

KUMAR ET AL. (2020)

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OUTLINE

- Background information and motivation
- Proposed solution and features
- Evaluation
- Questions



BACKGROUND INFO



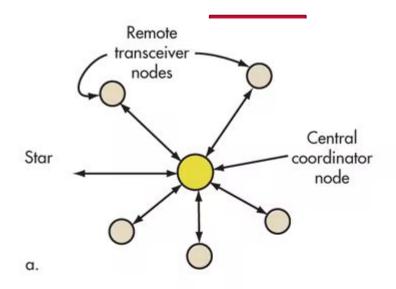
WIRELESS SENSOR NETWORKS

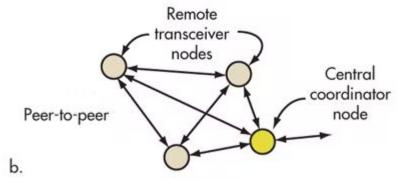


LOW-POWER AND LOSSY NETWORKS (LLNs)



IEEE 802.15.4 Network







SOMETIMES NEED RELIABLE COMMUNICATION



E.G. VOICE COMMANDS



SEVERAL EXISTING PROTOCOLS



FEW ARE TCP-BASED



WHY NOT TCP?



DEEMED TOO HEAVYWEIGHT



LLN TCP ISSUES

- 1. Memory usage non-deterministic
- 2. TCP headers take up half of IEEE 802.15.4 frames
- 3. Expected power usage poor



EXISTING TCP LACK STANDARD FEATURES



HENCE UDP-BASED PROTOCOLS



UNFORTUNATELY...



UDP-BASED HAS DRAWBACKS



SENSORS NOT FIRST-CLASS CITIZENS



APPLICATION SPECIFIC PROTOCOLS



REQUIRE DEDICATED BASESTATIONS



CUE TCPIp!



WHAT MAKES IT UNIQUE?



FULLY COMPATIBLE TCP



INTEROPERABLE W/ OTHER DEVICES



INCLUDES STANDARD TCP FEATURES



MUCH BETTER PERFORMANCE



HOW IS THIS ACCOMPLISHED?



BUILT USING FREEBSD TCP-STACK





BUILT INTO TWO RTOSes



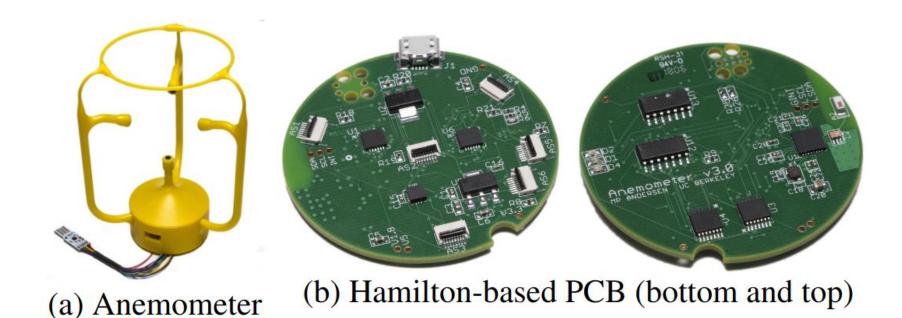


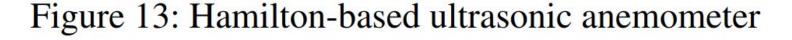


USES REPRESENTATIVE HARDWARE



USES REPRESENTATIVE HARDWARE







HOWEVER...



LLN TCP INEFFICIENT W/O MODIFICATIONS



TCP-STACK MODIFICATIONS



LLN TCP ISSUES

- 1. Memory usage non-deterministic
- 2. TCP headers take up half of IEEE 802.15.4 frames
- 3. Expected power usage poor



1. MEMORY USAGE



TCP SEND/RECV BUFFERS



STANDARD: BUFFERS SHRINK/GROW



DANGEROUS FOR EMBEDDED DEVICES



TCPIp: SIZES DEFINED AT COMPILE TIME



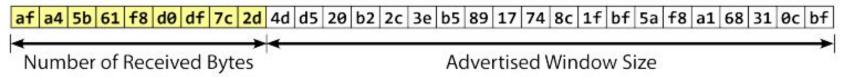
ZERO-COPY SEND BUFFER



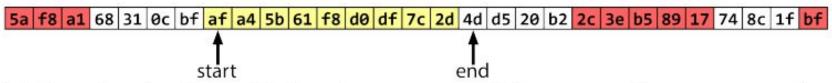
IN-PLACE REASSEMBLY FOR RECEIVE BUFFER



IN-PLACE REASSEMBLY



(a) Naïve receive buffer. Note that size of advertised window + size of buffered data = size of receive buffer.



