

Composite and product types

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Composite types

A composite object usually contains several fields, each describing the different aspect of the object. A composite object is like a record in Pascal and structure in C. A composite type provides a set of composite objects.

A composite type is constructed using the type constructor: **composed of ... end.**

The general form of a composite type declaration:

```
A = composed of  
    f_1: T_1  
    f_2: T_2  
    ...  
    f_n: T_n  
end
```

where f_i ($i = 1, \dots, n$) are variables called **fields** and T_i are their types. Each field represents an attribute of the composite object of the type. A value of a composite type is called **composite object** (or **composite value**).

A variable **co** of composite type **A** can be declared in one of the following forms:

(1) **co: A**

(2) **co: composed of**

f_1: T_1

f_2: T_2

...

f_n: T_n

end

Example: type declaration:

```
Account = composed of  
    account_no: nat  
    password: nat  
    balance: real  
end
```

Variable declaration:

```
account: Account
```

(1) Constructor

Only one constructor known as **make-function** is available for composite types. The format:

$$\text{mk_A}(v_1, v_2, \dots, v_n)$$

The make-function yields a composite value of composite type **A** whose field values are **v_i** (**i = 1, ..., n**) that corresponds to fields **f₁, f₂, ..., f_n**, respectively. For example,

$$\text{account} = \text{mk_Account}(1073548, 1234, 5000)$$

(2) Operators

(2.1) Field select

Let **co** be a variable of composite type **A**.
Then, we use

co.f_i

to represent the field **f_i** ($i = 1, \dots, n$) of the composite object **co**.

Examples:

account.password
account.balance

(2.2) Field modification (`modify`)

Given a composite value, say `co`, of type `A`, we can apply the field modification operator `modify` to create another composite value of the same type.

`modify(co, f_1 -> v_1, f_2 -> v_2, ..., f_n -> v_n)`

Example:

let `account = mk_Account(1073548, 1234, 5000)`

Then, we can have the expression:

`account1 = modify(account, password -> 4321)`

(3) Comparison

Two composite values can be compared to determine whether they are identical or not.

Examples:

```
mk_Account(1073548, 1234, 5000) =  
    mk_Account(1073548, 1234, 5000)
```

```
mk_Account(1073548, 1234, 5000) <>  
    mk_Account(1073548, 4321, 5000)
```

Product types

A product type defines a set of tuples with a fixed length.

Let T_1, T_2, \dots, T_n be n types.

Then, a product type T is declared as follows:

$$T = T_1 * T_2 * \dots * T_n$$

A value of T is created using the make-function:

$$\text{mk}_T(v_1, v_2, \dots, v_n)$$

Example:

Suppose type **Date** is declared as:

$$\text{Date} = \text{nat0} * \text{nat0} * \text{nat0}$$

Then

mk_Date(1999, 7, 25)

mk_Date(2000, 8, 30)

mk_Date(2001, 7, 10)

are the values of type **Date**.

Examples: the use of tuples:

d: Date;

d = mk_Date(1999, 7, 25)

d = mk_Date(2000, 8, 30)

d = mk_Date(2001, 7, 10)

There are two operators on tuples: **tuple application** and **tuple modification**.

(1) A tuple application yields an element of the given position in the tuple, whose general format is:

$$a(i): T * \text{nat} \rightarrow T_i$$

where **a** is a variable of product type **T**; **i** is a natural number indicating the position of the element referred to in tuple **a**; and **T_i** denotes the **ith** type in the declaration of **T**.

For example, let

`date1 = mk_Date(1999, 7, 25)`

`date2 = mk_Date(2000, 8, 30)`

Then, the following results can be derived:

`date1(1) = 1999`

`date1(2) = 7`

`date1(3) = 25`

`date2(1) = 2000`

`date2(2) = 8`

`date2(3) = 30`

A tuple can also be directly used in applications.

