

CS 1037
Fundamentals of Computer
Science II

Linked Data Structures

Ahmed Ibrahim

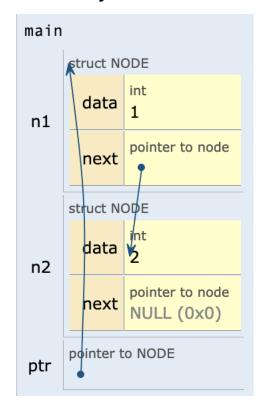
```
_modifier
  mirror object to mi
 peration == "MIRROR
mirror_mod.use_x = Tr
mirror_mod.use_y = Fa
 irror_mod.use_x = Fa
 !rror_mod.use_y = Tr
 irror_mod.use_z = Fa
  operation == "MIRRO
  rror_mod.use_x = Fa
  rror_mod.use_y = Fa
  rror_mod.use_z = Tr
  melection at the end
   ob.select= 1
   er ob.select=1
   ntext.scene.objects
  "Selected" + str(mo
    rror ob.select = 0
  bpy.context.select
   ata.objects[one.nam
  int("please select
  - OPERATOR CLASSES
  ext.active_object
```

Recall: Self-referential Structures

- Self-referential structures include a reference to data of their <u>own type</u>.
- In the struct node definition, the next variable acts as a pointer to a struct node type.
- The method of self-referential structures is the foundation for building linked data structures such as linked lists and trees.

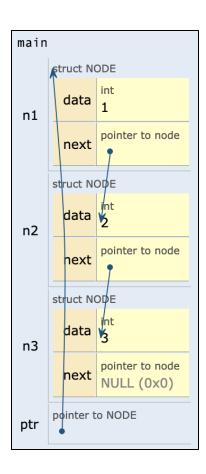
```
#include <stdio.h>
 2
   int main() {
       typedef struct node {
       int data;
 6
       struct node *next;
       } NODE;
        NODE n1, n2, *ptr = &n1;
10
11
        n1.data = 1;
        n1.next = &n2;
        n2.data = 2;
13
        n2.next = NULL;
14
        // output: 1
16
        printf("%d\n", ptr->data);
        // output: 2
20
        printf("%d\n", ptr->next->data);
21
22
        return 0;
23 }
```

Memory:



Recall: The concept of the Linked List

```
#include <stdio.h>
   int main() {
       typedef struct node {
 6
         int data;
         struct node *next;
       } NODE;
       NODE n1, n2, n3, *ptr = &n1;
       n1.data = 1;
       n1.next = &n2;
       n2.data = 2;
       n2.next = &n3;
       n3.data = 3;
       n3.next = NULL;
16
       return 0;
```



Advantages of Linked Lists

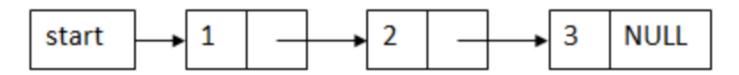
- The items do not have to be stored in consecutive memory locations; the successor can be anywhere physically.
 - So, you can insert and delete items without shifting data
 - Can increase the size of the data structure easily
- Linked lists can grow dynamically (i.e. at run time) –
 the amount of memory space allocated can grow
 and shrink as needed.

Real-World Applications

- Music and Video Playlists
- Undo/Redo Functionality in Text Editors
- Browser History
- Memory Management in Operating Systems
- Image Viewing Software
- Dynamic Data Structures in Gaming
- Simple To-do list applications

Linear Linked Data Structures

- A Linear or Singly linked list is a data structure in which each element (node) contains:
 - Data: The actual value stored in the node.
 - Next Pointer: A reference (or pointer) to the next node in the sequence.
- The last node's next pointer is set to null, indicating the end of the list.



Conceptual Diagram of a Singly-Linked List

Create a Simple ToDo List



- Imagine we want to create a simple program to represent a to-do list using a linked list, where each task is a node.
- Each node will <u>contain a description</u> of the task and a <u>reference to the next task</u> in the sequence.
- The last task will point to NULL, indicating the end of the list.

What do we need to do?

- 1. Define a Node Structure
- 2. Create a Function to Make a New Node
- 3. Add a New Task to the List
- 4. Print All Tasks in the List
- 5. Test you program

Linked List Operations

Insert Node to a Linked List

1. Create a new node:

Insert Node to a Linked List (cont.)

2. Insert the new node:

```
// New node points to the current front

tmp->next = start;

// Now the new node becomes the head

start = tmp;

tmp
```

Linked lists can grow and shrink dynamically (i.e., at run time).

Insert Node in the Middle

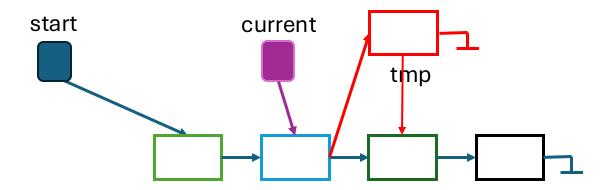
Find the Position to Insert the Node:

```
// assume that the tmp node has already been created and the position is given (#3)
    // Find the node before the desired position
   Node *current = start;
    for (int i = 0; i < position - 1 && current != NULL; i++) {
    // Traverse to the node before the insertion point
        current = current->next;
                       start
                                         current
                                current
                                                      tmp
```

Insert Node in the Middle

Insert the New Node:

