# CSC236 2017 Summer Assignment 1

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Due Date: Jun 7th, 2017

### 1. Proof by Simple Induction:

Define the Predicate P(n): f(0) + f(2) + ... + f(2n) = f(2n+1)

Want to prove:  $\forall n \in \mathbb{N}, P(n)$ 

Base Case: Let n = 0, We want to prove P(0) i.e, f(0) = f(2\*0+1) = f(1)

By the Fibonacci function, we know that f(0) = f(1) = 1. Then f(0) = f(2\*0 + 1)

Let n = 1, We want to prove P(1) i.e, f(0) + f(2) = f(2\*1+1) = f(3)

f(2) = f(0) + f(1) = 1 + 1 = 2 (By fibonacci function)

f(3) = f(1) + f(2) = 1 + 2 = 3

f(0) + f(2) = 1 + 2 = 3 = f(3)

**Induction Step**: Let  $k \in \mathbb{N}$ ,  $k \ge 0$ . Assume P(k) i.e, f(0) + f(2) + ... + f(2k) = f(2k+1)

Want to prove P(k + 1) i.e, f(0) + f(2) + ... + f(2k) + f(2k+2) = f(2(k+1)+1) = f(2k+3)

$$f(0) + f(2) + ... + f(2k) + f(2k+2) = f(2k+1) + f(2k+2)$$
 (By Induction hypothesis)  
=  $f(2k+3)$  (By Definition of Fibonacci function)  
=  $f(2(k+1)+1)$ 

Hence 
$$\forall n \in \mathbb{N}, f(0) + f(2) + \dots + f(2n) = f(2n+1) \square$$

## 2. Proof by Simple Induction:

Define the Predicate P(n):  $8 \mid (2n+1)^2 - 1$ 

Want to prove:  $\forall n \in \mathbb{N}, P(n)$ 

Base Case: Let n=0, We want to prove P(0) i.e,  $8\mid 0$ 

8\*0 = 0

**Induction Step**: Let  $k \in \mathbb{N}$ ,  $k \ge 0$ . Assume P(k) i.e,  $8 \mid (2k+1)^2 - 1$ 

Want to prove P(k + 1) i.e,  $8 | (2(k+1)+1)^2 - 1$ 

$$(2(k+1)+1)^{2} - 1 = (2k+3)^{2} - 1$$

$$= 4k^{2} + 12k + 8$$

$$= (4k^{2} + 4k + 1 - 1) + 8k + 8$$

$$= [(2k+1)^{2} - 1] + 8(k+1)$$

We know that  $8 \mid (2k+1)^2 - 1$  is true by Induction hypothesis.

And  $8 \mid 8(k+1)$  is true since  $k \in \mathbb{N}$ .

And then  $8 \mid [(2k+1)^2 - 1] + 8(k+1)$  is true.

Hence  $\forall n \in \mathbb{N}, 8 \mid (2n+1)^2 - 1 \square$ 

3. We cannot represent any amount with coins of denominations 3 and 5. But we can find a number 8 that any amount greater or equal to it we can represented it with the above coins.

Define 
$$P(n) = \exists a, b \in \mathbb{N}, n = 3a + 5b$$
.

We want to prove  $\forall n \in \mathbb{N}, n \geq 8 \Rightarrow P(n)$ .

Proof by simple induction.

Base Case: Let n = 8, We want to prove P(8), i.e, 8 = 3a + 5b.

Let a = 1, b = 5, Then 8 = 3a + 5b.

**Induction Step**: Let  $k \in N, k \ge 8$ . Assume P(k), i.e,  $\exists a, b \in N, k = 3a + 5b$ .

We want to prove P(k+1), i.e,  $\exists a', b' \in N, k+1 = 3a' + 5b'$ .

We break in into two cases.

Case1: **b** > **0**: 
$$k + 1 = 3a' + 5b'$$
  
  $= 3a + 5b + 1$ (By Induction Hypothesis)  
  $= 3a + 5b + (2 * 3 - 1 * 5)$   
  $= 3(a + 2) + 5(b - 1)$   
  $= 3a' + 5b'$   
Case2: **b** = **0**:  $k + 1 = 3a' + 5b'$   
  $= 3a + 5b + (2 * 5 - 3 * 3)$ (By Induction Hypothesis)  
  $= 3(a - 3) + 5(b + 2)$   
  $= 3a' + 5b'$ 

Hence,  $\forall n \in \mathbb{N}, n \geq 8 \Rightarrow \exists a, b \in N, n = 3a + 5b.$ 

4. Proof by Well-Ordering Principle.

Let set S:
$$\{m \in \mathbb{Z} | \exists k \in \mathbb{N}, n = 2^k * m\}$$

First we want to prove  $S \neq \emptyset$ 

Let 
$$k \in \mathbb{N}$$
, Let  $k = 0$ 

Then 
$$n = 2^0 * m = m$$

Then  $n \in S(\text{Since } m \in S)$ 

Then  $S \neq \emptyset$ 

We want to prove  $S \in \mathbb{N}$ 

We have  $2^k \ge 1$ .(Since  $k \in \mathbb{N}$ , the least natural number is 0)

And  $n = 2^k * m$ 

Then n > 1

We have  $m \in \mathbb{Z}, m \neq 0$ , otherwise, m = 0, then n = 0, which is contradict with  $n \geq 1$ .

