

Problem 1.a. Let expression $E: e \rightarrow 0 \mid 3 \mid e+e \mid e*e$, and a semantic function $\llbracket \cdot \rrbracket \in E \rightarrow \mathbb{Z}$. Then, $\llbracket 0 \rrbracket = 0$, $\llbracket 3 \rrbracket = 3$, $\llbracket e_1 + e_2 \rrbracket = \llbracket e_1 \rrbracket + \llbracket e_2 \rrbracket$, $\llbracket e_1 * e_2 \rrbracket = \llbracket e_1 \rrbracket \times \llbracket e_2 \rrbracket$.

Problem 1.b. Theorem: $\forall e \in E, \llbracket e \rrbracket = 3k$ where $k \in \mathbb{Z}$.

(1) Use structural induction on e .

(2) Base Case:

$$(2.1) e \rightarrow 0: \llbracket e \rrbracket = \llbracket 0 \rrbracket = 0 = 3 \cdot 0.$$

$$(2.2) e \rightarrow 3: \llbracket e \rrbracket = \llbracket 3 \rrbracket = 3 = 3 \cdot 1.$$

(3) $e \rightarrow e_1 + e_2$:

$$(3.1) \llbracket e \rrbracket = \llbracket e_1 + e_2 \rrbracket = \llbracket e_1 \rrbracket + \llbracket e_2 \rrbracket.$$

$$(3.2) \text{ By induction hypothesis, } \llbracket e_1 \rrbracket = 3k_1, \llbracket e_2 \rrbracket = 3k_2 \text{ where } k_1, k_2 \in \mathbb{Z}.$$

$$(3.3) \therefore \llbracket e \rrbracket = \llbracket e_1 \rrbracket + \llbracket e_2 \rrbracket = 3k_1 + 3k_2 = 3 \cdot (k_1 + k_2), \text{ and } (k_1 + k_2) \in \mathbb{Z}.$$

(4) $e \rightarrow e_1 * e_2$:

$$(4.1) \llbracket e \rrbracket = \llbracket e_1 * e_2 \rrbracket = \llbracket e_1 \rrbracket \times \llbracket e_2 \rrbracket.$$

$$(4.2) \text{ By induction hypothesis, } \llbracket e_1 \rrbracket = 3k_1, \llbracket e_2 \rrbracket = 3k_2 \text{ where } k_1, k_2 \in \mathbb{Z}.$$

$$(4.3) \therefore \llbracket e \rrbracket = \llbracket e_1 \rrbracket \times \llbracket e_2 \rrbracket = 3k_1 \times 3k_2 = 3 \cdot (3k_1k_2), \text{ and } 3k_1k_2 \in \mathbb{Z}.$$

Problem 1.c.
$$\frac{}{\langle 0 \rangle \rightarrow 0} \quad \frac{}{\langle 3 \rangle \rightarrow 3} \quad \frac{\langle e_1 \rangle \rightarrow v_1, \langle e_2 \rangle \rightarrow v_2}{\langle e_1 + e_2 \rangle \rightarrow v_1 + v_2} \quad \frac{\langle e_1 \rangle \rightarrow v_1, \langle e_2 \rangle \rightarrow v_2}{\langle e_1 * e_2 \rangle \rightarrow v_1 \times v_2}$$

Problem 2.a. Let $d_i = i$. Obviously, $\{d_i\}$ is a chain.

$$(1) f(\bigsqcup_i d_i) = f(1 \sqcup 2 \sqcup \dots) = f(\infty) = \top$$

$$(2) \bigsqcup_i f(d_i) = f(1) \sqcup f(2) \sqcup \dots \sqcup f(n) \sqcup \dots = \perp \sqcup \perp \sqcup \dots \sqcup \perp \sqcup \dots = \perp$$

$$(3) \therefore f(\bigsqcup_i d_i) \neq \bigsqcup_i f(d_i), \text{ hence } f \text{ is not continuous.}$$

Problem 2.b. f_k is continuous.

