

Please use the following QR code to check in and record your attendance.

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CS 1027

Fundamentals of Computer
Science II

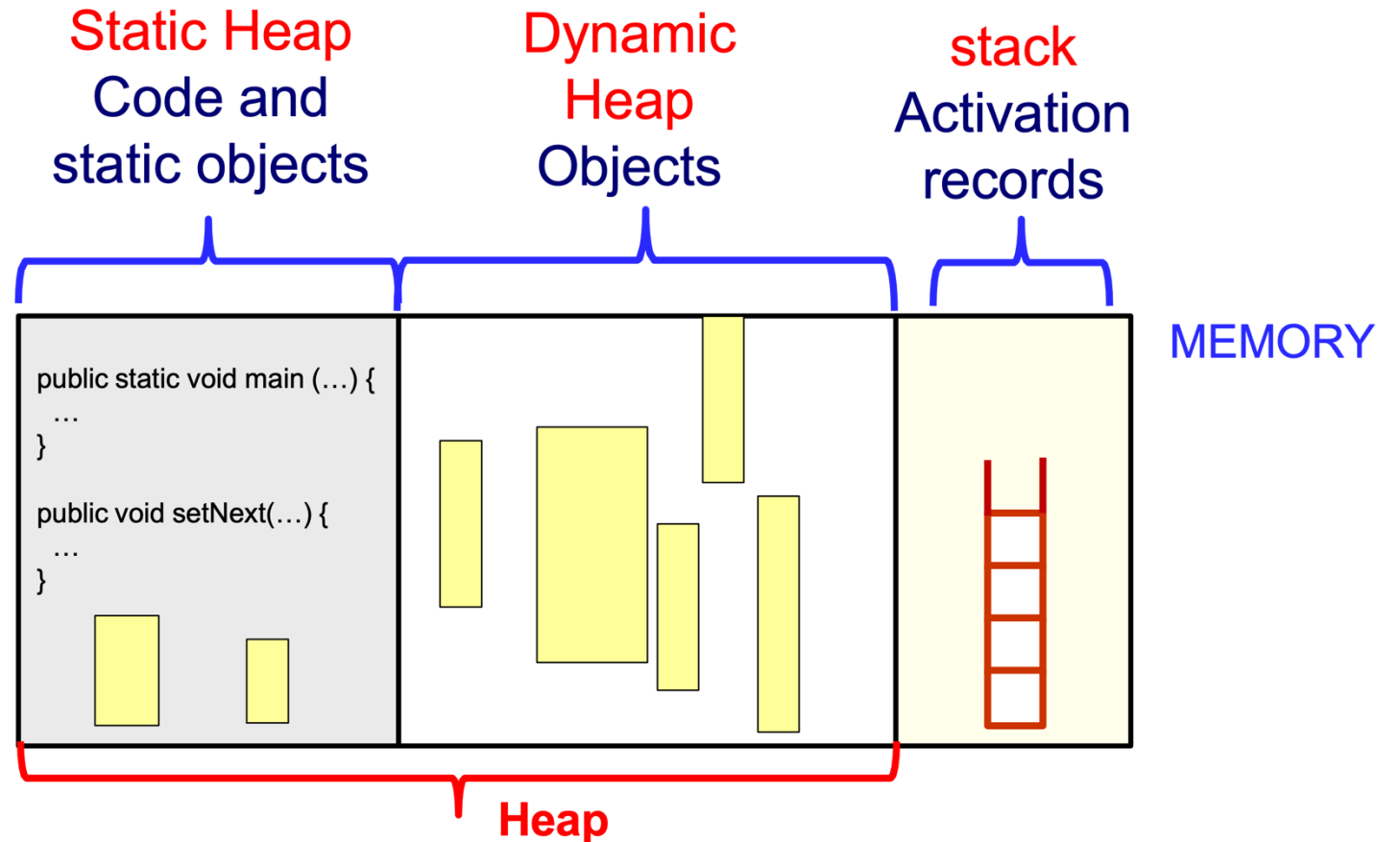
Memory Allocation in **Java** (cont.)

Ahmed Ibrahim



Memory Allocation in Java

- When a program is being executed, separate areas of memory are allocated:

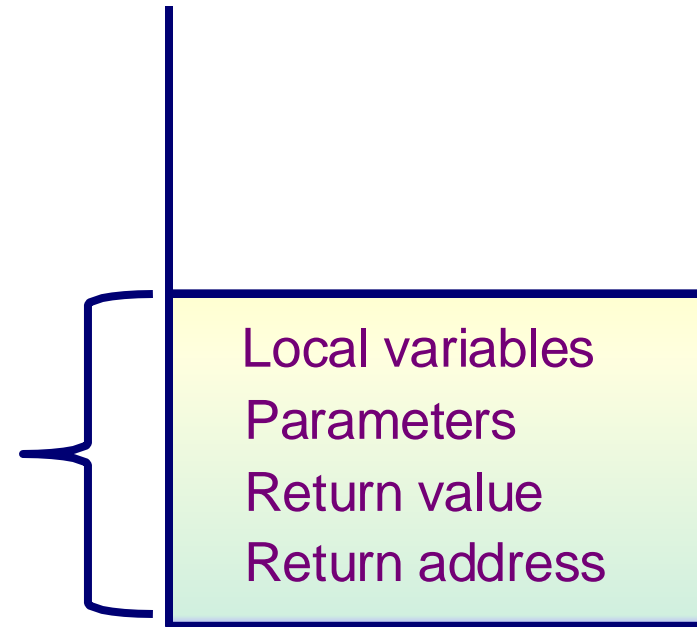


Recall: Execution Stack

- The **execution stack** (also called **runtime stack** or **call stack**) is used to store information needed while a method is being executed, like

- Local variables
- Formal parameters
- Return value
- Return address

Method information



Execution Stack

Recall: Execution of the Program

- When the **main** method is invoked:
 - An **activation record** for the **main** is created and pushed onto the execution stack
- When the **main** calls the method **m2**:
 - An **activation record for m2** is created and pushed onto the execution stack
- When **m2** calls **m3**:
 - An **activation record for m3** is created and pushed onto the execution stack
- When **m3** terminates, its activation record is popped off and control returns to **m2**

Execution of the Program

- When **m2** next calls **m4**:
 - What happens next?
 - What happens when **m4** terminates?
- What happens when **m2** terminates?
- What happens when the **main** function terminates?
- Its activation record is popped off, and control returns to the operating system

Activation Records

- Example of what is in the activation record for a method with primitive type variables, and non-primitive variables.

```
public static int m2 (int param2)
{int local2 = 1;
 Integer i = new Integer(3);
 m3(4);
 return local2 + param2 + m4(3);}
public static void m3 (int param3)
{int[] arr = new int[param3];}
public static int m4 (int param4) {return param4 * 2;}
public static void main (String[] args) {int local1 = m2(5);}
```

To execute this program, it is first compiled and translated to Java bytecode and stored in the static heap.

```
01001100011101001011010111
01001010110101010001111011
1010100101010100011100101
11101011101011101001110000
11010101010000100011100011
10111010011000111011000001
11010100100010101000101011
```

Static heap

Activation Records

- Example of what is in the activation record for a method with primitive type variables, and non-primitive variables.

```
public static int m2 (int param2)
```

```
{int local2 = 1;
```

```
Integer i = new Integer(3);
```

addr2 m3(4);

addr3 return local2 + param2 + m4(3);}

```
public static void m3 (int param3)
```

```
{int[] arr = new int[param3];}
```

```
public static int m4 (int param4) {return param4 * 2;}
```

```
public static void main (String[] args) {
```

addr1 int local1 = m2(5);}

Each instruction is assigned an address. We indicated the addresses of method invocations in the program.

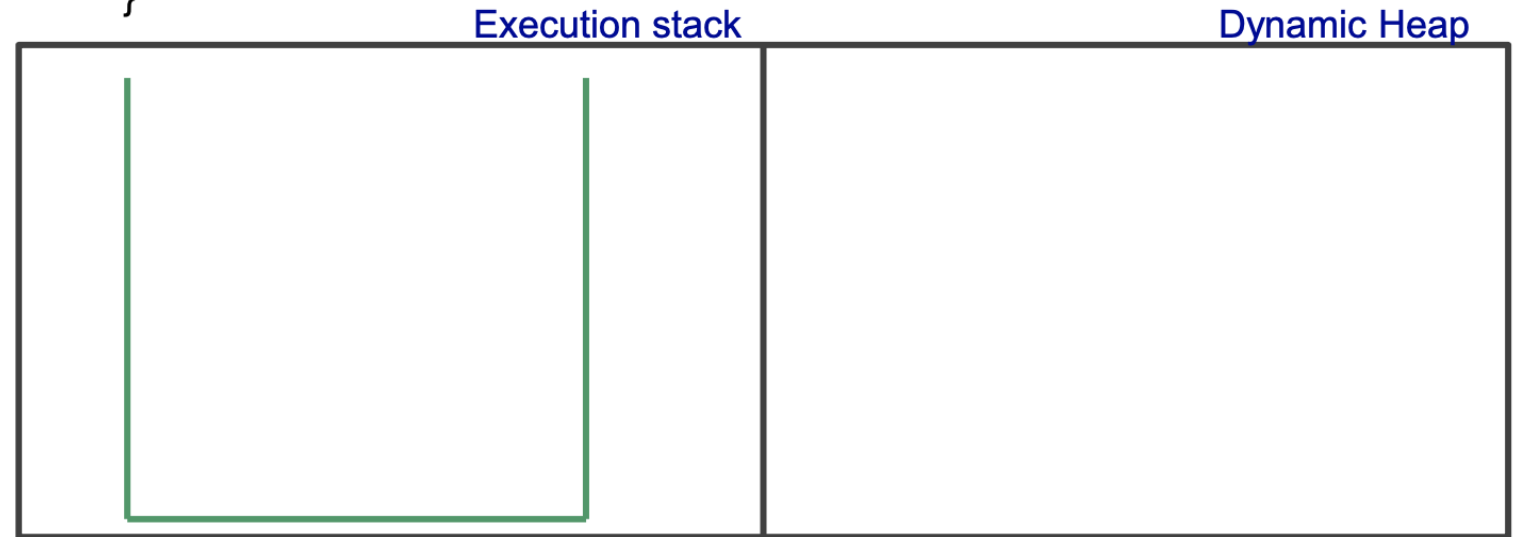
01001100011101001011010111
01001010110101010001111011
1010100101010100011100101
11101011101011101001110000
11010101010000100011100011
10111010011000111011000001
11010100100010101000101011

Static heap

Activation Records

```
public static int m2 (int param2) {  
    int local2 = 1;  
    Integer i = new Integer(3);  
    addr2 m3(4);  
    addr3 return local2 + param2 + m4(3);  
}  
  
public static void m3 (int param3) {  
    int[] arr = new int[param3];  
}
```