Then m > 0.(Since  $n \ge 1 \land 2^k \ge 1$ ).

Then  $S \in \mathbb{N}$ .

Then there exist a least element  $m' \in S$ .(By Well-Ordering Principle).

We want to prove m' is odd, i.e,  $\exists d \in \mathbb{N}, m' = 2d + 1$ 

Assume m' is even, i.e,  $\exists d \in \mathbb{N}, m' = 2d$ 

Then  $n = 2^k * m' = 2^{k+1} * d$ 

Then  $d \in S \land d < m'$  (Contradiction with m' is the least element in S).

Hence, given any natural number  $n \ge 1$ , there exists an odd integer m and a natural number k such that n = 2k m.

5. Define  $P(x,y): \exists k \in \mathbb{N}, (x,y) = (2^{k+1} + 1, 2^k + 1).$ 

Want to prove  $\forall (x,y) \in N, P(x,y)$ .

Proof by structural induction.

**Base Case**: Want to prove P(3, 2), i.e,  $\exists k \in \mathbb{N}, (x, y) = (2^{k+1} + 1, 2^k + 1)$ .

Let k = 0, Then  $(3, 2) = (2^1 + 1, 2^0 + 1)$ .

Induction Step:Assume  $H(\{x,y\})$ , i.e,  $\exists k \in \mathbb{N}, (x,y) = (2^{k+1} + 1, 2^k + 1)$ .

We want to prove P(3x - 2y, x), i.e.,  $\exists k' \in \mathbb{N}, (x, y) = (2^{k'+1} + 1, 2^{k'} + 1)$ .

Let 
$$k' = k + 1$$
, Then  $(3x - 2y, x) = (3(2^{k+1} + 1) - 2(2^k + 1), 2^{k+1} + 1)$ .  

$$= (2^{k+2} + 1, 2^{k+1} + 1)$$

$$= (2^{k'+1} + 1, 2^{k'} + 1)$$

Hence,  $\forall (x,y) \in \mathbb{N}, \exists k \in \mathbb{N}, (x,y) = (2^{k+1}+1,2^k+1)$  is true for all elements of M.

6.(a)Let n=2, then there are A and B are playing this game.

Then A has a unique nearest neighbour B

And B has a unique nearest neighbour A

Then A throw B, B throw A.

Then there may be no dry person.

(b) Define P(n): Suppose 2n+1 people are positioned such that each person has a unique nearest neighbour. Each person has a single water balloon that they

throw at their nearest neighbour. Then there is always at least one dry person.

Want to prove  $\forall n \in \mathbb{N}, n > 0 \Rightarrow P(n)$ .

Proof by simple induction.

Base Case: Let n = 0, want to prove P(1).

There is no unique nearest neighbour. (Since just one person are playing this game )

Then there is always at least one dry person.

Let n = 1, want to prove P(3).

Then we have A,B and C. If A throw B, then B throw A.

C will throw A or B, and no one will throw C.

Then C is the dry person.

Then there is always at least one dry person.

Induction Step: let  $n = k, k \in \mathbb{N}$ . Assume P(k), i.e, Suppose Suppose 2k+1 people are positioned such that each person has a unique nearest neighbour. Each person has a single water balloon that they throw at their nearest neighbour. Then there is always at least one dry person. Want to prove P(k+1), i.e, Suppose Suppose 2k+3 people are positioned such that each person has a unique nearest neighbour. Each person has a single water balloon that they throw at their nearest neighbour. Then there is always at least one dry person. We want to separate 2k+3 people into two group. Picking arbitrary two people A and B to formed Group 1. And the 2k+1 people formed Group2.

#### Case1: At least one person form Group2 throw A or B

We have one of the 2k+1 people is unique nearest A or B.

Then A or B is not dry

Then there are 2k+1 people have 2k water balloon.

Then there is always at least one dry person.

#### Case2: Every people from Group2 do not throw A and B

Then A will throw B, and B will throw A.

Then, we just consider Group 2.

Then Group must have one dry person (By Induction Hypothesis).

Then 2k+3 people is always have at least one dry person.

Put Case1 and Case2 together: 2k + 3 people always have one dry person. Hence,  $\forall n \in \mathbb{N}, P(n)$ .

