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Enriching a Verified Choreographic Language with a Simply Typed Lambda Calculus

— Honours project (S1/S2 2024)

A thesis submitted for the degree
Bachelor of Advanced Computing (Research and Development)

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Declaration:

I declare that this work:

- upholds the principles of academic integrity, as defined in the University Academic Misconduct Rules;
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October, Xin Lu

Acknowledgements

If you wish to do so, you can include some Acknowledgements here. If you don't want to, just comment out the line where this file is included.

There is absolutely no need to write an Acknowledgement section, so only do so when you *want* to – it's always important to stay sincere. One reason for including an acknowledgement could be to thank your supervisor for extraordinary supervision (or any other reason you deem noteworthy). Some supervisors sacrifice a lot, e.g., are always available, meet on weekends, provide multiple rounds of corrections for theses reports, or the like (keep in mind that writing a thesis is special for you, but not for them, so they do actually not have any reason to sacrifice their private time for this!). Seeing acknowledgements in this report can feel like a nice appreciation of this voluntary effort. For large works that form the end of some studies (like an Honours or Master thesis), it is also not uncommon to read acknowledgements to one's parents or partner. But again, completely optional!

Abstract

- choreography diagram; CC, Kalas - the meaning of deadlock freedom by design - mostly focus on interactions via message passing - contribution 1: richerLang: call by value, functional big step semantics with clock to be implemented in HOL4, richer data types, environment semantics; an environmental language model with type theory; and strong normalisation property - contribution 2: the enriched choreography, with a simple type theory, Kalas have the safety property of deadlock freedom

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Chapter 1

Introduction

Distributed systems consist of multiple endpoints that communicate by exchanging messages, operating with asynchrony and parallelism as these messages are sent and received between the various endpoints. But programming distributed system is notoriously error-prone as programmer has to implement the communication protocol by developing individual endpoint programs. Mismatched message sending and receiving can lead to errors such as *deadlock*, where the system is waiting forever for a message.

Choreography arises as a programming diagram to address this issue by providing a concrete global description of how the messages are exchanged between endpoints in a distributed system. A choreography program is written in a similar style to the "Alice and Bob" notation by Needham and Schroeder (1978):

1. $\texttt{Alice} \to \texttt{Bob} : \mathit{key}$
2. $\texttt{Bob} \to \texttt{Alice} : \mathit{message}$

Thus message mismatches are disallowed from the choreographic perspective. A property we refer to as *deadlock free by design*. The global choreography is then projected into process models for each endpoint via EndPoint Projection (EPP), with properties such as deadlock free by design preserved (Hallal et al., 2018).

While most choreography languages focus on the message exchange behaviours, few pay attention to the local computation happening in the individual endpoint. It is shown by Hirsch and Garg (2022) that as long as the local language exhibits type preservation and progress, the choreography inherits these properties as well. Thus, when analyzing message exchange behaviors in choreography—such as multiparty sessions, asynchrony, and parallelism—one can safely assume good behaviours of local computation (Montesi and Yoshida, 2013; Cruz-Filipe and Montesi, 2017; Carbone and Montesi, 2013). But when it comes to writing the choreography program that implements concrete system behaviours, if local computation is ever required by the system, one must describe the

inputs and outputs of local computations. Current work either delegate this part to an assumed well-behaved external implementation, for example, Kalas Pohjola et al. (2022) and Pirouette Hirsch and Garg (2022), or provide a basic framework where only natural numbers are considered, as in Core Choreography (CC) Cruz-Filipe and Montesi (2017). This makes writing choregraphy program with local computation implementation an unpleasant experience. For example, a choreography program where client performs local computation based on the input from server and then send its result back in Kalas will look like:

```
1. server.var -> client.x;
2. let v@client = fun(x) in
3. client.v -> server.result;
```

The processes in Kalas only communicate with strings, so in the external implementation of `fun`(x) one has to convert the string value stored in x to its desired data type first, and perform the computation based on the converted value. The result is then converted back to string before updating the client process variable v.

Thus this thesis extends Kalas, a verified choregraphy language with machine-checked end-to-end compilation, with a simple language of expressions over types such as integers, strings and booleans. We ensure that the extended Kalas has important properties such as type preservation and progress. Using the extended Kalas, the previous choregraphy where client computes the factorial of input integer can be written as:

```
1. server.var -> client.x;
2. let v@client = case NumOfx of in
3. client.v -> server.result;
4. client.v -> server.result;
5. client.v -> server.result;
```

This thesis provides three main contributions.