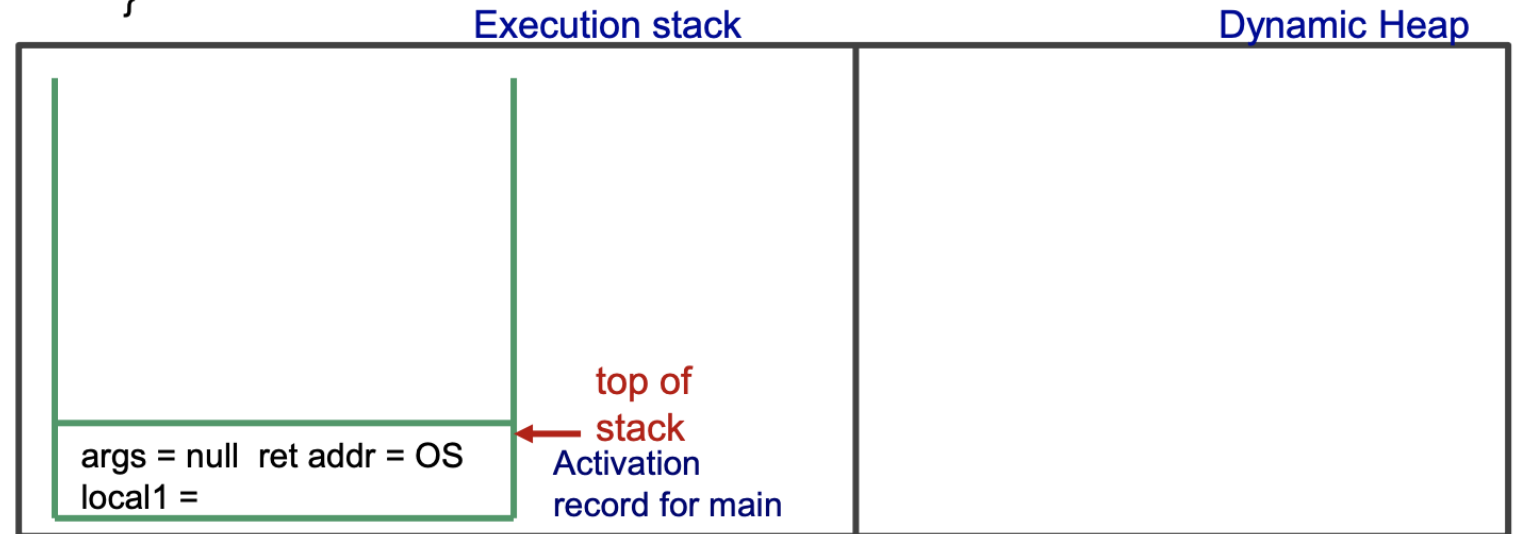
```
public static int m4 (int param4) {  
    return param4 * 2;  
}  
  
public static void main (String[] args) {  
    int local1 = m2(5); addr1  
}
```



Activation Records

```
public static int m2 (int param2) {  
    int local2 = 1;  
    Integer i = new Integer(3);  
    addr2 m3(4);  
    addr3 return local2 + param2 + m4(3);  
}  
  
public static void m3 (int param3) {  
    int[] arr = new int[param3];  
}
```

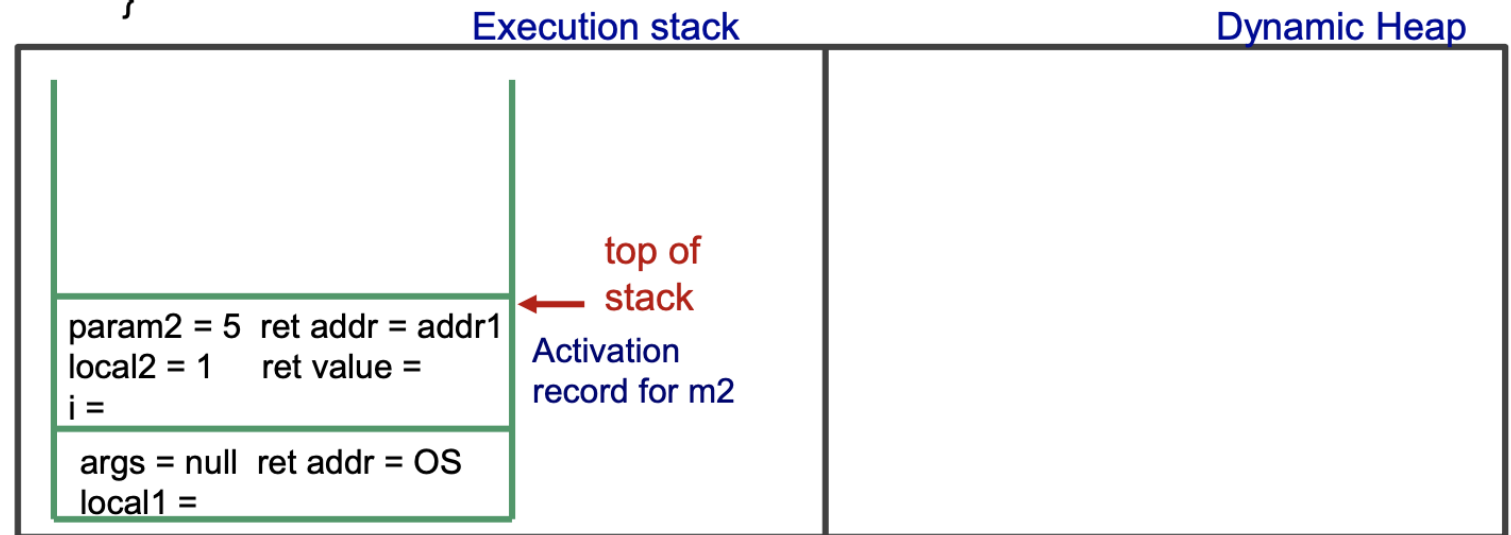
```
public static int m4 (int param4) {  
    return param4 * 2;  
}  
  
public static void main (String[] args) {  
    int local1 = m2(5); addr1  
}
```



Activation Records

```
public static int m2 (int param2) {  
    int local2 = 1;  
    Integer i = new Integer(3);  
    addr2 m3(4);  
    addr3 return local2 + param2 + m4(3);  
}  
  
public static void m3 (int param3) {  
    int[] arr = new int[param3];  
}
```

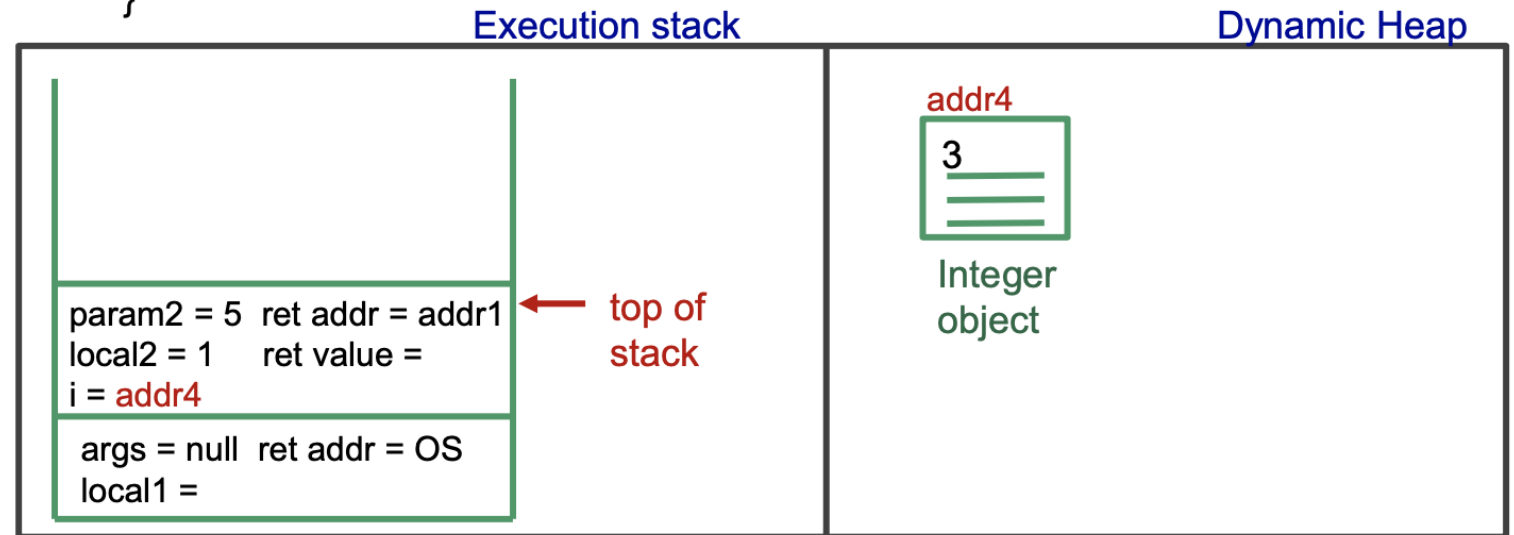
```
public static int m4 (int param4) {  
    return param4 * 2;  
}  
  
public static void main (String[] args) {  
    int local1 = m2(5); addr1  
}
```



Activation Records

```
public static int m2 (int param2) {  
    int local2 = 1;  
    Integer i = new Integer(3);  
addr2 m3(4);  
addr3 return local2 + param2 + m4(3);  
}  
  
public static void m3 (int param3) {  
    int[] arr = new int[param3];  
}  
  
public static int m4 (int param4) {  
    return param4 * 2;  
}  
  
public static void main (String[] args) {  
    int local1 = m2(5); addr1  
}
```

