one node to another, both nodes are in the same cycle.

Remark 1.5. The identity permutation has n+1 cycles. In fact, it is the only permutation with this many cycles.

Before we can find a formula for bid(p) we shall examine how single block interchanges can influence the block interchange distance of a permutation in the following two lemmas.

Definition 1.5. Canonical Block Interchange of p.

LEMMA 1.2. 9.10 Let $p \in S_n$ with $p \neq id$. Then there exists a block interchange which increases c(G(p)) by two.

Proof. 3 cases
$$\Box$$

Lemma 1.3. 9.11 A block interchange cannot increase c(G(p)) by more than two.

Theorem 1.1. 9.9 Let $p \in S_n$. Then $bid(p) = \frac{n+1-c(G(p))}{2}$

Proof. By the previous two lemmata we have for any $p \in S_n$ that

$$\frac{n+1-c(G(p))}{2} \le bid(p) \le \frac{n+1-c(G(p))}{2}$$

since we require this amount of block interchanges to attain n+1 cycles in p, thus culminating our desired equality.

REMARK 1.6. We have discovered that we need at most $\lfloor \frac{n}{2} \rfloor$ block interchanges to sort an n-permutation. Furthermore $bid(p) \in \mathbb{N}$ implies that (n+1) and c(G(p)) must either both be odd or even.

1.2. Average Number of Block Interchanges Required to Sort p.

Definition 1.6. Define the Hultman Number as

$$S_H(n,k) := |\{\pi \in S_n : c(G(p)) = k\}|$$

Example 1.3. $S_H(3,4) = 1$ and $S_H(3,2) = 5$. Draw sample graphs quickly? Maybe at least two. Continue example after 9.14 et al.

THEOREM 1.2. 9.14 and Corollary? Proof Sketch?

Remark 1.7. The previous theorem illustrates the connection between c(G(p)) and $c(\Gamma(p))$. It allows us to translate ????????? Skip all the other stuff, useless?

THEOREM 1.3. 9.20? 9.21, 9.22? just write theorems and quickly talk about at most.

Theorem 1.4. 9.23. The average number of block interchanges needed to sort an n-permutation is

$$\frac{1}{2}(n - \frac{1}{\left|\frac{n+2}{2}\right|} - \sum_{i=2}^{n} \frac{1}{i})$$