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Dynamic Event-Triggering Resilient Consensus for Heterogeneous Multi-Agent Systems (MASs) Against Malicious Attacks

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- **1. Research Background and Significance**
- 2. Literature Review
- 3. Methodology
- 4. Numerical Example
- 5. Conclusion and Future Work

1. Research Background and Significance

➤ Research Background



Swarm Performance

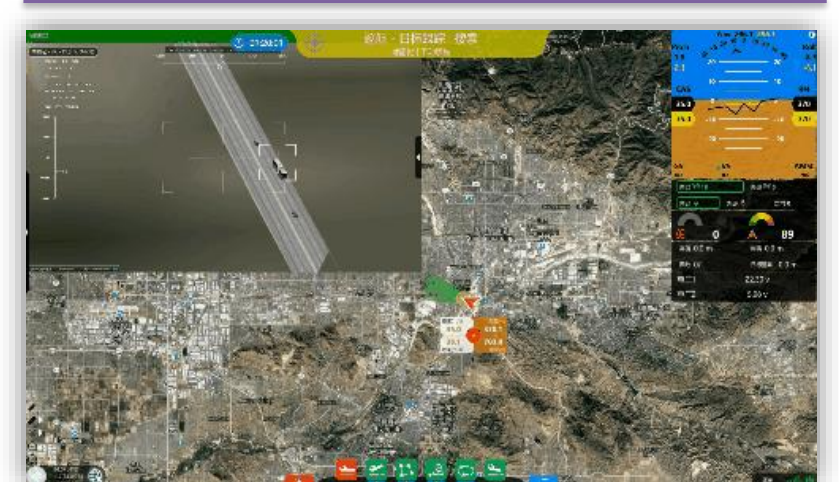


Agriculture



Cooperative Detection

Application of MASs in Various Fields



Target Tracking

1. Research Background and Significance

➤ Research Significance

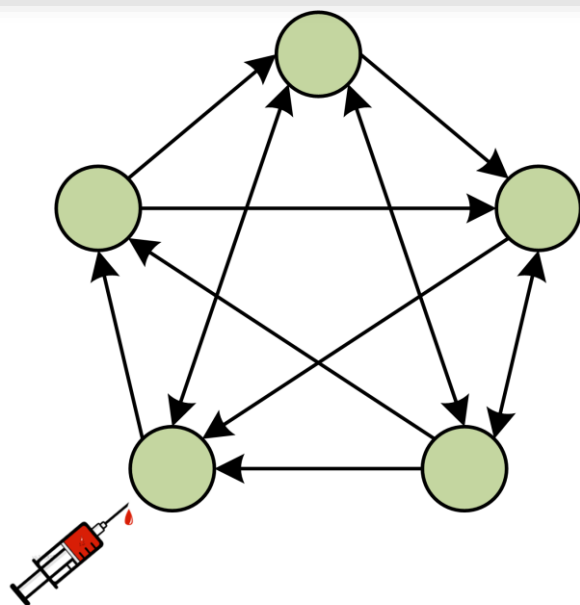
The swarm operates normally



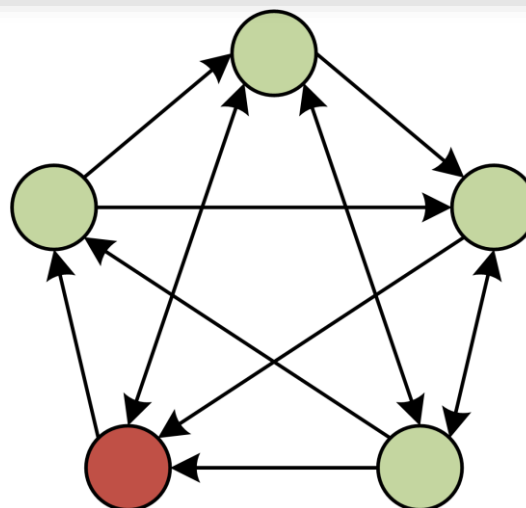
Some drones are under attack



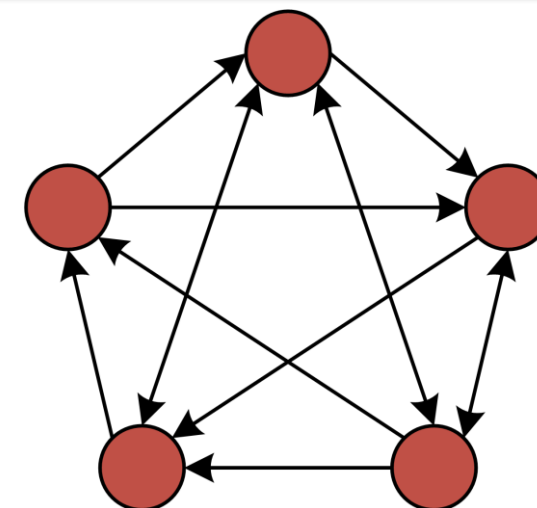
Swarm collapse



Malicious
Attack



Attack
Spread



1. Research Background and Significance



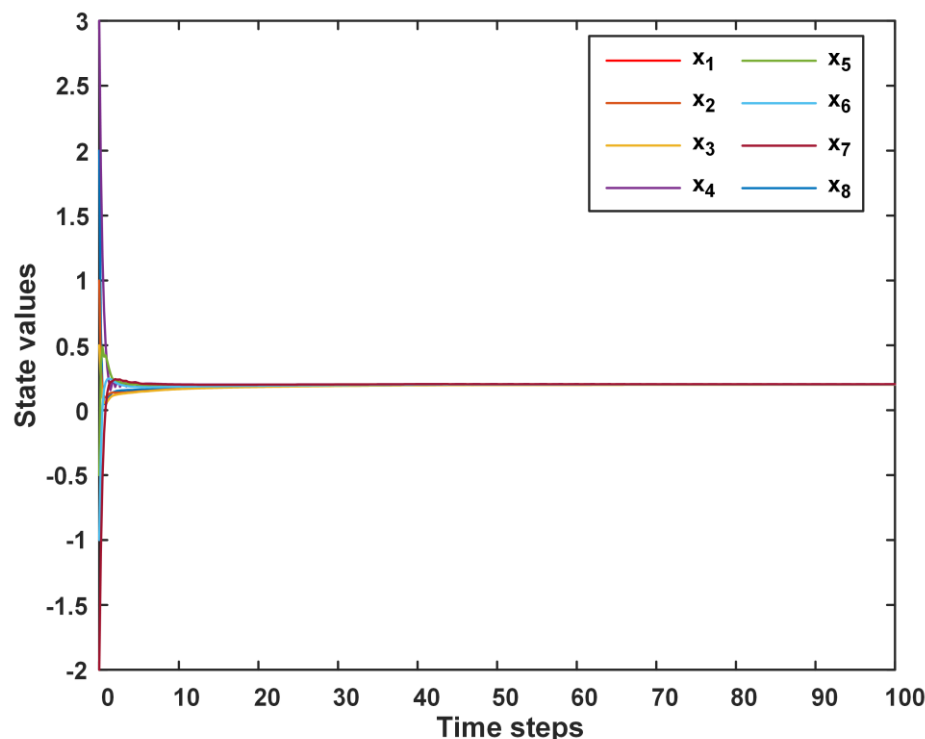
➤ Research Significance

- Consider a single-integrator MAS modelled by

$$x_i[k+1] = x_i[k] + u_i[k]$$

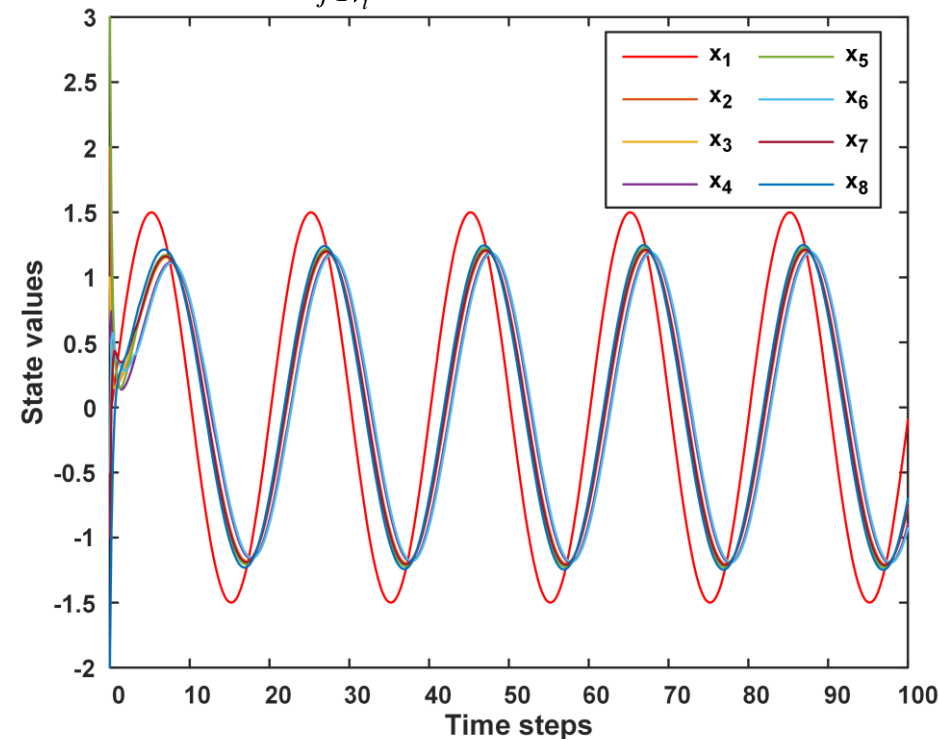
- A consensus-seeking protocol^[1] is introduced as

$$u_i[k] = \sum_{j \in \mathcal{V}_i} a_{ij}[k] (x_j[k] - x_i[k])$$



Without Malicious attack

Malicious
Attack
➔



Agent 1 is subject to malicious attack

Motivating Question: How to address resilient consensus problem when the MAS is under attack?

[1] Ren W, Cao L. Distributed coordination of multi-agent networks: emergent problems, models, and issues.

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2. Literature Review

➤ Review of Resilient Control Methods for MASs

Attack Detection and Isolation Method ^[2,3]	Attack Tolerant Method ^[4,5]
Detecting and isolating malicious agents	Deleting malicious state values
Agents are equipped with a bank of observers	More lightweight with less computational complexity
Massive data processing	Less data processing
High network connectivity required	Less prior information required
Unknown input observers	Mean Subsequent reduced (MSR) algorithm

[2] I. Shames, A. M. Teixeira, H. Sandberg, and K. H. Johansson, “Distributed fault detection for interconnected second-order systems,” *Automatica*, vol. 47, no. 12, pp. 2757–2764, 2011.

