



Resilient Consensus for Multi-Agent Systems (MASs) Through Dynamic Event-Triggered Mechanism

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Background



Swarm Performance



Agriculture



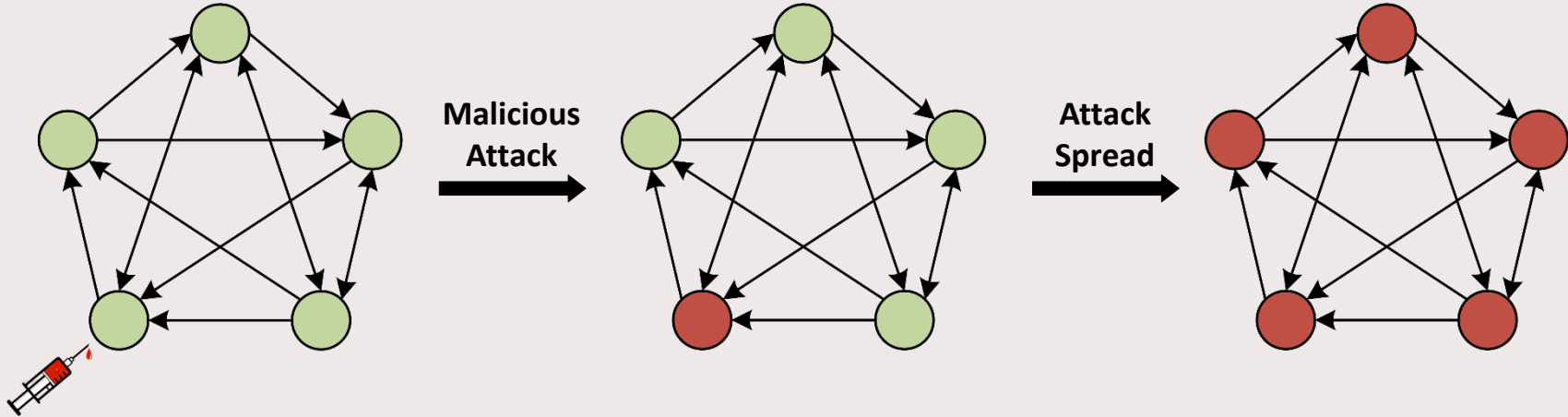
Area Surveillance

Application of MASs in Various Fields



Target Tracking

Background

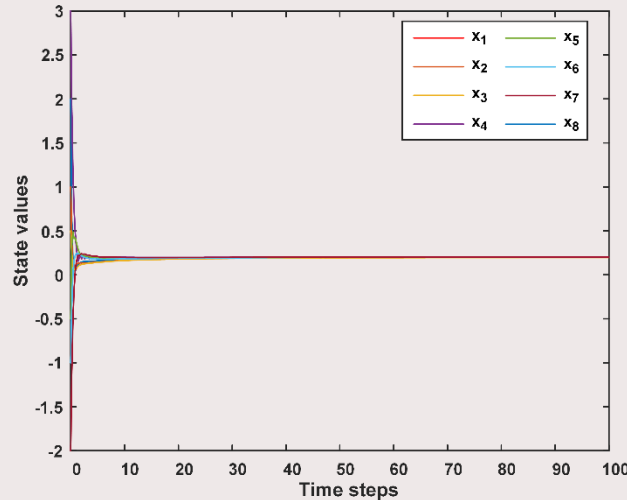


System Model

- Consider a single-integrator MAS modelled by
- A consensus-seeking protocol^[3] is introduced as

