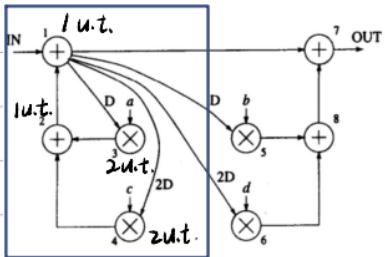
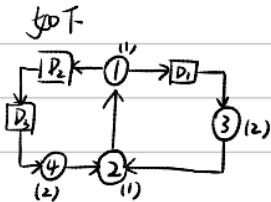


Chapter 2 迭代边界

3. 解：方框内才有环
仅考虑方框内的 DFG



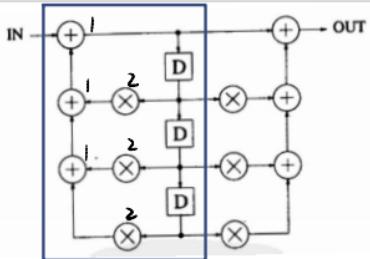
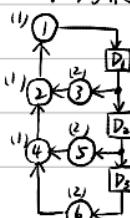
易知 $T_{\infty} = \max \left\{ \frac{4}{2}, \frac{4}{1} \right\} = 4$

下利用 LPM 算法

$$L^{(1)} = \begin{bmatrix} 4 & 4 & -\infty \\ -\infty & -\infty & 0 \\ 4 & 4 & -\infty \end{bmatrix} \rightarrow L^{(2)} = \begin{bmatrix} 8 & 8 & 4 \\ 4 & 4 & -\infty \\ 8 & 8 & 4 \end{bmatrix} \rightarrow L^{(3)} = \begin{bmatrix} 12 & 8 & 8 \\ 8 & 8 & 4 \\ 12 & 12 & 8 \end{bmatrix}$$

则迭代边界 $T_{\infty} = \max \left\{ \frac{8}{2}, \frac{4}{2}, \frac{4}{2}, \frac{12}{3}, \frac{8}{3}, \frac{8}{3} \right\} = 4$

4. 解 仅方框内有环
故仅考虑该部分 DFG



易知 $T_{\infty} = \max \left\{ \frac{4}{1}, \frac{5}{2}, \frac{5}{3} \right\} = 4$

下用 LPM 算法计算

$$L^{(1)} = \begin{bmatrix} 4 & 0 & 0 \\ 5 & -\infty & 0 \\ 5 & -\infty & -\infty \end{bmatrix} \rightarrow L^{(2)} = \begin{bmatrix} 8 & 4 & 4 \\ 9 & 5 & 5 \\ 9 & 5 & 5 \end{bmatrix} \rightarrow L^{(3)} = \begin{bmatrix} 12 & 8 & 8 \\ 13 & 9 & 9 \\ 13 & 9 & 9 \end{bmatrix}$$

\therefore 迭代边界 $T_{\infty} = \max \left\{ \frac{8}{2}, \frac{5}{2}, \frac{5}{2}, \frac{12}{3}, \frac{9}{3}, \frac{9}{3} \right\} = 4$

6. 解：仅图中标注部分有环。

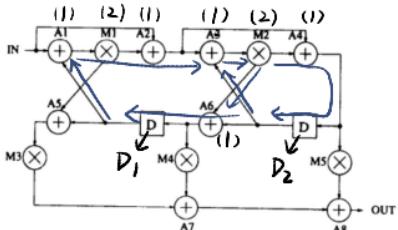
$$\text{易知 } T_{\infty} = \max \left\{ \frac{8}{1}, \frac{4}{1}, \frac{8}{2} \right\} = 8$$

按图示标注 D_1, D_2

则由 LPM 算法

$$L^{(1)} = \begin{bmatrix} 8 & 8 \\ 1 & 4 \end{bmatrix} \rightarrow L^{(2)} = \begin{bmatrix} 16 & 16 \\ 9 & 9 \end{bmatrix}$$

则迭代边界 $T_{\infty} = \max \left\{ \frac{16}{2}, \frac{9}{2} \right\} = 8$



Chapter 3 流水线与并行处理

2.

解：乘法 2 u.t. 加法 1 u.t.

(a) 关键路径为

$$M_1 \rightarrow A_2 \rightarrow M_2 \rightarrow A_1 \rightarrow M_3 \rightarrow A_3 \rightarrow A_4$$

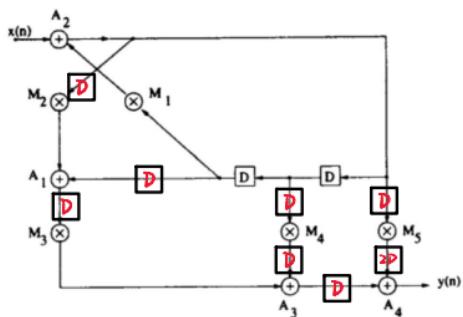
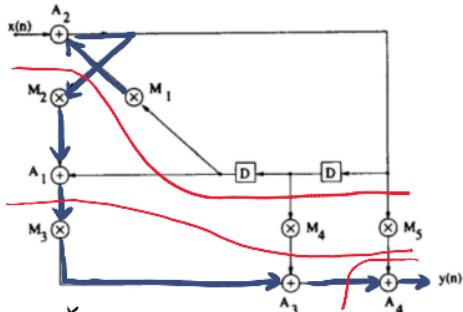
$$2+1+2+1+2+1+1 = 10$$

(b) 将关键路径降为 3 u.t.

在图中画红线位置各填加一个延时单元

结果如右图

共 9 个延时单元



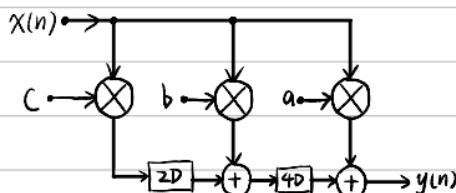
7. 考察 6 阶 FIR 滤波器

$$y(n) = ax(n) + bx(n-4) + cx(n-6) \quad (3-45)$$

(a) 画出该滤波器的拓扑结构，使时钟周期不超过一次乘加的时间。注意不要增加新的锁存器。

(b) 对(a)中结构，画出分组规模为3时的结构。重新分配这个分组结构，使时钟周期是一次乘加时间的四分之一。其中，假定乘法计算时间是加法计算时间的3倍。

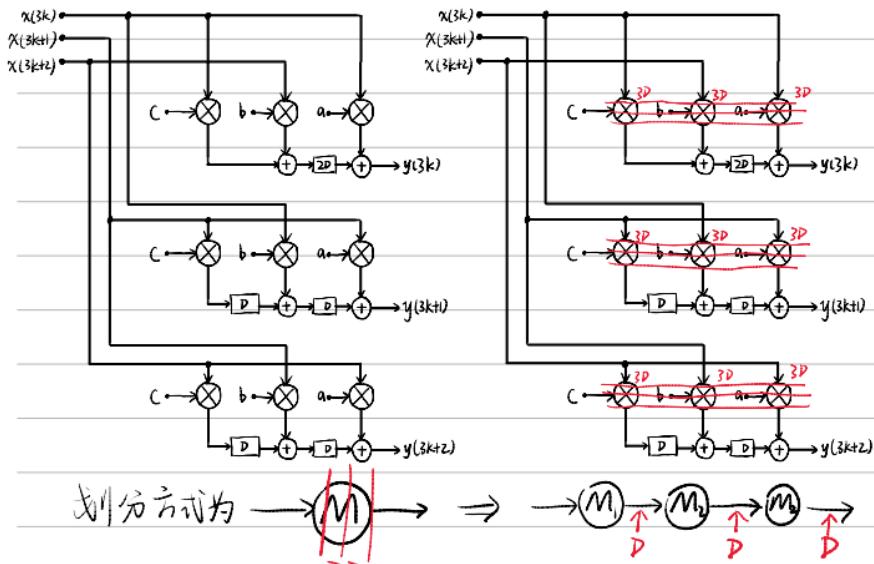
解：(a) 采用数据广播结构



$$\begin{aligned}
 (b) \quad & y(3k) = aX(3k) + bX(3k-4) + cX(3k-6) \\
 & = aX(3k) + bX(3(k-2)+2) + cX(3(k-2)) \\
 & y(3k+1) = aX(3k+1) + bX(3(k-1)) + cX(3(k-2)+1) \\
 & y(3k+2) = aX(3k+2) + bX(3(k-1)+1) + cX(3(k-2)+2)
 \end{aligned}$$

结构如下左图。 $T_m = 3T_A$ 需令 $T_{clk} = \frac{1}{3}(T_m + T_A) = T_A$

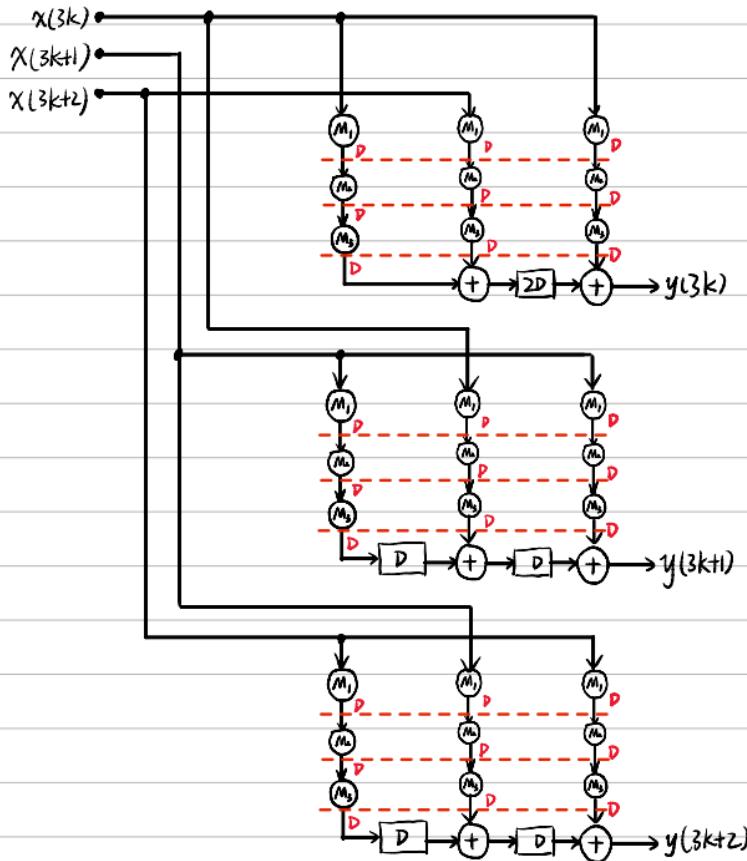
需使用细粒度流水线 / 将 M 度为 $M_1 \rightarrow M_2 \rightarrow M_3$



划分方式为 $\rightarrow \boxed{M} \rightarrow \Rightarrow \rightarrow (M_1) \uparrow D \rightarrow (M_2) \uparrow D \rightarrow (M_3) \uparrow D \rightarrow$