[3] Z. Tang, M. Kuijper, M. S. Chong, I. Mareels, and C. Leckie, “Linear system security-detection and correction of adversarial sensor attacks in the noise-free case,” *Automatica*, vol. 101, pp. 53–59, 2019.

[4] H. J. LeBlanc, H. Zhang, X. Koutsoukos, and S. Sundaram, “Resilient asymptotic consensus in robust networks,” *IEEE Journal on Selected Areas in Communications*, vol. 31, no. 4, pp. 766–781, 2013.

[5] Y. Zhai, Z. Liu, M. Ge, G. Wen, X. Yu, and Y. Qin, “Trusted-region subsequence reduction for designing resilient consensus algorithms,” *IEEE Transactions on Network Science and Engineering*, vol. 8, no. 1, pp. 259–268, 2020.

2. Literature Review

➤ Existing Resilient Control Methods based on MSR

Literature	Year	System Model	Communication Mode	Objective
[4]	2013	Single-integrator MASs	Time-triggered	Exact resilient consensus
[6]	2017	Double-integrator MASs	Time-triggered	Exact resilient consensus
[7]	2022	General linear MASs	Time-triggered	Exact resilient consensus
[8]	2020	Single-integrator MASs	Static event-triggered	Bounded resilient consensus
[9]	2021	Single-integrator MASs	Static event-triggered	Bounded resilient consensus
Our research	2024	Heterogeneous MASs	Dynamic event-triggered	Exact resilient consensus

[4] H. J. LeBlanc, H. Zhang, X. Koutsoukos, and S. Sundaram, “Resilient asymptotic consensus in robust networks,” *IEEE Journal on Selected Areas in Communications*, vol. 31, no. 4, pp. 766–781, 2013.

[6] S. M. Dibaji and H. Ishii, “Resilient consensus of second-order agent networks: Asynchronous update rules with delays,” *Automatica*, vol. 81, pp. 123–132, 2017.

[7] Y. Bai and J. Wang, “Resilient consensus of continuous-time linear networked systems,” *IEEE Transactions on Circuits and Systems II: Express Briefs*, vol. 69, no. 8, pp. 3500–3504, 2022.

[8] Y. Wang and H. Ishii, “Resilient consensus through event-based communication,” *IEEE Transactions on Control of Network Systems*, vol. 7, no. 1, pp. 471–482, 2020.

[9] N. Wang, and Y. Liu, “Resilient consensus of second-order multi-agent systems subject to malicious attacks,” in *Proc. of the 2021 Chinese Intelligent Systems Conference*. Springer, 2021, pp. 10-21.

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➤ Dynamic Modelling for Heterogeneous MASs

- Consider a heterogeneous MASs consisting of **first-order** and **second-order agents**.
- The **discrete system dynamics** are mathematically expressed as follows:

◆ Dynamic equations:

$$\begin{aligned} \dot{x}_i(t) &= u_i(t), & i \in \mathcal{N}_1 \\ \begin{cases} \dot{x}_i(t) = v_i(t) \\ \dot{v}_i(t) = u_i(t) \end{cases}, & i \in \mathcal{N}_2 \end{aligned}$$

Sampling Period T



Discretization

$$x_i[(k+1)T] = x_i[kT] + Tu_i[kT], \quad i \in \mathcal{N}_1$$

$$\begin{cases} x_i[(k+1)T] = x_i[kT] + Tv_i[kT] + \frac{T^2}{2}u_i[kT], \\ v_i[(k+1)T] = v_i[kT] + Tu_i[kT] \end{cases}, \quad i \in \mathcal{N}_2$$



Replace kT with k

$$x_i[k+1] = x_i[k] + Tu_i[k], \quad i \in \mathcal{N}_1$$

$$\begin{cases} x_i[k+1] = x_i[k] + Tv_i[k] + \frac{T^2}{2}u_i[k], \\ v_i[k+1] = v_i[k] + Tu_i[k] \end{cases}, \quad i \in \mathcal{N}_2$$

◆ Symbol meaning:

$x_i(t) \in \mathbb{R}$: Position of agent i at time t

$v_i(t) \in \mathbb{R}$: Velocity of agent i at time t

$u_i(t) \in \mathbb{R}$: Input of agent i at time t

➤ Problem Statement

The MAS is said to achieve **exact resilient consensus** if the following two conditions hold for any initial state values of agents:

◆ **Resilience:** For each normal agent, it holds $x_i[k] \in \mathcal{S}$, $\forall k \in \mathbb{Z}_+$, where \mathcal{S} denotes the safety interval.

◆ **Exact Consensus:** For each normal agent, it holds $\lim_{k \rightarrow \infty} |x_i[k] - x_j[k]| = 0$, $\forall i, j \in \mathcal{N}$.
For each normal second-order agent, it holds $\lim_{k \rightarrow \infty} v_l[k] = 0$, $\forall l \in \mathcal{N}_2$.

Research Question: How to design control protocols to protect the heterogeneous MAS from malicious attacks and achieve resilient consensus?

➤ Control Protocol Design

- For normal first-order agents, the control protocol relies on the **difference** between neighbor's positions and their own positions.
- For normal first-order agents, the control protocol considers both the **difference** and their own **velocities**.

