

缘桥

菜鸟-数据&规划部-数据技术专家

FLINK FORWARD # ASIA

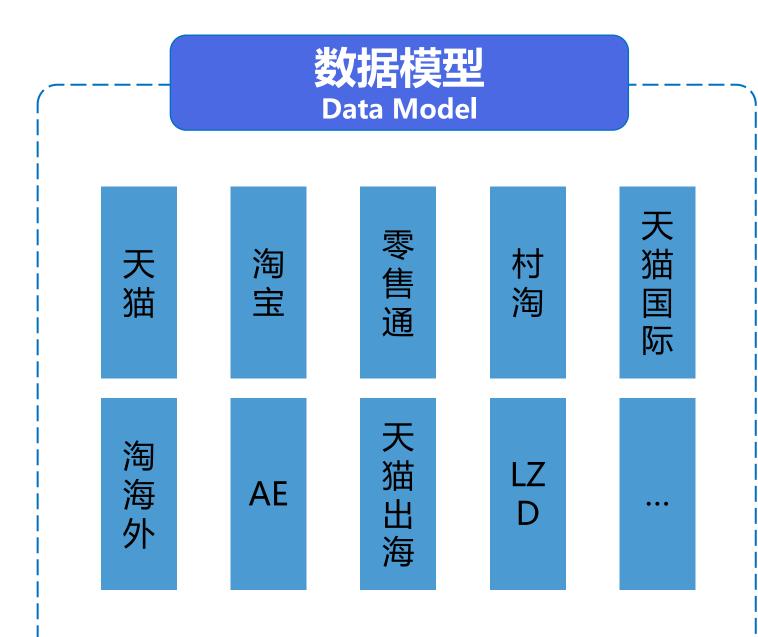
实时即未来 # Real-time Is The Future



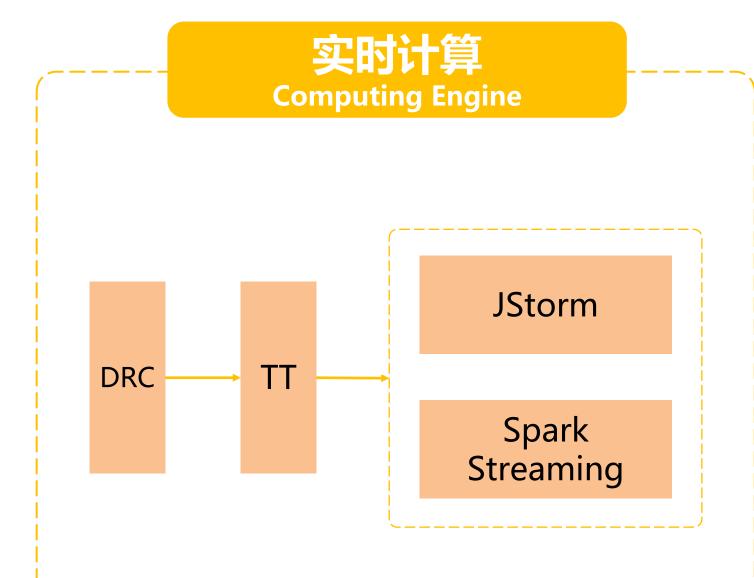
以前的实时数据技术架构



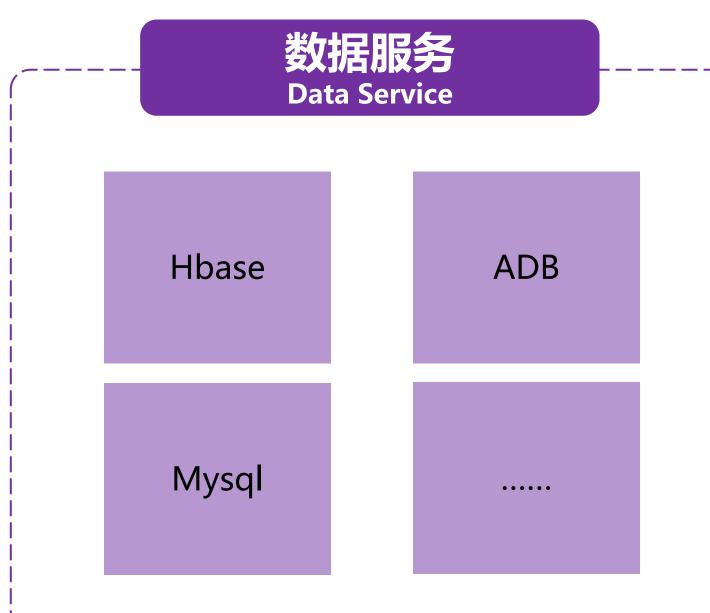
Real-time data warehouse and technology architecture for 2016



- 业务线内部模型层次混乱,数据使用 成本特别高
- The hierarchy of data model is disordered, and the cost of using data is very high.
- 需求驱动的烟囱式开发,完全没有复用的可能性,计算成本居高不下
- The chimney development model leads to excessive cost.



- 有些复杂的功能,实现起来不够简单,开发成本较高
- Some requirements are complex and costeffective.
- 很难同时兼顾功能、性能、稳定性以及快速的故障恢复能力
- It is difficult to simultaneously balance function, performance, stability.

