2019 KCTF 晋级赛Q1 | 第三题点评及解题思路

小雪 看雪学院 1周前



影分身之术是火影忍者中-鸣人的招牌忍术,火影迷们想必并不陌生了。与简单的分身术相比,这个由二代火影干手扉间开创的忍术,其分出来的分身不再是幻影,而是实打实的实体。

你能成功识破哪个是真身,哪个是幻影吗?接下来,让我们一起来看看这道《影分身之术》,是如何被大家破解的吧!

第三题 影分身之术

攻破	此题的战队			题目名称	第三题 影分身之	之术						
排名	战队名	破解时间	获取积分	出题战队	战神伽罗							
8	pizzatql	7965s	100	+		原来世界那么大,风景那么好! 亲						
	♦ 萌新队	23966s	100		身经历过后印象深刻,能提升的地方还有很多,远无止境,向大神们学习!							
8	kali-go	29998s	100	题目简介	CrackMe运行环境: Windows [公告]2019看雪CTF新賽季! 晋级赛每次6-15题,一次 性放題,賽期14天。战队必须通过晋级赛,才能参加年							
4.	計打酱油	37260s	100		底的总决赛!本比赛要求战队独立回答。在题目未结束 前,请勿在论坛、QQ群等公共场所讨论试题相关信 息。否则视为作弊。欢迎选手加比赛QQ群:8601428							
5.	tekkens	50447s	100	题目下载	rackme2019D7.rar							
6.	♣ A2	98519s	100	提交答案	请输入注册码(序列号)提交 提交							
7.	defxyz	105481s	100	解析文章								
8.	fade-vivi	121752s	100	WI WINST	notwolf	[原创]2019CTF晋级赛Q1第三题 分身之术分析						
9.	→ 雨落星沉	162263s	100		foyjog oooAooo	[原创]第三题 影分身之术 [原创] 2019看雪CTF 晋级赛Q1						
10.	♠ Pizzatql.	168175s	100		000A000	第3题						
		1301703			Cossack人人	[原创]影分身之术 WriteUp from W8C.MozhuCY						
11.	\$ _L_T_	195676s	100		kkHAIKE	[原创]CTF2019Q1 第三题 影分						

此题共有 35 支战队破解, 围观人数达 3443 人。

出题团队



个人主页: https://bbs.pediy.com/user-177594.htm

简介:

野生程序猿,临床医学专业转行,现为北京某网络安全公司研发,从事恶意软件动态行为分析产品的开发,对黑科技以及产品架构有很大兴趣,设计出了超稳定的大规模挂钩引擎,kctf中的战神伽罗系列题目某种程度上反应了本猿平时的工作特征,本系列赛题将会和你一起不断进化,成长,欢迎继续关注,感谢看雪!感谢有你参与!

看雪CTF 评委 crownless 点评

《影分身之术》这道题采用了javascript加密,并考察了简单的汇编。然而,不同的战队采用了不同的解题手法和工具,让人大开眼界,值得学习和借鉴。

题目设计思路

- 1. 首先将一部分密码封装在javascript中,通过javascript将自身进行加密。
- 2. 通过简单的汇编代码变形算法 (加减固定数值) , 将一部分密码代码编译到可执行区域,通过指令跳转和对硬编码的变形对这部分密码进行恢复比对。
- 3. 将恢复的代码注入到IE内核当中,并显示出来。

破解思路

1. 搜索内存可以搜到javascript代码,并将代码中的一部分密码获取出来。

```
0049422C
             24724900
                           dd crackme2.00497224
                                                          UNICODE "eval(function(p,a,c,k,e,d){e=function
                           dd crackme2.00492BA8
00494230
00494234
             832D 408F490 sub
                               Xcrackme2.00494247
0049423B
             73 OA
0049423D
             B8 28424900
                           mov eax, crackme2.00494228
00494242
             E8 0100F7FF
                                crackme2.00404248
00494247
             C3
00494248
             55
                                ebp
00494249
             8BEC
                               ebp, esp
00494248
             33C0
                           xor
                               eax,eax
0049424D
             55
                                ebp
0049424E
             68 67424900
                                crackme2.00494267
00494253
             64:FF30
                                [dword fs:eax]
00494256
             64:8920
                                dword fs:eax],esp
```

2. 解密后代码, 密码为: simpower91

```
function ckpswd() {
key="simpower91";
a = document.all.pswd.value;
if (a.indexOf(key) ==0) {
l=a.length;
i=key.length;
sptWBCallback(a.substring(i,1));
} else {
alert("wrong!<" + a + "> is not my GUID ;-)");
return "1234";
}
function ok() {
alert("congratulations!");
}
```