Control protocol for
heterogeneous MAS

$$u_i[k] = \begin{cases} \gamma \sum_{j \in \mathcal{R}_i[k]} a_{ij}[k] (x_j[k] - x_i[k]), & i \in \mathcal{N}_1 \\ \alpha \sum_{j \in \mathcal{R}_i[k]} a_{ij}[k] (x_j[k] - x_i[k]) - \beta v_i[k], & i \in \mathcal{N}_2 \end{cases}$$

Compared with the
control protocol in [1]

$$u_i[k] = \sum_{j \in \mathcal{V}_i} a_{ij}[k] (x_j[k] - x_i[k]) \quad \longrightarrow \quad u_i[k] = \gamma \sum_{j \in \mathcal{R}_i[k]} a_{ij}[k] (x_j[k] - x_i[k])$$

Control gain

regulate the system
state more accurately

Neighbor set

Do not use all neighbor
positions for update

3. Methodology

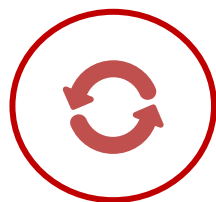
➤ MSR Algorithm



1. Sorting

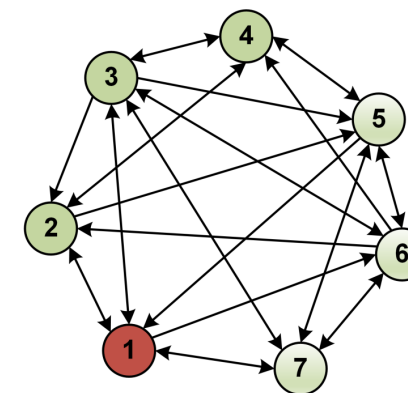


2. Deleting



3. Updating

➤ Example:



$$[x_1[0], \dots, x_7[0]] = [10, 8, 6, 7, 4, -5, 6]$$

$$[v_5[0], v_6[0], v_7[0]] = [3, -2, -1], \quad f = 2$$

Agent 2: $x_2[0] = 8, \quad x_j[0] = [10, 6, 7, -5]$



1. Sorting

$$S = [-5, 6, 7, 8, 10]^T$$



2. Deleting

$$S_{del} = [-5, 6, 7, 8, 10]^T, \quad \mathcal{R}_i[0] = \{4\}$$

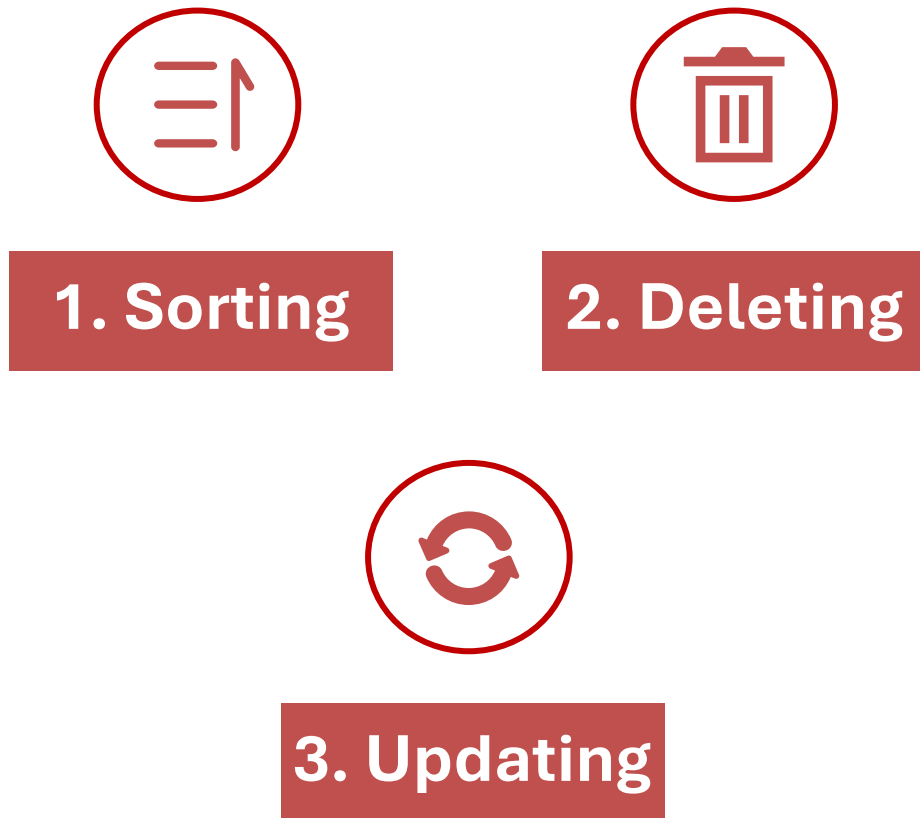
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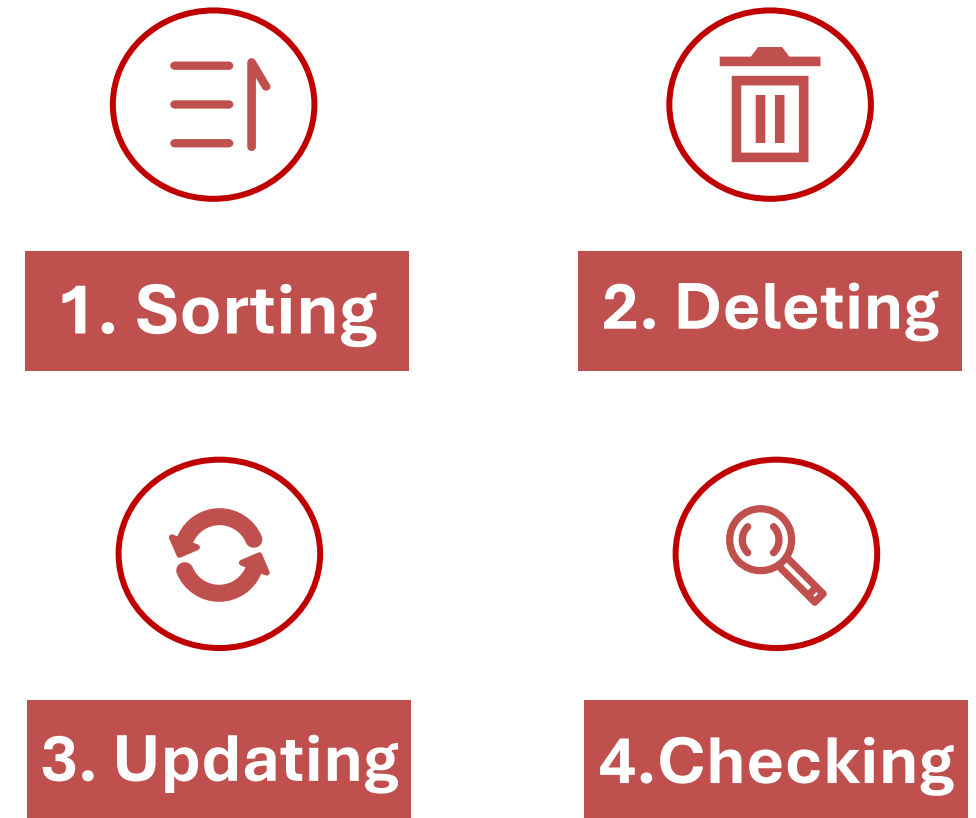
3. Updating