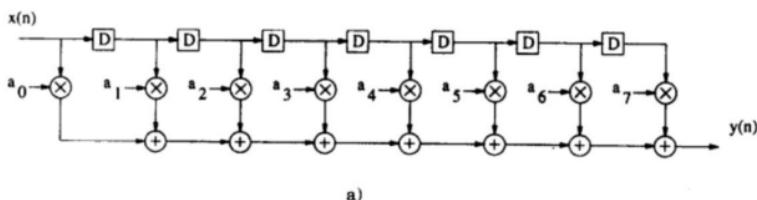
故结果为如上右图

最终行为

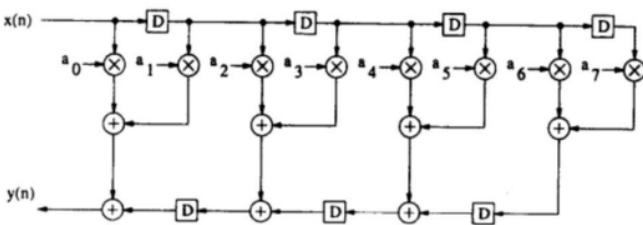


10 一个8阶FIR滤波器的两种实现如图3-24。假定乘法器的关键路径(或传播延时)为加法器的两倍，即 $T_m = 2T_a$ ，因而乘法器充电电容是加法器的两倍。进而假定乘法器的总电容是加法器的10倍，即 $C_m = 10C_a$ 。图3-24a中直接形式结构的关键路径是 $T_m + 7T_a = 9T_a$ 。图3-24b中的结构可以工作在较低的电源电压下，而保持时钟周期或采样周期 $9T_a$ 的约束。因此图3-24b的结构可以用来降低功耗。假如图3-24a的结构工作在4V电源电压下，工艺阈值电压为0.5V，电源电压大于1.2V以保证可容忍的噪声容限。

请问：当图3-24b所示结构达到所希望的 $9T_a$ 采样周期时，最小的工作电源电压是多少？与图3-24a的结构相比，计算图3-24b所示结构在功耗上降低的百分比。其中可以忽略在计算关键路径或功耗时延时元件的传播延时和电容。



a)



b)

解：对 a) 图 $(T_m + 7T_a) = \frac{(C_{cm} + 7C_{ca})V_o}{k(V_o - V_t)^2}$ 且 $T_{seq} = 9T_a = \frac{9C_{ca}V_o}{k(V_o - V_t)^2}$

对 b) 图，类似得 $T'_{seq} = \frac{4(C_{ca}\beta V_o)}{k(\beta V_o - V_t)^2}$ ($V_o' = \beta V_o$)

$$\therefore 9(\beta V_o - V_t)^2 = 4\beta(V_o - V_t)^2 \quad \text{其中 } \begin{cases} V_o = 4V \\ V_t = 0.5V \\ \beta V_o > 1.2V \end{cases}$$

解得 $\beta = \frac{9}{16} = 0.5625 > 0.3 \quad \therefore V_o' = \beta V_o = 2.25V$

故 $\eta = \frac{P'}{P} = \beta^2 \frac{C'_t}{C_t} = (\frac{9}{16})^2 = 31.64\% \quad (C'_t = C_t \text{ 为 } 8M7A)$

减少了 $1 - \eta = 68.36\%$ 的功耗

最小工作电压为 2.25V

12 计算某一计算系统的功耗降低。其中该系统采用了4级流水线和块尺寸为4的并行处理结构，并保持了与原始系统相同的采样速率。这里假定该原始系统工作在5V电源电压，CMOS工艺的阈值电压 V_t 为0.4V。计算这个并行-流水线系统与原始系统相比的功耗。该并行-流水线系统的电源电压是多少？

解：流水级数 $M=4$ 并行级数 $L=4$

$$\left\{ \begin{array}{l} T = \frac{C_{charge} V_o}{k(V_o - V_t)^2} \\ LT = \frac{C_{charge}/M \cdot \beta V_o}{k(\beta V_o - V_t)^2} \end{array} \right. \text{而} \left\{ \begin{array}{l} V_o = 5V \\ V_t = 0.4V \end{array} \right.$$

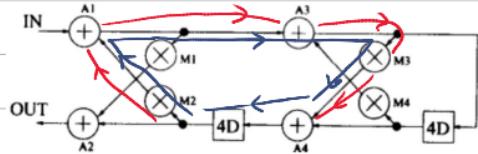
解得 $\beta = 0.177$ 功耗降低为原来的 $\beta^2 = 0.031$
 $V_o' = \beta V_o = 0.885V$

Chapter 4 重定时

3.

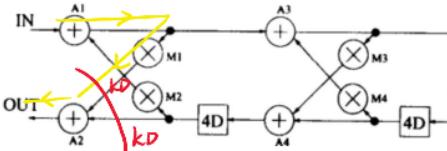
1) 迭代边界，由圈中环决定

$$T_{\infty} = \frac{3+4}{4} = \frac{7}{4}$$

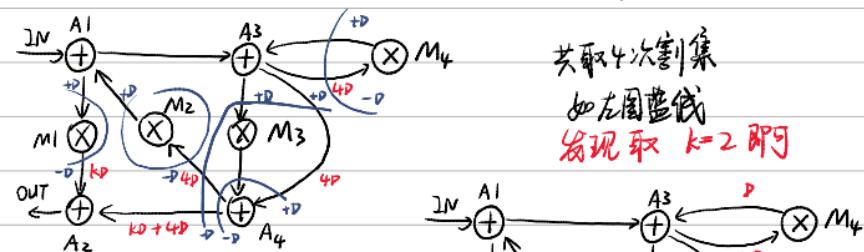
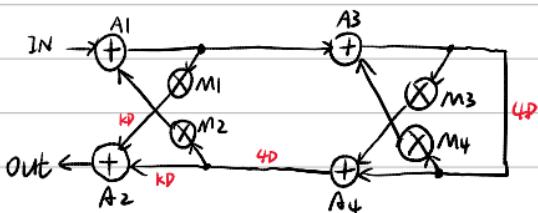


2) 关键路径，由圈中红箭头决定，为 $3+4=7$

3) 黄色路径，由输入至输出无延时，故需增加流水线
取红箭作前端割集
为方便，设割集增加 k 个延时

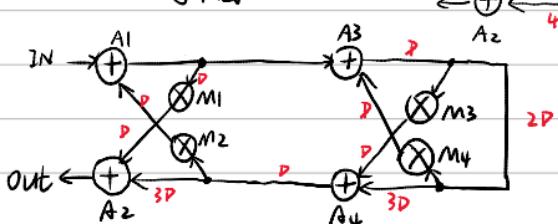


DFG 重绘为



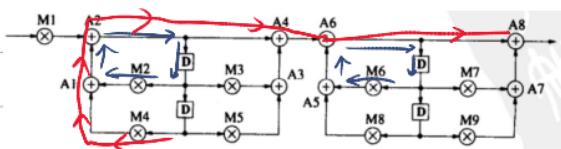
最终得到右图

与下图



5

解: (a) 关键路径如图红色标注 即 $M1 \rightarrow A1 \rightarrow A2 \rightarrow A4 \rightarrow A6 \rightarrow A8$ 为 7 u.t.

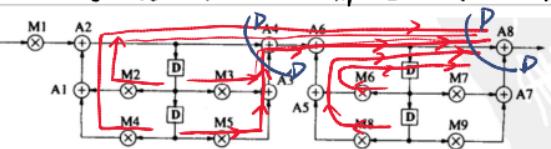


迭代边界为圆中两环

决策

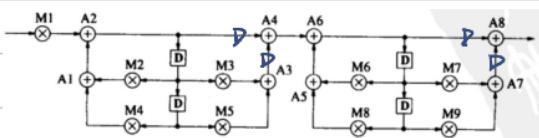
$$T_{\infty} = \frac{4}{1} = 4$$

(b) 由于 $T_{\infty} = 4$ 故仅需令关键路径降低为 4 及以下



圆中红色路径均大于 4

取割集为蓝线



如左图

最小时钟为 4 u.t.

Chapter 5 展开

2. (a) $J=2$ 展开

$$w=10 : i=0 \quad A_0 \rightarrow B_0 \quad \lfloor \frac{0+10}{2} \rfloor = 5$$

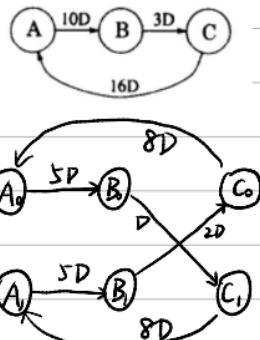
$$i=1 \quad A_1 \rightarrow B_1 \quad \lfloor \frac{1+10}{2} \rfloor = 5$$

$$w=3 : \quad i=0 \quad B_0 \rightarrow C_1 \quad \lfloor \frac{0+3}{2} \rfloor = 1$$

$$i=1 \quad B_1 \rightarrow C_0 \quad \lfloor \frac{1+3}{2} \rfloor = 2$$

$$w=16 : \quad i=0 \quad C_0 \rightarrow A_0 \quad \lfloor \frac{0+16}{2} \rfloor = 8$$

$$i=1 \quad C_1 \rightarrow A_1 \quad \lfloor \frac{1+16}{2} \rfloor = 8$$



$J=5$ 展开

$w=10 \quad i=0, 1, 2, 3, 4$ 均有 $(i+w)\%J=i$, $\lfloor \frac{i+10}{5} \rfloor = 2$ 故 $A_i \rightarrow B_i$ 上为 $2D$

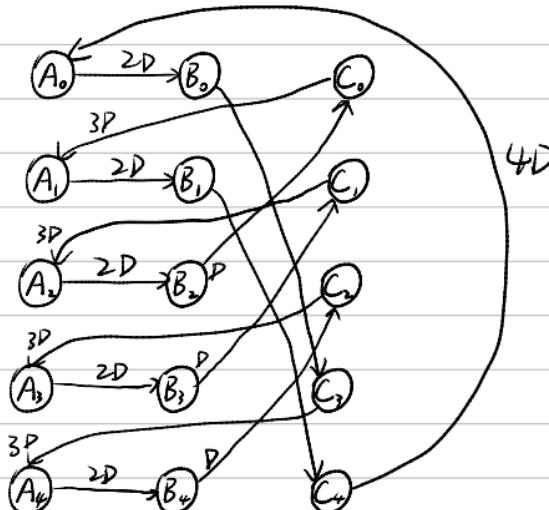
$w=3 \quad i=0, 1$ 有 $(i+w)\%J=i+3$, $\lfloor \frac{i+3}{5} \rfloor = 0$ 故 $B_i \rightarrow C_{i+3}$ 上为 $0D$

$i=2, 3, 4$ 有 $(i+w)\%J=i+2$, $\lfloor \frac{i+3}{5} \rfloor = 1$ 故 $B_i \rightarrow C_{i-2}$ 上为 D