3. 剩下4位很容易就可以跟踪到比对代码的那个call处,有4个很怪异的指令被跳过,将每个字节-7F就是剩余的ascii码。

```
8844316F
             55
00493F70
                                 ebp
00493F71
             8BEC
                           mov ebp,esp
00493F73
             81C4 D8FBFFF add esp, -0x428
             53
00493F79
                                 ebx
00493F7A
             56
                                 esi
                           xor ebx,ebx
00493F7B
             33DB
00493F7D
             899D D8FBFFF mov [dword ss:ebp-0x428],ebx
00493F83
             899D DCFBFFF mov [dword ss:ebp-0x424],ebx
00493F89
             899D EOFBFFF[mov [dword ss:ebp-0x420],ebx
00493F8F
             8955 EC
                           mov [dword ss:ebp-0x14],edx
00493F92
             8BD8
                           mov ebx,eax
                           mov eax,[dword ss:ebp-0x14]
00493F94
             8B45 EC
00493F97
             E8 8809F7FF
                            all crackme.00404924
00493F9C
             8B45 08
                           mov eax,[dword ss:ebp+0x8]
00493F9F
             E8 8009F7FF
                            call crackme.00404924
                           xor eax, eax
00493FA4
             33C0
00493FA6
             55
                                 ebp
00493FA7
             68 BA414900
                                 crackme.004941BA
00493FAC
             64:FF30
                                 [dword fs:eax]
00493FAF
             64:8920
                           mov [dword fs:eax],esp
                           mov eax,[dword ss:ebp-0x14]
00493FB2
             8B45 EC
00493FB5
             E8 7A07F7FF
                            call crackme.00404734
00493FBA
             83F8 04
                           cmp eax,0x4
00493FBD
              0F85 BE01000 jnz crackme.00494181
                            mp Xcrackme.00493FE6
00493FC3
              EB
00493FC5
             EØ
                           db E0
00493FC6
             BO
                           db B0
00493FC7
             B1 B2
                           mov cl, 0xB2
```

其中E0-7F='a'

B0-7F='1'

B1-7F='2'

B2-7F='3'

界面如下所示(注册成功后):



题目破解思路

本题破解思路由 oooAooo 提供



一、初探

初看程序是由dephi写的,从出题团队来看,可能是去年的加强版。用DEDE反编译一下看看,虽然注册了几个按钮事件,但并不是核心校验函数。可能是dephi+script形式。打开IDA可以发现类似如下字符串:

```
CODE:00493628 db 'function sptWBCallback(spt_wb_id, spt_wb_name, optionstr) {url=',27h; Te CODE:00493628 db '#sptWBCallback:id=',27h,';url=url+spt_wb_id+',27h,';eventName=',27h; CODE:00493628 db '+spt_wb_name;if(optionstr) url=url+',27h,';params=optionstr',27h,';'; CODE:00493628 db 'location=url;}',0 ; Text

CODE:00493868 db '<center><br/>
'code:00493868 db '<center><br/>
'code:00493868 db '><br>
'input type=button value="checkMyFlag" onclick="ckpswd()'; Text

CODE:00493868 db ';"></center>',0 ; Text

CODE:00493868 db ';"></center>',0 ; Text

CODE:0049386F align 10h
```

从字符串上看,存在一个html,里面有个onclick的处理函数 ckpswd函数,首先需要找到 ckpswd函数体。不想一步一步分析,启动程序,用ce在内存总搜索,很快可以发现如下脚本函数:

```
function ckpswd()
{
key="simpower91";
a = document.all.pswd.value;
if(a.indexOf(key) ==0)
{
l=a.length;
i=key.length;
sptWBCallback(a.substring(i,1));
}
else
{
alert("wrong!<" + a + "> is not my GUID;-)");
return "1234";
}
}
```

可以看出要求输入的sn开头部分必须是 "simpower91", 否则会提示 'worng!', 测试一下确实如此, 如果输入simpower91111, 没有任何提示, 看起来仍然有进一步的判断。从上面的脚本可知为sptWBCallback函数, 看起来像个回调函数。先找到sptWBCallback函数体。从前面的字符串中可以看到其影子, 在内存中搜索下, 发现如下:

```
function sptWBCallback(spt_wb_id, spt_wb_name, optionstr)
{
url='#sptWBCallback:id=';
url=url+spt_wb_id+';
eventName='+spt_wb_name;
```

```
if(optionstr)
url=url+';
params=optionstr';
location=url;
}
```

上面代码应该是WebBrowser JS回调delphi的方法。因此真正的检测应该仍然在dephi主程序中。需要找到注册时间的回调函数。分析IDA反汇编代码发下如下:

```
CODE:00491DCC sub_491DCC proc near; CODE XREF: _Tfrmcrackme_FormCreate+81 ↓ p
CODE:00491DCD push ebx

CODE:00491DCD mov ebx, eax

CODE:00491DCF mov eax, ebx

CODE:00491DD1 call sub_491B78; call function

CODE:00491DD6 mov eax, [ebx+40h]

CODE:00491DD9 mov [eax+2A4h], ebx

CODE:00491DDF mov dword ptr [eax+2A0h], offset sub_492088

CODE:00491DEA retn

CODE:00491DEA sub_491DCC endp
```