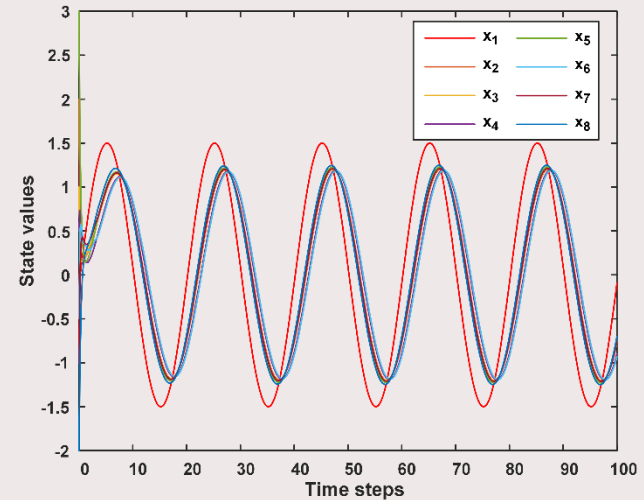
$$x_i[k+1] = x_i[k] + u_i[k]$$

$$u_i[k] = \sum_{j \in \mathcal{V}_i} a_{ij}[k] (x_j[k] - x_i[k])$$



- Without malicious attack

**Malicious
Attack**



- Agent 1 is subject to malicious attack

[3] W. Ren and Y. Cao, Distributed Coordination of Multi-Agent Networks: Emergent Problems, Models, and Issues. Springer, 2011.

Background

Definition 1 (Normal Agent). An agent is said to be normal if **adopts the predetermined rule for state update**.

Definition 2 (Malicious Agent). An agent is said to be malicious if it **adopts some other function at some time step**.

Challenges

- Resilient algorithm design
- Communication load
- Heterogeneity
- Time-varying networks
-

Existing Studies

Attack Detection and Isolation Method^[1]

Attack tolerant Method^[2]

- Detecting and isolating malicious agents
- Agents in the network are equipped with observers
- Massive data processing
- High network connectivity required

- Deleting malicious state values
- More **lightweight** with less **computational complexity**
- Less **data processing**
- Less **prior information** required

[1] I. Shames, A. M. Teixeira, H. Sandberg, and K. H. Johansson, “Distributed fault detection for interconnected second-order systems,” *Automatica*, vol. 47, no. 12, pp. 2757–2764, 2011.

[2] H. J. LeBlanc, H. Zhang, X. Koutsoukos, and S. Sundaram, “Resilient asymptotic consensus in robust networks,” *IEEE J. Sel. Areas Commun.*, vol. 31, no. 4, pp. 766–781, 2013.

Resilient Consensus

- **Resilience:** For each normal agent $i \in \mathcal{N}$, it holds $x_i[k] \in \mathcal{I}$, $\forall k \in \mathbb{Z}_{\geq 0}$, where

$$\mathcal{I} = [\min_{i \in \mathcal{N}} x_i[0], \max_{i \in \mathcal{N}} x_i[0]]$$

represents a safety interval.

- **(Exact) Consensus:** For each pair of normal agents $i, j \in \mathcal{N}$, it holds

$$\lim_{t \rightarrow \infty} |x_i[k] - x_j[k]| = 0, \quad \forall k \in \mathbb{Z}_{\geq 0}.$$

Mean Subsequence Reduced (MSR) Algorithm



1. Sorting



2. Deleting



3. Updating

MSR Algorithm



1. Sorting

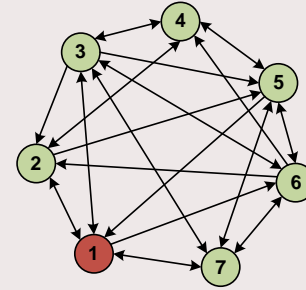


2. Deleting



3. Updating

Example:



$$[x_1[0], \dots, x_7[0]] = [10, 8, 6, 7, 4, -5, 6], \quad f=2$$

$$\text{Agent 2: } x_2[0] = 8 \quad x_j[0] = [10, 6, 7, -5]$$



1. Sorting

$$S = [-5, 6, 7, 8, 10]^T$$



2. Deleting

$$S_{del} = [-5, 6, 7, 8, 10]^T, \quad \mathcal{R}_i[k] = \{4\}$$

~~×~~ ~~×~~ ~~×~~

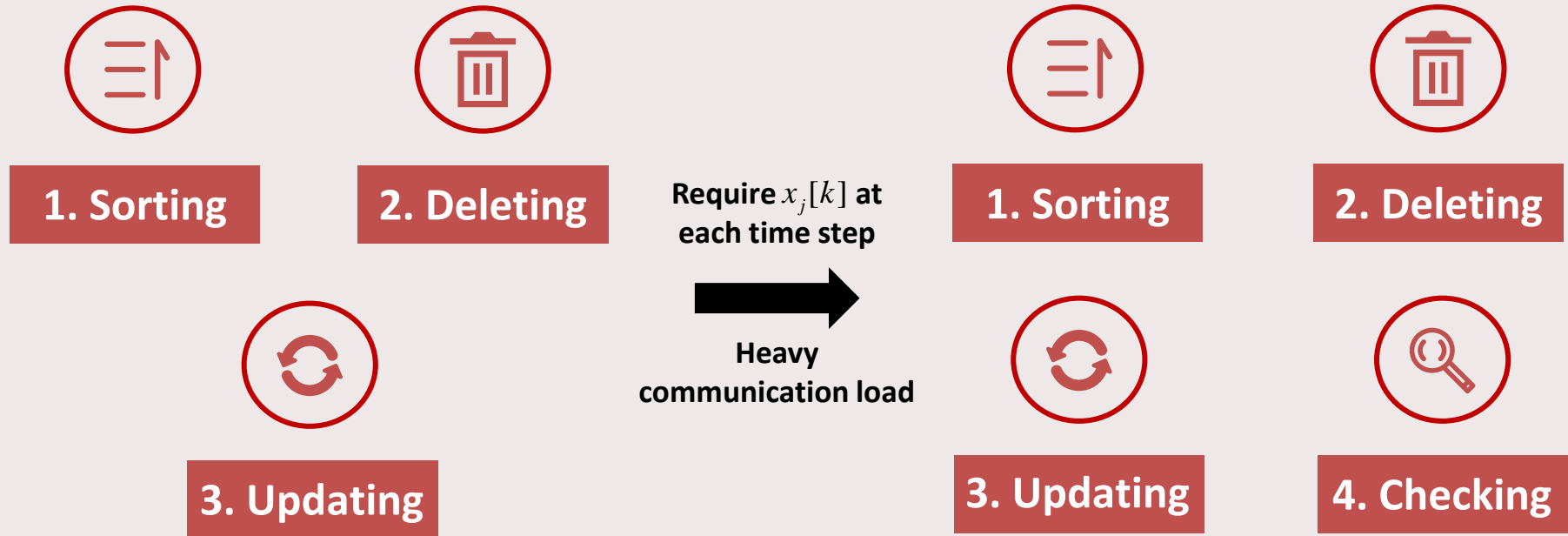


3. Updating

$$x_i[k+1] = x_i[k] + \sum_{j \in \mathcal{R}_i[k]} a_{ij}[k] (x_j[k] - x_i[k])$$

MSR Algorithm

Event-based MSR (E-MSR) Algorithm^[4]



[4] Y. Wang and H. Ishii, “Resilient consensus through event-based communication,” IEEE Trans. Control Netw. Syst., vol. 7, no. 1, pp. 471–482, 2020.

Static Event-Triggered Mechanism^[4]

- The control protocol is modified as

$$x_i[k+1] = x_i[k] + \sum_{j \in \mathcal{R}_i[k]} a_{ij}[k] \underbrace{(\hat{x}_j[k] - x_i[k])}_{\text{Auxiliary variable}}$$

- The update of auxiliary variable depends on

$$\hat{x}_i[k+1] = \begin{cases} x_i[k+1], & \text{if } f_i[k] > 0, \\ \hat{x}_i[k], & \text{otherwise.} \end{cases}$$

- The static event-triggered function (SETF) is designed as

$$f_i[k] = \underbrace{|x_i[k+1] - \hat{x}_i[k]|}_{\text{Error term } |e_i[k]|} - \underbrace{(c_0 + c_1 e^{-\alpha k})}_{\text{Threshold}}$$

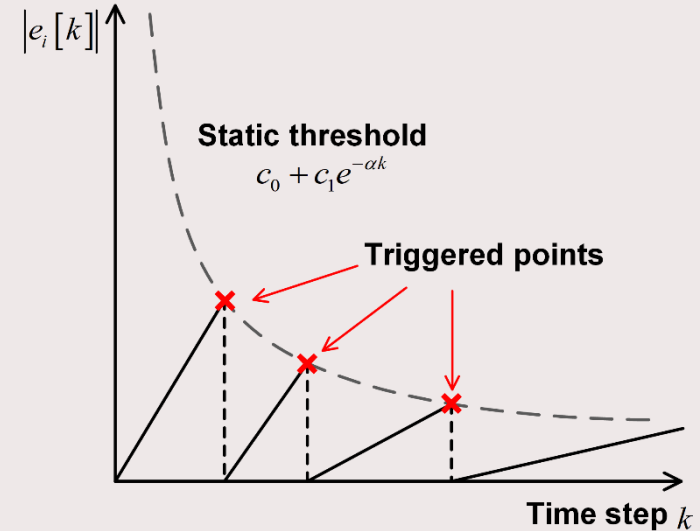


Illustration of the static event-triggered mechanism.