(b) Receive buffer with in-place reassembly queue. In-sequence data (yellow) is kept in a circular buffer, and out-of-order segments (red) are written in the space past the received data.

Figure 2: Naïve and final TCP receive buffers



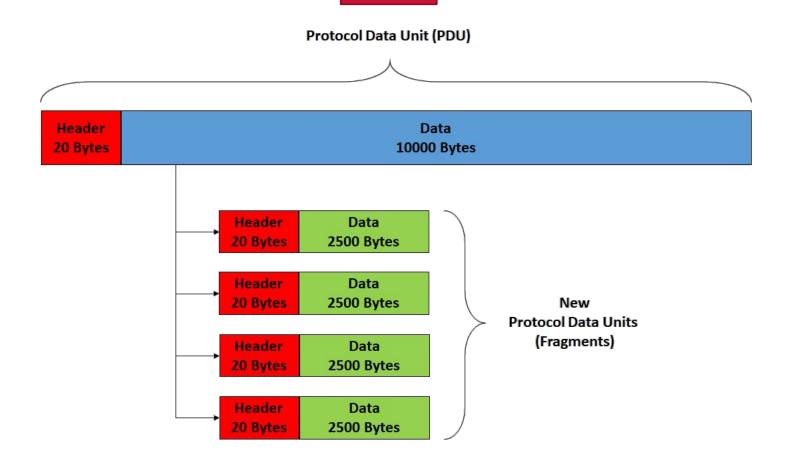
2. TCP HEADERS



STANDARD: TCP FRAGMENTS

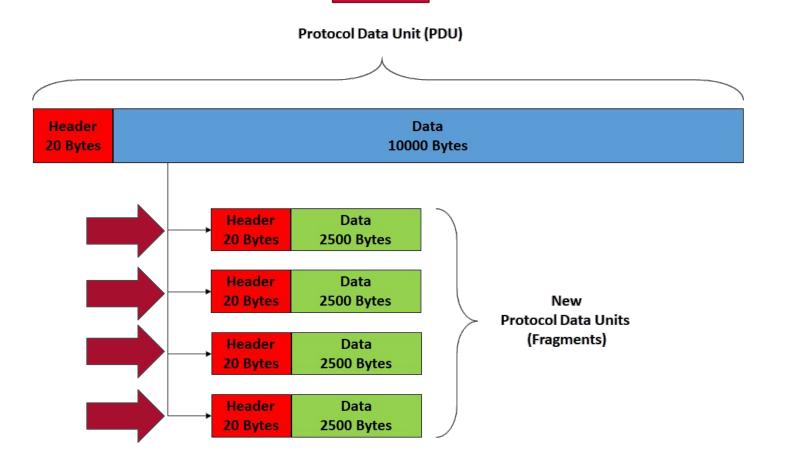


TCP FRAGMENTATION





TCP FRAGMENTATION





IN LLNs HALF EACH FRAME IS TCP HEADER



TCPIp: LOWER LAYER FRAGMENTS INSTEAD



TCP FRAGMENTATION

TCP: Data
20 Bytes 10000 Bytes

1 Packet



TCP FRAGMENTATION

TCP: Header 20 Bytes Data 10000 Bytes

1 Packet

6LoWPAN: Header 20 Bytes Data 100000 Bytes

1 2 3 4 N Frames



ONLY ONE TCP HEADER



HEADER ONLY IN FIRST FRAME, REST DATA



DECREASES RELIABILITY, BUT GENERALLY NOT MUCH



3. POWER USAGE



LLN DEVICES DUTY CYCLE



STANDARD DUTY CYCLING INEFFICIENT W/ TCP



LEAF NODES POLL FOR DATA



NO DATA => SLEEP, RETRY



E.G. ESTABLISHING CONNECTION



WAIT FULL DUTY CYCLE INCREASES LATENCY



TCPIp: MODIFY TO MAKE EFFICIENT FOR TCP



ADAPTIVE DUTY CYCLE



TCPIp: KEY POINTS



TCPIp KEY POINTS

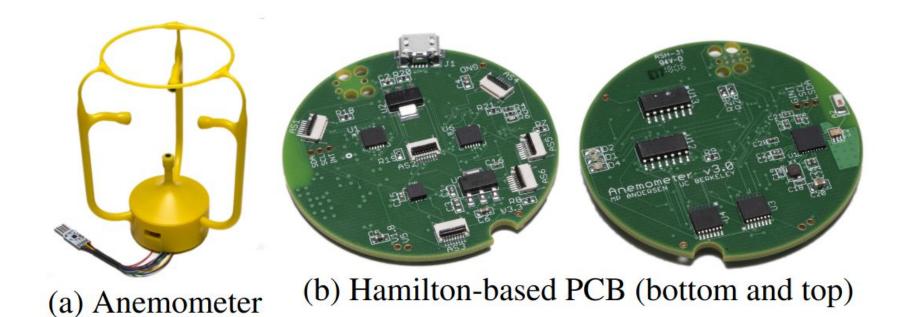
- Deterministic memory usage
- TCP headers not a problem
- More efficient power usage

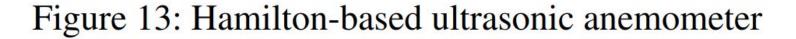


EXPERIMENTAL SETUP



HARDWARE







MULTI-HOP NETWORK