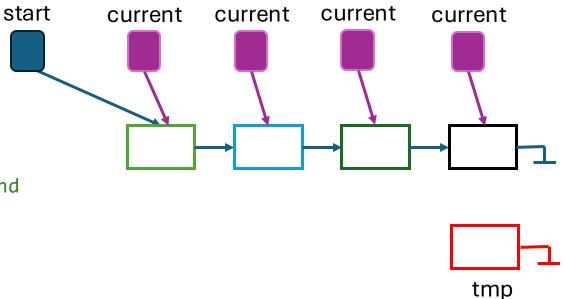
```
// If the current node is not NULL, insert the new node
   if (current != NULL) {
      tmp->next = current->next;
      current->next = tmp;
   }
```



Insert Node at the End

Find the Last Node in the List:

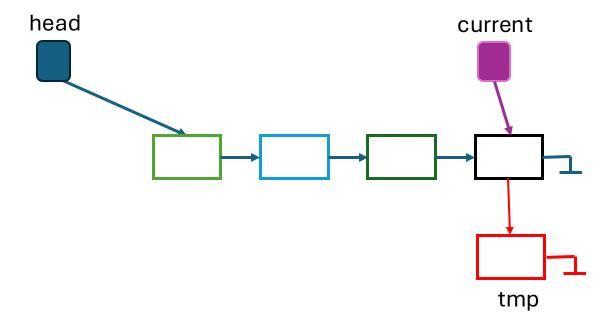
```
// Traverse to the last node in the list
Node *current = start;
while (current->next != NULL) {
// Move to the next node until we reach the end
    current = current->next;
}
```



Insert Node at the End

• Insert the New Node:

```
// Set the last node's next to the new node
  current->next = tmp;
```



Question!

What happens if we try to insert a node at a position larger than the current length of the list?

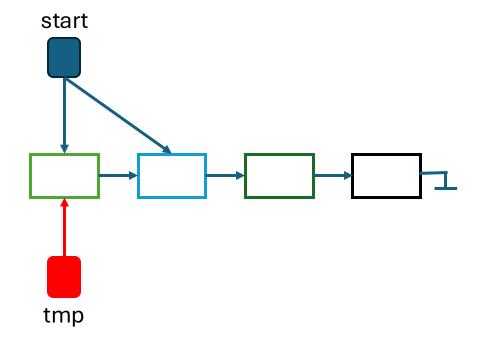
- A) Insert the node at the end of the list.
- B) Display an error message indicating the position is out of bounds.
- C) Do nothing.

What steps should we take to handle this?

Delete Node from the Front

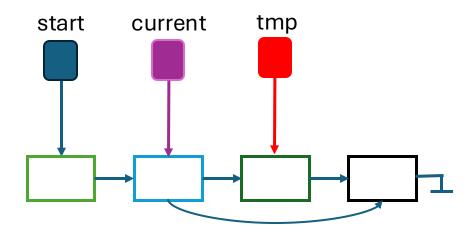
• Delete the First Node (Head):

```
// Move head to the next node
Node *tmp = start; // Store the current head
// Update head to the next node
start = start->next;
free(tmp); // Free the old head node
```



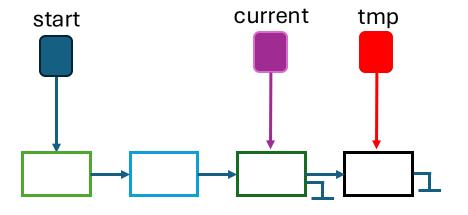
Delete Node from the Middle

```
// Traverse the list to find the node just
//before the desired position
Node *current = start;
for (int i = 0; i < position - 1 && current != NULL; i++) {
    current = current->next;}
// Delete the node at the desired position
Node *tmp = current->next;
current->next = current->next->next;
// Update the link to skip the deleted node
free(tmp); // Free the node
...
```



Delete Node from the End

```
// Traverse to the second-to-last node
Node *current = start;
while (current->next->next != NULL) {
    current = current->next;
// Free the last node and set the second-to-last
//node's next to NULL
Node *tmp = current->next;
current->next = NULL;
free(tmp);
```



CS 1037A | Fundamentals of Computer Science II

Question!

You are given the following C code snippet:

```
Node*start = createNode(10);
Start->next = createNode(20);
Start->next->next = createNode(30);
insertAtPosition(&start, 25, 2);
deleteAtPosition(&start, 1);
```

What will the final linked list look like after performing both operations, assuming the initial state and the insertion and deletion operations succeed without errors?

Memory Management in Linked Lists

- Linked lists in C often use dynamic memory allocation (e.g., malloc) to create new nodes.
- To avoid memory leaks, it is essential to deallocate memory using free() when a node is no longer needed.
- Do not call free() on the same pointer more than once; this can cause undefined behavior.

Linked List Traversal

 This operation involves visiting each node in the list and is essential for printing, searching, or processing each node's data.

```
void traverse(Node*start) {
    Node*current = start;
    while (current != NULL) {
        printf("%d\n", current->data);
        current = current->next;
    }
}
```

Searching allows you to find if a
 particular data element exists in the list.

```
int search(Node*start, char*data) {
   Node *current = start;
   while (current != NULL) {
   // Compare the current node's
   // data with the target data
       if (strcmp(current->data, data) == 0) {
       return 1; // Data found }
       // Move to the next node
       current = current->next;}
       return 0; // Data not found
}
```

Doubly Linked Data Structures

- A doubly linked list is a type of linked list where each node contains three fields:
 - **Data**: The value stored in the node.
 - **Next Pointer**: A reference to the next node in the sequence.
 - **Previous Pointer**: A reference to the previous node in the sequence.
- The first node's previous pointer is NULL, and the last node's next pointer is NULL,
 signifying the list's boundaries.

 Memory