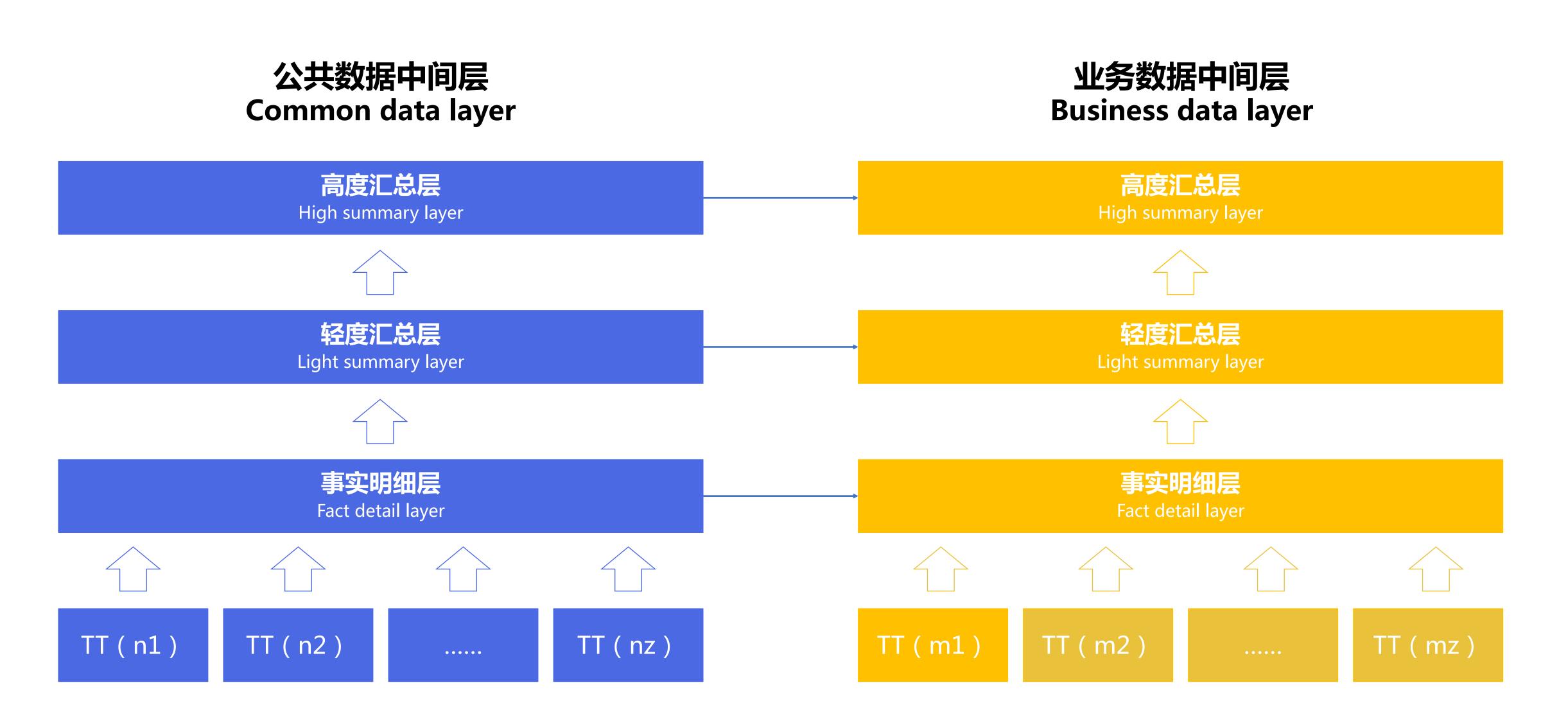


- 使用成本比较高,尤其是针对一些 NoSQL的数据库
- High cost of use, especially for some NoSQLbased database.
- 数据使用不可控,如热点阻断、权限 控制、全链路监控等
- Data security is not controllable. For example, the lack of full-link monitoring.

数据模型的升级 - 模型分层,充分复用公共中间层模型



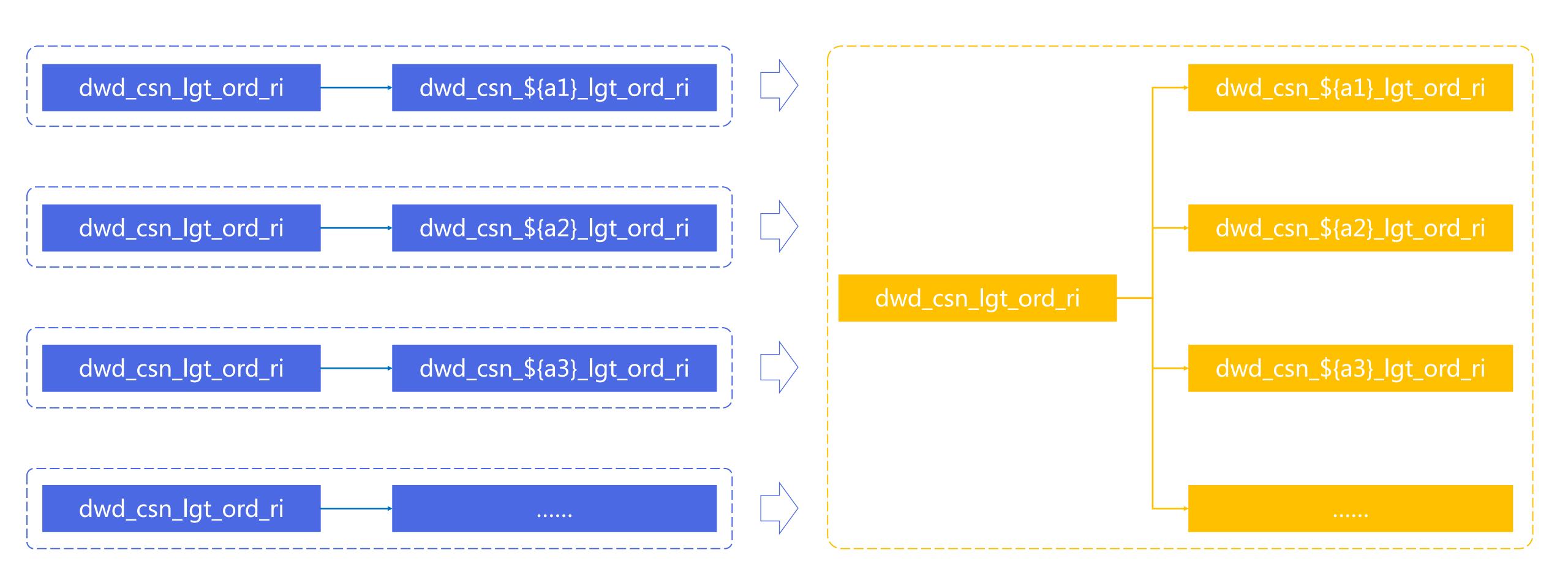
Upgrade of data model - Layering data model to improve the reusability of common data layer model