其中sub_492088函数比较可疑。设置个断点,输入"simpower91111"果然断了下了。

二、回调 sub_492088

```
int v27; // [esp+8h] [ebp-18h]
char v28[4]; // [esp+Ch] [ebp-14h]
int v29; // [esp+10h] [ebp-10h]
int System::AnsiString; // [esp+14h] [ebp-Ch]
int v31: // [esp+18h] [ebp-8h]
int v32; // [esp+1Ch] [ebp-4h]
int savedregs; // [esp+20h] [ebp+0h]
System::AnsiString = 0;
v29 = 0:
*(DWORD *)v28 = 0;
v27 = 0:
v26 = 0:
v25 = 0:
v24 = a4;
v23 = a5;
v31 = a3:
v32 = a2;
v12 = a1:
v22 = &savedregs;
v21 = \&1oc 4921D1;
v20 = readfsdword(0);
writefsdword(0, (unsigned int)&v20);
Variants:: linkproc VarToLStr(&v29, all);
v13 = sub 465C88((int)&str sptWBCallback [1], v29);
if (v13 > 0)
{
*a6 = -1:
Variants:: linkproc VarToLStr(&System::AnsiString, all);
v14 = GetJSLen(System::AnsiString);
v15 = System:: linkproc LStrCopy(System::AnsiString, v13 + 15, v14 - v13 - 14, (int)
LOBYTE(v15) = 1;
unknown_libname_161(System::AnsiString, (int)&str___41[1], (int)&str___19[1], v12, edi
System:: linkproc LStrLAsg(&System::AnsiString, *(signed int32 *)v28);
v16 = (Classes::TStrings *)TStreamCreate((int)cls Classes TStringList, 1);
(*(void ( fastcall **) (Classes::TStrings *, int)) (*( DWORD *) v16 + 44)) (v16, System::A
if (*(WORD *)(v12 + 50))
{
Classes::TStrings::GetValue(v16, (const int)&str params[1], (int)&v27);
v17 = v27;
Classes::TStrings::GetValue(v16, (const int)&str eventName[1], (int)&v26);
v18 = v26;
Classes::TStrings::GetValue(v16, (const int)&str id[1], (int)&v25);
(*(void (fastcall **) (DWORD, int, int, int)) (v12 + 48)) (*(DWORD *) (v12 + 52), v25,
```

```
}
if (*(_WORD *)(v12 + 58))
(*(void (__fastcall **)(_DWORD, int, int, int, int, int, int, int, _WORD *))(v12 + 0x38
*(_DWORD *)(v12 + 60),
v32,
v31,
a11,
a10,
a9,
a8,
a7,
a6);
__writefsdword(0, v20);
v22 = (int *)&loc_4921D8;
return System::__linkproc__ LStrArrayClr(&v25, 6);
}
```

核心处理函数为 V12+0X38, 地址为;493F70

三、493F70函数

```
loc_493F70: ; DATA XREF: _Tfrmcrackme_FormCreate+87 ↑ o
CODE: 00493F70 push ebp
CODE:00493F71 mov ebp, esp
CODE: 00493F73
CODE:00493F73 1oc 493F73: ; CODE XREF: CODE:00493FC5 ↓ j
CODE: 00493F73 add esp, 0FFFFFBD8h
CODE: 00493F79 push ebx
CODE: 00493F7A push esi
CODE: 00493F7B xor ebx, ebx
CODE: 00493F7D mov [ebp-428h], ebx
CODE: 00493F83 mov [ebp-424h], ebx
CODE:00493F89 mov [ebp-420h], ebx
CODE:00493F8F mov [ebp-14h], edx
CODE:00493F92 mov ebx, eax
CODE: 00493F94 mov eax, [ebp-14h]
CODE: 00493F97 call @System@@LStrAddRef$qqrpv; System:: linkproc LStrAddRef(void *)
CODE:00493F9C mov eax, [ebp+8]
CODE: 00493F9F call @System@@LStrAddRef$qqrpv; System:: linkproc LStrAddRef(void *)
CODE: 00493FA4 xor eax, eax
CODE: 00493FA6 push ebp
```

```
CODE: 00493FA7 push offset loc 4941BA
CODE:00493FAC push dword ptr fs:[eax]
CODE: 00493FAF mov fs: [eax], esp
CODE:00493FB2 mov eax, [ebp-14h]
CODE: 00493FB5 call GetJSLen; BDS 2005-2007 and Delphi6-7 Visual Component Library
CODE:00493FBA cmp eax, 4
CODE: 00493FBD jnz loc 494181
CODE:00493FC3 jmp short loc_493FE6
CODE:00493FC5 ; -----
CODE: 00493FC5 loopne near ptr loc 493F73+4
CODE: 00493FC7 mov cl, 0B2h
CODE:00493FC9 mov edx, offset byte 494045
CODE:00493FCE mov [ebp-4], edx
CODE:00493FD1 mov edx, offset loc 4940D0
CODE:00493FD6 mov [ebp-8], edx
CODE:00493FD9 mov edx, offset dword 494168
CODE:00493FDE mov [ebp-0Ch], edx
CODE: 00493FE1 jmp loc 4940DA
CODE: 00493FE6 : ----
CODE: 00493FE6
CODE:00493FE6 loc 493FE6: ; CODE XREF: CODE:00493FC3 ↑ j
CODE:00493FE6 mov eax, offset str data txt 1. Text
CODE:00493FEB call @Sysutils@FileExists$qqrx17System@AnsiString ; Sysutils::FileExists(
CODE: 00493FF0 test al, al
CODE:00493FF2 jz short loc 494010
CODE: 00493FF4 mov d1, 1
CODE:00493FF6 mov eax, off_416648
CODE: 00493FFB call TStreamCreate; BDS 2005-2007 and Delphi6-7 Visual Component Library
CODE:00494000 mov esi, eax
CODE:00494002 mov edx, offset str data txt 1. Text
CODE:00494007 mov eax, esi
CODE: 00494009 call @Ibsql@TIBXSQLVAR@LoadFromFile$qqrx17System@AnsiString; Ibsql::TIBX
CODE:0049400E jmp short loc 494025
CODE:00494010 ; -----
CODE: 00494010
CODE:00494010 loc 494010: ; CODE XREF: CODE:00493FF2 ↑ j
CODE: 00494010 or ecx, OFFFFFFFFh
CODE:00494013 mov edx, offset _str_data_txt_1.Text
CODE: 00494018 mov eax, [ebx+330h]
CODE:0049401E call GetVmpCode
CODE:00494023 mov esi, eax
CODE: 00494025
CODE:00494025 loc 494025: ; CODE XREF: CODE:0049400E ↑ j
CODE: 00494025 mov edx, offset byte 494045
CODE:0049402A \text{ mov } [ebp-4], edx
```

```
CODE:0049402D mov eax, [ebp-4]
CODE:00494030 push eax
CODE:00494031 mov eax, [esi+4]
CODE: 00494034 push eax
CODE: 00494035 mov eax, ds: dword 498F3C
CODE: 0049403A push eax
CODE:0049403B call sub 473A04
CODE: 0049403B ; -----
CODE: 00494040 dd 0
CODE: 00494044 db 0
CODE:00494045 byte 494045 db 0 ; DATA XREF: CODE:00493FC9 ↑ o
CODE:00494045 : CODE:1oc 494025 ↑ o
CODE:00494046 dd 0
CODE:0049404A align 4
CODE: 0049404C dd 1Ah dup (0)
CODE: 004940B4 db 2 dup (0)
CODE: 004940B6 dd 0
CODE:004940BA align 4
CODE: 004940BC dd 5 dup (0)
```

