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CS 1027

Fundamentals of Computer
Science II

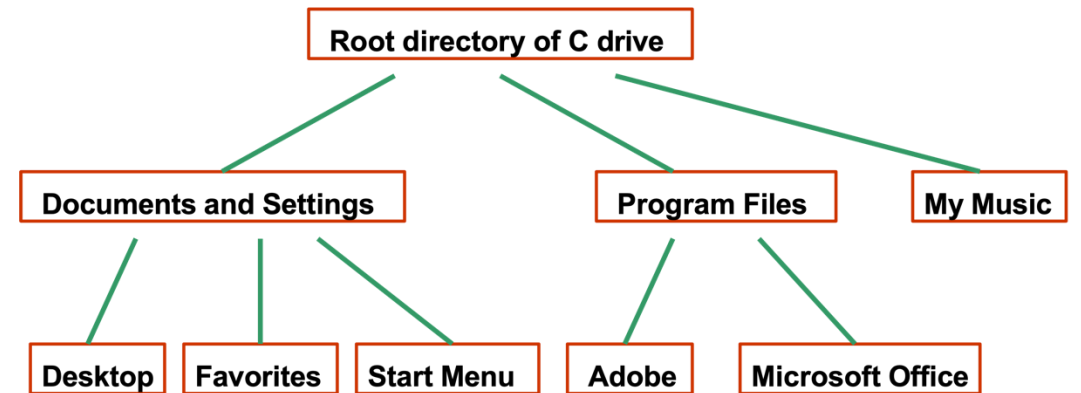
Trees **ADT**

Ahmed Ibrahim

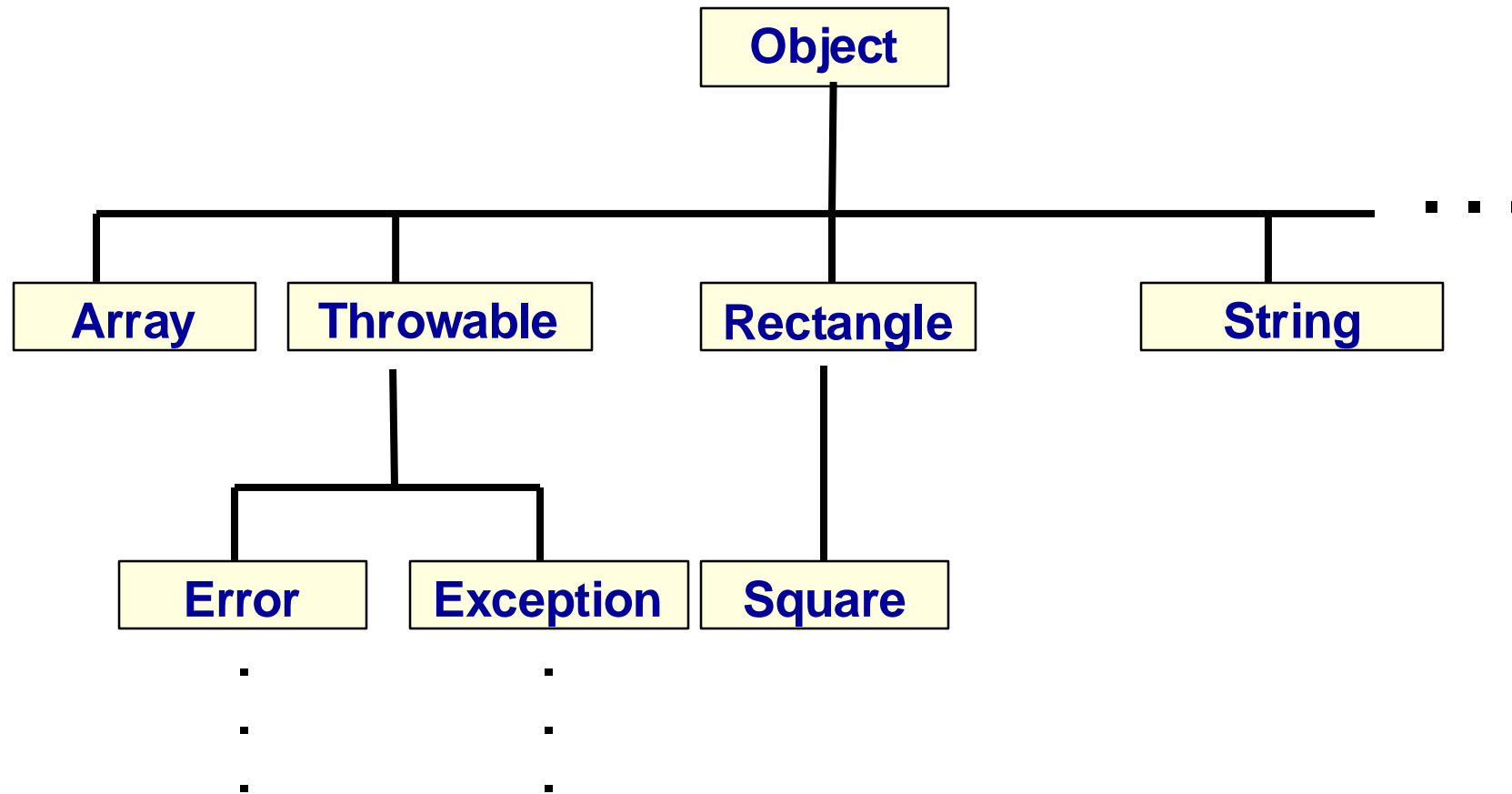


Trees

- A tree is a **non-linear abstract data type** that stores information in a hierarchy.
- Examples in real life:
 - Family tree
 - Table of contents of a book
 - Class Inheritance Hierarchy in Java
 - Computer file system (folders and subfolders)
 - Decision trees



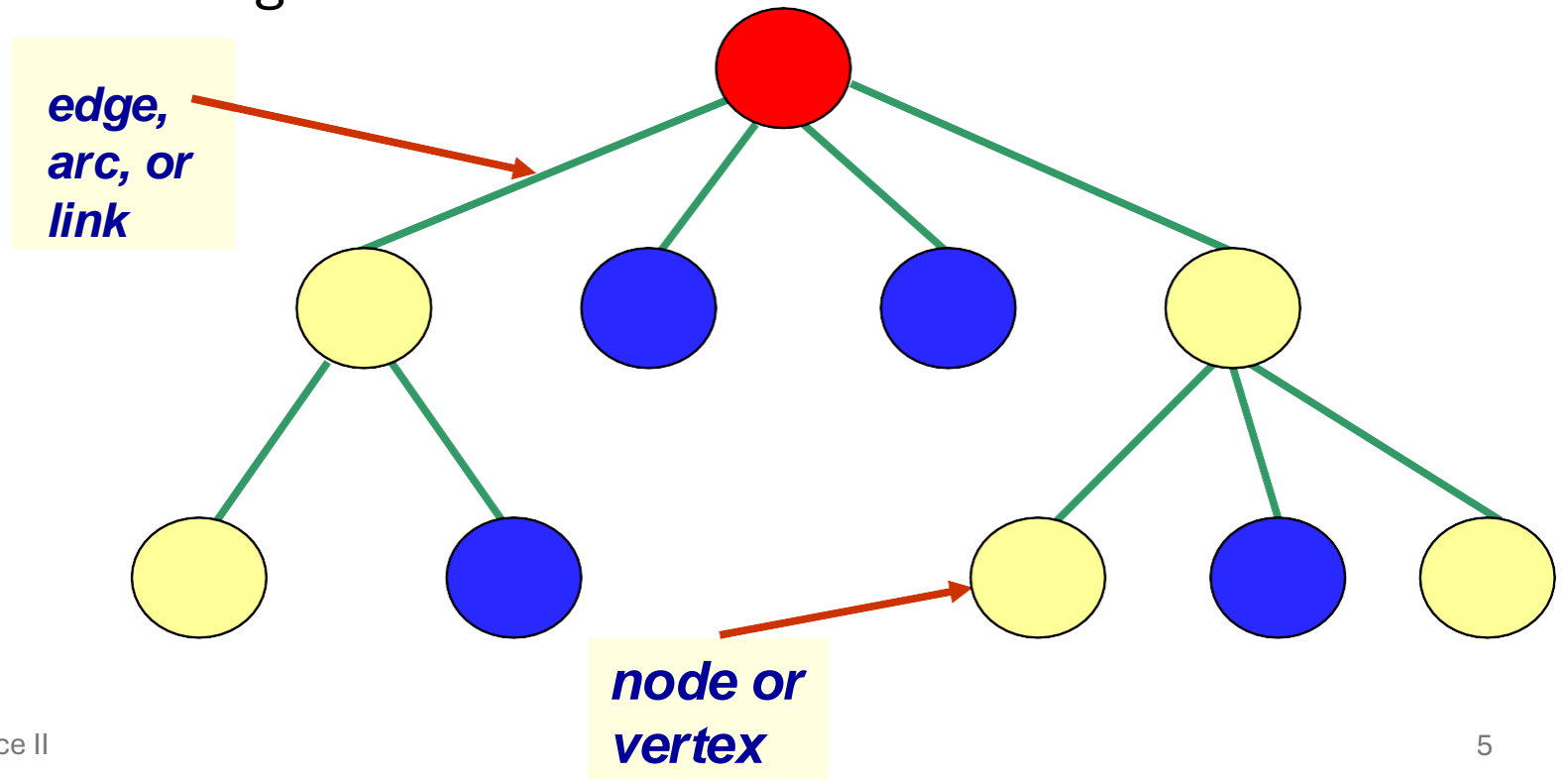
Example: Java's Class Hierarchy



Example: Java's Class Hierarchy

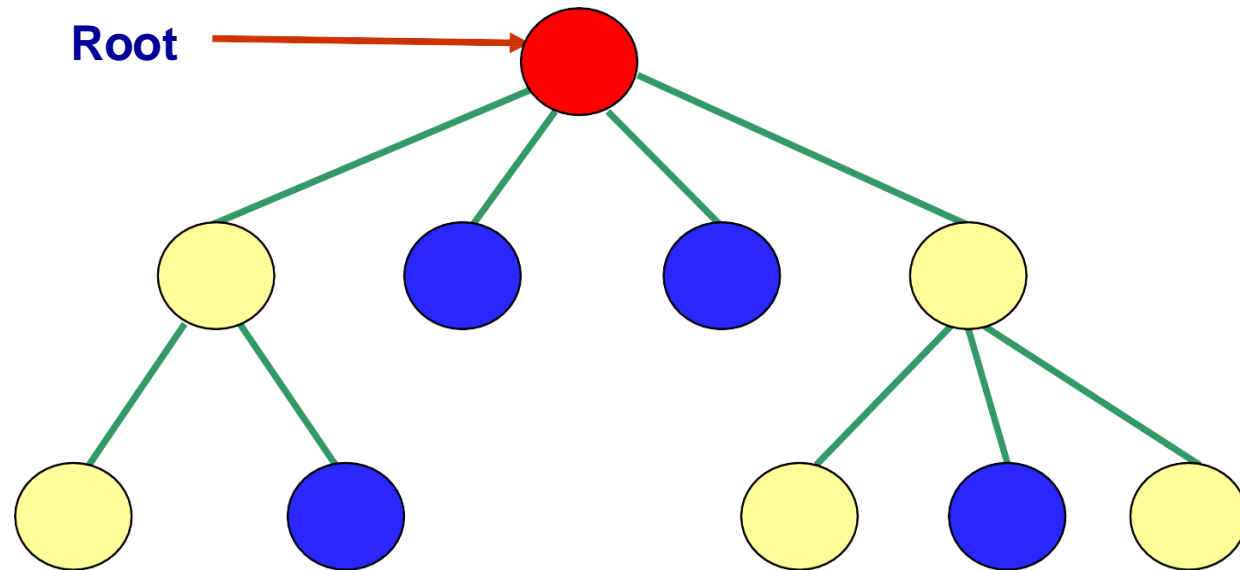
A **tree** consists of a set of

- **nodes** or vertices storing data and
- **edges**, links, or arcs connecting the nodes



