#### **NAME**

statistics - Shore Storage Manager performance information

### **SYNOPSIS**

```
class sm_stats_info_t;
static rc_t
ss_m::gather_stats(sm_stats_info_t& stats, bool reset=false);
```

### DESCRIPTION

The Shore Storage Manager keeps myriad statistics. The statistics are available to higher layers of software through the method **gather\_stats**. Each statistic kept has a name and a manifest constant of that name for use with the generic statistics-gathering programming interface described in **statistics(fc)**. The names and descriptions of the statistics are listed in sections below.

All the statistics are unsigned long integers except where noted below. Below each statistic is identified by the manifest constant for the statistic and also by its member name in the structure *sm\_stats\_info\_t*.

After gathering the storage manager statistics into a *sm\_stats\_info\_t*, the *sm\_stats\_info\_t* can be used with the *w\_statistics\_t* package for display:

```
sm_stats_info_t s;

rc = gather_stats(s);
if(!rc) {
    w_statistics_t stats << s;
    cout << stats;
}</pre>
```

Set the argument *reset* of **gather\_stats** to **true** if you want the statistics to be cleared (set to zero) after they are copied.

# **BUFFER POOL STATISTICS**

```
SM_rec_pin_cnt
```

*sm\_stats\_info\_t::rec\_pin\_cnt*. Records pinned in the buffer pool. (This counts the pinning events, not the distinct records pinned.)

```
SM_rec_unpin_cnt
```

```
sm_stats_info_t::rec_unpin_cnt. Records unpinned.
```

```
SM_page_fix_cnt
```

sm\_stats\_info\_t::page\_fix\_cnt. Times pages were fixed in the buffer pool.

```
SM_page_refix_cnt
```

*sm\_stats\_info\_t::page\_refix\_cnt*. Times pages were refixed in the buffer pool; refixing is cheaper than fixing, because it avoids a hash-table look-up.

```
SM_page_unfix_cnt
```

```
sm_stats_info_t::page_unfix_cnt. Times pages were unfixed.
```

The following group of statistics describes the operation of the buffer manager, which groups consecutive pages for writing to disk whenever possible.

```
SM_bf_one_page_write
```

*sm\_stats\_info\_t::bf\_one\_page\_write*. The number of times a single page was flushed from the buffer pool.

## SM\_bf\_two\_page\_write

*sm\_stats\_info\_t::bf\_two\_page\_write*. The number of times two consecutive pages were flushed from the buffer pool in one I/O request.

statistics (ssm)

### SM\_bf\_three\_page\_write

sm\_stats\_info\_t::bf\_three\_page\_write.

### SM\_bf\_four\_page\_write

sm\_stats\_info\_t::bf\_four\_page\_write.

## SM\_bf\_five\_page\_write

sm\_stats\_info\_t::bf\_five\_page\_write.

## SM\_bf\_six\_page\_write

*sm\_stats\_info\_t::bf\_six\_page\_write.* 

### SM\_bf\_seven\_page\_write

sm\_stats\_info\_t::bf\_seven\_page\_write.

# SM\_bf\_eight\_page\_write

*sm\_stats\_info\_t::bf\_eight\_page\_write.* 

# SM\_cleaner\_sweeps

*sm\_stats\_info\_t::bf\_cleaner\_sweeps*. The number of times the buffer-cleaner thread swept the buffer pool (a clock algorithm is used).

# SM\_bf\_log\_flush\_all

sm\_stats\_info\_t::bf\_log\_flush\_all. The number of times the entire log was flushed by the buffer manager.

# SM\_bf\_log\_flush\_lsn

*sm\_stats\_info\_t::bf\_log\_flush\_lsn*. The number of times the log was flushed up to a specific log sequence number by the buffer manager.

# SM\_bf\_write\_out

sm\_stats\_info\_t::bf\_write\_out. The number of calls to the buffer manager method that writes pages to disk.

# SM\_bf\_replace\_out

*sm\_stats\_info\_t::bf\_replace\_out*. The number of times a page was written in order to free the frame (during replacement) for a different page.

# SM replaced dirty

sm\_stats\_info\_t::bf\_replaced\_dirty. The number of times the frame replaced contained a dirty page.

### SM\_replaced\_clean

sm\_stats\_info\_t::bf\_replaced\_clean. The number of times the frame replaced contained a clean page.

