

Recurrences

Recurrence

Recurrence

Towers of Hanoi

Merge sort

Time complexity of algorithms

Recall that we used induction to prove statements like

$$\sum_{k=0}^n k = 0 + 1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

In problems like this, we used a common pattern:

$$\sum_{k=0}^0 k = 0$$

$$\sum_{k=0}^n k = \left(\sum_{k=0}^{n-1} k \right) + n, \quad \text{when } n > 0$$

That is, we can express the sum of natural numbers recursively in terms of a smaller sum.

Recurrence

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

Let $S(n)$ be the sum of all natural numbers not greater than n :

$$S(n) = \sum_{k=0}^n k,$$

It can be convenient to redefine the sum $S(n)$ as a *recurrence*:

$$S(0) = 0$$

$$S(n) = S(n-1) + n \quad (\forall n > 0)$$

This is just another way to express the same function S .

Recurrence

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

Exponentiation:

$$E(a, n) = a^n$$

Recursively:

$$E(a, 0) = 1$$

$$E(a, n) = E(a, n-1) \cdot a \quad (\forall n > 0)$$

Recurrence

Recurrence

Towers of Hanoi

Merge sort

Time complexity of algorithms

Factorial:

$$n! = 1 \cdot 2 \cdot \dots \cdot n$$

Recursively:

$$0! = 1$$

$$n! = (n-1)! \cdot n \quad (\forall n > 0)$$

The towers of Hanoi

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

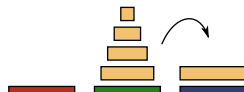
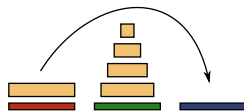
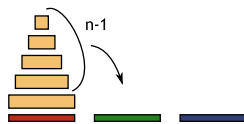


<http://www.mathsisfun.com/games/towerofhanoi.html>

The towers of Hanoi

Our recursive algorithm to move a tower of height n from #1 to #3:

1. Move an $(n-1)$ -tower from #1 to #2.
2. Move an 1-tower from #1 to #3.
3. Move an $(n-1)$ -tower from #2 to #3.



Recurrence

Towers of Hanoi

Merge sort

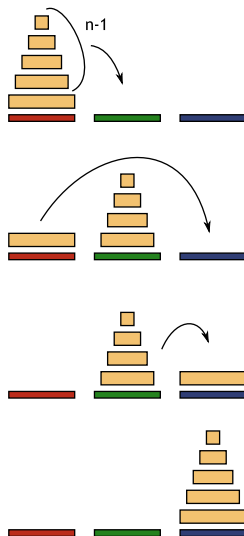
Time complexity of algorithms

The towers of Hanoi

Our recursive algorithm to move a tower of height n from #1 to #3:

1. Move an $(n-1)$ -tower from #1 to #2.
2. Move an 1-tower from #1 to #3.
3. Move an $(n-1)$ -tower from #2 to #3.

There is a way to find a recurrent formula for T_n , the total number of steps to move the tower from the peg 1 to the peg 3.



Recurrence

Towers of Hanoi

Merge sort

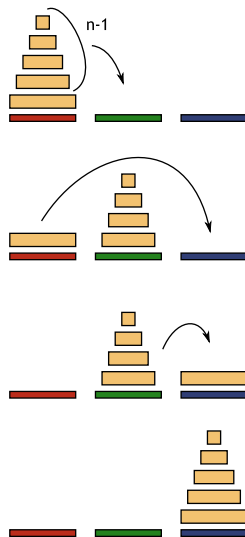
Time complexity of algorithms

The towers of Hanoi

T_n , the time to move a tower of height n :

$$T_1 = 1$$

$$T_n = T_{n-1} + 1 + T_{n-1} \quad (\forall n > 1)$$



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Towers of Hanoi

Merge sort

Time complexity of algorithms

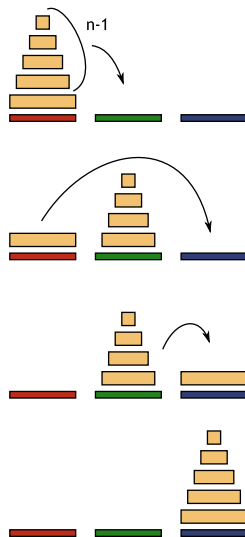
The towers of Hanoi

T_n , the time to move a tower of height n :

$$T_1 = 1$$

$$T_n = T_{n-1} + 1 + T_{n-1} \quad (\forall n > 1)$$

There is a proof by induction that this time is optimal for any algorithm.



