The Hessian of a Deep Net

Ali Rahimi

October 4, 2025

1 Introduction

The Hessian of a deep net is the matrix of second-order mixed partial derivatives of its loss with respect to the parameters. Equivalently, it's the derivative of the gradient with respect to the parameters. Over the past few decades, relying on this Hessian has come in and out of favor. In the 80s, when deep nets had only thousands of parameters, the Hessian could be used to implement second order optimizers that converged must faster than gradient descent [?]. As the number of parameters in deep nets grew, approximations of the Hessian were employed: low-rank estimates used in LBFGS [?], and diagonal approximations of the Hessian used in Adagrad [?]. Methods have also been proposed to approximate the spectrum of the Hessian to glean insights about the behavior of training algorithms [?]. However, as modern deep nets get larger, these methods have become increasingly less practical to apply: for a model with a billion parameters, the Hessian would have a quintillion entries, which far more than can be stored, multiplied, factorized, or inverted, even in the largest data centers.

We show that the Hessian of a deep net has a regular structure that makes it more amenable to these operations. Regardless of the specific operations in each layer, the Hessian can be represented as a matrix polynomial that involves the first order and second order mixed derivatives of each layer. This polynomial is also second order in the inverse of of a block-bi-diagonal operator that represents the backpropagatin algorithm. This structure allows us to multiply the Hessian by a vector, or solve linear systems involving the Hessian without ever forming or storing the full matrix. For an L-layer deep net where each layer has p parameters and produces a activations, naively storing the Hessian would require $O(L^2p^2)$ memory, and multiplying it by a vector or solving a linear system would require $O(L^2p^2 + L \max(a, p)^3)$ and $O(L^3p^3)$ operations, respectively. In contrast, we will show how to perform these operations using only $O(L \max(a, p)^3)$ computations. The dependence on the parameters on the number of activations and parameters in each layer is still cubic, but the dependence on the number of layers is only linear. This makes operating on the Hessian of tall and skinny networks more efficient than the Hessian of short and fat networks.

2 Overview

Our objective is to efficiently solve linear systems of the form Hx = g, where H is the Hessian of a deep neural network. Forming H explicitly is infeasible due to its size, and directly inverting or multiplying by H would require $O(L^3p^3)$ and $O(L^2p^2)$ operations, respectively—both prohibitively expensive for large networks. To overcome this, we employ the following strategy:

- 1. Write down the gradient of the deep net as a bi-diagonal system of linear equations. Solving this system uses back-substitution, which requires a forward and backward pass similar to those used in backpropagation. In fact, back-substitution and backpropagation are identical in this context.
- 2. Differentiate the components of this linear system, including the bi-diagonal matrix itself, to obtain the second-order derivatives.
- 3. Reformulate the resulting expressions so that the Hessian appears as a second-order polynomial in the inverse of the bi-diagonal matrix.
- 4. Derive a fast algorithm to apply the inverse of such polynomials.

3 Notation

We write a deep net as a pipeline of functions $\ell = 1, \dots, L$,

$$z_{1} = f_{1}(z_{0}; x_{1})$$
...
$$z_{\ell} = f_{\ell}(z_{\ell-1}; x_{\ell})$$
...
$$z_{L} = f_{L}(z_{L-1}; x_{L})$$
(1)

The vectors x_1, \ldots, x_L are the parameters of the pipeline. The vectors z_1, \ldots, z_L are its intermediate activations, and z_0 is the input to the pipeline. The last layer f_L computes the final activations and their training loss, so the scalar z_L is the loss of the model on the input z_0 . To make this loss's dependence on z_0 and the parameters explicit, we sometimes write it as $z_L(z_0;x)$. This formalization deviates slightly from the traditional deep net formalism in two ways: First, the training labels are subsumed in z_0 , and are propagated through the layers until they're used in the loss. Second, the last layer fuses the loss (which has no parameters) and the last layer (which does).

