

## Review

### ER Model -

- you should be able to look at and understand
- you should be able to draw one from a conceptual description
- 3 shapes to draw

Entities 

Relationships 

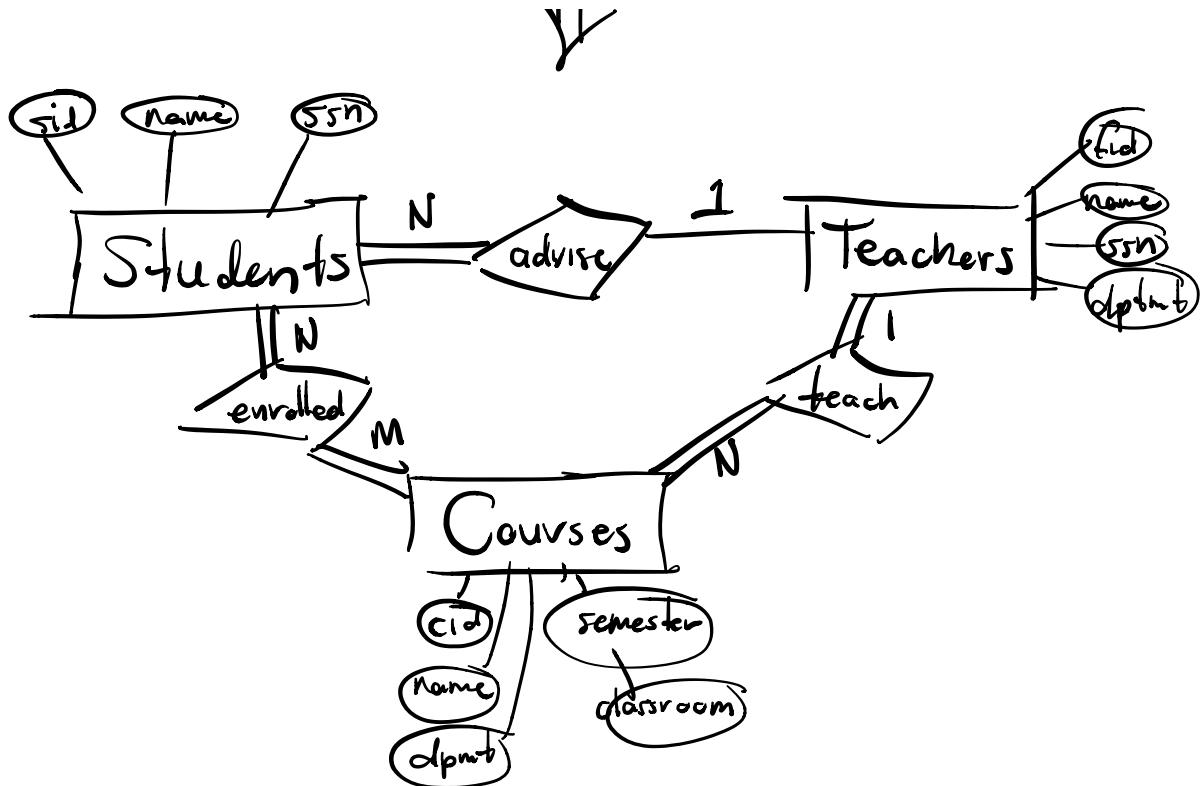
Attributes 

- the "arity" of relationships (1:1, 1:N, N:1, N:M)  
↳ should be able to understand and label

- participation constraints  
= (total) vs — (partial)

ex We want to draw an ER diagram for Students, teachers, and the courses they participate in:





### Relational Model -

- you should be able to convert from ER model to a relational model (and back)
- to express a relational model, just give the schema:
  - ex: student (sid, name, ssn, advisor-fid)
  - teachers (fid, name, ssn, dept)
  - courses (cid, name, dept, semester, classroom)
  - teach (fid, cid, how-long)
  - enrolled (sid, cid)
- you should be able to identify:  
Superkeys - set of attributes that can uniquely identify a tuple in a relation

Candidate keys - a minimal superkey (a superkey where no subset is a superkey)

Primary keys - know that a primary key is some selected candidate key

- know that a "relation" is a set of tuples
- tuples  $\approx$  rows
- relations  $\approx$  tables
- attributes  $\approx$  columns

Relational Algebra -

$\sigma_p$  - Selection operator

- selects tuples that match predicate 'p'

$\pi_{\text{attrs}}$  - Projection operator

- limits results to only the attributes in the 'attrs' list

$\times$  - Cartesian product operator

- match every tuple in one relation to every tuple in the other and results in a relation whose attributes are

the combined attributes of both participating relations

- $\cup$  - Union operator
- results in all tuples from both relations
  - can only be applied on relations w/ the exact same set of attributes

- $-$  - difference operator
- $(R_1 - R_2)$  results in tuples in  $R_1$  but not in  $R_2$
  - also only applies for relations w/ same attrs.

- $\cap$  - Intersection operator
- results in tuples that exist in both relations
  - also only applies for relations w/ same attrs.