$w=16 \quad i=0, 1, 2, 3$ 有 $(i+w)\%J=i+1$, $\lfloor \frac{i+16}{5} \rfloor = 3$ 故 $C_i \rightarrow A_{i+1}$ 上为 $3D$

$i=4$ 有 $(i+w)\%J=i-4$, $\lfloor \frac{i+16}{5} \rfloor = 4$ 故 $C_i \rightarrow A_{i-4}$ 上为 $4D$

故有



(b) $J=2$

$$w=1 \quad i=0 \quad A_0 \rightarrow B_{(0+1)\%2}(B_1) \text{ 上 } \lfloor \frac{0+1}{2} \rfloor = 0 D$$

$$i=1 \quad A_1 \rightarrow B_{(1+1)\%2}(B_0) \text{ 上 } \lfloor \frac{1+1}{2} \rfloor = 1 D$$

$$w=2 \quad i=0,1 \text{ 均有 } E_i \rightarrow D_i \text{ 上 } \lfloor \frac{i+2}{2} \rfloor = 1 D$$

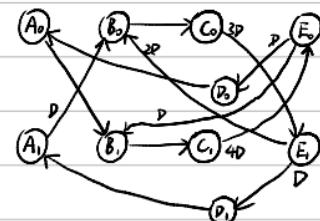
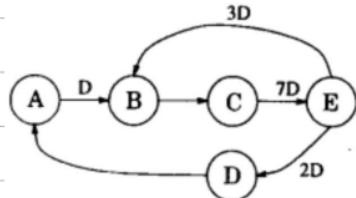
$$w=3 \quad i=0 \quad E_0 \rightarrow B_{(0+3)\%2}(B_1) \text{ 上 } \lfloor \frac{0+3}{2} \rfloor = 1 D$$

$$i=1 \quad E_1 \rightarrow B_{(1+3)\%2}(B_0) \text{ 上 } \lfloor \frac{1+3}{2} \rfloor = 2 D$$

$$w=7 \quad i=0 \quad C_0 \rightarrow E_{(0+7)\%2}(E_1) \text{ 上 } \lfloor \frac{0+7}{2} \rfloor = 3 D$$

$$i=1 \quad C_1 \rightarrow E_{(1+7)\%2}(E_0) \text{ 上 } \lfloor \frac{1+7}{2} \rfloor = 4 D$$

故最终为右图



$J=5$

$$w=1 \quad i=0,1,2,3 \text{ 有 } (i+w)\%J = i+1 \text{ 故 } A_i \rightarrow B_{i+1} \text{ 上 } \lfloor \frac{i+w}{J} \rfloor = 0 D$$

$$i=4 \quad \text{有 } (i+w)\%J = i+4 \text{ 故 } A_i \rightarrow B_{i+4} \text{ 上 } \lfloor \frac{i+w}{J} \rfloor = 1 D$$

$$w=2 \quad i=0,1,2 \text{ 有 } (i+w)\%J = i+2 \text{ 故 } E_i \rightarrow D_{i+2} \text{ 上 } \lfloor \frac{i+w}{J} \rfloor = 0 D$$

$$i=3,4 \quad \text{有 } (i+w)\%J = i+3 \text{ 故 } E_i \rightarrow D_{i+3} \text{ 上 } \lfloor \frac{i+w}{J} \rfloor = 1 D$$

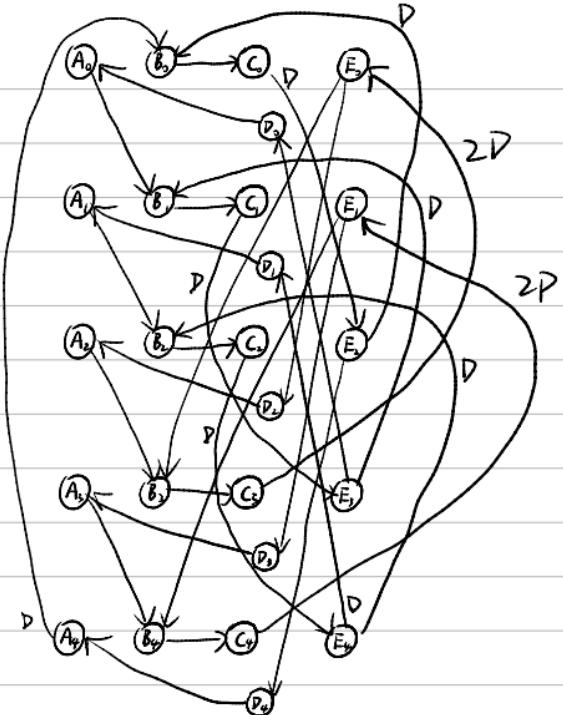
$$w=3 \quad i=0,1 \quad \text{有 } (i+w)\%J = i+3 \text{ 故 } E_i \rightarrow B_{i+3} \text{ 上 } \lfloor \frac{i+w}{J} \rfloor = 0 D$$

$$i=2,3,4 \quad \text{有 } (i+w)\%J = i+2 \text{ 故 } E_i \rightarrow B_{i+2} \text{ 上 } \lfloor \frac{i+w}{J} \rfloor = 1 D$$

$$w=7 \quad i=0,1,2 \quad \text{有 } (i+w)\%J = i+2 \text{ 故 } C_i \rightarrow E_{i+2} \text{ 上 } \lfloor \frac{i+w}{J} \rfloor = 1 D$$

$$i=3,4 \quad \text{有 } (i+w)\%J = i+3 \text{ 故 } C_i \rightarrow E_{i+3} \text{ 上 } \lfloor \frac{i+w}{J} \rfloor = 2 D$$

故最终得到

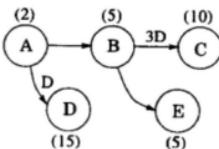


迭代算法：

7. 迭代，直到 $T_{crit} < JT$

$$J = \left\lceil \frac{T_{crit}}{T} \right\rceil$$

$T_{crit} = J$ 阶展开DFG的关键路径



解： $J=2$ 时 存在 $D_i \rightarrow A_0$ 路径为 1] $T_{crit} \geq 17$
 故 $T_{crit} \geq JT = 16$ 不成立

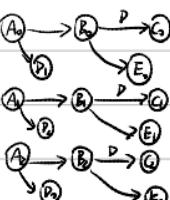
考虑 $J=3$ $w=1$ $i=0,1$ 有 $(i+w)\%J=i+1$ 故 $A_i \rightarrow D_{i+1}$ 上 $L^{\frac{i+w}{J}}=1D$

$i=2$ 有 $(i+w)\%J=i+2$ 故 $A_i \rightarrow D_{i+2}$ 上 $L^{\frac{i+w}{J}}=1D$

$w=3$ $\forall i$ 有 $(i+w)\%J=i$ 故 $B_i \rightarrow C_i$ 上 $L^{\frac{i+w}{J}}=1D$

故最终 $T_{crit} = 17 \leq 3 \times 8$

(D的不标准
作了修改)



故丁最小值为 3

20

解: $J=2$

$$[2, 3] = 6$$

则先扩展为右图
下做展开

$$6l+0 = 2(3l+0) + 0$$

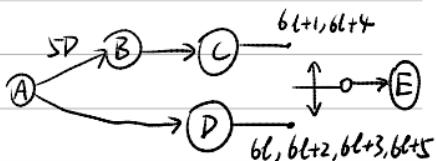
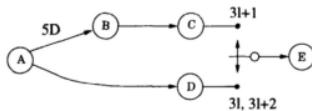
$$> 6l+1 = 2(3l+0) + 1$$

$$6l+2 = 2(3l+1) + 0$$

$$6l+3 = 2(3l+1) + 1$$

$$> 6l+4 = 2(3l+2) + 0$$

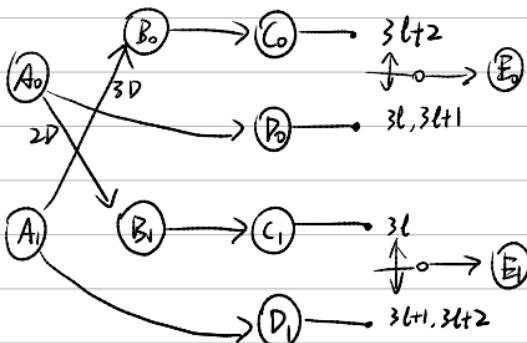
$$6l+5 = 2(3l+2) + 1$$

对 $A \rightarrow B$ 的 $w=5$

$$i=0 \quad (i+w)\%J=1$$

故 $A_0 \rightarrow B_1$ 有 $\lfloor \frac{0+5}{2} \rfloor = 2D$

$$i=1 \quad (i+w)\%J=0$$

故 $A_1 \rightarrow B_0$ 有 $\lfloor \frac{1+5}{2} \rfloor = 3D$ 

Chapter 7 脉动阵列

6.

$$\text{解: } d = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \quad p^T = \begin{pmatrix} 0 & 1 \end{pmatrix} \quad s^T = \begin{pmatrix} 2 & 1 \end{pmatrix}$$

表格

$$e^T \quad p^T e \quad s^T e$$

均满足 $s^T e \geq 0$

$$wt(1 \ 0) \quad 0 \quad 2$$

则 $\left\{ \begin{array}{l} \text{权重保持} \\ \text{输入右移} \end{array} \right.$

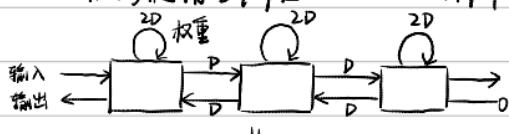
$$ip(0 \ 1) \quad 1 \quad 1$$

$\left\{ \begin{array}{l} \text{输出左移} \\ \text{输出左移} \end{array} \right.$

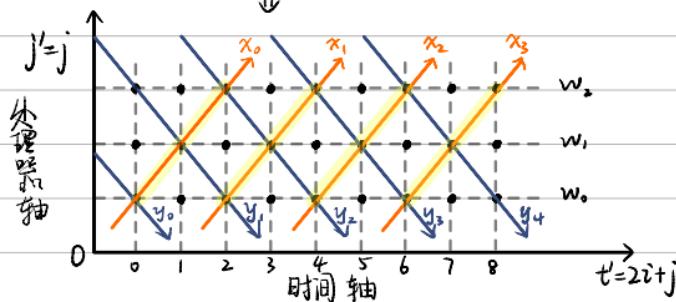
$$\text{result } (1 \ -1) \quad -1 \quad 1$$

$$\text{利用率 } HEU = \frac{1}{15 \cdot 1} = \frac{1}{2}$$

仍考虑用 3 个 PE



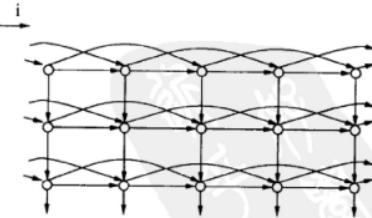
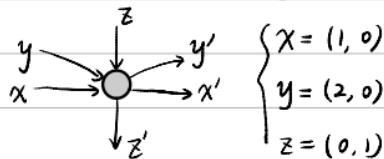
$$\begin{pmatrix} i' \\ j' \\ t' \end{pmatrix} = \begin{pmatrix} 0 & 0 & 1 \\ p^T & 0 \\ s^T & 0 \end{pmatrix} \begin{pmatrix} i \\ j \\ t \end{pmatrix} = \begin{pmatrix} t \\ j \\ 2i+j \end{pmatrix}$$



10.

