Theory of Computation

MIEIC, 2nd Year

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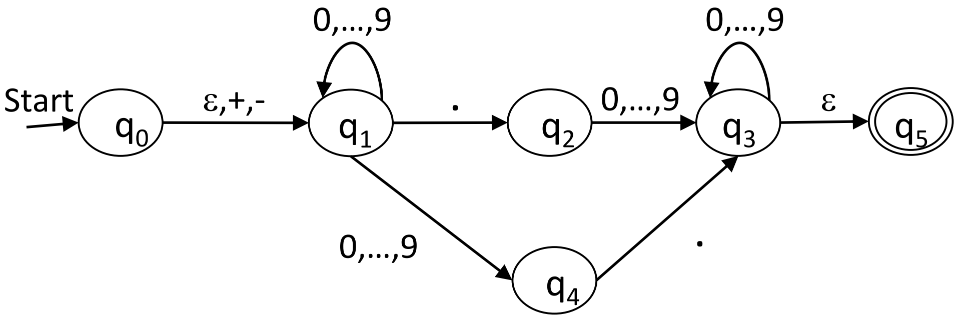
Outline

- Non-Deterministic Finite Automata with ε transitions (ε-NFAs)
- Conversion of ε-NFAs into DFAs

Finite Automata with ε Transitions

- Example: ε-NFA which recognizes decimal numbers
  - Signal + or – optional
  - Sequence of digits
  - A decimal point
  - Another sequence of digits (At least one of the sequences of digits is non-empty)



Exercise 6

- Modify the previous state diagram in order to unrecognize inputs like: 3, .5, and -.1; and to recognize inputs like: 1.
- More precise, this new definition of a decimal number is:
  - Signal + or – optional
  - A sequence of digits with length greater or equal 1
  - A decimal part consisting of a '.' followed by an optional sequence of digits x, such that $|x| \geq 0$.

Formal Notation ε-NFA

- ε-NFA $E = (Q, \Sigma, \delta, q_0, F)$
  - The major difference is in the transition function δ to deal with ε
    - $\delta(q, a)$: state $q \in Q$ and $a \in \Sigma \cup \{\varepsilon\}$
- Example: $E = (\{q_0, q_1, q_2, q_3, q_4, q_5\}, \{.,+,-,0,..,9\}, \delta, q_0, \{q_5\})$
- The symbol representing the empty-string, ε, is not visible in the sequence of digits
  - It represents spontaneous transitions
  - We deal with it in the same way as with the non-determinism, i.e., considering that the automaton can be in all the states before and after the ε transition
- To know which are the states we can reach from a state q with ε, we calculate the ε-close(q)
  - ε-close(q_0)= $\{q_0,q_1\}$; ε-close(q_3)= $\{q_3,q_5\}$

δ	ε	+,-	.	0,..,9
$\rightarrow q_0$	$\{q_1\}$	$\{q_1\}$	$\varnothing$	$\varnothing$
q_1	$\varnothing$	$\varnothing$	$\{q_2\}$	$\{q_1\ q_4\}$
q_2	$\varnothing$	$\varnothing$	$\varnothing$	$\{q_3\}$
q_3	$\{q_5\}$	$\varnothing$	$\varnothing$	$\{q_3\}$
q_4	$\varnothing$	$\varnothing$	$\{q_3\}$	$\varnothing$
$*q_5$	$\varnothing$	$\varnothing$	$\varnothing$	$\varnothing$

