Discrete Structures, CSCI-150.

Information

Monday and Wednesday 7:00 – 8:15 pm

Instructor: Alexey Nikolaev.

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Grading policy
No late homeworks accepted.
Expect to have homeworks every week.

Final grade:
HWs: 25%
Midterm exam: 35%
Final exam: 40%

Course content

Propositional Logic. Operators. Truth tables. Logical equivalence. Rules of inference. Satisfiability. Predicates and quantifiers. Proofs. Pigeonhole principle.

Counting. Sum and product rules. Permutations, n! Binomial coefficients, n choose k. Selection with replacement. Induction. Hanoi towers. Summation of series. Recurrence. Fibonacci numbers. Catalan numbers. Solving linear recurrence.

Sets. Operations, empty set, singleton set, powerset. Natural, rational, real numbers. Diagonalization. Relations and Functions. Graph of a function. Lambda-abstraction.

Counting 2. Inclusion-Exclusion. Counting and Bijection. Generating Functions.

Graphs. Bridges of Koenigsberg. Eulerian and Hamiltonian cycles. Trees, spanning trees. Travelling Salesman problem.

Number theory. Divisibility and primes. Modulo-arithmetics. GCD and Euclid's algorithm.

Literature

Primary books:

- *Rosen*
  "Discrete Mathematics and its Applications" edition 6 or 7.
  (you can find used or new 6th edition for $30–50)
- *Lehman and Leighton*
  Lecture notes "Mathematics for Computer Science" (2004).
  (free, but this is not a complete textbook)

Our first object

Something that is either

true or *false*

Propositions

Def. *A proposition is a declarative sentence that is either true or false, but not both.*

A good test for a proposition is to ask *"Is it true that . . . ?"* If that makes sense, it is a proposition.

- One plus two equals three.
- Washington, D.C., is the capital of the US.
- The Moon is a satellite of the Earth.

- Albany is the capital of Canada.
- The Sun is a planet.

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- One plus two equals three. *true* ✓
- Washington, D.C., is the capital of the US. *true* ✓
- The Moon is a satellite of the Earth. *true* ✓

- Albany is the capital of Canada. *false* ✓
- The Sun is a planet. *false* ✓ all are propositions

Propositions

Def. *A* proposition *is a declarative sentence that is either* true *or* false, *but not both.*

A good test for a proposition is to ask *"Is it true that . . . ?"* If that makes sense, it is a proposition.

- Three plus four.
- Consider these sentences.
- Does anyone have any questions?
- The largest planet in the Solar System.
- n in a prime number.

Propositions

Def. *A proposition is a declarative sentence that is either true or false, but not both.*

A good test for a proposition is to ask *"Is it true that . . . ?"* If that makes sense, it is a proposition.

- Three plus four. ✗ neither one is a proposition
- Consider these sentences. ✗
- Does anyone have any questions? ✗
- The largest planet in the Solar System. ✗
- n in a prime number. ✗

Propositions

Instead of writing sentences, we will abbreviate them by using *propositional variables*.

It is standard practice to use the lower-case letters: $p, q, r, \ldots$

Then, if

$$\begin{aligned} p &= \text{“It is raining”}, \\ q &= \text{“I have an umbrella”}, \end{aligned}$$

we can construct *compound propositions* using logical operators:

$$\begin{aligned} p \textit{ and } q &= \text{“It is raining, and I have an umbrella”}. \\ \textit{not } q &= \text{“I don't have an umbrella”}. \end{aligned}$$

Logical Operators

And (called Conjunction)

p *and* q
$p \land q$ is true when both p and q are true, otherwise false.

Or (called Disjunction)

p *or* q
$p \lor q$ is true when p or q or both are true, otherwise false.

Negation

not p
$\neg p$ is true when p is false, otherwise false.

Truth tables

And (Conjunction)

p	q	$p \wedge q$
T	T	T
F	T	F
T	F	F
F	F	F

Or (Disjunction)

p	q	$p \vee q$
T	T	T
F	T	T
T	F	T
F	F	F

Negation

p	$\neg p$
T	F
F	T

Think of the truth tables as our ultimate definition of the logical connectives (operators).

More Operators. Implication

Implication

if p *then* q

$p \to q$ is true if whenever p is true, so is q, otherwise false.

Truth table:

p	q	$p \to q$
T	T	T
F	T	T
T	F	F
F	F	T

An implication is true when the if-part *is false or the* then-part *is true.*

So, $p \to q$ is equivalent to $(\neg p) \vee q$.

"I need an umbrella, if it's raining".

"If the Earth is flat, my brother is a physicist".

More Operators. Implication

h	g	$h \to g$
T	T	T
F	T	T
T	F	F
F	F	T

"If he is hungry, he is grumpy".

$$h \to g = T$$

If an implication is true, we can make conclusions about g, if we know h.

If we know that he is indeed hungry,

$$h = T,$$

then

$$g = T.$$

More Operators. Implication

h	g	$h \rightarrow g$
T	T	T
F	T	T
T	F	F
F	F	T

"If he is hungry, he is grumpy".

$$h \rightarrow g = T$$

If an implication is true, we can make conclusions about g, if we know h.