[4] Y. Wang and H. Ishii, “Resilient consensus through event-based communication,” IEEE Trans. Control Netw. Syst., vol. 7, no. 1, pp. 471–482, 2020.

E-MSR Algorithm^[4]



1. Sorting



2. Deleting



3. Updating



4. Checking

4. Agent i checks whether SETF triggers or not and sets $\hat{x}_i[k+1]$ as

$$\hat{x}_i[k+1] = \begin{cases} x_i[k+1], & \text{if SETF triggers,} \\ \hat{x}_i[k], & \text{otherwise.} \end{cases}$$

Shortcomings:

- Only **bounded resilient consensus** can be achieved due to the existence of constant term c_0 , i.e.,

$$\lim_{k \rightarrow \infty} |x_i[k] - x_j[k]| \leq c, \quad \forall k \in \mathbb{Z}_{\geq 0}.$$
- The threshold **cannot dynamically adjust** as the error term changes.

[4] Y. Wang and H. Ishii, “Resilient consensus through event-based communication,” IEEE Trans. Control Netw. Syst., vol. 7, no. 1, pp. 471–482, 2020.

Dynamic Event-Triggered Mechanism

- The control protocol is also modified as

$$x_i[k+1] = x_i[k] + \sum_{j \in \mathcal{R}_i^{\text{in}}[k]} a_{ij}[k] (\hat{x}_j[k] - x_i[k])$$

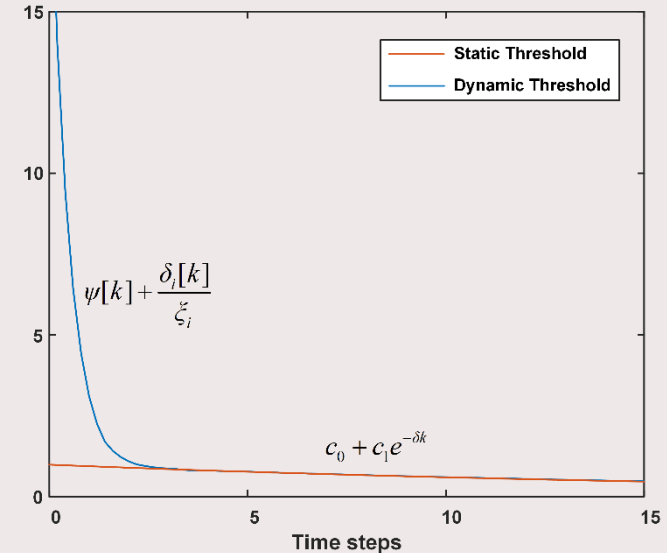
- A dynamic variable $\delta_i[k]$ is introduced, whose state update follows

$$\delta_i[k+1] = (1 - \theta_i) \delta_i[k] + \eta_i (\psi[k] - |e_i[k]|)$$

Always greater than zero and converges to zero.

- The dynamic event-triggered function (DETF) is designed as

$$f_i[k] = |e_i[k]| - \left(\psi[k] + \frac{\delta_i[k]}{\xi_i} \right)$$



Comparison between the static and dynamic event-triggered mechanism.

Dynamic Event-triggered Mean Subsequence Reduced (DE-MSR) Algorithm^[5]



1. Sorting



2. Deleting



3&5. Updating



4. Checking

4. Agent i checks whether DETF triggers or not and sets $\hat{x}_i[k+1]$ as

$$\hat{x}_i[k+1] = \begin{cases} x_i[k+1], & \text{if DETF triggers,} \\ \hat{x}_i[k], & \text{otherwise.} \end{cases}$$

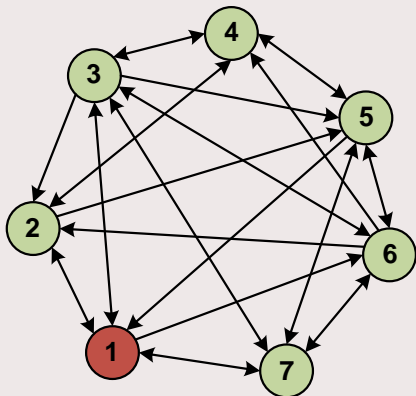
5. Agent i updates its interval dynamical variable according to

$$\delta_i[k+1] = (1 - \theta_i)\delta_i[k] + \eta_i(\psi[k] - |e_i[k]|)$$

[5] **Z. Liao**, J. Shi, S. Wang, Y. Zhang, and Z. Sun, “Resilient consensus through dynamic event-triggered mechanism,” IEEE Trans. Circuits Syst. II-Express Briefs, 2024, doi: 10.1109/TCSII.2024.3364524.