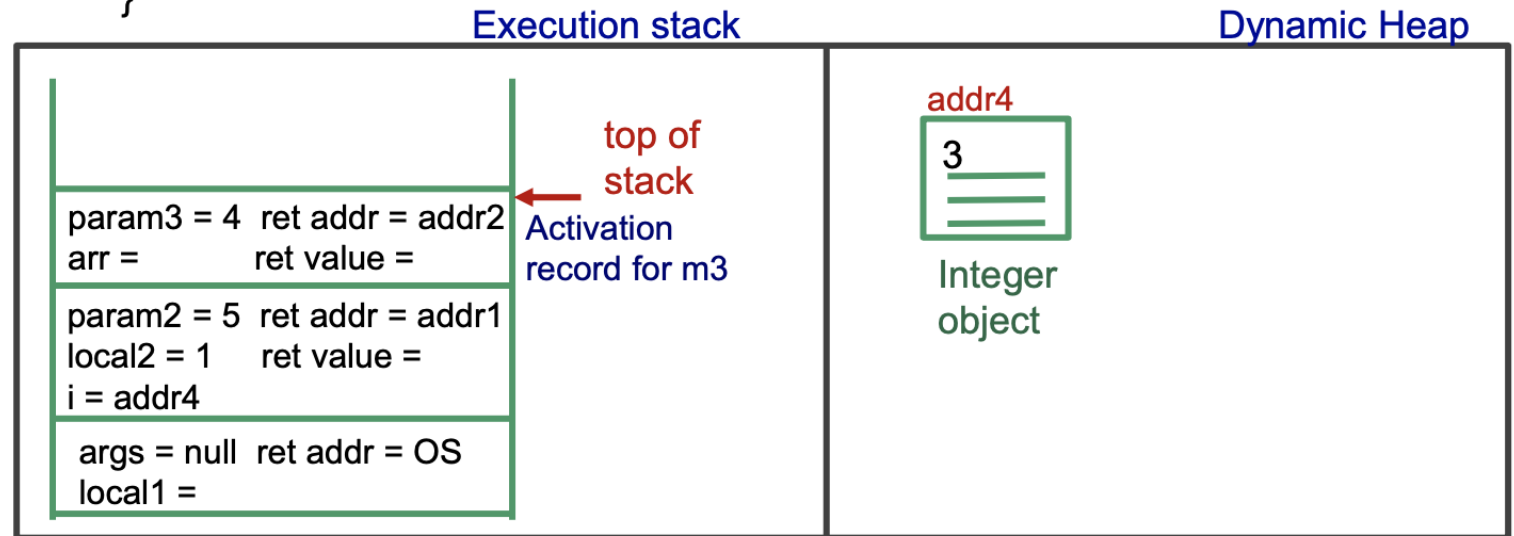
An object is created



Activation Records

```
public static int m2 (int param2) {  
    int local2 = 1;  
    Integer i = new Integer(3);  
    addr2 m3(4);  
    addr3 return local2 + param2 + m4(3);  
}  
  
public static void m3 (int param3) {  
    int[] arr = new int[param3];  
}
```

```
public static int m4 (int param4) {  
    return param4 * 2;  
}  
  
public static void main (String[] args) {  
    int local1 = m2(5); addr1  
}
```



Activation Records

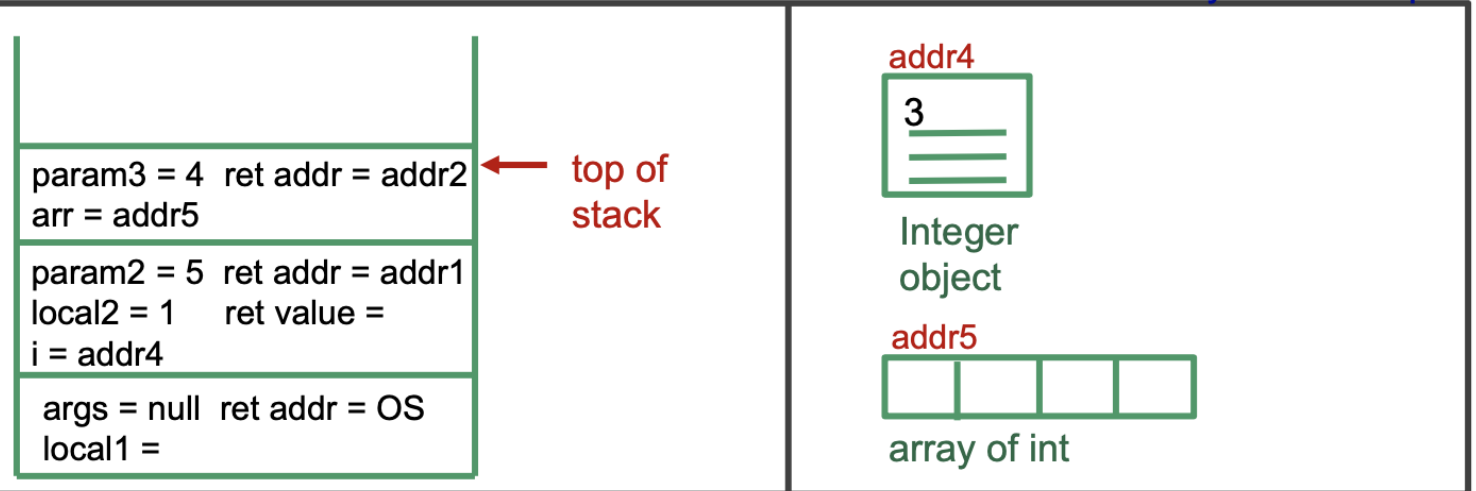
```
public static int m2 (int param2) {  
    int local2 = 1;  
    Integer i = new Integer(3);  
    addr2 m3(4);  
    addr3 return local2 + param2 + m4(3);  
}  
  
public static void m3 (int param3) {  
    int[] arr = new int[param3];  
}
```

```
public static int m4 (int param4) {  
    return param4 * 2;  
}  
  
public static void main (String[] args) {  
    int local1 = m2(5); addr1  
}
```

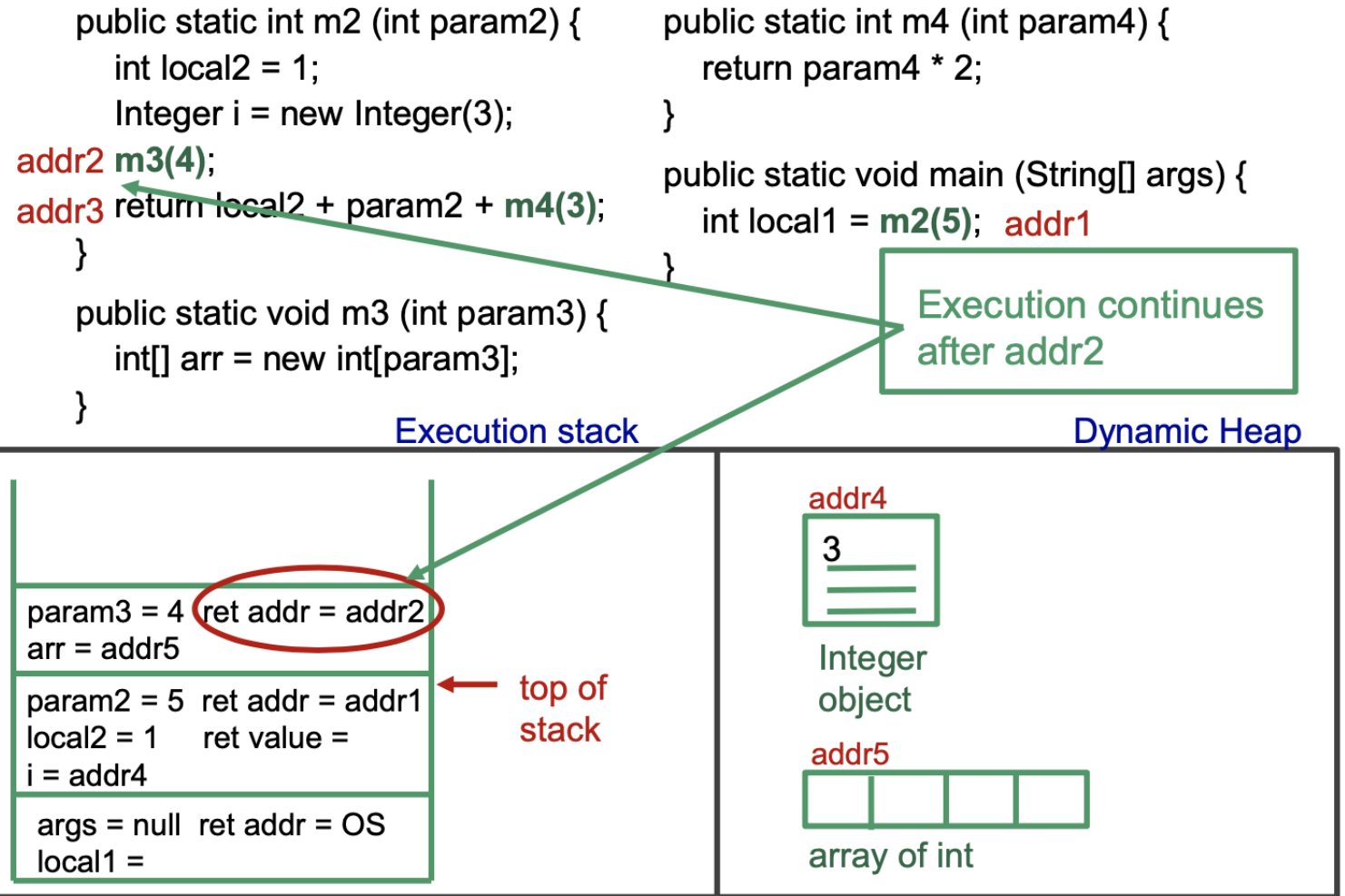
An object is created

Execution stack

Dynamic Heap



Activation Records



Activation Records

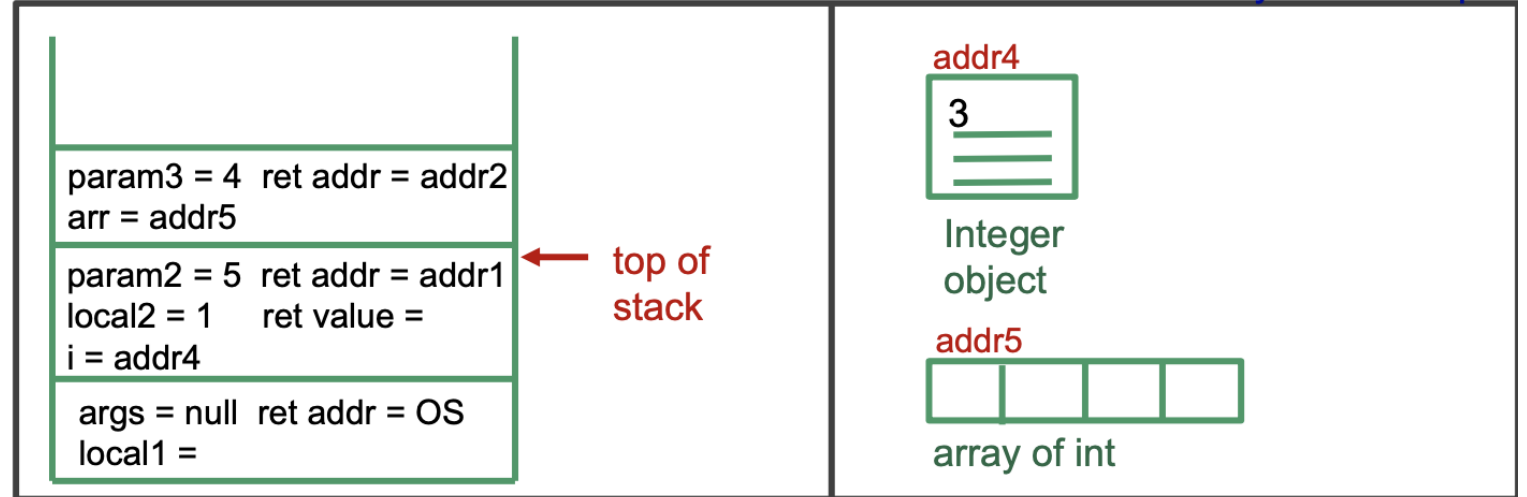
```
public static int m2 (int param2) {  
    int local2 = 1;  
    Integer i = new Integer(3);  
    addr2 m3(4);  
    addr3 return local2 + param2 + m4(3);  
}  
  
public static void m3 (int param3) {  
    int[] arr = new int[param3];  
}
```

```
public static int m4 (int param4) {  
    return param4 * 2;  
}  
  
public static void main (String[] args) {  
    int local1 = m2(5); addr1  
}
```