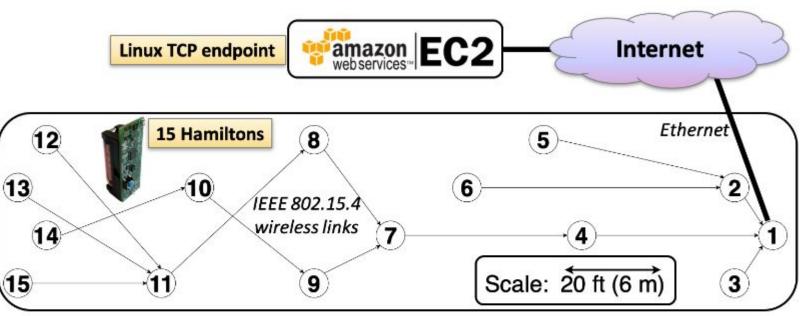


Figure 1: Snapshot of uplink routes in OpenThread topology at transmission power of -8 dBm (5 hops). Node 1 is the border router with Internet connectivity.



EVALUATION ___



"NETWORK STUDIES OVER IEEE 802.15.4 NETWORKS"



NETWORK STUDIES OVER IEEE 802.15.4 NETWORKS

	[144]	[22]	[67]	[86]	[69, 70]	This Paper (Hamilton Platform)
TCP Stack	uIP	uIP	uIP	BLIP	Arch Rock	TCPlp (RIOT OS, OpenThread)
Max. Seg Size	1 Frame	1 Frame	4 Frames	1 Frame	1024 bytes	5 Frames
Window Size	1 Seg.	1 Seg.	1 Seg.	1 Seg.	1 Seg.	1848 bytes (4 Seg.)
Goodput (One Hop)	1.5 kb/s	\approx 6.4 kb/s	$\approx 12 \text{ kb/s}$	$\approx 4.8 \text{ kb/s}$	15 kb/s	75 kb/s
Goodput (Multi-Hop)	$\approx 0.55 \text{ kb/s}$	$\approx 1.9 \text{ kb/s}$	$\approx 12 \text{ kb/s}$	$\approx 2.4 \text{ kb/s}$	9.6 kb/s	20 kb/s

Table 6: Comparison of *TCPlp* to existing TCP implementations used in network studies over IEEE 802.15.4 networks.⁶ Goodput figures obtained by reading graphs in the original paper (rather than stated numbers) are marked with the \approx symbol.



RECAP



RECAP

TCPIp

Fully-compatible TCP implementation for LLNs

Contributions:

- Performant TCP implementation for LLNs in two embedded OSes
- TCP optimizations for LLN implementations
- Evaluation of TCPIp in several common scenarios



QUESTIONS?

