SMT-based Test Program Generation for Cache Memory

Evgeny Kornikhin Moscow State University

Testing by Test Programs

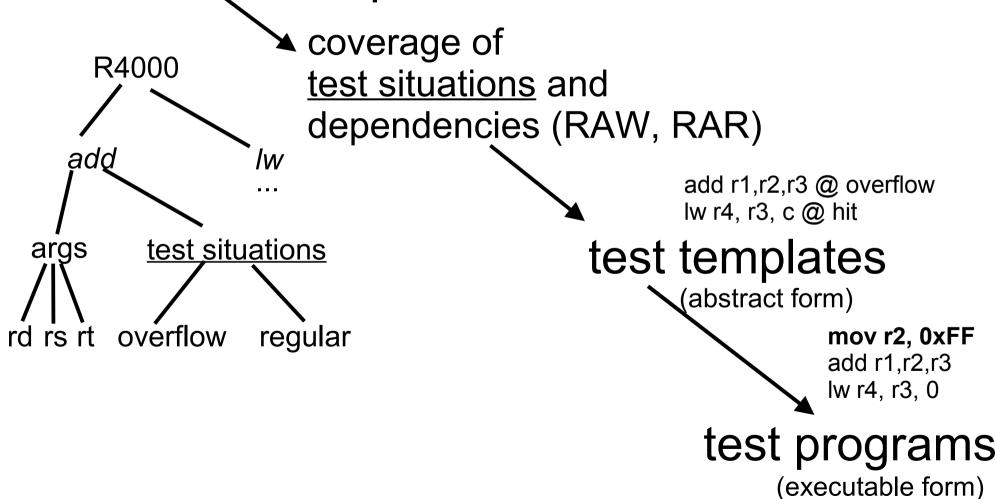
add r1,r2,r3 sub r4, r1, r2 lw r5, r1, 0 lui r2, r1, r4

assembler program microprocessor (test program)

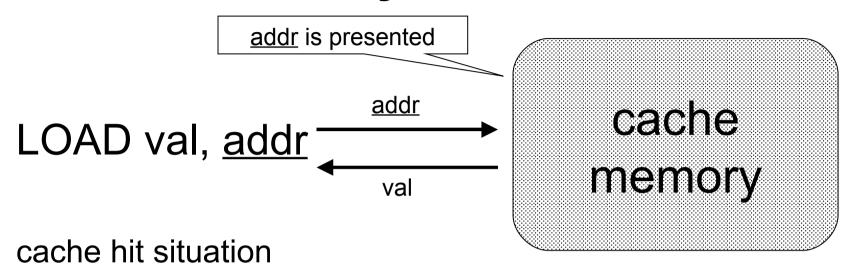
-

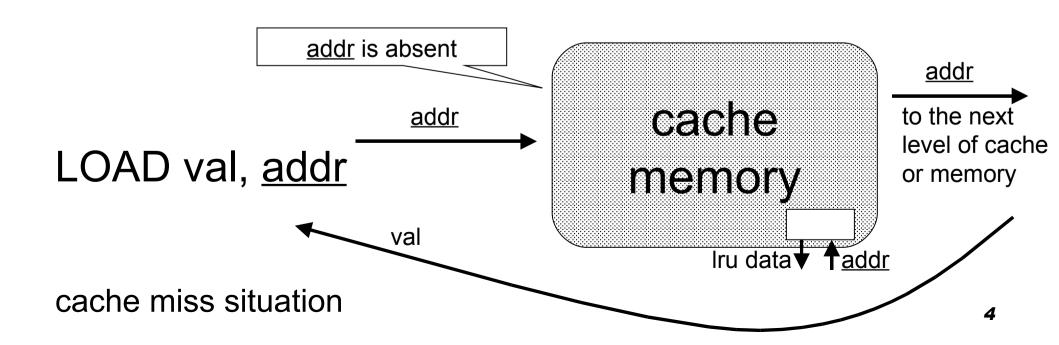
Test Programs Generation

model of microprocessor



Cache memory behavior







LOAD x, y @ hit STORE u, z @ miss LOAD z, y @ hit cache-hit with "y" cache-hit with "z" cache-hit with "y"

initial state of cache
(addresses of cached data)
and registers = ?