4 Backpropagation, the Matrix Way

We would like to fit the vector of parameters $x = (x_1, ..., x_L)$ given a training dataset, which we represent by a stochastic input z_0 to the pipeline. Training the model proceeds by gradient descent steps along the stochastic gradient $\partial z_L(z_0; x)/\partial x$. The components of this direction can be computed by the chain rule with a backward recursion:

$$\frac{\partial z_L}{\partial x_\ell} = \underbrace{\frac{\partial z_L}{\partial z_\ell}}_{b_\ell} \underbrace{\frac{\partial z_\ell}{\partial x_\ell}}_{\nabla_x f_\ell}$$

$$\frac{\partial z_L}{\partial z_\ell} = \underbrace{\frac{\partial z_L}{\partial z_{\ell+1}}}_{b_{\ell+1}} \underbrace{\frac{\partial z_{\ell+1}}{\partial z_\ell}}_{\nabla_x f_{\ell+1}}$$

The identification $b_\ell \equiv \frac{\partial z_L}{\partial z_\ell}$, $\nabla_x f_\ell \equiv \frac{\partial z_\ell}{\partial x_\ell}$, and $\nabla_z f_\ell \equiv \frac{\partial z_\ell}{\partial z_{\ell-1}}$ turns this recurrence into

$$\frac{\partial z_L}{\partial x_\ell} = b_\ell \cdot \nabla_x f_\ell$$
$$b_\ell = b_{\ell+1} \cdot \nabla_z f_{\ell+1},$$

with the base case $b_L = 1$, a scalar.

The above equations can be written in vector form as

$$\frac{\partial z_L}{\partial x} = \begin{bmatrix} \frac{\partial z_L}{\partial x_1} & \cdots & \frac{\partial z_L}{\partial x_L} \end{bmatrix} = \underbrace{\begin{bmatrix} b_1 & \cdots & b_L \end{bmatrix}}_{\equiv b} \underbrace{\begin{bmatrix} \nabla_x f_1 \\ & \ddots \\ & \nabla_x f_L \end{bmatrix}}_{\equiv D}$$
(2)

and

$$\begin{bmatrix} b_1 & b_2 & b_3 & \cdots & b_{L-1} & b_L \end{bmatrix} \underbrace{\begin{bmatrix} I & I & & & \\ -\nabla_z f_2 & I & & & \\ & -\nabla_z f_3 & I & & \\ & & \ddots & \ddots & \\ & & & -\nabla_z f_L & 1 \end{bmatrix}}_{\equiv M} = \underbrace{\begin{bmatrix} 0 & \cdots & 1 \end{bmatrix}}_{\equiv e_L}. \tag{3}$$

Solving for b and substituting back gives

$$\frac{\partial z_L}{\partial x} = e_L M^{-1} D. \tag{4}$$

The matrix M is block bi-diagonal. Its diagonal entries are identity matrices, and its off-diagonal matrices are the gradient of the intermediate activations with respect to the layer's parameters. The matrix D is block diagonal, with the block as the derivative of each layer's activations with respect to its inputs. M is invertible because the spectrum of a triangular matrix can be read off its diagonal, which in this case is all ones.

5 The Hessian

The gradient we computed above is the unique vector v such that $dz_L \equiv z_L(x+dx) - z_L(dx) \to v(x) \cdot dx$ as $dx \to 0$. In this section, we compute the Hessian H of z_L with respect to the parameters. This is the unique matrix H(x) such that $dv^{\top} \equiv v^{\top}(x+dx) - v^{\top}(x) \to H(x) dx$ as $dx \to 0$. We use the facts that $dM^{-1} = -M^{-1}(dM)M^{-1}$ and $b = e_L M^{-1}$ to write

$$dv = d(e_L M^{-1} D)$$

$$= e_L M^{-1} (dD) + e_L (dM^{-1}) D$$

$$= b \cdot dD - e_L M^{-1} (dM) M^{-1} D$$

$$= b \cdot dD - b \cdot (dM) M^{-1} D$$
(5)

We compute each of these terms separately.

