

Problem 1 (20pts). Consider the following linear program:

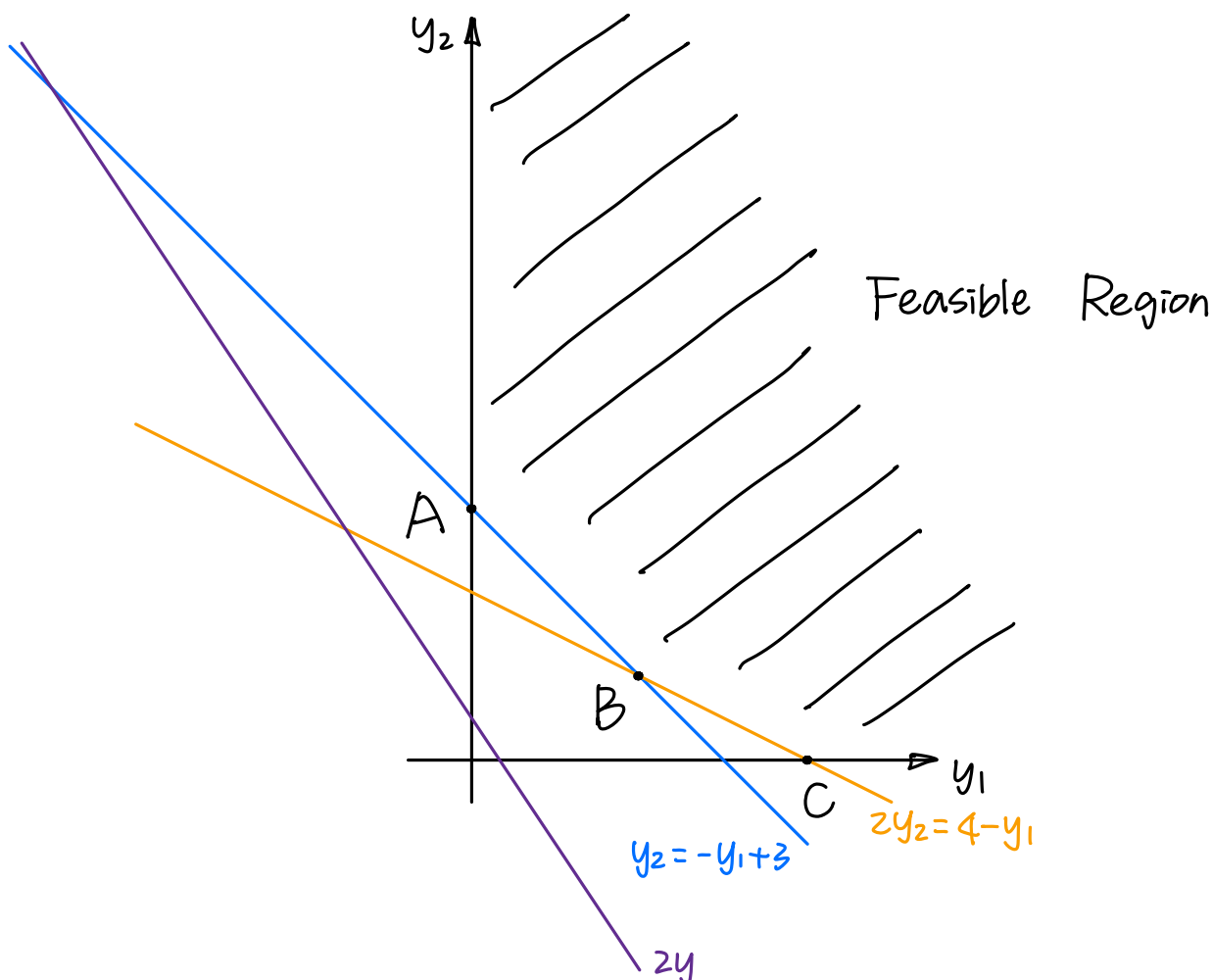
$$\begin{aligned} & \text{maximize} && 3x_1 + x_2 + 4x_3 \\ & \text{subject to} && x_1 + 3x_2 + x_3 \leq 5 \\ & && x_1 + 2x_2 + 2x_3 \leq 8 \\ & && x_1, x_2, x_3 \geq 0 \end{aligned}$$

- What is the corresponding dual problem?
- Solve the dual problem graphically.
- Use complementarity conditions for the primal-dual pair to solve the primal problem.