解：不关心数据含义，仅关心单节点与对应边

标末边(对应向量)为 x, y, z



(a) 投影矢量 $d^T = (0 \ 1)$ 则 $P^T = (k \ 0)$, 考虑 $k=1$

须使硬件利用率最大 \Rightarrow 考虑 $S^T = (t \ 1)$, 不妨 $t=1$

则 $\begin{cases} P^T = (1 \ 0) \\ S^T = (1 \ 1) \end{cases}$ 有 $HUE = \frac{1}{|S^T d|} = 1$ 表格 e p^Te s^Te

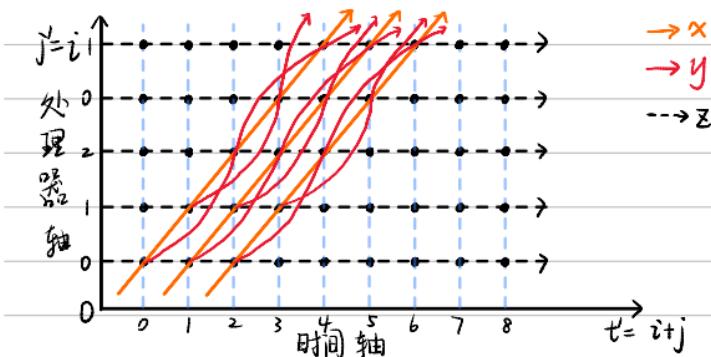
$\vec{s} \cdot \vec{p}$ 不平行 $x: \begin{pmatrix} 1 \\ 0 \end{pmatrix} \quad 1 \quad 1$
 $y: \begin{pmatrix} 2 \\ 0 \end{pmatrix} \quad 2 \quad 2$

显然 $S^T e \geq 0$ $z: \begin{pmatrix} 0 \\ 1 \end{pmatrix} \quad 0 \quad 1$

$$\binom{i'}{j'} = T \binom{i}{j} = \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \end{pmatrix} \binom{i}{j} = \binom{t}{i+j}$$

故 (i, j) 节点在处理器 i 上于 $(i+j)$ 时间执行

由于同一时间最多 3 个节点被执行，故最少需 3 个 PE



故 $d^T = (0 \ 1)$ 时 可取 $P = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$ $S = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$

硬件利用率 100% 需 3 个处理单元

(b) 投影矢量 $d^T = (1 \ 1)$ 则 $p^T = (k \ k)$, 考虑 $k=1$

须使硬件利用率最大 \Rightarrow 不妨考虑 $s^T = (1 \ 0)$

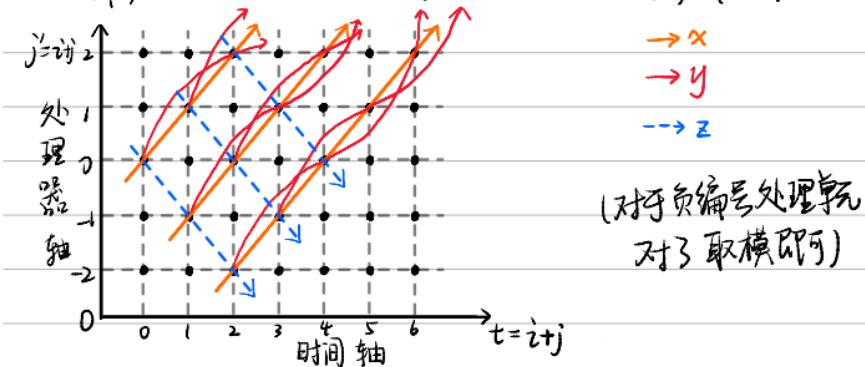
$$\text{则 } \begin{cases} p^T = (1 \ -1) \\ s^T = (1 \ 1) \end{cases} \text{ 有 } HUE = \frac{1}{\|s^T d\|} = \frac{1}{2} \text{ 表格 e } p^T e \quad s^T e$$

显然 $s^T e \geq 0$

$$\binom{i'}{j'} = T \binom{i}{j} = \begin{pmatrix} 0 & 0 & 1 \\ 1 & -1 & 0 \\ 1 & 1 & 0 \end{pmatrix} \binom{i}{j} = \binom{t}{i-j} \quad \binom{i}{i+j}$$

故 (i, j) 节点在处理器 $(i-j)$ 上于 $(i+j)$ 时间执行

同样，同一时间最多 3 个节点被执行，故最少需 3 个 PE



故 $d^T = (1 \ 1)$ 时 可取 $p = \begin{pmatrix} 1 \\ -1 \end{pmatrix}$ $s = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$

硬件利用率 50% 需 3 个处理单元

Chapter 8 快速卷积

5. $\beta_i = 0, 1, -1, 2, -2$

解：对 3×3 卷积 $S(p) = h(p) \otimes (p)$ 其中 $\begin{cases} h(p) = h_0 + h_1 p + h_2 p^2 \\ \otimes(p) = x_0 + x_1 p + x_2 p^2 \\ S(p) = s_0 + s_1 p + s_2 p^2 + s_3 p^3 + s_4 p^4 \end{cases}$

直接实现方式为 $\begin{bmatrix} s_0 \\ s_1 \\ s_2 \\ s_3 \\ s_4 \end{bmatrix} = \begin{bmatrix} h_0 & 0 & 0 \\ h_1 & h_0 & 0 \\ h_2 & h_1 & h_0 \\ 0 & h_2 & h_1 \\ 0 & 0 & h_0 \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \\ x_2 \end{bmatrix}$

下计算 $h(\beta_i)$ 与 $x(\beta_i)$

$$\beta_0 = 0$$

$$h(\beta_0) = h_0$$

$$x(\beta_0) = x_0$$

$$\beta_1 = 1$$

$$h(\beta_1) = h_0 + h_1 + h_2$$

$$x(\beta_1) = x_0 + x_1 + x_2$$

$$\beta_2 = -1$$

$$h(\beta_2) = h_0 - h_1 + h_2$$

$$x(\beta_2) = x_0 - x_1 + x_2$$

$$\beta_3 = 2$$

$$h(\beta_3) = h_0 + 2h_1 + 4h_2$$

$$x(\beta_3) = x_0 + 2x_1 + 4x_2$$

$$\beta_4 = -2$$

$$h(\beta_4) = h_0 - 2h_1 + 4h_2$$

$$x(\beta_4) = x_0 - 2x_1 + 4x_2$$

而 $S(\beta_i) = h(\beta_i) \cdot x(\beta_i)$ ($i=0, 1, 2, 3, 4$) 需 5 次乘法

$$\begin{aligned} \text{故 } S(p) &= \sum_{i=0}^4 S(\beta_i) \frac{p - \beta_i}{p - \beta_j} \\ &= S(\beta_0) \frac{(p-1)(p+1)(p-2)(p+2)}{(p-0)(p+1)(p-1)(p+2)} + S(\beta_1) \frac{(p-1)(p+1)(p-2)(p+2)}{(p-1)(p+1)(p-1)(p+2)} + S(\beta_2) \frac{(p-1)(p+1)(p-2)(p+2)}{(p-1)(p+1)(p-1)(p+2)} + S(\beta_3) \frac{(p-1)(p+1)(p-2)(p+2)}{(p-1)(p+1)(p-1)(p+2)} + S(\beta_4) \frac{(p-1)(p+1)(p-2)(p+2)}{(p-1)(p+1)(p-1)(p+2)} \\ &= \frac{s(\beta_0)}{4!} (p^4 - 5p^2 + 4) + \frac{s(\beta_1)}{-6} (p^4 + p^3 - 4p^2 - 4p) + \frac{s(\beta_2)}{-6} (p^4 - p^3 + 4p^2 + 4p) + \frac{s(\beta_3)}{24} (p^4 + 2p^3 - p^2 - 2p) + \frac{s(\beta_4)}{24} (p^4 - 2p^3 - p^2 + 2p) \\ &= s(\beta_0) + p \left(\frac{2s(\beta_1)}{3} - \frac{2s(\beta_2)}{3} - \frac{s(\beta_3)}{12} + \frac{s(\beta_4)}{12} \right) + p^2 \left(-\frac{5s(\beta_0)}{4} + \frac{2s(\beta_1)}{3} + \frac{2s(\beta_2)}{3} - \frac{s(\beta_3)}{24} - \frac{s(\beta_4)}{24} \right) \\ &\quad + p^3 \left(-\frac{s(\beta_0)}{6} + \frac{s(\beta_1)}{6} + \frac{s(\beta_2)}{12} - \frac{s(\beta_3)}{12} \right) + p^4 \left(\frac{s(\beta_0)}{4} - \frac{s(\beta_1)}{6} - \frac{s(\beta_2)}{6} + \frac{s(\beta_3)}{24} + \frac{s(\beta_4)}{24} \right) \end{aligned}$$

最后 转化为 矩阵-矢量形式

$$\begin{bmatrix} s_0 \\ s_1 \\ s_2 \\ s_3 \\ s_4 \end{bmatrix} = \begin{bmatrix} 4 & 0 & 0 & 0 & 0 \\ 0 & 4 & 4 & -2 & 2 \\ -5 & 4 & 4 & -1 & -1 \\ 0 & -1 & 1 & 2 & -2 \\ 1 & -1 & -1 & 1 & 1 \end{bmatrix} \begin{bmatrix} s(\beta_0)/4 \\ s(\beta_1)/6 \\ s(\beta_2)/12 \\ s(\beta_3)/24 \\ s(\beta_4)/48 \end{bmatrix} = \begin{bmatrix} 4 & 0 & 0 & 0 & 0 \\ 0 & 4 & -4 & -2 & 2 \\ -5 & 4 & 4 & -1 & -1 \\ 0 & -1 & 1 & 2 & -2 \\ 1 & -1 & -1 & 1 & 1 \end{bmatrix} \begin{bmatrix} \frac{h_0}{4} \\ \frac{h_1}{6} \\ \frac{h_2}{12} \\ \frac{h_3}{24} \\ \frac{h_4}{48} \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 1 \\ 1 & -1 & 1 \\ 1 & 2 & 4 \\ 1 & -2 & 4 \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \\ x_2 \end{bmatrix}$$

C

H

D

(1) $H = \text{diag}(H_0, H_1, H_2, H_3, H_4)$ 事先算出

(2) $X_i = x(\beta_i)$ 为 Dx 计算得到

(3) $S_i = H_i X_i$ 为对矩阵 H 计算得到

(4) 最后对 $H(Dx)$ 左乘 (得到) $\vec{S} = \begin{bmatrix} s_0 \\ s_1 \\ s_2 \\ s_3 \\ s_4 \end{bmatrix}$