这个函数与去年的类似,程序对地址 494045代码进行解密,里面存在一个反汇编引擎,对代码进行重定位,并一条一条执行,同时存在一个简单的VMP。

```
73 69 6D 76 6D 01 00 05 00 .....simvm....
01 00 07 00 00 F0 00 00 00 00 00 00 00 01 00 .....
00 00 01 05 00 00 F0 00 00 00 00 00 00 00 00 00 .....
00 00 00 02 01 01 02 00 00 F0 00 00 00 00 F3 FF ......
FF FF 00 00 00 00 00 07 00 00 F0 00 00 00 00 00 .....
00 00 00 00 00 00 00 03 01 00 07 00 00 F0 00 00 .....
00 00 00 00 00 00 01 00 00 01 05 00 00 F0 00 ......
00 00 00 01 00 00 00 00 00 00 00 03 01 01 02 00 .......
00 F0 00 00 00 00 F2 FF FF FF 00 00 00 00 07 ......
01 00 07 00 00 F0 00 00 00 00 00 00 00 01 00 .....
00 00 01 05 00 00 F0 00 00 00 00 02 00 00 00 00 ......
00 00 00 03 01 01 02 00 00 F0 00 00 00 00 F1 FF .....
FF FF 00 00 00 00 00 07 00 00 F0 00 00 00 00 00 ......
00 00 00 00 00 00 00 03 01 00 07 00 00 F0 00 00 ......
00 00 00 00 00 00 01 00 00 00 01 05 00 00 F0 00 ......
00 00 00 03 00 00 00 00 00 00 00 03 01 01 02 00 ......
```

00	00	F0	00	00	00	00	00	00	00	00	00	00	00	00	03	
01	00	07	00	00	F0	00	00	00	00	00	00	00	00	00	00	
00	00	04	02	00	00	F0	00	00	00	00	EC	FF	FF	FF	00	
00	00	00	03	01	00	07	00	00	F0	00	00	00	00	00	00	
00	00	00	00	00	00	01	07	00	00	F0	00	00	00	00	00	
00	00	00	00	00	00	00	03	6F	72	69	67	6E	83	CO	7F	orign啟.
33	D2	73	69	6D	76	6D	01	00	05	00	00	F0	00	00	00	3賴 imvm
00	00	00	00	00	01	00	00	00	01	02	00	00	F0	00	00	
00	00	F3	FF	FF	FF	00	00	00	00	03	6F	72	69	67	6E	orign
3B	C2	75	52	73	69	6D	76	6D	01	00	07	00	00	F0	00	;聈 Rsimvm
00	00	00	00	00	00	00	00	00	00	00	04	02	00	00	F0	
00	00	00	00	EC	FF	FF	FF	00	00	00	00	03	01	00	07	
00	00	F0	00	00	00	00	00	00	00	00	00	00	00	00	01	
07	00	00	F0	00	00	00	00	01	00	00	00	00	00	00	00	
04	6F	72	69	67	6E	83	CO	7F	33	D2	73	69	6D	76	6D	.orign兝 .3襰 imvm
01	00	05	00	00	F0	00	00	00	00	00	00	00	00	01	00	
00	00	01	02	00	00	F0	00	00	00	00	F2	FF	FF	FF	00	
00	00	00	03	6F	72	69	67	6E	3B	C2	75	3F	73	69	6D	orign;聈 ?sim
76	6D	01	00	07	00	00	F0	00	00	00	00	00	00	00	00	vm
00	00	00	00	04	02	00	00	F0	00	00	00	00	EC	FF	FF	
FF	00	00	00	00	03	01	00	07	00	00	F0	00	00	00	00	
00	00	00	00	00	00	00	00	01	07	00	00	F0	00	00	00	
00	02	00	00	00	00	00	00	00	04	6F	72	69	67	6E	83	orign.
CO	7F	33	D2	73	69	6D	76	6D	01	00	05	00	00	F0	00	3襰 imvm
00	00	00	00	00	00	00	01	00	00	00	01	02	00	00	F0	
																ori
67	6E	3B	C2	75	2C	73	69	6D	76	6D	01	00	07	00	00	gn;聈 ,simvm
F0	00	00	00	00	00	00	00	00	00	00	00	00	04	02	00	
00	F0	00	00	00	00	EC	FF	FF	FF	00	00	00	00	03	01	
00	07	00	00	F0	00	00	00	00	00	00	00	00	00	00	00	
00	00	04	6F	72	69	67	6E	83	CO	7F	33	D2	73	69	6D	orign兝 .3襋 im
76	6D	01	00	05	00	00	F0	00	00	00	00	00	00	00	00	vm
01																
FF																orign; 聈
85	ΕO	FB	FF	FF	50	33	С9	73	69	6D	76	6D	01	00	05	呧P3蓅 imvm
																錁 I.
																4
00	00	00	00	06	6F	72	69	67	6E	E8	27	DD	FF	FF	ЕВ	orign
09	73	69	6D	76	6D	65	6E	64	00	. S	imvr	nend	1			