Comparative example between SETF and DETF

Communication topology



Simulation Setting

- Agent 1 is a **malicious agent**, whose motion follows

$$x_1[k] = 5 \times \sin(k / 5) + 7$$
- Initial state values** of all agents are denoted as

$$[x_1[0], \dots, x_7[0]] = [7, 7, 6, 8, -5, 8, 6]$$
- SETF related parameters:

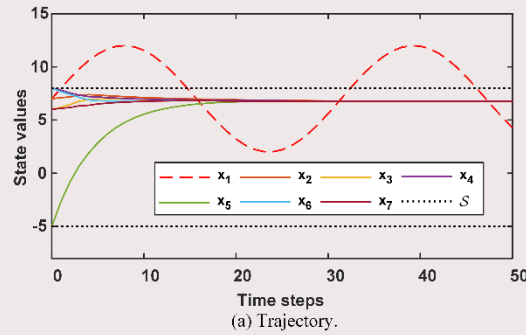
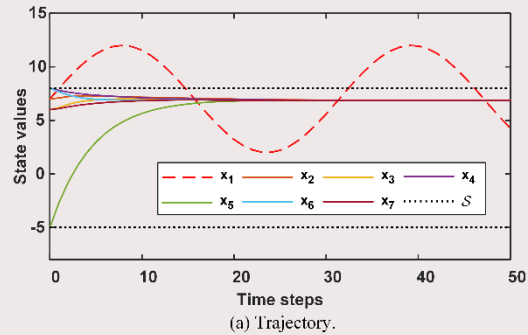
$$c_0 = 5 \times 10^{-4}, c_1 = 0.01, \alpha = 0.01$$
- DETF related parameters:

$$\zeta_i = 10, \eta_i = 0.6, \theta_i = 0.35, \delta_i[0] = 15, \psi[k] = e^{-\beta k}, \beta = 0.01$$

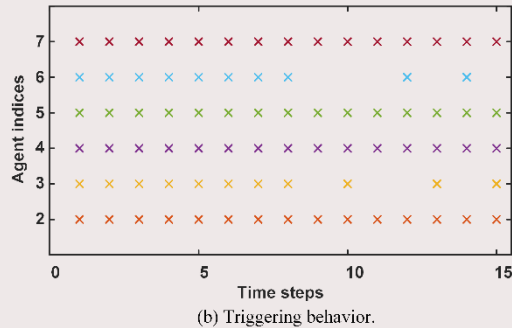
Purpose

- To validate the **effectiveness** of the proposed DE-MSR algorithm.
- To show the **superiority** of the DETF compared with the SETF.

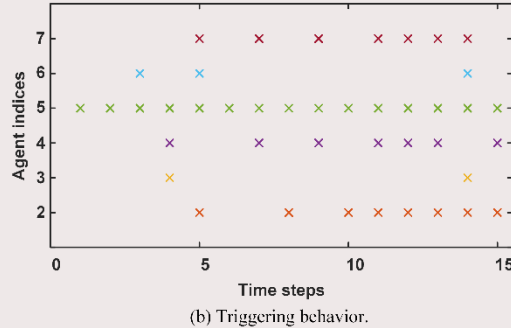
Simulation Results



The SETF and DETF show a **similar convergence rate**.



• SETF



• DETF

The DETF is **superior in reducing communication overheads**.

Event-triggered mechanism	Event counts (within 250 time steps)					
	ag.2	ag.3	ag.4	ag.5	ag.6	ag.7
SETF	33	24	29	26	18	22
DETF	26	15	20	26	10	14

Conclusion

- Study the resilient consensus problem based on the dynamic event-triggered mechanism.
- Propose the DE-MSR algorithm to reduce communication overheads.
- Validate the effectiveness of the proposed method through numerical simulations.

Future Work

- MASs with disturbances generated from the external environment.
- More general and complex tasks (e.g., resilient distributed optimization).



Thanks for Listening!