### SM\_await\_clean

*sm\_stats\_info\_t::bf\_await\_clean*. The number of times a page fix request awaited a frame to become clean by the cleaner thread, rather than forcing out a dirty page.

# SM\_cleaner\_signalled

*sm\_stats\_info\_t::bf\_cleaner\_signalled*. The number of times the cleaner was restarted by a signal from another thread. The following three counters distinguish the reasons for the signals.

# SM\_kick\_replacement

*sm\_stats\_info\_t::bf\_kick\_replacement*. A dirty page was replaced.

### SM\_kick\_full

*sm\_stats\_info\_t::bf\_kick\_full*. The buffer pool was full of dirty pages, noticed sometime other than on page replacement.

# SM\_kick\_threshhold

sm\_stats\_info\_t::bf\_kick\_threshhold. The buffer pool's dirty-page threshhold was met.

## SM\_sweep\_page\_hot

sm\_stats\_info\_t::bf\_sweep\_page\_hot. The buffer cleaner swept over a hot page, and left it in place.

### **B-TREE STATISTICS**

### SM\_bt\_find\_cnt

*sm\_stats\_info\_t::bt\_find\_cnt*. B-tree lookups (calls to ss\_m::find\_assoc());

#### SM bt insert cnt

sm\_stats\_info\_t::bt\_insert\_cnt. B-tree inserts (calls to ss\_m::create\_assoc());

### SM\_bt\_remove\_cnt

sm\_stats\_info\_t::bt\_remove\_cnt. B-tree removes (calls to ss\_m::destroy\_assoc());

#### SM bt scan ent

*sm\_stats\_info\_t::bt\_scan\_cnt*. Number of B-tree scans started.

### SM\_bt\_splits

*sm\_stats\_info\_t::bt\_splits*. B-tree pages split (interior and leaf).

### SM\_bt\_cuts

*sm\_stats\_info\_t::bt\_cuts*. B-tree pages removed (interior and leaf).

### SM\_bt\_grows

*sm\_stats\_info\_t::bt\_grows*. Times B-tree grew a level.

#### SM bt shrinks

sm\_stats\_info\_t::bt\_shrinks. Times B-tree shrunk a level.

### SM\_bt\_links

*sm\_stats\_info\_t::bt\_links*. Times B-tree sibling links were followed (while a structure modification operation was not yet propagated.)

## SM\_bt\_clr\_smo\_traverse

*sm\_stats\_info\_t::bt\_clr\_smo\_traverse*. Times SMO (structure-modification-operation-in-progress) bits cleared on traverse.

### SM\_bt\_posc

sm\_stats\_info\_t::bt\_posc. Times awaited a POSC (point of structural consistency).

### SM\_bt\_traverse\_cnt

*sm\_stats\_info\_t::bt\_traverse\_cnt*. Times b-trees were traversed from the root.

# SM\_bt\_upgrade\_fail\_retry

sm\_stats\_info\_t::bt\_upgrade\_fail\_retry. Failure to upgrade a latch without waiting caused a re-try.

# LOGICAL-ID STATISTICS

These will be zero if you do not use logical-IDs. Logical IDs are stored in a B-tree index, and recently used IDs are cached in a transient cache.

## SM lid lookups

sm\_stats\_info\_t::lid\_lookups. Times the logical-ID index was searched for an ID.

# SM\_lid\_remote\_lookups

sm\_stats\_info\_t::lid\_remote\_lookups. Times the index was searched for a second time for indirect references.

### SM lid inserts

*sm\_stats\_info\_t::lid\_inserts*. Logical-IDs added to the index.

#### SM lid removes

*sm\_stats\_info\_t::lid\_removes*. Logical-IDs removed from the index.

## SM\_lid\_cache\_hits

*sm\_stats\_info\_t::lid\_cache\_hits*. Times a logical-ID look-up request was satisfied by the cache of recent requests, avoiding a B-tree look-up.

## PAGE & EXTENT OPERATIONS

### SM\_page\_alloc\_cnt

sm\_stats\_info\_t::page\_alloc\_cnt. Pages allocated from free pages in allocated extents.

## SM\_page\_dealloc\_cnt

sm\_stats\_info\_t::page\_dealloc\_cnt. Pages deallocated. These pages are free for re-allocation in the same store.

### SM\_ext\_lookup\_hits

*sm\_stats\_info\_t::ext\_lookup\_hits*. Cache hits during extent-lookups.