数据模型的升级 - 预置分流,充分复用公共预置分流模型



Upgrade of data model - Preset dataflow is separated to improve the reusability of common data layer model

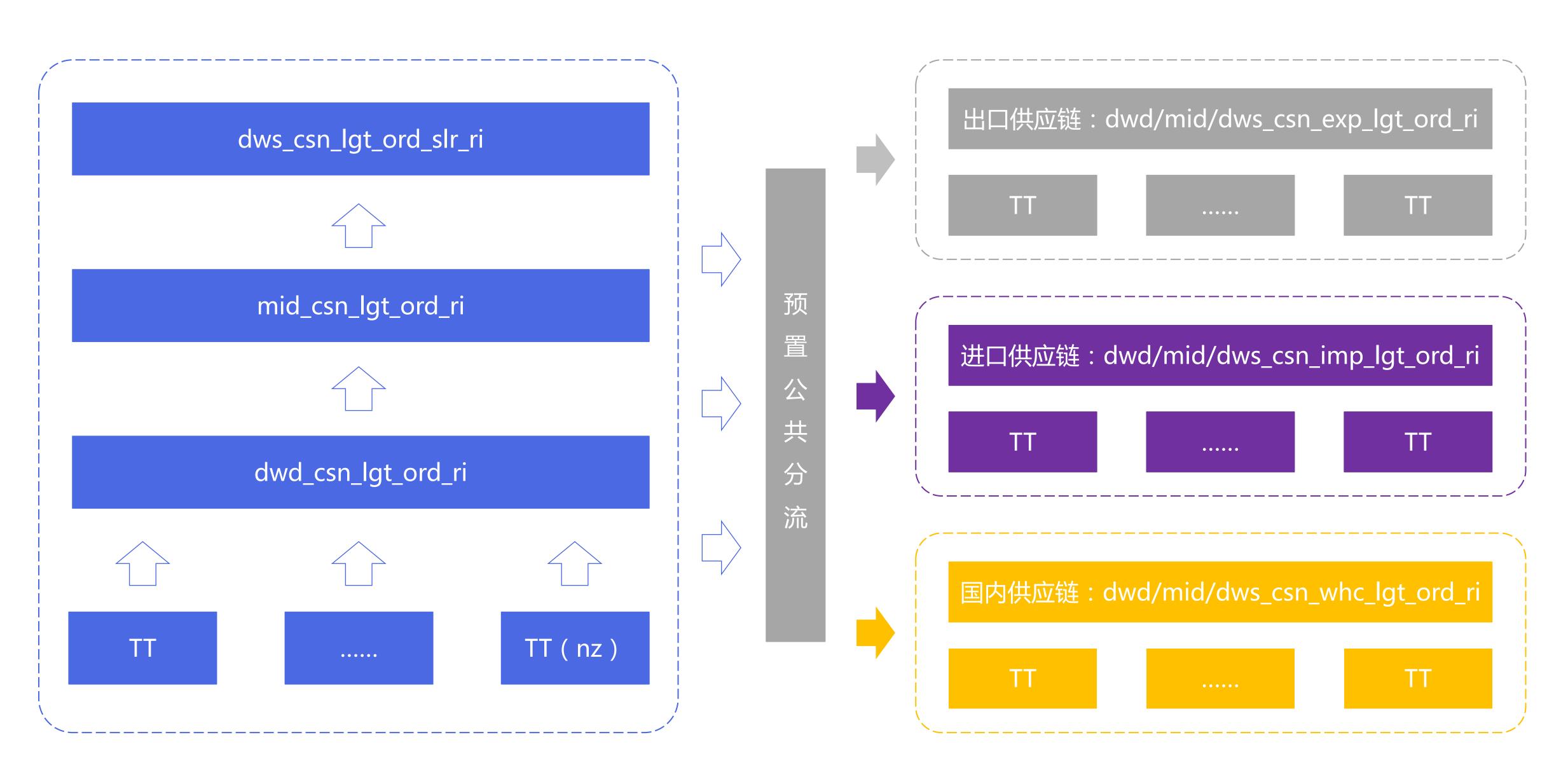


将原来下游做的分流作业,全部转移到上游一个公共分流作业来完成,可以大大节省计算资源。 Transfering from many sub-jobs to a common job can save many of computing resources.