Tree Definition

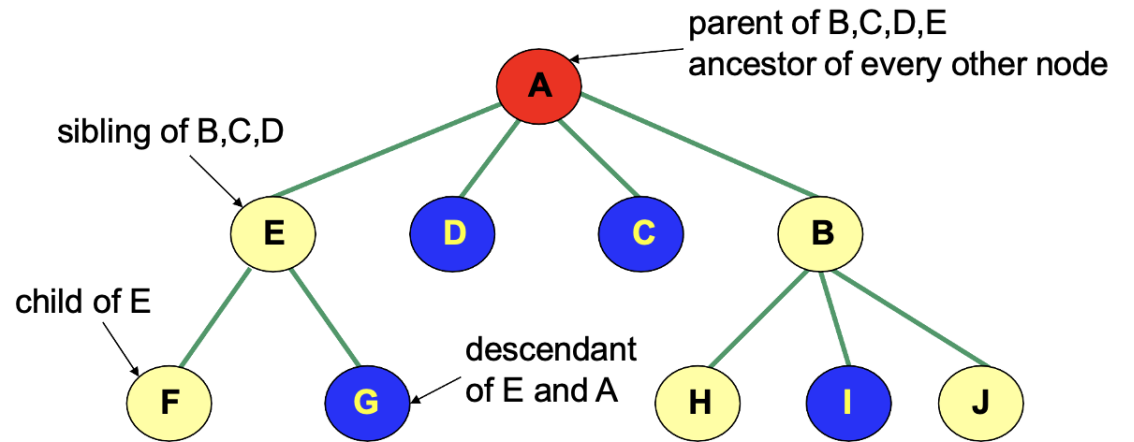
- There is a distinguished node called the root (usually drawn as the topmost node in the tree).



- An **empty** tree has no nodes or edges.

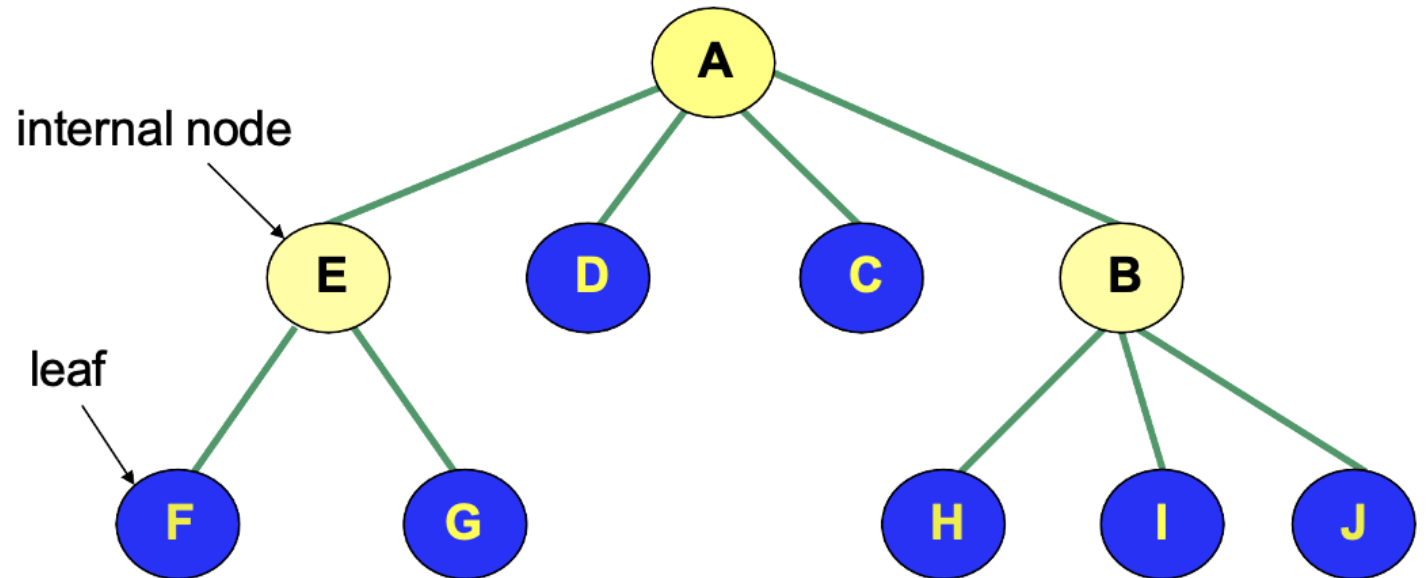
Tree Terminology

- **Parent:** the node directly above another node in the tree
- **Child:** a node directly below another node in the tree
- **Siblings:** nodes that have the same parent
- **Ancestors** of a node: its parent, the parent of its parent, etc.
- **Descendants** of a node: its children, the children of its children, etc.



Tree Terminology

- **Leaf node:** a node without children
- **Internal node:** a node that is not a leaf node



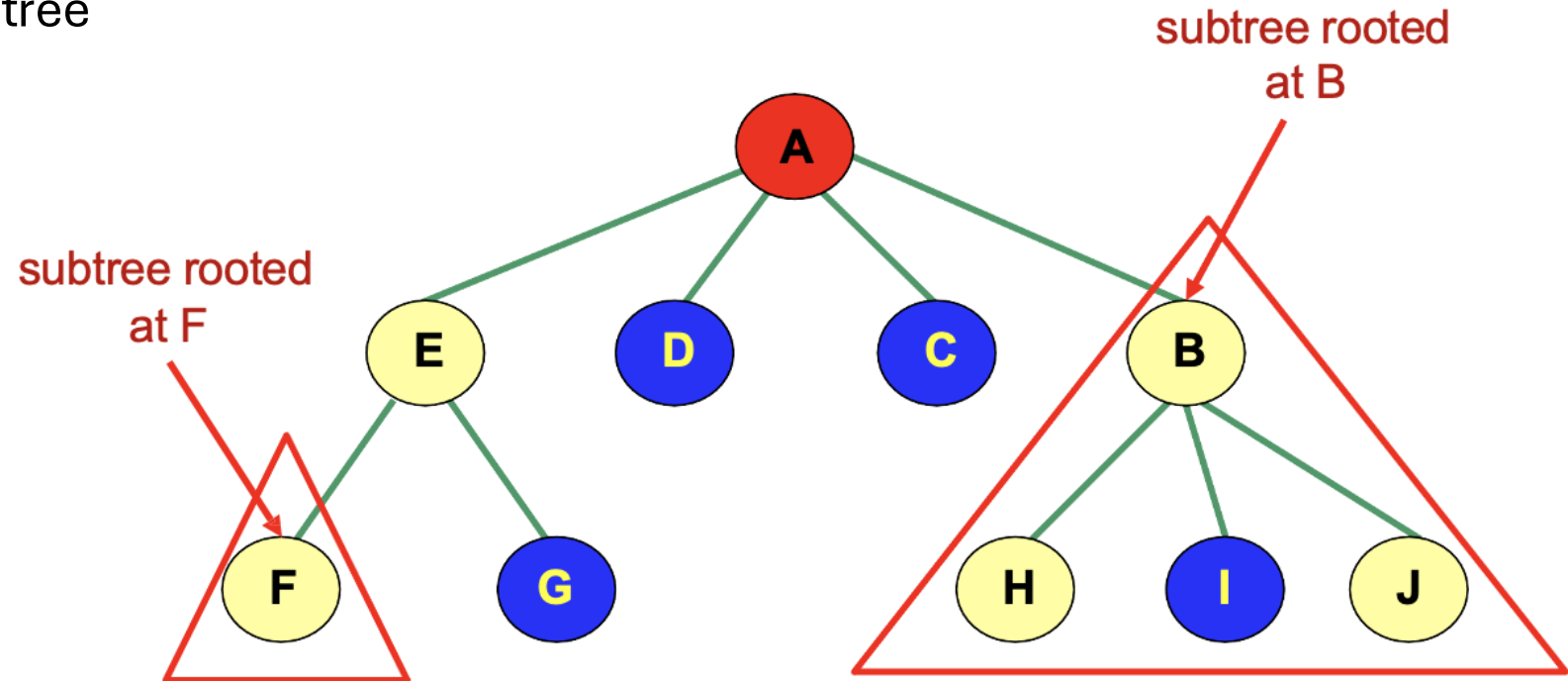
Stop & Think



- Does a leaf node have any children?
- Does the root node have a parent?
- How many children can a node have?

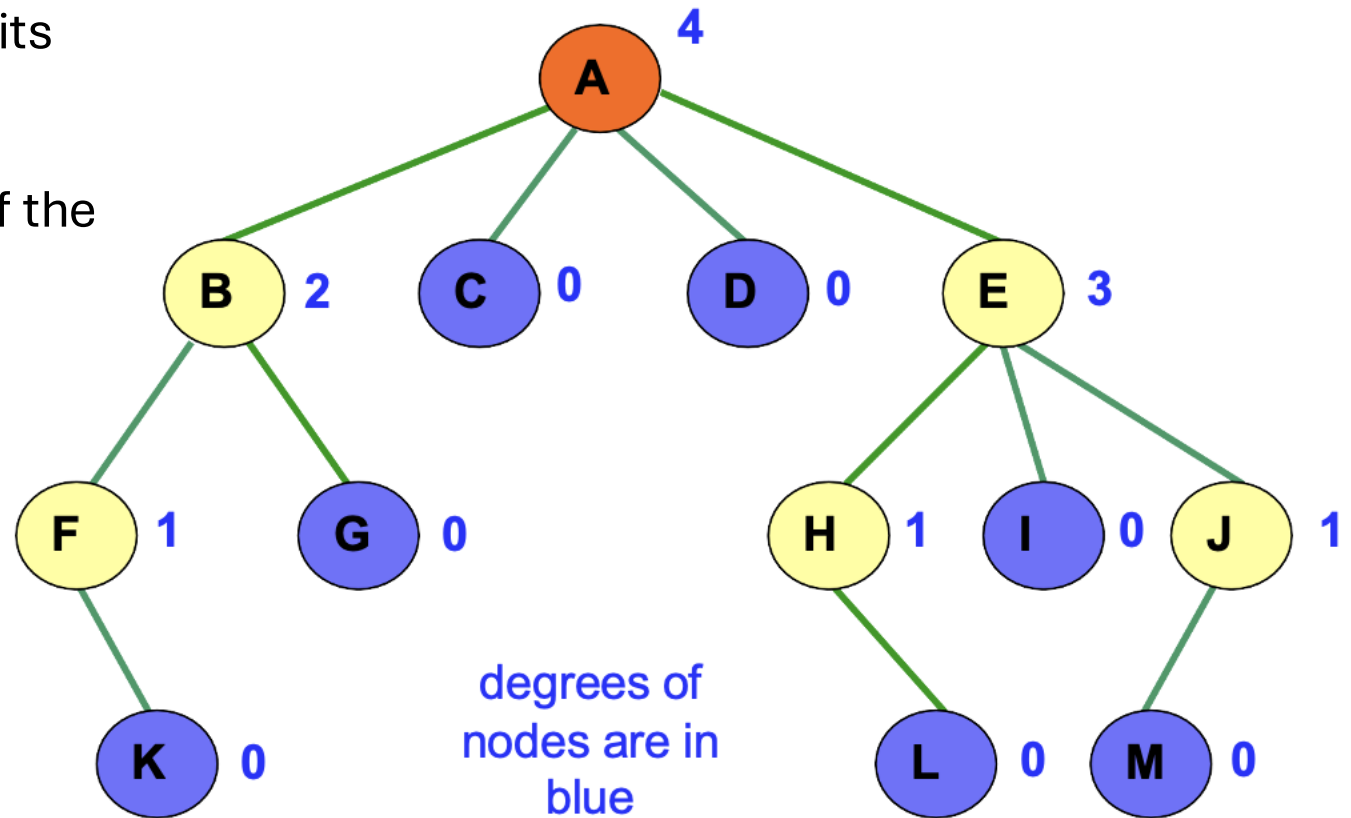
Subtrees

- The subtree rooted at a node consists of the node and all its descendants
- A subtree is itself a tree