### SM\_ext\_lookup\_hits

sm\_stats\_info\_t::ext\_lookup\_hits. Cache misses during extent-lookups.

# TRANSACTION STATISTICS

### SM\_begin\_xct\_cnt

*sm\_stats\_info\_t::begin\_xct\_cnt*. The number of transactions that were begun. This includes the number that resulted from chaining transactions.

# SM\_commit\_xct\_cnt

*sm\_stats\_info\_t::commit\_xct\_cnt*. The number of transactions that were committed. This includes the number that resulted from chaining transactions.

### SM abort xct cnt

sm stats info t::abort xct cnt. The number of transactions that were aborted.

# SM\_rollback\_savept\_cnt

*sm\_stats\_info\_t::rollback\_savept\_cnt*. The number of requests to roll back to a save point, without rolling back the entire transaction.

# SM\_mpl\_attach\_cnt

sm\_stats\_info\_t::mpl\_attach\_cnt. Times a thread attached to a transaction to which at least one other thread was already attached. This is for value-added servers that run transactions with multiple threads in parallel. (This is not a supported feature, as the circumstances in which this can be done are few.)

## SM anchors

*sm\_stats\_info\_t::anchors*. Times a transaction grabbed an anchor in the log (for the purpose of compensating around a top-level action).

### SM\_compensate\_in\_log

*sm\_stats\_info\_t::compensate\_in\_log*. Times a transaction wrote a compensation by snooping into the log.

### SM\_compensate\_in\_xct

sm\_stats\_info\_t::compensate\_in\_xct. Times a transaction wrote a compensation in its own log buffer.

### SM\_compensate\_records

sm\_stats\_info\_t::compensate\_records. Number of compensations-only log records written.

#### PARALLELISM STATISTICS

### SM\_await\_io\_monitor

sm\_stats\_info\_t::await\_io\_monitor. Times blocked on mutex for serializing access to the I/O monitor.

### LOCK STATISTICS

The following statistics measure activity in the lock table. Many of the statistics are of interest only the the developers of Shore, and are likely to be removed in future releases of Shore.

# SM\_unlock\_request\_cnt

*sm\_stats\_info\_t::unlock\_request\_cnt*. High-level unlock requests.

## SM\_lock\_request\_cnt

*sm\_stats\_info\_t::lock\_request\_cnt*. High-level lock requests (could have been satisfied with the lock cache or with the lock table.)

## SM\_lock\_cache\_hit\_cnt

*sm\_stats\_info\_t::lock\_cache\_hit\_cnt*. The number of cache hits (avoiding request to acquire locks through the lock table).

# SM\_lock\_acquire\_cnt

*sm\_stats\_info\_t::lock\_acquire\_cnt*. The number of times a request was made to acquire a lock through the lock table (as opposed to the lock cache).

# SM lock head t cnt

*sm\_stats\_info\_t::lock\_head\_t\_cnt*. The number of *lock\_head\_t* structures put in the table.

# SM\_lock\_request\_t\_cnt

*sm\_stats\_info\_t::lock\_request\_t\_cnt.* The number of *lock\_request\_t* structures chained from a *lock\_head\_t*.

### SM\_lock\_query\_cnt

sm\_stats\_info\_t::lock\_query\_cnt. The number of times the lock table was queried about a given lock.

The following statistics distinguish the locks acquired by lock-id:

## SM\_lk\_vol\_acq

*sm\_stats\_info\_t::lk\_vol\_acq*. Volume locks.

#### SM\_lk\_store\_acq

*sm\_stats\_info\_t::lk\_store\_acq*. Store locks. Stores are below volumes in the lock hierarchy.

### SM lk page acq

sm\_stats\_info\_t::lk\_page\_acq. Page locks. Pages are below stores in the lock hierarchy.

# SM\_lk\_kvl\_acq

*sm\_stats\_info\_t::lk\_kvl\_acq*. Key-value locks. Key-value pairs are below pages in the lock hierarchy.

### SM\_lk\_rec\_acq

sm\_stats\_info\_t::lk\_rec\_acq. Record locks. Records are below pages in the lock hierarchy.

### SM\_lk\_ext\_acq

sm\_stats\_info\_t::lk\_ext\_acq. Extent locks. Extents are not in the lock hierarchy.

### SM lock deadlock cnt

*sm\_stats\_info\_t::lock\_deadlock\_cnt*. Deadlocks detected by local deadlock detector.