Examples:

`mk_Date(2000, 8, 30)(2) = 8`

`mk_Date(2000, 8, 30)(3) = 30`

(2) A tuple modification is similar to a composite value modification. The same operator **modify** is also used for tuple modification, but with slightly different syntax:

$$\text{modify: } T * T_1 * T_2 * \dots * T_n \rightarrow T$$
$$\text{modify}(tv, 1 \rightarrow v_1, 2 \rightarrow v_2, \dots, n \rightarrow v_n)$$

where T is a product type, T_i ($i = 1, \dots, n, n \geq 1$) are the element types. This operation yields a tuple of the same type based on the given tuple tv , with the first element being v_1 , the second element being v_2 , and so on.

Examples:

$$\text{modify}(\text{mk_Date}(2000, 8, 30), 1 \rightarrow 2001, 3 \rightarrow 20) =$$
$$\text{mk_Date}(2001, 8, 20)$$
$$\text{modify}(\text{mk_Date}(2001, 8, 20), 2 \rightarrow 15) =$$
$$\text{mk_Date}(2001, 15, 20)$$

An example of specification

Suppose we want to build a table to record students' credits resulting from two courses:

personal data	course1	course2	total
Helen,0001,A3	2	2	4
John, 0002,A2	0	2	2
...

We aim to build several processes on this kind of table.

This table can be perceived as a sequence of student data. That is,

$$T = [\text{OneStudent1}, \text{OneStudent2}, \dots, \text{OneStudentn}]$$

type

CourseCredit = nat0;

TotalCredit = nat0;

```
PersonalData = composed of
    name: string
    id: nat0
    class: string
end;
```

```
OneStudent = PersonalData * CourseCredit * CourseCredit *
             TotalCredit;
```

```
StudentsTable = seq of OneStudent;
```

var

```
students_table: StudentsTable;
```

inv

```
forall[i, j: inds(students_table)] | i <> j => students_table(i)(1).id <>
students_table(j)(1).id);
```

```
process Search(search_id: nat0)
    info: OneStudent
ext rd students_table
pre exists[i: inds(students_table)] |
    students_table(i)(1).id = search_id
post exists[i: inds(students_table)] |
    students_table(i)(1).id = search_id and
    info = students_table(i)]
end_process;
```

```
process Update(one_student: OneStudent, credit1,  
               credit2: CourseCredit)  
ext wr students_table  
pre exists[i: inds(students_table)] |  
    students_table(i) = one_student  
post len(students_table) = len(~students_table) and  
    forall[i: inds(~students_table)] |  
        (~students_table(i) = one_student =>  
            students_table(i) =  
                modify(~students_table(i), 2 -> credit1,  
                    3 -> credit2,  
                    4 -> credit1 + credit2)) and  
        (~students_table(i) <> one_student =>  
            students_table(i) = ~students_table(i))  
end_process;  
end_module;
```

Class exercise 7

1. Let $a = \text{mk_Account}(010, 300, 5000)$, where the type Account is defined as follows:

```
Account = composed of
    account_no: nat1
    password: nat1
    balance: real
end
```

Then evaluate the expressions:

- a. $a.\text{account_no} = ?$
- b. $a.\text{password} = ?$
- c. $a.\text{balance} = ?$
- d. $\text{modify}(a, \text{password} \rightarrow 250) = ?$
- e. $\text{modify}(\text{mk_Account}(020, 350, 4050), \text{account_no} \rightarrow 100, \text{balance} \rightarrow 6000) = ?$

3. Let x be a variable of the type `Date` defined as follows:

$\text{Date} = \text{nat0} * \text{nat0} * \text{nat0}$

Let $x = \text{mk_Date}(2002, 2, 6)$.

Then evaluate the expressions:

a. $x(1) = ?$

b. $x(2) = ?$

c. $x(3) = ?$

d. $\text{modify}(x, 1 \rightarrow 2003)$

e. $\text{modify}(x, 2 \rightarrow 5, 3 \rightarrow 29)$

f. $\text{modify}(x, 1 \rightarrow x(1), 2 \rightarrow x(2))$

4. Define a composite type `Student` that has the fields: `name`, `date_of_birth`, `college`, and `grade`. Write specifications for the processes: `Register`, `Change_Name`, `Get_Info`. The `Register` takes a value of `Student` and adds it to the external variable `student_list`, which is a sequence of students. `Change_Name` updates the name of a given student with a new name in `student_list`. `Get_Info` provides all the available field values of a given student in `student_list`.