# 打

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# 字符串

### KMP 模版

```
//普通
void getnext1(char *s)
    int i = 0, j = -1, len = strlen(s);
    ne[0] = -1;
    while (i < 1en)
        if(j == -1 \mid \mid s[i] == s[j]) ne[++i] = ++j;
        else j = ne[j];
//加快预处理
void getnext2(char *s)
    int i = 0, j = -1, len = strlen(s);
    ne[0] = -1;
    while (i < 1en)
        if(j == -1 || s[i] == s[j])
            if(s[++i] == s[++j])
                                  ne[i] = ne[j];
            else ne[i] = j;
        else j = ne[j];
//返回 x 在 y 中出现的次数, 可以重叠
int kmp(char *x, char *y)
    getnext1(x);
    int i = 0, j = 0, ans = 0, leny = strlen(y), lenx = strlen(x);
    while (i < 1eny)
        if(j == -1 || x[j] == y[i])
            i++, j++,
            if(j == 1enx)
                ans++;
                j = ne[j];
                j = ne[j];
        else
    return ans;
```

### 浮动匹配

```
//求浮动匹配
//1[i]表示 i 结点之前(包括 i 结点),小于等于 a[i]的个数
//1e[i]表示 i 结点之前(包括 i 结点),小于等于 a[i]的个数
//1[i] == 1[j] && 1e[i] == 1e[j],则匹配
int n,k,s,a[100005],b[25005],ne[25005],tree[30],1[100005],1e[100005];
inline int lowbit(int x)
{
    return x&-x;
}
void update(int pos,int x)
```

```
while (pos \langle = s \rangle
         tree[pos] += x;
         pos += lowbit(pos);
int getsum(int pos)
    int sum = 0;
    while (pos)
        sum += tree[pos];
        pos -= lowbit(pos);
    return sum;
void getnext1()
    memset (tree, 0, sizeof (tree));
    int i = 0, j = -1;
    ne[0] = -1;
    while (i < k)
        if(j == -1 \mid | getsum(b[i]-1) == 1[j] \&\& getsum(b[i]) == 1e[j])
             ne[++i] = ++j;
             if(i < k) update(b[i], 1);
         else
             for(int t = i-j; t < i-ne[j]; t++) update(b[t],-1);
             j = ne[j];
vector<int> kmp()
    getnext1();
    memset (tree, 0, sizeof (tree));
    vector(int) ans;
    int i = 0, j = 0;
    update (a[0], 1);
    while (i < n)
         if(j == -1 \mid | getsum(a[i]-1) == 1[j] \&\& getsum(a[i]) == 1e[j])
             i++, j++;
             if(i < n)
                          update (a[i], 1);
             if(j == k)
                  ans. push back (i-k+1);
                  for (int \ \overline{t} = i-j; t < i-ne[j]; t++) update (a[t], -1);
                  j = ne[j];
         else
             for (int t = i-j; t < i-ne[j]; t++) update (a[t], -1);
             j = ne[j];
```

```
return ans;
int main()
    while ("scanf ("%d%d%d", &n, &k, &s))
        memset(tree, 0, sizeof(tree));
        memset(1, 0, sizeof(1));
        memset (1e, 0, sizeof(1e));
        memset (ne, 0, sizeof (ne));
        for (int i = 0; i < n; i++)
                                      scanf ("%d", &a[i]);
        for (int i = 0; i < k; i++)
            scanf("%d", &b[i]);
            update(b[i], 1);
            1[i] = getsum(b[i]-1);
            le[i] = getsum(b[i]);
        vector < int > v = kmp();
        printf("%d\n", v. size());
        for(int i = 0; i < v. size(); i++) printf("%d\n", v[i]);
    return 0;
                                           扩展 KMP
//ne[i]:x[i...m-1]与 x[0...m-1]的最长公共前缀
//ex[i]:y[i...n-1]与 x[0...m-1]的最长公共前缀
void getnext(char *s)
    int j = 0, len = strlen(s), k = 1;
    ne[0] = 1en;
    while (j+1 < len && s[j] == s[j+1]) j++;
    ne[1] = j;
for(int i = 2;i < len;i++)
        if(ne[i-k]+i < ne[k]+k) ne[i] = ne[i-k];
        else
            j = \max(0, ne[k]+k-i);
            while (i+j < len \&\& s[i+j] == s[j]) j++;
            ne[i] = j;
            k = i;
void ekmp(char *x, char *y)
    getnext(x);
    int j = 0, lenx = strlen(x), leny = strlen(y);
    while (j < lenx & j < leny & x[j] == y[j]) j++;
    ex[0] = j:
    int k = 0:
    for (int i = 1; i < leny; i++)
        if(ne[i-k]+i < ex[k]+k) ex[i] = ne[i-k];
        else
            j = \max(0, \exp[k] + k - i);
            while (i+j < leny && j < lenx && y[i+j] == x[j]) j++;
            ex[i] = j;
            k = i;
```

```
最长回文字串 Manacher
//abaa
//i: 0 1 3 4 5 6 7 8 9 10
//a[i]: $ # a # b # a # a #
         0 1 3 4 5 6 7 8 9 10
//p[i]: 1 1 2 1 4 1 2 3 2 1
char s[100005], a[200005];
int p[200005];
void manacher (int len)
     int mx = 0, id;
    for (int i = 1; i < len; i++)
         if(mx > i) p[i] = min(p[2*id-i], mx-i);
         else p[i] = 1;
while(a[i+p[i]] == a[i-p[i]]) p[i]++;
         if(p[i]+i > mx)
              mx = p[i]+i;
              id = i;
int main()
    scanf("%s", s);
a[0] = '$';
a[1] = '#';
    int 1en = 2;
    for (int i = 0; s[i]; i++)
         a[len++] = s[i];
a[len++] = '#';
    manacher (1en);
    int ans = 0;
for(int i = 0;i < len;i++) ans = max(ans,p[i]);</pre>
    printf("%d\n", ans-1);
    return 0:
```

最小表示

Trie 树

```
//添加
void add(char *s, int x)
    int now = 0;
    for (int i = 0; i < strlen(s); i++)
        int c = s[i]-'a';
        if (!ch[now][c])
            ch[now][c] = ++sz;
            cnt[sz] = 0;
        now = ch[now][c];
        cnt[now]++;
//查找数量
int getnum(char *s)
    int now = 0;
    for (int i = 0; i < strlen(s); i++)
        int c = s[i] - a';
        if(!ch[now][c]) return 0;
        now = ch[now][c];
    return cnt[now];
```

# AC 自动机

```
//用 n 个模式串建立自动机
//求目标中出现了几个模式串
struct Trie
    int next[500000][26], fail[500005], num[500005], root, cnt;
    int newnode()
        for (int i = 0; i < 26; i++)
                                    next[cnt][i] = -1;
        num[cnt++] = 0;
        return cnt-1;
    void init()
        cnt = 0;
        root = newnode();
    void insert(char *s)
        int now = root, len = strlen(s);
        for (int i = 0; i < len; i++)
            int c = s[i]-'a';
            if(next[now][c] == -1) next[now][c] = newnode();
            now = next[now][c];
        num[now]++;
    void build()
        queue<int> q;
        fail[root] = root;
        for (int i = 0; i < 26; i++)
```

```
if(next[root][i] == -1) next[root][i] = root;
             else
                 fail[next[root][i]] = root;
                 q. push (next[root][i]);
        while(!q. empty())
             int now = q. front();
             q. pop();
             for (int i = 0; i < 26; i++)
                 if(next[now][i] == -1) next[now][i] = next[fail[now]][i];
                 else
                      fail[next[now][i]] = next[fail[now]][i];
                      q. push (next[now][i]);
    int query(char *s)
        int now = root, ans = 0, len = strlen(s);
        for (int i = 0; i < len; i++)
             now = next[now][s[i]-'a'];
             int t = now;
             while(t != root)
                 ans += num[t];
                 num[t] = 0;
                 t = fail[t];
        return ans;
    void debug()
        for (int i = 0; i < cnt; i++)
             printf("id = %3d, fail = %3d, num = %3d, chi = [", i, fail[i], num[i]);
             for(int j = 0; j < 26; j++) printf("%2d", next[i][j]); printf("]\n");
int n;
char s[1000001];
Trie ac;
int main()
    int T;
scanf("%d", &T);
    while (T--)
        scanf ("%d", &n);
        ac.init();
        while (n--)
             scanf ("%s", s);
```

```
ac. insert(s);
        ac.build();
        scanf("%s", s);
printf("%d\n", ac. query(s));
    return 0;
                                        后缀数组 DA
//0 (nlogn)
/待排序数组长度 n, 放在 0^{\sim}n 中,最后补 0
//sa[i]:每个后缀串从小到大排第 i 小的位置
//rank[i]:i 位置的从小到大排序位置
//height[i]:sa[i]和 sa[i-1]对应后缀的最长公共前缀
//n = 8
//num[i]:
            1 1 2 1 1 1 1 2 0
                                      num[8]加0
              8 3 4 5 0 6 1 7 2
//sa|i|:
                                       num[0 n]有效
//rank[i]: 4 6 8 1 2 3 5 7 0
                                      num[0~n-1]有效
                                     num[2~n]有效
//height[i]:0 0 3 2 3 1 2 0 1
int t1[N], t2[N], c[N], sa[N], rank[N], height[N];
int mm[200005], best[20][200005], rmg[200005];
bool cmp(int *r, int a, int b, int 1)
    return r[a] == r[b] \&\& r[a+1] == r[b+1];
```

void da(int \*r, int \*sa, int \*rank, int \*height, int n, int m)

```
height[rank[i]] = k;
void initrmg(int n)
    mm[0] = -1;
    for (int i = 1; i < n; i++)
                                 mm[i] = (i&(i-1) == 0)?mm[i-1]+1:mm[i-1];
    for (int i = 1; i < n; i++)
                                 best[0][i] = i;
    for (int i = 1; i \le mm[n-1]; i++)
        for (int j = 1; j+(1 << i)-1 < n; j++)
            int a = best[i-1][j], b = best[i-1][j+(1 << (i-1))];
            best[i][j] = rmq[a] < rmq[b]?a:b;
int askrmq(int a, int b)
    int t = mm[b-a+1];
    b = (1 << t) -1;
    a = best[t][a];
    b = best[t][b];
    return rmq[a] < rmq[b]?a:b;</pre>
//求 a, b 位置开始的后缀的最长公共前缀
int lcp(int a, int b)
    a = rk[a];
    b = rk[b];
    if(a > b)
                swap(a, b);
    return height[askrmq(a+1, b)];
int main()
    gets(s);
    int len = strlen(s);
    for (int i = 0; i < len; i++) r[i] = s[i];
    da (r, sa, rk, height, 1en, 128);
    return 0;
                                        后缀数组 DC3
//0(n)
//所有数组开 3 倍
\#define F(x) ((x)/3+((x)\%3 == 1?0:tb))
#define G(x) ((x) < tb?(x)*3+1:((x)-tb)*3+2)
int wa[3*N], wb[3*N], wv[3*N], wss[3*N], r[3*N], sa[3*N], rk[3*N], height[3*N];
char s[3*N];
int c0(int *r, int a, int b)
    return r[a] == r[b] \&\& r[a+1] == r[b+1] \&\& r[a+2] == r[b+2];
int c12(int k, int *r, int a, int b)
    if (k == 2) return r[a] < r[b] || r[a] == r[b] && c12(1, r, a+1, b+1);
    return r[a] < r[b] \mid | r[a] == r[b] \&\& wv[a+1] < wv[b+1];
```

```
void sort(int *r, int *a, int *b, int n, int m)
                                    wv[i] = r[a[i]];
    for (int i = 0; i < n; i++)
                                    wss[i] = 0;
    for (int i = 0; i < m; i++)
    for (int i = 0; i < n; i++)
                                    wss[wv[i]]++;
    for (int i = 1; i < m; i++)
                                    wss[i] += wss[i-1];
    for (int i = n-1; i \ge 0; i--) b[--wss[wv[i]]] = a[i];
void dc3(int *r, int *sa, int n, int m)
    int *rn = r+n, *san = sa+n, ta = 0, tb = (n+1)/3, tbc = 0, i, j, p;
    r[n]=r[n+1]=0;
    for(i = 0; i < n; i++)
         if(i\%3) wa[tbc++] = i;
    sort(r+2, wa, wb, tbc, m);
    sort(r+1, wb, wa, tbc, m);
    sort (r, wa, wb, tbc, m);
    for (p = 1, rn[F(wb[0])] = 0, i = 1; i < tbc; i++)
        rn[F(wb[i])] = c0(r, wb[i-1], wb[i])?p-1:p++;
    if (p < tbc) dc3(rn, san, tbc, p);
    else
        for (i = 0; i < tbc; i++) san[rn[i]] = i;
    for (i = 0; i < tbc; i++)
         if(san[i] < tb) wb[ta++] = san[i]*3;
    if(n\%3 == 1)
                      wb[ta++] = n-1;
    sort (r, wb, wa, ta, m);
    for (i = 0; i < tbc; i++) wv [wb[i] = G(san[i])]=i;
    for (i = 0, j = 0, p = 0; i < ta && j < tbc; p++)
                                                          sa[p] =
c12 \text{ (wb[j]}\%3, r, wa[i], wb[j])?wa[i++]:wb[j++];
                         sa[p] = wa[i++];

sa[p] = wb[j++];
    for (; i < ta; p++)
    for (; j < tbc; p++)
void da(int *r, int *sa, int *rank, int *height, int n, int m)
    for (int i = n; i < n*3; i++) r[i] = 0;
    dc3(r, sa, n+1, m);
    int k = 0;
    for (int i = 0; i \le n; i++)
                                    rank[sa[i]] = i;
    for (int i = 0; i < n; i++)
         if(k)
                 k--;
         int t = sa[rank[i]-1];
         while (r[i+k] == r[t+k]) k++;
         height[rank[i]] = k;
```

# 后缀自动机

- 1.所有的子串都能够由 root 走到。
- 2.所有走到终止状态的路径都是后缀。
- 3.每个状态 s 代表的串的长度是区间(len<sub>pre</sub>,len<sub>s</sub>]。
- 4.对于每个状态 s,它代表的所有串在原串中出现次数和每次出现的右端点相同。

- 5.在后缀自动机的 Parent 树中,每个状态的 right 集合都是其父状态 right 集合的子集。
- 6.后缀自动机的 Parent 树是原串的反向前缀树。
- 7.两个串的最长公共后缀,位于这两个串对应状态在 Parent 树上的最近公共祖先状态。

# 后缀自动机模版

```
struct samnode
    int len, right, pre, sum, next[26];
    void clear()
        1en = 0;
        right = 0;
        sum = 0;
        pre = -1;
        memset(next, -1, sizeof(next));
}st[2*N];
int n, sz, root, last;
void saminit()
    sz = 0;
    root = 1ast = 0;
    st[root].clear();
void samadd(int w)
    int p = last, now = ++sz;
    last = now;
    st[now].clear();
    st[now].1en = st[p].1en+1;
    st[now].right = 1;
    while (p != -1 \&\& st[p]. next[w] == -1)
        st[p].next[w] = now;
        p = st[p].pre;
    if(p == -1)
        st[now].pre = root;
        return;
    int q = st[p].next[w];
    if(st[q].1en == st[p].1en+1)
        st[now].pre = q;
        return;
    int neww = ++sz;
    st[neww].clear();
    memcpy(st[neww].next, st[q].next, sizeof(st[q].next));
    st[neww].len = st[p].len+1;
    st[neww].pre = st[q].pre;
    st[q].pre = neww;
    st[now].pre = neww;
    while (p != -1 \&\& st[p]. next[w] == q)
        st[p].next[w] = neww;
        p = st[p].pre;
```

```
void sambuild(char *s)
    saminit();
    int len = strlen(s);
for(int i = 0;i < len;i++)</pre>
         samadd(s[i]-'a');
                                          求最长公共子串
int main()
    scanf ("%s", s);
    sambuild(s);
scanf("%s", s);
    int len = strlen(s), ans = 0;
    int p = root, 1 = 0;
    for (int i = 0; i < len; i++)
         int c = s[i]-'a';
         if (st[p]. next[c] != -1)
             1++;
             p = st[p].next[c];
         else
             while (p != -1 \&\&st[p].next[c] == -1)  p = st[p].pre;
             if(p == -1)
                  1 = 0;
                  p = root;
             else
                  1 = st[p].1en+1;
                  p = st[p].next[c];
         ans = \max(ans, 1);
    printf("%d\n", ans);
    return 0;
                                            求最小表示
int main()
```

```
now = st[now].next[j];
break;
}

printf("%d\n", st[now].len-len+1);
}
```

### 求出现次数大于 k 的子串数量

```
void sambuild(char *s)
    saminit();
    int len = strlen(s);
    for (int i = 0; i < len; i++)
        samadd(s[i]-'a');
char s[N];
int k, num[N], top[2*N];
int main()
    int T;
    scanf ("%d", &T);
    while (T--)
        scanf ("%d%s", &k, s);
        sambuild(s);
        int len = strlen(s);
        memset (num, 0, sizeof (num));
        for (int i = 1; i \le sz; i++)
                                      num[st[i].1en]++;
        for (int i = 1; i \le len; i++) num[i] += num[i-1];
        for (int i = sz; i >= 1; i--)
                                       top[num[st[i].len]--] = i;
        for (int i = sz; i >= 1; i-
             int p = top[i];
             if(st[p].pre != -1) st[st[p].pre].right += st[p].right;
        long long ans = 0;
        for (int i = 1; i \le sz; i++)
             if(st[i].right >= k)
                                      ans += st[i].len-st[st[i].pre].len;
        printf("%11d\n", ans);
    return 0;
```

### 求第 k 小的子串(重复算一个)

```
printf("%c", i+'a');
             dfs(p,k);
             return:
        k = st[p].sum;
int main()
    scanf ("%s%d%d", s, &t, &k);
    sambuild(s);
    int len = strlen(s);
    for(int i = 1; i \le sz; i++) num[st[i].len]++;
    for (int i = 1; i \le len; i++) num[i] += num[i-1];
    for(int i = sz; i \ge 1; i--) top[num[st[i].len]--] = i;
    for (int i = 1; i \le sz; i++)
                                  st[i].right = 1;
    st[0].right = 0;
    for (int i = sz; i >= 0; i--)
        int p = top[i];
        st[p].sum = st[p].right;
        for (int j = 0; j < 26; j++)
             int pp = st[p].next[j];
             if(pp == -1) continue;
             st[p].sum += st[pp].sum;
    if(k > st[0].sum)
                         printf("-1 \setminus n");
    else
        dfs(0, k);
        printf("\n");
    return 0;
```

# 求第 k 小的字串(重复算多个)

```
sambuild(s);
    int len = strlen(s);
    for (int i = 1; i \le sz; i++) num[st[i].len]++;
    for(int i = 1; i \le len; i++) num[i] += num[i-1];
    for(int i = sz; i \ge 1; i--) top[num[st[i].len]--] = i;
    for (int i = sz; i >= 1; i--)
        int p = top[i];
        if(st[p].pre != -1) st[st[p].pre].right += st[p].right;
    for (int i = sz; i >= 0; i--)
        int p = top[i];
        st[p].sum = st[p].right;
        for (int j = 0; j < 26; j++)
            int pp = st[p].next[j];
            if(pp == -1)
                          continue;
            st[p].sum += st[pp].sum;
    if (k > st[0].sum) printf ("-1 \setminus n");
    else
        dfs(0,k);
        printf("\n");
    return 0;
                                         字符串 hash
const int HASH = 10007;
const int N = 2010;
const int SEED = 13331;
char s[N];
struct HASHMAP
    int head[N], next[N], size, f[N];
    unsigned long long state[N];
    void init()
        size = 0:
        memset(state, -1, sizeof(state));
    int insert (unsigned long long x, int id)
        int h = x\%HASH;
        for (int i = head[h]; i != -1; i = next[i])
            if(x = state[i]) return f[i];
        f[size] = id;
```

state[size] = x;
next[size] = head[h];
head[h] = size++;

unsigned long long t = 0;

for (int  $i = 1; i \le len; i++$ ) t = t\*SEED+s[i-1];

return 0;

int main()

# 图论

# 最短路 dij 邻接矩阵

```
//0 (n^2)
void dij(int beg)
    memset(dis, 0x3f, sizeof(dis));
    memset (vis, 0, sizeof (vis));
    dis[beg] = 0;
    for (int i = 1; i \le n; i++)
         int k = -1, minn = INF;
         for (int j = 1; j \le n; j++)
             if(!vis[j] && dis[j] < minn)</pre>
                  minn = dis[j];
                  k = j;
         if (k == -1) break;
         vis[k] = 1;
         for (int j = 1; j \le n; j++)
             if(!vis[j] && dis[k]+a[k][j] < dis[j])</pre>
                  dis[j] = dis[k]+a[k][j];
```