Recurrence

Towers of Hanoi

Merge sort

Time complexity of algorithms

The towers of Hanoi

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

$$T_1 = 1$$

$$T_n = 2T_{n-1} + 1 \quad (\forall n > 1)$$

Our goal is to find a closed form expression for T_n as a function of n , without any recurrence.

The towers of Hanoi

Before we get a closed form formula for T_n ,
what are the numbers?

$$T_1 = 1$$

$$T_n = 2T_{n-1} + 1 \quad (\forall n > 1)$$

We can compute a list like this:

$$T_1 = 1$$

$$T_2 = 3$$

$$T_3 = 7$$

$$T_4 = 15$$

$$T_5 = 31$$

$$T_6 = 63$$

...

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

The towers of Hanoi

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

$$T_1 = 1$$

$$T_n = 2T_{n-1} + 1 \quad (\forall n > 1)$$

$$T_1 = 1$$

$$T_2 = 3$$

$$T_3 = 7$$

$$T_4 = 15$$

$$T_5 = 31$$

$$T_6 = 63$$

...

Guess and verify method... Let's try $T_n = 2^n - 1$?

The towers of Hanoi

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

$$T_1 = 1$$

$$T_n = 2T_{n-1} + 1 \quad (\forall n > 1)$$

Guess and verify method... Let's try $T_n = 2^n - 1$? We can show by induction that this formula is correct.

The base case, $n = 1$:

$$T_1 = 2^1 - 1 = 1.$$

Ok, the base case is true.

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Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

$$\begin{aligned}T_1 &= 1 \\T_n &= 2T_{n-1} + 1 \quad (\forall n > 1)\end{aligned}$$

We want to prove the closed form formula $T_n = 2^n - 1$.

The inductive step, $n > 1$:

Assume that $T_n = 2^n - 1$, and show that then $T_{n+1} = 2^{n+1} - 1$.

Proof. From the recurrence:

$$T_{n+1} = 2T_n + 1$$

By the inductive hypothesis:

$$2T_n + 1 = 2(2^n - 1) + 1 = 2^{n+1} - 2 + 1 = 2^{n+1} - 1$$

The towers of Hanoi

If it is difficult to guess the closed form expression for the recurrence, there is another technique:

$$T_1 = 1$$

$$T_n = 2T_{n-1} + 1$$

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

The towers of Hanoi

If it is difficult to guess the closed form expression for the recurrence, there is another technique:

$$T_1 = 1$$

$$T_n = 2T_{n-1} + 1$$

$$= 2(2T_{n-2} + 1) + 1$$

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

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$$= 2(2T_{n-2} + 1) + 1$$

$$= 2^2T_{n-2} + 2 + 1$$

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

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If it is difficult to guess the closed form expression for the recurrence, there is another technique:

$$T_1 = 1$$

$$T_n = 2T_{n-1} + 1$$

$$= 2(2T_{n-2} + 1) + 1$$

$$= 2^2T_{n-2} + 2 + 1$$

$$= 2^2(2T_{n-3} + 1) + 2 + 1$$

$$= 2^3T_{n-3} + 4 + 2 + 1$$

...

$$= 2^{n-1} \underbrace{T(n - (n-1))}_{=1} + \dots + 4 + 2 + 1$$

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

The towers of Hanoi

If it is difficult to guess the closed form expression for the recurrence, there is another technique:

$$T_1 = 1$$

$$T_n = 2T_{n-1} + 1$$

$$= 2(2T_{n-2} + 1) + 1$$

$$= 2^2 T_{n-2} + 2 + 1$$

$$= 2^2 (2T_{n-3} + 1) + 2 + 1$$

$$= 2^3 T_{n-3} + 4 + 2 + 1$$

...

$$= 2^{n-1} \underbrace{T(n - (n-1))}_{=1} + \dots + 4 + 2 + 1$$

$$= 2^{n-1} + \dots + 4 + 2 + 1 = \sum_{k=0}^{n-1} 2^k =$$

Recurrence

Towers of Hanoi

Merge sort

Time complexity of algorithms

The towers of Hanoi

If it is difficult to guess the closed form expression for the recurrence, there is another technique:

$$T_1 = 1$$

$$T_n = 2T_{n-1} + 1$$

$$= 2(2T_{n-2} + 1) + 1$$

$$= 2^2 T_{n-2} + 2 + 1$$

$$= 2^2 (2T_{n-3} + 1) + 2 + 1$$

$$= 2^3 T_{n-3} + 4 + 2 + 1$$

...

$$= 2^{n-1} \underbrace{T(n - (n-1))}_{=1} + \dots + 4 + 2 + 1$$

$$= 2^{n-1} + \dots + 4 + 2 + 1 = \sum_{k=0}^{n-1} 2^k = \frac{1-2^n}{1-2} = 2^n - 1.$$

Recurrence

Towers of Hanoi

Merge sort

Time complexity of algorithms

The towers of Hanoi

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

Why is it useful to know that the recurrence

$$T_1 = 1$$

$$T_n = 2T_{n-1} + 1 \quad (\forall n > 1)$$

is equivalent to the closed form formula $T_n = 2^n - 1$?