Huge exhaustive search

LOAD x, y @ hit STORE u, z @ miss LOAD z, y @ hit

select values

↓ ↑
run template

$$X \mid 0 ... 2^{64}-1$$

$$z \mid 0...2^{64}-1$$

$$u \mid 0 ... 2^{64}-1$$

initial registers' values

$$0...2^{64}$$
-1

$$0...2^{64}$$
-1

$$0...2^{64}-1$$

$$0...2^{64}-1$$

$$0...2^{64}$$
-1

$$0...2^{64}$$
-1

$$0...2^{64}$$
-1

$$0...2^{64}-1$$

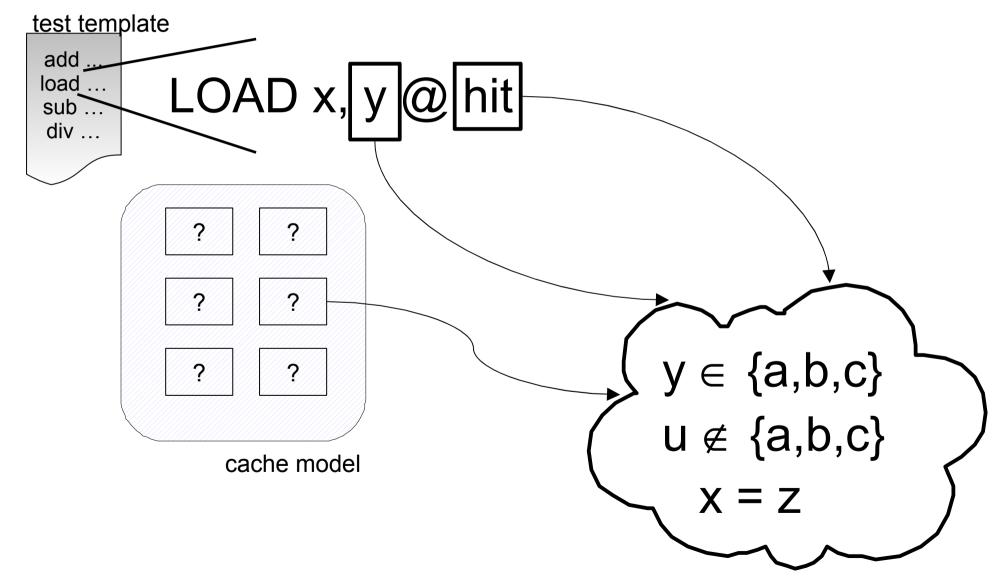
initial cache state

Hard behavior constraints

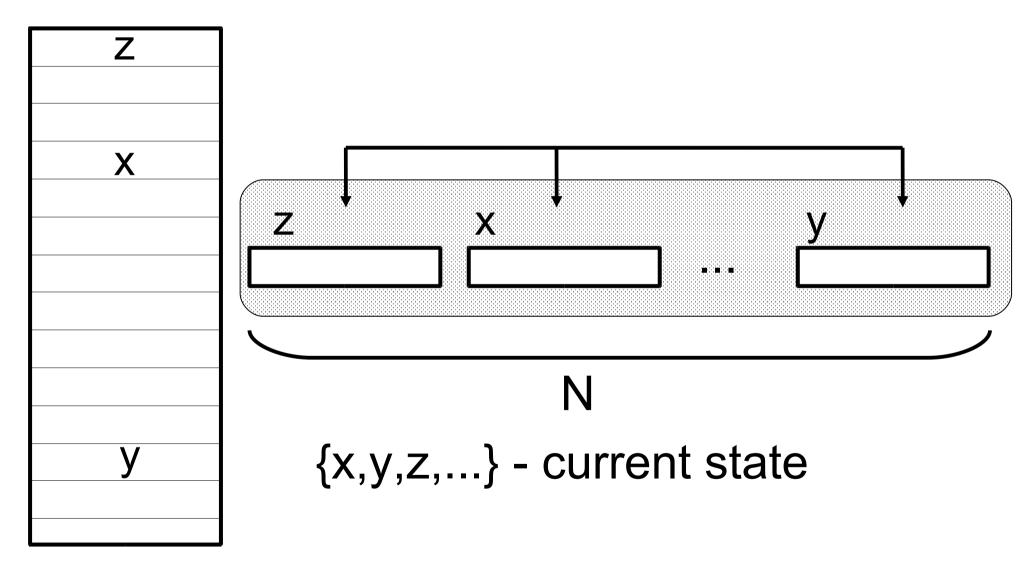
```
LOAD x, y @ hit \begin{cases} index(L,y) > 0 \\ L' == inc\_counter(L,y) \end{cases}
 \begin{array}{l} \text{STORE u,z @ miss} \rightarrow \text{z0} \\ \text{(cache-miss with "z")} \end{array} \begin{cases} \text{index(L',z) == 0} \\ \text{counter(L',z0)} \rightarrow \text{min} \\ \text{L" == remove(L',z0)} \\ \text{L"' == add(L",z)} \end{cases} 
                LOAD z, y @ hit \{ \text{index}(L''',y) > 0 \} (cache-hit with "y")
```