\downarrow
5 次乘法

$$6. \beta_i = 0, 1, -1, 2$$

解: $s(p) = h(p) \cdot x(p) = s_0 + s_1 p + s_2 p^2 + s_3 p^3 + s_4 p^4$ 且 $s_4 p^4 = h_2 x_2 p^4$
 $\therefore s(p) = s(p) - h_2 x_2 p^4$

下计算 $h(\beta_i)$ 、 $x(\beta_i)$ 与 $s'(\beta_i)$

$$\beta_0 = 0 \quad h(\beta_0) = h_0 \quad x(\beta_0) = x_0$$

$$\beta_1 = 1 \quad h(\beta_1) = h_0 + h_1 + h_2 \quad x(\beta_1) = x_0 + x_1 + x_2$$

$$\beta_2 = -1 \quad h(\beta_2) = h_0 - h_1 + h_2 \quad x(\beta_2) = x_0 - x_1 + x_2$$

$$\beta_3 = 2 \quad h(\beta_3) = h_0 + 2h_1 + 4h_2 \quad x(\beta_3) = x_0 + 2x_1 + 4x_2$$

而 $s'(\beta_i) = h(\beta_i) \cdot x(\beta_i) - h_1 x_1 \beta_i^4 \quad (i=0, 1, 2, 3)$

下先求 $s'(p) \quad (s(p) = s'(p) + h_1 x_1 p^2)$

$$\begin{aligned} s(p) &= \sum_{i=0}^3 s(\beta_i) \prod_{j \neq i} \frac{p - \beta_j}{\beta_i - \beta_j} \\ &= s(\beta_0) \frac{(p-0)(p+1)(p-2)}{(0-1)(0+1)(0-2)} + s(\beta_1) \frac{(p-1)(p+1)(p-2)}{(1-0)(1+1)(1-2)} + s(\beta_2) \frac{(p+1)(p-1)(p-2)}{(2+0)(2-1)(2-2)} + s(\beta_3) \frac{(p+2)(p-2)(p-1)}{(2-0)(2-1)(2-2)} \\ &= \frac{s(\beta_0)}{2} (p^3 - 3p^2 + 2p) + \frac{s(\beta_1)}{2} (p^3 - p^2 - 2p) + \frac{s(\beta_2)}{6} (p^3 - 3p^2 + 2p) + \frac{s(\beta_3)}{6} (p^3 - p) \\ &= s(\beta_0) + p \left(\frac{s(\beta_1)}{2} + s(\beta_2) - \frac{s(\beta_3)}{6} \right) + p^2 \left(-s(\beta_0) + \frac{s(\beta_1)}{2} + \frac{s(\beta_2)}{2} \right) + p^3 \left(\frac{s(\beta_0)}{2} - \frac{s(\beta_1)}{6} + \frac{s(\beta_3)}{6} \right) \end{aligned}$$

故

s_0	$2 \ 0 \ 0 \ 0 \ 0$
s_1	$-1 \ 2 \ -2 \ -1 \ 0$
s_2	$-2 \ 1 \ 3 \ 0 \ 0$
s_3	$1 \ -1 \ 1 \ 0 \ 0$
s_4	$0 \ 0 \ 0 \ 0 \ 1$

$s(\beta_0)/2$	$s(\beta_1)/2$	$s(\beta_2)/2$	$s(\beta_3)/6$
$s(\beta_0)/6$	$s(\beta_1)/6$	$s(\beta_2)/6$	$s(\beta_3)/6$
$s(\beta_0)/2$	$s(\beta_1)/2$	$s(\beta_2)/2$	$s(\beta_3)/6$
$s(\beta_0)/6$	$s(\beta_1)/6$	$s(\beta_2)/6$	$s(\beta_3)/6$
$h_2 x_2$	$h_2 x_2$	$h_2 x_2$	$h_2 x_2$

因而后置加法矩阵 $C = \begin{bmatrix} 2 & 0 & 0 & 0 & 0 \\ -1 & 2 & -2 & 0 & 0 \\ 2 & 1 & 3 & 0 & 0 \\ 1 & -1 & -1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & -\frac{1}{2} & 0 \\ 0 & 0 & 1 & 0 & -\frac{1}{6} \\ 0 & 0 & 0 & 1 & -\frac{1}{3} \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 2 & 0 & 0 & 0 & 0 \\ -1 & 2 & -1 & 2 & 0 \\ -2 & 1 & 3 & 0 & -1 \\ 1 & -1 & -1 & 2 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$

而 $H = \begin{bmatrix} \frac{h_0}{2} & 0 & 0 & 0 & 0 \\ 0 & \frac{h_1}{2} & 0 & 0 & 0 \\ 0 & 0 & \frac{h_2}{2} & 0 & 0 \\ 0 & 0 & 0 & \frac{h_3}{6} & 0 \\ 0 & 0 & 0 & 0 & h_4 \end{bmatrix} = \text{diag}(H_0, H_1, H_2, H_3, H_4)$ 后置加法矩阵 $D = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 1 \\ 1 & -1 & 1 \\ 1 & 2 & 4 \\ 0 & 0 & 1 \end{bmatrix}$

故最终

$$\tilde{s} = CHD\vec{x} = \begin{bmatrix} 2 & 0 & 0 & 0 & 0 \\ -1 & 2 & -2 & -1 & 2 \\ 2 & 1 & 3 & 0 & -1 \\ 1 & -1 & -1 & 1 & -2 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} \frac{h_0}{2} & 0 & 0 & 0 & 0 \\ 0 & \frac{h_1}{2} & 0 & 0 & 0 \\ 0 & 0 & \frac{h_2}{2} & 0 & 0 \\ 0 & 0 & 0 & \frac{h_3}{6} & 0 \\ 0 & 0 & 0 & 0 & h_4 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 1 \\ 1 & -1 & 1 \\ 1 & 2 & 4 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix}$$

(1) $H = \text{diag}(H_0, H_1, H_2, H_3, H_4)$ 事先算出

(2) 先对 \vec{x} 左乘 D 得到 $D\vec{x}$ ($X_i = x(\beta_i)$)

(3) $S_i = H_i X_i$ 为对 $D\vec{x}$ 左乘 H 计算得到

(4) 最后对 $H(D\vec{x})$ 左乘 C 得到 $\tilde{s} = \begin{bmatrix} s_0 \\ s_1 \\ s_2 \\ s_3 \\ s_4 \end{bmatrix}$

↓
5次乘法

11.

$$\text{解: } m(p) = p(p-1)(p+1)$$

$$\begin{cases} m^{(0)}(p) = p \\ m^{(1)}(p) = p-1 \\ m^{(2)}(p) = p+1 \end{cases} \quad \begin{cases} M^{(0)}(p) = \frac{m(p)}{m^{(0)}(p)} = p^2-1 \\ M^{(1)}(p) = \frac{m(p)}{m^{(1)}(p)} = p^2+p \\ M^{(2)}(p) = \frac{m(p)}{m^{(2)}(p)} = p^2-p \end{cases}$$

由 $N^{(2)}(p) M^{(2)}(p) + n^{(2)}(p) m^{(2)}(p) = 1$ 可计算并取 $N^{(2)}(p)$ 与 $n^{(2)}(p)$ 如下

	$m^{(2)}(p)$	$M^{(2)}(p)$	$n^{(2)}(p)$	$N^{(2)}(p)$
0	p	p^2-1	$-p$	-1
1	$p-1$	p^2+p	$-\frac{1}{2}(p+2)$	$\frac{1}{2}$
2	$p+1$	p^2-p	$-\frac{1}{2}(p-2)$	$\frac{1}{2}$

$$\text{得 } h^{(2)}(p) = h(p) \bmod m^{(2)}(p) \quad x^{(2)}(p) = x^{(2)}(p) \bmod m^{(2)}(p)$$

$$\text{有 } \begin{cases} h^{(0)}(p) = h_0 \\ h^{(1)}(p) = h_0 + h_1 \\ h^{(2)}(p) = h_0 - h_1 \end{cases} \text{ 与 } \begin{cases} x^{(0)}(p) = x_0 \\ x^{(1)}(p) = x_0 + x_1 \\ x^{(2)}(p) = x_0 - x_1 \end{cases} \text{ 而 } s^{(2)}(p) = h^{(2)}(p) \cdot x^{(2)}(p) \bmod m^{(2)}(p)$$

$$\begin{cases} s^{(0)}(p) = h_0 x_0 \\ s^{(1)}(p) = (h_0 + h_1)(x_0 + x_1) \\ s^{(2)}(p) = (h_0 - h_1)(x_0 - x_1) \end{cases} \Rightarrow \begin{bmatrix} s^{(0)}(p) \\ s^{(1)}(p) \\ s^{(2)}(p) \end{bmatrix} = \begin{bmatrix} h_0 & 0 & 0 \\ 0 & h_0 + h_1 & 0 \\ 0 & 0 & h_0 - h_1 \end{bmatrix} \begin{bmatrix} x_0 \\ x_0 + x_1 \\ x_0 - x_1 \end{bmatrix} = \begin{bmatrix} h_0 & 0 & 0 \\ 0 & h_0 + h_1 & 0 \\ 0 & 0 & h_0 - h_1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \end{bmatrix}$$

$$\begin{aligned} \text{由 } s(p) &= \sum_{i=0}^2 s^{(i)}(p) N^{(i)}(p) M^{(i)}(p) \bmod m(p) \\ &= -s^{(0)}(p^2-1) + \frac{s^{(1)}(p)}{2}(p^2+p) + \frac{s^{(2)}(p)}{2}(p^2-p) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & \frac{1}{2} & -\frac{1}{2} \\ -1 & \frac{1}{2} & \frac{1}{2} \end{bmatrix} \begin{bmatrix} s^{(0)}(p) \\ s^{(1)}(p) \\ s^{(2)}(p) \end{bmatrix} = \begin{bmatrix} 1 & 0 & p^2 \\ 1 & p & p^2 \end{bmatrix} \end{aligned}$$

从而得到

$$\vec{s} = \begin{bmatrix} s_0 \\ s_1 \\ s_2 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & \frac{1}{2} & -\frac{1}{2} \\ -1 & \frac{1}{2} & \frac{1}{2} \end{bmatrix} \begin{bmatrix} h_0 & 0 & 0 \\ 0 & h_0 + h_1 & 0 \\ 0 & 0 & h_0 - h_1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 1 & 1 \\ 1 & -1 \end{bmatrix} \vec{x} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & -1 \\ -2 & 1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & \frac{1}{2} & 0 \\ 0 & 0 & \frac{1}{2} \end{bmatrix} \begin{bmatrix} h_0 & 0 & 0 \\ 0 & h_0 + h_1 & 0 \\ 0 & 0 & h_0 - h_1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \end{bmatrix}$$

$$\text{即 } \vec{s} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & -1 \\ -2 & 1 & 1 \end{bmatrix} \begin{bmatrix} h_0 & 0 & 0 \\ 0 & \frac{1}{2} & 0 \\ 0 & 0 & \frac{1}{2} \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 1 & 1 \\ 1 & -1 \end{bmatrix} \vec{x} \quad \text{即 } c = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & -1 \\ -2 & 1 & 1 \end{bmatrix} H = \begin{bmatrix} 1 & 0 & 0 \\ 0 & \frac{1}{2} & 0 \\ 0 & 0 & \frac{1}{2} \end{bmatrix} = \text{diag}(H_0, H_1, H_2) \quad D = \begin{bmatrix} 1 & 0 \\ 0 & 1 \\ 1 & -1 \end{bmatrix}$$