The 7-disk puzzle will require $T_7 = 2^7 - 1 = 127$ moves to complete.

And the 100-disk puzzle will require

$$T_{100} = 2^{100} - 1 = 1267650600228229401496703205375 \text{ moves.}$$

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

function Merge:

Given two sorted lists, combine them into a single sorted list:

$$[1, 2, 4, 5] + [3, 4, 5, 6] \mapsto [1, 2, 3, 4, 4, 5, 5, 6]$$

function Sort:

Given a list: if it contains a single element, return it. Otherwise, split it in two halves sort them separately and merge the results:

$$S[5] \mapsto [5]$$

$$S[6, 7, 1, 8, 9, 7, 4, 3] \mapsto S[6, 7, 1, 8] + S[9, 7, 4, 3]$$

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

$S[1, 8, 3, 6, 5, 4, 7, 2] \mapsto$

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

$$\begin{aligned} &S[1, 8, 3, 6, 5, 4, 7, 2] \mapsto \\ &S[1, 8, 3, 6] + S[5, 4, 7, 2] \mapsto \end{aligned}$$

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

$$S[1, 8, 3, 6, 5, 4, 7, 2] \mapsto$$

$$S[1, 8, 3, 6] + S[5, 4, 7, 2] \mapsto$$

$$(S[1, 8] + S[3, 6]) + (S[5, 4] + S[7, 2]) \mapsto$$

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of algorithms

$$S[1, 8, 3, 6, 5, 4, 7, 2] \mapsto$$

$$S[1, 8, 3, 6] + S[5, 4, 7, 2] \mapsto$$

$$(S[1, 8] + S[3, 6]) + (S[5, 4] + S[7, 2]) \mapsto$$

$$((S[1] + S[8]) + (S[3] + S[6])) + ((S[5] + S[4]) + (S[7] + S[2])) \mapsto$$

Merge sort

Recurrence

Towers of Hanoi

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$$S[1, 8, 3, 6, 5, 4, 7, 2] \mapsto$$

$$S[1, 8, 3, 6] + S[5, 4, 7, 2] \mapsto$$

$$(S[1, 8] + S[3, 6]) + (S[5, 4] + S[7, 2]) \mapsto$$

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Merge sort

Recurrence

Towers of Hanoi

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$$S[1, 8, 3, 6, 5, 4, 7, 2] \mapsto$$

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$$([1] + [8]) + ([3] + [6]) + ([5] + [4]) + ([7] + [2]) \mapsto$$

$$[1, 8] + [3, 6] + [4, 5] + [2, 7] \mapsto$$

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
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$$S[1, 8, 3, 6, 5, 4, 7, 2] \mapsto$$

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$$[1, 3, 6, 8] + [2, 4, 5, 7] \mapsto$$

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
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$$S[1, 8, 3, 6, 5, 4, 7, 2] \mapsto$$

$$S[1, 8, 3, 6] + S[5, 4, 7, 2] \mapsto$$

$$(S[1, 8] + S[3, 6]) + (S[5, 4] + S[7, 2]) \mapsto$$

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$$([1, 8] + [3, 6]) + ([4, 5] + [2, 7]) \mapsto$$

