实验报告

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1.实验题目: 交换机学习实验

2.实验内容:

- 1 操作数据结构mac_port_map,实现查询lookup_port,插入insert_mac_port,删除老化节点 sweep_aged_mac_port_entry三个操作
- 2 实现广播操作
- 3 实现交换机学习的算法:
 - 。 在转发表里查找是否有源地址。
 - 如果有原地址目的地址。如果有,更新访问时间
 - 如果没有:把源地址加入到转发表,更新访问时间
 - 。 在转发表里查找是否有目的地址。
 - 如果有,则向这个目的地址发消息,并更新目的地址的访问时间
 - 如果没有,则广播这则消息
 - 。 用另一个线程检查转发表,删去一定时间内没有被访问的地址
- 4 比较广播与交换机学习两种机制的带宽利用情况

3.实验过程

mac_port_map的三个操作

1.删去老化的port

```
int sweep_aged_mac_port_entry()
       // TODO: implement the sweeping process here
       static int i = 0;
       int sleep_time = MAC_PORT_TIMEOUT;
       while(1){
                //fprintf(stdout, "%d sweeping\n",++i);
               pthread_mutex_lock(&mac_port_map.lock);
               mac_port_entry_t* entry;
               for(int i=0; i<HASH_8BITS ; ++i){</pre>
                       mac_port_entry_t *head,*tail;
                       tail=head=mac_port_map.hash_table[i];
                       for(entry=mac_port_map.hash_table[i]; entry!=NULL ; ){
                               //时间超过30s没访问
                               if(time(NULL)-entry->visited >= sleep_time){
                                       mac_port_entry_t* p;
                                       p=entry;
                                       //如果这个entry是头部: 需要把hash表项换成下个entry
                                       if(head == entry){
                                               tail=head=mac port map.hash table[i]=er
                                       }
                                       //如果不是头部,链表删掉这个节点
                                       else{
                                               tail->next = entry->next;
                                       entry=entry->next;
                                       free(p);
                               }
                               //跳过
                               else{
                                       tail = entry;
                                       entry=entry->next;
                               }
                        }
                pthread_mutex_unlock(&mac_port_map.lock);
                sleep(sleep_time);
       }
       return 0;
}
```

2.插入新的port

```
void insert_mac_port(u8 mac[ETH_ALEN], iface_info_t *iface)
{
       pthread mutex lock(&mac port map.lock);
       // TODO: implement the insertion process here
        //fprintf(stdout, "insert_mac_port\n");
       mac_port_entry_t* entry = mac_port_map.hash_table[hash8(mac,ETH_ALEN)];
        if(entry == NULL){
               entry = (mac_port_entry_t*)malloc(sizeof(mac_port_entry_t));
               memcpy(entry->mac,mac,ETH_ALEN*sizeof(u8));
               entry->iface = iface;
               entry->visited = time(NULL);
               entry->next = NULL;
               mac port map.hash table[hash8(mac,ETH ALEN)] = entry;
        }
        else{
               //沿着链表查找。如果找到,更新时间。如果没找到(查到链表尾部),插入新表
               int find = 0;
               mac_port_entry_t* p=entry;
               mac_port_entry_t* tail = NULL; //链表最后一个非空node
               for(p=entry ; p!=NULL ; p=p->next){
                       if(memcmp(p->mac,mac,ETH_ALEN*sizeof(u8))==0){
                               entry = p;
                               entry->visited = time(NULL);
                               find = 1;
                               break;
                       }
                       tail = p;
               if(!find){
                       entry = (mac_port_entry_t*)malloc(sizeof(mac_port_entry_t));
                       memcpy(entry->mac,mac,ETH_ALEN*sizeof(u8));
                       entry->iface = iface;
                       entry->visited = time(NULL);
                       entry->next = NULL;
                       tail->next = entry;
               }
        }
        pthread_mutex_unlock(&mac_port_map.lock);
}
```

3.在转发表中查找port

```
iface_info_t *lookup_port(u8 mac[ETH_ALEN])
        pthread mutex lock(&mac port map.lock);
        // TODO: implement the lookup process here
        //fprintf(stdout, "lookup_port\n");
        mac_port_entry_t* entry = mac_port_map.hash_table[hash8(mac,ETH_ALEN)];
        if(entry!= NULL){
                for( ; entry!=NULL ; entry=entry->next){
                        //查到
                        if(memcmp(entry->mac,mac,ETH_ALEN*sizeof(u8))==0){
                                entry->visited = time(NULL);
                                break;
                        }
                }
        }
        pthread_mutex_unlock(&mac_port_map.lock);
        return (entry==NULL)?NULL:entry->iface;
}
```