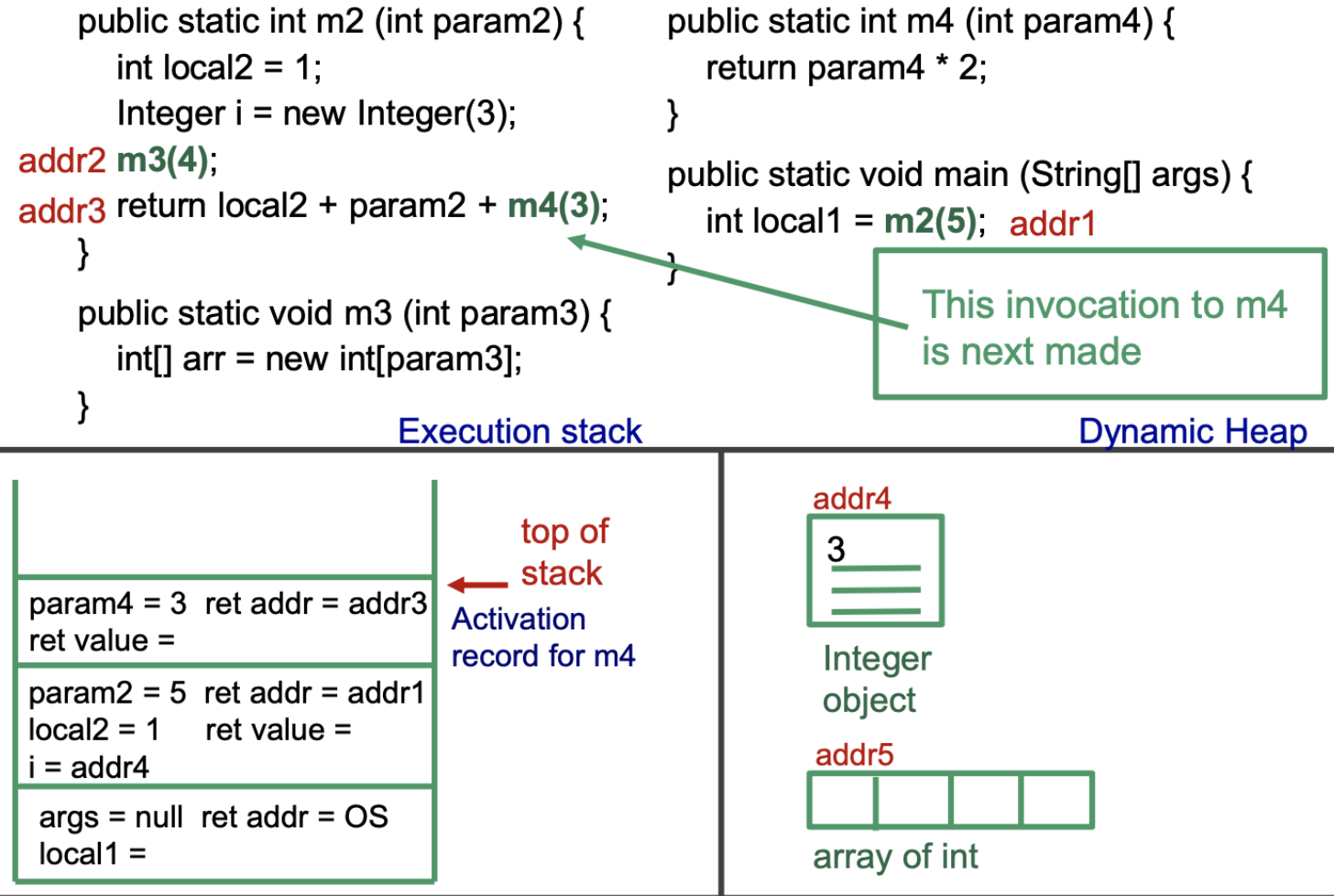
This invocation to m4
is next made

Execution stack

Dynamic Heap



Activation Records

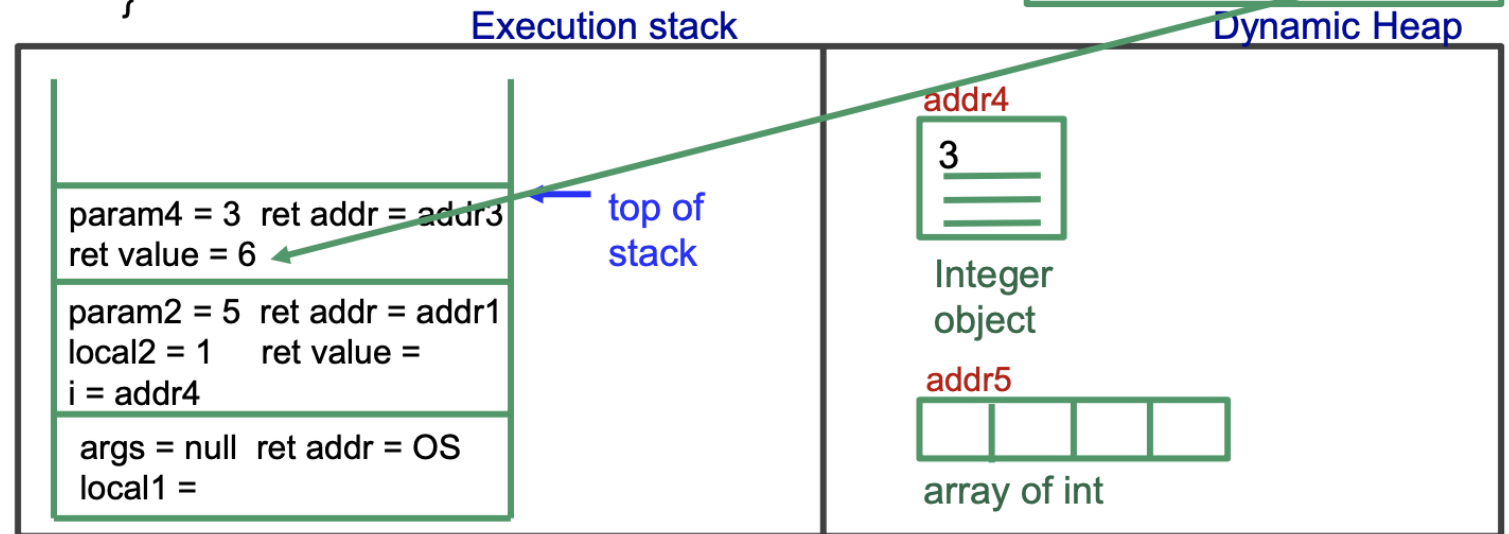


Activation Records

```
public static int m2 (int param2) {  
    int local2 = 1;  
    Integer i = new Integer(3);  
    addr2 m3(4);  
    addr3 return local2 + param2 + m4(3);  
}  
  
public static void m3 (int param3) {  
    int[] arr = new int[param3];  
}
```

```
public static int m4 (int param4) {  
    return param4 * 2;  
}  
  
public static void main (String[] args) {  
    int local1 = m2(5); addr1  
}
```

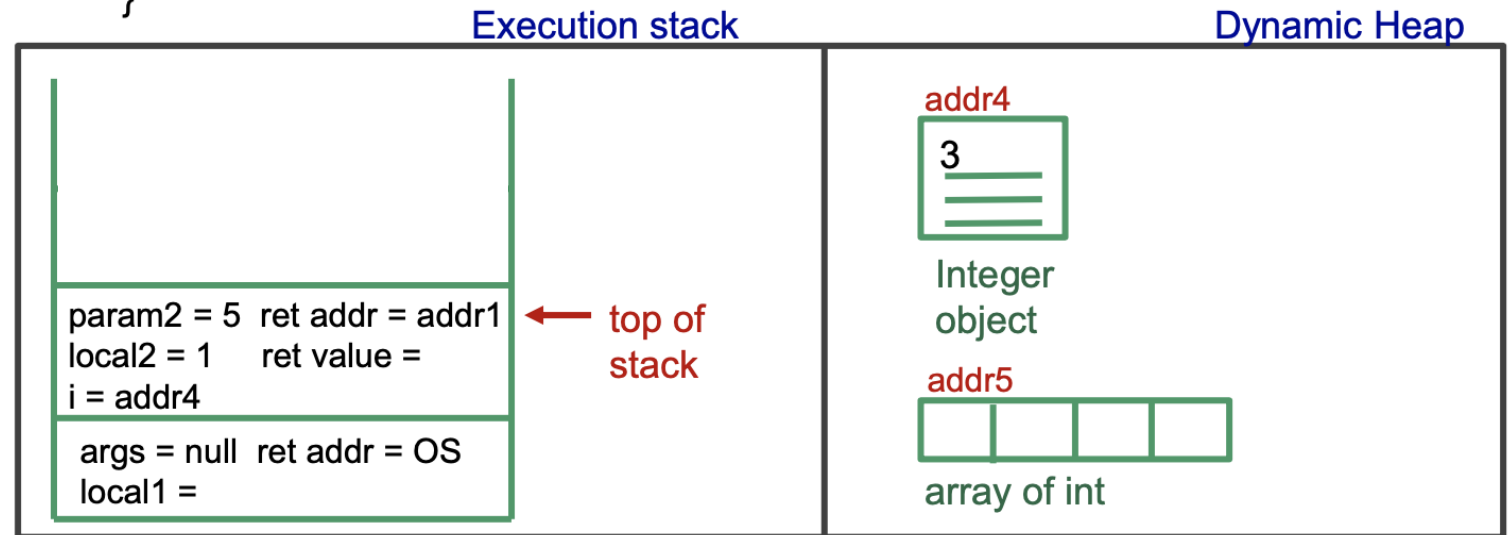
Return value is
computed and stored
in activation record