上面是指令码,其中带有 simvm 标识的为VMP代码,带有orgin的为原始代码,对于 simvm的代码走虚拟机执行引擎,对于orgin标识通过反编译引擎,并重定位执行。

493F70 会调用473A04 上面的逻辑主要在本函数执行。

```
CODE:00473A04 sub 473A04 proc near ; CODE XREF: CODE:0049403B ↓ p
CODE: 00473A04
CODE: 00473A04 \text{ var } 2C = dword ptr -2Ch
CODE: 00473A04 \text{ var } 8 = \text{dword ptr } -8
CODE: 00473A04 \text{ var } 4 = dword \text{ ptr } -4
CODE: 00473A04 \text{ var } s0 = dword ptr 0
CODE: 00473A04 \text{ arg } 4 = dword ptr OCh
CODE: 00473A04 \text{ arg } 8 = dword ptr 10h
CODE: 00473A04
CODE: 00473A04 push ebp
CODE:00473A05 mov ebp, esp
CODE: 00473A07 leave
CODE:00473A08 push dword ptr [esp+0]
CODE:00473AOB push [esp+var s0]
CODE: 00473A0E push [esp+4+var 4]
CODE:00473A11 pusha
CODE:00473A12 pushf
CODE:00473A13 call sub 4723EC
CODE: 00473A18 mov ebx, eax
CODE: 00473A1A xor eax, eax
CODE: 00473A1C mov [ebx+1050h], eax
CODE: 00473A22 mov eax, OFFFFFFFFh
CODE: 00473A27 mov [ebx+1054h], eax
CODE: 00473A2D mov [esp+2Ch+var 8], ebx
CODE: 00473A31 mov eax, [esp+2Ch+arg 8]; eax = encryptCodeAddr
CODE: 00473A35 \text{ mov } [esp+2Ch+var 4], eax
CODE:00473A39 mov eax, [esp+2Ch+arg 4]
CODE:00473A3D mov [esp+2Ch+var s0], eax
CODE:00473A41 mov eax, [ebx+1090h]
CODE:00473A47 cmp eax, 0
CODE: 00473A4A jnz short loc 473A57
CODE: 00473A4C mov eax, offset codeCrypt
CODE: 00473A51 mov [ebx+1090h], eax
CODE: 00473A57
CODE:00473A57 loc 473A57: ; CODE XREF: sub 473A04+46 ↑ j
CODE: 00473A57 popf
CODE: 00473A58 popa
CODE:00473A59 call HandleCode
CODE: 00473A5E pushf
```

```
CODE: 00473A5F pushf
CODE: 00473A60 pushf
CODE: 00473A61 pusha
CODE: 00473A62 pushf
CODE: 00473A63 jmp short loc 473A89
CODE: 00473A65 : ---
CODE: 00473A65
CODE:00473A65 loc_473A65: ; CODE XREF: sub_473A04+98 ↓ j
CODE: 00473A65 mov eax, [ebx+1058h]
CODE:00473A6B mov [esp+2Ch+var s0], eax
CODE: 00473A6F mov eax, [ebx+1050h]
CODE: 00473A75 mov [esp+2Ch+var 4], eax
CODE:00473A79 mov [esp+2Ch+var 8], ebx
CODE: 00473A7D popf
CODE: 00473A7E popa
CODE: 00473A7F call HandleCode
CODE: 00473A84 pushf
CODE: 00473A85 pushf
CODE: 00473A86 pushf
CODE: 00473A87 pusha
CODE: 00473A88 pushf
CODE: 00473A89
CODE:00473A89 1oc 473A89: ; CODE XREF: sub 473A04+5F↑j
CODE:00473A89 call sub 4723EC
CODE: 00473A8E mov ebx, eax
CODE: 00473A90 mov eax, [ebx+1050h]
CODE: 00473A96 cmp eax, [ebx+1054h]
CODE: 00473A9C jb short loc 473A65
CODE:00473A9E mov eax, [ebx+1050h]
CODE:00473AA4 mov [esp+30h], eax
CODE: 00473AA8 mov eax, [esp+2Ch+var 2C]
CODE:00473AAB mov [esp+2Ch+var s0], eax
CODE: 00473AAF popf
CODE: 00473ABO popa
CODE:00473AB1 add esp, 8
CODE: 00473AB4 popf
CODE:00473AB5 retn OCh
CODE: 00473AB5 sub 473A04 endp
CODE: 00473AB5
```