- The first contribution is an environment language with functional big step semantics. We also provide a typing system. The language is implemented using the HOL4 proof assistant Slind and Norrish (2008). We give type soundness proof for the proposed semantics and typing rules.
- The second contribution is the strong normalisation proof for our environment language model. It is essentially the strong normalisation proof for a simply-typed lambda calculus, but we define the logic relation based on environment semantics.
- The third contribution is the enriched Kalas. Besides common data types such as integer, string, and boolean, we also add function, pair, and sum types to the choregraphy. Common operators for our data types are included, such as addition, modulo, and negation. Integer-string convertor is implemented as well for the integration with Kalas. Last but not least, we prove the enriched Kalas exhibits type preservation and progress using HOL4.

Chapter 2

Background

2.1 Choreography as a Programming Diagram

- how choregraphy can be applied real examples - choreography semantics and basic results ? - a choregraphy diagram

2.2 Interactive Theorem Proving

- this work is done in HOL4 so maybe start from this and think how to introduce stuff
- Kalas is kind of both a choreography language and implemented in HOL

2.3 Kalas

Chapter 3

Related Work

- we discuss in terms of how local computation is considered and how important properties such as type preservbationa dn progress is influenced by those decisions

3.1 Choreography Models

3.1.1 Core Models

Choreography language model can ensure deadlock freedom without using a typing system. The property is typically proven through structural induction on the choreography's semantics, where an applicable rule always enables reduction, or reduction follows by inductive hypothesis. Typing system in this case might still be desired since it can discipline choreography in other ways such as correct protocol implementation or ensuring deadlock freedom for EPP.

3.1.2 Typing System for Choreography

Choreography language model can ensure deadlock freedom without using a typing system. The property is typically proven through structural induction on the choreography's semantics, where an applicable rule always enables reduction, or reduction follows by inductive hypothesis. Typing system in this case might still be desired since it can discipline choreography in other ways such as correct protocol implementation or ensuring deadlock freedom for EPP.

When the local computation is involved in choreography, since we cannot solve the halting problem in semantics by deciding whether the local computation will terminate, a typing system for the choreography is desired to re-establish the deadlock freedom property. To the best of our knowledge, the closet work to this discussion is Pirouette by Hirsch and Garg (2022). They assume a substitution model with small-step semantics for

the local language, and based on a set of admissible typing rules for the sake of providing a reasoning ground, their results show that type preservation and progress property for local language implies the preservation and progress for the choreography language adopting it. Their progress result aligns with our results, but our local language adopts big-step semantics and thus preservation for choreography requires strong normalisation from local language rather than type preservation.

Pirouette is a higher-order functional choreography language, where it has three forms of values: local value, local function that takes local value inputs and outputs a choregraphy, and global function that takes choregraphy and outputs choregraphy. This makes choregraphy has return value on which its typing system is built on. On the other hand, since neither Kalas or the enriched Kalas has return values, our typing system mainly checks for the well-formedness of a choreography within the typing environment and if any local computation is involved, we discipline it with its own typing system in a localised typing environment. Pirouette types its local values in a similar projected typing environment, which contains bindings from variables to types in a specific location.

There is also an interesting observation between Pirouette and Kalas in terms of how asynchronous behaviours are controlled. Pirouette uses a set of blocked locations to disallow a branching from current choregraphy in which multiple attempts to write to one target local variable may occur in different orders, whereas Kalas uses an action label to record previous communication to disallow a branching where the target local variable is appear in the label.

Another state-of-art functional choregraphy language Chorλ by **?** uses a different approach than Pirouette, where choreographies are interpreted as terms in λ-calculus. They are not concerned with how the local values are computed and assume the presence of a local value for communications happening in the choregraphy. This contributes to a very different typing system than Pirouette: their local environment is not obtained from global environment by any means of projection, instead local value types are annotated with the associated roles and are embedded in types bound to choregraphy names in a typing environment. Type preservation and progress follow typing and semantics rules of Chorλ.

- compositonal choregraphy ..

3.1.3 Undecided Title But There Will Be One

Other choreography languages ...

- (deadlock freedom/ type soundness) typing of choreography - CC, Kalas, CC M do not have typing system, deadlock freedom for closed choreography idirectly follows the semantics, where the local computation is always assumed terminated - When deadlock freedom does not directly follows from the choreography semantics, such as the compositional choreography proposed in (Yoshida), typing system is used to ensure progress - others, Channel Choreographies (ChC) in (Fabrizio Montesi thesis) uses typing system

for checking protocal Correctness

- local procedure in choreography

- exceptions

3.2 Undecided Title But There Will Be One

- functional big-step semantics

- SN for STLC; strong normalisation for languages of environmental semantics

Chapter 4

The Environment Model

4.1 Syntax

- binary operators - unary operators (StrOf, NumOf) - value script?