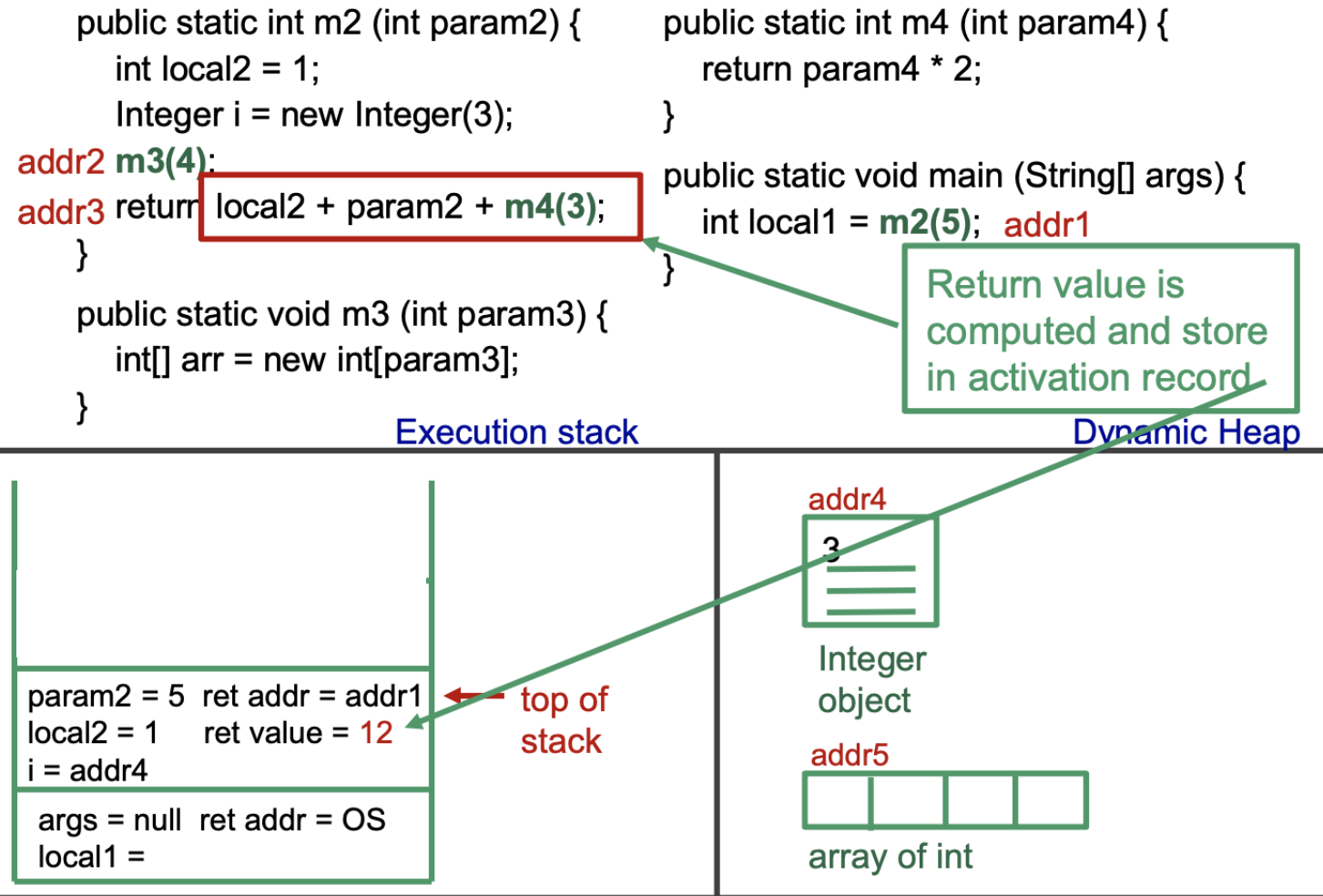
Activation Records

```
public static int m2 (int param2) {  
    int local2 = 1;  
    Integer i = new Integer(3);  
    addr2 m3(4);  
    addr3 return local2 + param2 + m4(3);  
}  
  
public static void m3 (int param3) {  
    int[] arr = new int[param3];  
}  
  
public static int m4 (int param4) {  
    return param4 * 2;  
}  
  
public static void main (String[] args) {  
    int local1 = m2(5); addr1  
}
```

Execution continues
after addr3



Activation Records



Activation Records

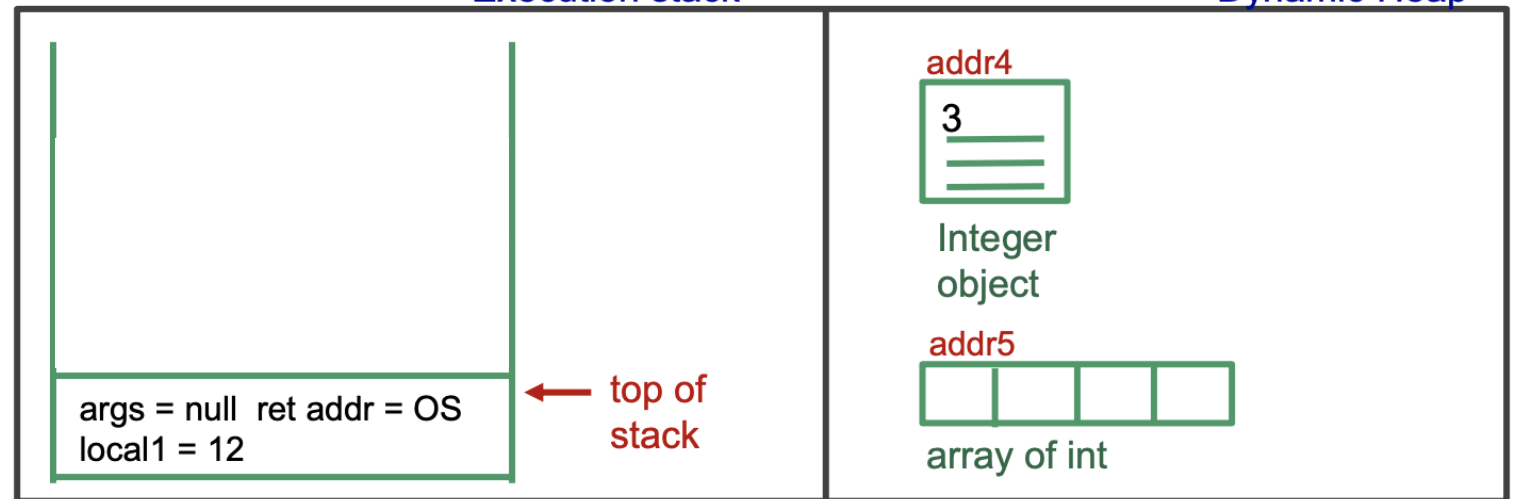
```
public static int m2 (int param2) {  
    int local2 = 1;  
    Integer i = new Integer(3);  
    addr2 m3(4);  
    addr3 return local2 + param2 + m4(3);  
}  
  
public static void m3 (int param3) {  
    int[] arr = new int[param3];  
}
```

```
public static int m4 (int param4) {  
    return param4 * 2;  
}  
  
public static void main (String[] args) {  
    int local1 = m2(5); addr1  
}
```

Execution continues
after addr1

Execution stack

Dynamic Heap

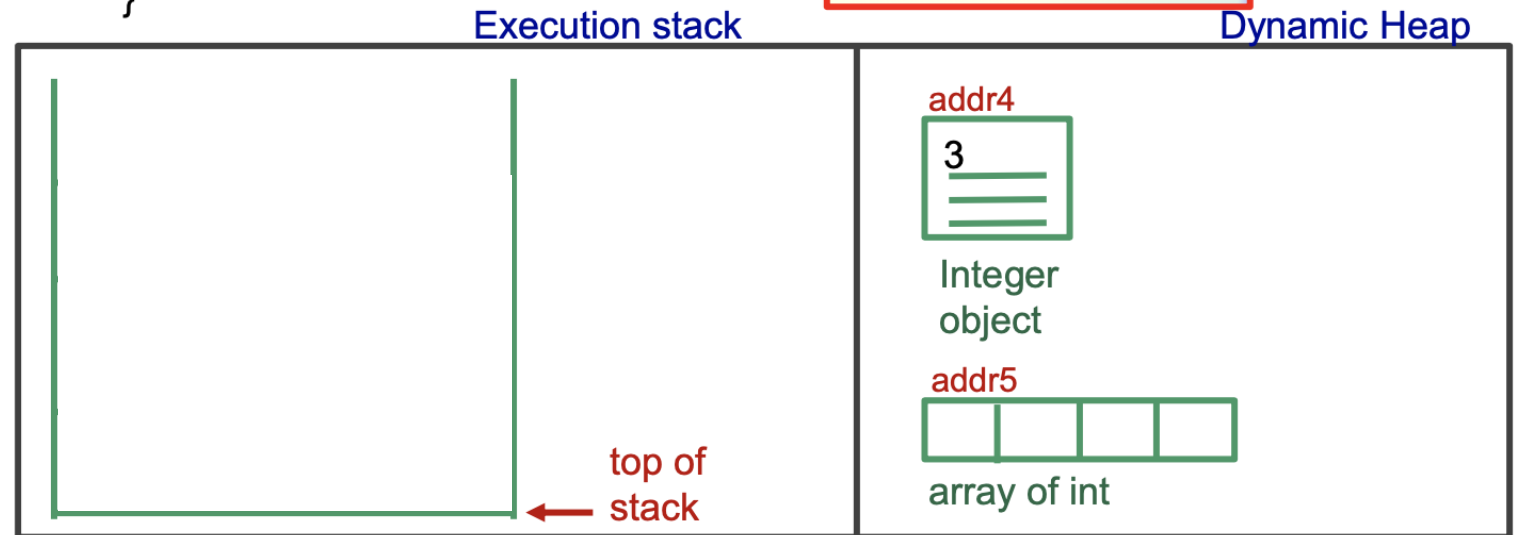


Activation Records

```
public static int m2 (int param2) {  
    int local2 = 1;  
    Integer i = new Integer(3);  
    addr2 m3(4);  
    addr3 return local2 + param2 + m4(3);  
}  
  
public static void m3 (int param3) {  
    int[] arr = new int[param3];  
}
```

```
public static int m4 (int param4) {  
    return param4 * 2;  
}  
  
public static void main (String[] args) {  
    int local1 = m2(5); addr1  
}
```

Program ends and
control goes back to
Operating System



Activation Records – Example 2

```
public class CallFrameDemo2 {  
    private static void printAll (String s1, String  
        s2, String s3) {  
        System.out.println(s1.toString( ));  
        System.out.println(s2.toString( ));  
        System.out.println(s3.toString( ));  
    }  
  
    public static void main (String args[ ]) {  
        String str1, str2, str3;  
        str1 = new String(" string 1 ");  
        str2 = new String(" string 2 ");  
        str3 = new String(" string 3 ");  
        printAll(str1, str2, str3);  
    }  
}
```

Activation Records – Example 2

Draw a picture of the execution stack and of the heap as the above program executes:

1. Activation record for main
2. Activation record for String constructor for str1 – then popped off
3. Activation record for String constructor for str2 – then popped off
4. Activation record for String constructor for str3 – then
5. popped off
6. Activation record for printAll
7. Activation record for toString for str1 – then popped off
8. Activation record for System.out.println – then popped off
9. etc.

Activation Records – Example 2

- What will be stored in the activation record for **main**?
 - Address to return to operating system
 - Variable **args**
 - Variable **str1**
 - Initial value?
 - Value after return from **String constructor**?
 - Variable **str2**
 - Variable **str3**
- What will be in the activation record for **printAll**?