而其核心函数为HandleCode (472EAC)

```
int usercall HandleCode@<eax>(int a1@<ebx>, int a2@<edi>, int a3@<esi>, int a4, int a
unsigned int v15; // et0
int ( fastcall *v16) (unsigned int *); // ST10 4
unsigned int v18; // [esp+10h] [ebp-30h]
int v19; // [esp+30h] [ebp-10h]
void **v20; // [esp+34h] [ebp-Ch]
int *v21; // [esp+38h] [ebp-8h]
int v22; // [esp+3Ch] [ebp-4h]
unsigned int vars0; // [esp+40h] [ebp+0h]
void *retaddr; // [esp+44h] [ebp+4h]
vars0 = a3:
v22 = a2:
v21 = (int *) & vars0:
v20 = &retaddr;
v19 = a1;
v15 = readeflags();
*(DWORD *) (a13 + 4188) = retaddr;
DecryptData(a13, a14, a12, (char *)a15, &v21, (int *)&v20, &v19);
vars0 = v18;
writeeflags(v18);
writeeflags(v18);
return v16(&vars0);
                                                                                      •
```

其中 "16(&vars0);"语句为执行重定位后的代码。而虚拟机以及原始代码重定位在 DecryptData函数 (472EAC) 如下:

```
DWORD *_stdcall DecryptData(int al, int orgAddr, int a3, char *dataOfFile, _DWORD *a5,
{
    _DWORD *global; // ebx
int v8; // esi
int v10; // eax
int v11; // ecx
int v12; // esi
int v13; // esi
int v14; // eax
int v15; // ST18_4
unsigned int v17; // [esp+1Ch] [ebp-77Ch]
void *v18; // [esp+20h] [ebp-778h]
int *v19; // [esp+24h] [ebp-7774h]
```

```
int v20; // [esp+34h] [ebp-764h]
int v21; // [esp+38h] [ebp-760h]
char data[1024]; // [esp+3Fh] [ebp-759h]
int v23; // [esp+440h] [ebp-358h]
char v24; // [esp+444h] [ebp-354h]
char v25; // [esp+544h] [ebp-254h]
unsigned __int8 orgCodeLen; // [esp+77Ah] [ebp-1Eh]
char v27; // [esp+77Bh] [ebp-1Dh]
char *nextData; // [esp+77Ch] [ebp-1Ch]
char *data1; // [esp+780h] [ebp-18h]
BYTE *v30; // [esp+784h] [ebp-14h]
int outDataLen; // [esp+788h] [ebp-10h]
int v32; // [esp+78Ch] [ebp-Ch]
int opCodeLen; // [esp+790h] [ebp-8h]
_DWORD *v34; // [esp+794h] [ebp-4h]
int vars0; // [esp+798h] [ebp+0h]
v21 = 0;
v20 = 0:
v30 = 0;
v19 = &vars0;
v18 = \&1oc 4732B2;
v17 = readfsdword(0);
writefsdword(0, (unsigned int)&v17);
global = sub 4723EC();
v34 = (DWORD *) (vars0 + 8);
data1 = data;
v8 = *(DWORD *) off 497644;
*(DWORD *) (*(DWORD *) off 497644 + 52) = 0;
v27 = 0;
while (1)
callCodeDecrypt((int)global, dataOfFile, data, &outDataLen);
while (1)
while ( *( DWORD *) (v8 + 52) == 1 )
{
orgCodeLen = 0;
nextData = callVmpHandle(v8, (BYTE *)orgAddr, data, (int)&orgCodeLen);
opCodeLen = nextData - data;
GetNextValueByLen(&dataOfFile, nextData - data);
GetNextValueByLen(&orgAddr, orgCodeLen);
callCodeDecrypt((int)global, dataOfFile, data, &outDataLen);
if (ifCode simvm orign(v8, &data1)!= 1)// 不是 simvm
```

```
break;
data1 = data;
v27 = 1;
simvm Handle(global, v34);
data1 = data;
if (!v27)
break;
v27 = 0;
orgin Handle (global, v34);
outDataLen = off 498CAO(data, outDataLen, 0x400000, &v23, 4, v17);// 执行原始code
*a5 = outDataLen;
v9 = unknown_1ibname_29(32 - outDataLen - 6);
data1 = (char *) globa1 + v9 + 4144;
sub_464FA4(data1, data, outDataLen);
*((BYTE *)global + v9 + outDataLen + 4144) = 104;
GetCharFormTstring((int)&v30, &v25);
if ((unsigned int8) sub 472520(v30, global[1047]))
{
GetCharFormTstring((int)&v30, &v24);
opCodeLen = sub 465C88((int)&str 29[1], (int)v30);
v10 = GetJSLen(v30);
System:: linkproc LStrCopy((int)v30, opCodeLen + 1, v10 - opCodeLen, (int)&v30);
opCodeLen = sub 465CEC(v30);
if (outDataLen -1 \ge 0)
v11 = outDataLen;
v32 = 0;
do
*((BYTE *)global + v9 + v32++ + 4144) = -112;
--v11;
while ( v11 );
opCodeLen += outDataLen;
v12 = sub 471F04(v8, data, opCodeLen);
else if (sub_465C88((int)\&str_CALL_0[1], (int)v30) == 1)
GetCharFormTstring((int)&v30, &v24);
if ( sub 465C88((int)&str FF 0[1], (int)v30) != 1 )
opCodeLen = sub_465C88((int)&str___29[1], (int)v30);
```

```
v32 = opCodeLen - 1;
v13 = (opCodeLen - 1) / 2;
data1 = (char *) globa1 + v9 + v13 + 4144;
v14 = GetJSLen(v30);
System:: linkproc LStrCopy((int)v30, opCodeLen + 1, v14 - opCodeLen, (int)&v30);
opCodeLen = sub 465CEC(v30);
v32 = opCodeLen + orgAddr - ((_DWORD)global + v9 + 4144);
sub_464FA4(data1, &v32, outDataLen - v13);
opCodeLen = outDataLen;
v12 = outDataLen;
else
opCodeLen = outDataLen;
v12 = outDataLen;
*a7 = opCodeLen;
opCodeLen = a3:
data1 = (char *) globa1 + v9 + outDataLen + 4145;
sub 464FA4 (data1, &opCodeLen, 4u);
*((BYTE *)global + v9 + outDataLen + 4149) = 0xC3u;
*a6 = (int) global + v9 + 0x1030;
global[1044] = orgAddr + *a7;
global[1046] = &dataOfFile[v12];
GetCharFormTstring((int)&v21, &v25);
v15 = v21:
GetCharFormTstring((int)&v20, &v24);
sub 4736D4((int)global, *a6, orgAddr, v15, v20, outDataLen);
writefsdword(0, v17);
v19 = (int *) \& 1oc 4732B9;
System::__linkproc__ LStrArrayClr(&v20, 2);
return System::__linkproc__ LStrClr(&v30);
```