#### Q7. Proof by Complete Induction

**Define Predicate P(n):** A convex polygon with consecutive vertices  $v_1, v_2, ..., v_n$  is triangulated into n-2 triangles, the n-2 triangles can be numbered 1, 2, ..., n-2 so that  $v_i$  is a vertex of triangle i for i=1,2,...,n-2.

**Base Case:** Let n = 3, Want to prove P(3), i.e. A convex polygon with consecutive vertices v1, v2, v3 is triangulated into 1 triangle, this triangle can be numbered as 1 so that  $v_i$  is a vertex of triangle i for i=1.

The convex polygon with three vertices v1, v2, v3 must be a triangle and the triangle itself can be numbered as 1 which v1 is a vertex of this triangle. Therefore, P(3) is true.

#### **Induction Step:** Let $n \in \mathbb{N}$ , and $n \geq 3$ .

Assume P(k), i.e. A convex polygon with consecutive vertices  $v_1, v_2, ..., v_k$  is triangulated into k-2 triangles, the k-2 triangles can be numbered 1, 2, ..., k-2 so that  $v_i$  is a vertex of triangle i for i = 1, 2, ..., k - 2, for  $\forall k \in \mathbb{N}$ , and  $3 \leq k < n$ .

Want to prove P(n) i.e. A convex polygon with consecutive vertices  $v_1, v_2, ..., v_n$  is triangulated into n-2 triangles, the n-2 triangles can be numbered 1, 2, ..., n-2 so that  $v_i$  is a vertex of triangle i for i = 1, 2, ..., n - 2.

There are two cases we need to consider:

 $Case1: v_n$  is connected with one of the vertices  $v_m$ 

And the diagonal between  $v_n$  and  $v_m$  will separate the convex polygon into two smaller polygons, we named them A1 and A2.

Since  $v_m$  can not be the neighbour with  $v_1$  or  $v_n$ , we know that 1 < m < n - 1.

We consider A1 first, there are m+1 vertices such as  $v_1, v_2...v_m, v_n$  in A1.

Since we need consecutive vertices, and we rename  $v_n$  to  $v_m + 1$ .

From the condition above: 1 < m < n - 1, we can get  $3 \le m + 1 < n$  by some simple manipulations. And this condition satisfies what we assumed in induction hypothesis. Therefore, A1 satisfy the predicate.

And then, consider A2, there are n-m+1 vertices such as  $v_m$ ,  $v_m + 1...v_n - i$ ,  $v_n$  in A2. Since we need consecutive vertices, and we rename  $v_i$  to  $v_i - m + 1$ .

Since we know that 1 < m < n - 1, and 1-n < -m < -1.

From 1-n < -m < -1, we can get  $3 \le n$  - m + 1 < n by some simple manipulations. And this condition satisfies what we assumed in induction hypothesis.

Therefore, A2 satisfy the predicate.

Combine A1 and A2, and rename the vertices in P2 as  $v_i - 1 + m$  in order to fit the condition of consecutive vertices, and then P(n) is true for the whole convex polygon.

Case2:  $v_n - 1$  is connected with one of the vertices  $v_m$ 

And the diagonal between  $v_n - 1$  and  $v_m$  will separate the convex polygon into two smaller polygons, we named them A1 and A2.

Since  $v_m$  can not be the neighbour with  $v_1$  or  $v_n - 1$ , we know that  $1 \leq m < n - 2$ .

We consider A1 first, there are m+2 vertices such as  $v_1, v_2...v_m, v_n - 1, v_n$  in A1.

Since we need consecutive vertices, and we rename  $v_n - 1$  to  $v_m + 1$  and  $v_n$  to  $v_m + 2$ .

From the condition above:  $1 \leq m < n - 2$ , we can get  $3 \leq m + 2 < n$  by some simple manipulations. And this condition satisfies what we assumed in induction hypothesis.

Therefore, A1 satisfy the predicate.

And then, consider A2, there are n-m vertices such as  $v_m$ ,  $v_m + 1...v_n - i$  in A2.

Since we need consecutive vertices, and we rename  $v_i$  to  $v_i - m + 1$ .

Since we know that  $1 \leq m < n - 2$ , and  $2-n < -m \leq -1$ .

From 2-n < -m  $\leqslant$  -1, we can get 2 < n - m  $\leqslant$  n - 1 by some simple manipulations. And this condition satisfies what we assumed in induction hypothesis.

Therefore, A2 satisfy the predicate.

Combine A1 and A2, and rename the vertices in P2 as  $v_i - 1 + m$  in order to fit the condition of consecutive vertices, and then P(n) is true for the whole convex polygon.

Finally, we have proved P(n) for both two cases.  $\square$