数据模型的升级-案例:菜鸟供应链实时数据模型



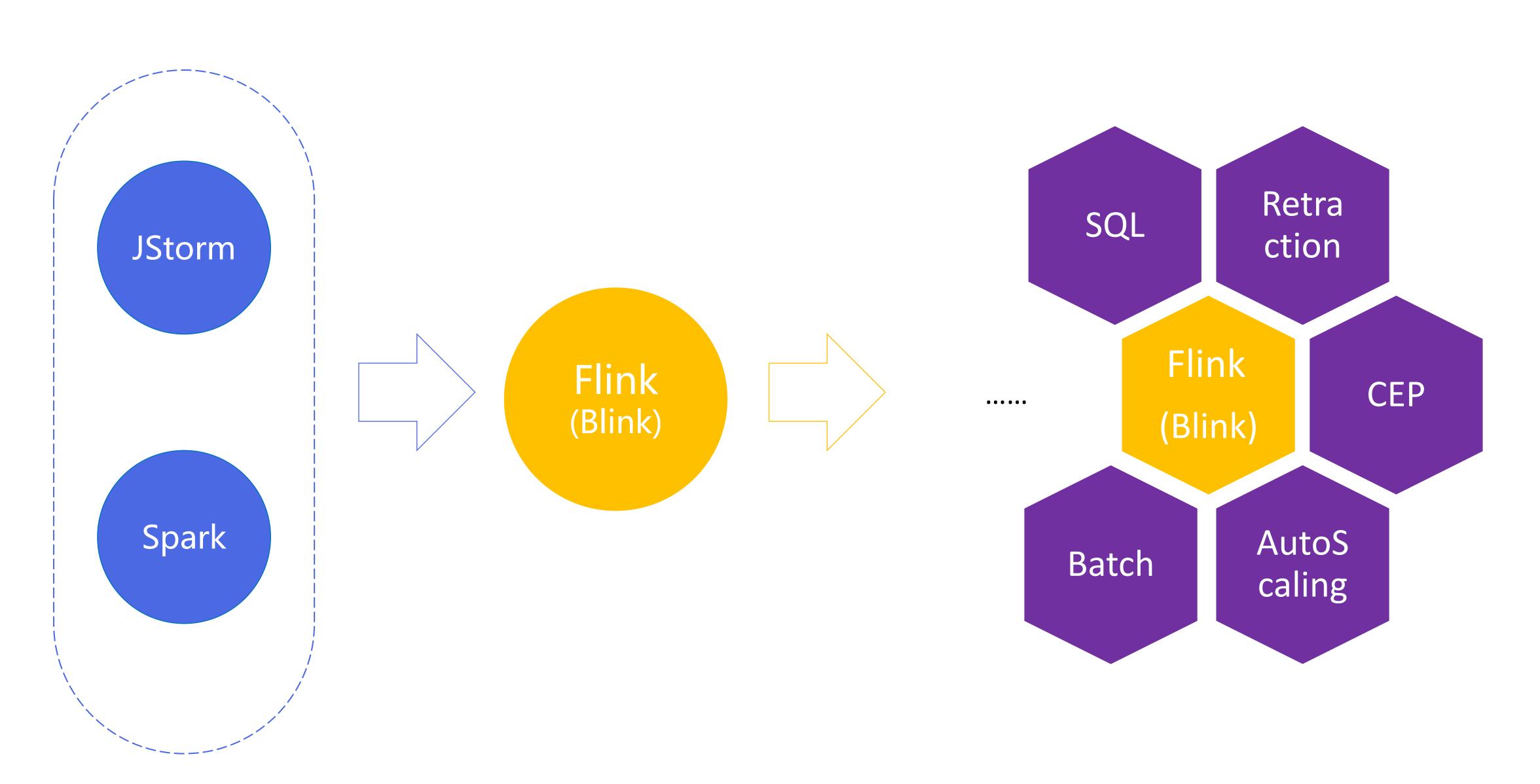
Upgrade of data model – Cainiao SCM data model



计算引擎的升级 – 基于Flink的实时计算引擎



Upgrade of computing engine - Real-time computing engine based on Flink



计算引擎的升级 – 案例1:神奇的Retraction

Upgrade of computing engine - Amazing retraction



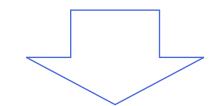
物流订单号 lg_order_code	创建时间 gmt_create	是否取消 is_cancel	计划配送公司 plan_tms
LP1	2019-10-01 00:01:00	Υ	tmsA
LP2	2019-10-01 00:05:00	Υ	tmsA
LP1	2019-10-01 00:01:00	Υ	tmsA
LP1	2019-10-01 00:01:00	Υ (tmsB
LP2	2019-10-01 00:05:00	Υ	tmsA
LP3	2019-10-01 00:18:00	Υ	tmsA
LP2	2019-10-01 00:05:00	Υ	tmsA
LP1	2019-10-01 00:01:00	Υ (tmsC
LP3	2019-10-01 00:18:00	Υ	tmsA
LP2	2019-10-01 00:05:00	Υ	tmsA
LP3	2019-10-01 00:18:00	Υ	tmsA
LP3	2019-10-01 00:18:00	N	tmsA

计划发货仓 plan_store	创建物流订单量 create_order_cnt
tmsA	3 (LP1+LP2+LP3)
tmsB	1 (LP1)
tmsC	1 (LP1)



计划发货仓 plan_store	创建物流订单量 create_order_cnt
tmsA	1 (LP2)
tmsC	1 (LP1)





利用Flink内置的Retraction机制可以轻松实现流式消息的回撤统计

It's very easy to use the flink's build-in retraction mechanism.

如何统计每个配送公司计划履行多少有效单量?

For each tms, how to count the number of plan orders?

计算引擎的升级 – 案例1:神奇的Retraction

FLINK FORWARD

Upgrade of computing engine - Amazing retraction

```
--临时视图
create view dws_csn_whc_lgt_ord_tms_ri_v1 as
select lg_order_code
        , last_value(gmt_create) as gmt_create
        , last_value(plan_tms ) as plan_tms
        , last_value(is_cancel) as is_cancel
        dwd_csn_whc_lgt_ord_ri_v1
from
group by Ig_order_code
--最终结果
insert into dws_csn_whc_lgt_ord_tms_ri
select substr(gmt_create, 1, 10)
        ,plan_tms
        ,count(Ig_order_code) as
plan_lgtord_cnt
        dws_csn_whc_lgt_ord_tms_ri_v1
from
        coalesce(is_cancel, 'N') = 'N'
where
group by substr(gmt_create, 1, 10)
        ,plan_tms
```

利用Flink SQL内置函数last_value,获取聚合key的最后一条非空的数值

一旦gmt_create、plan_tms、is_cancel中的任何一个字段发生变化,都会触发Flink的retraction机制

计算引擎的升级 - 案例2:实时超时统计的福音

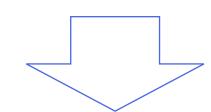
Upgrade of computing engine - Real-time timeout statistics



日志时间 gmt_modified	物流订单号 lg_order_code	出库时间 storeout_time	揽收时间 tmsgot_time
2019-10-01 00:01:00	LP1		
2019-10-01 00:05:00	LP1	2019-10-01 00:05:00	
2019-10-01 00:10:00	LP2		

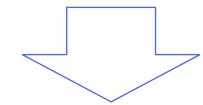
业务需求:仓出库超6小时配未揽收的单量?