# dij 优先队列

```
///0 \text{ (mlogm)}
struct xxx
    int to, w;
    xxx(int a, int b):to(a), w(b) {};
    friend bool operator <(xxx X, xxx Y)
        return X.w > Y.w;
void dij(int beg)
    priority_queue<xxx> q;
    memset(dis, 0x3f, sizeof(dis));
    memset(vis, 0, sizeof(vis));
    dis[beg] = 0;
    q. push(xxx(beg, 0));
    while(!q.empty())
        int now = q. top(). to;
        q. pop();
        if (vis[now])
                         continue;
        vis[now] = 1;
        for (int i = 0; i < v[now]. size (); i++)
             int t = v[now][i]. to, w = v[now][i]. w;
             if(!vis[t] && dis[now]+w < dis[t])</pre>
```

```
dis[t] = dis[now]+w;
                 q. push(xxx(t, dis[t]));
                                             Floyd
//0 (n^3)
void floyd()
    for (int k = 0; k < n; k++)
        for (int i = 0; i < n; i++)
                                           a[i][j] = min(a[i][j], a[i][k]+a[k][j]);
             for (int j = 0; j < n; j++)
                                          bellman_ford
//0 \, (nm)
struct xxx
    int u, v, w;
    xxx(int a, int b, int c):u(a), v(b), w(c) {};
vector <xxx> v;
//存在负环返回0,否则返回1
bool bellman ford(int beg)
    memset (dis, 0x3f, sizeof (dis));
    dis[beg] = 0;
    for (int i = 1; i < n; i++)
        for (int j = 0; j < v. size(); j++)
             int x = v[j].u, y = v[j].v, w = v[j].w;
             dis[y] = min(dis[y], dis[x]+w);
    for (int i = 0; i < v. size(); i++)
        int x = v[i].u, y = v[i].v, w = v[i].w;
        if(dis[y] > dis[x]+w) return 0;
    return 1;
                                             Spfa
//0 \, (km)
struct xxx
    int to, w;
    xxx(int a, int b):to(a), w(b) {};
vector \langle xxx \rangle v[205];
//存在负环返回0,否则返回1
bool spfa(int beg)
    queue<int> q;
    memset(c, 0, sizeof(c));
    memset(vis, 0, sizeof(vis));
    memset(dis, 0x3f, sizeof(dis));
```

```
dis[beg] = 0;
q. push (beg);
vis[beg] = 1;
c[beg] = 1;
while(!q.empty())
    int now = q. front();
    q. pop();
    vis[now] = 0;
    for (int i = 0; i < v[now]. size(); i++)
         int t = v[now][i]. to, w = v[now][i]. w;
         if (dis[now]+w < dis[t])</pre>
             dis[t] = dis[now]+w;
             vis[t] = 1;
             q.push(t);
if(++c[t] > n) return 0;
return 1;
```

### 分层图最短路

```
struct xxx
    int to, w;
    xxx(int a, int b):to(a), w(b) {};
    friend bool operator (xxx x, xxx y)
         return x.w > y.w;
vector \langle xxx \rangle v[11005];
int n, m, k, a[1005][1005], dis[11005], vis[11005];
void dij(int beg)
    memset (dis, 0x3f, sizeof (dis));
    memset(vis, 0, sizeof(vis));
    priority_queue<xxx> q;
dis[beg] = 0;
    q. push (xxx(beg, 0));
    while (!q. empty())
         int now = q. top(). to;
         q. pop();
         if (vis[now])
                          continue;
         vis[now] = 1;
         for (int i = 0; i < v[now]. size (); i++)
              int t = v[now][i]. to, w = v[now][i]. w;
              if (w+dis[now] < dis[t])</pre>
                  dis[t] = w+dis[now];
                  q. push (xxx (t, dis[t]));
int main()
```

```
while ("scanf ("%d%d%d", &n, &m, &k))
     for (int i = 1; i \le 11000; i++)
                                              v[i].clear();
     memset (a, 0x3f, sizeof(a));
     for (int i = 1; i \le m; i++)
          int u, v, w;
         scanf("%d%d%d", &u, &v, &w);
a[u][v] = min(a[u][v], w);
a[v][u] = min(a[v][u], w);
    for (int i = 1; i \leq n; i++)
          for (int j = 1; j \le n; j++)
               if(i == j \mid \mid a[i][j] == INF)
                                                        continue;
               for (int 1 = 0; 1 \le k; 1++)
                    v[1*n+i]. push_back(xxx(1*n+j, a[i][j]));
                    if(1 != k)
                         v[1*n+i]. push back(xxx((1+1)*n+j,0));
     dij(1);
     printf("%d\n", dis[(k+1)*n]);
return 0;
```

# 最小生成树 prim 邻接矩阵

```
//不连通返回-1,否则返回最小花费
int prim()
    int ans = 0;
    memset(vis, 0, sizeof(vis));
    vis[1] = 1;
    for (int i = 2; i \le n; i++)
                                 cost[i] = a[1][i];
    for (int i = 2; i \le n; i++)
        int minn = INF;
        int k = -1;
        for (int j = 1; j \le n; j++)
            if(!vis[j] && minn > cost[j])
                minn = cost[j];
                k = j;
        if (minn == INF) return -1;
        ans += minn;
        vis[k] = 1;
        for (int j = 1; j \le n; j++)
            if(!vis[j] \&\& cost[j] > a[k][j]) cost[j] = a[k][j];
    return ans;
```

# prim 优先队列

```
struct xxx
    int to, w;
    xxx(int a, int b):to(a), w(b) {}
    friend bool operator (xxx a, xxx b)
        return a.w > b.w;
//非连通返回-1,否则返回最小花费
int prim()
    memset(vis, 0, sizeof(vis));
    for (int i = 0; i < n; i++)
                                 dis[i] = INT MAX;
    int ans = 0;
    priority_queue<xxx> q;
    q. push (xxx(1, 0));
    while(!q.empty())
         while(!q.empty() && vis[q.top().to]) \\ q.pop(); 
        if(q.empty()) break;
        ans += q. top().w;
        int now = q. top(). to;
        q. pop();
        vis[now] = 1;
        for (int i = 0; i < v[now]. size(); i++)
            int t = v[now][i]. to, w = v[now][i]. w;
            if(vis[t]) continue;
            if (dis[t] <= w) continue;
            dis[t] = w;
            q. push(xxx(t, w));
    return ans;
                                           Kruskal
```

```
struct xxx
    int from, to, w;
    xxx(int a, int b, int c):from(a), to(b), w(c) {};
    friend bool operator <(xxx a, xxx b)
        return a.w < b.w;
vector (xxx) v;
int findd(int x)
    return pre[x] == x?x:findd(pre[x]);
//不连通返回-1,否则返回最小花费
bool kruskal()
                                 pre[i] = i;
    for (int i = 1; i \le n; i++)
    int ans = 0, cnt = 0;
    sort(v.begin(), v.end());
    for (int i = 0; i < v. size(); i++)
        int x = findd(v[i].from), y = findd(v[i].to);
```

# 矩阵树定理

- 1.G 的度数矩阵 D[G]是一个 n\*n 的矩阵,并且满足: 当 i≠j 时,dij=0; 当 i=j 时,dij 等于 vi 的度数。
- 2.G 的邻接矩阵 A[G]也是一个 n\*n 的矩阵, 并且满足: 如果 vi、vj 之间有边直接相连,则 aij=1,否则为 0。
- 3.G 的 Kirchhoff 矩阵(也称为拉普拉斯算子)C[G]=D[G]-A[G]。

则 G 的生成树个数等于 C[G]任何一个 n-1 阶主子式的行列式的绝对值。

### 矩阵树

```
int n, m, k, a [55] [55];
long long g[55][55];
long long calc()
    long long ans = 1;
    for (int i = 1; i < n; i++)
        for (int j = i+1; j < n; j++)
             while (g[j][i] != 0)
                 long long t = g[i][i]/g[j][i];
                 for (int k = i; k < n; k++)
                                                g[i][k] = t*g[j][k];
                                                swap(g[i][k], g[j][k]);
                 for (int k = i; k < n; k++)
                 ans = -ans;
        if(g[i][i] == 0)
                              return 0;
        ans = ans*g[i][i];
    return abs(ans);
int main()
    ios::sync_with_stdio(false);
    while (cin \gg n \gg m \gg k)
        memset(a, 0, sizeof(a));
        memset(g, 0, sizeof(g));
        while(m--)
             int x, y;
             cin \gg x \gg y;
             a[x][y] = a[y][x] = 1;
        for (int i = 1; i \le n; i++)
             int t = 0;
             for (int j = 1; j \le n; j++)
                 if(i != j && !a[i][j])
```

```
t^{++};
g[i][j] = -1;
            g[i][i] = t;
        cout << calc() << endl;
    return 0;
                                         次小生成树
int n, vis[505], used[505], pre[505], cost[505], a[505][505], maxd[505][505];
//不连通返回-1,否则返回次小花费
int prim()
    int ans = 0;
    memset(vis, 0, sizeof(vis));
    memset (used, 0, sizeof (used));
    memset (maxd, 0, sizeof (maxd));
    vis[1] = 1;
                                 cost[i] = a[1][i];
    for (int i = 2; i \le n; i++)
    for (int i = 2; i \le n; i++)
        int minn = INF;
        int k = -1;
        for (int j = 1; j \le n; j++)
            if(!vis[j] && minn > cost[j])
                minn = cost[j];
                k = j;
        if (minn == INF) return -1;
        ans += minn;
        vis[k] = 1;
        for (int j = 1; j \le n; j++)
            if(vis[j]) maxd[j][k] = maxd[k][j] = max(maxd[j][pre[k]], cost[k]);
            if(!vis[j] && cost[j] > a[k][j])
                cost[j] = a[k][j];
                pre[j] = k;
    return ans;
                                         Tarjan 缩点
//0 (n+m)
//缩点, belong 保存每个点属于的强联通集合编号
int n, m, vis[10005], dfn[10005], low[10005], belong[10005], cnt, ans;
vector < int > v[10005];
```

void tarjan(int now)

stack<int> s;
s. push(now);
vis[now] = 1;

dfn[now] = 1ow[now] = ++cnt;

for (int i = 0; i < v[now]. size(); i++)

```
int t = v[now][i];
        if(dfn[t] == 0)
            tarjan(t);
            low[now] = min(low[now], low[t]);
                            low[now] = min(low[now], dfn[t]);
        else if (vis[t])
    if(dfn[now] == low[now])
        ans++;
        while(!s.empty())
            int t = s. top();
            s. pop();
            vis[t] = 0;
            belong[t] = ans;
            if(t == now) break;
int main()
    for (int i = 1; i \le n; i++)
        if(dfn[i] == 0) tarjan(i);
                                     Kosaraju 求强联通分量
int n, m, vis1[10005], vis2[10005], ans;
vector < int > v1[10005], v2[10005];
stack<int> s;
void dfs1(int now)
    vis1[now] = 1;
    for (int i = 0; i < v1[now]. size(); i++)
        int t = v1[now][i];
        if(vis1[t]) continue;
        dfs1(t);
    s. push (now);
void dfs2(int now)
    vis2[now] = 1;
    for (int i = 0; i < v2[now]. size (); i++)
        int t = v2[now][i];
        if(vis2[t]) continue;
        dfs2(t);
int main()
    ios::sync_with_stdio(false);
    while (scanf ("%d%d", &n, &m) && (n+m))
        for (int i = 1; i \le n; i++)
```

```
v1[i].clear();
         v2[i]. clear();
    memset (vis1, 0, sizeof (vis1));
    memset (vis2, 0, sizeof (vis2));
    while(!s.empty())
                           s. pop();
    ans = 0;
    while (m--)
         int x, y;
scanf("%d%d", &x, &y);
         v1[x]. push back(y);
         v2[y]. push back (x);
    for (int i = 1; i \le n; i++)
         if(!vis1[i]) dfs1(i);
    while (!s. empty())
         if(!vis2[s. top()])
             ans++:
             dfs2(s.top());
         s. pop();
    if (ans == 1) printf ("Yes\n");
             printf("No\n");
    else
return 0;
```

### 无向图中一些基本概念

```
割点集合: 删除这个顶点集合,以及这个集合中所有顶点关联的边以后,原图变成多个连通块。
点连通度: 最小割点集合中的顶点数。
割边集合: 删除这个边集合以后,原图变成多个连通块。
边连通度: 最小割边集合中的边数。
点双连通(双连通或重连通): 无向连通图的点连通度大于 1。
割点(关节点): 当且仅当这个图的点连通度为 1,则割点集合的唯一元素被称为割点。
边双连通(双连通或重连通): 无向连通图的边连通度大于 1。
桥(关节边): 当且仅当这个图的边连通度为 1,则割边集合的唯一元素被称为桥。
边双连通分支: 在连通分支中去掉一条边,连通分支仍连通。
点双连通分支(块): 在连通分支中去掉一个点,连通分支仍连通。
```

### 求割点和桥

```
struct bridge
{
    int u, int v;
    bridge(int a, int b):u(a),v(b) {}
};
vector<bridge> ansb;
vector<int> absc;

//ansb 保存桥, ansc 保存割点
void tarjan(int now, int pre)
{
    low[now] = dfn[now] = ++cnt;
    int son = 0;
    for(int i = 0;i < v[now].size();i++)
    {
        int t = v[now][i];
    }
}</pre>
```

```
if(t == pre)
                 continue;
    if(!dfn[t])
       son++:
       tarjan(t, now);
       low[now] = max(low[now], low[t]);
       //桥, (u, v)为树枝, low[v] > dfn[u]
       if(low[t] > dfn[now])
           ansb. push back(bridge(now, t));
           ansb. push back(bridge(t, now));
       //割点, (u, v)父子边, u 不为树根, low[v] >= dfn[u]
       if (now != root && low[t] >= dfn[now]) ansc. push back(t);
           low[now] = max(low[now], low[t]);
   else
//割点, u 为根, 分支数大于 1
if(now == root && son > 1) ansc.push_back(now);
```

### 求边双连通分支

去掉桥, 其余的连通分支就是边双连通分支了。

一个有桥的连通图要变成双连通图的话,把双连通子图缩成一个点,形成一棵树。

再加上(leaf+1)/2条边。

# 求点双连通分支

```
//cutcnt 表示个数
//block 储存每个点连通分支
stack<int> s;
vector<int> block[10005];
void tarjan(int now, int pre)
    low[now] = dfn[now] = ++cnt;
    s. push (now);
    for (int i = 0; i < v[now]. size(); i++)
        int t = v[now][i];
        if(t == pre)
                       continue;
        if (!dfn[t])
            tarjan(t, now);
            low[now] = max(low[now], low[t]);
            if(low[t]) = dfn[now]
                cutcnt++;
                block[cutcnt].clear();
                while(!s.empty())
                    int tt = s. top();
                    s. pop();
                    block[cutcnt].push back(tt);
                    if(tt = t) break;
                block[cutcnt].push back(now);
                low[now] = max(low[now], low[t]);
        else
```

### 最小树形图

```
struct xxx
    int u, v, w;
}e[10005];
int n, m, a[55], sum[55], cnt, in[800], pre[800], vis[800], id[800];
void addedge(int u, int v, int w)
    e[cnt].u = u;
    e[cnt].v = v;
    e[cnt++].w = w;
//已处理重边,不存在最小树形图返回-1,否则返回最小值
int zhuliu(int root, int n, int m)
    int ans = 0;
    while (1)
        //找最小入边
       for (int i = 0; i < n; i++)
                                 in[i] = INF;
        for (int i = 0; \langle m; i++ \rangle
            int u = e[i].u;
            int v = e[i].v;
            if(e[i].w < in[v] \&\& u != v)
                pre[v] = u;
                in[v] = e[i].w;
                 //去除自环和重边
        for (int i = 0; i < n; i++)
            if(i == root) continue;
if(in[i] == INF) retur
                                                                           //除了跟以外
                              return -1;
有点没有入边,则根无法到达它
                                                 //找环
        int cnt = 0;
        memset (id, -1, sizeof(id));
        memset(vis, -1, sizeof(vis));
        in[root] = 0;
        for (int i = 0; i < n; i++)
            ans += in[i];
            int v = i:
            while(vis[v] != i && id[v] == -1 && v != root) //每个点寻找其前序点,要么
最终寻找至根部,要么找到一个环
               vis[v] = i;
                v = pre[v];
            if(v != root && id[v] == -1)//缩点
                for(int u = pre[v]; u != v; u = pre[u]) id[u] = cnt;
                id[v] = cnt++;
        if(cnt == 0) break; //无环
                                        则 break
        for (int i = 0; i < n; i++)
            if(id[i] == -1) id[i] = cnt++;
```

```
for(int i = 0; i < m;)
{
    int u = e[i].u;
    int v = e[i].v;
    e[i].u = id[u];
    e[i].v = id[v];
    if(id[u]!=id[v]) e[i++].w -= in[v];
    else e[i]=e[--m];
}
    n = cnt;
    root = id[root];
}
return ans;
}</pre>
```

二分图

最大匹配 == 最小点覆盖 最小路径覆盖 == 最大独立集 == 顶点数-最大匹配 最大团 == 补图最大独立集

```
匈牙利
```

```
int n, m, 1;
int linker[505];
bool used[505];
vector<int> v[505];
bool dfs(int u)
    for (int i = 0; i < v[u]. size(); i++)
        int t = v[u][i];
        if(!used[t])
             used[t] = 1:
             if(linker[t] == -1 \mid | dfs(linker[t]))
                 linker[t] = u;
                 return 1;
    return 0;
int MaxMatch()
    int ans = 0;
    memset(linker, -1, sizeof(linker));
    for (int i = 1; i \le n; i++)
        memset (used, 0, sizeof (used));
        if(dfs(i))
                      ans++;
    return ans;
```

Hopcroft-Carp

```
//O(V^0.5 E)
//适用于数据较大的二分匹配
int g[MAXN], Mx[MAXN], My[MAXN], Nx, Ny;
int dx[MAXN], dy[MAXN], dis;
bool vst[MAXN];
bool searchP()
```

```
queue<int>Q;
    dis=INF;
    memset (dx, -1, sizeof(dx));
    memset (dy, -1, sizeof(dy));
    for (int i=0; i < Nx; i++)
         if(Mx[i]==-1)
              Q. push(i);
              dx[i]=0;
    while(!Q. empty())
         int u=Q. front();
         Q. pop();
         if(dx[u]>dis) break;
         for(int v=0; v<Ny; v++)
if(g[u][v]&&dy[v]==-1)
                   dy[v]=dx[u]+1;
                  if (My \lfloor v \rfloor ==-1) dis=dy\lfloor v \rfloor;
                  else
                       dx[My[v]]=dy[v]+1;
                       Q. push (My[v]);
    return dis!=INF;
bool DFS (int u)
    for (int v=0; v<Ny; v++)
        if(!vst[v]\&\&g[u][v]\&\&dy[v]==dx[u]+1)
             vst[v]=1;
             if (My[v]!=-1\&dy[v]==dis) continue;
             if(My[v] == -1 \mid DFS(My[v]))
                 My[v]=u;
                 Mx[u]=v;
                 return 1;
    return 0;
int MaxMatch()
    int res=0;
    memset (Mx, -1, sizeof(Mx));
    memset (My, -1, sizeof(My));
    while(searchP())
         memset(vst, 0, sizeof(vst));
         for (int i=0; i<Nx; i++)
           if(Mx[i]==-1\&\&DFS(i)) res++;
    return res;
                                            多重匹配匈牙利
bool dfs(int u)
```

for (int i = 0; i < v[u]. size(); i++)

```
int t = v[u][i];
         if(used[t]) continue;
        used[t] = 1;
if(linker[t][0] < capacity[t])</pre>
             linker[t][++linker[t][0]] = u;
             return 1;
        for (int j = 1; j \le capacity[t]; j++)
             if(dfs(linker[t][j]))
                 linker[t][j] = u;
                 return 1:
        return 0;
int MaxMatch()
    int ans = 0;
    memset(linker, 0, sizeof(linker));
    for (int i = 1; i \le n; i++)
        memset (used, 0, sizeof (used));
        if(dfs(i))
                       ans++;
    return ans;
                                          二分图最大权 KM
int n, m, g[305][305], linker[305], lx[305], ly[305], slack[305], visx[305], visy[305];
bool dfs(int u)
    visx[u] = 1;
    for (int i = 1; i \le m; i++)
        if(visy[i]) continue;
         int t = 1x[u]+1y[i]-g[u][i];
         if(t == 0)
             visy[i] = 1;
             if(!linker[i]|| dfs(linker[i]))
                 linker[i] = u;
                 return 1;
                 slack[i] = min(slack[i], t);
        else
    return 0;
int KM()
    memset(linker, 0, sizeof(linker));
    memset (1y, 0, sizeof(1y));
    for (int i = 1; i \le n; i++)
        1x[i] = -INF;
```

```
for (int j = 1; j \le m; j++)
         if(g[i][j] > 1x[i])  1x[i] = g[i][j];
for (int i = 1; i \le n; i++)
    memset(slack, 0x3f, sizeof(slack));
    while (1)
         memset(visx, 0, sizeof(visx));
         memset(visy, 0, sizeof(visy));
         if(dfs(i))
                      break:
         int d = INF;
         for (int j = 1; j \le m; j++)
             if(!visy[j]) d = min(d, slack[j]);
         for (int j = 1; j \le n; j++)
             if(visx[j])  1x[j] = d;
         for (int j = 1; j \le m; j++)
             if(visy[j]) ly[j] += d;
else slack[j] -= d;
int ans = 0;
for (int i = 1; i \le m; i++)
    if(linker[i]) ans += g[linker[i]][i];
return ans;
```