Tree Terminology

- **Degree of a node:** the number of its children
- **Degree of a tree:** the maximum of the degrees of the tree's nodes

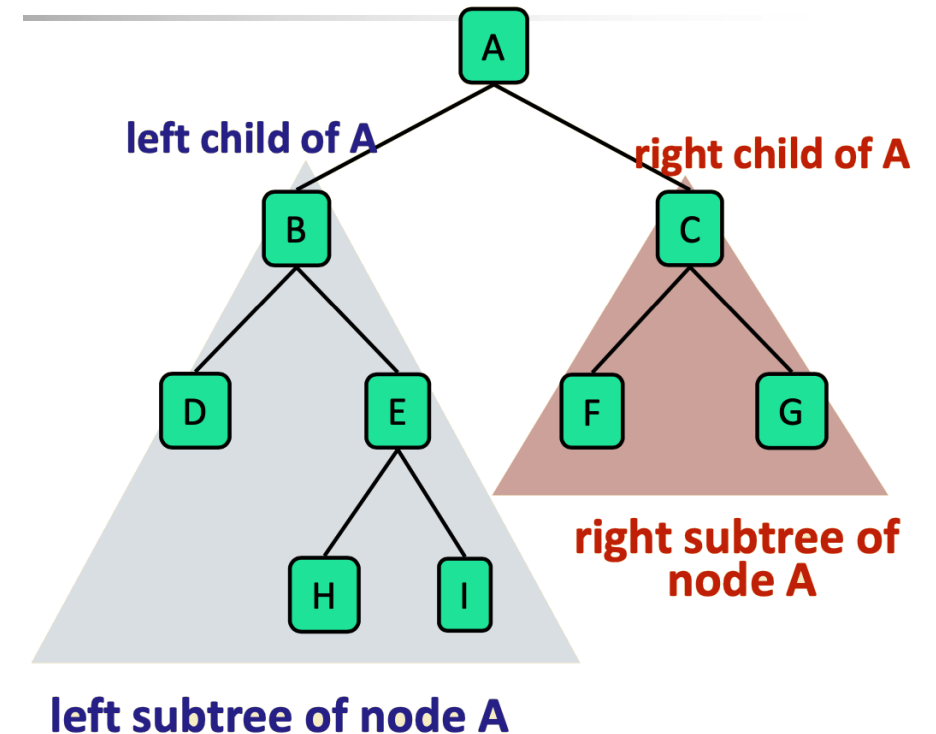


Classification of Trees

- Trees can be classified into many categories by their **properties** and **applications**. We will look into the following categories.
 - General trees – no restriction
 - Binary trees – each node has at most two children.
 - Binary search trees – binary trees for efficient searching
 - Ex. **AVL trees**: height-balanced binary search trees
 - Multi-way search trees – a generalization to binary search trees
 - Ex. **B-trees** – balanced multi-way search trees

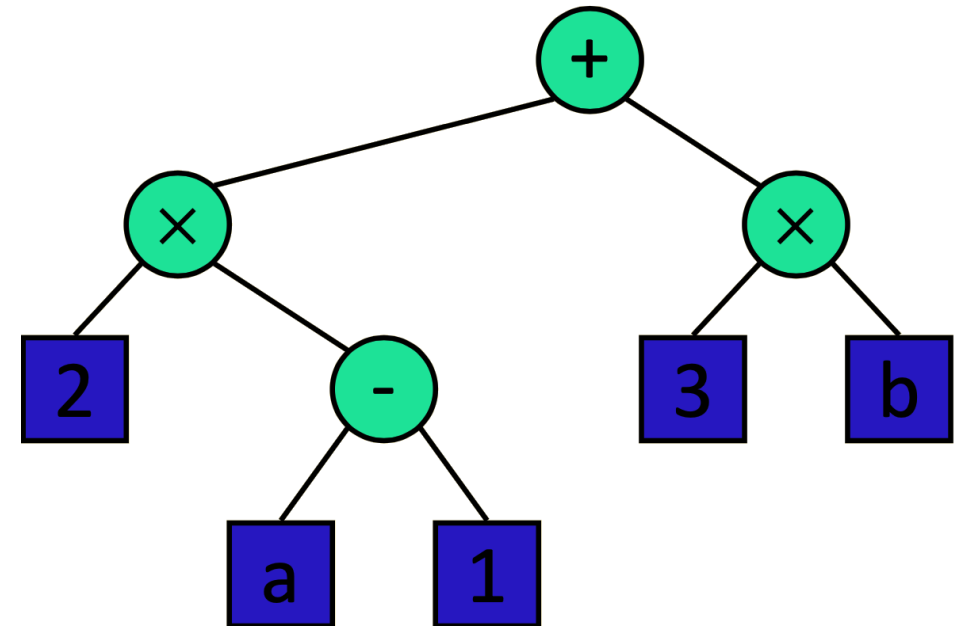
Binary Trees (BTs)

- In a **Binary tree**, a node has at most two children
- Children are an ordered pair
 - **left child** and **right child**
 - corresponding subtrees are the **left subtree** and **right subtree**
- In a binary tree, each internal node has exactly two children
- Applications
 - Arithmetic expressions
 - Decision processes



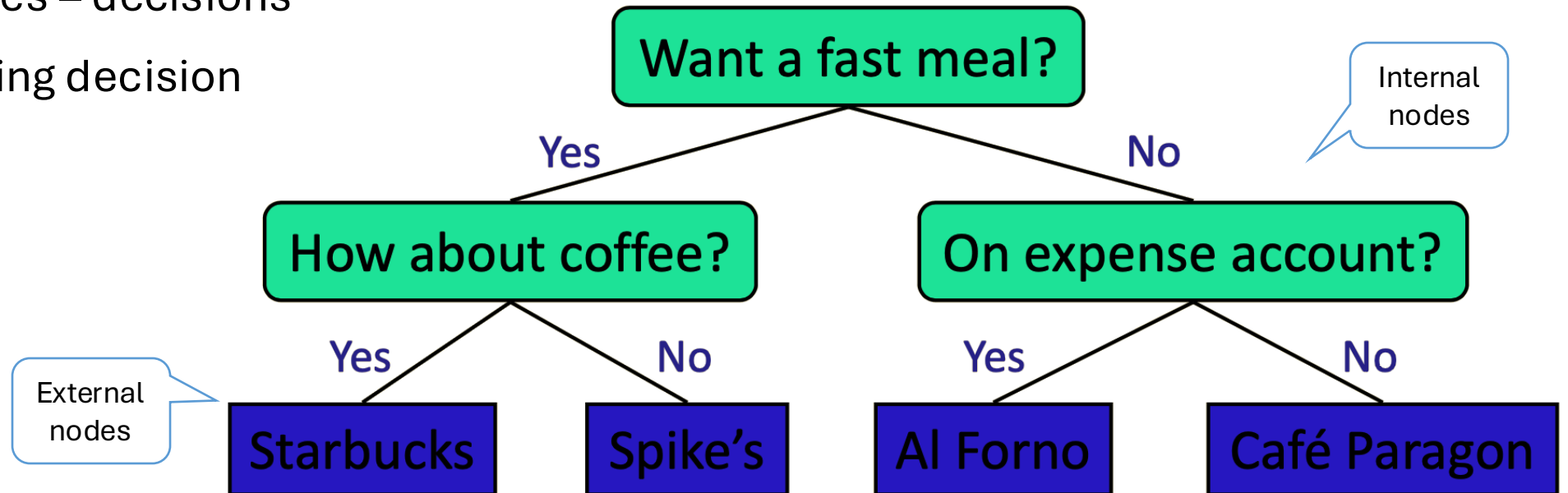
Arithmetic Expression Tree

- The binary tree associated with an arithmetic expression
 - internal nodes store operators
 - external nodes store operands
- Example: arithmetic expression tree for the expression: $(2 \times (a - 1) + (3 \times b))$



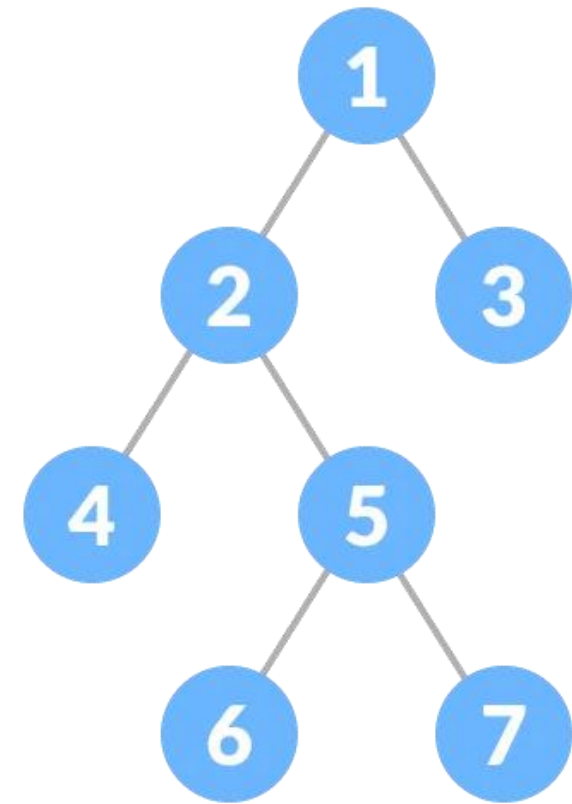
Decision Tree

- Binary tree associated with a decision process
- Internal nodes – questions with yes/no answer
- External nodes – decisions
- Example dining decision



Types of BTs: Full Binary Tree

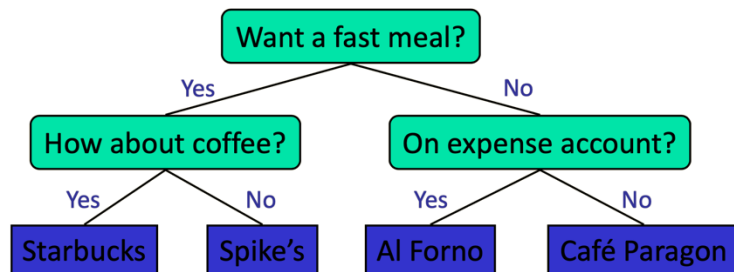
- A **full** Binary tree is a special type of binary tree in which every parent node/internal node has either **two** or **no** children.
- It is also known as a **proper** binary tree.



Properties of Proper BT

- Notations

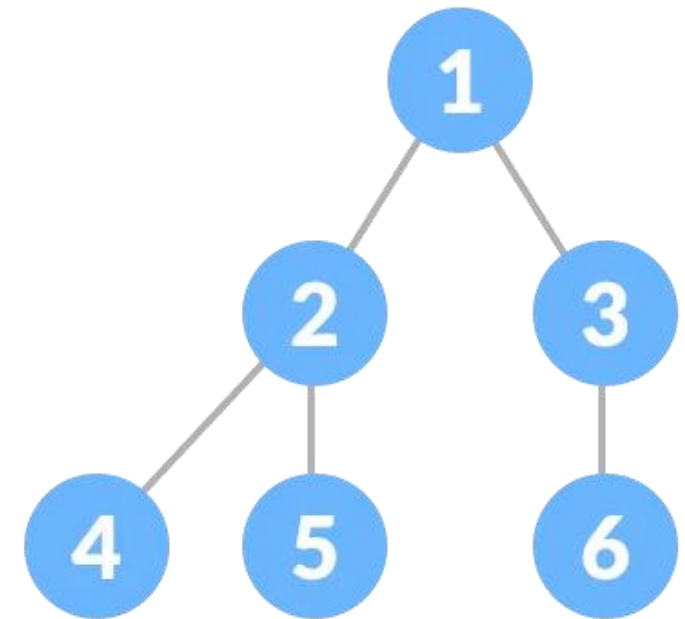
- **n** number of all nodes
- **e** number of external nodes
- **i** number of internal nodes
- **d** depth
- **h** height



1. The number of leaves is $i + 1$.
2. The total number of nodes is $2i + 1$.
3. The number of internal nodes is $(n - 1) / 2$.
4. The number of leaves is $(n + 1) / 2$.
5. The total number of nodes is $2e - 1$.
6. The number of internal nodes is $e - 1$.
7. The number of leaves is at most 2^{h-1} .

