COMP 737011 - Memory Safety and Programming Language Design

Lecture 9: Rust Concurrency

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Outline

- 1. Multi-Thread Rust
- 2. Basic Concurrency APIs
- 3. Marker Trait for Concurrency

1. Multi-Thread Rust

Spawn A New Thread

Spawn A Group of Threads

save thread handlers in a vector

Access Objects In the New Thread

- Access the same object from multiple threads is risky:
 - race condition
 - the thread may outlive the lifetime of the object

```
spawn: x = 10
main: x= 1
```

Object Access: Cont'd

main: x=1

```
let mut x = Box::new(1);
let tid = thread::spawn(move|| {
                                       move the ownership of x to the thread
    *x = 10;
    println!("spawn: x = \{\}", x);
});
tid.join().unwrap();
println!("main: x= {}", x);
                                    Illegal to access x
let mut x = Box::new(1);
let mut y = x.clone();
                                       make a clone of x as y
let tid = thread::spawn(move|| {
    *v = 10;
                                       access y
    println!("spawn: y = \{\}", y);
});
tid.join().unwrap();
println!("main: x= {}", x);
spawn: y = 10
```

Share Data Among Threads

```
let x = Box::new(1);
for i in 0..3 {
    let r = &x;
    thread::spawn(move || {
        println!("{:?}",r);
    });
}
the thread may live longer than x
```

```
let x = RC::new(Box::new(1));
for i in 0..3 {
    let cl = x.clone();
    thread::spawn(move || {
        println!("{:?}",cl);
    });
}
```

RC is not thread safe

We Need Concurrency-Safety APIs

- Basic APIs
 - Atomicity or lock
 - Synchronization or memory Barrier
- Advanced features

2. Basic Concurrency APIs

Atomic Types

- Several atomic types
 - AtomicBool,
 - Atomiclsize,
 - AtomicUsize,
 - •
- Similar to C++ std::<atomic>

Sample Mutex Lock with Memory Barrier

```
pub struct Mutex { flag: AtomicBool, }
impl Mutex {
    pub fn new() -> Mutex {
        Mutex { flag: AtomicBool::new(false), }
    pub fn lock(&self) {
        while self.flag.compare_exchange_weak(
                  false, true,
                  Ordering::Relaxed, Ordering::Relaxed)
                  .is err() {}
                                         all subsequent loads will see
        fence(Ordering::Acquire);
                                         the stored data
    pub fn unlock(&self) {
        self.flag.store(false, Ordering::Release);
```

Arc<T>: Atomically Ref Counted

- Similar to RC<T>, but is thread safe
- Use atomic operations for reference counting
- Mutating through an Arc generally use Mutex, RwLock, etc.

```
let x = Arc::new(Box::new(1));
for i in 0..3 {
    let cl = x.clone();
    thread::spawn(move || {
        println!("{:?}",cl);
    });
}
```

Mutex

- Use lock() or try_lock() to access the data
 - Returns Result<T>
 - lock() is blocking mode
 - most usage simply unwrap() the result, why?
 - try_lock() is nonblocking mode
 - returns Err() if fails

```
let x = Arc::new(Mutex::new(0));
for i in 0..3 {
    let cl = x.clone();
    thread::spawn(move || {
        let mut data = cl.lock().unwrap();
        *data += 1;
        println!("{:?}", data);
    });
}
```

Synchronizing Primitive: Condition Variable

- Do not consume CPU when threads need to wait for a resource to become available
- How to implement the feature?

```
let x = Arc::new((Mutex::new(0), Condvar::new()));
let cl = Arc::clone(&x);
thread::spawn(move|| {
    let (1, c) = &*c1;
    let mut t = 1.lock().unwrap();
    *t = 100;
    c.notify one();
});
let (1, c) = &*x;
let mut t = 1.lock().unwrap();
while *t == 0 {
    t = c.wait(t).unwrap();
    println!{"while: t = {}", t};
}
```

Mutex: Poison Strategy

- What if a thread holding the lock panics?
- Using a poison flag to detect/recover from the bad state

```
let arc = Arc::new(Mutex::new(0));
let cl = arc.clone();
let _ = thread::spawn(move | | -> () {
    let mut data = cl.lock().unwrap();
    panic!();
                                                Panic the thread
}).join();
assert_eq!(arc.is_poisoned(), true); -----
                                                The lock is poisoned
let mut guard = match arc.lock() {
                                                Release the locked data
    Ok(guard) => guard,
    Err(poisoned) => poisoned.into inner(),
*guard += 1;
```