网络流

### 最大流 == 最小割

### EdmondsKarp

```
//s 为源点, t 为汇点
//0 \, (nm^2)
int n, m, s, t, a[20][20], pre[20], vis[20];
bool bfs()
    memset (pre, 0, sizeof (pre));
    memset(vis, 0, sizeof(vis));
    vis[s] = 1;
    queue (int) q;
    q. push(s);
    while(!q.empty())
        int now = q. front();
        q. pop();
        if(now == t)
                        return 1;
        for (int i = 1; i \le n; i++)
             if(vis[i]) continue;
             if (a[now][i])
                 q. push(i);
```

```
pre[i] = now;
                 vis[i] = 1;
    return 0;
int EdmondsKarp()
   int ans = 0;
   while(bfs())
       int minn = 1e9;
       for(int i = t;i != s;i = pre[i]) minn = min(minn, a[pre[i]][i]);
       for(int i = t;i != s;i = pre[i])
           a[pre[i]][i] = minn;
           a[i][pre[i]] += minn;
       ans += minn;
   return ans;
                                            Dinic
struct edge
    int u, v, c, next;
e[20*N];
int n, m, sp, tp, cnt, head[N], d[N];
//源点汇点全局定义
void init()
    cnt = 0;
    memset (head, -1, sizeof (head));
void addedge(int u, int v, int c)
    e[cnt].u = u;e[cnt].v = v;e[cnt].c = c;
    e[cnt].next = head[u]; head[u] = cnt++;
    e[cnt].u = v;e[cnt].v = u;e[cnt].c = 0;
    e[cnt].next = head[v]; head[v] = cnt++;
int bfs()
    queue (int) q;
    memset (d, -1, sizeof(d));
    d[sp]=0;
    q. push (sp);
    while (!q. empty())
        int cur = q. front();
        for (int i = head[cur]; i != -1; i = e[i]. next)
             int u = e[i].v;
            if(d[u] == -1 \&\& e[i].c > 0)
                 d[u] = d[cur]+1;
                 q. push (u);
```

```
return d[tp] != -1;
int dfs(int a, int b)
    int r = 0;
    if (a == tp) return b;
    for(int i = head[a]; i != -1 \&\& r < b; i = e[i].next)
        int u = e[i].v;
        if(e[i].c > 0 \&\& d[u] == d[a]+1)
             int x = min(e[i].c,b-r);
             x = dfs(u, x);
            r_{-}+= x;
             e[i].c = x;
             e[i^1].c += x;
    if(!r) d[a] = -2;
    return r;
int dinic()
    int ans = 0, t;
    while(bfs())
        while (t = dfs(sp, INF))
        ans += t;
    return ans;
                                             Isap
struct Edge
    int to, next, cap, flow;
} edge[M];
int n, m, cnt, head[N], gap[N], d[N], cur[N], que[N], p[N];
void init()
    cnt = 0;
    memset (head, -1, sizeof (head));
void addedge(int u, int v, int c)
    edge[cnt].to = v;
    edge[cnt].cap = c;
    edge[cnt].flow = 0;
    edge[cnt].next = head[u];
    head[u] = cnt++;
    edge[cnt].to = u;
    edge[cnt].cap = 0;
    edge[cnt].flow = 0;
    edge[cnt].next = head[v];
    head[v] = cnt++;
```

```
void bfs(int source, int sink)
    memset (d, -1, sizeof(d));
    memset(gap, 0, sizeof(gap));
    gap[0] = 1;
    int front = 0, rear = 0;
    d[sink] = 0;
    que[rear++] = sink;
    while(front != rear)
        int u = que[front++];
        for (int i = head[u]; i != -1; i = edge[i]. next)
            int v = edge[i].to;
            if (d[v] != -1) continue;
            que[rear++] = v;
            d[v] = d[u]+1;
            gap[d[v]]++;
int isap(int source, int sink, int N)
                                            //源点,汇点,总点数
    bfs(source, sink);
    memcpy (cur, head, sizeof (head));
    int top = 0, x = source, flow = 0;
    while(d[source] < N)</pre>
        if(x == sink)
             int Min = INF, inser;
            for (int i = 0; i < top; i++)
                 if (Min > edge[p[i]]. cap-edge[p[i]]. flow)
                     Min = edge[p[i]]. cap-edge[p[i]]. flow;
                     inser = i;
            for (int i = 0; i < top; i++)
                 edge[p[i]]. flow += Min;
                 edge[p[i]^1]. flow = Min;
            flow += Min;
            top = inser;
            x = edge[p[top]^1].to;
            continue;
        for (int i = cur[x]; i != -1; i = edge[i]. next)
            int v = edge[i].to;
            if(edge[i]. cap > edge[i]. flow && d[v]+1 == d[x])
                 ok = 1;
                 cur[x] = i;
                 p[top++] = i;
                 x = edge[i].to;
                 break;
        if (!ok)
```

```
int Min = N;
            for (int i = head[x]; i != -1; i = edge[i]. next)
                 if(edge[i].cap > edge[i].flow && d[edge[i].to] < Min)</pre>
                     Min = d[edge[i].to];
                     cur[x] = i;
            if(-gap[d[x]] == 0) break;
            gap[d[x] = Min+1]++;
            if (x != source) x = edge[p[--top]^1]. to;
    return flow;
                                        最小费用最大流
int n, m, s, t, pre [1005], vis [1005], dis [1005], head [1005], cnt = 0;
struct ee
    int u, v, cap, cost, next;
e[2005]:
void add(int u, int v, int cap, int cost)
    e[cnt].u = u;
    e[cnt].v = v;
    e[cnt].cap = cap;
    e[cnt].cost = cost;
    e[cnt].next = head[u];
    head[u] = cnt++;
    e[cnt].u = v;
    e[cnt].v = u;
    e[cnt]. cap = 0;
    e[cnt].cost = -cost;
    e[cnt].next = head[v];
    head[v] = cnt++;
bool spfa()
    memset(vis, 0, sizeof(vis));
    memset(pre, -1, sizeof(pre));
    for (int i = s; i \le t; i++) dis[i] = 1e9;
    vis[s] = 1;
    dis[s] = 0;
    queue <int> q;
    q. push(s);
    while(!q.empty())
        int u = q. front();
        q. pop();
        vis[u] = 0;
        for (int i = head[u]; i != -1; i = e[i]. next)
            int v = e[i].v;
            if (e[i].cap && dis[v] > dis[u]+e[i].cost)
                 dis[v] = dis[u]+e[i].cost;
                 pre[v] = i;
                 if(vis[v]) continue;
                 q. push (v);
```

#### 2-sat

#### 染色法

```
//可得到最小字典序的解
int n, m, vis[16005];
vector \langle int \rangle v[16005];
stack(int) s;
bool dfs(int u)
    if (vis[u<sup>1</sup>])
                    return 0;
    if(vis[u]) return 1;
    vis[u] = 1;
    s. push (u);
    for (int i = 0; i < v[u]. size(); i++)
         int t = v[u][i];
        if(!dfs(t)) return 0;
    return 1;
bool twosat(int n)
    memset(vis, 0, sizeof(vis));
    for (int i = 0; i < n; i += 2)
        if(vis[i] || vis[i^1]) continue;
```

```
while(!s.empty()) s.pop();
         if(!dfs(i))
             while(!s.empty())
                 vis[s.top()] = 0;
                 s. pop();
             if(!dfs(i^1)) return 0;
    return 1;
int main()
    ios::sync_with_stdio(false);
    while (cin >> n >> m)
        for (int i = 0; i < 2*n; i++) v[i]. clear();
        memset(vis, 0, sizeof(vis));
        while(m--)
             int a, b;
             cin \gg a \gg b;
             b-
             v[a].push_back(b^1);
             v[b]. push_back(a^1);
         if (twosat (2*n))
             for (int i = 0; i < 2*n; i++)
                  if(vis[i]) cout \langle\langle i+1 \langle\langle end1;
                 cout << "NIE" << endl;</pre>
        else
    return 0;
                                       tarjan 缩点+拓扑排序
int n, m, vis[2*N], dfn[2*N], low[2*N], belong[2*N], pos[2*N], deg[2*N], ok[2*N], cnt, color;
vector < int > v[2*N], v2[2*N], ans;
stack(int) st;
//求任意解
void top(int n)
    ans. clear();
    memset (ok, 0, sizeof (ok));
    queue <int> q;
    for (int i = 1; i \le color; i++)
        if(deg[i] == 0) q. push(i);
    while (!q. empty())
         int now = q. front();
        q. pop();
        if(ok[now] == 0)
             ok[now] = 1;
```

```
ok[pos[now]] = 2;
        for (int i = 0; i < v2[now]. size (); i++)
             int t = v2[now][i];
                                   q.push(t);
             if(--deg[t] == 0)
    for (int i = 0; i < n; i++)
        if(ok[belong[i]] == 1) ans. push back(i);
void tarjan(int now)
    st. push (now);
    vis[now] = 1;
dfn[now] = low[now] = ++cnt;
    for (int i = 0; i < v[now]. size(); i++)
        int t = v[now][i];
        if(dfn[t] == 0)
             tar jan(t);
             low[now] = min(low[now], low[t]);
                          low[now] = min(low[now], dfn[t]);
        else if (vis[t])
    if(dfn[now] == low[now])
        color++;
        while(!st.empty())
             int t = st. top();
             st. pop();
             vis[t] = 0;
             belong[t] = color;
             if(t == now) break;
bool solve(int n)
    for (int i = 0; i < n; i++)
        if(dfn[i] == 0) tarjan(i);
    for (int i = 0; i < n; i += 2)
        if(belong[i] == belong[i^1])
                                            return 0;
        pos[belong[i]] = belong[i^1];
pos[belong[i^1]] = belong[i];
    for (int i = 0; i < n; i++)
        for (int j = 0; j < v[i]. size(); j++)
             int t = v[i][j];
             if (belong[i] != belong[t])
                 deg[belong[i]]++;
                 v2[belong[t]].push back(belong[i]);
```

```
top(n);
    return 1;
                                       曼哈顿最小生成树
struct point
    int x, y, id;
    friend bool operator <(point a, point b)
        if (a. x != b. x) return a. x < b. x;
        return a. y < b. y;
}p[10005];
struct Bit
    int minn, pos;
    void init()
        minn = INF;
        pos = -1;
}bit[10005];
struct edge
    int u, v, d;
    friend bool operator (edge a, edge b)
        return a.d < b.d;
e[40005]:
int n, k, cnt, pre[10005], a[10005], b[10005];
int findd(int x)
    return pre[x] == x?x:findd(pre[x]);
void addedge(int u, int v, int d)
    e[++cnt].u = u;
    e[cnt].v = v;
    e[cnt].d = d;
int lowbit(int x)
    return x&-x;
void update(int i, int x, int pos)
    while (i > 0)
        if (x < bit[i].minn)</pre>
            bit[i].minn = x;
            bit[i].pos = pos;
        i = lowbit(i);
```

```
int ask(int i, int m)
    int minn = INF, pos = -1;
    while (i \le m)
        if(bit[i].minn < minn)</pre>
             minn = bit[i].minn;
             pos = bit[i].pos;
        i += lowbit(i);
    return pos;
int dist(point a, point b)
    return abs (a. x-b. x) +abs (a. y-b. y);
void manhattan()
    cnt = 0:
    for (int dir = 0; dir < 4; dir++)
        if(dir == 1 \mid \mid dir == 3)
             for (int i = 1; i \le n; i++)
                                           swap(p[i].x,p[i].y);
        else if (dir == 2)
             for (int i = 1; i \le n; i++) p[i].x = -p[i].x;
        sort(p+1, p+1+n);
        for (int i = 1; i \le n; i++)
                                       a[i] = b[i] = p[i].y-p[i].x;
        sort(b+1, b+1+n);
        int t = unique(b+1, b+n+1)-b-1;
                                     bit[i].init();
        for (int i = 1; i \le t; i++)
        for (int i = n; i >= 1; i--)
             int pos = 1 ower bound (b+1, b+1+t, a[i]) - b+1;
             int ans = ask(pos, t);
             if (ans != -1) addedge (p[i].id, p[ans].id, dist <math>(p[i], p[ans]));
             update(pos, p[i]. x+p[i]. y, i);
//返回第 k 小的边长
int solve(int k)
    manhattan():
    for(int i = 1; i \le n; i++) pre[i] = i;
    sort(e+1, e+1+cnt);
    for (int i = 1; i \le cnt; i++)
        int u = e[i].u, v = e[i].v;
        int x = findd(u), y = findd(v);
        if(x != y)
             pre[x] = y;
             if(--k == 0) return e[i].d;
```

```
一般图匹配带花树
//g[i][j]存放关系图: i, j是否有边, match[i]存放 i 所匹配的点
//最终匹配方案为 match
//复杂度 0 (n^3)
//点是从 1 到 n 的
bool g[MAXN][MAXN], inque[MAXN], inblossom[MAXN], inpath[MAXN];
int match[MAXN], pre[MAXN], base[MAXN];
queue < int > q;
//找公共祖先
int findancestor(int u, int v)
    memset (inpath, false, sizeof (inpath));
    while (1)
        u = base[u];
        inpath[u] = 1;
        if (\text{match}[u] == -1) break;
        u = pre[match[u]];
    while (1)
        v = base[v];
        if (inpath[v])
                       return v;
        v = pre[match[v]];
//压缩花
void reset(int u, int anc)
    while (u != anc)
        int v = match[u];
        inblossom[base[u]] = 1;
        inblossom[base[v]] = 1;
        v = pre[v];
        if (base[v] != anc) pre[v] = match[u];
        u = v;
void contract(int u, int v, int n)
    int anc = findancestor(u, v);
    memset(inblossom, 0, sizeof(inblossom));
    reset (u, anc);
    reset (v, anc);
    if(base[u] != anc) pre[u] = v;
    if (base[v] != anc) pre[v] = u;
    for (int i = 1; i \le n; i++)
        if(inblossom[base[i]])
            base[i] = anc;
            if(!inque[i])
                q. push back(i);
                inque[i] = 1;
```

```
bool bfs(int S, int n)
    for (int i = 0; i \le n; i++) pre[i]=-1, inque[i]=0, base[i]=i;
    q. clear();
    q. push_back(S);
    inque[\overline{S}] = 1;
    while(!q. empty())
        int u = q. front();
        q. pop();
        for (int v = 1; v \le n; v++)
             if(g[u][v] && base[v] != base[u] && match[u] != v)
                 if(v == S \mid | match[v] !=-1 \&\& pre[match[v]] != -1) contract(u, v, n);
                 else if (pre[v] == -1)
                     pre[v] = u;
                     if(match[v] != -1)
                          q.push_back(match[v]);
                          inque[match[v]] = 1;
                     else
                          u = v;
                          while (u! = -1)
                              v = pre[u];
                              int w = match[v];
                              match[u] = v;
                              match[v] = u;
                              u = w;
                         return 1;
    return 0;
int solve(int n)
    memset (match, -1, sizeof (match));
    int ans = 0;
    for (int i=1; i<=n; i++)
        if(match[i] == -1 \&\& dfs(i,n)) ans++;
    return ans;
                                      最近公共祖先倍增 Ica
void dfs(int now, int pre)
```

```
void dfs(int now, int pre)
{
    dep[now] = dep[pre]+1;
    for(int i = 0; i < v[now]. size(); i++)
    {
}</pre>
```

```
int t = v[now][i];
        if(t == pre)
                       continue;
        fa[t][0] = now;
        dfs(t, now);
int lca(int x, int y)
    if(dep[x] < dep[y]) swap(x, y);
    int t = dep[x]-dep[y];
    for (int i = 0; i \le 20; i++)
        if((1 << i) &t) x = fa[x][i];
    if(x == y) return x;
    for (int i = 20; i >= 0; i--)
        if (fa[x][i] != fa[y][i])
             x = fa[x][i];
             y = fa[y][i];
    return fa[x][0];
int main()
    dfs(1,0);
    for (int i = 1; i \le 20; i++)
        for (int j = 1; j \le n; j++)
             fa[j][i] = fa[fa[j][i-1]][i-1];
                                            lca+rmq
int n, m, f[N*2], rmq[N*2];
vector<int> v[N];
struct ST
    int mm[2*N], dp[2*N][20];
    void init(int n)
        mm[0] = -1;
        for (int i = 1; i \le n; i++)
             mm[i] = ((i&(i-1)) == 0)?mm[i-1]+1:mm[i-1];
             dp[i][0] = i;
        for (int j = 1; j \le mm[n]; j++)
             for (int_{i} = 1; i+(1 \le j)-1 \le n; i++) dp[i][j] = rmq[dp[i][j-1]] \le rmq[dp[i][j-1]]
rmq[dp[i+(1<<(j-1))][j-1]]?dp[i][j-1]:dp[i+(1<<(j-1))][j-1];
    int query (int a, int b)
        if(a > b) swap (a, b);
```

```
int k = mm[b-a+1];
        return rmq[dp[a][k]] \leq rmq[dp[b-(1<<k)+1][k]]?dp[a][k]:dp[b-(1<<k)+1][k];
}st;
void dfs (int now, int pre, int dep)
    f[++cnt] = now;
    rmq[cnt] = dep;
    p[now] = cnt;
    for (int i = 0; i < v[now]. size(); i++)
        int t = v[now][i];
        if(t == pre)
                         continue;
        dfs(t, now, dep+1);
        f[++cnt] = now;
        rmq[cnt] = dep;
void initlca(int root, int n)
    cnt = 0:
    dfs(root, root, 0);
    st. init (2*n-1);
int querylca(int x, int y)
    return f[st.query(p[x], p[y])];
```

## 欧拉路径

- ①无向图:连通,每个顶点度数都为偶数或者仅有两个点的度数为奇数。
- ②有向图: 底图连通,所有顶点的出度 == 入度 或者 有且仅有一个顶点 出度 == 入度+1,有且仅有一个顶点 入度 == 出度+1。
- ③混合图:底图连通,存在欧拉回路,则一定存在欧拉路径。否则如果有且仅有两个点的(出度-入度)是奇数,那么给这两个点加边,判断是否存在欧拉回路。

#### 欧拉回路

- ①无向图:连通,每个顶点度数都为偶数。
- ②有向图: 底图连通,每个顶点出度等于入度。
- ③混合图:底图连通,网络流判断。

## 无向图欧拉路径

```
vector<int> ans;
struct edge
{
    int to, next, id, dir, flag;
}e[20005];

void addedge(int u, int v, int id)
{
    e[cnt]. to = v;
    e[cnt]. id = id;
    e[cnt]. dir = 0;
    e[cnt]. flag = 0;
    e[cnt]. next = head[u];
    head[u] = cnt++;
    e[cnt]. to = u;
    e[cnt]. id = id;
```

```
#define MAXN 250
vector(int) ans;
struct edge
    int to, next, id, flag;
}e[20005];
void addedge(int u, int v, int id)
    e[cnt]. to = v;
    e[cnt].id = id;
    e[cnt].flag = 0;
    e[cnt].next = head[u];
    head[u] = cnt++;
vector < edge > ans;
void dfs(int u)
    for (int i = head[u]; i != -1; i = e[i]. next)
        if(!e[i]. flag)
             e[i]. flag = 1;
             dfs(e[i].to);
             ans. push_back(i);
```

#### 混合图欧拉回路判断

```
struct edge
{
    int u, v, cap, next;
}e[20005];
int n, m, cnt, head[N], dep[N], gap[N], cur[N], st[N], que[N], ind[N], outd[N];
int main()
{
    int T;
    scanf("%d", &T);
    while(T--)
    {
        scanf("%d%d", &n, &m);
    }
}
```

```
cnt = 0;
    memset (head, -1, sizeof (head));
    memset (ind, 0, sizeof (ind));
    memset(outd, 0, sizeof(outd));
    while (m--)
         int x, y, z;
scanf ("%d%d%d", &x, &y, &z);
outd[x]++;
         ind[y]++;
if(z == 0) addedge(x, y, 1);
    int sum = 0, flag = 1;
    for (int i = 1; i \le n; i++)
         if (outd[i]-ind[i] > 0)
              addedge(0, i, (outd[i]-ind[i])/2);
              sum += (outd[i]-ind[i])/2;
                  if(ind[i]-outd[i] > 0) addedge(i, n+1, (ind[i]-outd[i])/2);
         if((outd[i]-ind[i])%2) flag = 0;
    if(flag \&\& isap(0, n+1, n+2) == sum)
                                                printf("possible\n");
             printf("impossible\n");
return 0;
```

# 数据结构

```
并查集
void init()
   for (int i = 1; i \le n; i++)
                             pre[i] = i;
int findd(int x)
                      递归
   return pre[x] == x?x:pre[x] = findd(pre[x]);
void join(int a, int b)
   int x = findd(a), y = findd(b);
   if(x != y) pre[x] = y;
                                     带权值并查集
const int N = 50005;
int pre[N], rk[N];
   解题思路: father[y] = x
    如果 rank[y]为 0 代表 x 和 y 属于同一种
    如果 rank[y]为1代表x吃y
    如果 rank[y]为 2 代表 y 吃 x
   本题的关键就是如何在路径压缩的过程中更新关系权值
    需要总结更新公式
int findd(int x)
   if(pre[x] == x) return x;
   int xx = findd(pre[x]);
   rk[x] = (rk[x]+rk[pre[x]])%3;
   pre[x] = xx;
   return xx;
bool ok(int n, int type, int x, int y)
   if (x > n \mid y > n) return 0;
   int fx = findd(x), fy = findd(y);
   if(fx == fy)
       if((rk[y]-rk[x]+3)\%3 != type-1) return 0;
       return 1;
   else
       pre[fy] = fx;
       rk[fy] = (rk[x]-rk[y]+type-1+3)%3;
       return 1;
int main()
   int n, k;
scanf("%d %d", &n, &k);
   for (int i = 1; i \le n; i++)
       pre[i] = i;
       rk[i] = 0;
```

```
int cnt = 0;
           for (int i = 0; i < k; i++)
                      int d, x, y;
scanf("%d %d %d", &d, &x, &y);
                      if(!ok(n, d, x, y))
                                                                       cnt++;
           printf("%d\n", cnt);
           return 0;
                                                                                                                      ·维 rmq
struct ST
           int mm[100005], dp[100005][20];
           void init()
                      mm \lceil 0 \rceil = -1;
                      for (int i = 1; i \le n; i++)
                                  mm[i] = ((i&(i-1)) == 0)?mm[i-1]+1:mm[i-1];
                                  dp[i][0] = a[i];
                      for (int j = 1; j \le mm[n]; j++)
                                  for (int i = 1; i+(1 << j)-1 <= n; i++) dp[i][j] = min(dp[i][j-1], dp[i+(1 << (j-1)) dp[i+(1 << ((j-1)) dp[i+(1 << ((j-1)
1))][j-1]);
           int query (int a, int b)
                      if(a > b)
                                                      swap (a, b);
                      int k = mm[b-a+1]:
                      return min(dp[a][k], dp[b-(1<<k)+1][k]);
}st;
                                                                                                                   二维 rmq
//同一维一样 用 dp[row][col][i][j]表示(row, col)到(row+2^i, col+2^j)矩形内的最大值
int a[305][305], mm[305], dp[305][305][9][9];
void init(int n, int m)
           for (int i = 1; i \le 300; i++) mm[i] = ((i&(i-1)) == 0)?mm[i-1]+1:mm[i-1];
           for (int i = 1; i \le n; i++)
                                                                                                    dp[i][j][0][0] = a[i][j];
                      for (int j = 1; j \le m; j++)
           for (int i = 0; i \le mm[n]; i++)
                      for (int j = 0; j \le mm[m]; j++)
                                  if (i==0 && j==0)
                                                                                         continue;
                                  for (int row = 1; row+(1<i)-1 <= n; row++)
                                             for (int col = 1; col+(1<<j)-1 <= m; col++)
                                                                               dp[row][col][i][j] = max(dp[row][col][i-
                                                         if(i)
1][j], dp[row+(1 << (i-1))][co1][i-1][j]);
                                                                              dp[row][col][i][j] = max(dp[row][col][i][j-
                                                         else
1], dp[row][col+(1 << (j-1))][i][j-1]);
```

```
int rmq(int x1, int y1, int x2, int y2)
{
    int k1 = mm[x2-x1+1], k2 = mm[y2-y1+1];
    x2 = x2-(1<<k1)+1, y2 = y2-(1<<k2)+1;
    return
max(max(dp[x1][y1][k1][k2], dp[x1][y2][k1][k2]), max(dp[x2][y1][k1][k2], dp[x2][y2][k1][k2]));
}
</pre>
```