(1) $x \sqsubseteq x' \Rightarrow f(x) \sqsubseteq f(x')$ (i.e f is monotonic)

$$(1.1) x' \leq k \Rightarrow f(x) = f(x') = \perp \quad (\because x \leq x')$$

$$(1.2) x \leq k \wedge k < x' \Rightarrow f(x) = \perp \sqsubseteq \top = f(x')$$

$$(1.3) k < x \Rightarrow f(x) = f(x') = \top$$

(2) $\forall \{d_i\}, f(\bigsqcup_i d_i) = \bigsqcup_i f(d_i)$.

(2.1) Assume $k \neq \infty$. (Otherwise, $\forall x f_\infty(x) = \perp$ and obviously satisfies (2))

(2.2) Let $x \in \{d_i\}$ and assume that $x > k$.

$$(2.3) \text{ If } x \text{ exists, } \bigsqcup_i d_i > k \text{ so } f_k(\bigsqcup_i d_i) = \top, \bigsqcup_i f_k(d_i) = \dots \sqcup f_k(x) \sqcup \dots = \dots \sqcup \top \sqcup \dots = \top.$$

$$(2.4) \text{ Otherwise, } \bigsqcup_i d_i \leq k \text{ so } f_k(\bigsqcup_i d_i) = \perp, \bigsqcup_i f_k(d_i) = \perp \sqcup \perp \sqcup \dots \sqcup \perp = \perp.$$

Problem 2.c. f_0, f_1, \dots constitute a chain, and f_∞ is the least element.

(1) $n \leq m \Rightarrow f_m \sqsubseteq f_n$

$$(1.1) x \leq n \Rightarrow f_m(x) = f_n(x) = \perp$$

$$(1.2) n < x \leq m \Rightarrow f_m(x) = \perp \sqsubseteq \top = f_n(x)$$

$$(1.3) m < x \Rightarrow f_m(x) = f_n(x) = \top$$

$$(1.4) \therefore \forall x \in P f_m(x) \sqsubseteq f_n(x)$$

(2) By (1), $\forall n \in P f_\infty \sqsubseteq f_n$. So, f_∞ is the least element.

Problem 3. Use the fixed point theorem.

$$(0) \text{ Let } f^\perp(x) = \begin{cases} \perp & \text{if } x = \perp \\ f(x) & \text{otherwise} \end{cases} \text{ and } \llbracket \cdot \rrbracket : \mathbb{Z}_\perp \rightarrow \mathbb{Z}_\perp.$$

(1) In plain English, f is a function returns 1 if $z \geq 0$, otherwise returns $-z + 1$.

(2) $\llbracket f \rrbracket = fix(F)$ where $F = (\lambda f. \lambda z. \text{if } z \geq 0 \text{ then } 1 \text{ else } (1 + f^\perp(z + 1)))$.

(3) $\llbracket f \rrbracket = fix(F) = \bigsqcup \{F^0(\perp_f), F^1(\perp_f), \dots\}$ where $\perp_f = \lambda z. \perp$.

$$(3.1) F^0(\perp_f) = \lambda z. \perp.$$

$$(3.2) F^1(\perp_f) = \lambda z. \text{if } z \geq 0 \text{ then } 1 \text{ else } (1 + (\lambda z. \perp)^\perp(z + 1))) = \begin{cases} 1 & \text{if } z \geq 0 \\ \perp & \text{otherwise} \end{cases}.$$

$$(3.3) F^2(\perp_f) = \lambda z. \text{if } z \geq 0 \text{ then } 1 \text{ else } (1 + (F^1(\perp_f))^\perp(z + 1))) = \begin{cases} 1 & \text{if } z \geq 0 \\ 2 & \text{if } z = -1 \\ \perp & \text{otherwise} \end{cases}.$$

$$(3.4) \therefore F^n(\perp_f) = \lambda z. \text{if } z \geq 0 \text{ then } 1 \text{ else } (1 + (F^{n-1}(\perp_f))^\perp(z + 1))) = \begin{cases} 1 & \text{if } z \geq 0 \\ -z + 1 & \text{if } 0 > z \geq -n + 1 \\ \perp & \text{otherwise} \end{cases}.$$

$$(4) \therefore \llbracket f \rrbracket = \begin{cases} 1 & \text{if } z \geq 0 \\ -z + 1 & \text{otherwise} \end{cases} \text{ and matches to the plain meaning that is described in (1).}$$