How to count the number of orders which is consigned but not collected more than 6 hours.



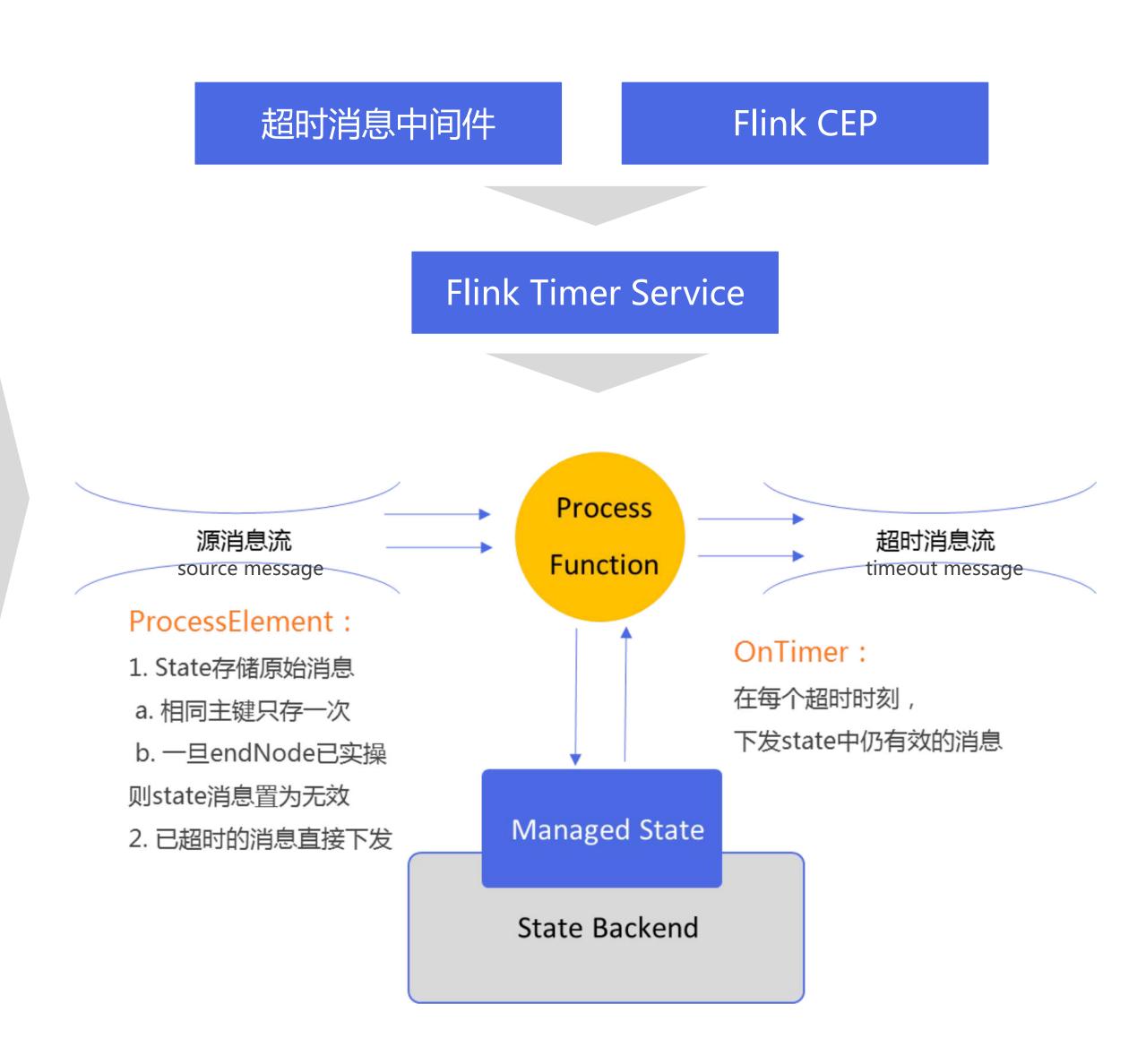
难点:如果仓出库后配未揽收,意味着没有新的消息流入,没有消息,如何进行超时消息的计算呢?

If the order is consigned but not collected, it means that there is no new message inflow. In this case, how to count the num of timeout message?



方案:没有消息,通过Flink制造消息!

Making timeout messages through flink!