As part of this agenda, we need to define the gradient of tensor-valued functions $g: \mathbb{R}^d \to \mathbb{R}^{o_1 \times \cdots \times o_k}$. We define this gradient $\nabla_x g(x) \in \mathbb{R}^{(o_1 \cdots o_k) \times d}$ as the unique matrix-valued function that satisfies vec $(g(x+dx)-g(x)) \to \nabla_x g(x)$ dx as $dx \to 0$. This convention allows us to readily define, for example, the Hessian of a vector-valued function: If $g: \mathbb{R}^d \to \mathbb{R}^o$, then $\nabla_{xx} g(x) \in \mathbb{R}^{o \times d^2}$ is the unique matrix such that vec $(\nabla_x g(x+dx) - \nabla_x g(x)) \to \nabla_{xx} g(x)$ dx. This convention also supports the chain rule as expected: For example, the gradient of $h(x) \equiv f(g(x))$ for a matrix-valued f and g can be written as $\nabla f \nabla g$ as expected. The chain rule in turn lets us define mixed partial derivatives, like $\nabla_{yz} g$ for $g: \mathbb{R}^{|x|} \to \mathbb{R}^{|g|}$. For example, if $g \in \mathbb{R}^{|y|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|x|}$ to some $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|z|}$ to some $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|z|}$ to some $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|z|}$ to some $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|z|}$ to some $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$ are restrictions of $g \in \mathbb{R}^{|z|}$ and $g \in \mathbb{R}^{|z|}$

5.1 The Term Involving dD

The matrix D is block-diagonal with its ℓ th diagonal block containing the matrix $D_{\ell} \equiv \nabla_x f_{\ell}$. Using the facts that $\operatorname{vec}(ABC) = (C^{\top} \otimes A) \operatorname{vec}(B)$, and $(A \otimes B)^{\top} = A^{\top} \otimes B^{\top}$, we get

$$b \cdot (dD) = \begin{bmatrix} b_1 & \cdots & b_L \end{bmatrix} \begin{bmatrix} dD_1 & & & \\ & \ddots & & \\ & & dD_L \end{bmatrix}$$

$$= \begin{bmatrix} b_1 \cdot dD_1 & \cdots & b_L \cdot dD_L \end{bmatrix}$$

$$= \begin{bmatrix} \operatorname{vec} (dD_1)^\top \left(I \otimes b_1^\top \right) & \cdots & \operatorname{vec} (dD_L)^\top \left(I \otimes b_L^\top \right) \end{bmatrix}$$

$$= \begin{bmatrix} \operatorname{vec} (dD_1) \\ \vdots \\ \operatorname{vec} (dD_L) \end{bmatrix}^\top \begin{bmatrix} I \otimes b_1^\top & & \\ & \ddots & \\ & & I \otimes b_L^\top \end{bmatrix}$$

$$(6)$$

Observe that $\operatorname{vec}(dD_{\ell}) = d \operatorname{vec} \nabla_x f_{\ell}(z_{\ell-1}; x_{\ell})$ varies with dx through both its arguments x_{ℓ} and $z_{\ell-1}$. Using mixed partials of vector-valued functions described above, we get

$$\operatorname{vec}(dD_{\ell}) = d\operatorname{vec}(\nabla_x f_{\ell}) = (\nabla_{xx} f_{\ell}) \ dx_{\ell} + (\nabla_{zx} f_{\ell}) \ dz_{\ell-1}. \tag{7}$$

Stacking these equations gives

$$\begin{bmatrix} \operatorname{vec} (dD_1) \\ \vdots \\ \operatorname{vec} (dD_L) \end{bmatrix} = \begin{bmatrix} \nabla_{xx} f_1 \\ & \ddots \\ & & \nabla_{xx} f_L \end{bmatrix} dx + \begin{bmatrix} \nabla_{zx} f_1 \\ & \ddots \\ & & \nabla_{zx} f_L \end{bmatrix} \begin{bmatrix} dz_0 \\ \vdots \\ dz_{L-1} \end{bmatrix}. \tag{8}$$