(1) $H = \text{diag}(H_0, H_1, H_2)$ 事先算出

(2) 先对 \vec{x} 左乘 D 得 $D\vec{x}$

(3) 再对 $D\vec{x}$ 左乘 H 得到 $H(D\vec{x})$, 此时有3个系数

(4) 再左乘后置加法矩阵 C 得到了

Chapter 9

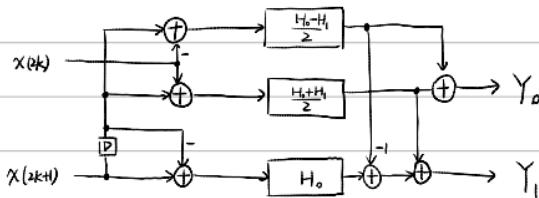
5.

$$\begin{aligned} h_0x_0 &= h_0x_0 \\ h_0x_1 + h_1x_0 &= \frac{h_0+h_1}{2}(x_0+x_1) - \frac{h_0-h_1}{2}(x_0-x_1) \\ h_1x_1 &= -h_0x_0 + \frac{h_0+h_1}{2}(x_0+x_1) + \frac{h_0-h_1}{2}(x_0-x_1) \end{aligned}$$

解：先将差分方程写为矩阵形式（题到） $\vec{s} = CHA\vec{x}$

$$\begin{bmatrix} S_2 \\ S_1 \\ S_0 \end{bmatrix} = \begin{bmatrix} h_1 & 0 \\ h_0 & h_1 \\ 0 & h_0 \end{bmatrix} \begin{bmatrix} X_1 \\ X_0 \\ X_0 \end{bmatrix} = \begin{bmatrix} 1 & 1 & -1 \\ -1 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \text{diag} \begin{bmatrix} \frac{h_0+h_1}{2} \\ \frac{h_0-h_1}{2} \\ h_0 \end{bmatrix} \begin{bmatrix} 1 & -1 \\ 1 & 1 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} X_1 \\ X_0 \\ X_0 \end{bmatrix}$$

而 $Y = (CHA)^T X$ 故 $\begin{bmatrix} Y_0 \\ Y_1 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 0 \\ -1 & 1 & 1 \end{bmatrix} \text{diag} \begin{bmatrix} \frac{h_0+h_1}{2} \\ \frac{h_0-h_1}{2} \\ H_0 \end{bmatrix} \begin{bmatrix} 1 & -1 & 0 \\ 1 & 1 & 0 \\ -1 & 0 & 1 \end{bmatrix} \begin{bmatrix} z^2X_1 \\ X_0 \\ X_1 \end{bmatrix}$

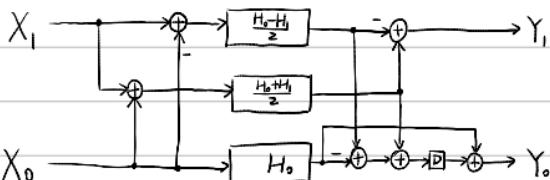


6

解：5中结构可进一步表示为 $Y = \begin{bmatrix} 1 & 1 & 0 \\ -1 & 1 & 1 \end{bmatrix} \text{diag} \begin{bmatrix} \frac{h_0-h_1}{2} \\ \frac{h_0+h_1}{2} \\ H_0 \end{bmatrix} \begin{bmatrix} 1 & z^{-2} \\ 1 & z^{-2} \\ 0 & 1-z^{-2} \end{bmatrix} X$

$$\text{则 } Q^T = \begin{bmatrix} 1 & -1 \\ 1 & 1 \\ 0 & 1 \end{bmatrix} \quad P^T = \begin{bmatrix} -1 & 1 & 0 \\ z^2 & z^2 & 1-z^2 \end{bmatrix} \quad Y_F = P^T H^T Q^T X_F$$

$$\begin{bmatrix} Y_1 \\ Y_0 \end{bmatrix} = \begin{bmatrix} -1 & 1 & 0 \\ z^2 & z^2 & 1-z^2 \end{bmatrix} \text{diag} \begin{bmatrix} \frac{H_0-H_1}{2} \\ \frac{H_0+H_1}{2} \\ H_0 \end{bmatrix} \begin{bmatrix} 1 & -1 \\ 1 & 1 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} X_1 \\ X_0 \\ X_0 \end{bmatrix}$$



9.

$$Y_0 = X_0 H_0 + z^{-2} X_1 H_1$$

$$Y_1 = X_0 H_0 + X_1 H_1 - (X_0 - X_1)(H_0 - H_1)$$

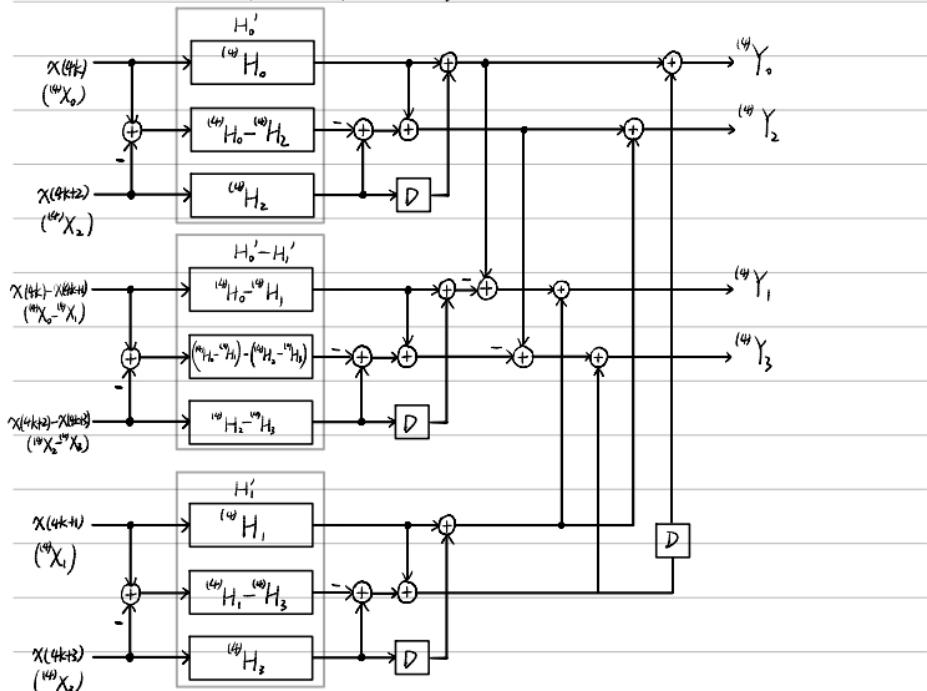
$$\begin{bmatrix} Y_0 \\ Y_1 \end{bmatrix} = \begin{bmatrix} 1 & z^{-2} & 0 \\ 1 & 1 & -1 \end{bmatrix} \text{diag} \begin{bmatrix} H_0 \\ H_1 \\ H_0 - H_1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} X_0 \\ X_1 \end{bmatrix}$$

解：为方便，设4并行FFA输出为 ${}^{(4)}Y = [{}^{(4)}Y_0, {}^{(4)}Y_1, {}^{(4)}Y_2, {}^{(4)}Y_3]^T$ 相应有 ${}^{(4)}X_i$ 与 ${}^{(4)}H_i$

$$\begin{cases} {}^{(4)}Y = (X'_0 + z^{-2} X'_1) (H'_0 + z^{-2} H'_1) \\ \text{其中 } \begin{cases} X'_0 = {}^{(4)}X_0 + z^{-2} {}^{(4)}X_1 \\ X'_1 = {}^{(4)}X_1 + z^{-2} {}^{(4)}X_2 \end{cases} \quad \begin{cases} H'_0 = {}^{(4)}H_0 + z^{-2} {}^{(4)}H_1 \\ H'_1 = {}^{(4)}H_1 + z^{-2} {}^{(4)}H_2 \end{cases} \end{cases}$$

则由该2并行FFA 有 ${}^{(4)}Y = X'_0 H'_0 + z^{-2} [X'_0 H'_0 + X'_1 H'_1 - (X'_0 - X'_1)(H'_0 - H'_1)] + z^{-2} X'_1 H'_1$

$$\begin{aligned} X'_0 H'_0 &= ({}^{(4)}X_0 + z^{-2} {}^{(4)}X_1) ({}^{(4)}H_0 + z^{-2} {}^{(4)}H_1) = {}^{(4)}X_0 {}^{(4)}H_0 + z^{-2} [{}^{(4)}X_0 {}^{(4)}H_0 + {}^{(4)}X_1 {}^{(4)}H_1 - ({}^{(4)}X_0 - {}^{(4)}X_1)({}^{(4)}H_0 - {}^{(4)}H_1)] + z^{-2} {}^{(4)}X_1 {}^{(4)}H_0 \\ X'_1 H'_1 &= ({}^{(4)}X_1 + z^{-2} {}^{(4)}X_2) ({}^{(4)}H_1 + z^{-2} {}^{(4)}H_2) = {}^{(4)}X_1 {}^{(4)}H_1 + z^{-2} [{}^{(4)}X_1 {}^{(4)}H_1 + {}^{(4)}X_2 {}^{(4)}H_2 - ({}^{(4)}X_1 - {}^{(4)}X_2)({}^{(4)}H_1 - {}^{(4)}H_2)] + z^{-2} {}^{(4)}X_2 {}^{(4)}H_1 \\ (X'_0 - X'_1)(H'_0 - H'_1) &= [({}^{(4)}X_0 - {}^{(4)}X_1) + z^{-2} ({}^{(4)}X_1 - {}^{(4)}X_2)][({}^{(4)}H_0 - {}^{(4)}H_1) + z^{-2} ({}^{(4)}H_1 - {}^{(4)}H_2)] \\ &= ({}^{(4)}X_0 - {}^{(4)}X_1) ({}^{(4)}H_0 - {}^{(4)}H_1) + z^{-2} \{ ({}^{(4)}X_0 - {}^{(4)}X_1) ({}^{(4)}H_0 - {}^{(4)}H_1) + ({}^{(4)}X_1 - {}^{(4)}X_2) ({}^{(4)}H_1 - {}^{(4)}H_2) - [({}^{(4)}X_0 - {}^{(4)}X_1) - ({}^{(4)}X_1 - {}^{(4)}X_2)] [({}^{(4)}H_0 - {}^{(4)}H_1) - ({}^{(4)}H_1 - {}^{(4)}H_2)] \} \\ &\quad + z^{-4} ({}^{(4)}X_2 - {}^{(4)}X_3) ({}^{(4)}H_2 - {}^{(4)}H_3) \end{aligned}$$



33.