4.2 Semantics

- functional big step semantics: intepreter style with clock, total function. for being implemented in HOL4 - properties: clock increment (cases on result) (and clock decrement) - closure: issues with dynamic environment, so we use lexical environment; do we explain restrciting to fv(e) and how? - exceptions

4.3 Typing

4.3.1 Syntax

- rich data types: int, string, boolean, sum type, pair type, closure, ... - typecheck - uoptype, boptype - valuetype - envtype

4.3.2 Typing Rules

- typecheck - uoptype, boptype - valuetype: closure - no type uniqueness

4.3.3 Main Properties

- value invertability; envtype used for fnT case of it - operators: uoptype soundness and boptype soundness (use the invertability) - typecheck: reducing typing environment (for soundness fn case)

4.3.4 Type Soundness

Chapter 5

Strong Normalisation

Chapter 6

The Enriched Choreography

- exceptions

Chapter 7

Concluding Remarks

If you wish, you may also name that section *"Conclusion and Future Work"*, though it might not be a perfect choice to have a section named "A & B" if it has subsections "A" and "B". Also note that you don't necessarily have to use these subsections; that also depends on how much content you have in each. (E.g., having a section header might be odd if it contains just three lines.)

7.1 Conclusion

This section usually summarizes the entire paper including the conclusions drawn, e.g., did the developed techniques work? Maybe add why or why not. Also don't hold back on limitations of your work; it shows that you understood what you have done. And science isn't about claiming how great something is, but about objectively testing hypotheses. Also note that every single scientific paper has such a section, so you can check out many examples, preferably at top-tier venues, e.g., by your supervisor(s).

7.2 Future Work

- asynchronus messages; confluence property - Progress for EPP

Chapter 8

Test

We define what it is for a choreograph to be well-formed with the $G, Th \vdash c\Box$ relation.

This is a theorem:

$$\vdash \emptyset, \Theta \vdash c\Box \Rightarrow \exists \tau\ l\ s'\ c'.\ \emptyset \rhd c \xrightarrow[l]{\tau} s' \rhd c' \vee \neg \texttt{not\_finish}\ c$$

The transition relation looks like `eval_exp` *clk E exp*

Appendix A

Appendix: Explanation on Appendices

You may use appendices to provide additional information that is in principle relevant to your work, though you don't want *every reader* to look at the entire material, but only those interested.

There are many cases where an appendix may make sense. For example:

- You developed various variants of some algorithm, but you only describe one of them in the main body, since the different variants are not that different.
- You may have conducted an extensive empirical analysis, yet you don't want to provide *all* results. So you focus on the most relevant results in the main body of your work to get the message across. Yet you present the remaining and complete results here for the more interested reader.
- You developed a model of some sort. In your work, you explained an excerpt of the model. You also used mathematical syntax for this. Here, you can (if you wish) provide the actual model as you provided it in probably some textfile. Note that you don't have to do this, as artifacts can be submitted separately. Consult your supervisor in such a case.
- You could also provide a list of figures and/or list of tables in here (via the commands `\listoffigures` and `\listoftables`, respectively). Do this only if you think that this is beneficial for your work. If you want to include it, you can of course also provide it right after the table of contents. You might want to make this dependent on how many people you think are interested in this.

Appendix B

Appendix: Explanation on Page Borders

What you find here is an explanation of why the border width keeps flipping from left to right – which you might have spotted and wondered why that's the case.

Firstly, that is *intended* and thus correct, so there is no reason to worry about this. The reason is that this document is configured as a two-sided book, which means:

- We assume the document will be printed out,
- that this will be done in a two-sided mode (i.e., the document will be printed on both sides of each page), and
- that the bookbinding will be in the middle, just like in every book.

When you open the book, there are three borders of equal size n. This however requires that even pages have a border of n on their left and $\frac{n}{2}$ on their right, and odd pages have a border of $\frac{n}{2}$ on their left and n on their right. This is illustrated in Figure B.1.

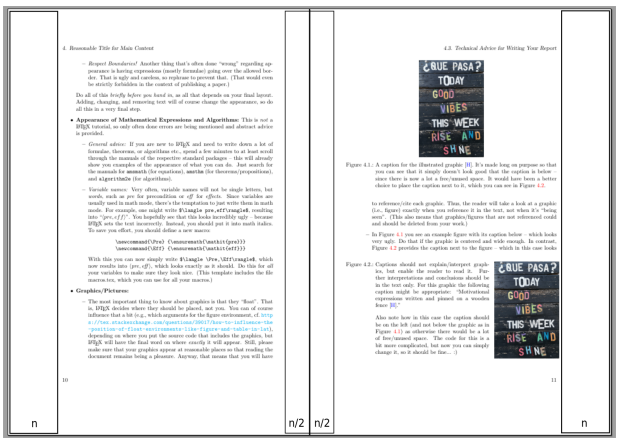

Figure B.1: Illustration showing why page borders flip.

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