计算引擎的升级 - 案例2:实时超时统计的福音

Upgrade of computing engine - Real-time timeout statistics



```
--创建执行环境
StreamExecutionEnvironment env = StreamExecutionEnvironment.getExecutionEnvironment();
DataStream<Row> pds = env.addSource(tt4Source).....process(new TimeOutEmit())....
-- TimeOutEmit函数
public void processElement(Tuple2<String, TimingMsg> value, ProcessFunction<Tuple2<String,
TimingMsg>, TimingMsg>.Context context, Collector<TimingMsg> collector) throws Exception {
 for (String timingHour : timingHours) {
   long registerTime = startNodeTimeStamp + Long.valueOf(timingHour) * 3600000L;
   if (registerTime > context.timerService().currentProcessingTime()) {
     context.timerService().registerProcessingTimeTimer(registerTime);
 this.state.update(currentMsg);
public void on Timer (long timestamp, Process Function < Tuple 2 < String, Timing Msg >,
TimingMsg>.OnTimerContext ctx, Collector<TimingMsg> out) throws Exception {
 TimingMsg result = this.state.value();
  out.collect(result);
```

构造Process Function (访问keyed state 和 timers)

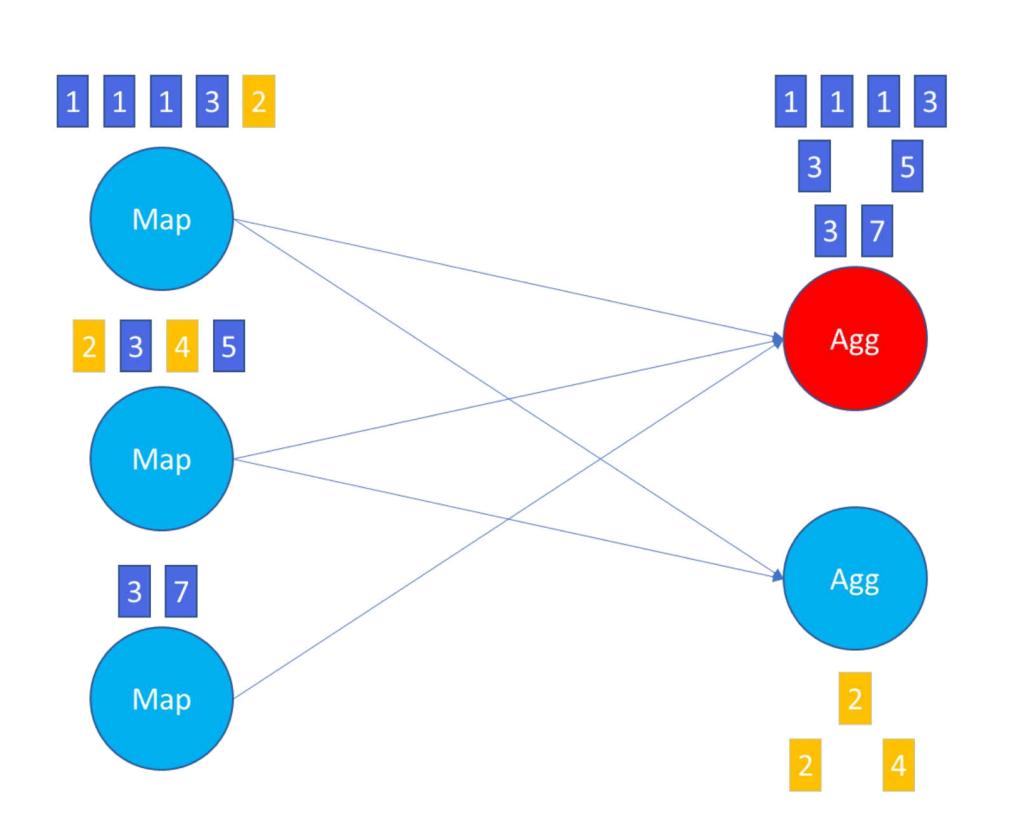
processElement,告诉state存储什么样的数据,并为每个超时事件注册一个timerService

onTimer,在超时的时刻去state读取数据,并将超时消息下发

计算引擎的升级 - 案例3:从手动优化到智能优化



Upgrade of computing engine - From manual optimization to intelligent optimization



数据热点

Hotspot

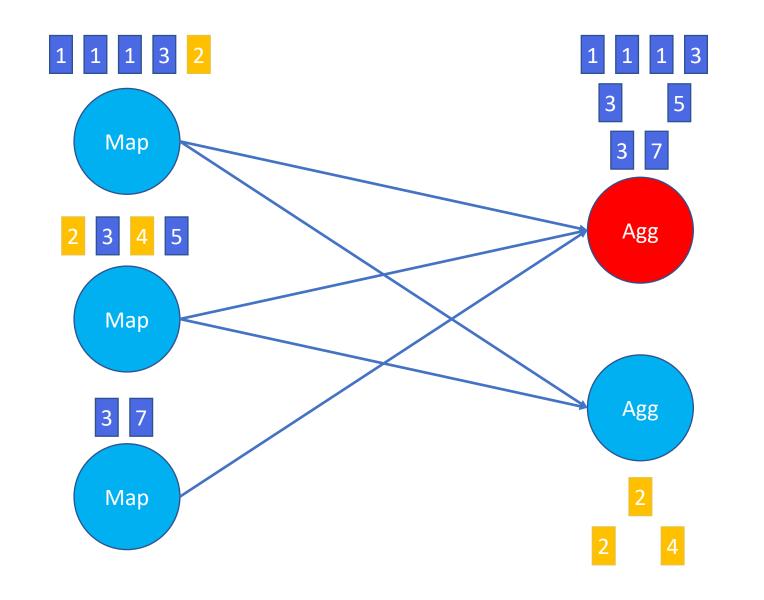
```
--hash散列
create view dws_csn_whc_lgt_ord_si_ri_v1 as
select mod(hash_code(lg_order_code),256)
        , substr(gmt_create, 1, 10)
        ,service_item_id
        ,count(distinct Ig_order_code) as
mid_crt_lgtord_cnt
from source_dwd_csn_whc_lgt_fl_ord_ri
group by mod(hash_code(lg_order_code),256)
        ,substr(gmt_create, 1, 10)
        ,service_item_id
--汇总结果
create view dws_csn_whc_lgt_ord_si_ri as
select substr(gmt_create, 1, 10) as stat_date
        ,service_item_id
        ,sum(mid_crt_lgtord_cnt) as
crt_lgtord_cnt
         dws_csn_whc_lgt_ord_si_ri_v1
from
group by substr(gmt_create, 1, 10)
        ,service_item_id
```