解：(a) 证明：将(9-53)部分列写如下

$$\begin{aligned} X(1) &= M_0c_1 + M_1c_7 + M_2c_3 + M_3c_5 \\ X(7) &= M_0c_7 - M_1c_1 - M_2c_5 + M_3c_3 \\ X(3) &= M_0c_3 - M_1c_5 - M_2c_7 - M_3c_1 \\ X(5) &= M_0c_5 + M_1c_3 - M_2c_1 + M_3c_7 \\ X(2) &= M_{10}c_2 + M_{11}c_6 \\ X(6) &= M_{10}c_6 - M_{11}c_2 \end{aligned} \left. \begin{array}{l} \text{均为1次乘，3次加} \\ \text{共需 } 4 \times 4 = 16 \text{ 个 CLK} \end{array} \right\}$$

$$\left. \begin{array}{l} \text{均为1次乘，1次乘加} \\ \text{共需 } 2 \times 2 = 4 \text{ 个 CLK} \end{array} \right\}$$

$$\left. \begin{array}{l} \text{共23次} \\ \text{乘法} \\ \text{共23个 CLK} \end{array} \right\}$$

考虑 $X(0)$ 与 $X(4)$

$$\begin{bmatrix} X(0) \\ X(4) \end{bmatrix} = \begin{bmatrix} c_4 & c_4 \\ c_4 & -c_4 \end{bmatrix} \begin{bmatrix} P_{10} \\ P_{11} \end{bmatrix}$$

$$X(0) = c_4 P_{10} + c_4 P_{11} \quad \left. \begin{array}{l} \text{需 1次乘，1次乘加} \\ \text{共3个 CLK} \end{array} \right\}$$

$$X(1) = X(0) - 2c_4 P_{11} \quad \text{需 1次乘加}$$

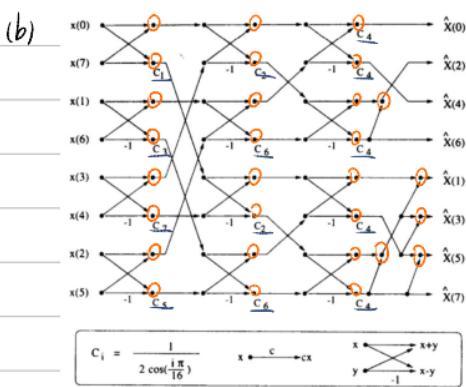
其中，相隔数

由(9-54)知

$$\begin{aligned} M_0 &= x_0 - x_7 & P_0 &= x_0 + x_7 \\ M_1 &= x_3 - x_4 & P_1 &= x_3 + x_4 \\ M_2 &= x_1 - x_6 & P_2 &= x_1 + x_6 \\ M_3 &= x_2 - x_5 & P_3 &= x_2 + x_5 \\ M_{10} &= P_0 - P_1 & P_{10} &= P_0 + P_1 \\ M_{11} &= P_2 - P_3 & P_{11} &= P_2 + P_3 \end{aligned} \left. \begin{array}{l} \text{共12次加法} \\ \text{共需12个 CLK} \end{array} \right\}$$

综上 共 $16 + 4 + 3 + 12 = 35$ 个 CLK

□



共13次乘法 需13个 CLK

共29次加法 需29个 CLK

故总计 42个 CLK

图9-17 8点快速DCT流图

精度高

(c) 对于该DSP，仅考虑时间，则(a)更好因为总耗时较少；但实现(b)中更好，因为一般

乘法耗时更长，虽然(b)中加法更多，但乘法比(a)中少了 $\frac{1}{3}$ ，在一般的DSP上执行得更快

Chapter 10

$$4. \quad H(z) = \frac{1}{1 - \frac{4}{3}z^{-1} + \frac{5}{12}z^{-2}}$$

解：a) 聚类超前 注意到

$$(1 - \frac{4}{3}z^{-1} + \frac{5}{12}z^{-2})(1 + \frac{4}{3}z^{-1} + \frac{49}{36}z^{-2} + \frac{34}{27}z^{-3}) = 1 - \frac{1441}{1296}z^{-4} + \frac{85}{162}z^{-5}$$

$$\therefore H(z) = \frac{1 + \frac{4}{3}z^{-1} + \frac{49}{16}z^{-2} + \frac{34}{27}z^{-3}}{1 - \frac{1441}{1296}z^{-4} + \frac{85}{162}z^{-5}}$$

b) 离散超前 显然 $-\frac{4}{3}z^{-1} + \frac{5}{12}z^{-2}$ 对应 $1 + \frac{4}{3}z^{-1} + \frac{5}{12}z^{-2}$

$$(1 - \frac{4}{3}z^{-1} + \frac{5}{12}z^{-2})(1 + \frac{4}{3}z^{-1} + \frac{5}{12}z^{-2}) = 1 - \frac{17}{18}z^{-2} + \frac{25}{144}z^{-4}$$

同样 $1 - \frac{17}{18}z^{-2} + \frac{25}{144}z^{-4}$ 对应 $1 + \frac{17}{18}z^{-2} + \frac{25}{144}z^{-4}$

$$(1 - \frac{17}{18}z^{-2} + \frac{25}{144}z^{-4})(1 + \frac{17}{18}z^{-2} + \frac{25}{144}z^{-4}) = 1 - \frac{353}{648}z^{-4} + \frac{625}{20736}z^{-8}$$

$$\text{而 } (1 + \frac{4}{3}z^{-1} + \frac{5}{12}z^{-2})(1 + \frac{17}{18}z^{-2} + \frac{25}{144}z^{-4}) = 1 + \frac{4}{3}z^{-1} + \frac{49}{36}z^{-2} + \frac{34}{27}z^{-3} + \frac{245}{432}z^{-4} + \frac{25}{108}z^{-5} + \frac{125}{1728}z^{-6}$$

$$\therefore H(z) = \frac{1 + \frac{4}{3}z^{-1} + \frac{49}{36}z^{-2} + \frac{34}{27}z^{-3} + \frac{245}{432}z^{-4} + \frac{25}{108}z^{-5} + \frac{125}{1728}z^{-6}}{1 - \frac{353}{648}z^{-4} + \frac{625}{20736}z^{-8}}$$

□

$$5. H(z) = \frac{1}{1 - 0.6z^{-1} + 0.25z^{-2}}$$

$$\text{解: } H(z) = \frac{1}{(1 - 0.3z^{-1})^2 + (0.2z^{-2})^2} = \frac{1}{[1 - (0.3 + j0.4)z^{-1}][1 - (0.3 - j0.4)z^{-1}]}$$

为方便记 $P = \frac{1}{2}e^{j\theta} = 0.3 + j0.4$, $\theta = \arctan \frac{0.4}{0.3} \approx 53^\circ$, 基本幅角 $\tilde{\theta} = \frac{\theta}{2} = 0.3 - j0.4$ \boxed{P} $H(z) = \frac{1}{(1 - Pz^{-1})(1 - \tilde{P}z^{-1})}$

对 $M_1 \times M_2 = 3 \times 3$ 分解, 第一级实现 $N(M_1 - 1) = 4$ 个零点, 第二级实现 $NM_1(M_2 - 1) = 12$ 个零点

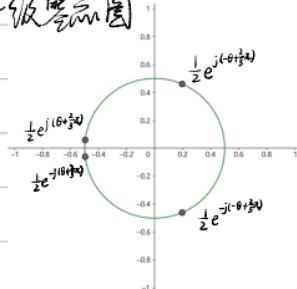
$$H(z) = \frac{1}{(1 - Pz^{-1})(1 - \tilde{P}z^{-1})} = \frac{(1 + Pz^{-1} + P^2z^{-2})(1 + \tilde{P}z^{-1} + \tilde{P}^2z^{-2})}{1 - P^2z^{-2}} \cdot \frac{(1 + 2z^{-1} + 2^2z^{-2})(1 + 2\tilde{P}z^{-1} + 2^2\tilde{P}^2z^{-2})}{1 - 2^2z^{-2}}$$

$$= \underbrace{[(1 + Pz^{-1} + P^2z^{-2})(1 + 2z^{-1} + 2^2z^{-2})]}_{\text{为第一级}} \underbrace{[(1 + \tilde{P}z^{-1} + \tilde{P}^2z^{-2})(1 + 2\tilde{P}z^{-1} + 2^2\tilde{P}^2z^{-2})]}_{\text{为第二级}} \frac{1}{(1 - P^2z^{-2})(1 - 2^2z^{-2})}$$

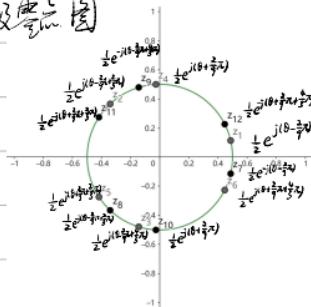
注意到 $\begin{cases} 1 + Pz^{-1} + P^2z^{-2} = (Pz^{-1} - e^{j\frac{2\pi}{3}})(Pz^{-1} - e^{-j\frac{2\pi}{3}}) \\ 1 + 2z^{-1} + 2^2z^{-2} = (2z^{-1} - e^{j\frac{2\pi}{3}})(2z^{-1} - e^{-j\frac{2\pi}{3}}) \end{cases}$ 故算一级零点为 $\left\{ e^{j\frac{\pi}{3}} = \frac{1}{2}e^{j(\frac{2\pi}{3}+0)} \right.$ 与 $\left. e^{-j\frac{\pi}{3}} = \frac{1}{2}e^{j(\frac{2\pi}{3}-0)} \right.$

类似, 得到第二级的零点为 $\frac{1}{2}e^{j[\theta \pm \frac{2}{3}\pi \pm \frac{2k}{3}\pi]} = \frac{1}{2}e^{-j[\theta \pm \frac{2}{3}\pi \pm \frac{2k}{3}\pi]}$ $k = 0, 1, 2$ 共 12 个

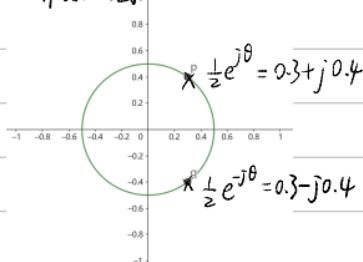
第一级零点图



第二级零点图



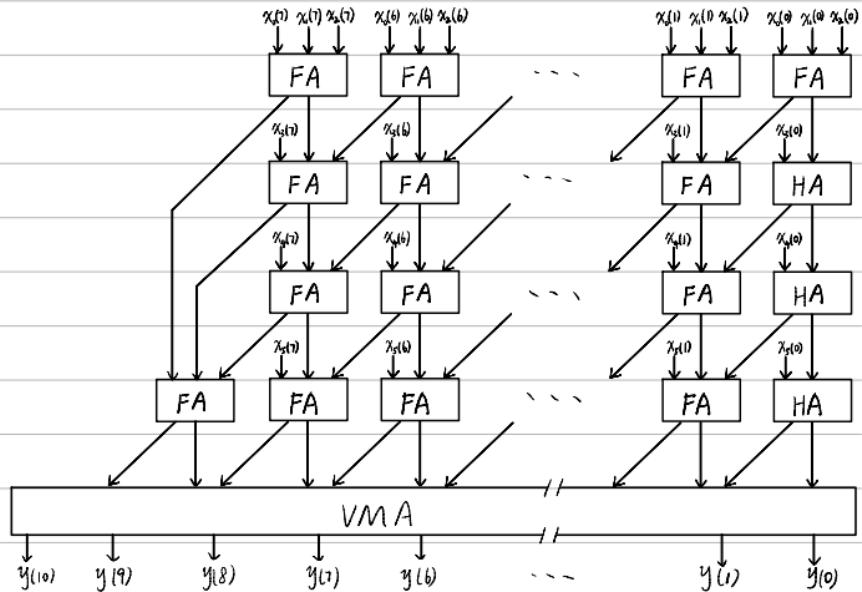
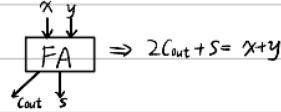
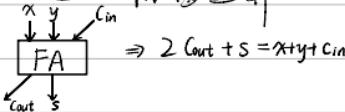
极点图



Chapter 13

3.