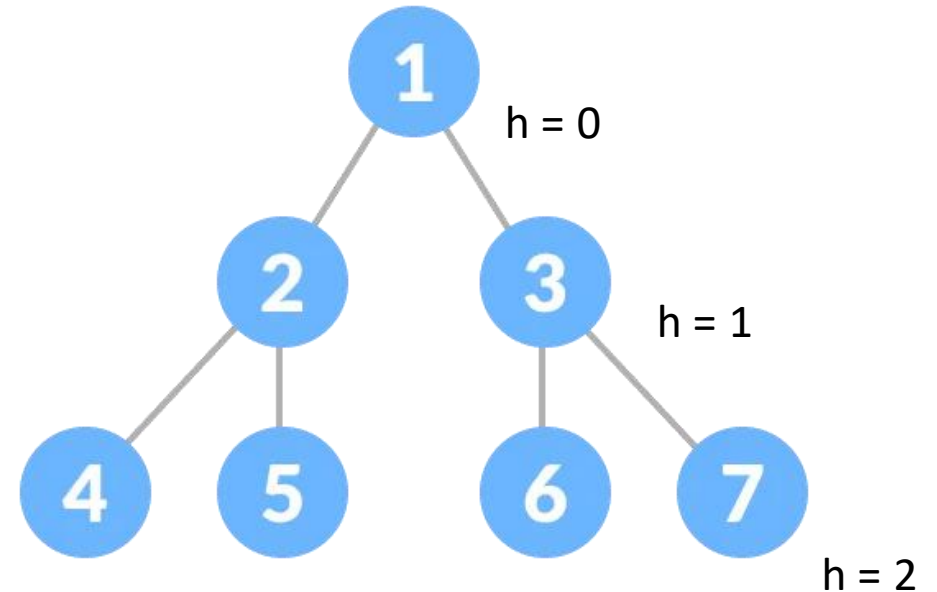
Types of BTs: Complete Binary Tree

- A complete binary tree is a binary tree in which every level, except possibly the last, is **completely filled**, and all nodes are as far left as possible.
- This means that:
 - All levels above the last level are fully filled.
 - The last level may not be fully filled, but if it has missing nodes, those nodes are only on the right side (i.e., **all leaf nodes lean to the left**).



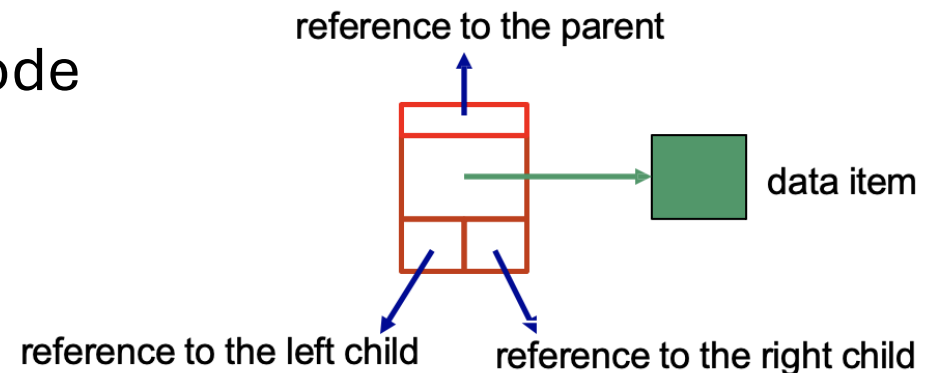
Types of BTs: Perfect Binary Tree

- A perfect binary tree is a binary tree in which every internal node has exactly **two child nodes** and all the leaf nodes are at the same level.
- A perfect binary tree of height h has $2^{h+1} - 1$ node.
- A perfect binary tree of height h has 2^h leaf nodes.



Linked Binary Tree Implementation

- To represent a binary tree, we will use a **linked** structure of nodes
 - **root**: reference to the node that is the **root of the tree**
 - **count**: keeps track of the number of nodes in the tree
- First, how will we represent a *node of a binary tree*?
- A **binary tree node** will contain
 - a **reference** to the data stored in the node
 - **references** to its **left** and **right** children
 - [**optionally**] a reference to its parent

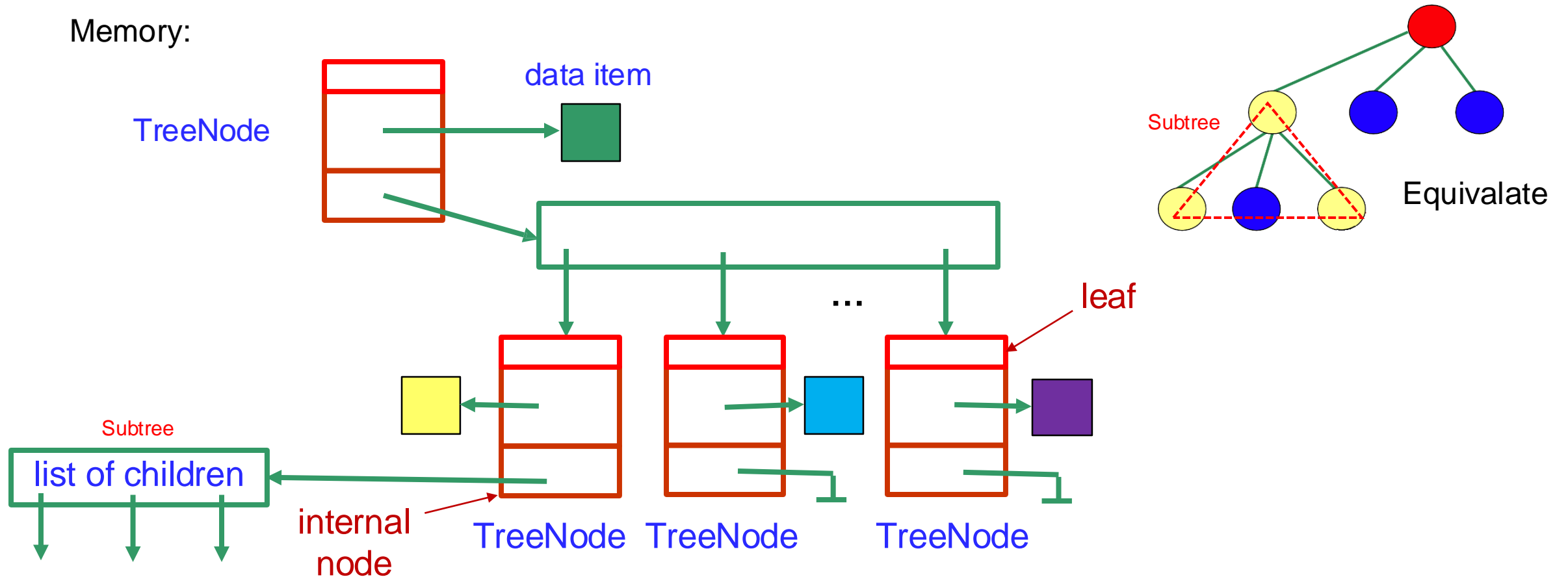


Linked Binary Tree Implementation

```
public class BinaryTreeNode<T> {  
    private T dataItem;  
    private BinaryTreeNode<T> parent, leftChild, rightChild;  
    /* Creates a new tree node with the specified data. */  
    BinaryTreeNode (T newData) {  
        dataItem = newData;  
        leftChild = null; rightChild = null; parent = null;  
    }  
    // Getter and setter methods  
    public BinaryTreeNode<T> getParent() {  
        ...  
    }  
}
```

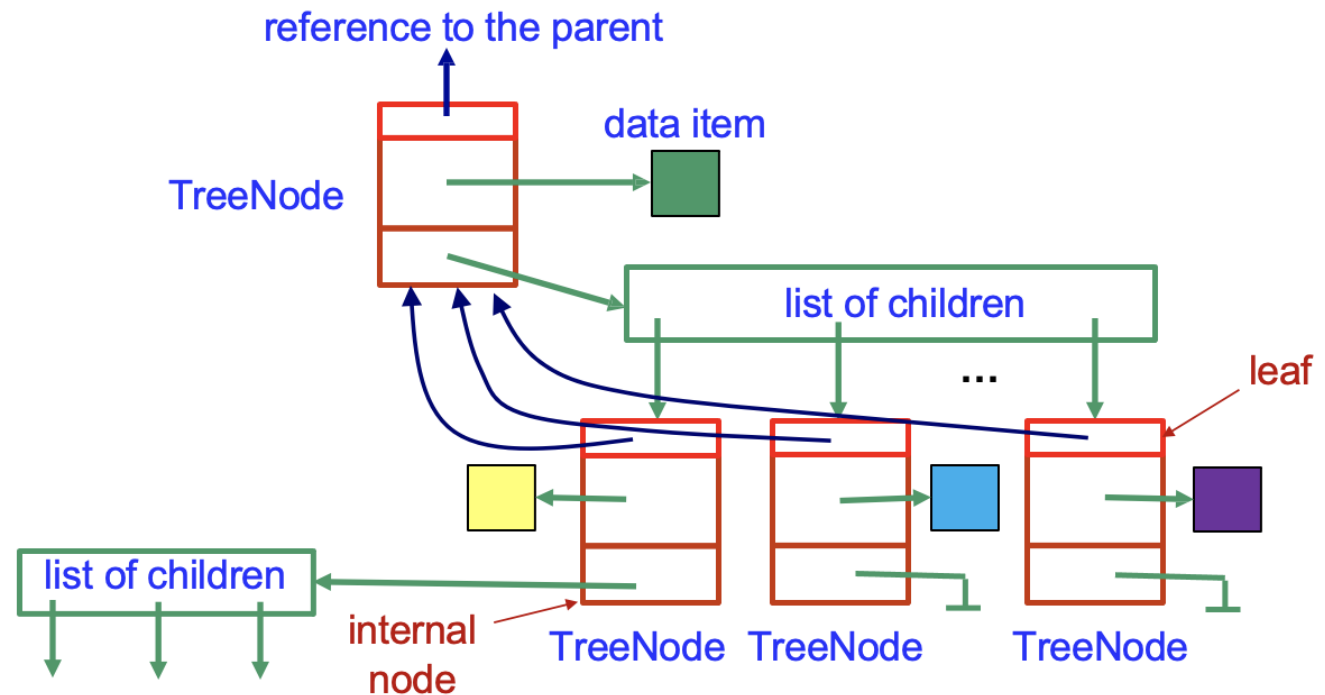
A **TreeNode** Implementation (v1)

Memory:



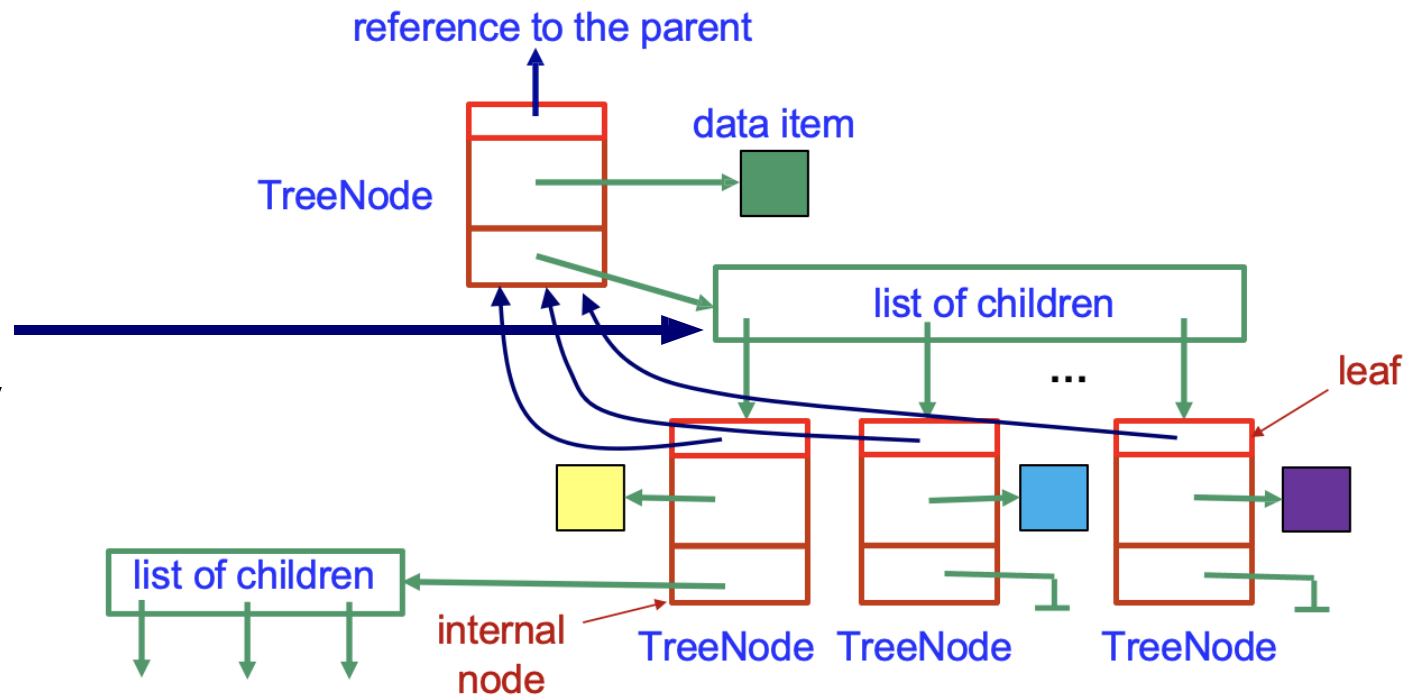
A **TreeNode** Implementation (v2)

- If the tree is not binary, then each node will have a **reference** to the data item it stores, a **reference** to its parent, and a **reference** to a list of its children.



A **TreeNode** Implementation (v2)

- The children of a node can be stored in an array, a **circular array**, a **singly linked list**, a **doubly linked list**, or any other data structure implementing a list.



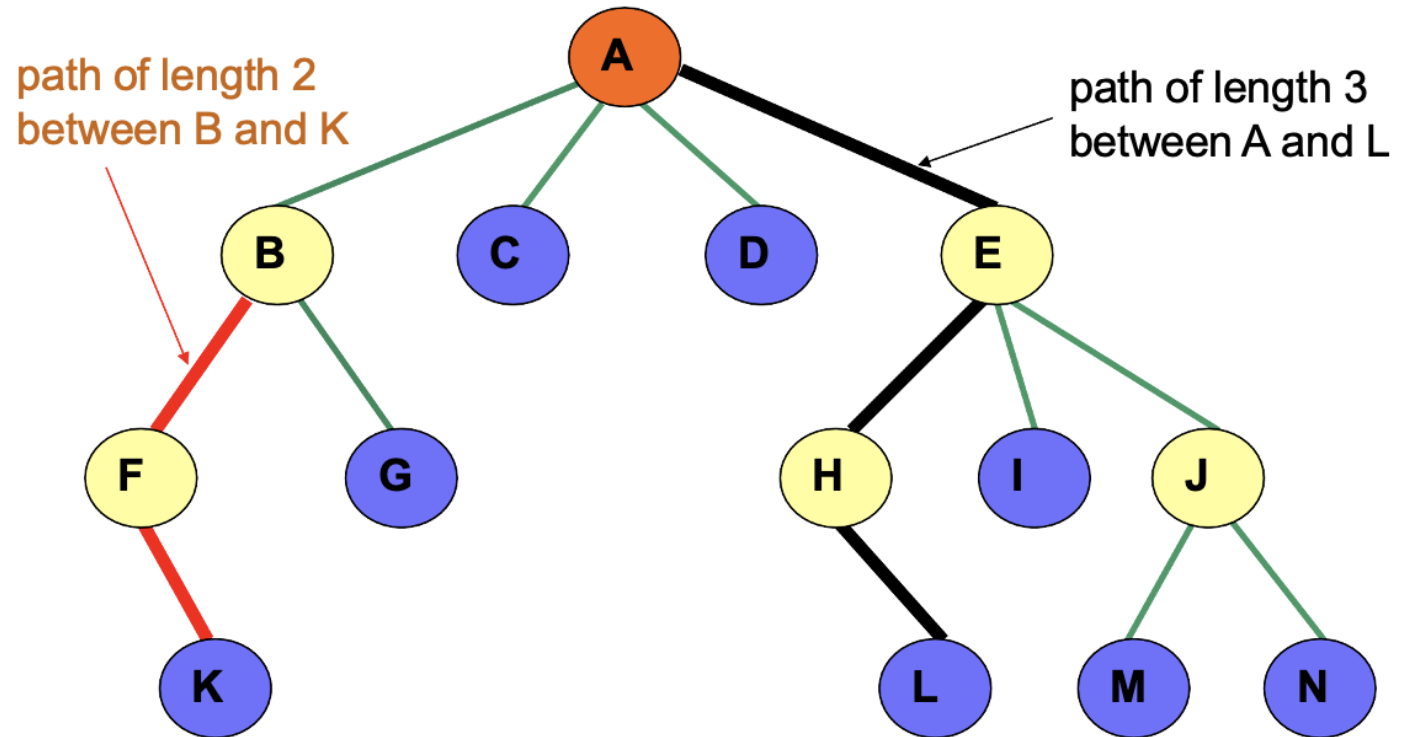
Linked Binary Tree Implementation

- This implementation uses an array to store the list of children:

```
public class TreeNode<T> {  
    private final int DEFAULT_CAPACITY = 10;  
    private T dataItem;  
    private TreeNode<T> parent; //optional  
    private TreeNode<T>[] children;  
    private int numChildren;  
  
    /* Creates a new tree node with the specified data. */  
    TreeNode (T newData) {  
        dataItem = newData;  
        parent = null; //optional  
        children = new TreeNode<T>[DEFAULT_CAPACITY];  
        numChildren = 0;  
    }  
}
```

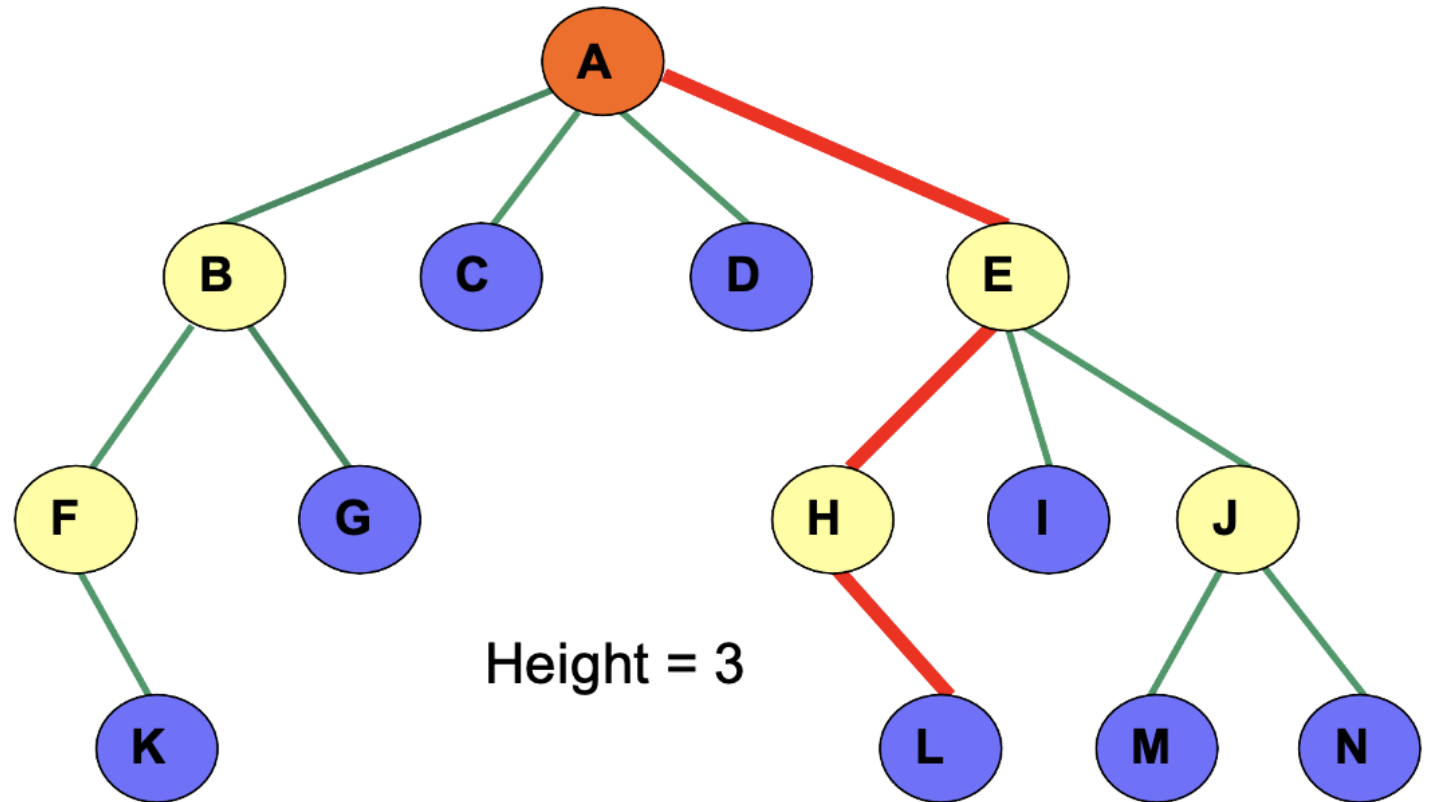
Tree Terminology

- A path is a sequence of edges leading from one node to another
- Length of a path:
number of edges on the path



Tree Terminology

- Height of a tree: length of the longest path from the root to a leaf
- What is the height of a tree that has only a root node?