其中如下代码为执行VMP代码:

```
while ( 1 )
{
callCodeDecrypt((int)global, dataOfFile, data, &outDataLen);
while ( 1 )
{
while ( *( DWORD *) (v8 + 52) == 1 )
```

```
orgCodeLen = 0;
nextData = callVmpHandle(v8, (BYTE *)orgAddr, data, (int)&orgCodeLen);
opCodeLen = nextData - data;
GetNextValueByLen(&dataOfFile, nextData - data);
GetNextValueByLen(&orgAddr, orgCodeLen);
callCodeDecrypt((int)global, dataOfFile, data, &outDataLen);
if (ifCode_simvm_orign(v8, &data1)!= 1)// 不是 simvm
break:
data1 = data;
v27 = 1;
simvm_Handle(global, v34);
data1 = data;
if (!v27)
break;
v27 = 0;
orgin Handle (global, v34);
```

后半部分为执行原始代码,进行反编译从定位。

下面是虚拟程序相关函数:

```
data_2 = vmpHandle(unknowGlogal_1, data_1, (char *)global);
else
*(DWORD *) (unknowGlogal 1 + 52) = 2;
return data_2;
char *__fastcall vmpHandle(int unknowGlogal, char *data, char *a3)
char *v3; // ST1C_4
int unknowGlogal 1; // edi
char numOData; // b1
int num6Data; // ebp
char *dataEnd; // [esp+4h] [ebp-38h]
char num5Data; // [esp+8h] [ebp-34h]
int num7Data; // [esp+Ch] [ebp-30h]
int num8Data; // [esp+10h] [ebp-2Ch]
int num9Data; // [esp+14h] [ebp-28h]
int numlData; // [esp+18h] [ebp-24h]
int num2Data; // [esp+1Ch] [ebp-20h]
int num3Data; // [esp+20h] [ebp-1Ch]
int num4Data; // [esp+24h] [ebp-18h]
char *data 1; // [esp+28h] [ebp-14h]
v3 = a3:
unknowGlogal_1 = unknowGlogal;
dataEnd = data + 0x24;
data 1 = data;
GetNextValueByLen(&data 1, 1);
numOData = *data 1;
GetNextValueByLen(&data_1, 1);
num1Data = *(DWORD *) data 1;
GetNextValueByLen(&data 1, 4);
num2Data = *(DWORD *) data 1;
GetNextValueByLen(&data 1, 4);
num3Data = *(DWORD *) data 1;
GetNextValueByLen(&data 1, 4);
num4Data = *(DWORD *) data 1;
GetNextValueByLen(&data 1, 4);
num5Data = *data 1;
GetNextValueByLen(&data 1, 1);
num6Data = *(DWORD *) data 1;
GetNextValueByLen(&data 1, 4);
num7Data = *(DWORD *) data 1;
```

```
GetNextValueByLen(&data_1, 4);
num8Data = *(DWORD *) data 1;
GetNextValueByLen(&data 1, 4);
num9Data = *(DWORD *) data 1;
GetNextValueByLen(&data 1, 4);
*v3 = *data_1;
if ( numOData || num5Data )
if ( numOData && num5Data )
sub_471CDC (numOData, num9Data, num8Data, num7Data, num6Data, num4Data, num3Data);
else if (!numOData | num5Data)
if (!numOData && num5Data)
call_vmp_add_reg_value(
unknowGlogal 1,
num1Data,
num2Data,
num3Data,
num4Data,
num6Data,
num7Data,
num8Data,
num9Data,
num5Data):
else
{
call_vmp_memcpy_regAndReg_num(
unknowGlogal 1,
numlData,
num2Data.
num3Data,
num4Data,
num6Data,
num7Data,
num8Data,
num9Data,
numOData);
else
sub 471C08(unknowGlogal 1, num1Data, num2Data, num3Data, num4Data, num6Data, num7Data,
```

```
return dataEnd;
int __fastcall call_vmp_add_reg_value(int unknowGlogal, int num1Data, int num2Data, int
if ( num1Data && num6Data )
return vmp_add_reg_value(unknowGlogal, num1Data, num3Data, num6Data, num8Data, num5Data
if ( num1Data && num7Data )
return vmp_add_reg_value(unknowGlogal, num1Data, num3Data, num7Data, num8Data, num5Data
if ( num1Data && num9Data )
return vmp_add_reg_value(unknowGlogal, num1Data, num3Data, num9Data, 0, 0);
if ( num2Data && num6Data )
return vmp add reg value(unknowGlogal, num2Data, num3Data, num6Data, num8Data, num5Data
if ( num2Data && num7Data )
return vmp add reg value(unknowGlogal, num2Data, num3Data, num7Data, num8Data, num5Data
if ( num2Data )
if ( num9Data )
unknowGlogal = vmp_add_reg_value(unknowGlogal, num2Data, num3Data, num9Data, 0, 0);
return unknowGlogal;
```