计算引擎的升级 - 案例3:从手动优化到智能优化



Upgrade of computing engine - From manual optimization to intelligent optimization

数据热点 Hotspot



MiniBatch

减轻对state的查询压力

Relief the pressure of state

LocalGlobal

轻松应对count热点

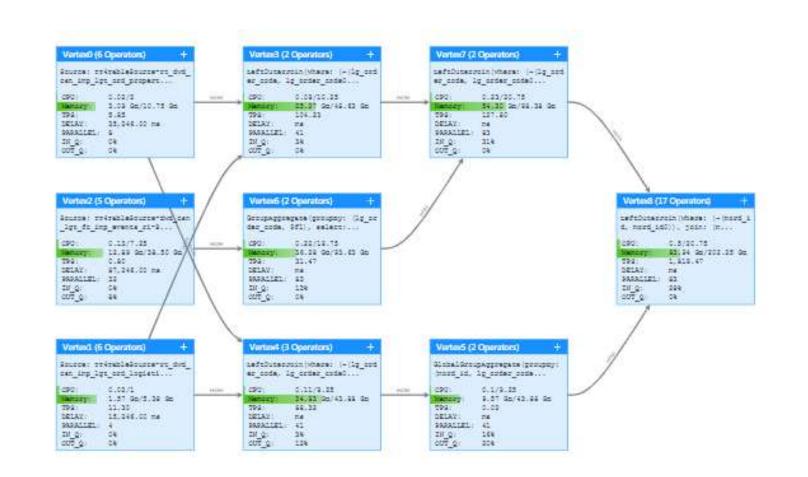
Easy to deal with hot issues from count

PartialFinal

轻松应对count_distinct热点

Easy to deal with hot issues from count distinct

资源配置 Resource config

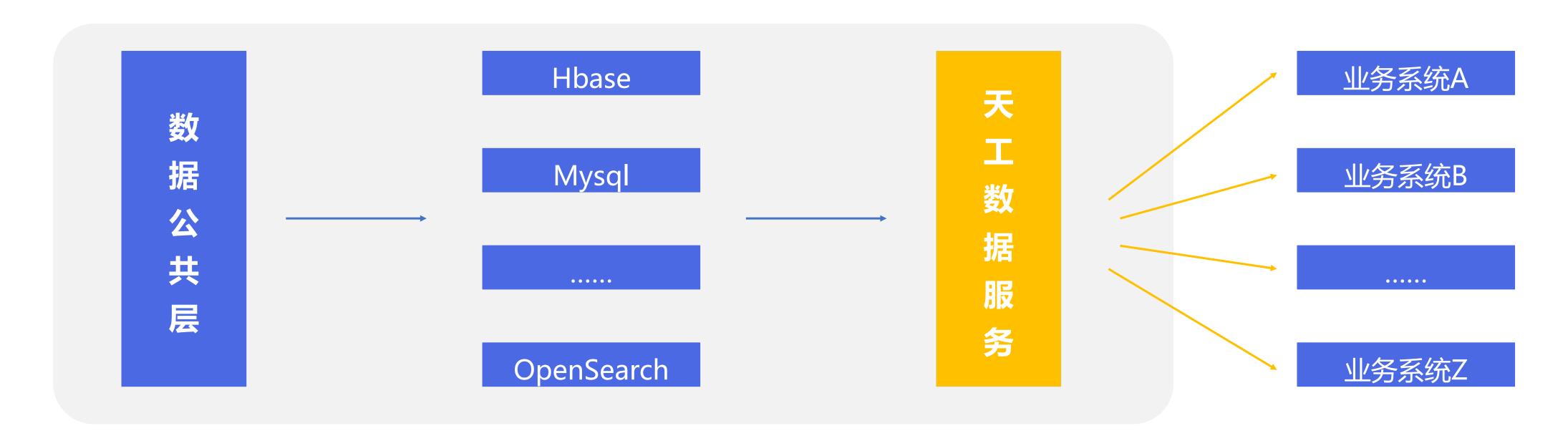


AutoScaling智能调优
AutoScaling intelligent optimization

数据服务的升级 - 统一数据服务中间件

Upgrade of data service - Unified data service middleware





统一数据库接入
Unified database access

统一的权限管理
Unified authority management

统一的数据保障
Unified data guarantee

统一全链路监控
Full link monitoring

标准

将 SQL 做为一等公民,作为数据服务的 DSL

Using SQL as the DSL of data service

提供标准化的服务接入方式 (HSF)