单调栈

```
stack<long long> s;
long long a[100005];
int main()
    int n;
    while (scanf ("%d", &n) && n)
        while(!s.empty())
                             s. pop();
        long long ans = 0;
        for(int i = 1; i <= n; i++) scanf("%11d", &a[i]);
        a[n+1] = 0;
        for (int i = 1; i \le n+1; i++)
             while (!s. empty() && a[s. top()] > a[i])
                 int temp = s. top();
                 s. pop();
                 int len = s. empty()?i-1:i-s. top()-1;
                 ans = max(ans, len*a[temp]);
             s. push(i);
        printf("%11d\n", ans);
    return 0;
```

单调队列

```
//XD0J1156
struct st
    int num, t;
}s;
deque(st) q;
int n, x;
int main()
    while ("scanf ("%d", &n))
         while(!q.empty()) q.pop_back();
         int out = 1, cnt = 1;
         for (int i = 1; i \le n; i++)
             scanf ("%d", &x);
             switch(x)
                  case 1:
                      int tt;
scanf("%d", &tt);
                      s.t = tt-i;
                      s.num = cnt++;
                      while(!q.empty() && q.back().t <= s.t)</pre>
                                                                      q. pop back();
                      q. push back(s);
```

```
break;
                 case 2:
                     if (q. front(). num == out++) q. pop front();
                     break:
                 case 3:
                     printf("%d\n", q. front(). t+i);
    return 0;
                                           维树状数组
int tree[50005], n;
inline int lowbit(int x)
    return x & (-x);
void update(int pos, int x)
    while (pos \leq n)
        tree[pos] += x;
        pos += lowbit(pos);
int getsum(int pos)
    int sum = 0;
    while (pos > 0)
        sum += tree[pos];
        pos -= lowbit(pos);
    return sum;
                                           维树状数组
int tree[1005][1005] = \{0\}, a[1005][1005] = \{0\};
int lowbit(int x)
    return x&(-x);
void update(int x, int y, int w)
    for (int i = x; i \le 1001; i += 1 \text{ owbit } (i))
        for(int j = y; j \le 1001; j += lowbit(j)) tree[i][j] += w;
int getsum(int x, int y)
    int ans = 0;
    for(int i = x; i > 0; i = lowbit(i))
        for(int j = y; j > 0; j = lowbit(j)) ans += tree[i][j];
    return ans;
```

划分树查询区间第 k 大

```
int n, m, tree [20] [N], sorted [N], to left [20] [N];
void build(int 1, int r, int dep)
    if(1 == r)return;
    int mid = (1+r)/2, same = mid-1+1;
    for (int i = 1; i \le r; i++)
         if(tree[dep][i] < sorted[mid])</pre>
                                            same--:
    int lpos = 1, rpos = mid+1;
    for (int i = 1; i \le r; i++)
         if(tree[dep][i] < sorted[mid]) tree[dep+1][1pos++] = tree[dep][i];</pre>
         else if(tree[dep][i] == sorted[mid] && same > 0)
             tree[dep+1][lpos++] = tree[dep][i];
             same--;
                 tree[dep+1][rpos++] = tree[dep][i];
         toleft[dep][i] = toleft[dep][1-1]+1pos-1;
    build (1, mid, dep+1);
    build (mid+1, r, dep+1);
int query (int 1, int r, int q1, int qr, int dep, int k)
    if(q1 == qr)
                      return tree[dep][q1];
    int mid = (1+r)/2, cnt = toleft[dep][qr]-toleft[dep][q1-1];
    if(cnt >= k)
         int 11 = 1 + \text{toleft[dep][q1-1]} - \text{toleft[dep][1-1]}, \text{rr} = 11 + \text{cnt-1};
         return query (1, mid, 11, rr, dep+1, k);
    else
         int rr = qr+toleft[dep][r]-toleft[dep][qr], 11 = rr-(qr-ql-cnt);
         return query (mid+1, r, 11, rr, dep+1, k-cnt);
int main()
    while (scanf ("%d%d", &n, &m)!=EOF)
         for (int i = 1; i \le n; i++)
             scanf("%d", &tree[0][i]);
             sorted[i] = tree[0][i];
         sort (sorted+1, sorted+n+1);
         build(1, n, 0);
         while(m--)
             int 1, r, k;
             scanf("%d%d%d", &1, &r, &k);
             printf(\sqrt[m]{d} n, query(1, n, 1, r, 0, k));
    return 0;
```

点分治求边权和等于 k 的数量

```
int n, k, root, son[10005], vis[10005], minn, ans, sum, dis[10005], dep[10005];
struct xx
    int to, w;
    xx(int a, int b):to(a), w(b) {};
vector \langle xx \rangle v[10005];
void dfsroot(int now, int pre)
    son[now] = 1;
    int maxx = 0;
    for (int i = 0; i < v[now]. size(); i++)
         int t = v[now][i]. to, w = v[now][i]. w;
         if(t == pre || vis[t]) continue;
         dfsroot(t, now);
         son[now] += son[t];
         \max x = \max(\max x, son[t]);
    \max = \max(\max, \text{sum-son}[\text{now}]);
    if (maxx < minn)</pre>
        minn = maxx;
        root = now;
void dfsdep(int now, int pre)
    dep[++dep[0]] = dis[now];
    for (int i = 0; i < v[now]. size(); i++)
         int t = v[now][i]. to, w = v[now][i]. w;
         if(t == pre || vis[t]) continue;
         dis[t] = dis[now]+w;
         dfsdep(t, now);
int calc(int now, int w)
    dis[now] = w;
    dep[0] = 0;
    dfsdep(now, 0);
    sort(dep+1, dep+dep[0]+1);
    int 1 = 1, r = dep[0], cnt = 0;
    while (1 < r)
        if(dep[1]+dep[r] \le k)
             cnt += r-1;
             1++;
         else r--;
    return cnt;
void solve(int now)
    ans += calc(now, 0);
```

```
vis[now] = 1;
    for (int i = 0; i < v[now]. size (); i++)
         int t = v[now][i]. to, w = v[now][i]. w;
         if(vis[t]) continue;
         ans -= calc(t, w);
         minn = 1e9;
         sum = son[t];
         dfsroot(t, 0);
         solve (root);
int main()
    while (scanf ("%d%d", &n, &k))
         if(n == 0 | k == 0)
                                    break:
         ans = 0, sum = n;
         memset(vis, 0, sizeof(vis));
         for (int i = 1; i \le n; i++)
                                         v[i].clear();
         for (int i = 1; i < n; i++)
             int x, y, z;
scanf("%d%d%d", &x, &y, &z);
             v[x]. push_back(xx(y, z));
             v[y]. push_back(xx(x, z));
         minn = 1e9;
         df sroot (1, 0);
         solve(root);
printf("%d\n", ans);
    return 0;
                                        点分治求点权积等于 k
int n, k, root, son[100005], vis[100005], dep[100005], id[1000005], minn, ans1, ans2, sum;
long long a[100005], inv[1000005], mp[1000005];
vector<int> v[100005];
void dfsroot(int now, int pre)
    son[now] = 1;
    int maxx = 0;
    for (int i = 0; i < v[now]. size(); i++)
         int t = v[now][i];
         if(t == pre || vis[t]) continue;
         dfsroot(t, now);
         son[now] += son[t];
         \max x = \max(\max x, \operatorname{son}[t]);
    \max = \max(\max, \text{sum-son}[\text{now}]);
    if (maxx < minn)</pre>
         minn = maxx;
         root = now;
```

void dfsdep (int now, int pre, long long x)

dep[++dep[0]] = x\*a[now]%MOD;

```
id[dep[0]] = now;
    for (int i = 0; i < v[now]. size (); i++)
        int t = v[now][i];
        if(t == pre || vis[t]) continue;
        dfsdep(t, now, x*a[now]%MOD);
void update(int now, int x, int y)
    int t = mp[inv[x*a[now]%MOD]*k%MOD];
    if(!t) return;
    if(t > y) swap(t, y);
    if (t < ans1 \mid | t == ans1 \&\& y < ans2)
        ans1 = t;
        ans2 = y;
void solve(int now)
    vis[now] = 1;
    mp[1] = now;
    for (int i = 0; i < v[now]. size(); i++)
        int t = v[now][i];
        if(vis[t])
                   continue;
        dep[0] = 0;
        dfsdep(t, now, 1);
       for (int j = 1; j \le dep[0]; j++) update (now, dep[j], id[j]);
        for (int j = 1; j \le dep[0]; j++)
            if(!mp[dep[j]] || mp[dep[j]] > id[j])     mp[dep[j]] = id[j];
    mp[1] = 0;
    for (int i = 0; i < v[now]. size (); i++)
        int t = v[now][i];
        if(vis[t]) continue;
        dep[0] = 0;
        dfsdep(t, now, 1);
        for (int j = 1; j \le dep[0]; j++) mp[dep[j]] = 0;
    for (int i = 0; i < v[now]. size(); i++)
        int t = v[now][i];
        if(vis[t]) continue;
        minn = 1e9;
        sum = son[t];
        dfsroot(t, 0);
        solve (root):
int main()
    inv[1] = 1;
    ans1 = 1e9, ans2 = 0;
```

```
memset(vis, 0, sizeof(vis));
                                          v[i].clear();
scanf("%11d",&a[i]);
         for (int i = 1; i \le n; i++)
         for (int i = 1; i \le n; i++)
         for (int i = 1; i < n; i++)
              int x, y;
              scanf ("%d%d", &x, &y);
              v[x]. push_back(y);
              v[y]. push_back(x);
         minn = 1e9;
         sum = n;
         dfsroot(1,0);
         solve (root);
         if(ans1 == 1e9) printf("No solution\n");
                  printf("%d %d\n", ans1, ans2);
    return 0;
                                          cdq 分治求三维偏序
int n, maxx, tree[100005], ans[100005];
struct xx
    int x, y, z, id;
    xx(int a, int b, int c, int d):x(a), y(b), z(c), id(d) {};
    xx() \{\};
}a[100005],b[100005];
inline lowbit(int x)
    return x&-x;
void add(int pos, int x)
    while (pos \leq maxx)
         tree[pos] += x;
         pos += lowbit(pos);
int getsum(int pos)
    int ans = 0;
    while (pos > 0)
         ans += tree[pos];
         pos -= lowbit(pos);
    return ans;
bool cmp1 (xx a, xx b)
    if(a. x != b. x)     return a. x < b. x;
if(a. y != b. y)     return a. y < b. y;</pre>
    return a. z < b. z;
bool cmp2 (xx a, xx b)
    if (a. y != b. y) return a. y < b. y;
```

```
void cdq(int 1, int r)
    if(1 == r) return;
    int mid = (1+r)/2;
    int cnt = 0;
    for (int i = 1; i \le mid; i++) b[++cnt] = xx(0, a[i], y, a[i], z, 0);
    for (int i = mid+1; i \le r; i++) b[++cnt] = xx(0, a[i], y, a[i], z, a[i], id);
    sort(b+1, b+1+cnt, cmp2);
    for (int i = 1; i \le cnt; i++)
        if (b[i]. id == 0) add (b[i]. z, 1);
        else ans[b[i].id] += getsum(b[i].z);
    for (int i = 1; i \le cnt; i++)
        if(b[i].id == 0) add(b[i].z, -1);
    cdq(1, mid);
    cdq(mid+1,r);
int main()
    int T;
    scanf ("%d", &T);
    while (T--)
        scanf ("%d", &n);
        \max x = 0;
        for (int i = 1; i \le n; i++)
            scanf("%d%d%d", &a[i].x, &a[i].y, &a[i].z);
            a[i].id = i;
            \max x = \max(\max x, a[i].z);
        sort(a+1, a+1+n, cmp1);
        memset (ans, 0, sizeof (ans));
        int cnt = 0;
        for (int i = n-1; i >= 1; i--)
            if(a[i].x == a[i+1].x && a[i].y == a[i+1].y && a[i].z == a[i+1].z) cnt++;
            else cnt = 0;
            ans[a[i].id] += cnt;
        cdq(1, n);
        for (int i = 1; i \le n; i++) printf ("%d\n", ans [i]);
    return 0;
                                     树链剖分线段树(点权)
//点权,修改路径点权,查询单点权
int n, m, q, cnt, a[N], top[N], fa[N], dep[N], num[N], p[N], fp[N], son[N];
vector<int> v[N];
char s[10];
struct segtree
    int 1, r;
    long long x, lazy;
tree[4*N];
```

return a. id < b. id;

```
inline void pushup (int pos)
    tree[pos].x = tree[pos << 1].x + tree[pos << 1|1].x;
inline void pushdown (int pos)
    if (tree[pos]. lazy)
         tree[pos <<1].x += (tree[pos <<1].r-tree[pos <<1].1+1)*tree[pos].1azy;
         tree[pos <<1|1].x += (tree[pos <<1|1].r-tree[pos <<1|1].1+1)*tree[pos].1azy;
         tree[pos<<1].lazy += tree[pos].lazy;
         tree[pos<<1|1].lazy += tree[pos].lazy;
         tree[pos].1azy = 0;
void build(int pos, int 1, int r)
    tree[pos].1 = 1;
    tree[pos].r = r;
    tree[pos].1azy = 0;
    if(1 == r)
         tree[pos].x = a[fp[1]];
         return;
    int mid = (1+r)/2;
    build(pos<<1, 1, mid);
    build (pos <<1 | 1, mid+1, r);
    pushup (pos);
void update(int pos, int 1, int r, long long x)
    if(r < tree[pos].1 || tree[pos].r < 1) return;</pre>
    if(1 \le tree[pos].1 \&\& tree[pos].r \le r)
         tree[pos].x += (tree[pos].r-tree[pos].1+1)*x;
         tree[pos].lazy += x;
         return;
    pushdown (pos);
    update (pos <<1, 1, r, x);
    update (pos<<1 | 1, 1, r, x);
    pushup (pos);
long long getsum(int pos, int 1, int r)
    if (r < tree[pos].1 \mid | tree[pos].r < 1) return 0;
    if (1 \le \text{tree}[pos].1 \&\& \text{tree}[pos].r \le r)
                                                    return tree[pos].x;
    pushdown (pos):
    return getsum(pos\langle\langle 1, 1, r \rangle+getsum(pos\langle\langle 1|1, 1, r \rangle);
void dfs (int now, int pre, int d)
    dep[now] = d;
    fa[now] = pre;
    num[now] = 1;
    for (int i = 0; i < v[now]. size(); i++)
```

```
int t = v[now][i];
        if(t == pre)
                        continue;
        dfs(t, now, d+1);
        void getpos(int now, int sp)
    top[now] = sp;
    p[now] = ++cnt;
    fp[p[now]] = now;
    if (son now == -1) return;
    getpos(son[now], sp);
    for (int i = 0; i < v[now]. size(); i++)
        int t = v[now][i];
                                             getpos(t,t);
        if(t != son[now] && t != fa[now])
void change(int u, int v, int x)
    int f1 = top[u], f2 = top[v];
    while (f1 != f2)
        if(dep[f1] < dep[f2])
            swap(f1, f2);
            swap(u, v);
        update(1, p[f1], p[u], x);
        u = fa[f1];
        f1 = top[u];
    if(dep[u] > dep[v]) swap(u, v);
    update (1, p[u], p[v], x);
int main()
    while ("scanf ("%d%d%d", &n, &m, &g))
        memset (son, -1, sizeof(son));
        memset(tree, 0, sizeof(tree));
        for (int i = 1; i \le n; i++) v[i]. clear();
        cnt = 0;
                                     scanf ("%d", &a[i]);
        for (int i = 1; i \le n; i++)
        while (m--)
            int x, y;
            scanf ("%d%d", &x, &y);
            v[x]. push_back(y);
            v[y]. push_back(x);
        dfs(1,0,0);
        getpos (1, 1);
        build(1, 1, cnt);
        while (q--)
            scanf("%s", s);
if(s[0] == 'Q')
```

```
int x;
                   scanf("%d", &x);
printf("%11d\n", getsum(1, p[x], p[x]));
              else
                   int x, y, z;
scanf("%d%d%d", &x, &y, &z);
                   if(s[0] == 'D') z = -z;
                   change (x, y, z);
    return 0;
                                          树链剖分线段树(边权)
int n, m, q, cnt, a[N], val[N], top[N], fa[N], dep[N], num[N], p[N], fp[N], son[N], e[N][3];
struct segtree
    int 1, r, x;
tree[4*N];
vector<int> v[N];
char s[10];
void pushup(int pos)
    tree[pos]. x = max(tree[pos << 1]. x, tree[pos << 1|1]. x);
void build(int pos, int 1, int r)
    tree[pos].1 = 1;
    tree[pos].r = r;
    if(1 == r)
         tree[pos].x = val[1];
         return;
    int mid = (1+r)/2;
    build(pos<<1, 1, mid);
    build(pos<<1|1, mid+1, r);
    pushup (pos);
void update(int pos, int 1, int r, int x)
    if(r < tree[pos].1 || tree[pos].r < 1) return;</pre>
    if(1 \le tree[pos].1 \&\& tree[pos].r \le r)
         tree[pos].x = x;
         return;
    update (pos <<1, 1, r, x);
    update (pos<<1 | 1, 1, r, x);
    pushup(pos);
int query (int pos, int 1, int r)
    if(r < tree[pos].1 || tree[pos].r < 1) return 0;</pre>
    if(1 \le tree[pos].1 \&\& tree[pos].r \le r)
                                                        return tree[pos].x;
    return \max(\text{query}(\text{pos}\langle\langle 1, 1, r), \text{query}(\text{pos}\langle\langle 1|1, 1, r));
```

```
void dfs(int now, int pre, int d)
    dep[now] = d;
    fa[now] = pre;
    num[now] = 1;
    for (int i = 0; i < v[now]. size(); i++)
         int t = v[now][i];
         if(t == pre)
                            continue;
         dfs(t, now, d+1);
         \begin{array}{lll} num[now] & += & num[t]; \\ if(son[now] & == & -1 & || & num[t] > & num[son[now]]) & son[now] & = & t; \end{array}
void getpos(int now, int sp)
    top[now] = sp;
    p[now] = ++cnt;
    fp[p[now]] = now;
    if(son[now] == -1) return;
    getpos(son[now], sp);
    for (int i = 0; i < v[now]. size(); i++)
         int t = v[now][i];
         if(t != son[now] \&\& t != fa[now]) getpos(t,t);
int solve(int u, int v)
    int f1 = top[u], f2 = top[v], t = 0;
    while (f1 != f2)
         if(dep[f1] < dep[f2])
              swap(f1, f2);
              swap(u, v);
         t = max(t, query(1, p[f1], p[u]));
         u = fa[f1];
         f1 = top[u];
    if(u == v) return t;
    if(dep[u] > dep[v]) swap(u, v);
return max(t, query(1, p[son[u]], p[v]));
int main()
    int T;
    scanf ("%d", &T);
    while(T--)
         scanf ("%d", &n);
         memset (son, -1, sizeof (son));
         for (int i = 1; i \le n; i++) v[i]. clear();
         cnt = 0;
         for (int i = 1; i < n; i++)
              scanf ("%d%d%d", &e[i][0], &e[i][1], &e[i][2]);
              v[e[i][0]]. push_back(e[i][1]);
              v[e[i][1]]. push back(e[i][0]);
```

