Problem	Points	Lost	Gained	Running Total	TA
1	1 (+1)				
2	18				
3	22				
4	20				

Name:_____GT Number: ____

You may	ask	for	clar	ific	cation	but	you	are	ulti	mately	respons	sible	for	the
answer	you	write	on	the	paper	. I:	E you	ı mak	ce an	v assu	mptions	state	the	em.

Please look through the entire test before starting. WE MEAN IT!!!

NOTE: $M = 10^6$ K = 10^3 Mi = 2^{20} Ki = 2^{10} Illegible answers are wrong answers.

20 19

100

Show your work in the space provided to get any credit for problem-oriented questions.

Good luck!

1. (1 point, 1 min) (don't worry you get 1 point regardless of your answer!
You get 1 extra credit point if you get it right!!)

Given:

5

6 Total

> 5+3+2 = 151022 9+2+4 = 183652 8+6+3 = 4824665+4+5 = 202541

What is 7+2+5 = _____

Name:	_GT Number:			
Cache				
2.(18 points, 20 min)				
a) (6 points) Given:		C0	C1	
 8 total cache blocks 				[

- 2-way set-associative organization
- Cache initially empty
- LRU replacement policy
- Memory blocks L, M, N and O all map to the **same** cache line.
- The processor performs a total of eight accesses, to memory blocks O, N, M, L, O, O, N, and M in that order. For each of these accesses, specify (by filling in the table below) whether it is a cache hit or a cache miss, type of miss (cold/capacity/conflict), and the memory block evicted (if any). Note: capacity miss dominates over conflict; cold dominates over capacity.

Memory Access	Hit/miss	Type of miss	Block evicted from cache
0			
N			
M			
L			
0			
0			
N			
M			_

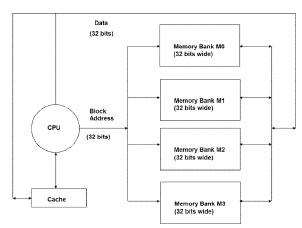
(Area for rough work)

- b) (2 points) In a 4-way set associative cache with 64K bytes of data, 64 bytes per block and with a t-bit tag (circle the right choice)
 - (i) There are four t-bit tag comparators
 - (ii) There are 64 t-bit tag comparators
 - (iii) There are 1K t-bit tag comparators
 - (iv) There is one t-bit tag comparator for the entire cache

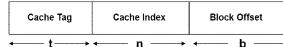
Name: GT Number:

- c) (2 points) Virtually indexed and physical tagged cache is attractive because (circle the correct choice)
 - (i) it results in a better cache hit ratio for a given cache organization
 - (ii) it eliminates the memory aliasing problem with physically indexed physically tagged caches
 - (iii) it enables building bigger first level caches than physically indexed and physically tagged caches
 - (iv) it removes the address translation through the TLB out of the critical path of the cache access
- d) (4 points) Consider a
 - 4-way interleaved memory system
 - DRAM access time: 70 cycles
 - bus cycle time for address or data transfer from/to the CPU/memory: 5 cycles
 - Block size: 4 words (each 32 bits)

Compute the block transfer time.



e) (4 points) Given



Let \mathbf{a} be the number of bits in the memory address, \mathbf{S} be the total size of the cache in bytes, and \mathbf{B} the block size in bytes. Assuming a direct-mapped cache, answer the following:

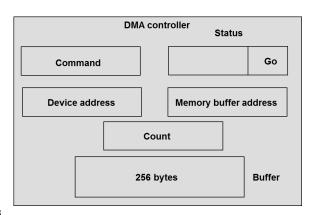
- (i) b = ____ (ii) n = ____
- (iii) t = _____

Name: GT Number:	
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Input/Output and Disk

- 3. (22 points, 20 min)
- a) (2 points) To transfer N bytes of data from a memory buffer at address M to the device at address D, the CPU executes the following instructions to the DMA controller:
 - 1. Store N into the Count register.
 - 2. Store M into the Memory buffer address register.
 - 3. Set the Go bit in the Status register.
 - 4. Store D into the Device address register.
 - 5. Store "write to the device" command into the Command register.