Each vector dz_{ℓ} in turn varies with dx via $dz_{\ell} = (\nabla_x f_{\ell}) dx_{\ell} + (\nabla_z f_{\ell}) dz_{\ell-1}$, with the base case $dz_0 = 0$, since the input z_0 does not vary with dx. Stacking up this recurrence gives

$$\begin{bmatrix} I \\ -\nabla_z f_2 & I \\ & \ddots \\ & -\nabla_z f_L & 1 \end{bmatrix} \begin{bmatrix} dz_1 \\ \vdots \\ dz_{L-1} \\ dz_L \end{bmatrix} = \begin{bmatrix} \nabla_x f_1 \\ & \ddots \\ & & \nabla_x f_L \end{bmatrix} dx. \tag{9}$$

We can solve for the vector $\begin{bmatrix} dz_1 \\ \vdots \\ dz_L \end{bmatrix} = M^{-1}Ddx$ and use the downshifting matrix

$$P \equiv \begin{bmatrix} 0 \\ I & 0 \\ & \ddots \\ & I & 0 \end{bmatrix} \tag{10}$$

to plug back the vector $\begin{bmatrix} dz_0 \\ \vdots \\ dz_{L-1} \end{bmatrix} = PM^{-1}Ddx$:

$$\begin{bmatrix} \operatorname{vec} (dD_1) \\ \vdots \\ \operatorname{vec} (dD_L) \end{bmatrix} = \begin{pmatrix} \begin{bmatrix} \nabla_{xx} f_1 & & \\ & \ddots & \\ & & \nabla_{xx} f_L \end{bmatrix} + \begin{bmatrix} \nabla_{zx} f_1 & & \\ & \ddots & \\ & & \nabla_{zx} f_L \end{bmatrix} PM^{-1}D dx. \tag{11}$$

5.2 The Term Involving dM

The matrix dM is lower-block-diagonal with dM_2, \ldots, dM_L , and $dM_\ell \equiv d\nabla_z f_\ell$. Similar to the above, we can write

$$b \cdot (dM)M^{-1}D = \begin{bmatrix} b_1 & \cdots & b_{L-1} & b_L \end{bmatrix} \begin{bmatrix} 0 \\ -dM_2 & 0 \\ & \ddots & \\ & -dM_L & 0 \end{bmatrix} M^{-1}D$$

$$= -\begin{bmatrix} b_2 \cdot dM_2 & \cdots & b_L \cdot dM_L & 0 \end{bmatrix} M^{-1}D$$

$$= -\begin{bmatrix} \operatorname{vec}(dM_2)^{\top} \left(I \otimes b_2^{\top} \right) & \cdots & \operatorname{vec}(dM_L)^{\top} \left(I \otimes b_L^{\top} \right) & 0 \end{bmatrix} M^{-1}D$$

$$= -\begin{bmatrix} \operatorname{vec}(dM_1) \\ \vdots \\ \operatorname{vec}(dM_L) \end{bmatrix}^{\top} \begin{bmatrix} 0 \\ I \otimes b_2^{\top} & 0 \\ & \vdots \\ I \otimes b_L^{\top} & 0 \end{bmatrix} M^{-1}D$$

$$= -\begin{bmatrix} \operatorname{vec}(dM_1) \\ \vdots \\ \operatorname{vec}(dM_L) \end{bmatrix}^{\top} \begin{bmatrix} I \otimes b_1^{\top} \\ & \ddots \\ & I \otimes b_L^{\top} \end{bmatrix} PM^{-1}D. \tag{12}$$

Each matrix $dM_{\ell} = d\nabla_z f_{\ell}(z_{\ell-1}; x_{\ell})$ varies with dx through both x_{ℓ} and $z_{\ell-1}$ as d vec $(M_{\ell}) = (\nabla_{xz} f_{\ell}) dx_{\ell} + (\nabla_{zz} f_{\ell}) dz_{\ell-1}$. Following the steps of the previous section gives

$$\begin{bmatrix} \operatorname{vec} (dM_1) \\ \vdots \\ \operatorname{vec} (dM_L) \end{bmatrix} = \begin{pmatrix} \begin{bmatrix} \nabla_{xz} f_1 & & \\ & \ddots & \\ & & \nabla_{xz} f_L \end{bmatrix} + \begin{bmatrix} \nabla_{zz} f_1 & & \\ & \ddots & \\ & & \nabla_{zz} f_L \end{bmatrix} PM^{-1}D \end{pmatrix} dx. \tag{13}$$