## 线段树模版

```
struct segtree
    int 1, r;
    long long x, lazy;
} tree [400005];
inline void pushup(int pos)
    tree[pos]. x = tree[pos <<1]. x+tree[pos <<1|1]. x;
inline void pushdown (int pos)
    if (tree[pos]. lazy)
        tree[pos <<1].x += (tree[pos <<1].r-tree[pos <<1].1+1)*tree[pos].1azy;
        tree[pos <<1|1].x += (tree[pos <<1|1].r-tree[pos <<1|1].1+1)*tree[pos].lazy;
        tree[pos<<1].lazy += tree[pos].lazy;
        tree[pos<<1|1].lazy += tree[pos].lazy;
        tree[pos]. lazy = 0;
void build(int pos, int 1, int r)
    tree[pos].1 = 1;
    tree[pos].r = r;
    tree[pos]. lazy = 0;
    if(1 == r)
        tree[pos]. x = a[1];
        return:
    int mid = (1+r)/2;
    build(pos<<1,1,mid);
build(pos<<1|1,mid+1,r);
    pushup(pos);
void update(int pos, int 1, int r, long long x)
    if (r < tree[pos]. 1 | tree[pos]. r < 1) return;
```

```
if(1 \le tree[pos].1 \&\& tree[pos].r \le r)
         tree[pos]. x = (tree[pos]. r-tree[pos]. 1+1)*x;
         tree[pos].1azy = x;
         return;
    pushdown (pos);
    update (pos <<1, 1, r, x);
    update (pos <<1 | 1, 1, r, x);
    pushup(pos);
long long getsum(int pos, int 1, int r)
    if(r < tree[pos].1 || tree[pos].r < 1) return 0;</pre>
    if(1 <= tree[pos].1 && tree[pos].r <= r)</pre>
                                                    return tree[pos].x;
    pushdown (pos);
    return getsum(pos <<1, 1, r)+getsum(pos <<1 | 1, 1, r);
                                          动态开点线段树
int n, k, q, cnt = 0, a[100005];
struct segtree
    int 1, r, x, lazy;
    segtree *1son, *rson;
}*root;
segtree *newnode(int 1, int r)
    segtree *t = new segtree;
    t->1 = 1;
    t->r = r;
    t\rightarrow 1son = NULL;
    t\rightarrow rson = NULL;
    t\rightarrow 1azy = 0;
    t-\rangle_X = a[1];
    return t;
segtree *newlson(segtree *pos)
    int mid = (pos->1+pos->r)/2;
    return newnode (pos->1, mid);
segtree *newrson(segtree *pos)
    int mid = (pos->1+pos->r)/2;
    return newnode (mid+1, pos->r);
void pushup(segtree *pos)
    if(!pos->1son) pos->1son = new1son(pos);
    if (!pos-\rangle rson) pos-\rangle rson = newrson(pos);
    pos->x = min(pos->1son->x, pos->rson->x);
void pushdown(segtree *pos)
    if(!pos->1son)
                      pos \rightarrow 1son = new1son(pos);
    if (!pos->rson)
                      pos->rson = newrson(pos);
    if(pos \rightarrow lazy)
```

```
pos \rightarrow 1son \rightarrow x = pos \rightarrow 1azy;
          pos->rson->x = pos->lazy;
          pos \rightarrow 1son \rightarrow 1azy = pos \rightarrow 1azy;
          pos \rightarrow rson \rightarrow 1azy = pos \rightarrow 1azy;
          pos \rightarrow lazy = 0;
void update(segtree *pos, int 1, int r, int x)
     if (r < pos \rightarrow 1 \mid pos \rightarrow r < 1)
                                              return;
     if (1 \leq pos \rightarrow 1 \&\& pos \rightarrow r \leq r)
          pos->_X = x;
          pos \rightarrow lazy = x;
          return;
     pushdown(pos);
     update (pos\rightarrow1son, 1, r, x);
     update (pos\rightarrowrson, 1, r, x);
     pushup(pos);
int getmin(segtree *pos, int 1, int r)
     if (r < pos \rightarrow 1 \mid pos \rightarrow r < 1)
                                             return 1e9;
     if (1 \leq pos-)1 \&\& pos-)r \leq r return pos-)x;
     pushdown (pos);
     return min(getmin(pos\rightarrow1son, 1, r), getmin(pos\rightarrowrson, 1, r));
                                       区间取模,单点修改,询问区间和
int n, m;
long long a[100005];
struct segtree
     int 1, r;
     long long x, sum;
} tree[400005];
inline void pushup (int pos)
     tree[pos]. x = max(tree[pos << 1]. x, tree[pos << 1|1]. x);
     tree[pos].sum = tree[pos<<1].sum+tree[pos<<1|1].sum;
void build(int pos, int 1, int r)
     tree[pos].1 = 1;
     tree[pos].r = r;
     if(1 >= r)
          tree[pos].x = a[1];
          tree[pos].sum = a[1];
          return;
     int mid = (1+r)/2;
     build(pos<<1, 1, mid);
     build (pos <<1 | 1, mid+1, r);
     pushup (pos);
void edit(int pos, int 1, int r, long long x, int t)
```

```
if(t < 1 \mid | t > r) return;
    if(1 == r)
         tree[pos].x = x;
         tree[pos].sum = x;
         return;
    int mid = (1+r)/2;
    edit(pos<<1,1,mid,x,t);
edit(pos<<1|1,mid+1,r,x,t);
    pushup(pos);
void mod (int pos, int 1, int r, long long x)
    if (r < tree[pos].1 \mid | tree[pos].r < 1 \mid | tree[pos].x < x)
    if(tree[pos].1 == tree[pos].r)
         tree[pos].x \%= x;
         tree[pos].sum = tree[pos].x;
         return;
    mod(pos << 1, 1, r, x);
    mod(pos <<1 | 1, 1, r, x);
    pushup (pos);
long long getsum(int pos, int 1, int r)
    if(r < tree[pos].1 || tree[pos].r < 1) return 0;</pre>
    if(1 \le tree[pos].1 \&\& tree[pos].r \le r)
                                                       return tree pos sum;
    return getsum(pos\langle\langle 1, 1, r \rangle+getsum(pos\langle\langle 1 | 1, 1, r \rangle;
                                   区间加减,区间开根,询问区间和
int n, m;
long long a[100005];
struct segtree
    int 1, r;
    long long ma, mi, sum, lazy;
} tree [400005]:
inline void pushup (int pos)
    tree[pos]. ma = max(tree[pos << 1]. ma, tree[pos << 1|1]. ma);
    tree[pos].mi = min(tree[pos<<1].mi, tree[pos<<1|1].mi);
    tree[pos].sum = tree[pos<<1].sum+tree[pos<<1|1].sum;
inline void pushdown (int pos)
    if (tree[pos]. lazy)
         long long t = tree[pos].lazy;
         tree[pos<<1].lazy += t;
tree[pos<<1]1].lazy += t;</pre>
         tree[pos << 1].mi += t;
         tree[pos << 1|1].mi += t;
         tree[pos << 1]. ma += t;
         tree[pos<<1|1].ma += t;
         tree[pos << 1].sum += (tree[pos << 1].r-tree[pos << 1].1+1)*t;
         tree[pos <<1 | 1]. sum += (tree[pos <<1 | 1]. r-tree[pos <<1 | 1]. 1+1)*t;
```

```
tree[pos].1azy = 0;
void build(int pos, int 1, int r)
    tree|pos|.1 = 1;
    tree[pos].r = r;
    tree[pos].1azy = 0;
    if(1 == r)
         tree[pos].mi = a[1];
         tree[pos].ma = a[1];
         tree [pos]. sum = a[1];
         return;
    int mid = (1+r)/2;
    build(pos << 1, 1, mid);
    build (pos <<1 | 1, mid+1, r);
    pushup(pos);
long long getsum(int pos, int 1, int r)
    if(r < tree[pos].1 || tree[pos].r < 1) return 0;</pre>
    if(1 <= tree[pos].1 && tree[pos].r <= r) return tree[pos].sum;
    pushdown (pos);
    return getsum(pos\langle\langle 1, 1, r \rangle+getsum(pos\langle\langle 1|1, 1, r \rangle);
void update(int pos, int 1, int r, long long x)
    if (r < tree[pos].1 \mid | tree[pos].r < 1)
                                                    return;
    if(1 \le tree[pos].1 \&\& tree[pos].r \le r)
         tree pos . mi += x;
         tree[pos].ma += x;
         tree[pos].sum += (tree[pos].r-tree[pos].1+1)*x;
         tree[pos].lazy += x;
         return;
    pushdown(pos);
    update (pos<<1, 1, r, x);
    update (pos<<1 | 1, 1, r, x);
    pushup (pos);
void sq(int pos, int 1, int r)
    if(r < tree[pos].1 \mid | tree[pos].r < 1)
                                                   return;
    if(1 \le tree[pos].1 \&\& tree[pos].r \le r \&\& (long long) sqrt(tree[pos].ma) - (long long) sqrt(tree[pos].ma)
long)sqrt(tree[pos].mi) == tree[pos].ma-tree[pos].mi)
         long long t = tree[pos].ma-(long long)sqrt(tree[pos].ma);
         tree[pos].mi -= t;
         tree[pos].ma -= t;
         tree[pos].sum -= (tree[pos].r-tree[pos].1+1)*t;
         tree[pos]. lazy -= t;
         return;
    pushdown (pos);
    sq(pos <<1, 1, r);
    sq(pos << 1 | 1, 1, r);
    pushup(pos);
```

### 区间最小值,询问区间最大值,询问区间和

```
int n, m;
long long a[1000005];
struct segtree
    int 1, r, max1, max2, num;
    long long sum;
} tree [4000005];
inline void pushup (int pos)
    tree[pos]. max1 = max(tree[pos << 1]. max1, tree[pos << 1]1]. max1);
    tree[pos].sum = tree[pos<<1].sum+tree[pos<<1|1].sum;
    tree[pos].num = 0;
    tree[pos]. max2 = max(tree[pos << 1]. max2, tree[pos << 1|1]. max2);
    if(tree[pos].max1 == tree[pos << 1].max1) tree[pos].num += tree[pos << 1].num;
             tree[pos]. max2 = max(tree[pos]. max2, tree[pos<<1]. max1);
    if(tree[pos]. max1 == tree[pos << 1 | 1]. max1) tree[pos]. num += tree[pos << 1 | 1]. num;
             tree[pos]. max2 = max(tree[pos]. max2, tree[pos<<1|1]. max1);
inline void pushdown (int pos)
    if(tree[pos].max1 < tree[pos << 1].max1)
         tree[pos<<1].sum -= (long long) tree[pos<<1].num*(tree[pos<<1].max1-
tree[pos].max1);
         tree[pos << 1]. max1 = tree[pos]. max1;
    if (\text{tree}[\text{pos}]. \text{max}1 < \text{tree}[\text{pos} << 1 | 1]. \text{max}1)
         tree[pos <<1|1]. sum == (long long) tree[pos <<1|1]. num*(tree[pos <<1|1]. max1=
tree[pos].max1);
         tree[pos << 1 | 1]. max1 = tree[pos]. max1;
void build(int pos, int 1, int r)
    tree[pos].1 = 1;
    tree[pos].r = r;
    if(1 == r)
         scanf ("%d", &tree[pos]. max1);
         tree[pos].sum = tree[pos].max1;
         tree[pos]. \max 2 = -1;
         tree[pos].num = 1;
         return;
    int mid = (1+r)/2;
    build(pos<<1, 1, mid);
    build (pos<<1|1, mid+1, r);
    pushup (pos);
void update(int pos, int 1, int r, int x)
    if(r < tree[pos].1 \mid | tree[pos].r < 1 \mid | x >= tree[pos].max1)
    if (1 \le \text{tree[pos]}.1 \&\& \text{tree[pos]}.r \le r \&\& \text{tree[pos]}.max2 < x)
         tree[pos].sum -= (long long) tree[pos].num*(tree[pos].max1-x);
         tree[pos].max1 = x;
```

```
return;
    pushdown (pos);
    update (pos <<1, 1, r, x);
    update (pos<<1 | 1, 1, r, x);
    pushup (pos);
long long getsum(int pos, int 1, int r)
    if (r < tree[pos].1 \mid | tree[pos].r < 1) return 0;
    if(1 \le tree[pos].1 \&\& tree[pos].r \le r)
                                                    return tree[pos].sum;
    pushdown (pos);
    return getsum(pos\langle\langle 1, 1, r \rangle)+getsum(pos\langle\langle 1|1, 1, r \rangle);
int getmax (int pos, int 1, int r)
    if(r < tree[pos].1 || tree[pos].r < 1) return 0;</pre>
    if(1 \le tree[pos].1 \&\& tree[pos].r \le r)
                                                   return tree[pos].max1;
    pushdown (pos);
    return \max(\text{getmax}(\text{pos}\langle\langle 1, 1, r), \text{getmax}(\text{pos}\langle\langle 1|1, 1, r))):
                      区间加减,区间覆盖,询问区间最大值,询问历史区间最大值
int n, m;
char s[10];
long long a[1000005];
struct segtree
    int 1, r, nmax, pmax, nlazy, plazy, ncov, pcov;
} tree [4000005];
inline void pushup (int pos)
    tree[pos]. nmax = max(tree[1s]. nmax, tree[rs]. nmax);
    tree[pos].pmax = max(tree[1s].pmax, tree[rs].pmax);
inline void pushdown (int pos)
    tree[1s].pmax =
max(tree[1s].pmax, max(tree[pos].pcov, tree[pos].plazy+tree[1s].nmax));
    tree[rs].pmax =
max(tree[rs].pmax, max(tree[pos].pcov, tree[pos].plazy+tree[rs].nmax));
    if (tree pos ncov == -INF)
         tree[1s].nmax += tree[pos].nlazy;
         tree[rs].nmax += tree[pos].nlazy;
         if(tree[1s].ncov == -INF)
             tree[ls].plazy = max(tree[ls].plazy, tree[ls].nlazy+tree[pos].plazy);
             tree[1s]. nlazy += tree[pos]. nlazy;
         else
             tree[1s].pcov = max(tree[1s].pcov, tree[1s].ncov+tree[pos].plazy);
             tree[1s].ncov = tree[1s].nmax;
         if(tree[rs].ncov == -INF)
             tree[rs].plazy = max(tree[rs].plazy, tree[rs].nlazy+tree[pos].plazy);
             tree[rs].nlazy += tree[pos].nlazy;
```

```
else
            tree[rs].pcov = max(tree[rs].pcov, tree[rs].ncov+tree[pos].plazy);
            tree[rs].ncov = tree[rs].nmax;
    else
        max(tree[1s].plazy, tree[1s].nlazy+tree[pos].plazy);
               tree[1s].pcov = max(tree[1s].pcov, tree[1s].nmax+tree[pos].plazy);
        if(tree[rs].ncov == -INF)
                                   tree[rs].plazy =
max(tree[rs].plazy, tree[rs].nlazy+tree[pos].plazy);
                tree[rs].pcov = max(tree[rs].pcov, tree[rs].nmax+tree[pos].plazy);
        tree[1s].nmax = tree[pos].ncov;
        tree[rs].nmax = tree[pos].ncov;
        tree[1s].ncov = tree[pos].ncov;
        tree[rs].ncov = tree[pos].ncov;
        tree[1s].pcov = max(tree[1s].pcov, tree[pos].pcov);
        tree[rs].pcov = max(tree[rs].pcov, tree[pos].pcov);
    tree[pos].nlazy = 0;
    tree[pos].plazy = 0;
    tree[pos].ncov = -INF;
    tree [pos]. pcov = -INF;
void build(int pos, int 1, int r)
    tree[pos].1 = 1;
    tree[pos].r = r;
    tree[pos].ncov = -INF;
    tree[pos].pcov = -INF;
    tree[pos].nlazy = 0;
    tree[pos].plazy = 0;
    if(1 == r)
        scanf ("%d", &tree[pos]. nmax);
        tree[pos].pmax = tree[pos].nmax;
        return;
    int mid = (1+r)/2;
    build(1s, 1, mid);
    build(rs, mid+1, r);
    pushup (pos);
long long getans (int pos, int 1, int r, int t)
    if (r < tree[pos].1 | tree[pos].r < 1) return -INF;
    if(1 \le tree[pos].1 \&\& tree[pos].r \le r)
                                                return t?tree[pos].pmax:tree[pos].nmax;
    pushdown(pos);
    return \max(\text{getans}(1s, 1, r, t), \text{getans}(rs, 1, r, t));
void update1(int pos, int 1, int r, int x)
    if (r < tree[pos]. 1 | tree[pos]. r < 1) return;
    if(1 \le tree[pos].1 \&\& tree[pos].r \le r)
        tree[pos].nmax += x;
        tree[pos].pmax = max(tree[pos].pmax, tree[pos].nmax);
        if (tree[pos].ncov == -INF)
```

```
tree[pos].nlazy += x;
             tree[pos].plazy = max(tree[pos].plazy, tree[pos].nlazy);
        else
             tree[pos].ncov = tree[pos].nmax;
             tree[pos].pcov = max(tree[pos].pcov, tree[pos].ncov);
        return;
    pushdown (pos);
    updatel(ls, l, r, x);
    update1(rs, 1, r, x);
    pushup (pos);
void update2(int pos, int 1, int r, int x)
    if(r < tree[pos].1 || tree[pos].r < 1) return;</pre>
    if (1 \le \text{tree}[pos].1 \&\& \text{tree}[pos].r \le r)
        tree[pos].nmax = x;
        tree[pos].pmax = max(tree[pos].pmax, tree[pos].nmax);
        tree[pos].ncov = x;
        tree[pos].pcov = max(tree[pos].pcov, tree[pos].ncov);
        return;
    pushdown (pos);
    update2(1s, 1, r, x);
    update2(rs, 1, r, x);
    pushup(pos);
                                           块状数组
int n, m, a[200005], cnt[200005], next[200005], belong[200005], 1[1005], r[1005];
void build()
    int len = sqrt(n);
    int num = (n+1en-1)/1en;
    if (n%len) num++;
    for (int i = 1; i \le num; i++)
        1[i] = (i-1)*len+1;
        r[i] = i*len;
    r[num] = n;
    for (int i = 1; i \le n; i++)
                                  belong[i] = (i-1)/len+1;
    for (int i = n; i >= 1; i--)
        int t = i+a[i];
        if(t > n)
             cnt[i] = 1;
             next[i] = -1;
        else if(belong[i] == belong[t])
             cnt[i] = cnt[t]+1;
             next[i] = next[t];
        else
             cnt[i] = 1;
```

```
next[i] = t;
void update(int pos, int x)
    a[pos] = x;
    for (int i = pos; i >= 1[belong[pos]]; i--)
        int t = i+a[i];
        if(t > n)
             cnt[i] = 1;
             next[i] = -1;
        else if(belong[i] == belong[t])
             cnt[i] = cnt[t]+1;
             next[i] = next[t];
        else
             cnt[i] = 1;
next[i] = t;
int getsum(int x)
    int ans = cnt[x];
    while (next[x] > 0)
        x = next[x];
        ans += cnt[x];
    return ans;
                                              莫队
int a[30005], num[30005], n, m;
long long ans[30005];
struct node
    int 1, r, num, belong;
    friend bool operator (node a, node b)
        if (a. belong == b. belong)
                                     return a.r < b.r;
        return a. 1 < b. 1;
} query [30005];
int main()
    while ("scanf ("%d%d", &n, &m))
        int size = sqrt(n);
        for (int i = 1; i \le n; i++)
                                         scanf ("%d", &a[i]);
        for (int i = 1; i \le m; i++)
             scanf("%d%d", &query[i].1, &query[i].r);
query[i].num = i;
             query[i].belong = (query[i].1-1)/size+1;
```

```
sort (query+1, query+1+m);
    memset (num, 0, sizeof (num));
    int 11 = 1, rr = 1;
    LL now = 1;
    num[a[1]]++;
    for (int i = 1; i \le m; i++)
         while(rr < query[i].r)</pre>
             rr++;
             num[a[rr]]++;
             //\text{now} = ;
         while (11 > query[i].1)
             11--;
             num[a[11]]++;
             //now = ;
         while (11 < query[i].1)
             //\text{now} = ;
             num[a[11]]--;
             11++;
         while(rr > query[i].r)
             //\text{now} = ;
             num[a[rr]]--;
             rr--;
         ans[query[i].num] = now;
    for (int i = 1; i \le m; i++) printf ("%11d\n", ans [i]);
return 0;
```