What (if any) is wrong with the above code sequence?



- b) (2 points) What is the role of the buffer in a DMA controller? (select ALL that apply: +1 for correct choice; -1 for incorrect choice)
 - (i) Asynchronously transfer data to/from the device
 - (ii) Synchronously transfer data to/from the device
 - (iii) Asynchronously transfer data to/from the memory
 - (iv) Synchronously transfer data to/from the memory
- c) (2 points) Memory mapped I/O is a ... (circle the right choice)
 - (i) technique for interfacing slow speed devices to the CPU
 - (ii) technique for interfacing high speed devices to the CPU
 - (iii) technique that allows the CPU to use Load/Store instructions to access the device registers
 - (iv) technique that allows the CPU to quickly find the location of the handler code for a device

Na	ıme:_	GT Number:
d)	Given	the following specification for a disk drive: Average seek time = 4 ms Rotational speed = 6000 RPM Platters = 3 Surface per platter = 2 Tracks per surface = 5000 Sectors per track = 1024 Recording density = 256 bytes per sector
	(i) (2 disk?	2 point)How much time is needed to get to a random sector on the
		(2 point) How much time is needed to read one random sector from the when the head is already positioned on the desired sector?
		(2 points) If the disk gets a request to read 6 random sectors, how total time will that request take to complete?
		(2 points) If the disk gets a request to read 6 consecutive sectors, ach total time will that request take to complete?
	(v) (2	2 points) What is the transfer rate of the disk?

Name:	_GT Number:
e) (4 points) Consider a disk using parameters:	g zoned-bit recording with the following
p - number of platters, n - number of surfaces per planets, z - number of zones, t_{zi} - number of tracks at a s_{zi} - number of sectors per b - number of bytes per s	zone z_i , z_i track at zone z_i ,
Give an expression for the total	l capacity of the disk.
f) (1 point) Which of the following starvation? (circle all the right	g disk scheduling algorithms are prone to nt choices)
(i) FCFS	
(ii) SCAN	
(iii) SSTF	
(iv) LOOK	
(1 point) Explain why.	
File Systems (see last page for a control of the current state of a file "foo" in the	is as follows: O Apr 27 21:01 foo ands:
(i) -rwxrx	
(ii) -rwxr-x	
(iii) -rwxrwx	
(iv) -rwxrw-	

Name:	GT Number:
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b) (8 points) In the following table, assume none of the files exist to start with in the current directory. Fill in the table. The reference count in the table pertains to the i-node that is affected by the command in that row. If a new i-node is created, show the old reference count for that i-node as 0.

Command	New i-node created	Reference	ce count
	(yes/no)	old	new
touch f1			
touch f2			
ln f2 f3			
ln -s f1 f4			
ln f4 f5			
ln -s f4 f6			
rm f6			
rm f3			

Use this area for rough work for this question

- c) (2 points) Hybrid indexed allocation results in (circle ONE choice that captures ALL the TRUE statements in the following list)
 - (i) External Fragmentation
 - (ii) Internal fragmentation
 - (iii) Ability to grow the file easily
 - (iv) Ability to handle small files efficiently
 - (v) {i and ii}
 - (vi) {ii and iii}
 - (vii) {ii, iii, and iv}
 - (viii) None of the above

Name:	GT Number:	
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d) (**NOTE Mi = 2^{20}**)

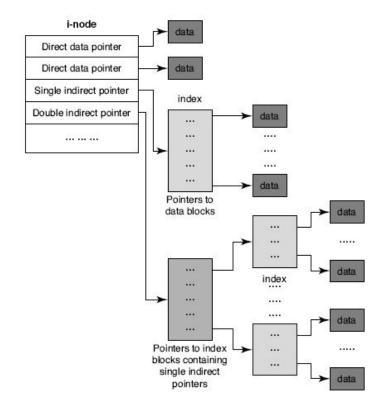
Given the following:

Size of index block = 128 bytes Size of Data block = 2048 bytes Size of pointer = 8 bytes (to index or data blocks)

The i-node consists of

- 2 direct data block pointers,
- 1 single indirect pointer, and
- 1 double indirect pointer.