广播

```
void broadcast_packet(iface_info_t *iface, const char *packet, int len)
{
    // TODO: implement the broadcast process here
    //fprintf(stdout, "TODO: implement the broadcast process here.\n");
    //mine
    iface_info_t *ift = NULL;
    list_for_each_entry(ift, &instance->iface_list, list) {
        if(ift != iface){
            iface_send_packet(ift,packet,len);
        }
    }
}
```

交换机学习

```
void handle_packet(iface_info_t *iface, char *packet, int len)
        struct ether_header *eh = (struct ether_header *)packet;
        //log(DEBUG, "the dst mac address is " ETHER_STRING ".\n", ETHER_FMT(eh->ether_
        // TODO: implement the packet forwarding process here
        //fprintf(stdout, "handle.\n");
        iface_info_t* d_ift=lookup_port(eh->ether_dhost),
                                * s_ift=lookup_port(eh->ether_shost);
        if(d_ift != NULL){
                //fprintf(stdout, "find.\n");
                iface_send_packet(d_ift,packet,len);
        }
        else{
                //fprintf(stdout, "broadcast****.\n");
                broadcast_packet(iface,packet,len);
        }
        if(s_ift==NULL){
                insert_mac_port(eh->ether_shost,iface);
        }
}
```

4.实验结果

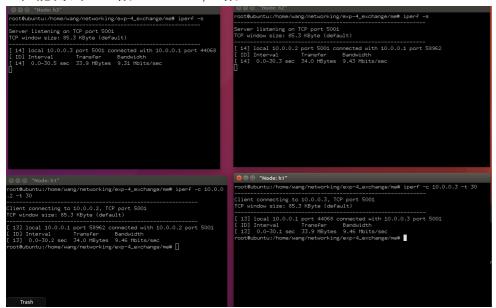
1.互相ping

```
**CottAbbuntus/Phome/wang/networking/exp-4_exchange/me# ping 10.0.0.2 -c 4
PING 10.0.0.2 (10.0.0.2) 56(84) bytes of data.
64 bytes from 10.0.0.2: icmp_seapt titled times 0.16 ms
64 bytes from 10.0.0.2: icmp_seapt titled times 0.16 ms
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64 bytes from 10.0.0.3: icmp_seapt titled times 0.16 ms
64 bytes from 10.0.0.3: icmp_sea
```

2.性能测试: h1做service, h2,h3做clinet

```
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```

3.性能测试: h1做clinet, h2,h3做service



5.结果分析

广播算法性能测试: h1做clinet, h2,h3做service

```
root@ubuntu:/home/wang/networking/exp-3_broadcast/04-broadcast/me1# iperf =8
root@ubuntu:/home/wang/networking/exp-3_broadcast/me1# iperf =8
root@ubuntu:/home/wang/networking/exp-3_broadcast
```

h1,h2,h3最大带宽为20Mb/s,10Mb/s,10Mb/s

• h1 server, h2,h3 clinet

结果:

h2,h3不管是同时向h1发消息,还是分别发送,都是接近10Mb/s.学习算法和广播都是这样的结果。

分析:

对于学习算法,h2向h1发消息时,交换机不向h3发消息。h3亦然。由于h1的最大带宽为h2,h3最大带宽之和,所以即使h2,h3同时向h1发消息,h1的带宽也能容纳,因此h2,h3都接近最大发送带宽。

对于广播算法,虽然h2向h1发消息时也向h3发消息,但占用的是h3的接收带宽而不占用h3发送带宽。因而h2,h3同时发送也接近最大发送带宽

• h1 clinet, h2,h3 server

结果:

对于学习算法,h1无论是单独还是同时向h2、h3发消息,都接近最大带宽。对于广播算法,h1单独向h2、h3发消息时接近h2、h3最大带宽,同时向h2、h3发消息时带宽大约减半

分析:

对于学习算法,h1向h2发消息时,交换机不向h3发消息,并且h1带宽为h2,h3之和,因此h1同时向h2,h3发消息,两边互不干扰,能接近h2、h3最大带宽。

对于广播算法,h1向h2发消息时h3也会收到消息,占用h3接收带宽。h1同时向h2、h3发消息时,h2的接收带宽内,除了收到发给h2的消息,还有发给h3的消息。h3亦然。因而最终结果是h2、h3的最大带宽的大约一半

结论:

学习算法可以防止占用除了源、目的节点以外的节点的带宽,从而在多个节点同时发出多个请求时,提高了有效带宽,因而优于广播算法

6.bug以及原因:

现象1: ping不通

结果: 交换机的学习算法一开始设计的不对。最初版我是在没有查到目的地址时广播并把目的地址插进转发表。正确的做法应该是在没有查到源地址时把源地址插入到转发表,在没有查到目的地址时广播(不是把目的地址插入转发表)

现象2: 交换机学习算法在iperf时带宽只有3.几M

结果: 屏幕打印的信息太多。去掉打印信息就达到9.几M了