Message Passing

- Multi-producer, single-consumer FIFO queue
 - Asynchronous or synchronous mode

```
use std::sync::mpsc;
use std::thread;
fn main() {
    let (tx, rx) = mpsc::channel();
    let tx = tx.clone();
    let tid = thread::spawn(move|| {
        for i in 0..10 {
            tx.send(i).unwrap();
    });
    while let Ok(msg) = rx.recv(){
        println!{"receive: {}",msg};
```

Synchronizing Primitive: Once

- Run global initialization only one time
 - access 'static mut' variables

```
static mut VAL: usize = 0;
static INIT: Once = Once::new();

fn get_cached_val() -> usize {
    unsafe {
        INIT.call_once(|| {
            VAL = expensive_computation();
            });
        VAL
        }
}
```

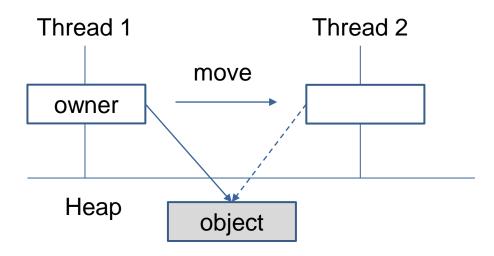
3. Marker Trait for Concurrency

Marker Traits for Concurrency

- Marker Traits have no methods to implement
- They are compiler intrinsic and auto derived
 - Send/!Send
 - Sync/!Sync
- Other marker traits
 - Copy/!Copy
 - Sized/!Sized
 - Unpin/!Unpin

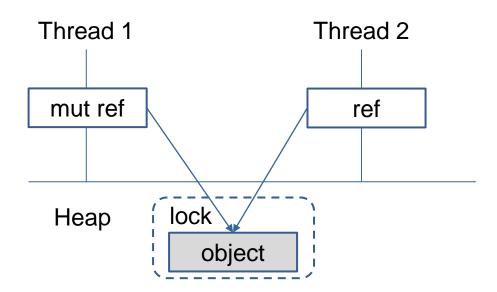
Send

- The type of Send can be transferred between threads
- Use the move operator, which is similar as =
 - For types of Copy trait, make a copy of the object
 - For non-copy, transfer the ownership
- Almost all primitive types are Send
- Any struct composed of Send types is automatically marked as Send



Sync

- The type of Sync is safe to be referenced from multiple threads
- Any type T is Sync if &T is Send
- Sync is usually more rigid than Send. Why?



Raw pointers are neither Send nor Sync

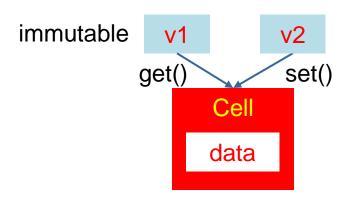
- Possible to create shared objects (although unsafe)
- Should be manually implemented as unsafe

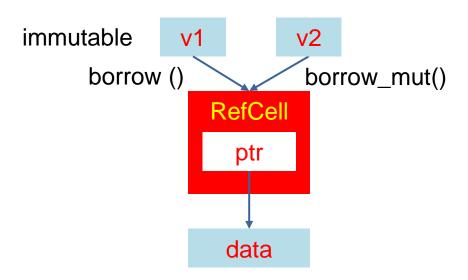
```
struct Unsend{ ptr: *mut i64, }
impl Unsend{
    fn add(&self, i:i64){
        unsafe{*(self.ptr) = *self.ptr + i};
unsafe impl Send for Unsend{}
unsafe impl Sync for Unsend{}
let mut var = 0i64;
let mut v = Unsend{ptr:&mut var as *mut i64};
let tid = thread::spawn(move | {
    for i in 1..100001{ v.add(i); }
});
for i in 1..100001{ var+=i; }
tid.join();
println!("{}",var);
```

Implement Send/Sync is unsafe

Can Cell/RefCell Be Send/Sync?

- unsynchronized interior mutability
- Send but not Sync





Code to Verify Send/Sync Properties

Rc<T> and Arc<T>

- Rc<T> is neither Send nor Sync, why?
 - !Sync: atomicity in reference counter update
 - !Send: cloned Rc exist in multiple threads
- Does Arc<T> have bound on T to be thread-safe?
 - The compiler checks the wrapped data during compilation

Can Mutex be Send/Sync?

Require T is Send

Can Mutex be Send/Sync? Cont'd

Sync but not Send?

- Cases are rare
- Exceptions may relate to thread-local features,
 - e.g., MutexGuard

In-Class Practice

- Rewrite your program (binary search tree or doublelinked list) to be thread-safe
 - Support Sync/Send
- Show that your program is thread safe