$$x_i[1] = x_i[0] + \sum_{j \in \mathcal{R}_i[0]} a_{ij}[0] (x_j[0] - x_i[0])$$

➤ MSR Algorithm



➤ Event-based MSR (E-MSR) Algorithm^[8]



Require $x_j[k]$ at
each time step



Heavy
communication load

[8] Y. Wang and H. Ishii, “Resilient consensus through event-based communication,” *IEEE Transactions on Control of Network Systems*, vol. 7, no. 1, pp. 471–482, 2020.

➤ Static Event-Triggered Mechanism^[8]

- The control protocol is also modified as

$$u_i[k] = \begin{cases} \gamma \sum_{j \in \mathcal{R}_i[k]} a_{ij}[k] (\hat{x}_j[k] - x_i[k]), & i \in \mathcal{N}_1 \\ \alpha \sum_{j \in \mathcal{R}_i[k]} a_{ij}[k] (\hat{x}_j[k] - x_i[k]) - \beta v_i[k], & i \in \mathcal{N}_2 \end{cases}$$

Auxiliary variable

- The update of auxiliary variable depends on

$$\hat{x}_i[k+1] = \begin{cases} x_i[k+1], & \text{if } f_i[k] > 0, \\ \hat{x}_i[k], & \text{otherwise.} \end{cases}$$

- The static event-triggered function (SETF) is designed as

$$f_i[k] = \underbrace{|x_i[k+1] - \hat{x}_i[k]|}_{\text{Error term } |e_i[k]|} - \underbrace{(c_0 + c_1 e^{-\alpha k})}_{\text{Threshold}}$$

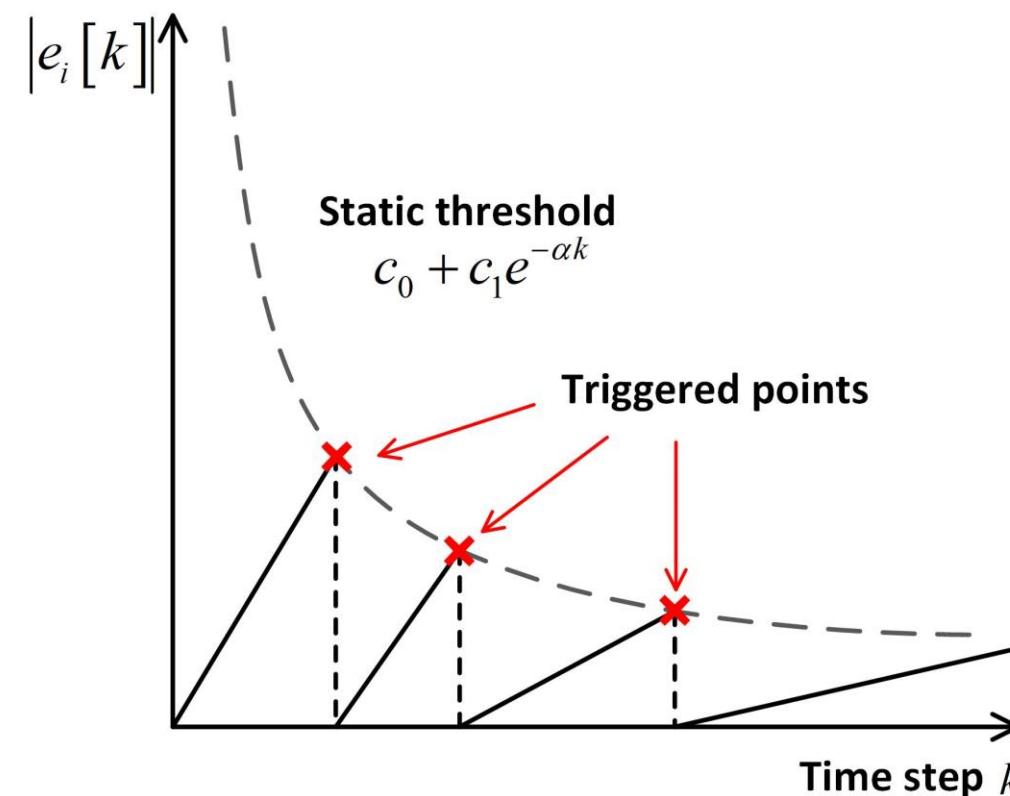
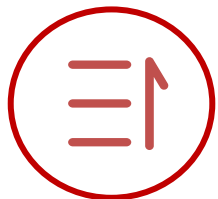