Extended Transitions

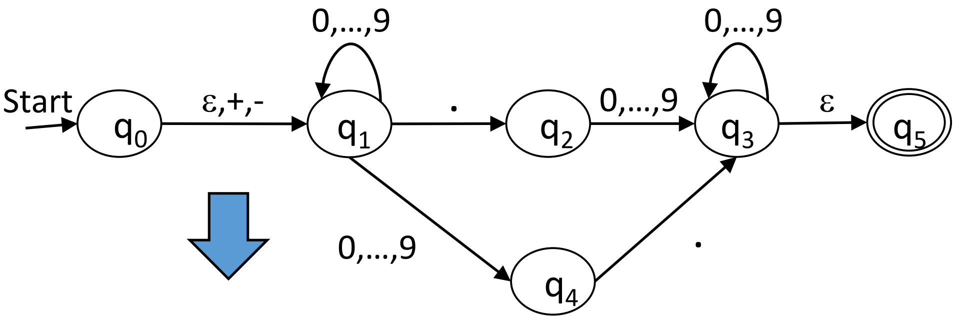
- ε-*close(q) or Eclose(q)*
  - Basis: State *q* is in *EClose(q)*
  - Induction: if *p* is in *EClose(q)* and exists an ε transition from *p* to *r* com with label ε, then *r* is also in *EClose(q)*
- Extended transition $\hat{\delta}$
  - Basis: $\hat{\delta}(q,\varepsilon)$ =*EClose(q)*
  - Induction: *w=xa*, $a \in \Sigma$ (so, $a \neq \varepsilon$)
    - 1. let's $\hat{\delta}(q,x)=\{p_1, p_2, \ldots, p_k\}$
    - 2. $\bigcup_{i=1}^{k} \delta(p_{i,} a) = \{r_1,\ldots,r_m\}$
    - 3. $\hat{\delta}(q,w) = \bigcup_{j=1}^{m} Eclose(r_j)$
  - (1) gives the states reached from *q* following a path representing **x** that can include (and/or terminate) one or more ε

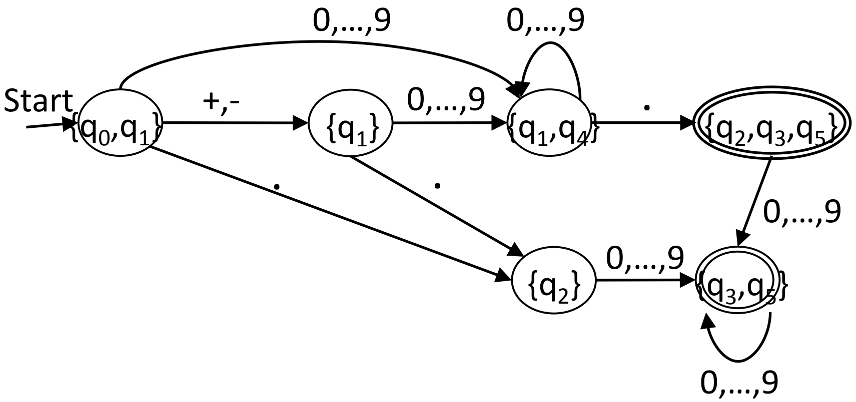
Eliminating ε Transitions

- Given an ε-NFA E there exists always an equivalent DFA D
  - E and D accept the same language
- Technique of subsets construction
  - ε-NFA E = $(Q_E, \Sigma, \delta_E, q_0, F_E)$ → DFA D = $(Q_D, \Sigma, \delta_D, q_D, F_D)$
- Q_D is the set of subsets of Q_E closed in ε
  - S= EClose(S)
- State of start: q_D = EClose(q_0)
- $F_D = \{S \mid S$ is in Q_D and $S \cap F_E \neq \emptyset\}$
- Transition $\delta_D(S,a)$, with a in Σ and S in Q_D
  - $S=\{p_1, p_2, \ldots, p_k\}$
  - Calculate $$\bigcup_{i=1}^{k} \delta_E(p_i, a) = \{r_1, \ldots, r_m\}$$
  - Terminate with $$\delta_D(S,a) = \bigcup_{j=1}^{m} EClose(r_j)$$

Example of the Recognizer of Decimals

- Convert to a DFA:





Exercise 7

- Consider the following ε-NFA:

	ε	a	b	c
→p	∅	{p}	{q}	{r}
q	{p}	{q}	{r}	∅
*r	{q}	{r}	∅	{p}

- Calculate the ε-close for each state
- Indicate all the strings with length ≤ 3 accepted by the automaton
- Convert the NFA into an ε-DFA