5.3 Putting it all Together

We have just shown that the Hessian of the deep net has the form

$$H \equiv \frac{\partial^2 z_L}{\partial x^2} = D_D \left(D_{xx} + D_{zx} P M^{-1} D_x \right) + D_x^{\top} M^{-T} P^{\top} D_M \left(D_{xz} + D_{zz} P M^{-1} D_x \right)$$
(14)
$$= D_D D_{xx} + D_D D_{zx} P M^{-1} D_x + D_x^{\top} M^{-T} P^{\top} D_M D_{xz} + D_x^{\top} M^{-T} P^{\top} D_M D_{zz} P M^{-1} D_x.$$
(15)

The definitions below annotate the size of the various matrices in this expression assuming all but the last L layers have a-dimensional activations ($z_{\ell} \in \mathbb{R}^{a}$) and p-dimensional parameters ($x_{\ell} \in \mathbb{R}^{p}$):

$$D_D \equiv \begin{bmatrix} I \otimes b_1 \\ \vdots \\ P \equiv \begin{bmatrix} 0 \\ I & 0 \\ \vdots \\ I & 0 \end{bmatrix}, \qquad D_M \equiv \begin{bmatrix} I \otimes b_1 \\ \vdots \\ -\nabla_z f_2 \\ \vdots \\ -\nabla_z f_3 \end{bmatrix} I \\ -\nabla_z f_3 \quad I \\ \vdots \\ -\nabla_z f_L \quad 1 \end{bmatrix},$$

$$D_x \equiv \begin{bmatrix} \nabla_x f_1 \\ \vdots \\ \nabla_x f_L \\ \vdots \\ \nabla_x f_L \end{bmatrix}, \qquad D_{xx} \equiv \begin{bmatrix} \nabla_x f_1 \\ \vdots \\ \nabla_x f_1 \\ \vdots \\ \nabla_x f_L \end{bmatrix}, \qquad D_{xz} \equiv \begin{bmatrix} \nabla_x f_1 \\ \vdots \\ \nabla_x f_L \\ \vdots \\ \nabla_x f_L \end{bmatrix}, \qquad D_{xz} \equiv \begin{bmatrix} \nabla_x f_1 \\ \vdots \\ \nabla_x f_L \\ \vdots \\ \nabla_x f_L \end{bmatrix}, \qquad D_{xz} \equiv \begin{bmatrix} \nabla_x f_1 \\ \vdots \\ \nabla_x f_L \\ \vdots \\ \nabla_x f_L \end{bmatrix},$$

$$D_{xx} \equiv \begin{bmatrix} \nabla_x f_1 \\ \vdots \\ \nabla_x f_L \\ \vdots \\ \nabla_x f_L \end{bmatrix}, \qquad D_{xz} \equiv \begin{bmatrix} \nabla_x f_1 \\ \vdots \\ \nabla_x f_L \\ \vdots \\ \nabla_x f_L \end{bmatrix}.$$

5.4 Mulitplying a vector by the Hessian

Given a vector $g \in \mathbb{R}^{Lp}$, one the formula above allows us to compute Hg in $O\left(Lap^2 + La^2p + La^3\right)$ operations without forming H. This cost is dominated by multiplying by the D_{xx} , D_{zx} , and D_{zz} matrices.

It's tempting to use this insight to use Krylove methods to solve systems of the form Hx = b without forming H. This would require compute Hg some number of times that depends on the condition number of H. However, the next section shows how to compute $H^{-1}b$ with roughly only as many operations as are needed to compute Hg.