# SM\_lock\_esc\_to\_page

sm\_stats\_info\_t::lock\_esc\_to\_page. Lock escalations from record to page.

#### SM\_lock\_esc\_to\_store

*sm\_stats\_info\_t::lock\_esc\_to\_store*. Lock escalations from page to store.

# SM\_lock\_esc\_to\_volume

*sm\_stats\_info\_t::lock\_esc\_to\_volume*. Lock escalations from store to volume.

The following statistics measure collisions from the lock table hash function. They are generally not of interest to the user, but they might be of interest to someone who is building a value-added server and might consider using a different hash function. Only buckets of non-zero length are counted when these statistics are computed. They are all zero when there is no transaction active, since the lock table is empty at that time. These statistics are computed when **gather\_stats** is called; it traverses the entire lock table, so it should not be used habitually when transactions are active.

#### SM lock conversion cnt

sm\_stats\_info\_t::lock\_conversion\_cnt. Lock requests requiring a conversion of the lock mode.

# SM\_lock\_extraneous\_req\_cnt

sm\_stats\_info\_t::lock\_extraneous\_req\_cnt. Requests already granted.

sm locktablesize.

# SM\_lock\_max\_bucket\_len

*sm\_stats\_info\_t::lock\_max\_bucket\_len*. The largest bucket in use.

## SM\_lock\_min\_bucket\_len

*sm\_stats\_info\_t::lock\_min\_bucket\_len*. The smallest bucket in use.

## SM\_lock\_mode\_bucket\_len

*sm\_stats\_info\_t::lock\_mode\_bucket\_len*. The mode of the lengths of the buckets used.

# SM\_lock\_mean\_bucket\_len

float sm\_stats\_info\_t::lock\_mean\_bucket\_len. The mean bucket length (of the buckets used).

# SM\_lock\_var\_bucket\_len

float sm\_stats\_info\_t::lock\_var\_bucket\_len. The variance of the bucket length (of the buckets used).

#### SM\_lock\_std\_bucket\_len

float sm\_stats\_info\_t::lock\_std\_bucket\_len. The Standard deviation of the bucket length (of the buckets used).

### OPERATIONS ON LOCAL DATA VOLUMES

## SM\_vol\_reads

*sm\_stats\_info\_t::vol\_reads*. Lowest-level read requests made to the 'diskrw' process, which effects non-blocking disk I/O for data volumes.

## SM\_vol\_writes

*sm\_stats\_info\_t::vol\_writes*. Lowest-level write requests made to the 'diskrw' process, which effects non-blocking disk I/O for data volumes.

## SM\_vol\_blks\_written

*sm\_stats\_info\_t::vol\_blks\_written*. Data volume pages written to disk.

### **OPERATIONS ON THE LOG**

### SM\_log\_records\_generated

*sm\_stats\_info\_t::log\_records\_generated*. The number of log records written.

## SM\_log\_bytes\_generated

*sm\_stats\_info\_t::log\_bytes\_generated*. The number of bytes written to the log.

### SM\_log\_sync\_cnt

*sm\_stats\_info\_t::log\_sync\_cnt*. The number of times the log was flushed to disk.

## SM\_log\_dup\_sync\_cnt

*sm\_stats\_info\_t::log\_dup\_sync\_cnt*. The number of times the log was flushed superfluously (for debugging).

# SM\_log\_fsync\_cnt

*sm\_stats\_info\_t::log\_fsync\_cnt*. The number of times the fsync(2) system call was used to flush a log that is a Unix file.

## SM\_log\_sync\_nrec\_max

sm\_stats\_info\_t::log\_sync\_nrec\_max. Maximum number log records buffered between flushes.

### SM\_log\_sync\_nbytes\_max

sm\_stats\_info\_t::log\_sync\_nbytes\_max. Maximum number bytes buffered between log flushes.

## SM\_log\_chkpt\_cnt

*sm\_stats\_info\_t::log\_chkpt\_cnt*. Number of checkpoints taken.

## **MISCELLANEOUS**

# SM\_idle\_yield\_return

sm\_stats\_info\_t::idle\_yield\_return. Times the idle thread returned from yield(). (For debugging.)

#### SM\_idle\_wait\_return

Times the idle thread returned from wait(). (For debugging.) sm\_stats\_info\_t::idle\_wait\_return.

#### VERSION

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## SEE ALSO

options(svas), statistics(svas), and statistics(fc)