Memory Deallocation

- What happens when a method returns?
 - On the execution stack:
 - The activation record is popped off when the method returns
 - So, that memory is deallocated

Memory Deallocation

- What happens to **objects** on the heap?
 - An object stays in the heap even if no variable is referencing it!
 - So, Java has automatic garbage collection
 - When memory runs low, objects that no longer have a variable referencing them are identified, and their memory is deallocated.

Recursive Definitions

A thick, hand-drawn style orange line that underlines the title "Recursive Definitions".

Recursive Definition

- Defining something in terms of a smaller or simpler version of itself.
- A recursive definition consists of two parts:
 - The *base case*: this defines the *simplest* case or starting point
 - The *recursive part* is the general case that describes all the other cases in terms of smaller versions of itself.



- Example:

```
// Recursive method to calculate factorial
public static int factorial(int n) {
    // Base case: if n is 1, return 1 (factorial of 1 is 1)
    if (n == 1) {return 1;}
    // Recursive part: n * factorial of (n - 1)
    else {return n * factorial(n - 1);}
}
```

base

recursive

Why is a base case needed in any recursive algorithm?

Recursion vs. Iteration

- What is iteration? Repetition, as in a loop
- What is **recursion**? Defining something in terms of a smaller or simpler version of itself (why smaller/simpler?)
- Recursion is a very powerful problem-solving technique.
- Many complex problems would be very difficult to solve without the use of recursion.

Example of Recursive Problem

- Consider the problem of computing the sum of all the numbers between 1 and n :

$$1 + 2 + 3 + 4 + \dots + n-1 + n$$

- Here is a simple iterative algorithm for this problem:

```
Algorithm sum (n)
total = 0
for i = 1 to n do
    total = total + i
return total
```

Example of Recursive Problem

- Consider the problem of computing the sum of all the numbers between 1 and n :

$$1 + 2 + 3 + 4 + \dots + n-1 + n$$

- Recursive definition:

sum of 1 to 1 \Rightarrow 1 (*base case*)

sum of 1 to $n \Rightarrow n +$ the sum of 1 to $n-1$, for $n > 1$

$$\sum_{k=1}^n k = n + \sum_{k=1}^{n-1} k \quad (\text{recursive case})$$

Recursive Algorithm

- Recursive definition:

sum of 1 to 1 \Rightarrow 1 (*base case*)

sum of 1 to n \Rightarrow n + the sum of 1 to n-1, for n > 1

- Recursive algorithm for this problem:

Algorithm *sum*(n)

In: Positive value n

Out: 1 + ... + n

if n = 1 then return 1 // base case

else return n + *sum*(n-1) // recursive case

How Recursion Works

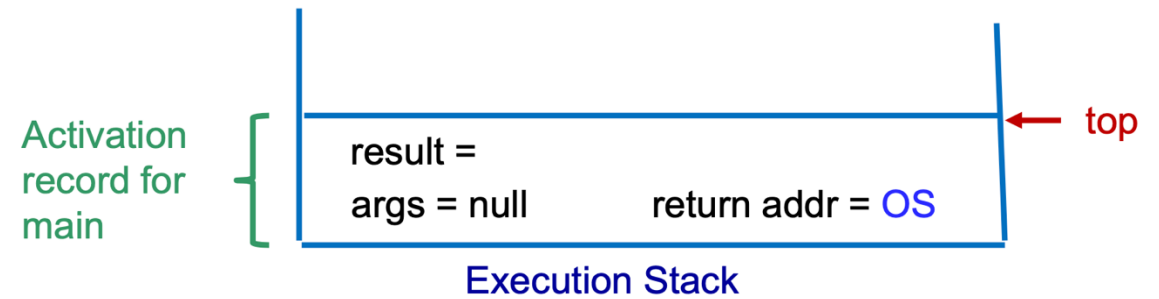
- Consider the following program

```
public static void main (String[] args) {int result = sum(4);} // addr 1
public static int sum (int n) {if (n == 1) return 1;
                               else return n + sum(n-1);} // addr2
```

- An **activation record** is created for the method main when the program is executed. This activation record stores:
 - The return address addr1
 - The variable `result`
 - The parameter `args`

How Recursion Works

```
public static void main (String[] args)
{int result = sum(4);} // addr 1
public static int sum (int n) {
if (n == 1) return 1;
else return n + sum(n-1);} // addr2
```

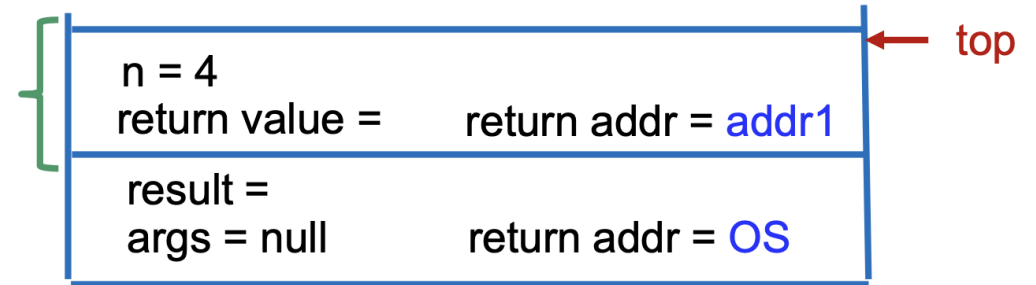


- At this point, the **execution stack** looks like the following figure. We assume no parameter is passed to the **main** function, so **args** is **null**.
- The **result** variable has no value assigned to it yet, so we left its value blank. **OS** denotes the address of the virtual machine's instruction where the main method was invoked.

How Recursion Works

```
public static void main (String[] args)
{int result = sum(4);} // addr 1
public static int sum (int n) {
if (n == 1) return 1;
else return n + sum(n-1);} // addr2
```

Activation
record for
sum

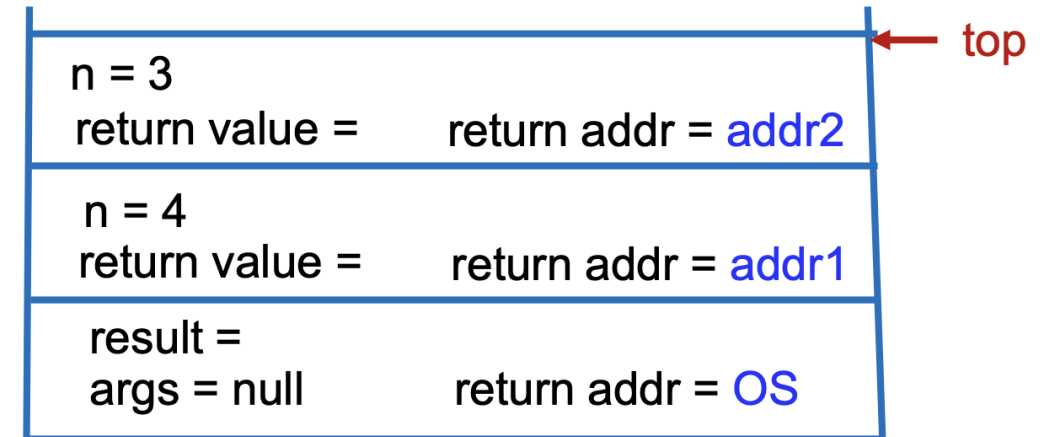


Execution Stack

- Once the activation record for the main function has been created and the values of the parameters and return address have been stored, the execution of the method main starts.
- The first and only statement of the main function invokes method **sum**(4). This creates another **activation record**, which is pushed into the **execution stack**, as shown above.

How Recursion Works (cont.)

```
public static void main (String[] args)
{int result = sum(4);} // addr 1
public static int sum (int n) {
if (n == 1) return 1;
else return n + sum(n-1);} // addr2
```

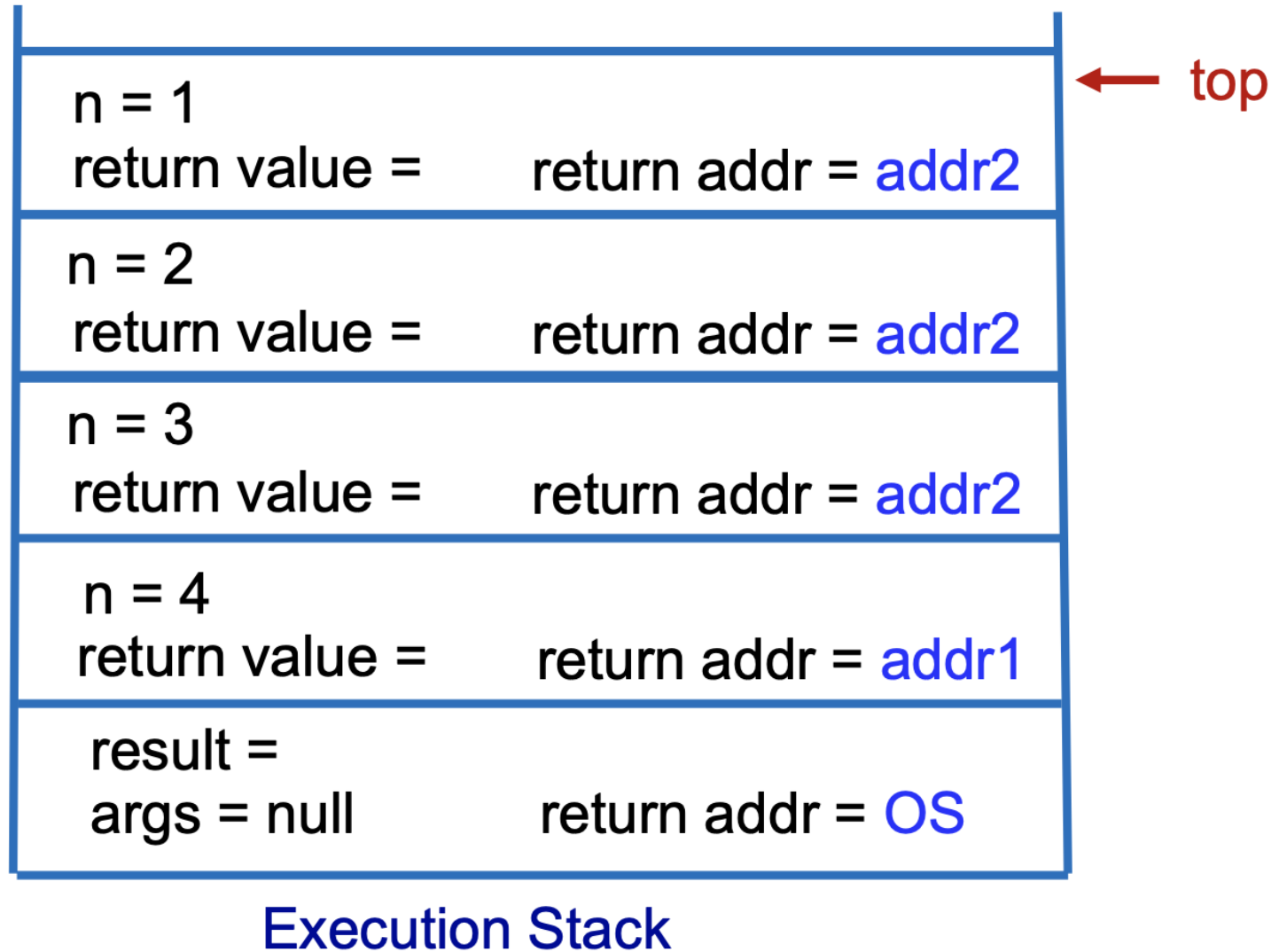


Execution Stack

- Once the **activation record** has been created, the execution of the method **sum** starts.
- Since $n > 1$, method **sum (n-1)** is invoked.
- A new **activation record** is created and pushed into the stack