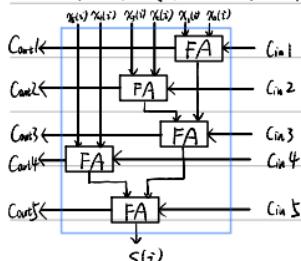
解-(a) 进位保留进算



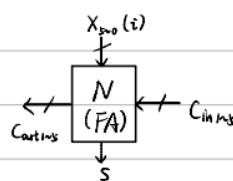
延迟为 $4t_{FA}$

(b) 树高度减小技术

为方便，设每个节点如下. 记为节点 $N(FA)$ 或 $N(HA)$.



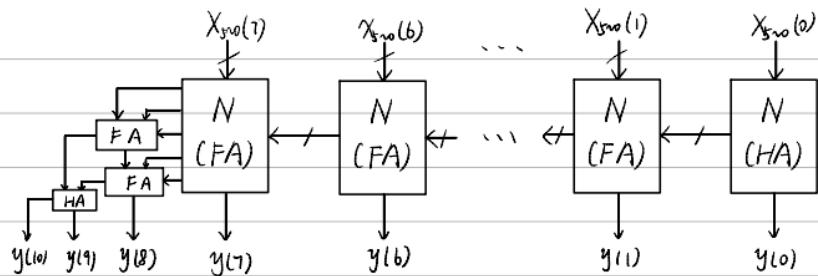
当无 C_{in5} 时用 HA 代替 FA 即可



$$X_{S0}(i) \triangleq \{x_5(i), x_4(i), x_3(i), x_2(i), x_1(i), x_0(i)\}$$

$$C_{out5} \triangleq \{C_{out1}, C_{out2}, C_{out3}, C_{out4}, C_{out5}\}$$

$$C_{in1s} \triangleq \{C_{in1}, C_{in2}, C_{in3}, C_{in4}, C_{in5}\}$$



延迟为 $2t_{HA} + 10t_{FA}$

20.

解: $C_0 = 1.111111\overline{11}$ 将从右向左第一次出现的两位“11”变为01 并向高1位加1

得到 $C_0: 0.000000\overline{01}$ 无“11” 故最终结果即为 0.000000\overline{01}

利用 $\theta_i = \hat{\alpha}_i \oplus \hat{\alpha}_{i+1}$ $\gamma_i = \overline{\gamma_{i+1}}\theta_i$ $\alpha_i = (1 - 2\hat{\alpha}_{i+1})\gamma_i$

i	W	$W-1$	$0-1$
-----	-----	-------	-------	-------

$\hat{\alpha}_i$	1	1	1111111	0
------------------	---	---	---------	---

θ_i	0	0	0000001	
------------	---	---	---------	--

γ_i	0	0	00000001	0
------------	---	---	----------	---

α_i	0	0000000\overline{1}		\Rightarrow 相一致
------------	---	---------------------	--	-------------------

$C_1 = 1.\overline{11}\underline{00}010$ 类似操作得到 0.000\overline{1}0010

$\hat{\alpha}_i$	1	1111100	100	0
------------------	---	---------	-----	---

θ_i	0	0000101	10
------------	---	---------	----

γ_i	0	0000100	100
------------	---	---------	-----

α_i	0000\overline{1}	000	10
------------	------------------	-----	----

类似 $C_2 = 0.01100\underline{1}0 = 0.01\underline{1}010\overline{1}0 = 0.10\overline{1}010\overline{1}0$

$C_3 = 0.01001100\underline{1}1 = 0.0100\underline{1}10100\overline{1} = 0.01010\overline{1}0100\overline{1}$

Chapter 14

4. 设计关键路径为3个全加器的基4最小冗余加法器，假设基4数位 x_i 用3条线编码为 $-2x_i^{-2} + x_i^+ + x_i^{++}$ 。所有输入和输出操作数均用最小冗余表示法表示。只使用4个全加器和一些延迟元件，给出相应的最高数位优先和最低数位优先架构。(提示：对于第3级加法器，将 $-a$ 替换为 $-2a+a$ 。)

解：需计算 $S_{<4,2>} = X_{<4,2>} + Y_{<4,2>} \quad \text{将 } Y_{<4,2>} \text{ 拆为 } Y_{<4,2>} = Y^+ - Y^-$

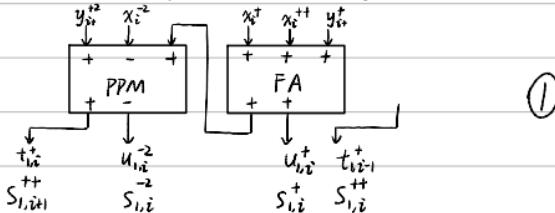
这样 先后计算 $S'_{<4,2>} = X_{<4,2>} + Y^+$ ， $S_{<4,2>} = S'_{<4,2>} - Y^-$

$$\text{而 } y_i = -2y_i^{-2} + y_i^+ + y_i^{++} = -2y_i^+ - y_i^+ + 2y_i^+ + y_i^{++} = (2y_i^+ + y_i^{++}) - (2y_i^{-2} + y_i^+)$$

$$\text{可记为 } y_i = y_{i+} - y_{i-} = (2y_{i+}^{-2} + y_{i+}^+) - (2y_{i-}^{-2} + y_{i-}^+) \quad \text{则 } \begin{cases} y_{i+}^{-2} - y_{i+}^+ \\ y_{i-}^{-2} - y_{i-}^+ \\ y_{i+}^+ = y_{i-}^{++} \\ y_{i-}^+ = y_{i+}^+ \end{cases}$$

若计算 $x_i + y_{i+} = 4t_{i+} + u_{i+}$ 结合 表 14-5 与 式 14-13

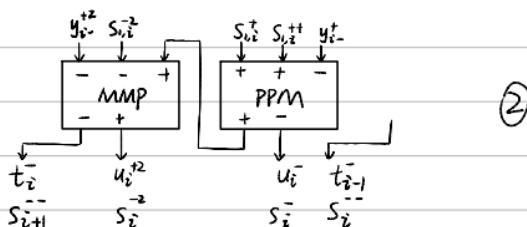
$$(-2x_i^{-2} + 2y_i^{-2}) + (x_i^+ + x_i^{++} + y_i^+) = 4t_{i+}^+ - 2u_{i+}^{-2} + u_{i+}^+$$



$$\text{之后计算 } S_{i+1} - y_{i-} = 4t_{i-} + u_{i-}$$

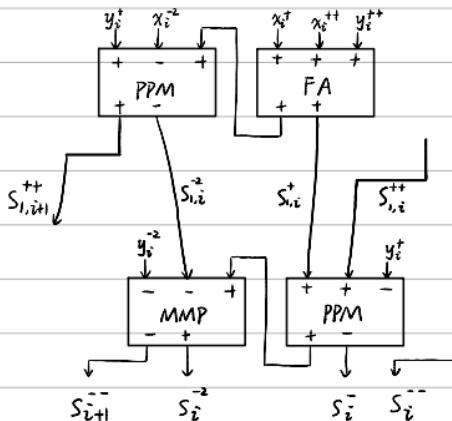
数位	基4数位集合	二进制码
S_{i+1}	{2, 1, 0, 1, 2}	$-2S_{i+1}^{-2} + S_{i+1}^1 + S_{i+1}^{++}$
y_{i-}	{0, 1, 2, 3}	$-2y_{i-}^{-2} + y_{i-}^+$
$P_i = S_{i+1} - y_{i-}$	{5, 6, 3, 2, 7, 0, 1, 2}	...
t_{i-}	{7, 0}	$-t_{i-}^+$
u_{i-}	{7, 0, 1, 2}	$2u_{i-}^{-2} - u_{i-}^+$
$S_i = u_i + t_{i-}$	{2, 1, 0, 1, 2}	$S_i^{-2} - S_i^1 - S_i^{++}$

$$(-2S_{i+1}^{-2} - 2y_{i-}^{-2}) + (x_i^+ + x_i^{++} - y_i^+) = -4t_{i-}^+ + 2u_{i-}^{-2} - u_{i-}^+$$



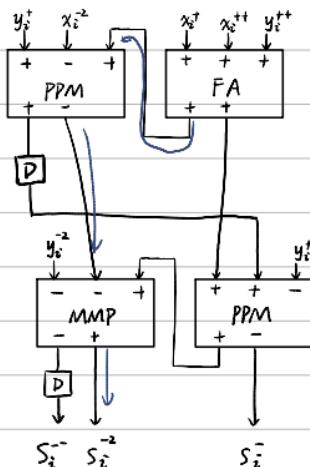
将 ① ② 级联 结合 $\begin{cases} y_{i+1}^{++} = y_i^+ \\ y_{i+1}^+ = y_i^{++} \\ y_{i+1}^{--} = y_i^+ \end{cases}$

有

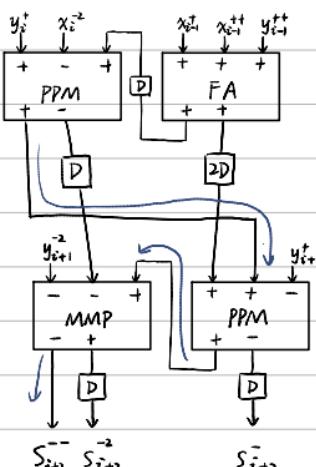


不难得出最高数位优先与最低数位优先架构如下

最低数位优先



最高数位优先 (已插入流水线)



→ 标注出了关键路径，均不超过 3 个 FA