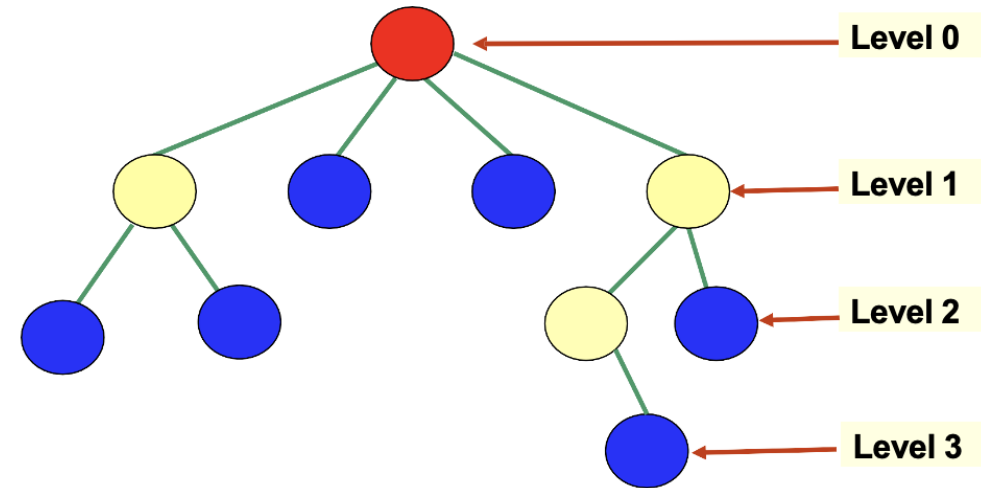


Level of a Node

- Level of a node: number of edges between the root and the node
- It can be defined recursively:
 - The level of the **root** node is 0
 - The level of a node that is not the **root** is the level of its parent + 1.

What is a tree's height (h) in terms of levels?

In terms of levels, the tree has a height of **3**, because there are 3 levels.



Algorithm level (node)

Input: node of a tree

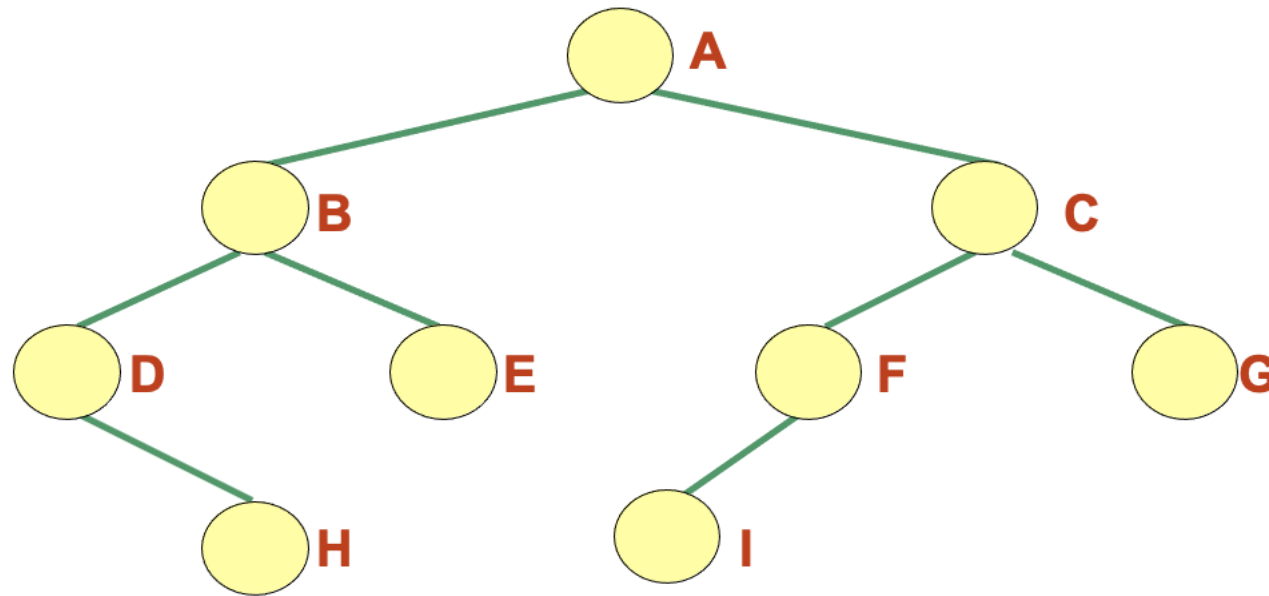
Output: level of the node

```
public int level (TreeNode<T> node) {  
    // Input: node of a tree  
    // Output: level of the node  
    if (node.getParent() == null) return 0;  
    else return 1 + level(node.getParent());}
```

Tree Traversals

- Given the root node of a tree, a traversal requires visiting each node once.
- Note that the only node we know of in a tree is its root. Using a tree traversal, we must be able to access all the other nodes in the tree from the root node.
- Common tree traversals:
 - preorder
 - postorder
 - level-order
- For binary trees, there is another traversal:
 - inorder

Binary Tree Traversals



We will consider **only traversals of binary trees**. We will study the different tree traversals using this and other trees.

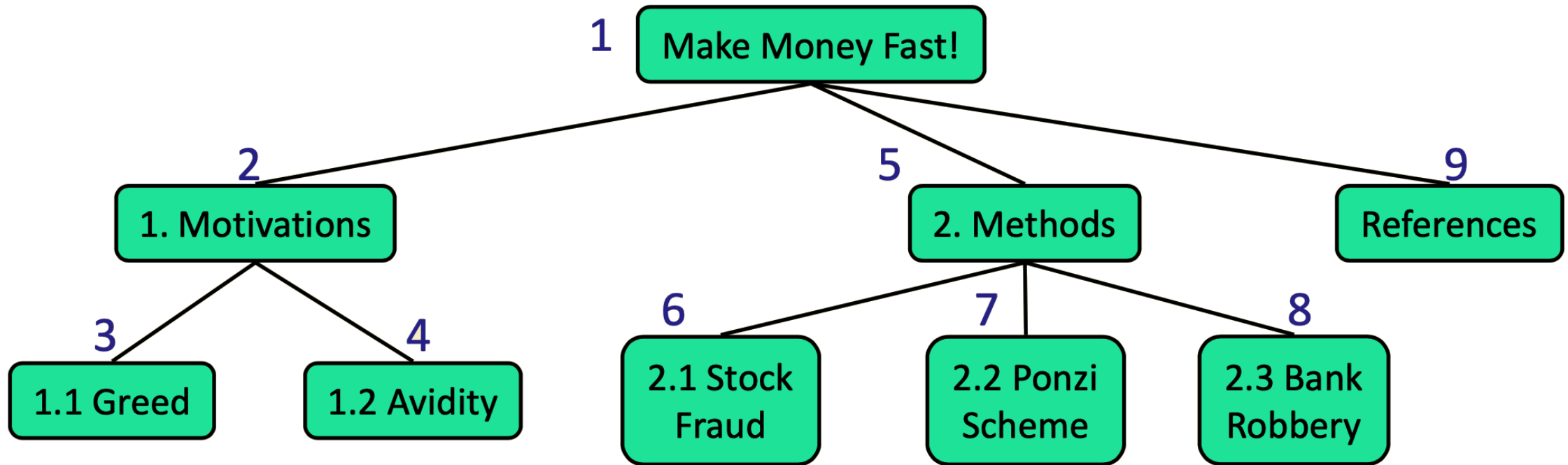
Pre-order Traversal

- If the tree is not empty:
 - Visit the **root** node of the tree
 - Perform **pre-order traversal** of the left subtree
 - Perform **pre-order traversal** of the right subtree
- This is a **recursive algorithm** for performing a pre-order traversal of a tree.
 - *What is the base case?*
 - *What is the recursive case?*

```
public void preorder (BinaryTreeNode<T> r)
{ if (r != null) {
    visit(r); // This method depends on the
              // application traversing the tree
    preorder (r.getLeftChild());
    preorder (r.getRightChild());
  }
}
```

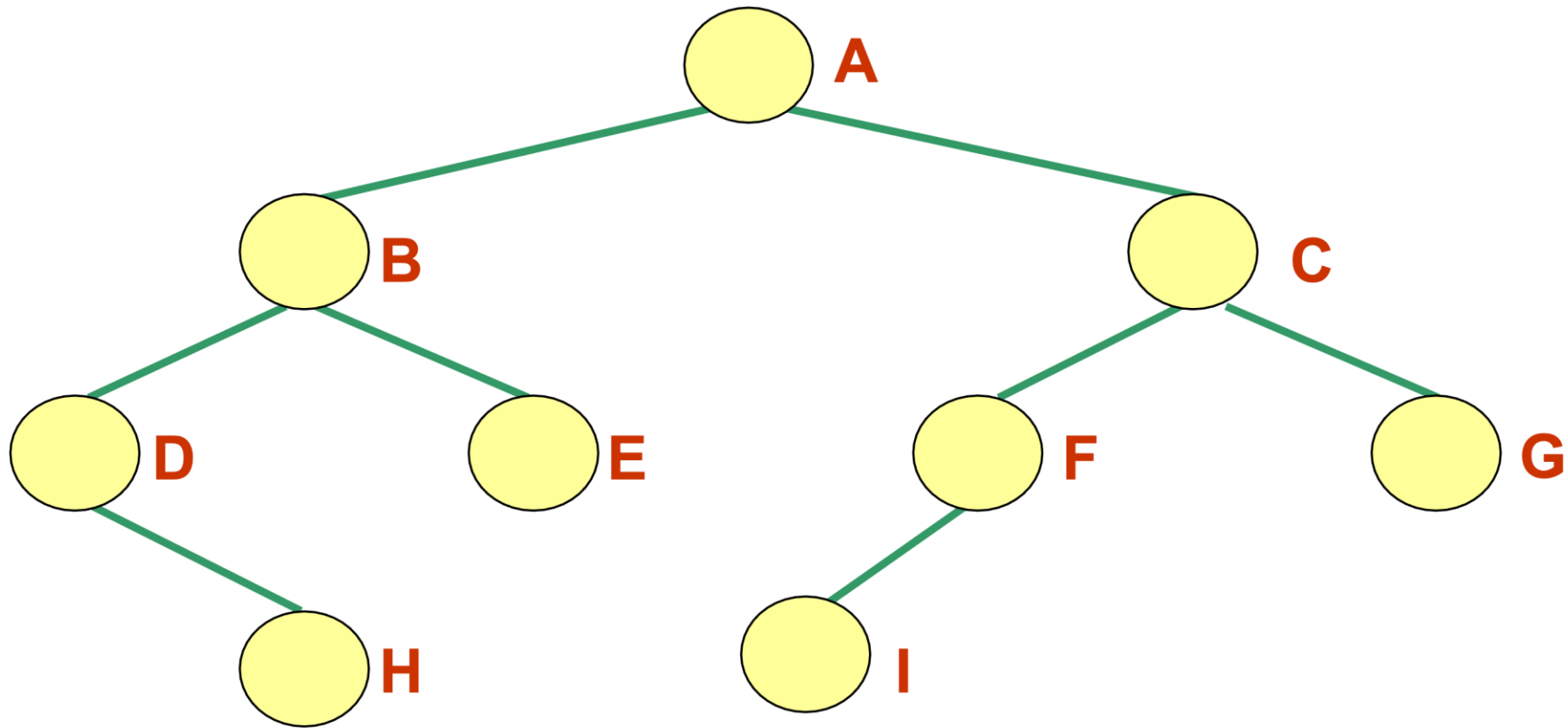
Pre-order Traversal

- A node is visited before its descendants.
- When is it applied?
 - Use when must perform computations for a node before any computations for its descendants