Illustration of the static event-triggered mechanism.

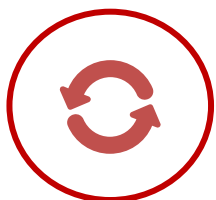
➤ E-MSR Algorithm



1. Sorting



2. Deleting



3. Updating



4. Checking

Step 4: Agent i checks whether SETF triggers or not and sets $\hat{x}_i[k+1]$ as

$$\hat{x}_i[k+1] = \begin{cases} x_i[k+1], & \text{if SETF triggers,} \\ \hat{x}_i[k], & \text{otherwise.} \end{cases}$$

Shortcomings:

- Only **bounded resilient consensus** can be achieved due to the existence of constant term c_0 , i.e.,

$$\lim_{k \rightarrow \infty} |x_i[k] - x_j[k]| \leq c, \quad \forall k \in \mathbb{Z}_{\geq 0}.$$

- The threshold **cannot dynamically adjust** as the error term changes.

➤ Dynamic Event-Triggering Mechanism^[10]

- The control protocol is also modified as

$$u_i[k] = \begin{cases} \gamma \sum_{j \in \mathcal{R}_i[k]} a_{ij}[k] (\hat{x}_j[k] - x_i[k]), & i \in \mathcal{N}_1 \\ \alpha \sum_{j \in \mathcal{R}_i[k]} a_{ij}[k] (\hat{x}_j[k] - x_i[k]) - \beta v_i[k], & i \in \mathcal{N}_2 \end{cases}$$

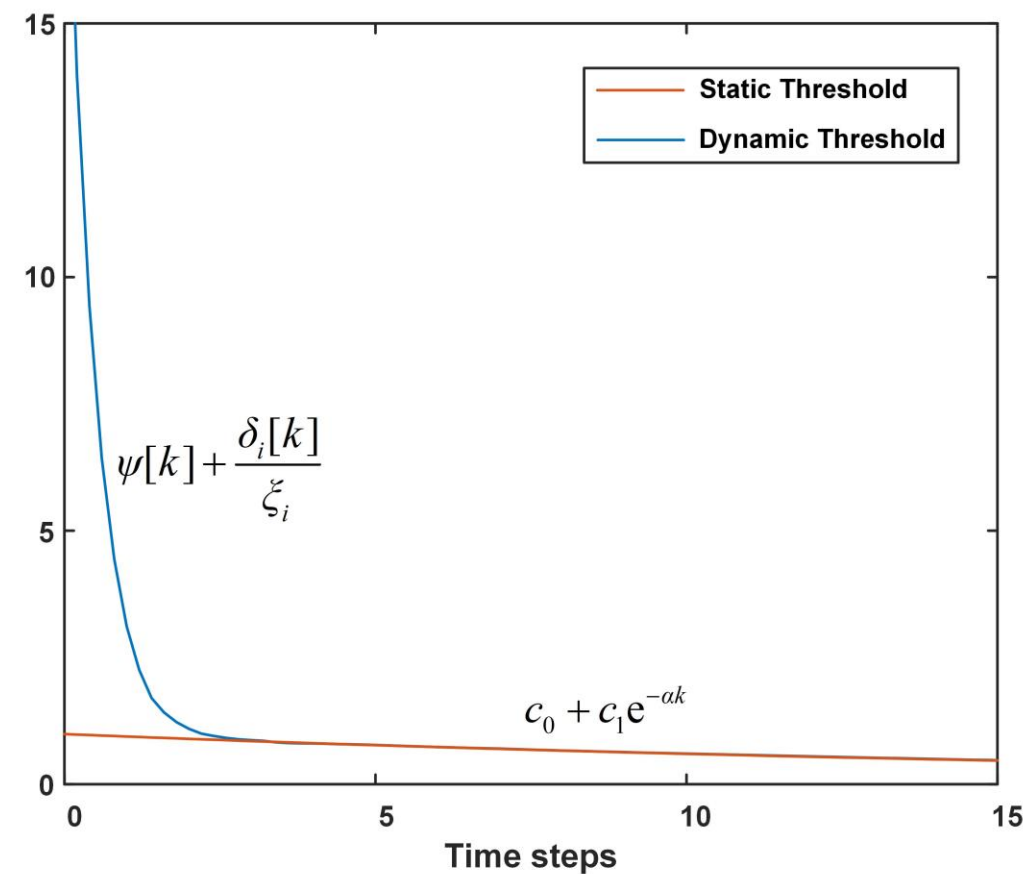
- A dynamic variable $\delta_i[k]$ is introduced, whose state update follow

$$\delta_i[k+1] = (1 - \theta_i) \delta_i[k] + \eta_i (\psi[k] - |e_i[k]|)$$

Always greater than zero and exponentially converges to zero.