$$[1, 3, 6, 8] + [2, 4, 5, 7] \mapsto$$

$$[1, 2, 3, 4, 5, 6, 7, 8]$$

Merge sort

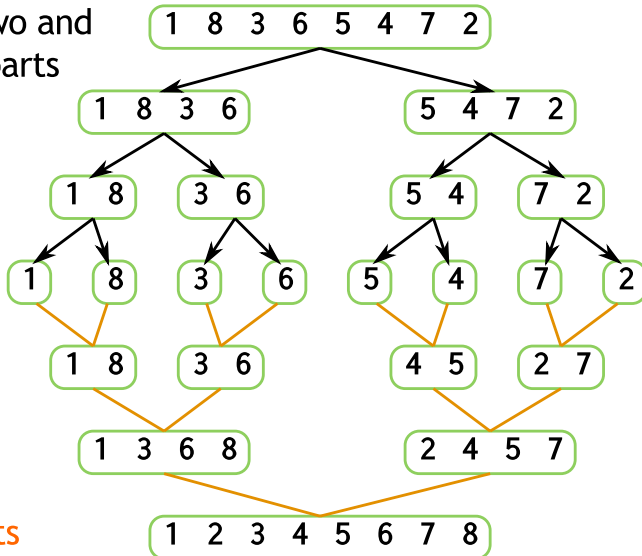
Recurrence

Towers of Hanoi

Merge sort

Time complexity of algorithms

Split in two and
sort the parts



Merge
sorted lists

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

How much time does it take to sort a list of n elements?

To estimate the time complexity, we are going to *count the number of comparisons* between the elements.

We assume that the size of the given list is a power of 2. It makes the analysis easier, but does not affect the result.

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

- (a) To merge two lists of size $n/2$, we need to do at most $n - 1$ comparisons.
- (b) To sort a list, we have to split it in two, sort both halves, and merge them.

Therefore,

$$T(1) = 0$$

$$T(n) = 2T(n/2) + n - 1 \quad (\forall n > 1)$$

Merge sort

Given

$$T(n) = 2T(n/2) + n - 1 \quad (\forall n > 1)$$

Since we assume that $n = 2^k$,

$$T(n) = T(2^k) =$$

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

Merge sort

Given

$$T(n) = 2T(n/2) + n - 1 \quad (\forall n > 1)$$

Since we assume that $n = 2^k$,

$$T(n) = T(2^k) = 2T(2^k/2) + 2^k - 1 = 2T(2^{k-1}) + (2^k - 1)$$

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

Given

$$T(n) = 2T(n/2) + n - 1 \quad (\forall n > 1)$$

Since we assume that $n = 2^k$,

$$\begin{aligned} T(n) &= T(2^k) = 2T(2^k/2) + 2^k - 1 = 2T(2^{k-1}) + (2^k - 1) \\ &= 2(2T(2^{k-2}) + 2^{k-1} - 1) + (2^k - 1) \\ &= 2^2T(2^{k-2}) + (2^k - 2) + (2^k - 1) \end{aligned}$$

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of algorithms

Given

$$T(n) = 2T(n/2) + n - 1 \quad (\forall n > 1)$$

Since we assume that $n = 2^k$,

$$\begin{aligned} T(n) &= T(2^k) = 2T(2^k/2) + 2^k - 1 = 2T(2^{k-1}) + (2^k - 1) \\ &= 2(2T(2^{k-2}) + 2^{k-1} - 1) + (2^k - 1) \\ &= 2^2T(2^{k-2}) + (2^k - 2) + (2^k - 1) \\ &= 2^2(2T(2^{k-3}) + 2^{k-2} - 1) + (2^k - 2) + (2^k - 1) \\ &= 2^3T(2^{k-3}) + (2^k - 4) + (2^k - 2) + (2^k - 1) \\ &= 2^3(2T(2^{k-4}) + 2^{k-3} - 1) + (2^k - 4) + (2^k - 2) + (2^k - 1) \\ &= 2^4T(2^{k-4}) + (2^k - 8) + (2^k - 4) + (2^k - 2) + (2^k - 1) \\ &= \dots = 2^k \underbrace{T(2^{k-k})}_{T(1)=0} + \sum_{i=0}^{k-1} (2^k - 2^i) = \sum_{i=0}^{k-1} (2^k - 2^i). \end{aligned}$$

Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of algorithms

$$T(n) = T(2^k) = \sum_{i=0}^{k-1} (2^k - 2^i) = \sum_{i=0}^{k-1} (n - 2^i) = n \cdot k - \sum_{i=0}^{k-1} 2^i.$$