Note that the index blocks and data blocks are allocated on a need basis. An index block is used for the top-level i-node as well as for the index blocks that store pointers to other index blocks and data blocks (see Figure).



- i. (2 points) How many pointers does each index block contain?
- ii. (2 points) How many data blocks are used to store a 40 KiB file?
- iii. (2 points) How many index blocks (including the i-node for the file)
 are needed to store a 256 KiB file?

iv. (2 points) What is the largest file size that can be supported in this file system?

Name:	GT Number:
Parallel Systems	
threads are scheduled on its thread-id and termina	re are four threads in a process. Assume that the a uniprocessor. Once spawned, each thread printates. Assuming a non-preemptive thread scheduler tions are possible? (circle the right choice)
(i) 4	
(ii) 1	
(iii) 16	
(iv) 4!	
(1 point) Explain why.	
	re two threads and their execution history on an cution of each instruction is atomic. Assume tha
Thread 1 (T1) Time 0: R1 <- Mem[x] Time 2: R1 <- R1+2 Time 4: Mem[x] <- R1	Thread 2 (T2) Time 1: R2 <- Mem[x] Time 3: R2 <- R2+1 Time 5: Mem[x] <- R2
what is the final value in m	memory location \mathbf{x} ?
(i) 0	
(ii) 1	
(iii) 2	
(iv) 3	
(2 point) Explain why.	

Name:GT Number:
c) (5 points) The internal representation in the thread library for a LOCK is shown below:
Name of lock Who has it Who are waiting
Assume that the following events happen in the order shown (T1-T5 are threads of the same process): • T1 executes thread_mutex_lock(L1); • T2 executes thread_mutex_lock(L1); • T3 executes thread_mutex_lock(L2); • T4 executes thread_mutex_lock(L2); • T5 executes thread_mutex_lock(L1);
Assuming that there have been no other calls to the threads library prior to this, show the state of the internal data structures of L1 and L2 after the above five calls (by filling in the figures below).
L1
L2
d) (4 points) Shown in the figure below are the points of execution (indicated by the arrows) of four threads (T1-T4) of the same process. Assume that the critical sections are mutually exclusive (i.e., they are governed by the same lock).
T1 T2 T3 T4

For each thread, state if it is active or blocked, and why.

Т1	is	because
		,
1.5	is	because
Т3	is	because
т4	is	hecause

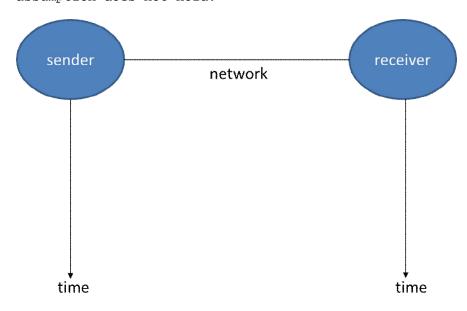
Name:	GT Number:
e) (2 points) Deadlock (circle	the right choice)
(i) Is a condition where the	reads are not using mutex locks
(ii) Is a condition where al	ll the locks variables are in use
(iii) A lock variable that i	is dead
(iv) Is a condition where or that will never happen	ne or more threads are waiting for an event
shared variables. Each of t	g code, buflock, bufavail, and frame_buf are the functions (digitizer and tracker) are reads. What is the problem (if any) with the
<pre>digitizer() { image_type dig_image; int tail = 0;</pre>	tracker() { image_type track_image; int head = 0;
<pre>loop { grab(dig_image); thread_mutex_lock(buflock); while (bufavail == 0) do nothing; thread_mutex_unlock(buflock); frame_buf[tail mod MAX] =</pre>	<pre>loop { thread_mutex_lock(buflock); while (bufavail == MAX) do nothing; thread_mutex_unlock(buflock); track_image = frame_buf[head mod</pre>
}	}

Name: GT Number:	Name:	GT Number:
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Networking

- 6. (19 points, 20 min)
- a) (3 points) The one-bit sequence number requirement for the "Stop-and-wait protocol (aka alternating bit protocol)" is predicated on the following assumption:
 - packets do not get reordered in the network OR arbitrarily delayed

Show using the space-time diagram below what can go wrong if this assumption does not hold.



b) (2 points) Why you need MAC address when each node on the Internet has an IP address?

