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QNo1:- Knapsack problem with weight of knapsack of even, odd.

You can keep two tables with  $n$  rows and  $C$  columns:

$DP_{\text{even}}$  that saves the best knapsack solution with an even number of items.

$DP_{\text{odd}}$  that saves the best knapsack solution with odd number of items.

To fill  $DP_{\text{even}}$  you look at the previous best solution of  $DP_{\text{odd}}$  plus a new item, or of  $DP_{\text{even}}$  if you don't take the item.

$$DP_{\text{Even}}[i, j] = \max(DP_{\text{Odd}}[i-1, j-c[i]] + v[i], DP_{\text{Even}}[i-1, j])$$

The same idea goes for  $DP_{\text{Odd}}$

$$DP_{\text{Odd}}[i, j] = \max(DP_{\text{Even}}[i-1, j-c[i]] + v[i], DP_{\text{Odd}}[i-1, j])$$

Your solution will be in  $DP_{\text{Even}}[n, c]$

**QNo2-**

## Huffman Codes

i) :- Data can be encoded efficiently using Huffman codes.

ii) :- it is a widely used and beneficial technique for compressing data.

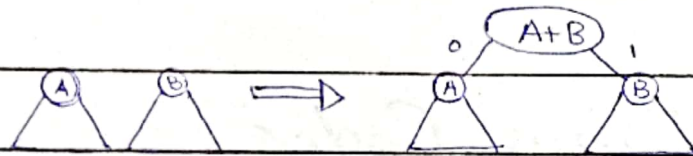
iii) :- Huffman's greedy algorithm uses a table of frequencies of occurrences of each character to build up an optimal way of representing each character as

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a binary string.

## Greedy Algorithm for Constructing a Huffman Code:-

Huffman invented a greedy algorithm that creates an optimal prefix called a Huffman Code.



The algorithm builds the tree  $T$  analogous to the optimal code in a bottom-up manner.

## Algorithm of Huffman Code

### Huffman( $C$ )

- 1):-  $n = |C|$
- 2):-  $Q \leftarrow C$
- 3):- for  $i = 1$  to  $n - 1$
- 4):- do
- 5):-  $Z = \text{allocate-Node}()$
- 6):-  $X = \text{left}[Z] = \text{Extract-Min}(Q)$
- 7):-  $Y = \text{right}[Z] = \text{Extract-Min}(Q)$



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$$8):- P[Z] = P[X] + P[Y]$$

$$9):- \text{Insert}(\mathcal{Q}, z)$$

$$10):- \text{return Extract-Min}(\mathcal{Q})$$

Example:-

$$a: 50, b: 25, c: 15, d: 40, e: 75$$

Solution

$$\text{Given that } C = \{a, b, c, d, e\}$$

$$P(C) = \{50, 25, 15, 40, 75\}$$

$$n = 5$$

$$\mathcal{Q} \leftarrow C$$

$$\text{For } i \leftarrow 1 \text{ to } 4$$

$$i = 1 \quad Z \leftarrow \text{Allocate node}$$

$$X \leftarrow \text{Extract-Min}(\mathcal{Q})$$

$$Y \leftarrow \text{Extract-Min}(\mathcal{Q})$$

$$\text{Left}[Z] \leftarrow X$$

$$\text{Right}[Z] \leftarrow Y$$

$$P(Z) \leftarrow P(X) + P(Y) = 15 + 25$$

$$P(Z) = 40$$

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## Greedy algorithm of Huffman coding:-

~~Greedy~~ A

1):-  $n \leftarrow \text{length}[S]$

2):-  $A \leftarrow \{1\}$

3):-  $j \leftarrow 1$

4):- for  $i \leftarrow 2$  to  $n$

5):- do if  $s_i \geq p_i$

6):- then  $A \leftarrow A \cup \{i\}$

7):-  $j \leftarrow i$

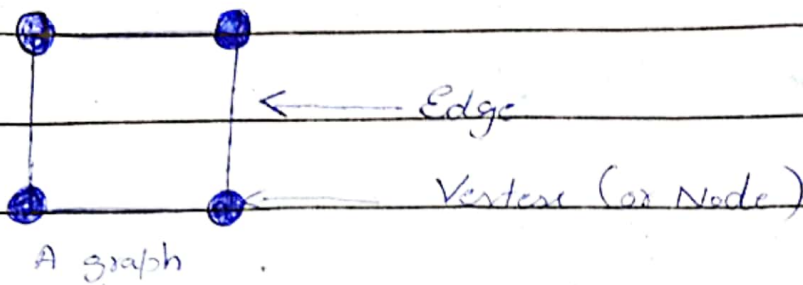
8):- return A

"Yes it is related to activity selection process."