The sum of the geometric progression is

$$\sum_{i=0}^{k-1} 2^i = \frac{2^k - 1}{2 - 1} = 2^k - 1 = n - 1.$$

Thus $T(n) = n \cdot k - n + 1$. And since $n = 2^k$, $k = \log_2 n$, so

$$T(n) = n \log_2 n - n + 1.$$

Merge sort

Recurrence

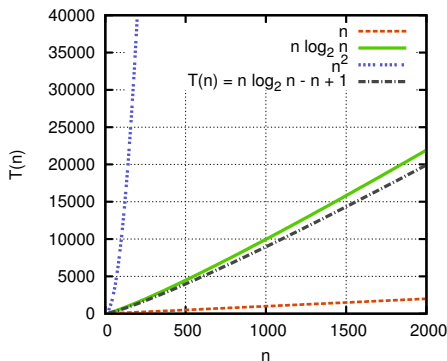
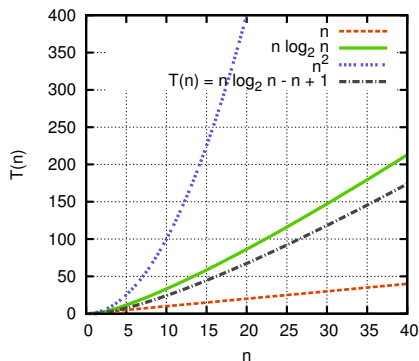
Towers of Hanoi

Merge sort

Time complexity of algorithms

To sort a list of length n , takes time (the number of comparisons)

$$T(n) = n \log_2 n - n + 1 \approx n \log_2 n.$$



Merge sort

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

In merge sort, we had a recurrence:

$$T(n) = 2T(n/2) + n - 1$$

In general, if the time complexity of an algorithm is expressed by a recurrence:

$$T(n) = a \cdot T(n/b) + f(n)$$

To solve such recurrences, there is a so called *Master theorem*:

https://en.wikipedia.org/wiki/Master_theorem

It covers different forms of the function f , as well as difference values of the constants a and b .

Time complexity, big-O

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

Let's say that we've got this function as an estimation of the time complexity of an algorithm:

$$T(n) = 6n \log_2 n + 100n + \log_2 n + 50$$

Informally:

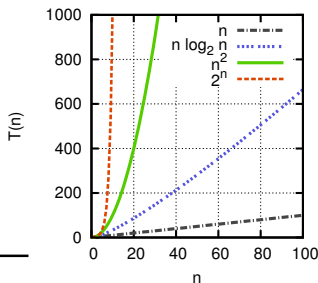
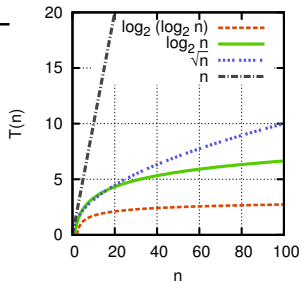
- (a) If $T(n)$ is a sum, we take the fastest growing term only.
- (b) We don't really care about constant factors.

$$T(n) = O(n \log_2 n)$$

Time complexity, big-O

Some common time complexities, from the slowest to the fastest:

Running time	Name
$O(1)$	constant
$O(\log(\log n))$	log-logarithmic
$O(\log n)$	logarithmic
$O(\sqrt{n})$	square root (sub-linear)
$O(n)$	linear
$O(n \log n)$	n-log-n
$O(n^2)$	quadratic
$O(2^n)$	exponential
$O(n!)$	factorial
$O(2^{(2^n)})$	double exponential



Recurrence

Towers of Hanoi

Merge sort

Time complexity of algorithms

Time complexity, big-O

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Towers of Hanoi

Merge sort

Time complexity of
algorithms

Faster than linear algorithms, for example $O(\log n)$, cannot go through the whole input. They are cherry-picking in some sense, knowing where to search for the answer. Usually the input is structured in some way.

Example: *Binary search in a sorted array.*

Linear time algorithms, $O(n)$, usually have to read the whole input.

Example: *Search for the largest element in an unsorted array.*

Time complexity, big-O

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

Algorithms *slower than linear*, $O(n \log n)$ or $O(n^2)$, $O(n^7)$. Not only read the whole input, but also perform some extra work, but in a reasonably efficient way.

Example: *Sorting algorithms.*

Algorithms *much slower than linear*, exponential, for example, $O(2^n)$, are doing some non-trivial work.

Example: *Satisfiability of a statement in propositional logic.*

Time complexity, big-O

Recurrence

Towers of Hanoi

Merge sort

Time complexity of
algorithms

Formally:

We say that

$$T(n) = O(f(n))$$

if there are constants C and k such that

$$|T(n)| \leq C|f(n)| \quad \text{for all } n > k$$

This definition says that after $n > k$, all slowly-growing terms don't really matter, and $T(n)$ behaves similarly to $f(n)$. To be more exact, $T(n)$ never exceeds $C \cdot f(n)$ when n is large enough.

$$T(n) = 6n \log_2 n + 100n + \log_2 n + 50$$

$$T(n) = O(n \log_2 n)$$