## 主席树 静态区间第 k 大

```
int n, q, m, tot, a[N], b[N];
int tree[N], 1son[N*30], rson[N*30], c[N*30];
void init hash()
    for (int i = 1; i \le n; i++) b[i]=a[i];
    sort(b+1, b+n+1);
    m = unique(b+1, b+n+1)-b-1;
int gethash(int x)
    return lower bound (b+1, b+1+m, x)-b;
void build(int &now, int 1, int r)
    now = tot++;
    c[now] = 0;
    if(1 == r) return;
    int mid = (1+r)/2;
    build(lson[now], 1, mid);
    build(rson[now], mid+1, r);
void update (int &now, int last, int l, int r, int pos, int x)
    now = tot++;
    lson[now] = lson[last];
    rson[now] = rson[last];
    c[now] = c[1ast] + x;
    if(1 == r) return;
    int mid = (1+r)/2;
    if (pos \leq mid) update (1son[now], 1son[1ast], 1, mid, pos, x);
             update(rson[now], rson[last], mid+1, r, pos, x);
int query (int 1t, int rt, int 1, int r, int k)
    if(1 == r) return 1;
    int mid = (1+r)/2;
    if(c[1son[rt]]-c[1son[1t]] >= k)
                                           return query (lson[lt], lson[rt], l, mid, k);
    return query(rson[lt], rson[rt], mid+1, r, k-c[lson[rt]]+c[lson[lt]]);
int main()
    ios::sync with stdio(0);
    while ("scanf ("%d%d", &n, &q))
                                       scanf ("%d", &a[i]);
        for (int i = 1; i \le n; i++)
        init_hash();
        tot = 0;
        build(tree[0], 1, m);
                                        update(tree[i], tree[i-1], 1, m, gethash(a[i]), 1);
        for (int i = 1; i \le n; i++)
        while (q--)
             int 1, r, k;
             scanf ("%d%d%d", &1, &r, &k);
             printf("%d\n", b[query(tree[1-1], tree[r], 1, m, k)]);
```

```
return 0;
                                   静态区间查询不同的数个数
int n, q, m, tot, a[N], b[N];
int tree[N], lson[N*100], rson[N*100], c[N*100];
map<int, int> mp;
int build(int 1, int r)
    int now = tot++;
    c[now] = 0;
    if(1 == r) return now;
    int mid = (1+r)/2;
    lson[now] = build(1, mid);
    rson[now] = build(mid+1, r);
    return now;
int update(int last, int pos, int x)
    int now = tot++, t = now;
    c[now] = c[1ast] + x;
    int 1 = 1, r = n;
    while (1 < r)
        int mid = (1+r)/2;
        if(pos \le mid)
            lson[now] = tot++;
            rson[now] = rson[last];
            now = 1son[now];
            last = lson[last];
            r = mid;
        else
            rson[now] = tot++;
            lson[now] = lson[last];
            now = rson[now];
            last = rson[last];
            1 = mid+1:
        c[now] = c[1ast] + x;
    return t;
int query (int pos, int now, int 1, int r)
    if(1 == r) return c[now];
    int mid = (1+r)/2;
    if (pos <= mid) return query (pos, 1son[now], 1, mid) +c[rson[now]];
    return query(pos, rson[now], mid+1, r);
int main()
    while ("scanf ("%d", &n))
        for (int i = 1; i \le n; i++) scanf ("%d", &a[i]);
        tot = 0;
```

```
tree[0] = build(1, n);
        for (int i = 1; i \le n; i++)
            if(mp. find(a[i]) == mp. end())
                                             tree[i] = update(tree[i-1], i, 1);
            else
                int t = update(tree[i-1], mp[a[i]], -1);
                tree[i] = update(t, i, 1);
            mp[a[i]] = i;
        scanf ("%d", &q);
        while (q--)
            int 1, r;
            scanf("%d%d", &1, &r);
printf("%d\n", query(1, tree[r], 1, n));
    return 0;
                                    树上路径点权第 k 大
int n, q, m, tot, cnt, a[N], b[N], p[N], f[N*2], rmq[N*2];
int tree[N], 1 son[N*40], rson[N*40], c[N*40];
vector<int> v[N];
struct ST
    int mm[2*N], dp[2*N][20];
    void init(int n)
        mm[0] = -1;
        for (int i = 1; i \le n; i++)
            mm[i] = ((i&(i-1)) == 0)?mm[i-1]+1:mm[i-1];
            dp[i][0] = i;
        for (int j = 1; j \le mm[n]; j++)
            for (int i = 1; i+(1 << j)-1 <= n; i++) dp[i][j] = rmq[dp[i][j-1]] <
rmq[dp[i+(1<<(j-1))][j-1]]?dp[i][j-1]:dp[i+(1<<(j-1))][j-1];
    int query (int a, int b)
        if(a > b)
                    swap (a, b);
        int k = mm[b-a+1];
        }st;
void inithash()
    for (int i = 1; i \le n; i++) b[i] = a[i]:
    sort(b+1, b+1+n);
    m = unique(b+1, b+1+n)-b-1;
int gethash(int x)
    return lower bound (b+1, b+1+m, x)-b;
void dfs (int now, int pre, int dep)
```

```
f[++cnt] = now;
    rmq[cnt] = dep;
    p[now] = cnt;
    for (int i = 0; i < v[now]. size (); i++)
        int t = v[now][i];
        if (t == pre) continue;
        dfs(t, now, dep+1);
        f[++cnt] = now;
        rmq[cnt] = dep;
void initlca(int root, int n)
    cnt = 0;
    dfs(root, root, 0);
    st. init (2*n-1);
int querylca(int x, int y)
    return f[st.query(p[x], p[y])];
int build(int 1, int r)
    int now = tot++;
    c[now] = 0;
    if(1 == r) return now;
    int mid = (1+r)/2;
    lson[now] = build(1, mid);
rson[now] = build(mid+1, r);
    return now;
int update(int last, int pos, int x)
    int now = tot++, t = now;
    c[now] = c[1ast] + x;
    int 1 = 1, r = m;
    while (1 < r)
        int mid = (1+r)/2;
        if(pos \le mid)
             lson[now] = tot++;
             rson[now] = rson[last];
             now = 1son[now];
             last = lson[last];
             r = mid;
        else
             rson[now] = tot++;
             lson[now] = lson[last];
             now = rson[now];
             last = rson[last];
             1 = mid+1;
        c[now] = c[last] + x;
    return t;
```

```
void dfsbuild(int now, int pre)
    tree[now] = update(tree[pre], gethash(a[now]), 1);
    for (int i = 0; i < v[now]. size(); i++)
         int t = v[now][i];
        if (t == pre) continue;
        dfsbuild(t, now);
int query (int 1t, int rt, int 1cat, int pos, int k, int 1, int r)
    if(1 == r) return 1;
    int mid = (1+r)/2;
    int t = c[lson[lt]]+c[lson[rt]]-2*c[lson[lcat]];
    if(pos >= 1 && pos <= mid)</pre>
                                   t++;
    if(t >= k) return query(1son[1t], 1son[rt], 1son[1cat], pos, k, 1, mid);
    return query (rson[1t], rson[rt], rson[lcat], pos, k-t, mid+1, r);
int main()
    scanf ("%d%d", &n, &q);
    for (int i = 1; i \le n; i++) scanf ("%d", &a[i]);
    inithash();
    tot = 0;
    for (int i = 1; i < n; i++)
        int x, y;
scanf("%d%d", &x, &y);
        v[x]. push_back(y);
        v[y]. push back(x);
    initlca(1, n);
    tree[0] = build(1, m);
    dfsbuild(1,0);
    while (q--)
         int x, y, k;
        scanf ("%d%d%d", &x, &y, &k);
        int t = querylca(x, y);
        printf(\frac{n}{d}, b[query(tree[x], tree[y], tree[t], gethash(a[t]), k, 1, m)]);
    return 0;
                                           动态第 k 大
int n, q, m, tot, a[N], b[N], use[N];;
int tree[N], tree2[N], 1son[N*50], rson[N*50], c[N*50];
struct Query
    int kind;
    int 1, r, k;
} query[10010];
void Init_hash(int k)
    sort(b, b+k):
    m = unique(b, b+k) - b;
```

```
int hash(int x)
    return lower bound (b, b+m, x)-b;
int build(int 1, int r)
    int root = tot++;
    c[root] = 0;
    if(1 != r)
        int mid = (1+r)/2;
lson[root] = build(1, mid);
rson[root] = build(mid+1, r);
    return root;
int Insert (int root, int pos, int val)
    int newroot = tot++, tmp = newroot;
    int 1 = 0, r = m-1;
    c[newroot] = c[root] + val;
    while (1 < r)
         int mid = (1+r) >> 1;
         if(pos \le mid)
             lson[newroot] = tot++; rson[newroot] = rson[root];
             newroot = 1son[newroot]; root = 1son[root];
             r = mid;
         else
             rson[newroot] = tot++; lson[newroot] = lson[root];
             newroot = rson[newroot]; root = rson[root];
             1 = mid+1;
         c[newroot] = c[root] + val;
    return tmp;
int lowbit(int x)
    return x&(-x);
void add(int x, int pos, int val)
    while (x \le n)
         tree2[x] = Insert(tree2[x], pos, val);
         x += 1 \text{ owbit } (x);
int sum(int x)
    int ret = 0;
    while (x > 0)
        ret += c[1son[use[x]]];
         x = lowbit(x);
    return ret;
```

```
int Query(int left, int right, int k)
    int left root = tree[left-1];
    int right_root = tree[right];
    int 1 = 0, r = m-1;
    for (int i = left-1; i : i -= lowbit(i)) use [i] = tree2[i];
    for(int i = right; i ; i -= lowbit(i)) use[i] = tree2[i];
    while (1 < r)
        int mid = (1+r)/2;
         int tmp = sum(right) - sum(left-1) + c[lson[right root]] - c[lson[left root]];
        if(tmp >= k)
             r = mid;
             for (int i = left-1; i ; i = lowbit(i))
                  use[i] = 1son[use[i]];
             for(int i = right; i; i -= lowbit(i))
    use[i] = lson[use[i]];
             left_root = lson[left_root];
             right_root = lson[right_root];
        else
             1 = mid+1:
             k = tmp;
             for (int i = left-1; i; i = lowbit(i))
                 use[i] = rson[use[i]];
             for(int i = right; i ; i -= lowbit(i))
                 use[i] = rson[use[i]];
             left_root = rson[left_root];
             right_root = rson[right_root];
    return 1;
void Modify(int x, int p, int d)
    while (x \le n)
        tree2[x] = Insert(tree2[x], p, d);
        x += 1 \text{ owbit } (x);
int main()
    int T;
scanf("%d", &T);
    while (T--)
        scanf ("%d%d", &n, &q);
        tot = 0;
        \mathbf{m} = 0:
        for (int i = 1; i \le n; i++)
             scanf("%d", &a[i]);
             b[m++] = a[i];
        char op[10];
        for (int i = 0; i < q; i++)
             scanf("%s", op);
             if(op[0] == 'Q')
```

```
query[i].kind = 0;
               scanf("%d%d%d", &query[i]. 1, &query[i]. r, &query[i]. k);
          else
               query[i].kind = 1;
               scanf("%d%d", &query[i].1, &query[i].r);
b[m++] = query[i].r;
     Init_hash(m);
     tree[0] = build(0, m-1);
     for (int i = 1; i \le n; i++)
          tree[i] = Insert(tree[i-1], hash(a[i]), 1);
     for (int i = 1; i \le n; i++)
          tree2[i] = tree[0];
     for (int i = 0; i < q; i++)
          if(query[i].kind == 0)
               printf("%d\n", b[Query(query[i].1, query[i].r, query[i].k)]);
          else
               Modify(query[i].1, hash(a[query[i].1]),-1);
Modify(query[i].1, hash(query[i].r),1);
a[query[i].1] = query[i].r;
return 0;
```

## 动态规划

```
最长公共子序列 LCS
int dp[1005][1005] = \{0\}, 1en1, 1en2;
char a[1005], b[1005];
void 1cs()
    for (int i = 1; i \le len1; i ++)
        for (int j = 1; j \le 1en2; j++)
             if(a[i-1] == b[j-1]) 	 dp[i][j] = dp[i-1][j-1]+1;
                     dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
void printflcs(int x, int y)
    if(x == 0 | y == 0)
                              return;
    if(a[x-1] == b[y-1])
        printflcs (x-1, y-1);
        printf("%c", a[x-1]);
    else if (dp[x][y-1] > dp[x-1][y])
                                           printflcs (x, y-1);
    else
             printflcs (x-1, y);
                                       最长上升子序列 LIS
int n, a [50005], b [50005] = \{0\};
int main()
    scanf("%d", &n);
for(int i = 1;i <= n;i++)</pre>
                                   scanf ("%d", &a[i]);
    b[1] = a[1];
    int len = 1;
    for (int i = 2; i \le n; i++)
        int t = lower bound(b+1, b+len+1, a[i])-b;
        b[t] = a[i];
        if(t == 1en+1)  1en++;
    printf("%d\n", len);
    return 0;
                                            01 背包
int n, m, w[3500], v[3500], dp[13000] = \{0\};
//数量 n, 背包 m, 重量 w, 价值 v
int main()
    scanf ("%d%d", &n, &m);
    for (int i = 1; i \le n; i++)
                                    scanf ("%d%d", &w[i], &v[i]);
    for (int i = 1; i \le n; i++)
        for (int j = m; j \ge w[i]; j--)
                                           dp[j] = max(dp[j], dp[j-w[i]]+v[i]);
    printf("%d\n", dp[m]);
    return 0;
```

01 背包(容量很大)

int dp[50005], v[1005], w[1005], n, W;

int n, m, dp[105];

int main()

```
//数量 n, 背包 W, 重量 w, 价值 v
//w 很大,考虑每个价值的最小重量
int main()
    ios::sync_with_stdio(0);
    while (cin >> n >> W)
        memset(dp, 0x3f, sizeof(dp));
        dp[0] = 0;
        int sum = 0;
        for (int i = 1; i \le n; i++)
            cin >> w[i] >> v[i];
            sum += v[i];
        for (int i = 1; i \le n; i++)
            for (int j = sum; j \ge v[i]; j--) dp[j] = min(dp[j-v[i]]+w[i], dp[j]);
        for (int i = sum; i >= 0; i--)
            if(dp[i] \leftarrow W)
                cout << i << end1;
                break;
    return 0;
                                         完全背包
//1、如果背包要求正好装满则初始化 dp[0] = 0, dp[1~m] = INF。
//2、如果不需要正好装满 dp[0m] = 0。
int n, m, e, f, w[505], v[505], dp[10005];
//数量 n, 背包 m, 重量 w, 价值 v
int main()
    int T;
    scanf("%d", &T);
    while (T--)
        scanf ("%d%d%d", &e, &f, &n);
        m = f - e;
        memset (dp, 0x3f, sizeof(dp));
        dp[0] = 0;
                                      scanf ("%d%d", &v[i], &w[i]);
        for (int i = 1; i \le n; i++)
        for (int i = 1; i \le n; i++)
            for (int j = w[i]; j \le m; j++)
                                             dp[j] = min(dp[j], dp[j-w[i]]+v[i]);
        if(dp[m] == INF)
                            printf("This is impossible.\n");
               printf ("The minimum amount of money in the piggy-bank is %d.\n", dp[m]);
    return 0;
```

多重背包二进制

```
int T;
    cin >> T;
    while (T--)
         memset (dp, 0, sizeof (dp));
         cin >> n >> m;
         for (int i = 1; i \le m; i++)
              int p, h, c; //花费, 价值, 数量
              cin \gg p \gg h \gg c;
              int now = 1;
              while (1)
                  int t = min(now, c);
                  c = t:
                  int pp = p*t, hh = h*t;
                  for (int j = n; j \ge pp; j--) dp[j] = max(dp[j], dp[j-pp]+hh); if (c == 0) break;
                  now *= 2;
         cout \ll dp[n] \ll end1;
    return 0;
                                          多重背包优先队列
int n, m, dp[105], a[105], b[105];
int main()
    int T;
    cin \gg T:
    while (T--)
         memset(dp, 0, sizeof(dp));
         cin >> n >> m;
         for (int i = 1; i \le m; i++)
              int p, h, c;
              cin \gg p \gg h \gg c;
              if(n/p < c) c = n/p;
              for (int d = 0; d < p; d++)
                  int 1 = 1, r = 0;
                  for (int j = 0; j \le (n-d)/p; j++)
                       while (1 \le r \& b[r] \le dp[j*p+d]-j*h) r--;
                       a[++r] = j;
                       b[r] = dp[j*p+d]-j*h;
                       if(a[1] < j-c) 1++;
dp[j*p+d] = b[1]+j*h;
         cout \ll dp[n] \ll end1;
    return 0;
                                               数位 dp
//求 1-n 中不包含 m 字段的数量
int a[25], b[25], x[25], cnt1, cnt2;
```

long long n, m, dp[25][25];

```
long long dfs (int now1, int now2, int limit)
    if(now2 == 0)
                       return 0;
    if(now1 == 0)
                       return 1;
    if(!limit && dp[now1][now2] != -1) return dp[now1][now2];
    int endd = limit?a[now1]:9;
    long long ans = 0;
    for (int i = endd; i >= 0; i--)
         int t = now2;
         while (b[t] != i \&\& t != cnt2) t = x[t];
         if(b[t] == i) t--;
         ans += dfs(now1-1, t, limit && i == endd);
    if(!limit) dp[now1][now2] = ans;
    return ans;
int main()
    ios::sync_with_stdio(0);
    int T;
    cin >> T;
    while (T--)
         cin \gg n \gg m;
         cnt1 = 0;
         cnt2 = 0;
         while (n)
              a[++cnt1] = n\%10;
              n /= 10;
         while (m)
              b[++cnt2] = m\%10;
              m /= 10;
         int i = cnt2, j = cnt2+1;
x[cnt2] = cnt2+1;
         while (i >= 1)
              if(j = cnt2+1 \mid b[i] = b[j])  x[--i] = --j;
              else j = x[j];
         x[cnt2] = cnt2;
         memset (dp, -1, sizeof(dp));
         \operatorname{cout} << \operatorname{dfs}(\operatorname{cnt1}, \operatorname{cnt2}, 1) - 1 << \operatorname{end1};
    return 0;
                                               状压 dp
//裁玻璃 2*2
```

```
int a[15][15], sta[100], dp[15][100], sum[15][100], n, m;
int main()
{
    ios::sync_with_stdio(false);
    int T;
    cin >> T;
    while(T--)
    {
        memset(dp, 0, sizeof(dp));
        memset(sum, 0, sizeof(sum));
    }
}
```

```
cin >> n >> m;
         for (int i = 1; i \le n; i++)
             for (int j = 1; j \le m; j++)
                                             cin \gg a[i][j];
         int cnt = 0, endd = 1 << (m-1);
         for (int i = 0; i < endd; i++)
             if (i & (i << 1)) continue;
             sta[++cnt] = i;
         for (int i = 1; i < n; i++)
             for (int j = 1; j \le cnt; j++)
                  for (int k = 1; k \le m-1; k++)
                       if (1<<(k-1) & sta[j])
                           if(a[i][k] && a[i][k+1] && a[i+1][k] && a[i+1][k+1])
sum[i][j]++;
         for (int i = 1; i \le cnt; i++) dp[1][i] = sum[1][i];
         for (int i = 1; i < n; i++)
             for (int j = 1; j \leftarrow cnt; j++)
                  for (int k = 1; k \le cnt; k++)
                       dp[i+2][k] = max(dp[i+2][k], dp[i][j]+sum[i+2][k]);
                      if(sta[j] & sta[k])
                                                       continue;
                      if(sta[j] & (sta[k]<<1))
if(sta[j] & (sta[k]>>1))
                                                       continue:
                                                      continue:
                      dp[i+1][k] = max(dp[i+1][k], dp[i][j]+sum[i+1][k]);
         int ans = 0;
         for (int i = 1; i \le cnt; i++) ans = max(ans, dp[n-1][i]);
         cout << ans << end1;</pre>
    return 0;
```

## 数论

```
int cnt, prime [MAXN+1], mi [MAXN+1], vis [MAXN+1] = \{0\};
//cnt 表示素数个数
//prime 存放每个素数
//mi 存放每个数的最小素数因子
void getprime()
    cnt = 0;
    memset(prime, 0, sizeof(prime));
    for (int i = 2; i \le MAXN; i++)
        if(!vis[i])
            prime[++cnt] = i;
            mi[i] = i;
        for(int j = 1; j \le cnt \&\& (long long) i*prime[j] \le MAXN; j++)
            vis[prime[j]*i] = 1;
            mi[prime[j]*i] = prime[j];
            if(!(i%prime[j])) break;
int cnt2, notprime2[1000005], prime2[1000005];
//区间 1, r 之间的素数
void getprime2(int 1, int r)
    cnt2 = 0;
    memset (notprime2, 0, sizeof (notprime2));
    if(1 < 2)   1 = 2;
    for(int i = 1; i \le cnt \&\& (long long)prime[i]*prime[i] <= r; i++)
        int t = 1/prime[i];
        if(1\%prime[i]) t++;
        if(t == 1) t = 2;
        for(int j = t; (long long) j*prime[i] <= r; j++)</pre>
            if((long\ long)\ j*prime[i] >= 1) notprime2[j*prime[i]-1] = 1;
    for (int i = 0; i \le r-1; i++)
        if(!notprime2[i]) prime2[++cnt2] = i+1;
                                        合数分解
long long factor[105][2];
//返回素因子个数
//factor[i][0]表示第 i 个素因子, factor[i][1]表示第 i 个素因子的个数
int getfac (long long x)
```

for(int i = 1; (long long)prime[i]\*prime[i] <= x;i++)</pre>

factor[++cnt][0] = prime[i];

while (x%prime[i] == 0)