- The dynamic event-triggered function (DETF) is designed as

$$f_i[k] = |e_i[k]| - \left(\psi[k] + \frac{\delta_i[k]}{\xi_i} \right)$$



Comparison of the two triggering thresholds.

➤ Dynamic Event-triggered Mean Subsequence Reduced (DE-MSR) Algorithm^[10]



1. Sorting



2. Deleting



3&5. Updating



4. Checking

Step 4: Agent i checks whether DETF triggers or not and sets $\hat{x}_i[k+1]$ as

$$\hat{x}_i[k+1] = \begin{cases} x_i[k+1], & \text{if SETF triggers,} \\ \hat{x}_i[k], & \text{otherwise.} \end{cases}$$

Step 5: Agent i updates its interval dynamic variable according to

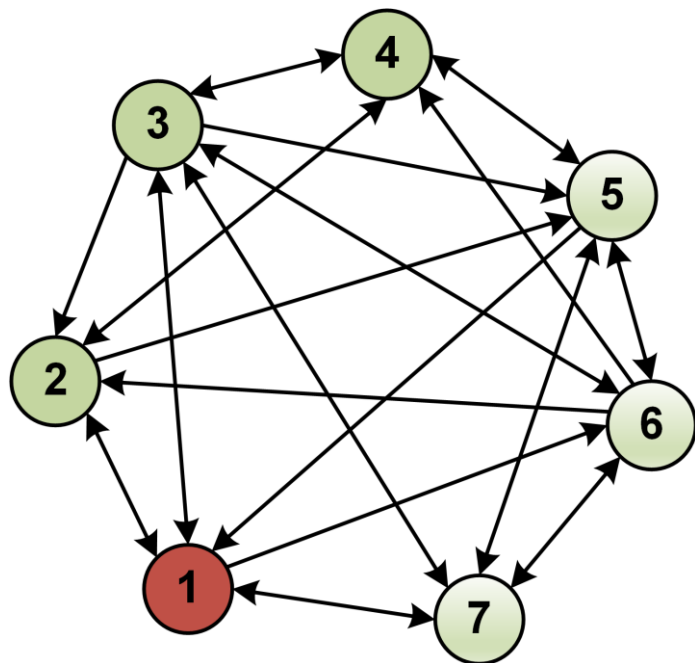
$$\delta_i[k+1] = (1 - \theta_i)\delta_i[k] + \eta_i(\psi[k] - |e_i[k]|)$$

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➤ A Comparative Numerical Example Between SETF and DETF

◆ Communication Topology



◆ Simulation Setting

- Agents 1,2,3,4: first-order; Agents 5,6,7: second-order

- Agent 1 is a **malicious agent**, whose motion follows

$$x_1[k] = 5 \times \sin(0.2 \times k) + 5$$

- Initial state values** of all agents are denoted as

$$X[0] = [7, 7, 6, 8, -5, 8, 6], \quad V[0] = [0, 0, 0, 0, 3, -2, -1]$$

- SETF related parameters:

$$c_0 = 0.005, \quad c_1 = 0.01, \quad \alpha = 0.01$$

- DETF related parameters:

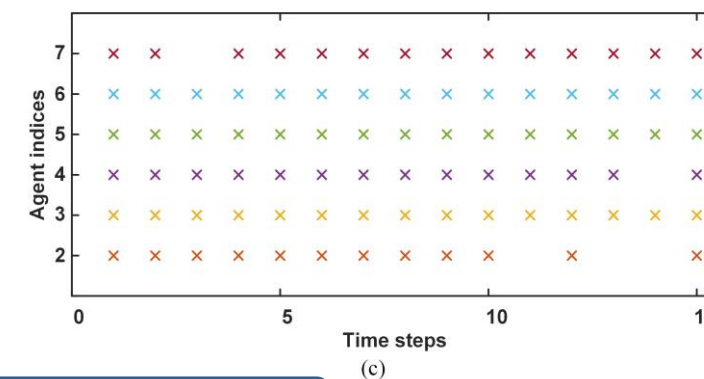
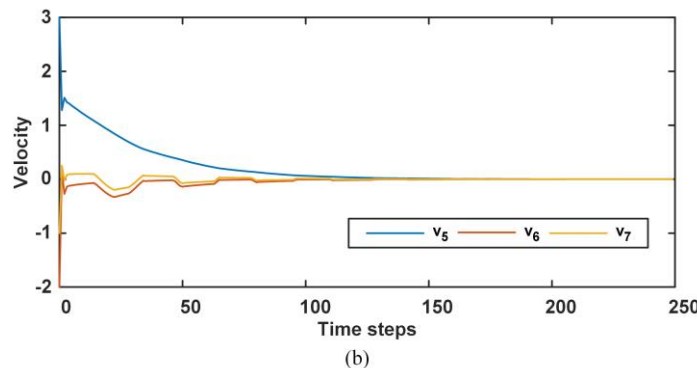
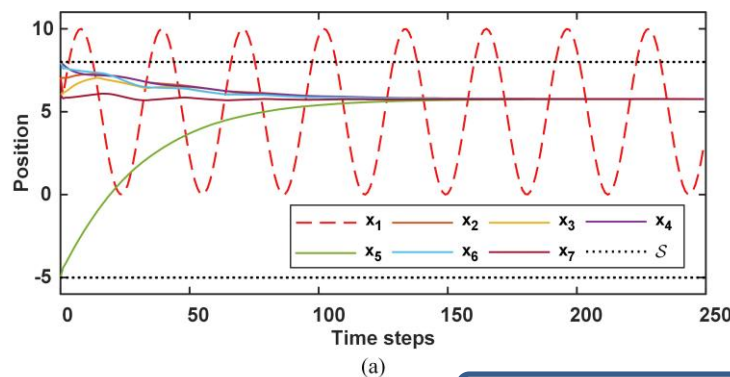
$$\xi_i = 10, \quad \eta_i = 0.6, \quad \theta_i = 0.35, \quad \delta_i[0] = 15, \quad \psi[k] = e^{-\beta k}, \quad \beta = 0.01$$