Name:GT Number:									
c) (2 points) The sequence number in a packet (circle the right ch							cle the right choice)		
(i) Gives the destination address									
(ii) Is needed for message reconstruction at the destination							estination		
	(iii) Assures the integrity of the packet								
(iv) Is computed using cyclic redundancy check (CRC) algorithm							C) algorithm		
	(v) Is the same for every packet in a given message							2	
d)	(2 points) In t	he sli	ding v	window	protoc	col,		
	(i) when	is the	e windo	w size	e decre	eased?			
	(ii) when is the window size increased?								
e)	e) (4 points) Show the wave form for the following packet with Manchester encoding (the space between the vertical lines represent time per bit):								
	1	1	1	0	0	1	0	1	

Nar	ne:	GT Number:
f) (•	ints) Given: Message size = 100,000 bytes; Header size per packet = 100 bytes Packet size = 1100 bytes Packet loss = 10%
		umber of packets to transmit the message (ignoring fractional packet: (circle the correct choice)
(i)	100
(ii)	91
(iii)	1000
(iv)	111
(-	v)	110
(-	vi)	Cannot be determined with the data given
g) (de: soi nui sec	ints) A packet header consists of the following fields: stination_address urce_address m_packets quence_number cket_size ecksum
Assu	me t	hat each of these fields occupies 4 bytes. Given that the packet siz

ze is 1500 bytes, compute the payload in each packet.

h) (2 points)

Given the following:

Compute the throughput.

Name:GT Number:	
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Unix command	Semantics	Elaboration
touch <name></name>	Create a file with the name <name></name>	Creates a zero byte file with the name <name> and a creation time equal to the current wall clock time</name>
mkdir <sub-dir></sub-dir>	Create a sub-directory <sub-dir></sub-dir>	The user must have write privilege to the current working directory (if <sub-dir> is a relative name) to be able to successfully execute this command</sub-dir>
rm <name></name>	Remove (or delete) the file named <name></name>	Only the owner of the file (and/or superuser) can delete a file
rmdir <sub-dir></sub-dir>	Remove (or delete) the sub-directory named <sub-dir></sub-dir>	Only the owner of the <sub-dir> (and/or the superuse) can remove the named sub-directory</sub-dir>
ln -s <orig> <new></new></orig>	Create a name <new> and make it symbolically equivalent to the file <orig></orig></new>	This is name equivalence only; so if the file <orig> is deleted, the storage associated with <orig> is reclaimed, and hence <new> will be a dangling reference to a non-existent file</new></orig></orig>
ln <orig> <new></new></orig>	Create a name <new> and make it physically equivalent to the file <orig></orig></new>	Even if the file <orig> is deleted, the physical file remains accessible via the name <new></new></orig>
chmod <rights> <name></name></rights>	Change the access rights for the file <name> as specified in the mask <rights></rights></name>	Only the owner of the file (and/or the superuser) can change the access rights
chown <user> <name></name></user>	Change the owner of the file <name> to be <user></user></name>	Only superuser can change the ownership of a file
chgrp <group> <name></name></group>	Change the group associated with the file <name> to be <group></group></name>	Only the owner of the file (and/or the superuser) can change the group associated with a file
cp <orig> <new></new></orig>	Create a new file <new> that is a copy of the file <orig></orig></new>	The copy is created in the same directory if <new> is a file name; if <new> is a directory name, then a copy with the same name <orig> is created in the directory <new></new></orig></new></new>
mv <orig> <new></new></orig>	Renames the file <orig> with the name <new></new></orig>	Renaming happens in the same directory if <new> is a file name; if <new> is a directory name, then the file <orig> is moved into the directory <new> preserving its name <orig></orig></new></orig></new></new>
cat/more/less <name></name>	View the file contents	