$$x, y, z, u, L = ?$$

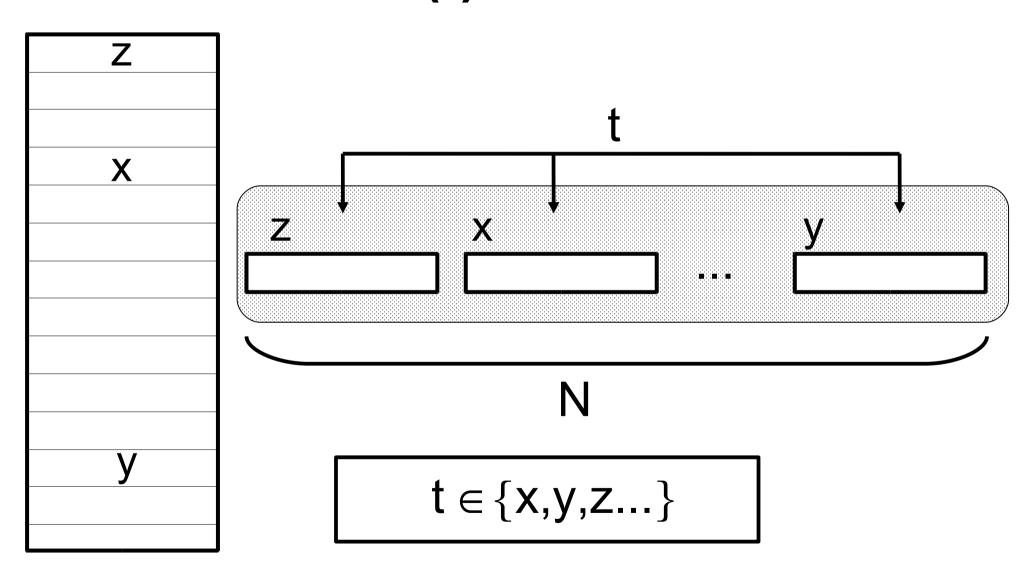
Key Idea



Fully Associative Cache



Cache-hit hit(t)



Cache-miss miss(t)

X $t \notin \{x,y,z...\}$

new cache $\{x,y,z...\} \cup \{t\} \setminus \{?\}$

Cache-miss miss(t)→u

X $t \notin \{x,y,z...\} | | u \in \{x,y,z...\} | | Iru(u) |$ new cache $\{x,y,z...\} \cup \{t\} \setminus \{u\}$

Iru(u)

hit x1

hit x2

miss x3 \rightarrow x4

hit x5 u = x2 $\{x3, x5\} = L\setminus\{u\}$

Iru(u)

hit x1
hit x2
hit x2
miss x3
$$\rightarrow$$
x4
$$\begin{cases} u = x1 \\ \{x2, x3, x5\} = L \setminus \{u\} \end{cases}$$
hit x5
miss t \rightarrow u

there are another cases...

Example

initial state:

LOAD x, y @ hit

STORE u, z @ miss →z0

LOAD z, y @ hit

$$N = 3$$

$$y \in \{\alpha, \beta, \gamma\}$$

$$z\notin\{\alpha,\beta,\gamma\}$$

$$z0 \in \{\alpha, \beta, \gamma\}$$

$$z0=\beta$$

$$\{\alpha,\beta,\gamma\}\setminus\{z0\}=\{\gamma,y\}$$

$$y \in \{\alpha, \beta, \gamma, z\} \setminus \{z0\}$$

Example

$$y = \alpha$$

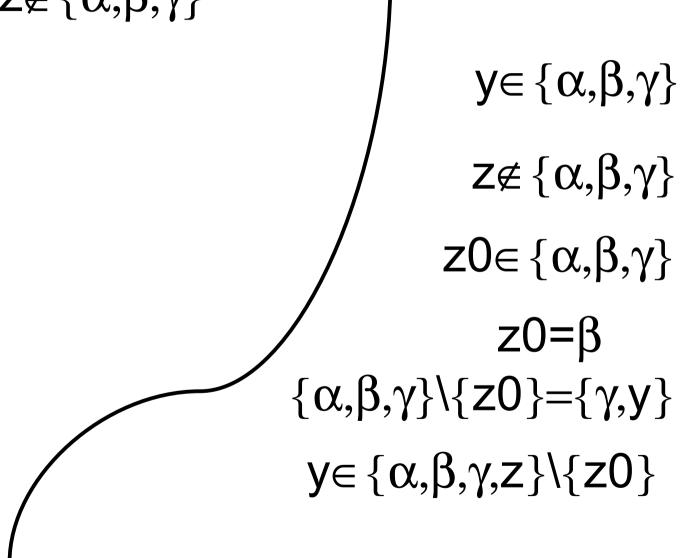
$$z\notin\{\alpha,\beta,\gamma\}$$

$$y=\alpha=0$$

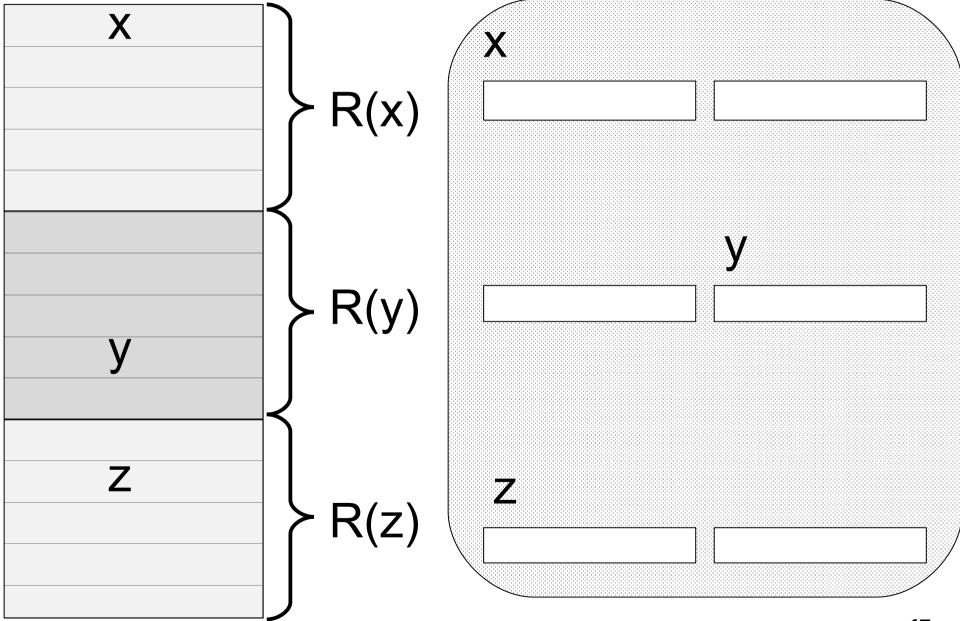
$$\beta=1$$

$$\gamma = 2$$

$$z=3$$



Common Cache



Common Cache

hit(t)
$$t \in L$$

```
u \in L

t \notin L

L \cup \{t\} \setminus \{u\} is the next cache

R(t) = R(u)

Iru(u)
```

Iru(u)

```
hit x1

hit x2

miss x3\rightarrowx4

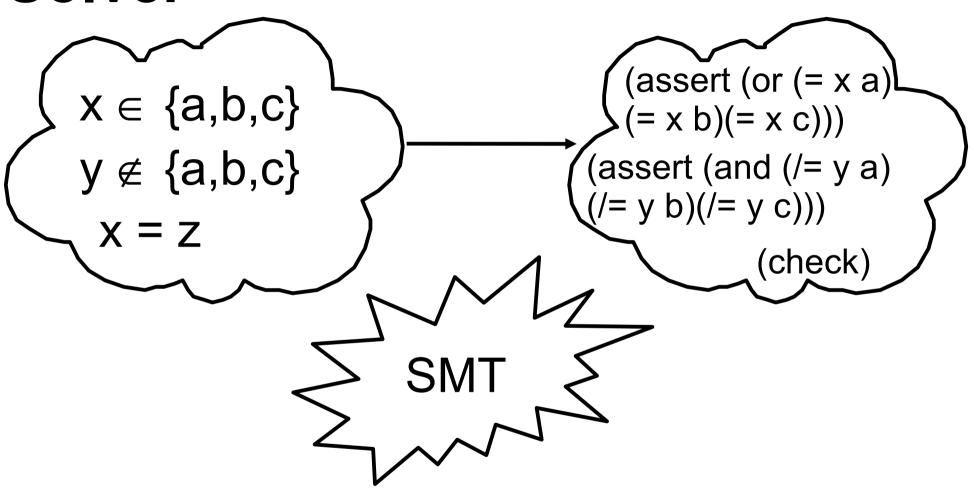
hit x5

miss t\rightarrowu
\begin{cases} u = x2\\ \{x3, x5\} \cap R(u)\\ = (L\setminus\{y\}) \cap R(u) \end{cases}
```

Example

```
x1,x2 \in \{a1,a2,b1,b2,c1,c2\}
x3 \notin \{a1,a2,b1,b2,c1,c2\}
R(x3) = R(y3)
x4 \in \{a1,...,c2,x3\}\setminus\{y3\}
x5 \notin \{a1,...,c2,x3\}\setminus\{y3\}
v3 = c2
{y3} = ({a1,...,c2} \setminus {x1,x2, y3}) \cap R(y3)
y5 = x2
\{y5\} = (\{a1...c2,x3\}\setminus\{y3,y5,x3,x4\})\cap R(y5)
```

Solver



SAT modulo theories
Yices



 Cache behavior can be expressed using equations on addresses' set

Equations can be solved by SMT-solver

Contacts

http://tesla-project.googlecode.com

http://hardware.ispras.ru

kornevgen@gmail.com

Thank you for attention!