6 The Inverse of the Hessian

The above shows that the Hessian is a second order matrix polynomial in M^{-1} . While M itself is block-bidiagonal, M^{-1} is dense, so H is dense. Nevertheless, this polynomial can be lifted into a higher order

object whose inverse is easy to compute:

$$\begin{split} H &= D_D D_{xx} + D_D D_{zx} P M^{-1} D_x + D_x^\top M^{-\top} P^\top D_M D_{xz} + D_x^\top M^{-\top} P^\top D_M D_{zz} P M^{-1} D_x \\ &= \begin{bmatrix} M^{-1} D_x \\ I \end{bmatrix}^\top \begin{bmatrix} P^\top D_M D_{zz} P & P^\top D_M D_{xz} \\ D_D D_{zx} P & D_D D_{xx} \end{bmatrix} \begin{bmatrix} M^{-1} D_x \\ I \end{bmatrix} \\ &= I + \begin{bmatrix} D_x \\ I \end{bmatrix}^\top \underbrace{\begin{bmatrix} M^{-\top} \\ D_D D_{zx} P & D_D D_{zx} P \end{bmatrix}}_{\hat{M}^{-\top}} \underbrace{\begin{bmatrix} D_x \\ D_D D_{zx} P & D_D D_{xx} - I \end{bmatrix}}_{\equiv Q} \underbrace{\begin{bmatrix} M^{-1} \\ I \end{bmatrix}}_{\hat{M}^{-1}} \underbrace{\begin{bmatrix} D_x \\ I \end{bmatrix}}_{\hat{M}^{-1}}. \end{split}$$

The Woodbury formula gives

$$H^{-1} = I - \begin{bmatrix} D_x \\ I \end{bmatrix}^{\top} \left(\left(\hat{M}^{-\top} Q \hat{M}^{-1} \right)^{-1} + \begin{bmatrix} D_x \\ I \end{bmatrix} \begin{bmatrix} D_x \\ I \end{bmatrix}^{\top} \right)^{-1} \begin{bmatrix} D_x \\ I \end{bmatrix}$$
$$= I - \begin{bmatrix} D_x \\ I \end{bmatrix}^{\top} \left(\underbrace{\hat{M} Q^{-1} \hat{M}^{\top} + \begin{bmatrix} D_x D_x^{\top} & D_x \\ D_x^{\top} & I \end{bmatrix}}_{=A} \right)^{-1} \begin{bmatrix} D_x \\ I \end{bmatrix}. \tag{16}$$

The matrix Q^{-1} can be computed explicitly using the partitioned matrix inverse formula. Define the Schur complement $S = Q_{11} - Q_{12}Q_{22}^{-1}Q_{21}$, where Q_{ij} denote the i, jth block of Q as defined above. Then

$$Q^{-1} = \begin{bmatrix} S^{-1} & -S^{-1}Q_{12}Q_{22}^{-1} \\ -Q_{22}^{-1}Q_{21}S^{-1} & Q_{22}^{-1} + Q_{22}^{-1}Q_{21}S^{-1}Q_{12}Q_{22}^{-1} \end{bmatrix}.$$
 (17)

The matrices Q_{11} , Q_{12} , Q_{21} , and Q_{22} are all block-diagonal. S is also block diagonal because Q_{11} and $Q_{12}Q_{22}^{-1}Q_{21}$ are both block-diagonal. Since all the terms involved in the blocks of Q^{-1} are block-diagonal, Q^{-1} has the same banded structure as Q.

The inverse of $A \equiv \hat{M}Q^{-1}\hat{M}^{\top} + \begin{bmatrix} D_x D_x^{\top} & D_x \\ D_x^{\top} & I \end{bmatrix}$ can be applied efficiently. Instead of applying the

Woodbury formula again, we compute its LDL^{\top} decomposition and apply the inverse of that decomposition. The LDL^{\top} decomposition of A is

$$A = \begin{bmatrix} I & A_{12}A_{22}^{-1} \\ 0 & I \end{bmatrix} \begin{bmatrix} A_{11} - A_{12}A_{22}^{-1}A_{12}^{\top} & 0 \\ 0 & A_{22} \end{bmatrix} \begin{bmatrix} I & A_{12}A_{22}^{-1} \end{bmatrix}^{\top}$$

$$A_{11} = M \begin{bmatrix} Q^{-1} \end{bmatrix}_{11} M^{\top} + D_x D_x^{\top}$$