if(x%prime[i] == 0)

```
factor[cnt][1]++;
                 x \neq prime[i];
    if(x != 1)
        factor[++cnt][0] = x;
        factor[cnt][1] = 1;
    return cnt;
                                            快速乘
long long qmul (long long a, long long b, long long c)
    long long ans = 0;
    a = a\%c;
    b = b\%c;
    while (b)
        if (b%2)
             ans += a;
             if (ans > c) ans -= c;
        a = a+a;
        if(a > c)
                     a = c;
        b /= 2;
    return ans;
                                            快速幂
long long qpower (long long a, long long b, long long c)
    long long ans = 1;
    a = a\%c;
    while (b)
        if (b%2)
                    ans = ans*a\%c;
        a = a*a%c;
        b /= 2;
    return ans;
                                         string 快速幂
string mul(string a, string b)
    int arr[200], len = a. length() + b. length();
    memset (arr, 0, sizeof (arr));
    reverse (a. begin (), a. end ());
    reverse (b. begin (), b. end ());
    for (int i = 0; i < a. length(); i++)
        for(int j = 0; j < b.length(); j++) arr[i+j] += (a[i]-'0')*(b[j]-'0');
    for (int i = 0; i < len; i++)
        arr[i+1] += arr[i]/10;
        arr[i] %= 10;
    string ans = string(len, '0');
    for (int i = 0; i < len; i++) ans [i] += arr[i];
```

```
reverse (ans. begin (), ans. end ());
    cout << ans << end1;
    return ans;
string strpow(string x, int b)
    string ans = "1";
    while(b)
        if(b\%2) ans = mul(ans, x);
        x = mu1(x, x);
        b /= 2;
    return ans;
                                         矩阵快速幂
struct matrix
    long long m[2][2];
matrix one = \{ 1, 0, \}
                 0, 1 };
matrix mul (matrix a, matrix b)
    matrix tmp;
    for (int i = 0; i < 2; i++)
        for (int j = 0; j < 2; j++)
            tmp. m[i][j] = 0;
            for (int_k = 0; k \le 2; k++)
                                          tmp.m[i][j] =
(tmp. m[i][j]+a. m[i][k]*b. m[k][j])%MOD;
    return tmp;
matrix magower (matrix a, int b)
    matrix ans = one;
    while (b)
        if(b\%2) ans = mul(ans, a);
        a = mul(a, a);
        b /= 2;
    return ans;
                                     miller_rabin 判断素数
//用 a 来检验 n 是否为素数
//是素数返回1,不是返回0
bool check(long long a, long long n, long long x, long long t)
    long long ans = qpower (a, x, n);
    if (ans == 1 \mid | ans == n-1) return 1;
    while (t--)
        ans = qmul(ans, ans, n);
        if (ans == n-1) return 1;
```

```
return 0;
//判断 n 是否为素数
//是素数返回1,不是返回0
bool miller rabin (long long n)
    if (n < 2) return 0;
    if (n == 2) return 1;
    if(n\%2 == 0) return 0;
    long long x = n-1;
    long long t = 0;
    while (x\%2 == 0)
        x /= 2;
        t++;
    srand(time(0));
    for (int i = 1; i \le 10; i++)
        long long a = rand()\%(n-1)+1;
        if(!check(a, n, x, t)) return 0;
    return 1;
                                     pollard_rho 分解质因数
long long factor[105];
int cnt;
long long pollard rho(long long n, long long c)
    long long i = 1, k = 2;
    srand(time(0));
    long long x = rand()\%(n-2)+1, y = x;
    while(1)
        i++;
        x = (qmu1(x, x, n) + c) %n;
        long long d = __gcd(y-x, n);
if(1 < d && d < n) return d;</pre>
        if(x == y) return n;
        if(i == k)
            y = x;
            k = 2;
// 求 n 的质因数, k 设为 107
void findfac(long long n, int k)
    if(n == 1) return;
    if (miller_rabin(n))
        factor[++cnt] = n;
        return;
    long long p = n;
    int c = k;
    while (p \ge n)  p = pollard_rho(p, c--);
    findfac(p,k);
    findfac (n/p, k);
```

```
gcd、lcm
LL gcd(LL x, LL y) //最大公约数
    return y?gcd(y, x%y):x;
                         //最小公倍数
LL 1cm(LL x, LL y)
    return x*y/gcd(x, y);
                                       lcm(1,2,3,...n)
//1cm(C(n, 0), C(n, 1), ..., C(n, n)) = 1cm(1, 2, 3, ...n+1)/(n+1)
//计算 1cm(1, 2, 3, ···n)
const int N=100000007;
int visit [N/32+50];
unsigned int data[5800000];
int prime [5800000], np=0;
             //筛素数,数组从 0 开始
void Prime()
    prime[0] = data[0] = 2;
    np = 1;
    for (int i = 3; i < N; i += 2)
        if(!(visit[i/32]&(1<<((i/2)%16))))</pre>
            prime[np] = i;
            data[np] = data[np-1]*i;
            for (int j = 3*i; j < N; j += 2*i) visit[j/32] = (1 << ((j/2)\%16));
long long Deal(int n)
    int p = upper bound(prime, prime+np, n)-prime-1;
    long long ans = data[p];
    for (int i = 0; i < np && prime[i]*prime[i] <= n; i++)
        int mul = prime[i], tmp = prime[i] * prime[i]; ;
        while (tmp/mul == prime[i] && tmp <= n)</pre>
            tmp *= prime[i];
            mul *= prime[i];
        ans = ans*(mul/prime[i])%MOD;
    return ans;
                                         欧拉函数
小于 N 且与 N 互质的数的个数(包括 1)。
对于一个正整数 N 的素数幂分解 N = P1^q1*P2^q2*...*Pn^qn。
\phi(N)=N*(1-1/P1)*(1-1/P2)*...*(1-1/Pn)
//先分解质因数
int euler(int n)
```

for (int i = 1;  $i \le cnt$ ; i++) ans = ans/factor[i][0]\*(factor[i][0]-1);

```
long long euler (long long n)
    long long ans = 1;
    for (int i = 2; (long long) i*i \le n; i++)
        if(n\%i == 0)
             n /= i;
             ans *= i-1;
             while (n\%i == 0)
                 n /= i;
                 ans *= i;
    if(n > 1)
                   ans *= n-1;
    return ans;
```

线性筛

```
int cnt, phi[N+10], prime[N+10], vis[N+10];
//同时得到欧拉函数和素数表
void geteuler()
    memset (vis, 0, sizeof (vis));
    phi[1] = 1;
    cnt = 0;
    for (int i = 2; i \le N; i++)
        if(!vis[i])
            prime[++cnt] = i;
            phi[i] = i-1;
        for (int j = 1; j \leftarrow cnt; j++)
            if((long long)i*prime[j] > N)
                                              break;
            vis[i*prime[j]] = 1;
            if(i\%prime[j] == 0)
                 phi[i*prime[j]] = phi[i]*prime[j];
                 break:
                     phi[i*prime[j]] = phi[i]*(prime[j]-1);
            e1se
```

### lucas 定理

```
A、B 是非负整数, p 是质数。AB 写成 p 进制: A=a[n]a[n-1]...a[0], B=b[n]b[n-1]...b[0]。
则组合数 C(A,B)与 C(a[n],b[n])*C(a[n-1],b[n-1])*...*C(a[0],b[0]) modp 同余。
即: Lucas(n,m,p)=c(n%p,m%p)*Lucas(n/p,m/p,p)。
//用与 n, m 很大, MOD 不是很大的情况
long long qmod(long long a, long long b, long long c)
    long long ans = 1;
    a = a\%c;
    while (b)
       if(b\%2) ans = (ans*a)\%c;
```

```
a = (a*a)%c;
        b /= 2;
    return ans;
long long c (long long m, long long n)
    if(m < n)
                 return 0;
    if(m == n) return 1;
    if(n > m-n) n = m-n;
    long long mm = 1, nn = 1;
    for (long long i = 0; i < n; i++)
        mm = mm*(m-i)%MOD;
        nn = nn*(n-i)%MOD;
    return mm*qmod(nn, MOD-2, MOD) %MOD;
long long lucas (long long m, long long n)
    long long ans = 1;
    while (m && n && ans)
        ans = ansMOD*c (mMOD, nMOD) MOD;
        n /= MOD;
        m /= MOD;
    return ans;
                                       扩展 lucas(模合数)
long long inv (long long a, long long MOD)
    if(!a) return 0;
    long long b = MOD, x = 0, y = 0;
    e_{gcd}(a, b, x, y);
    x = ((x\%b)+b)\%b;
    if(!x) x += b;
    return x;
long long mul (long long n, long long pi, long long pk)
    if(!n) return 1;
    long long ans=1;
    for (long long i = 2; i \leq pk; i++)
        if(i%pi)
                     ans = ans*i\%pk;
    ans = qpower(ans, n/pk, pk);
    for (long long i = 2; i \le n\%pk; i++)
        if(i%pi)
                     ans = ans*i\%pk;
    return ans*mul(n/pi, pi, pk)%pk;
long long C(long long n, long long m, long long MOD, long long pi, long long pk)
    if (m > n) return 0:
    long long a = mul(n, pi, pk), b = mul(m, pi, pk), c = mul(n-m, pi, pk), k = 0, ans = 0;
    for (long long i = n; i; i /= pi) k += i/pi;
```

```
for (long long i = m; i; i /= pi) k -= i/pi;
    for (long long i = n-m; i; i \neq pi) k = i/pi;
    ans = a*inv(b, pk)%pk*inv(c, pk)%pk*qpower(pi, k, pk)%pk;
    return ans*(MOD/pk)%MOD*inv(MOD/pk, pk)%MOD;
long long lucas (long long n, long long m, long long MOD)
    long long ans = 0, x = MOD;
    for (long long i = 2; i*i \le x; i++)
        if(x\%i == 0)
            long long t = 1;
            while (x\%i == 0)
                t *= i;
                x /= i:
            ans = (ans+C(n, m, MOD, i, t)) %MOD;
                ans = (ans+C(n, m, MOD, x, x)) %MOD;
    if(x > 1)
    return ans;
                                           逆元
费马小定理: 假如 p 是素数,且 a 与 p 互质,那么 a^{(p-1)} = 1 (mod p)。
费马大定理: 当整数 n>2 时,关于 x,y,z 的方程 x^n+y^n=z^n 没有正整数解。
欧拉定理: 若 n,a 为正整数,且 n,a 互质,则:a^(phi(n)) = 1(mod n),可得 a^b%p = a^(b%phi(p))%p
另外,a,b 不互质且 b < phi(n)时,有 a^b = a^(b \cdot phi(p) + phi(p))%p
逆元: ax = 1 \pmod{m}的 x 值(0 < a < m)。
当 p 是质数的时候 a/x \mod p == a*x^(p-2) \mod p。
当 p 不是质数的时候 a/x \mod p == a*x^(phi(p)-1) \mod p。
                                          逆元 1
long long inv (long long a, long long m)
    return qmod(a, m-2, m);
                                          逆元 2
long long inv (long long a, long long m)
    if(a == 1) return 1;
    return inv (m\%a, m)*(m-m/a)\%m;
                                          逆元 3
//扩展欧几里德
long long e gcd(long long a, long long b, long long &x, long long &y)
    if(!b)
        X = 1;
        y = 0;
        return a;
    long long d = e_{gcd}(b, a\%b, y, x);
```

```
y = a/b*x;
    return d;
long long inv (long long a, long long m)
    long long x, y;
    long long gcd = e_gcd(a, m, x, y);
    if(gcd == 1)  return (x\%m+m)\%m;
    return -1;
                                       中国剩余定理
//求解方程组 x%m = a, x 的最小值
//无解返回-1
long long e_gcd(long long a, long long b, long long &x, long long &y)
    if(!b)
        x = 1;
        y = 0;
        return a;
    long long d = e \gcd(b, a\%b, y, x);
    y = a/b*x;
    return d;
long long solve (long long *m, long long *a, long long n)
    long long M = m[1], A = a[1], x, y;
    for (int i = 2; i \le n; i++)
        long long d = e_{gcd}(M, m[i], x, y);
        if ((a[i]-A)\%d) return -1;
        x = (a[i]-A)/d*x\%(m[i]/d);
        A += x*M;
        M = M/d*m[i];
        A \% = M;
    if(A < 0)
                A += M;
    return A;
                                      浮点数高斯消元
//浮点型只有唯一解时可计算
//返回 0 表示无解
#define eps 1e-9
double a[N][N], x[N];
                            //左边矩阵和右边值,结果存在 x 数组内
int Gauss (int equ, int var) //方程个数和未知数个数
    for (int row = 0, col = 0; col \langle var&&row \langle equ; col++, row++ \rangle
        int max r=row;
        for (int i = row+1; i < equ; i++)
            if(fabs(a[i][col])-fabs(a[max r][col]) > eps)
                                                              max r=i;
        if(fabs(a[max r][col]) < eps)
                                        return 0;
        if(max r != row)
            for (int j = 0; j \le var; j++) swap (a[row][j], a[max r][j]);
        for (int i = row+1; i < equ; i++)
```

01 异或高斯消元

```
//有 equ 个方程, var 个变元。增广矩阵行数为 equ, 分别为 0 到 equ-1, 列数为 var+1, 分别为 0 到
var.
int equ, var;
int a[250][250];
                    //增广矩阵
int x[250]:
                    //解集
int free_x[250];
                    //标记是否不确定的变元
int free num;
                    //不确定变元个数
int n;
char s[30];
//-1 表示无解, 0 表示唯一解, 大于 0 表示无穷解, 并返回自由变元的个数)
void init()
    memset(a, 0, sizeof(a));
    memset(x, 0, sizeof(x));
    equ = n*n;
    var = n*n;
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
            int t = i*n+j;
            a[t][t] = 1;
            if (i > 0)  a[(i-1)*n+j][t] = 1;
            if (i < n-1) a[(i+1)*n+j][t] = 1;
            if(j > 0) a[i*n+j-1][t] = 1;
            if(j < n-1) \ a[i*n+j+1][t] = 1;
void Debug()
    for (int i = 0; i < equ; i++)
        for(int j = 0; j < var+1; j++) cout \langle \langle a[i][j] \langle \langle "";
        cout << end1;
    cout << end1;
int Gauss()
    int max_r, col, k;
    free_num = 0;
    for (k = 0, col = 0; k < equ && col < var; k++, col++)
        \max r = k;
        for(int i = k+1; i < equ; i++)
```

```
if(abs(a[i][col]) > abs(a[max r][col])) max r = i;
        if(max r != k)
             for (int i = col; i < var+1; i++) swap (a[k][i], a[max r][i]);
        if(a[k][col] == 0)
             free_x[free_num++] = co1;
             continue;
        for (int i = k+1; i < equ; i++)
             if(a[i][col] == 0) continue;
             for(int j = col; j < var+1; j++) a[i][j] = a[k][j];
    for (int i = k; i < equ; i++)
        if(a[i][col] != 0) return -1;
    return var-k;
void solve()
    int t = Gauss();
    if (t == -1) printf ("inf \setminus n");
    else
        int ans = INT_MAX, tot = (1 << t);
        for (int i = 0; i < tot; i++)
             int cnt = 0;
             for (int j = 0; j < t; j++)
                 if (i& (1<< j))
                     x[free_x[j]] = 1;
                     cnt++;
                 else x[free x[j]] = 0;
             for (int j = var-t-1; j >= 0; j--)
                 int t = a[j][var];
                 for (int k = j+1; k < var; k++)
                     if(a[j][k]) t = x[k];
                 x[j] = t;
cnt += x[j];
            ans = min(ans, cnt);
        printf("%d\n", ans);
                                      同余方程组高斯消元
```

int a[MAXN][MAXN]; //增广矩阵 int x[MAXN]; //解集

```
bool free x[MAXN]; //标记是否是不确定的变元
int gcd(int a, int b)
    return b?gcd(b, a%b):a;
int lcm(int a, int b)
    return a/gcd(a, b)*b;
//-1 表示无解,0 表示唯一解,大于 0 表示无穷解,并返回自由变元的个数)
//有 equ 个方程, var 个变元。增广矩阵行数为 equ, 分别为 0 到 equ-1, 列数为 var+1, 分别为 0 到
int Gauss(int equ, int var)
    int i, j, k;
    int max_r;// 当前这列绝对值最大的行.
int col;//当前处理的列
    int ta, tb;
    int LCM;
    int temp;
    int free x num;
    int free_index;
for(int i = 0;i <= var;i++)</pre>
        x[i]=0;
        free_x[i]=true;
    col=0;
    for (k = 0; k < equ \&\& col < var; k++, col++)
        max_r=k;
        for(i = k+1; i < equ; i++)
             if(abs(a[i][col]) > abs(a[max_r][col])) max r=i;
        if(max_r != k)
             for (j = k; j < var+1; j++) swap (a[k][j], a[max_r][j]);
        if(a[k][co1] == 0)
             k--:
             continue;
        for (i = k+1; i < equ; i++)
             if (a[i][col] != 0)
                 LCM = 1cm(abs(a[i][col]), abs(a[k][col]));
                 ta = LCM/abs(a[i][col]);
tb = LCM/abs(a[k][col]);
if(a[i][col]*a[k][col] < 0) tb = -tb;</pre>
                 for (j = col; j < var+1; j++)
                     a[i][j] = ((a[i][j]*ta-a[k][j]*tb)%7+7)%7;
    for (i = k; i < equ; i++)
        if (a[i][col] != 0) return -1;
```

```
if(k < var)
    for (i = k-1; i >= 0; i--)
        free x num = 0;
        for (j = 0; j < var; j++)
             if (a[i][j] != 0 \&\& free_x[j]) free_x_num++, free_index = j;
        if(free_x_num > 1) continue;
        temp = a[i][var];
        for (j = 0; j < var; j++)
             if(a[i][j] != 0 \&\& j != free index) temp -= a[i][j]*x[j]%7;
             temp = (temp\%7+7)\%7;
        x[free\_index] = (temp/a[i][free\_index])%7;
        free_x[free_index] = 0;
    return var-k;
for (i = var-1; i \ge 0; i--)
    temp = a[i][var];
    for (j = i + 1; j < var; j++)
        if(a[i][j] != 0) temp -= a[i][j]*x[j];
        temp = (temp\%7+7)\%7;
    while (\text{temp}\%a[i][i] != 0)
                                  temp += 7;
    x[i] = (temp/a[i][i])\%7;
return 0;
```

### FFT 递归

typedef complex<double> C;

```
int n, m;
C a[N], b[N];
int gi()
    int res = 0, fh = 1;
    char ch = getchar();
    while((ch > '9' | | ch < '0') && ch != '-') ch = getchar();
if(ch == '-')</pre>
         fh=-1;
         ch=getchar();
    while (ch >= '0' && ch <= '9')
         res = res*10+ch-'0';
         ch=getchar();
    return fh*res;
void fft(C *a, int n, int f)
    if(n == 1) return;
    C \text{ wn}(\cos(2.0*PI/n), \sin(f*2.0*PI/n)), w(1,0), t, a0[n>>1], a1[n>>1];
    for (int i = 0; i < n >> 1; i++)
```

```
a0[i] = a[i << 1];
         a1[i] = a[i << 1]1;
    fft(a0, n >> 1, f);
    fft(a1, n >> 1, f);
    for (int i = 0; i < n >> 1; i++, w *= wn)
         t = w*a1[i];
        a[i] = a0[i]+t;
        a[i+(n)>1)] = a0[i]-t;
int main()
    n = gi();
    m = gi();
    for(int i = 0; i \le n; i++)
                                    a[i]=gi();
    for (int i = 0; i \le m; i++)
                                    b[i]=gi();
    m += n;
    for (n = 1; n \le m; n *= 2);
    fft(a, n, 1);
    fft(b, n, 1);
                                    a[i] *= b[i];
    for (int i = 0; i \le n; i++)
    fft(a, n, -1);
                                    printf("%d", int(a[i]. real()/n+0.5));
    for (int i = 0; i \le m; i++)
    return 0;
```

### FFT 非递归

```
typedef complex<double> C;
int n, m, L, R[N];
C a[N], b[N];
int gi()
    int res = 0, fh = 1;
    char ch = getchar();
    while ((ch > '9' \mid | ch < '0') \&\& ch != '-') ch = getchar();
    if (ch == '-')
        fh=-1;
        ch=getchar();
    while (ch >= '0' && ch <= '9')
        res = res*10+ch-'0';
        ch=getchar();
    return fh*res;
void fft(C *a, int f)
    for (int i = 0; i < n; i++)
        if(i < R[i])
                       swap(a[i], a[R[i]]);
    for (int i = 1; i < n; i *= 2)
        C \text{ wn}(\cos(PI/i), \sin(f*PI/i)), x, y;
        for (int j = 0; j < n; j += i << 1)
```

```
C w(1, 0);
             for (int k = 0; k < i; k++, w *= wn)
                 x = a[j+k];
                 y = w*a[j+i+k];
                  a[j+k] = x+y;
                  a[j+i+k] = x-y;
        }
int main()
    n = gi();
    m = gi();
                                   a[i] = gi();
b[i] = gi();
    for (int i = 0; i \le n; i++)
    for (int i = 0; i \le m; i++)
    m += n;
    for (n = 1; n \le m; n *= 2)
                                   L++:
                                    R[i] = (R[i>>1]>>1) | ((i&1)<<(L-1));
    for (int i = 0; i < n; i++)
    fft(a, 1);
    fft(b, 1);
    for (int i = 0; i \le n; i++)
                                    a[i] *= b[i];
    fft(a,-1);
                                    printf("%d", int(a[i]. real()/n+0.5));
    for (int i = 0; i \le m; i++)
    return 0;
                                            约数和定理
//需要 getprime, qpower, getfac, 先调用 getprime
//计算 1+p+p^2+^
                  ^{\text{h}}+p^{\text{h}}n
long long sum(long long p, long long n)
    if(p == 0) return 0;
    if(n == 0) return 1;
    if (n\%2) return (1+qpower(p, n/2+1, MOD))\%MOD*sum(p, n/2)\%MOD;
           return ((1+qpower(p, n/2+1, MOD)) MOD*sum(p, n/2-1)
1) +qpower (p, n/2, MOD) \% MOD) \% MOD;
//求 a b 的约数和对 MOD 取模
long long solve (long long a, long long b)
```

### 莫比乌斯反演

$$F(n) = \sum_{d|n} f(d) => f(n) = \sum_{d|n} \mu(d) F(\frac{n}{d})$$

int cnt = getfac(a);
long long ans = 1;

ans %= MOD;

return ans;

for (int i = 1;  $i \le cnt$ ; i++)

$$F(n) = \sum_{n|d} f(d) => f(n) = \sum_{n|d} \mu(\frac{d}{n}) F(d)$$

```
const int MAXN = 100000;
bool check[MAXN+5];
```

ans \*= sum(factor[i][0], b\*factor[i][1])%MOD;

```
int prime[MAXN+5], mu[MAXN+5], sum[MAXN+5];
//求莫比乌斯函数
void Moblus()
    memset (check, 0, sizeof (check));
    mu[1] = 1;
    int cnt = 0;
    for (int i = 2; i \le MAXN; i++)
         if(!check[i])
             prime[cnt++] = i;
             mu[i] = -1;
         for (int j = 0; j < cnt && i*prime[j] <= MAXN; <math>j++)
             check[i*prime[j]] = 1;
             if(i\%prime[j] == 0)
                  mu[i*prime[j]] = 0;
                  break;
                      mu[i*prime[j]] = -mu[i];
             else
    sum[0] = 0;
                                        sum[i] = sum[i-1]+mu[i];
    for (int i = 1; i \le MAXN; i++)
//求[1, n], [1, m]内互质的数的对数
int solve(int n, int m)
    int ans = 0;
    if(n > m) swap (n, m);
    for (int i = 1, last; i \le n; i = last+1)
         last = \min(n/(n/i), m/(m/i));
         ans += (sum[1ast]-sum[i-1])*(n/i)*(m/i);
    return ans;
int main()
    Moblus();
    int a, b, c, d, k, T;
scanf("%d", &T);
    while (T--)
         scanf ("%d%d%d%d%d", &a, &b, &c, &d, &k);
         c-
         a \neq k;
         b /= k;
         c /= k;
         d /= k;
         printf ("%d \setminus n", solve (b, d) -solve (a, d) -solve (b, c) +solve (a, c));
```

```
int Inval(int a, int b, int n)
    long long x, y, e;
    e \gcd(a, n, x, y);
    e = (long long) x*b%n;
    return e < 0?e+n:e;
//求 A^x = B(mod C)(C 可以为非素数)
int BSGS (int A, int B, int C)
    map <int, int > H;
    long long buf = 1\%C, D = buf, K;
    int d = 0, tmp;
    for (int i = 0; i \le 100; buf = buf*A%C, i++)
    if (buf == B) return i;
    while ((tmp = gcd(A, C)) != 1)
        if (B % gcd(A, C) != 0) return -1;
        d++;
        C /= tmp;
        B /= tmp;
        D = D*A/tmp%C;
    H. clear();
    int M = (int)ceil(sqrt(C));
    buf = 1\%C;
    for (int i = 0; i \le M; buf = buf*A%C, i++)
        if(H. find((int)buf) == H. end()) H[(int)buf] = i;
    K = \text{qpower}(A, M, C);
    for (int i = 0; i \le M; D = D*K%C, i++)
        tmp = Inval((int)D, B, C);
        if(tmp \ge 0 \&\& H.find(tmp) != H.end()) return i*M+H[tmp]+d;
    return -1;
                                     自适应 simpson 积分
const double eps = 1e-6;
// 三点 simpson 法。这里要求 F 是一个全局函数
double simpson (double a, double b)
    double c = a+(b-a)/2;
    return (F(a) + 4*F(c) + F(b))*(b-a)/6;
// 自适应 Simpson 公式(递归过程)。已知整个区间[a,b]上的三点 simpson 值 A
double asr (double a, double b, double eps, double A)
    double c = a+(b-a)/2;
    double L = simpson(a, c), R = simpson(c, b);
    if (fabs(A-L-R) \le 15*eps) return L+R+(A-L-R)/15;
    return asr (a, c, eps/2, L) + asr(c, b, eps/2, R);
// 自适应 Simpson 公式 (主过程)
double asr (double a, double b, double eps)
    return asr(a, b, eps, simpson(a, b));
```