◆ Purpose

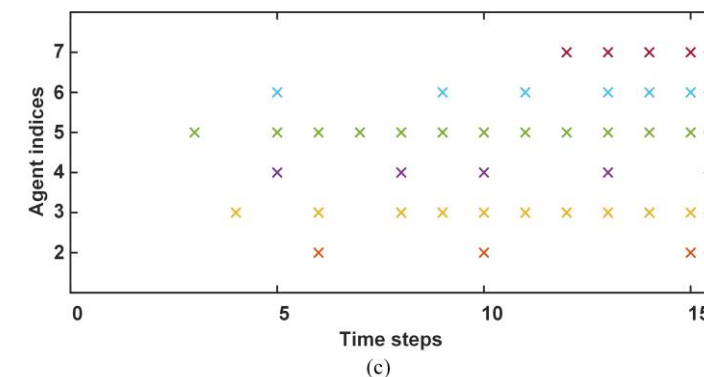
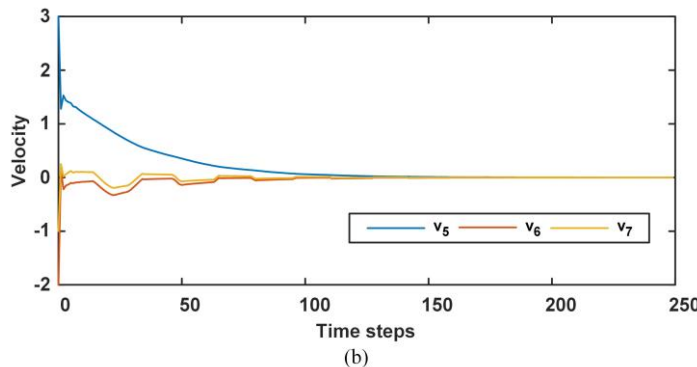
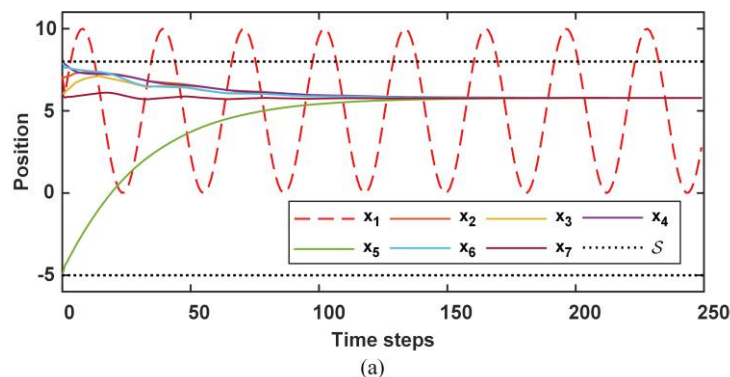
- To validate the **effectiveness** of the proposed DE-MSR algorithm.
- To show the **superiority** of DETF compared with SETF.

4. Numerical Example

➤ Simulation Results



Simulation Results with SETF: (a) Position; (b) Velocity; (c) Triggering behavior.



Simulation Results with DETF: (a) Position; (b) Velocity; (c) Triggering behavior.

Event-triggering mechanism	Event counts (within 250 time steps)					
	ag.2	ag.3	ag.4	ag.5	ag.6	ag.7
SETF	150	127	149	178	113	102
DETF	140	122	138	174	106	90



- The SETF and DETF show a **similar convergence rate**.
- The DETF is superior in **reducing communication times**.

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5. Conclusion and Future Work

➤ Conclusion

- Study the resilient consensus problem for **heterogeneous MASs**.
- Propose the **DE-MSR algorithm** to reduce communication overheads.
- Validate the **effectiveness** of the proposed method through numerical simulations.

➤ Prospection

Future work includes:

- MASs with **disturbances** generated from the external environment.
- More **general and complex tasks** (e.g., resilient distributed optimization).

Thank you so much for your attention!

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Full Paper



WeChat