$$A_{12} = M \begin{bmatrix} Q^{-1} \end{bmatrix}_{12} + D_x$$

$$A_{22} = \begin{bmatrix} Q^{-1} \end{bmatrix}_{22} + I.$$
(18)

so

$$A^{-1} = \begin{bmatrix} I & -A_{12}A_{22}^{-1} \\ 0 & I \end{bmatrix}^{\top} \begin{bmatrix} A_{11} - A_{12}A_{22}^{-1}A_{12}^{\top} & 0 \\ 0 & A_{22} \end{bmatrix}^{-1} \begin{bmatrix} I & -A_{12}A_{22}^{-1} \\ 0 & I \end{bmatrix}$$
(19)

Since A_{11} is block tri-diagonal, A_{12} is block-bidiagonal, and A_{22} is block-diagonal, applying A^{-1} to a vector is fast.

Summary: Algorithm to compute $H^{-1}g$

Given a vector $g \in \mathbb{R}^{Lp}$, compute $H^{-1}g$ as follows:

1. Compute the auxiliary vector:

$$g' \in \mathbb{R}^{La+Lp} \equiv \begin{bmatrix} D_x \\ I \end{bmatrix} g.$$

 D_x has L $a \times p$ blocks on its diagonal, so it takes Lap multiplications to compute v.

2. Form the banded matrix:

$$A \in \mathbb{R}^{L(a+p) \times L(a+p)} \equiv \hat{M} Q^{-1} \hat{M}^{\top} + \begin{bmatrix} D_x D_x^{\top} & D_x \\ D_x^{\top} & I \end{bmatrix}.$$

To compute Q^{-1} , we first compute the blocks of Q. These take La^3 multiplications for $Q_{11} \in \mathbb{R}^{La \times La}$, La^2p for $Q_{12} \in \mathbb{R}^{La \times Lp}$ and $Q_{21} \in \mathbb{R}^{Lp \times La}$, and La^2p for $Q_{22} \in \mathbb{R}^{Lp \times Lp}$. Computing $S \in \mathbb{R}^{La \times La}$ takes Lp^3 to compute Q_{22}^{-1} , and $2Lap^2$ to compute the product $Q_{12}Q_{22}^{-1}Q_{21}$. Given these quantities, for the blocks of Q^{-1} , it takes an additional La^3 to compute the upper left block, $L(a^2p+ap^2)$ to compute the off-diagonal blocks, and somewhat less than that to compute the bottom diagonal block since the matrices involved have already been computed. In all, it takes less than $9L \max(a,p)^3$ multiplications to compute Q^{-1} .

To compute $\hat{M}Q^{-1}\hat{M}^{\top}$ requires an additional $2La^3$ operations for a total of $11L \max(a,p)^3$ multiplications.

Finally, computing and adding the second term requires La^2p multiplications, bringing the tally to at most $12L \max(a, p)^3$ multiplications to compute A.

3. Apply A^{-1} to g':

$$g'' = \begin{bmatrix} I & -A_{12}A_{22}^{-1} \\ 0 & I \end{bmatrix}^{\top} \begin{bmatrix} A_{11} - A_{12}A_{22}^{-1}A_{12}^{\top} & 0 \\ 0 & A_{22} \end{bmatrix}^{-1} \begin{bmatrix} I & -A_{12}A_{22}^{-1} \\ 0 & I \end{bmatrix} g'.$$

This computation requires $2L \max(a, p)^3$ multiplications to compute A_{22}^{-1} and $\left[A_{11} - A_{12}A_{22}^{-1}A_{12}\right]^{-1}$. The remaining operations are matrix multiplications that take at most $3L \max(a, p)^2$, which is smaller than Lp^3 when p > 3. This brings the tally to at most $15L \max(a, p)^3$ multiplications.

4. Compute the final result:

$$y = g - \begin{bmatrix} D_x \\ I \end{bmatrix}^\top g''.$$

These are again matrix-vector multiplications that take at most $L \max(a, p)^2$ when p > 1, bringing the tally to at most $16L \max(a, p)^3$.