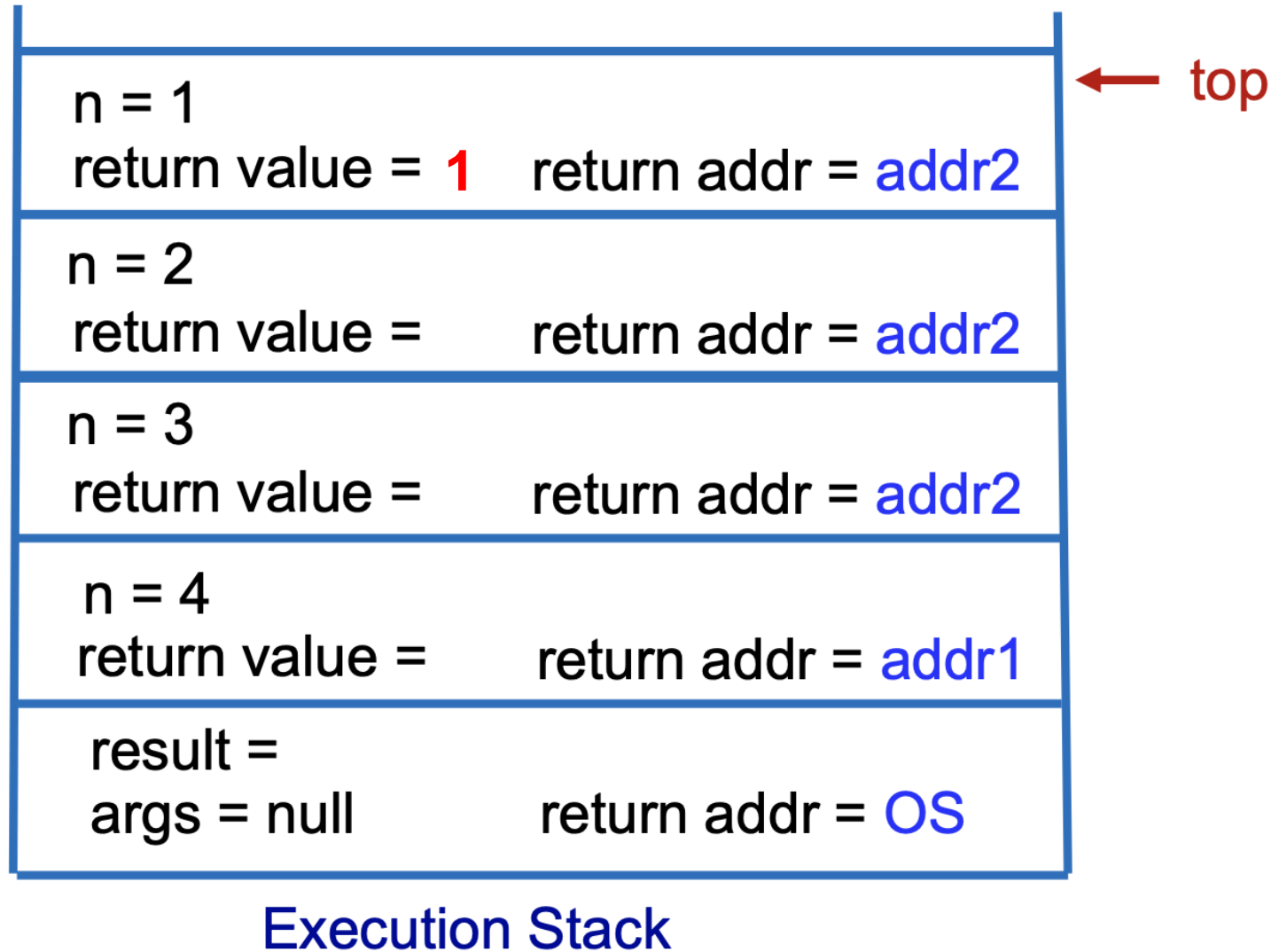
How Recursion Works (cont.)

- Then, two more invocations to the method `sum` with parameters 2 and 1 are made.
- After the last invocation, the execution stack looks like the figure given.



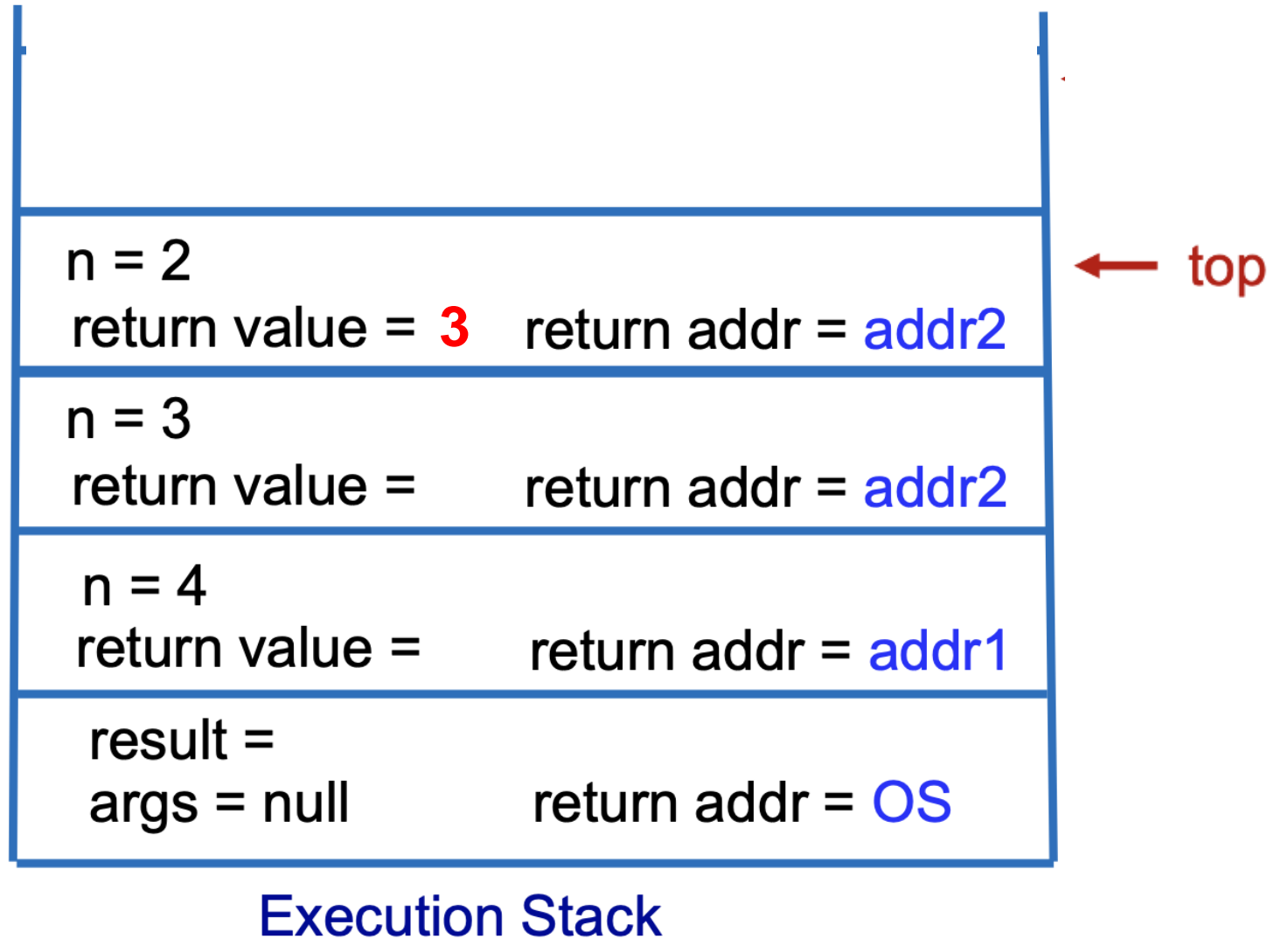
How Recursion Works (cont.)

- Since the value of n is 1 in the last invocation of the method `sum`, the statement `return 1` (base case) is executed.
- The value 1 is stored in the **return value**.