就是一些简单的寄存器读取以及加减乘除运算。

下面的函数是针对代码解密的,就是异或0xFF.

```
DWORD *__fastcall codeCrypt(int al, int dataOfFile, char *outData, _DWORD *a4)
{
   int v4; // ecx
   char v5; // zf
   _DWORD *v6; // eax
   unsigned int v8; // [esp+8h] [ebp-20h]
   void *v9; // [esp+Ch] [ebp-1Ch]
   int *v10; // [esp+10h] [ebp-18h]
   int v11; // [esp+20h] [ebp-8h]
   char *outData1; // [esp+24h] [ebp-4h]
   int savedregs; // [esp+28h] [ebp+0h]
```

```
v10 = &savedregs;
v9 = \&1oc 473839;
v8 = readfsdword(0);
writefsdword(0, (unsigned int)&v8);
*a4 = 0x400:
outData1 = outData;
v4 = 0:
do
outData1[v4] = ^{\sim}*(BYTE *)(data0fFile + <math>v4);
++v4:
}
while (v4 != 0x400);
GetCharFormTstring((int)&v11, outData1 + 2);
System::__linkproc__ LStrCmp(v11, &str_simvmend[1]._top);
if (v5) // 判断代码是否结束
v6 = sub 4723EC();
v6[0x415] = v6[0x414];
writefsdword(0, v8);
v10 = (int *) \& 1oc 473840;
return System::__linkproc__ LStrClr(&v11);
```

三、sn

如下vm指令的主要功能就是解密最后的字符串"a123"

00	00	01	05	00	00	F0	00	00	00	00	02	00	00	00	00	
00	00	00	03	01	01	02	00	00	F0	00	00	00	00	F1	FF	
FF	FF	00	00	00	00	00	07	00	00	F0	00	00	00	00	00	
00	00	00	00	00	00	00	03	01	00	07	00	00	FO	00	00	
00	00	00	00	00	00	01	00	00	00	01	05	00	00	F0	00	
00	00	00	03	00	00	00	00	00	00	00	03	01	01	02	00	
00	F0	00	00	00	00	F0	FF	FF	FF	00	00	00	00	00	07	
00	00	F0	00	00	00	00	00	00	00	00	00	00	00	00	03	
01	00	07	00	00	F0	00	00	00	00	00	00	00	00	00	00	
00	00	04	02	00	00	F0	00	00	00	00	EC	FF	FF	FF	00	
00	00	00	03	01	00	07	00	00	F0	00	00	00	00	00	00	
00	00	00	00	00	00	01	07	00	00	F0	00	00	00	00	00	
00	00	00	00	00	00	00	03	6F	72	69	67	6E	83	CO	7F	orign兝.
33	D2	73	69	6D	76	6D	01	00	05	00	00	F0	00	00	00	3襰 imvm
00	00	00	00	00	01	00	00	00	01	02	00	00	F0	00	00	
00	00	F3	FF	FF	FF	00	00	00	00	03	6F	72	69	67	6E	orign
3B	C2	75	52	73	69	6D	76	6D	01	00	07	00	00	F0	00	;聈 Rsimvm
00	00	00	00	00	00	00	00	00	00	00	04	02	00	00	F0	
00	00	00	00	EC	FF	FF	FF	00	00	00	00	03	01	00	07	
00	00	F0	00	00	00	00	00	00	00	00	00	00	00	00	01	
07	00	00	F0	00	00	00	00	01	00	00	00	00	00	00	00	
04	6F	72	69	67	6E	83	CO	7F	33	D2	73	69	6D	76	6D	.orign兝 .3襰 imvm
01	00	05	00	00	F0	00	00	00	00	00	00	00	00	01	00	
00	00	01	02	00	00	F0	00	00	00	00	F2	FF	FF	FF	00	
00	00	00	03	6F	72	69	67	6E	3B	C2	75	3F	73	69	6D	orign;聈 ?sim
76	6D	01	00	07	00	00	F0	00	00	00	00	00	00	00	00	vm
00	00	00	00	04	02	00	00	F0	00	00	00	00	EC	FF	FF	
FF	00	00	00	00	03	01	00	07	00	00	F0	00	00	00	00	
00	00	00	00	00	00	00	00	01	07	00	00	F0	00	00	00	
																orign.
																3襰 imvm
00	00	00	00	00	00	00	01	00	00	00	01	02	00	00	F0	
																ori
																gn;聈 ,simvm
																orign兝 .3襰 im
76																vm
01																
																orign; 聈
																呧 P3蓅 imvm
00	00	00	00	00	00	00	00	00	00	00	00	E8	41	49	00	鐰 I.

sn为: simpower91a123

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