Provide standardized service access mode — HSF

数据服务的升级 – 案例1: NoSQL To TgSQL

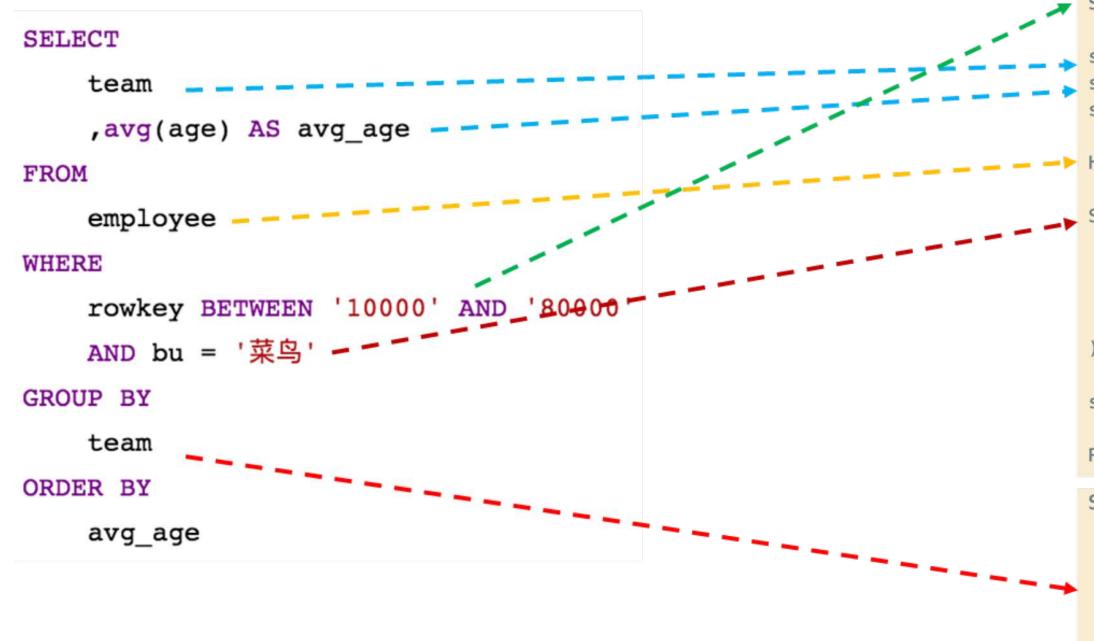


Upgrade of data service – NoSQL to TgSQL

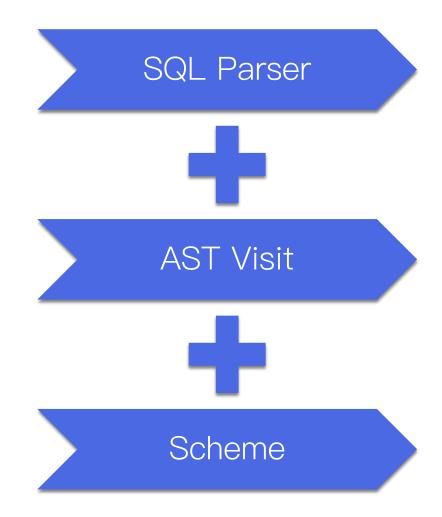
Employe	e	物理表
rowkwy	CF : base	CF: extends
666666	name:xiaoming;age:31;	team:菜鸟数据部;bu:菜鸟
•••••	•••••	•••••



Employee _.	Employee_hb 逻辑表			
rowkwy	name	age	team	bu
666666	xiaoming	32	菜鸟数据部	菜鸟
•••••				



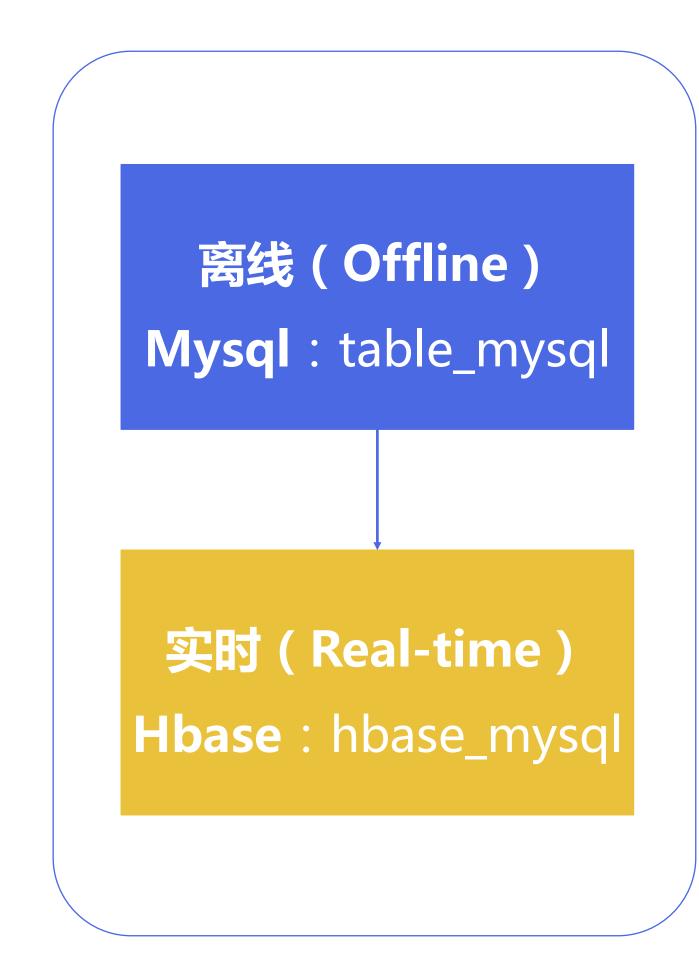


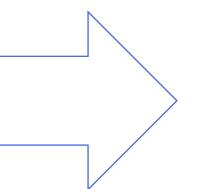


数据服务的升级 - 案例2:跨源数据查询

Upgrade of data service - Cross-source data query







```
a.store_code
select
        ,a.plan_lgtord_cnt
        ,b.real_lgtord_cnt
from
         (select
          from
                  mysql_table
                   cate_id = 101
          where
          ) a
left outer join
         (select
          from
                   hbase_table
                   cate_id = 101
          where
          ) b
         a.store_code = b.store_code
on
```



数据服务的升级 - 案例3:服务保障升级

Upgrade of data service – Upgrade of Service guarantee





主备切换 Rapid switchover 主备双活 Dual active 动态负载 Dynamic load balancing

热点服务阻断 Blocking-up hotspot 白名单限流 Whitelist current limit

• • • • •

其他技术工具的探索和创新

FLINK FORWARD

Exploration and innovation of other technical tools

实时压测

Pressure Measurement

过程监控

Process Monitor

资源管理

Resource Management

数据质量

Data Quality

菜乌实时数仓-未来发展与思考

FLINK

Upgrade of data service - Future development and thinking



联系方式

Contact information



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