The activity selection problem is a mathematical optimization problem."

Q No 3:-

## Graph



In simplest terms, a graph is a combination of vertices (or nodes) and edges.

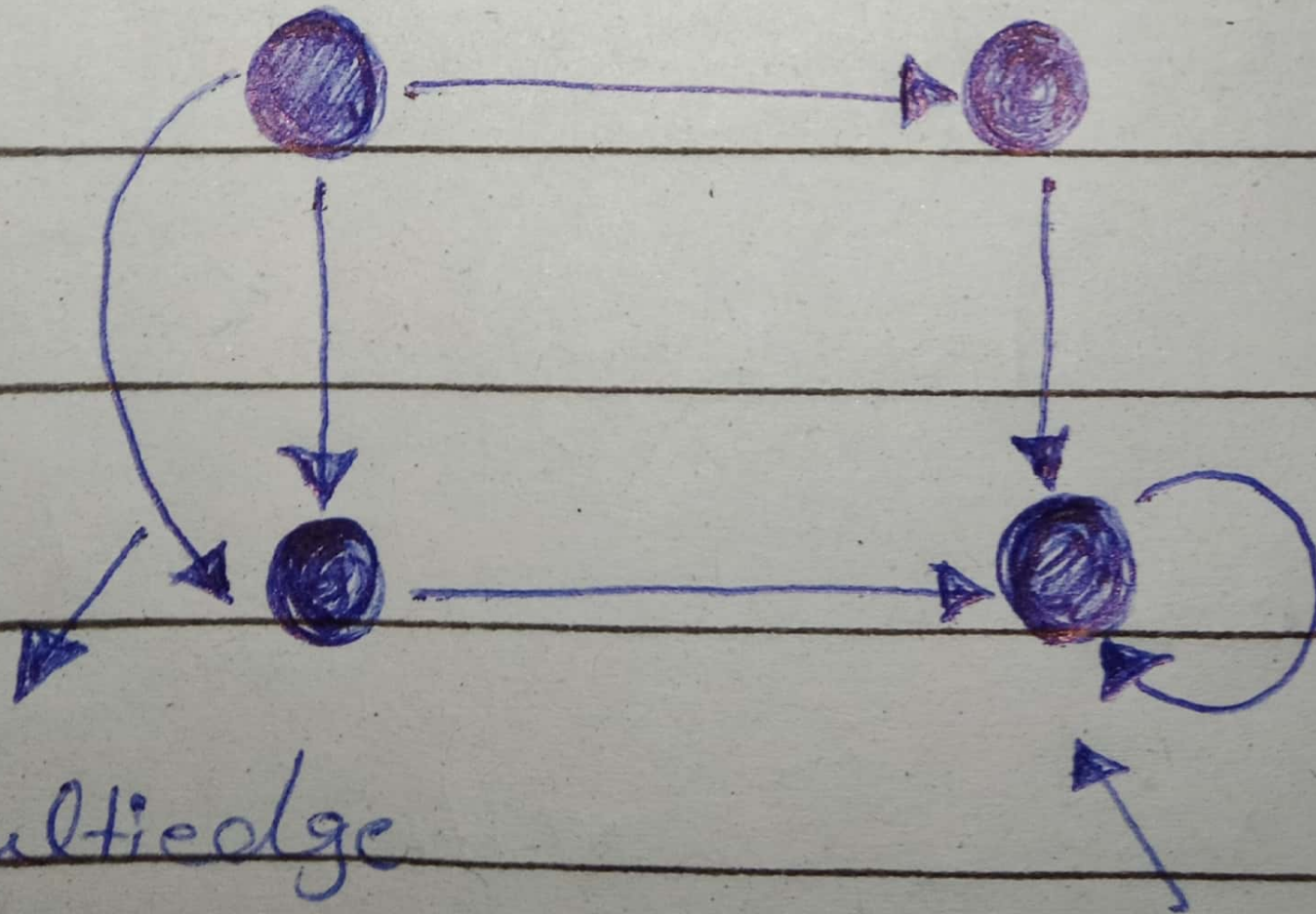
We have 4 nodes and 4 edges and it is a graph.

Graph is a very important data structure to store data which are connected to each other.

## Digraph with multiedge.

The graph which has self-loops or an edge  $(i, j)$  occurs more than once (also called multiedge and graph is called multigraph) is a non-simple graph.





Multiedge

Self loop