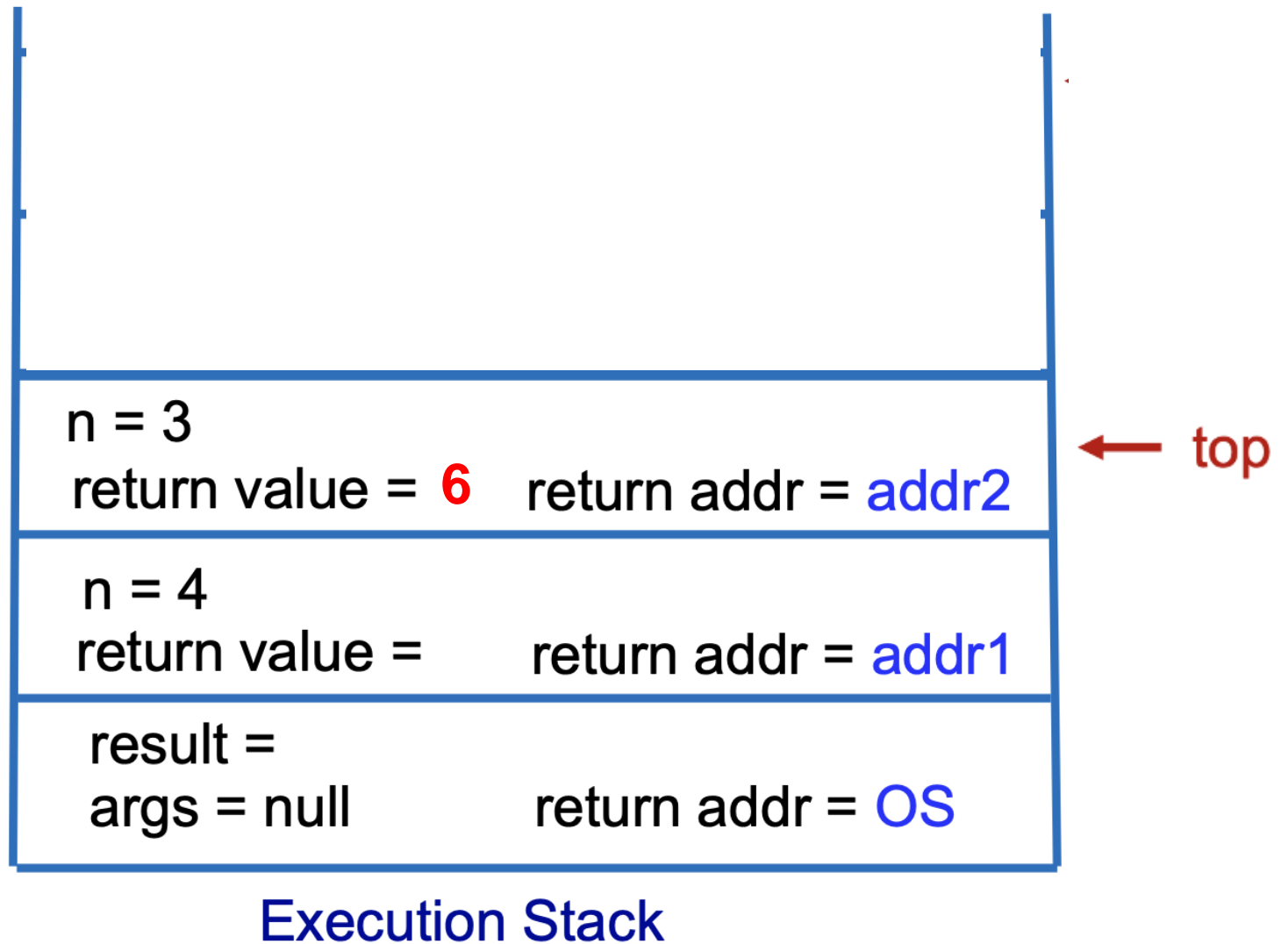
How Recursion Works (cont.)

- The method `sum` ends, and hence, an **activation record** is popped off the execution stack.
- The return address `addr2` is recovered, and execution continues at the statement in that address: This call just finished, and it returned the value `1`.
- Hence, $n + \text{sum}(n-1) = 2 + 1 = 3$ will be returned.

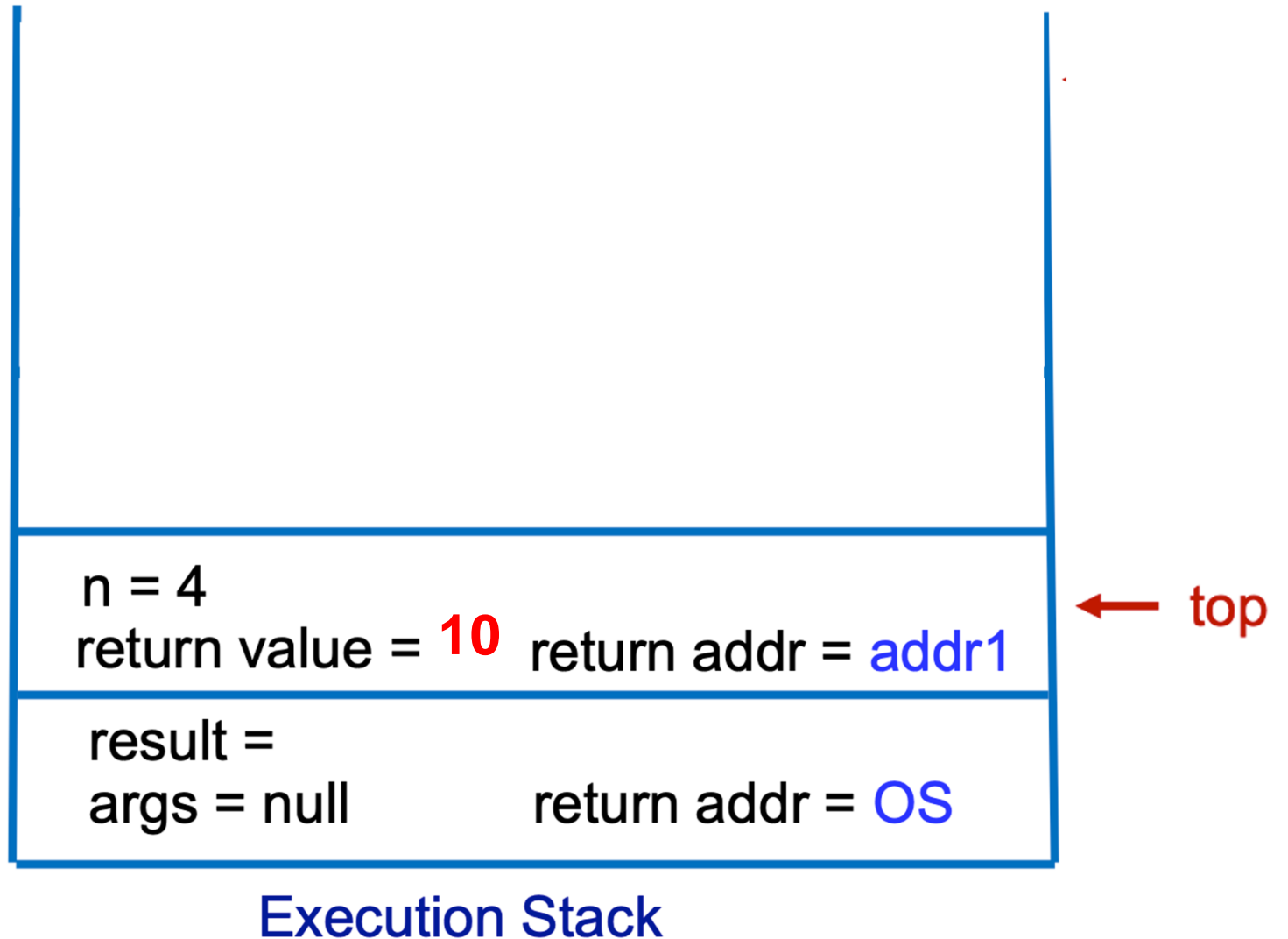


How Recursion Works (cont.)

- The next call returns the value **3**, and an activation record is popped off the **execution stack**. The return address **addr2** is recovered, and execution continues with the statement in that address.
- The value $n + \text{sum}(n-1) = 3 + 3 = \mathbf{6}$ will be returned.

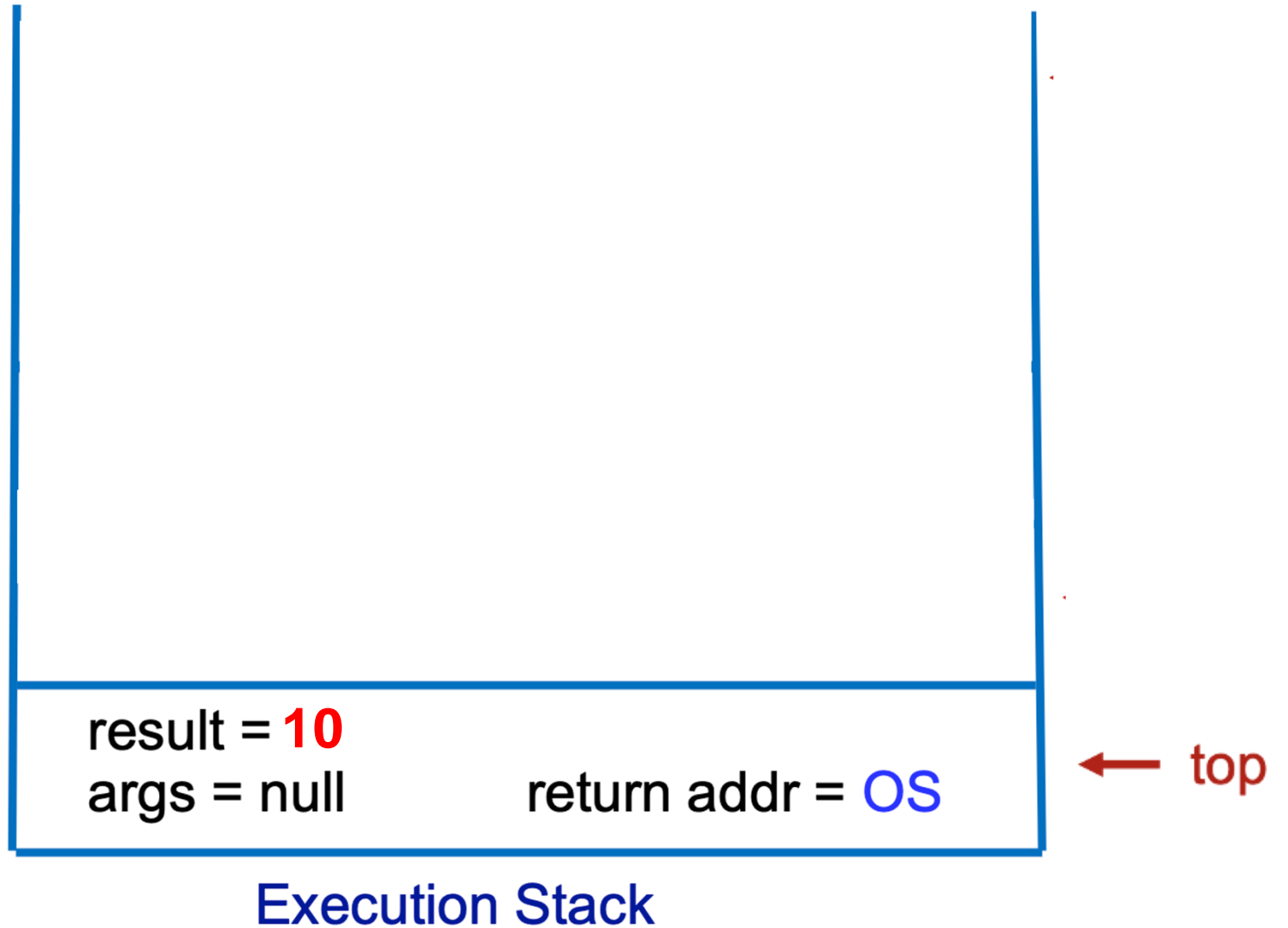


How Recursion Works (cont.)



How Recursion Works (cont.)


```
public static void main  
(String[] args) {int result =  
sum(4);} // addr 1
```



Recursion vs. Iteration

- Every **recursive** algorithm can also be written as an **iterative** algorithm. However, the algorithm could be much more complex and require the use of an auxiliary stack or other data structures to simulate the **execution stack**.
- Thus, just because we can use recursion to solve a problem does not mean we should!
- Would you use iteration or recursion to compute the sum of 1 to n ? Why?

Recursion vs. Iteration



- **Recursion** often uses more memory and can lead to stack overflow errors if the recursive depth is too large. Sometimes, an iterative (loop-based) solution is faster and more memory-efficient.



Thank
you