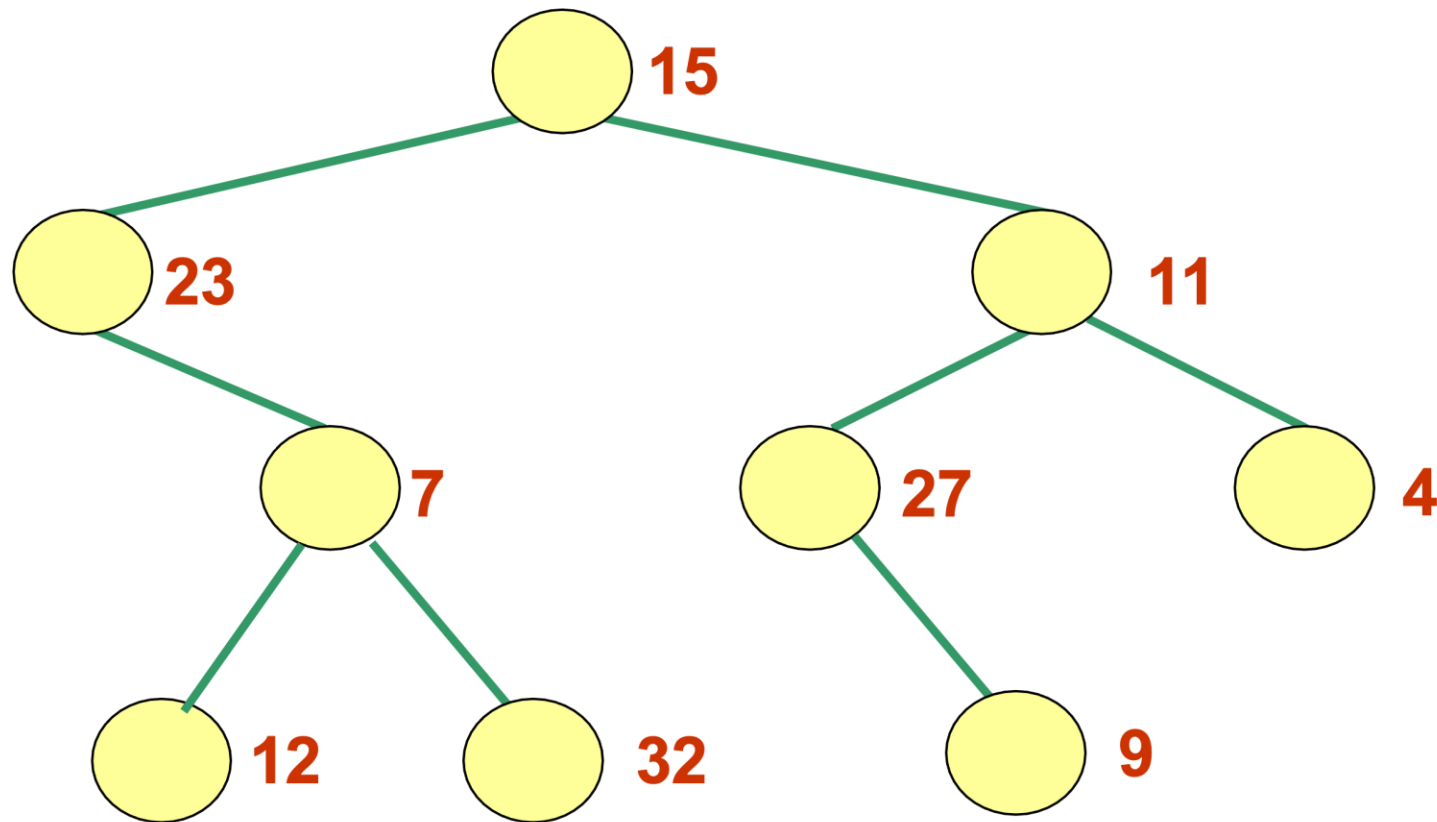
Pre-order Traversal Example 1

- Pre-order traversal: A B D H E C F I G



Pre-order Traversal Example 2

- Pre-order traversal: 15 23 7 12 32 11 27 9 4

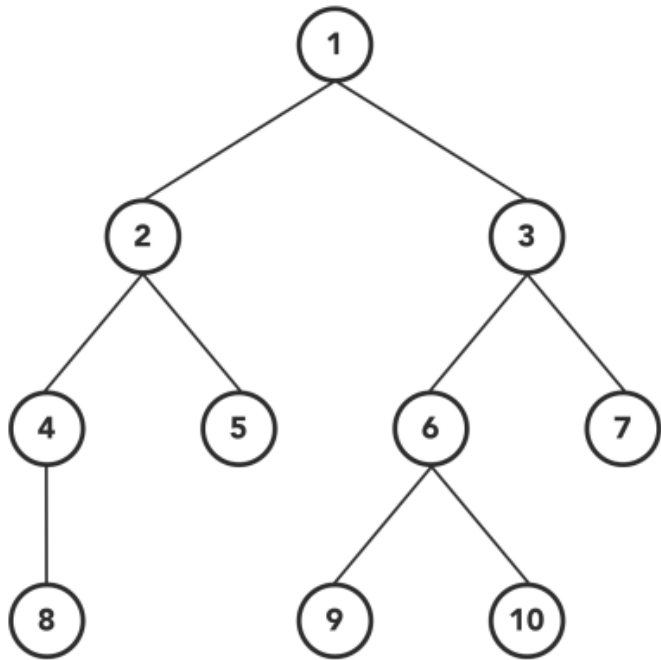


In-order Traversal

- If the tree is not empty,
 - Perform **in-order traversal** of the left subtree
 - Visit the root node of the tree
 - Perform **in-order traversal** of the right subtree
- This is a **recursive algorithm** for performing an in-order traversal of a tree.
 - *What is the base case?*
 - *What is the recursive case?*

```
public void inorder
(BinaryTreeNode<T> r) {
    if (r != null) {
        inorder (r.getLeftChild());
        visit(r);
        inorder (r.getRightChild());
    }
}
```

In-order Traversal

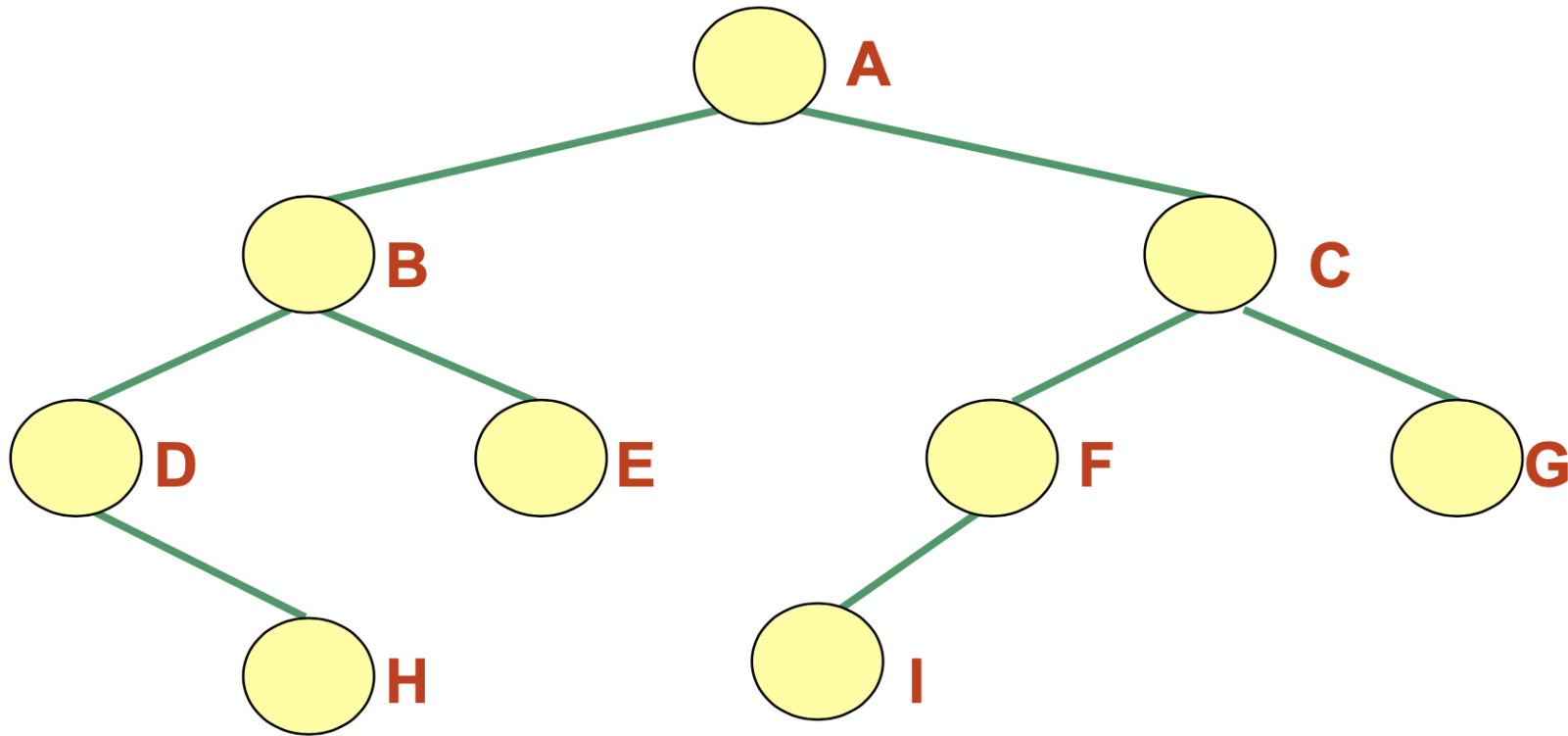


```
public void inorder (BinaryTreeNode<T> r)
{
    if (r != null) {
        inorder (r.getLeftChild());
        visit(r);
        inorder (r.getRightChild());
    }
}
```