- $\bowtie_\theta$  - Theta join operator
- combination of  $\times$  and  $\sigma_\theta$
  - $R_1 \bowtie_\theta R_2 = \sigma_\theta(R_1 \times R_2)$

- $\bowtie_{\text{attrs}}$  - Equijoin operator
- join on only the attributes in 'attrs'
  - ex: Student  $\bowtie_{\text{sid}} \text{Enrolled}$   
Student  $\bowtie_{s.sid=e.sid} \text{Enrolled}$

- Natural join
- join on all matching attributes between two relations
- ex. Student  $\bowtie$  Teacher

||

Student  $\bowtie$ <sub>s.name = t.name and s.ssn = t.ssn</sub> Teacher

## SQL Queries -

### SELECT select-list

- same as  $\Pi$  from relational algebra
- narrows down columns
- instead of select-list you can put \* to get all columns

### [DISTINCT]

- goes after SELECT to remove duplicate rows in the result

### FROM from-list

- same as  $\times$  operator between all tables in from-list
- know how to abbreviate table names  
ex: (table1 f1, table2 f2 ... etc)

- know how to use those abbreviations elsewhere (dot notation)  
ex:  $(f1.attr \dots f2.attr)$
- NOTE: remember that this only does cartesian product, so to perform a join, you need to include constraints in the WHERE clause

## WHERE conditions

- same as the  $\sigma$  selection operator
- selects only rows that match 'conditions'
- you can build complex expressions in the WHERE clause
- ex:  $(\text{and}, \text{OR}, \text{or}, \text{||}, =, !=, <, >, \text{etc})$   
 $(+, -, *, /, \text{CONCAT}(a, b))$

## $\langle \text{expr} \rangle \text{ AS } \langle \text{column-name} \rangle$

- goes inside the SELECT clause and lets you use expressions to overwrite columns in the result, or create new columns

ex:  $(\text{SELECT rating} + 2 \text{ AS rating\_plus_2},$   
 $\text{rating}/2 \text{ AS half\_rating}\dots)$

- again, the 'expr' can use any of the basic operators as with the WHERE clause

`<string-column> LIKE <string-matching-expression>`

- goes in the WHERE clause to do string matching beyond simple equality (=)
- we learned two matching characters for the matching expressions:

`%` - matches 0 or more arbitrary characters  
`_` - matches exactly 1 arbitrary character

ex:

`( WHERE name LIKE "B%" )`

- Not case-sensitive

`<query> UNION <query>`  
`<table> <table>`

- this goes between two full queries (or something resulting in a table) and gives the union ( $\cup$ ) of the results.

ex:

`( SELECT ... FROM .... WHERE ... )`  
`UNION [ALL]`  
`SELECT ... FROM ... WHERE ... )`

`[ALL]` - if you want to retain duplicates

## Example problems

Schema: Sailors (sid, name, rating)  
boats (bid, name, color)  
reserv (sid, bid, date)

1) Find names of sailors who have reserved boat #2?

$\pi_{\text{name}} \left( \sigma_{\text{bid} = 2} \left( \sigma_{\text{s.sid} = \text{r.sid}} (\text{Sailors} \times \text{Reserv}) \right) \right)$

$\pi_{\text{name}} \left( \sigma_{\text{bid} = 2} (\text{Sailors} \bowtie \text{Reserv}) \right)$

$\pi_{\text{name}} \left( \text{Sailors} \bowtie_{\text{s.sid} = \text{r.sid} \wedge \text{r.bid} = 2} \text{Reserv} \right)$

SELECT name  
FROM Sailors, reserv  
WHERE sailors.sid = reserv.sid and reserv.bid = 2;

}

SELECT s.name  
FROM Sailors s, reserv r  
WHERE s.sid = r.sid and r.bid = 2;

2) Find names of sailors who have reserved a blue boat?

$\pi_{s.name} \left( \sigma_{b.color = \text{Blue}} \left( \sigma_{s.sid = r.sid \wedge b.bid = r.bid} \left( \text{Sailors} \times \text{Reserv} \times \text{Boats} \right) \right) \right)$

$\pi_{s.name} \left( \sigma_{b.color = \text{Blue}} \left( \text{Sailors} \underset{s.id}{\bowtie} \text{Reserv} \underset{r.bid}{\bowtie} \text{Boat} \right) \right)$

SELECT s.name  
FROM Sailors s, reserv r, boats b  
WHERE b.color = "Blue" AND s.sid = r.sid  
AND b.bid = r.bid;

---

3) Find colors of boats reserved by sailors named "Christie Weeks"?

$\pi_{b.color} \left( \sigma_{s.name = \text{"Christie Weeks"}} \left( \sigma_{s.sid = r.sid \wedge r.bid = b.bid} \left( \text{Sailors} \times \text{Reserv} \times \text{Boats} \right) \right) \right)$

SELECT b.color  
FROM Sailors s, reserv r, boats b  
WHERE s.name = "Christie" AND s.sid = r.sid AND  
"Weeks" r.bid = b.bid;

SELECT b.color

From sailors s INNER JOIN reserv r ON s.sid =  
INNER JOIN boats b ON r.bid =

WHERE s.name = "Christie Weeks"; b.bid

---

4) Find the names of sailors who have reserved at least one boat?

$\pi_{s.name} (\text{Reserv} \bowtie_{sid} \text{Sailors})$

SELECT s.name  
FROM reserv r, sailors s  
WHERE r.sid = s.sid;

---

5) Find names of sailors who reserved both a red and a green boat?

$\pi_{s.name} \left[ \left( \text{Sailors} \bowtie_{sid} \text{reserv} \bowtie_{bid} \left( \sigma_{\text{color}=\text{Red}}(\text{boats}) \right) \right) \cap \left( \text{Sailors} \bowtie_{sid} \text{reserv} \bowtie_{bid} \left( \sigma_{\text{color}=\text{green}}(\text{boats}) \right) \right) \right]$

SQL?  $\Rightarrow$  we don't know how to do intersect  $\cap$  yet  
... so don't try