(a) The corresponding dual problem is:

$$\begin{aligned} & \min && 5y_1 + 8y_2 \\ & \text{s.t.} && y_1 + y_2 \geq 3 \\ & && 3y_1 + 2y_2 \geq 1 \\ & && y_1 + 2y_2 \geq 4 \\ & && y_1, y_2 \geq 0 \end{aligned}$$

(b) The graph is as followed:



Three extreme points $A(0, 3)$; $B(2, 1)$; $C(4, 0)$

$$A(0, 3): 5y_1 + 8y_2 = 24$$

$$B(2,1): 5y_1 + 8y_2 = 18$$

$$C(4,0): 5y_1 + 8y_2 = 20$$

So the optimal value is 18.

The optimal solution of the dual problem is $(2,1)$.

(c) By the complementarity condition, we have:

$$x_1(3 - y_1 - y_2) = 0$$

$$x_2(1 - 3y_1 - 2y_2) = 0$$

$$x_3(4 - y_1 - 2y_2) = 0$$

$$y_1(5 - x_1 - 3x_2 - x_3) = 0$$

$$y_2(8 - x_1 - 2x_2 - 2x_3) = 0$$

By (b), we already know that both the primal and dual have optimal finite value, and $y_1 = 2, y_2 = 1$.

Solving these equations, we have optimal solutions:

$$\text{primal problem: } (x_1, x_2, x_3) = (2, 0, 3)$$

$$\text{dual problem: } (y_1, y_2) = (2, 1)$$

Problem 2 (25pts). Consider the following table of food and corresponding nutritional values:

	Protein, g	Carbohydrates, g	Calories	Cost
Bread	4	7	130	3
Milk	6	10	120	4
Fish	20	0	150	8
Potato	1	30	70	2

The ideal intake for an adult is at least 30 grams of protein, 40 grams of carbohydrates, and 400 calories per day. The problem is to find the **least** costly way to achieve those amounts of nutrition by using the four types of food shown in the table.

- Formulate this problem as a linear optimization problem (specify the meaning of each decision variable and constraint).
- Solve it using MATLAB, find an optimal solution and the optimal value.
- Formulate the dual problem. Interpret the dual problem. (Hint: Suppose a pharmaceutical company produces each of the nutrients in pill form and sells them each for a certain price.)
- Use MATLAB to solve the dual problem. Find an optimal solution and the optimal value.

(a) Denote the number of bread, milk, fish and potato separately as x_1, x_2, x_3, x_4 .

$$\min \quad 3x_1 + 4x_2 + 8x_3 + 2x_4$$

$$\text{s.t.} \quad 4x_1 + 6x_2 + 20x_3 + x_4 \geq 30$$

$$7x_1 + 10x_2 + 30x_4 \geq 40$$

$$130X_1 + 120X_2 + 150X_3 + 70X_4 \geq 400$$

$$X_1, X_2, X_3, X_4 \geq 0$$

(b) The Code is as followed:

```

编辑器 - D:\大二上\MAT3007\assignment\homework4\HW4_2.m
HW4_2.m  x +
1      clear;
2      clc;
3
4      cvx_begin
5          variables x1 x2 x3 x4
6          minimize 3*x1 + 4*x2 + 8*x3 + 2*x4
7          subject to
8              4*x1 + 6*x2 + 20*x3 + x4 >= 30;
9              7*x1 + 10*x2 + 30*x4 >= 40;
10             130*x1 + 120*x2 + 150*x3 + 70*x4 >= 400;
11             x1 >= 0;
12             x2 >= 0;
13             x3 >= 0;
14             x4 >= 0;
15             cvx_end
16
17             % print outcome
18             x1
19             x2
20             x3
21             x4
22             3*x1 + 4*x2 + 8*x3 + 2*x4

```

outcome:

```

命令行窗口
x1 =

    1.0732

x2 =

    1.1560e-08

x3 =

    1.2312

x4 =

    1.0829

ans =

    15.2351

fx >>

```

(c) Denote y_1, y_2, y_3 as the number of pills that separately contain protein, Carbohydrates and Calories.

Each pill is sold with certain price.

The dual problem is:

$$\begin{aligned}
 \max \quad & 30y_1 + 40y_2 + 400y_3 \\
 \text{s.t.} \quad & 4y_1 + 7y_2 + 130y_3 \leq 3 \\
 & 6y_1 + 10y_2 + 120y_3 \leq 4 \\
 & 20y_1 + 150y_3 \leq 8 \\
 & y_1 + 30y_2 + 70y_3 \leq 2 \\
 & y_1, y_2, y_3 \geq 0
 \end{aligned}$$

(d) The Code is as followed:

```
25 % Homework4 Problem2 (d)
26 cvx_begin
27 variables y1 y2 y3
28 maximize 30*y1 + 40*y2 + 400*y3
29 subject to
30 4*y1 + 7*y2 + 130*y3 <= 3;
31 6*y1 + 10*y2 + 120*y3 <= 4;
32 20*y1 + 150*y3 <= 8;
33 y1 + 30*y2 + 70*y3 <= 2;
34 y1 >= 0;
35 y2 >= 0;
36 y3 >= 0;
37 cvx_end
38
39 % print outcome
40 y1
41 y2
42 y3
43 30*y1 + 40*y2 + 400*y3
```

outcome:

编辑器 - D:\大二上\MAT3007\assignment\homework4\HW4_2.m

命令行窗口

```
Status: Solved
Optimal value (cvx_optval): +15.2351
```

```
y1 =
    0.3099
```

```
y2 =
    0.0283
```

```
y3 =
    0.0120
```

```
ans =
    15.2351
```

fx>>

The optimal solution: $(x_1, x_2, x_3, x_4) = (1.0732, 1.1560 \times 10^{-8}, 1.2312, 1.0829)$
 $(y_1, y_2, y_3) = (0.3099, 0.0283, 0.0120)$

The optimal value: 15.23151

Problem 3 (25pts). Consider the max flow problem on the graph in figure 1 with the orange node being the source node and the green node being the terminal node (the number on each edge is its capacity, see the lecture slide 13). Do the following based on the lecture slides.

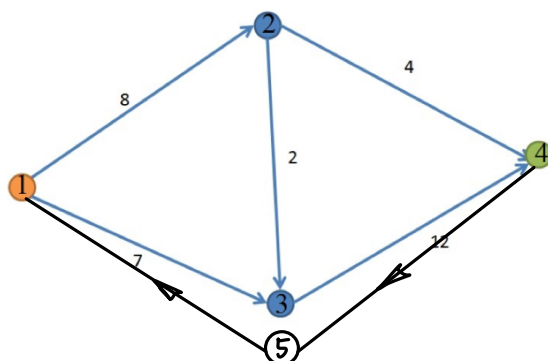


Figure 1: Max flow problem

- (a). Formulate it as a linear program and solve it using MATLAB.
 (b). Formulate the dual of this problem and solve it using MATLAB.
 (c). Find the corresponding maximum flow and minimum cut of the graph.(Please draw the cut on figure 1).

(a) Assume there is an imaginary node 5 with edges (5,1) and (4,5)

Denote the flow from 5 to 1 as m .

The flow from $1 \rightarrow 2$ as x_1 , $1 \rightarrow 3$ as x_2 , $2 \rightarrow 3$ as x_3 ,

$2 \rightarrow 4$ as x_4 , $3 \rightarrow 4$ as x_5

The problem is:

maximize Δ

x, Δ

s.t. $\sum_{j: (j,i) \in E} x_{ji} - \sum_{j: (i,j) \in E} x_{ij} = 0, \forall i \neq 1, 5$

$\sum_{j: (j,1) \in E} x_{j1} - \sum_{j: (1,j) \in E} x_{1j} + \Delta = 0$

$\sum_{j: (j,5) \in E} x_{j5} - \sum_{j: (5,j) \in E} x_{5j} - \Delta = 0$

$x_{ij} \leq w_{ij} \quad \forall (i,j) \in E$

$x_{ij} \geq 0$

The code is as followed:

```

4 % Homework4 Problem4 (a)
5 w=[0 8 7 0;
6     0 0 2 4;
7     0 0 0 12;
8     0 0 0 0];
9
10 cvx_begin
11     variable x(4,5);
12
13     maximize sum(x(:,5));
14     subject to
15         sum(x(2,1:4))-sum(x(1:4,2))==0;
16         sum(x(3,1:4))-sum(x(1:4,3))==0;
17         sum(x(1:4,1))-sum(x(1,1:4))+sum(x(:,5))==0;
18         sum(x(1:4,4))-sum(x(4,1:4))-sum(x(:,5))==0;
19         x>=0;
20         for i = 1:4
21             for j=1:4
22                 x(i,j)<=w(i,j)
23             end
24         end
25     cvx_end
26     x, sum(x(:,5))

```

Outcome: 命令行窗口

```

-----
Status: Solved
Optimal value (cvx_optval): +13

x =

    0.0000    6.0000    7.0000    0.0000    3.2500
    0.0000    0.0000    2.0000    4.0000    3.2500
    0.0000    0.0000    0.0000    9.0000    3.2500
    0.0000    0.0000    0.0000    0.0000    3.2500

```

```

ans =

    13.0000

```

fx >>

(b) The dual problem is:

$$\text{minimize } \sum_{(i,j) \in E} w_{ij} z_{ij}$$

$$\text{s.t. } z_{ij} \geq y_i - y_j$$

$$y_1 - y_4 = 1$$

$$z_{ij} \geq 0$$

The code is as followed:

```

29 % Homework4 Problem4 (b)
30 cvx_begin
31     variable z(4,4);
32     variable y(4);
33     minimize sum(sum(w.*z));
34     for i = 1:4
35         for j = 1:4
36             if w(i,j)~=0
37                 z(i,j) >= y(i)-y(j)
38             end
39         end
40     end
41     y(1)-y(4) == 1
42     z >= 0
43
44 cvx_end
45 z,y
46