If we know that he is not hungry,

$$h = F,$$

then

$$g \text{ can be } T \text{ or } F.$$

More Operators. Implication

p	q	$p \to q$
T	T	T
F	T	T
T	F	F
F	F	T

Big Al told us that dogs can't look up.

Our thought was that:
$p \to q =$ "If dogs can look up, Big Al is a liar".

$p =$ "Dogs can look up"
$q =$ "Big Al is a liar"

$(\neg p) \vee q =$ "Dogs can't look up, or Big Al is a liar".

More Operators. Biconditional

Biconditional

p *if and only if* q
$p \leftrightarrow q$
is true when p and q have the same truth values, otherwise false.

p	q	$p \leftrightarrow q$
T	T	T
F	T	F
T	F	F
F	F	T

Often, "if and only if" is abbreviated to *iff*:

$$p \textit{ iff } q$$

"You can take the flight if and only if you buy a ticket."

Theorems are often formulated as implications or biconditionals.

Combined truth tables for connectives $\neg$, $\wedge$, $\vee$, $\rightarrow$, and $\leftrightarrow$

p	q	$\neg p$	$p \wedge q$	$p \vee q$	$p \rightarrow q$	$p \leftrightarrow q$
T	T	F	T	T	T	T
F	T	T	F	T	T	F
T	F	F	F	T	F	F
F	F	T	F	F	T	T

Complex compound propositions?

Let's take a complex compound proposition:

$$q \vee ((\neg q) \wedge r)$$

$$q \text{ or } ((\text{not } q) \text{ and } r)$$

q	r	$\cdots$
T	T	$\cdots$
F	T	$\cdots$
T	F	$\cdots$
F	F	$\cdots$

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q	r	$\neg q$	$\cdots$
T	T	F	$\cdots$
F	T	T	$\cdots$
T	F	F	$\cdots$
F	F	T	$\cdots$

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Let's take a complex compound proposition:

$$q \vee ((\neg q) \wedge r)$$

$$q \text{ or } ((\text{not } q) \text{ and } r)$$

q	r	$\neg q$	$(\neg q) \wedge r$	$\cdots$
T	T	F	F	$\cdots$
F	T	T	T	$\cdots$
T	F	F	F	$\cdots$
F	F	T	F	$\cdots$

Complex compound propositions?

Let's take a complex compound proposition:

$$q \lor ((\neg q) \land r)$$

$$q \text{ or } ((\text{not } q) \text{ and } r)$$

q	r	$\neg q$	$(\neg q) \land r$	$q \lor ((\neg q) \land r)$
T	T	F	F	T
F	T	T	T	T
T	F	F	F	T
F	F	T	F	F

Complex compound propositions?

The number of rows in the truth table of a compound proposition is equal to 2^n, where n is the number of used propositional variables.

$$(\neg p) \vee ((q \rightarrow r) \wedge p)$$

p	q	r	$\cdots$
T	T	T	$\cdots$
F	T	T	$\cdots$
T	F	T	$\cdots$
F	F	T	$\cdots$
T	T	F	$\cdots$
F	T	F	$\cdots$
T	F	F	$\cdots$
F	F	F	$\cdots$

Each of the three variables can take two possible values, so the system has $2 \cdot 2 \cdot 2 = 8$ possible states.

Equivalence

Two compound propositions are equivalent if they have the same truth values for all possible cases (have the same truth tables).

q	r	$q \vee ((\neg q) \wedge r)$	$q \vee r$
T	T	T	T
F	T	T	T
T	F	T	T
F	F	F	F

Therefore, these two propositions are logically equivalent!

We write it as follows

$$q \vee ((\neg q) \wedge r) \quad \equiv \quad q \vee r$$

Note that the statement of the equivalence of two compound propositions, $a \equiv b$, is not a proposition itself.

Equivalent formulae

$$
\begin{array}{rl}
(a \wedge b) \equiv (b \wedge a) & \text{commutativity of } \wedge \\
(a \vee b) \equiv (b \vee a) & \text{commutativity of } \vee \\
((a \wedge b) \wedge c) \equiv (a \wedge (b \wedge c)) & \text{associativity of } \wedge \\
((a \vee b) \vee c) \equiv (a \vee (b \vee c)) & \text{associativity of } \vee \\
\neg(\neg a) \equiv a & \text{double-negation elimination} \\
(a \rightarrow b) \equiv (\neg b \rightarrow \neg a) & \text{contraposition} \\
(a \rightarrow b) \equiv (\neg a \vee b) & \text{implication elimination} \\
(a \leftrightarrow b) \equiv (a \rightarrow b) \wedge (b \rightarrow a) & \text{biconditional elimination} \\
\neg(a \wedge b) \equiv (\neg a \vee \neg b) & \text{De Morgan's Law} \\
\neg(a \vee b) \equiv (\neg a \wedge \neg b) & \text{De Morgan's Law} \\
(a \wedge (b \vee c)) \equiv (a \wedge b) \vee (a \wedge c) & \text{distributivity of } \wedge \text{ over } \vee \\
(a \vee (b \wedge c)) \equiv (a \vee b) \wedge (a \vee c) & \text{distributivity of } \vee \text{ over } \wedge
\end{array}
$$