In-order Traversal

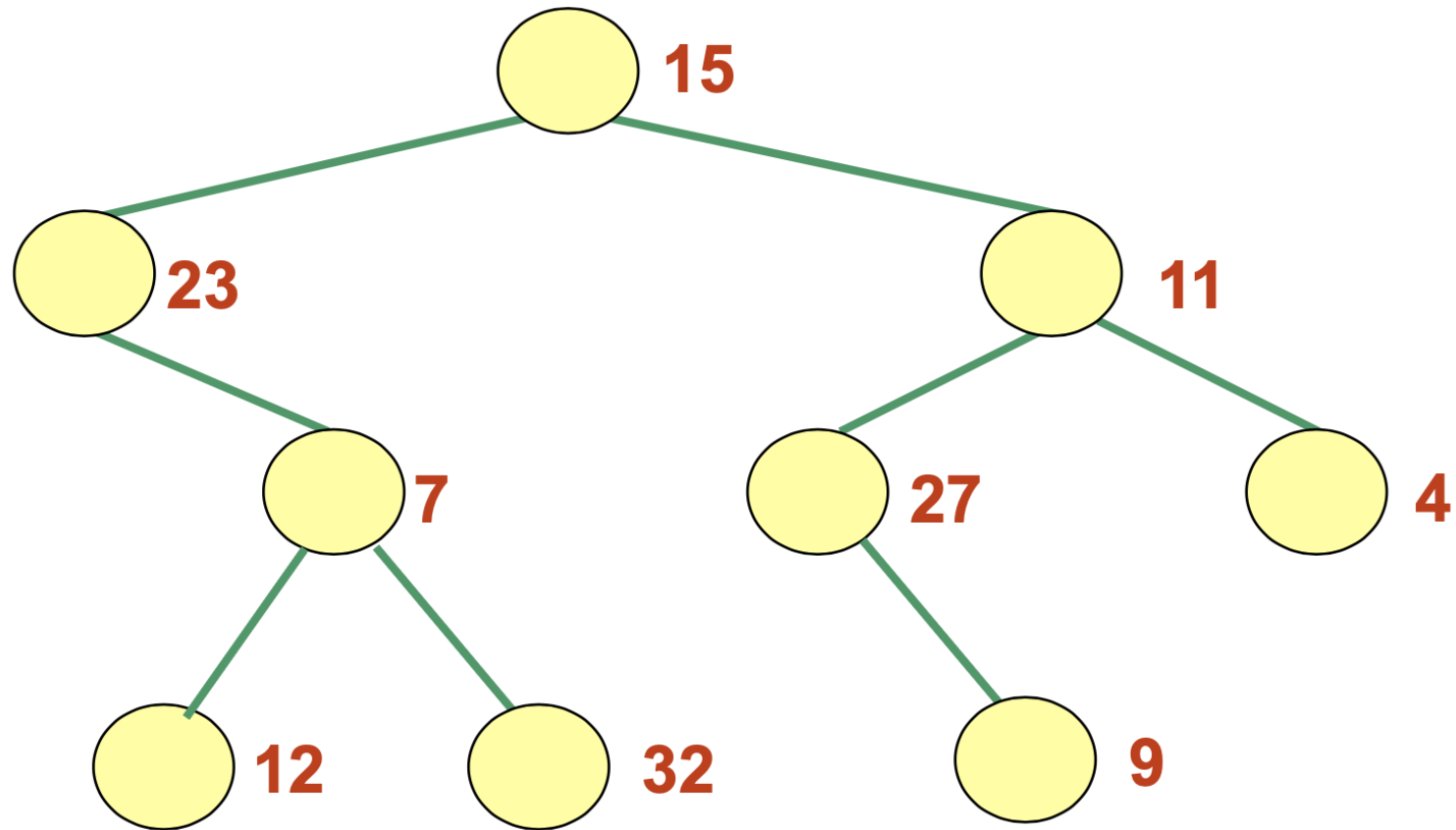
Example 1

- In-order traversal: D H B E A I F C G



In-order Traversal Example 2

- In-order traversal: 23 12 7 32 15 27 9 11 4



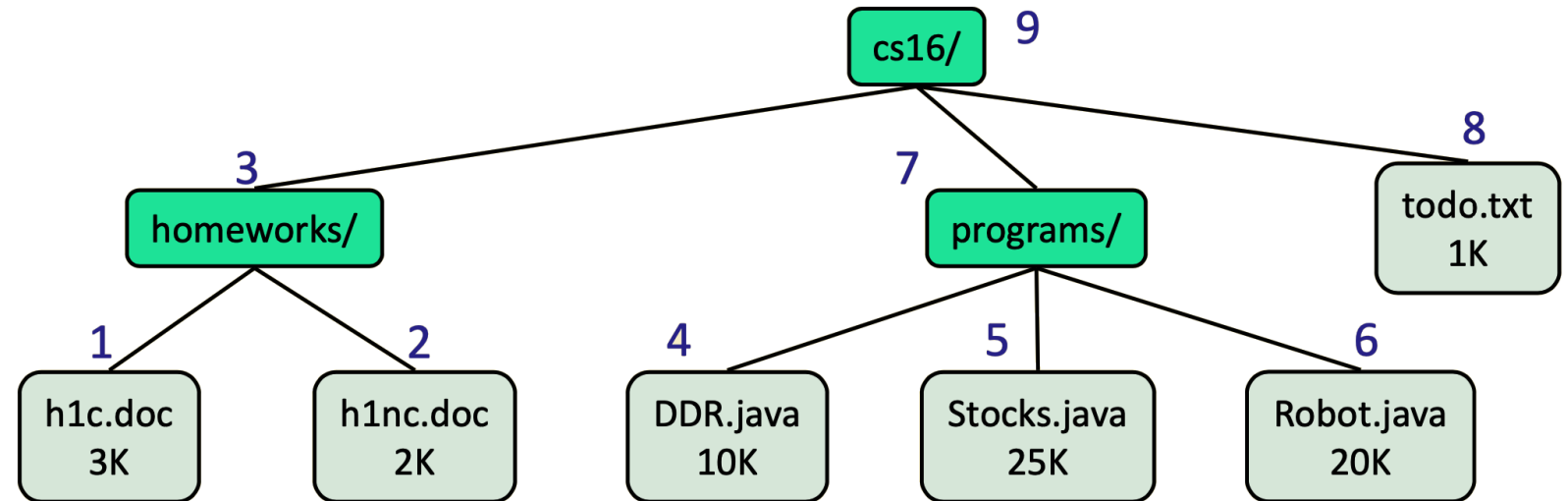
Post-order Traversal

- If the tree is not empty,
 - Perform **post-order traversal** of the left subtree
 - Perform **post-order traversal** of the right subtree
 - Visit the root node of the tree
- This is a **recursive algorithm** for performing a post-order traversal of a tree.
 - *What is the base case?*
 - *What is the recursive case?*

```
public void postorder (BinaryTreeNode<T> r)
{
    if (r != null) {
        postorder (r.getLeftChild());
        postorder (r.getRightChild());
        visit(r);
    }
}
```

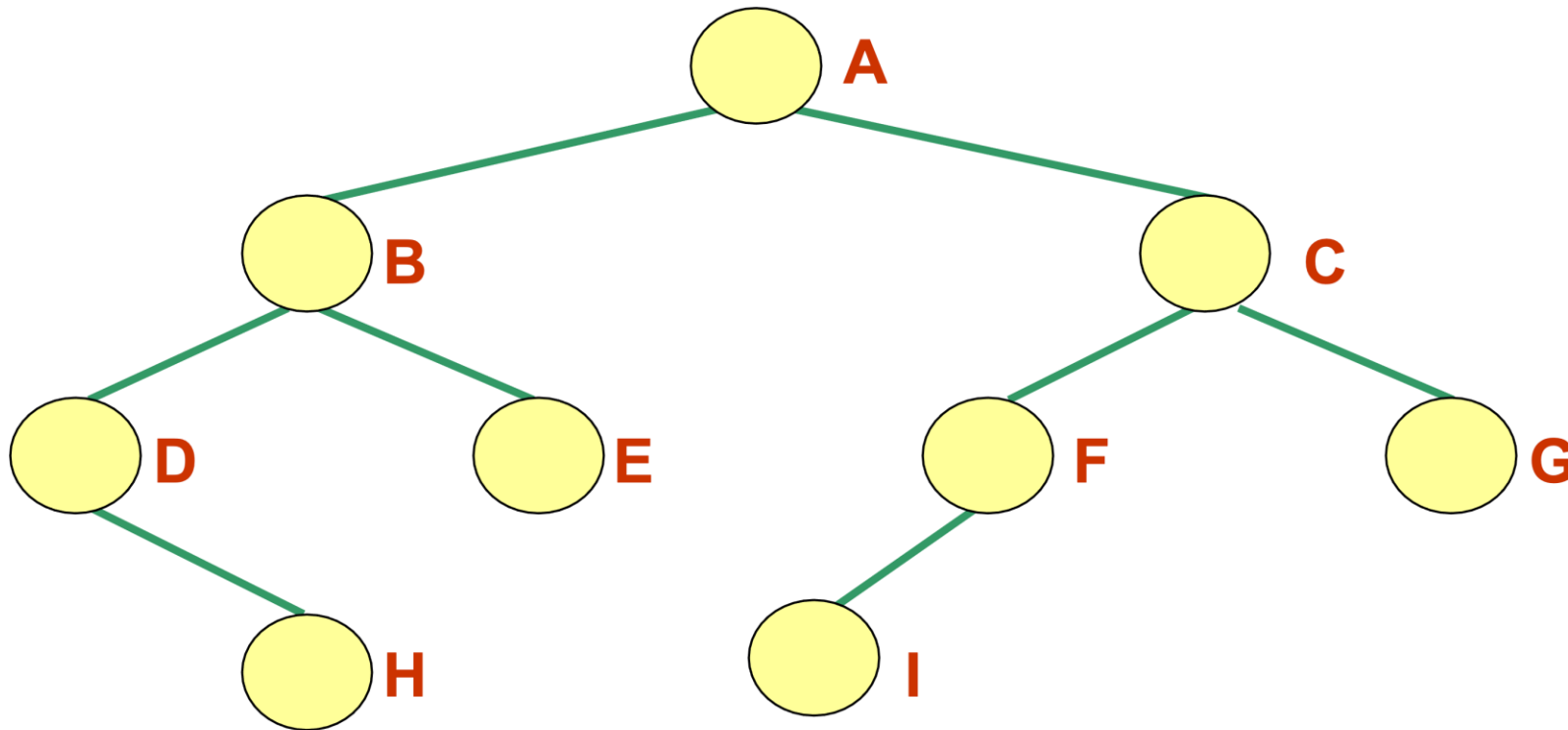
Post-order Traversal

- Node is visited after its descendants
- When is it applied?
 - Visit leaf nodes first
 - trying to delete a tree



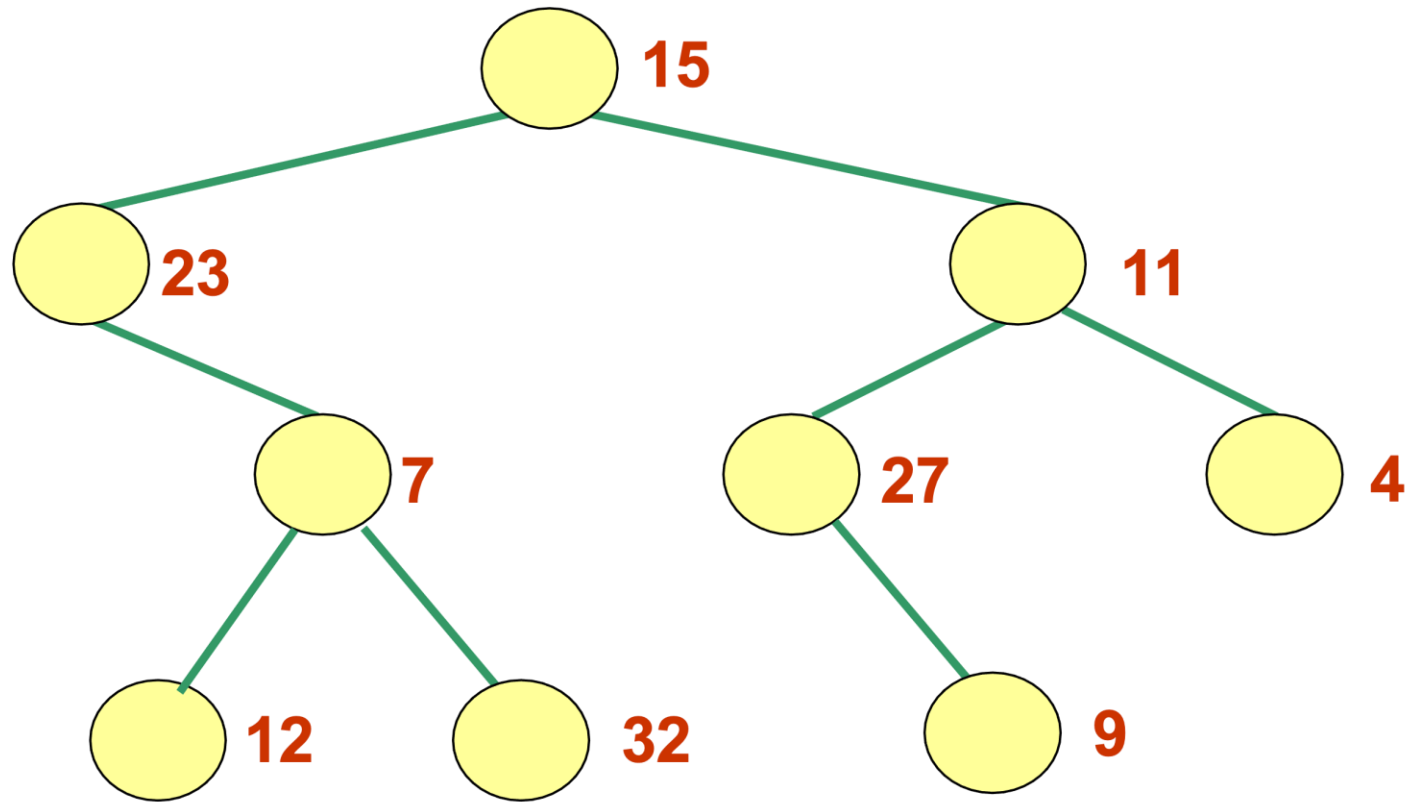
Post-order Traversal Example 1

- Post-order traversal: H D E B I F G C A



Post-order Traversal Example 2

- Post-order traversal: 12 32 7 23 9 27 4 11 15





Thank
you