```

Outcome:

命令行窗口

Status: Solved
Optimal value (cvx_optval): +13

z =

(1,2)	0.0000
(1,3)	1.0000
(2,3)	1.0000
(2,4)	1.0000
(3,4)	0.0000

y =

1.0000
1.0000
-0.0000
-0.0000

fx >>

(c)

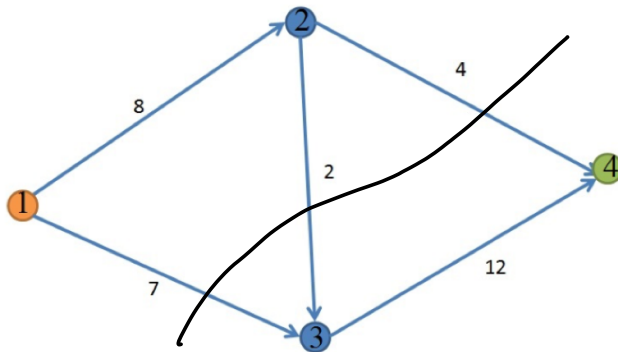


Figure 1: Max flow problem

Problem 4 (15pts). Use linear program duality to show that exactly one of the following systems has a solution

1. $Ax \leq b$

2. $y^T A = 0, b^T y < 0, y \geq 0$

Hint: You can first show that they can't both have solutions. Then you show that if the second one is infeasible, the first one must be feasible.

Primal Problem:

$$\begin{aligned} \min \quad & c^T x \\ \text{s.t.} \quad & Ax = b, x \geq 0 \end{aligned}$$

Dual Problem:

$$\begin{aligned} \max \quad & b^T y \\ \text{s.t.} \quad & A^T y \leq c \end{aligned}$$

In this problem, $c = 0$.

Argue by contradiction.

Suppose they can both have solutions.

$$\text{Then } 0 \cdot x = (y^T A) \cdot x = y^T (Ax) = y^T b = (b^T y)^T < 0$$

Since $0 \cdot x = 0$, the above statement is wrong.

\therefore They can't both have solutions.

If 2 is infeasible, then $y^T A \neq 0$

Suppose 1 is feasible.

$$y \cdot y^T A \cdot x \leq y \cdot y^T b = y \cdot (b^T y)^T < 0$$

$$\Rightarrow y y^T A x < 0$$

$$\text{Since } y > 0, y \cdot y^T = \sqrt{y^T} > 0$$

$$\Rightarrow Ax < 0$$

① If $b \geq 0$, then $Ax \leq b$ holds.

② If $b < 0$, then the solution to $Ax \leq b$ is a subset of the solution to $Ax \leq 0$

So $\exists x$ s.t. $Ax \leq b$ holds.

$\Rightarrow Ax \leq b$ is feasible.

So if the second one is infeasible,

the first one must be feasible.

\therefore Exactly one of the systems has solution.

Problem 5 (15pts).

Suppose M is a square matrix such that $M = -M^T$, for example,

$$M = \begin{pmatrix} 0 & 1 & 2 \\ -1 & 0 & -4 \\ -2 & 4 & 0 \end{pmatrix}$$

Consider the following optimization problem:

$$\begin{array}{ll} \text{minimize } & \mathbf{c}^T \mathbf{x} \\ \text{subject to } & M\mathbf{x} \geq -\mathbf{c} \\ & \mathbf{x} \geq 0 \end{array}$$

- (a). Show that the dual problem of it is equivalent to the primal problem.
 (b). Argue that the problem has optimal solution if and only if there is a feasible solution.

(a) In this problem, $b = -c$

The dual problem is:

$$\begin{array}{ll} \text{maximize } & \mathbf{b}^T \mathbf{y} \\ \text{s.t. } & M^T \mathbf{y} \leq \mathbf{c} \\ & \mathbf{y} \geq 0 \end{array}$$

Since $b = -c$, $M = -M^T$

The dual problem becomes:

$$\begin{array}{ll} \text{minimize } & \mathbf{c}^T \mathbf{y} \\ \text{s.t. } & M \mathbf{y} \geq -\mathbf{c} \\ & \mathbf{y} \geq 0 \end{array}$$

Which is equivalent to the primal problem.

(b) Know that the only possible cases for LP are:

P D	Finite Optimum	Unbounded	Infeasible
Finite Optimum	✓		
Unbounded			✓
Infeasible		✓	✓

Since the primal problem and dual problem are equivalent,
 either they are both feasible or they are both infeasible.

When they are both infeasible, the problem will not have
 optimal solution.

So the problem has optimal solution if and only if
 there is a feasible solution.