## 康托展开

```
X = an*(n-1)!+an-1*(n-2)!+...+ai*(i-1)!+...+a2*1!+a1*0!

其中, ai 为当前未出现的元素中是排在第几个(从 0 开始)。

int f(vector<int> v)
{
    int sum = 0;
    for(int i = 0; i < 8; i++)
    {
        int t = 0;
        for(int j = i+1; j < 8; j++)
        {
            if(v[i] > v[j]) t++;
        }
        sum += t*h[7-i];
    }
    return sum;
}
```

sg 函数

```
//f[]: 可以取走的石子个数
//sg[]:0 n 的 SG 函数值
//hash[]:mex{}
int n, m, p, f[N], sg[N], hashh[N];

void getSG(int n)
{
    memset(sg, 0, sizeof(sg));
    for(int i = 1; i <= n; i++)
    {
        memset(hashh, 0, sizeof(hashh));
        for(int j = 1; f[j] <= i; j++)
        for(int j = 0; j <= n; j++)
        {
            if (hashh[j] == 0)
            {
                  sg[i] = j;
                 break;
            }
        }
     }
}
```

### 整数划分分类

1. 若划分的多个整数可以相同

设 dp[i][j]为将 i 划分为不大于 j 的划分数。

- (1) 当 i < j 时, i 不能划分为大于 i 的数, dp[i][j] = dp[i][i]。
- (2) 当 i>j 时,可以根据划分中是否含有 j 分为两种情况。

若划分中含有j,划分方案数为dp[i-j][j]。

若划分数中不含 j,相当于将 i 划分为不大于 j-1 的划分数,为 dp[i][j-1]。

所以当 i>j 时, dp[i][j]=dp[i-j][j]+dp[i][j-1]。

- (3) 当 i=j 时,若划分中含有 j 只有一种情况,若划分中不含 j 相当于将 i 划分为不大于 j-1 的划分数。此时 dp[i][j]=1+dp[i][j-1]。
- 2.将 n 划分为 k 个整数的划分数

dp[i][j]表示将 i 划分成 j 个正整数的划分数。

- (1) 当 i < j 时, i 不能划分成 j 个正整数, dp[i][j]=0。
- (2) 当 i=j 时, 只有一种情况, dp[i][j]=1。
- (3) 当 i>j 时, 此时根据含 1 和不含 1, 故 dp[i][j]=dp[i-1][j-1]+dp[i-j][j]。
- 3. 若划分的正整数必须不同

设 dp[i][j]为将 i 划分为不超过 j 的不同整数的划分数。

- (1) 当 i < j 时, i 不能划分为大于 i 的数, 所以 dp[i][j]=dp[i][i]。
- (2) 当 i>j 时,可以根据划分中是否含有 j 分为两种情况。

若划分中含有 j,则其余的划分中最大只能是 j-1,方案数为 dp[i-j][j-1]。

若划分中不含 j,相当于将 i 划分为不大于 j-1 的划分数,为 dp[i][j-1]。

所以当 i>j 时 dp[i][j]=dp[i-j][j-1]+dp[i][j-1]。

- (3) 当 i=j 时,若划分中含有 j 只有一种情况,若划分中不含 j 相当于将 i 划分为不大于 j-1 的划分数。此时 dp[i][j]=1+dp[i][j-1]。
- 4.将 n 划分为若干正奇数之和的划分数

设 f[i][j]为将 i 划分为 j 个奇数之和的划分数,g[i][j]为将 i 划分为 j 个偶数之和的划分数。

使用截边法,将g[i][j]的j个划分都去掉1,可得g[i][j] = f[i-j][j]。

f[i][j]中有包含 1 的划分方案和不包含 1 的划分方案。

对于包含 1 的划分方案,可以将 1 的划分除去,转化为"将 i-1 划分为 j-1 个奇数之和的划分数",即 f[i-1][j-1]。 对于不包含 1 的划分方案,可以使用截边法对 j 个划分每一个都去掉一个 1,转化为"将 i-j 划分为 j 个偶数之和的划分数",即 g[i-j][j]。

所以 f[i][j]=f[i-1][j-1]+g[i-j][j]。

f[n][0]+f[n][1]+.....+f[n][n]为将 n 划分为若干奇数的划分数。

#### Polya 定理

设  $G = \{a1, a2, ..., ag\}$ 是 N 个对象的置换群,用 M 种颜色给这 N 个对象着色,

则不同的着色 方案数为:  $|G|^{(-1)} * \{M^c(a1) + M^c(a2) + ... + M^c(ag)\}$ 。其中 c(ai)为置换 ai 的循环节数,( i=1,2,...,g )。

### 原根

设 m 是正整数, a 是整数, 若 a 模 m 的阶等于 φ(m), 则称 a 为模 m 的一个原根。

如果  $g \in P$  的原根,那么 g 的(1...P-1)次幂 mod P 的结果一定互不相同。

如果  $g \in P$  的原根,就是  $g^{(P-1)} = 1 \pmod{P}$  当且仅当指数为 P-1 的时候成立.(这里  $P \in P$  是素数).

```
int n, a[35000], cnt = 0;
int main()
{
    scanf("%d", &n);
    int endd = sqrt(n-1);
    for(int i = 2; i <= endd; i++)
    {
        if((n-1)%i == 0)      a[++cnt] = i;
    }
    for(int i = 2; i <= n-1; i++)
    {
        int j;
    }
}</pre>
```

### 卡特兰数

#### 1, 2, 5, 14, 42

```
h(0) = 1, h(1) = 1
```

h(n) = h(0)\*h(n-1)+h(1)\*h(n-2) + ... + h(n-1)h(0) (n>=2)

h(n) = h(n-1)\*(4\*n-2)/(n+1)

h(n) = C(2n,n)/(n+1)

h(n) = C(2n,n)-C(2n,n-1)

- 1.n 对括号正确匹配组成的字符串数。
- 2.一个栈(无穷大)的进栈序列为 1, 2, 3, ..., n, 有多少个不同的出栈序列。
- 3.在一个凸多边形中,通过若干条互不相交的对角线,把这个多边形划分成了若干个三角形,求不同划分的方案数 f (n)。
- 4.一位大城市的律师在她住所以北 n 个街区和以东 n 个街区处工作。
- 5.长度为 2n 的 Dyck words 的数量。 Dyck words 是由 n 个 X 和 n 个 Y 组成的字符串,并且从左往右数, Y 的数量不超过 X。
- 6.拥有 n+1 个叶子节点的二叉树的数量。例如 4个叶子节点的所有二叉树形态。
- 7.圆桌握手问题: 圆桌周围有 2n 个人,他们两两握手,但没有交叉的方案数。

## 调和级数

f(n)≈ln(n)+C+1/2\*n (n 很大时)

C≈0.57721566490153286060651209

# 数学结论

排列组合

$$A_n^m = n(n-1)\cdots(n-m+1) = \frac{n!}{(n-m)!}$$

$$C_n^m = \frac{A_n^m}{m!} = \frac{n!}{m!(n-m)!} = C_n^{n-m}$$

$$\frac{P_n^m}{\text{圆排列}} = \frac{n!}{(n-m)! \times m}, (1 \le m \le n)$$

有限多重集的排列 n! /(n1!\*n2!\*...\*nk!)

n 元无限集可重-r 组合 C(n+r-1,r)种。

n 元无限集取 r 个,n 中每个至少出现一次 C(r-1,n-1)(r≥n)

直线分平面

$$f(1) = 2$$

$$f(n) = f(n-1)+n = n(n+1)/2+1$$

折线分平面

$$f(1) = 2$$

$$f(n) = f(n-1)+4(n-1)+1 = 2n^2-n+1$$

圆分平面

$$f(1) = 2$$

$$f(n) = f(n-1)+2(n-1) = n^2-n+2$$

三角形分平面

$$f(1) = 2$$

$$f(n) = f(n-1)+6*(n-1)$$

平面分空间

$$f(1) = 2$$

$$g(n) = n(n+1)/2+1$$

$$f(n) = f(n-1)+g(n-1) = (n^3+5n)/6+1$$

求 Ax+By+C

已知 p1(x1,x1),p2(x2,y2),求 Ax+By+C=0

$$A = y2 - y1$$

$$B = x1 - x2$$

$$C = x2*y1-x1*y2$$

海伦公式

$$S = \sqrt{p(p-a)(p-b)(p-c)}$$
, p 为半周长

四边形最大面积

 $S = \sqrt{(p-a)(p-b)(p-c)(p-d)}$ , p 为半周长

## 斯特林公式

$$n! pprox \sqrt{2\pi n} \left(rac{n}{e}
ight)^n$$
 .

## 笛卡尔定理

4 个圆相切,外切 k = 1/r,内切 k = -1/r,圆退化成直线 k = 0。

$$(k_1+k_2+k_3+k_4)^2=2\,(k_1^2+k_2^2+k_3^2+k_4^2). 
onumber \ k_4=k_1+k_2+k_3\pm2\sqrt{k_1k_2+k_2k_3+k_3k_1}.$$

## Purfer 序列

一个长度为 n-2 的 Purfer 序列唯一对应一个 n 个点的树,且 Purfer 序列中 i 出现的次数就是节点 i 的度数减一。

## 圆锥最大面积

全面积为 па2 的圆锥最大面积为 sqrt(2)/12\*па^3。

## 正弦、余弦、正切

sinA / a = sinB / b = sinC/c

 $a^2 = b^2 + c^2 - 2bc \cdot cosA$ 

(a+b)/(a-b) = tan[(A+B)/2]/tan[(A-B)/2]

### 多边形面积

如果逆时针给出点坐标, 值为正,

$$\frac{1}{2} \left| \sum_{i=0}^{n-1} (x_i y_{i+1} - x_{i+1} y_i) \right|$$

如果顺时针给出点坐标, 值为负。

## Bash 游戏

有一堆石子共有 N 个。A B 两个人轮流拿, A 先拿。每次最少拿 1 颗, 最多拿 K 颗, 拿到最后 1 颗石子的人获胜。 若 n 不足 k+1,则 A 第一次取完,A 胜。

若 n E k+1 的倍数,每次 A 取 x, B 都能取 k+1-x, B 胜。否则, A 胜。

### 威佐夫游戏

有 2 堆石子。A B 两个人轮流拿, A 先拿。每次可以从一堆中取任意个或从 2 堆中取相同数量的石子, 但不可不取。 两堆石头 x < y,若(y-x)\*(sqrt(5)+1) / 2+1) == x,则 B 胜。否则 A 胜。

### 1/a 循环节长度

a 先约去 2 和 5 的因子->b,然后球欧拉函数值 φ,求一个最小的 x,使得 x|φ 且  $10^x$ % b=1.

# 计算几何

## 计算几何基本函数

```
#define eps 1e-8
#define PI acos(-1)
int sgn(double x)
    if(fabs(x) < eps) return 0;
    if(x < 0)
               return -1;
    return 1;
struct point
    double x, y;
    point() {};
    point (double a, double b):x(a), y(b) {};
                                                   //向量加法
    friend point operator+(point a, point b)
        return point (a. x+b. x, a. y+b. y);
    friend point operator—(point a, point b)
                                                   //向量减法
        return point (a. x-b. x, a. y-b. y);
    friend double operator*(point a, point b)
                                                   //点积
        return a. x*b. x+a. y*b. y;
    friend double operator (point a, point b)
                                                   //叉积
        return a. x*b. y-a. y*b. x;
                           //绕原点旋转弧度 B
    void trans (double B)
        double tx = x, ty = y;
        x = tx*cos(B)-ty*sin(B);
        y = tx*sin(B)+ty*cos(B);
struct line
    point s, e;
    line() {};
    line (point a, point b): s(a), e(b) {};
    friend pair (int, point) operator (line a, line b) //重合 0, 平行 1, 相交 2
        point ans = a.s;
        if(sgn((a. s-a. e)^(b. s-b. e)) == 0)
            if(sgn((a. s-b. e)^{\hat{}}(b. s-b. e)) == 0)
                                                   return make pair(0, ans);
                                                                                 //重合
            return make pair (1, ans);
        double t = ((a. s-b. s) (b. s-b. e))/((a. s-a. e) (b. s-b. e));
        ans. x += (a. e. x-a. s. x)*t;
        ans. y += (a. e. y-a. s. y)*t;
        return make_pair(2, ans);
```

#### 两点距离

```
double dis(point a, point b) //两点距离
         return sqrt((a-b)*(a-b));
                                                                                        点到直线距离
point dis2(point p, line 1)
                                                                          //点到直线距离,返回垂点
         point ans;
         double t = ((p-1.s)*(1.e-1.s))/((1.e-1.s)*(1.e-1.s));
         ans. x = 1. s. x+(1. e. x-1. s. x)*t;
         ans. y = 1. s. y+(1. e. y-1. s. y)*t;
         return ans;
                                                                                         点到线段距离
point dis3(point p, line 1)
                                                                         //点到线段距离,返回线段上最近的点
         point ans;
         double t = ((p-1.s)*(1.e-1.s))/((1.e-1.s)*(1.e-1.s));
         if(t) = 0 \&\& t <= 1
                  ans. x = 1. s. x+(1. e. x-1. s. x)*t;
                  ans. y = 1. s. y+(1. e. y-1. s. y)*t;
         else if (dis(p, 1.s) < dis(p, 1.e)) ans = 1.s;
         else
                           ans = 1.e;
         return ans;
                                                                                        判断线段相交
bool inter(line a, line b) //判断线段相交
         double x1 = a. s. x, y1 = a. s. y, x2 = a. e. x, y2 = a. e. y, x3 = b. s. x, y3 = b. s. y, x4 =
b. e. x, y4 = b. e. y;
         double t1 = (x2-x1)*(y3-y2)-(x3-x2)*(y2-y1);
         double t2 = (x2-x1)*(y4-y2)-(x4-x2)*(y2-y1);
         double t3 = (x4-x3)*(y1-y4)-(x1-x4)*(y4-y3);
         double t4 = (x4-x3)*(y2-y4)-(x2-x4)*(y4-y3);
         return t1*t2 <= 0 && t3*t4 <= 0;
                                                                                       直线和线段相交
bool inter2(line a, line b) //判断直线 a 和线段 b 相交
         return sgn((b. s-a. e)^{(a. s-a. e)})*sgn((b. e-a. e)^{(a. s-a. e)}) <= 0;
                                                                                       |判断点在直线上|
bool online (point p, line 1) //判断点在直线上
         return sgn((1. s-p)^{(1. e-p)}) == 0;
                                                                                       判断点在线段上
bool onseg(point p, line 1) //判断点在残段上
         return sgn((1. s-p)^{(1. e-p)}) == 0 \&\& sgn((p. x-1. s. x)*(p. x-1. e. x)) <= 0 \&\& sgn((p. y-1. s. x)) <= 0 && s
1. s. y)*(p. y-1. e. y)) \langle = 0;
                                                                                     判断点在凸多边形
bool onconvex(point a, point *p, int n) //判断点在凸多边形(凸包,点逆时针,若顺时针改
为>0)内,-1 在外,0 在边上,1 在内
         for (int i = 1; i \le n; i++)
```

```
int j = i+1;
          \begin{array}{ll} \textbf{if}(j > n) & \textbf{j} = 1; \\ \textbf{if}(\textbf{sgn}((\textbf{p[i]}-\textbf{a}) \hat{\ }(\textbf{p[j]}-\textbf{a})) < 0) & \textbf{return} = 1; \end{array}
          if (onseg(a, line(p[i], p[j])))     return 0;
     return 1;
                                                判断点在多边形内
                                                  //判断点在多边形内,-1 在外,0 在边上,1 在内
bool onpoly (point a, point *p, int n)
     int cnt = 0;
     line ray = line (a, point (-1e9, a, y));
     for (int i = 1; i \le n; i++)
          int j = i+1;
          if(j > n) j = 1;
          line side = line(p[i], p[j]);
          if (onseg(a, side)) return 0;
          if (sgn(side. s. y-side. e. y) == 0) continue;
          if (onseg (side. s, ray))
               if (sgn(side. s. y-side. e. y) > 0) cnt++;
          else if (onseg (side. e, ray))
               if (sgn(side. e. y-side. s. y) > 0) cnt++;
          else if (inter (ray, side))
                                            cnt++;
     if (cnt%2)
                    return 1:
     return -1;
                                                 计算多边形面积
                                          //计算多边形面积,可顺时针或逆时针
double calcarea (point *p, int n)
     double ans = 0;
     for (int i = 1; i \le n; i++)
          int j = i+1;
          if(j > n) j = 1;
ans += (p[i] p[j])/2;
     return fabs (ans);
                                                判断是否凸多边形
                                               //判断是否凸多边形,允许共线边,可顺时针或逆时针
bool isconvex(point *p, int n)
     boo1 s[3] = \{0\};
     for (int i = 1; i \le n; i++)
          int j = i+1, k = j+2;
          \begin{array}{cccc} \mathbf{if} (\mathbf{j} > \mathbf{n}) & \mathbf{j} = \mathbf{n}; \\ \mathbf{if} (\mathbf{k} > \mathbf{n}) & \mathbf{k} = \mathbf{n}; \end{array}
          s[sgn((p[j]-p[i]) ^(p[k]-p[i]))+1] = 1;
if(s[0] && s[2]) return 0;
     return 1;
                                                  判断四点共面
bool isface (int x1, int y1, int z1, int x2, int y2, int z2, int x3, int y3, int z3, int x4, int
y4, int z4)
```

int all = x1-x2, al2 = x2-x3, al3 = x3-x4;

```
int a21 = y1-y2, a22 = y2-y3, a23 = y3-y4;
                      int a31 = z1-z2, a32 = z2-z3, a33 = z3-z4;
                      return a11*a22*a33+a12*a23*a31+a13*a21*a32-a13*a22*a31-a11*a23*a32-a12*a21*a33 ==
 0:
                                                                                                                                                                                           判断三角形和圆相交
bool trianglecircle(int XX, int YY, int R, int X1, int Y1, int X2, int Y2, int X3, int Y3)
                      double a1 = (X1-XX)*(X1-XX)+(Y1-YY)*(Y1-YY);
                      double a2 = (X2-XX)*(X2-XX)+(Y2-YY)*(Y2-YY):
                      double a3 = (X3-XX)*(X3-XX)+(Y3-YY)*(Y3-YY);
                      double rr = R*R;
                      if (a1 < rr-1e-6 && a2 < rr-1e-6 && a3 < rr-1e-6)
                                                                                                                                                                                                                                                                                                   return 1;
                      if (a1 > rr+1e-6 && a2 > rr+1e-6 && a3 > rr+1e-6)
                                          double t = ((XX-X1)*(X2-X1)+(YY-Y1)*(Y2-Y1))*((XX-X2)*(X1-X2)+(YY-Y2)*(Y1-Y2));
                                          if(t > 1e-6)
                                                               t = abs((X1-XX)*(Y2-YY)-(X2-XX)*(Y1-YY))/sqrt((X1-X2)*(X1-X2)+(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2)*(Y1-Y2
 Y2));
                                                               if(t*t < rr+1e-6)
                                                                                                                                                             return 0;
                                           t = ((XX-X3)*(X2-X3)+(YY-Y3)*(Y2-Y3))*((XX-X2)*(X3-X2)+(YY-Y2)*(Y3-Y2));
                                          if(t > 1e-6)
                                                               t = abs((X3-XX)*(Y2-YY)-(X2-XX)*(Y3-YY))/sqrt((X3-X2)*(X3-X2)+(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2)*(Y3-Y2
 Y2));
                                                               if(t*t < rr+1e-6)
                                                                                                                                                            return 0;
                                          t = ((XX-X3)*(X1-X3)+(YY-Y3)*(Y1-Y3))*((XX-X1)*(X3-X1)+(YY-Y1)*(Y3-Y1));
                                          if(t > 1e-6)
                                                               t = abs((X3-XX)*(Y1-YY)-(X1-XX)*(Y3-YY))/sqrt((X3-X1)*(X3-X1)+(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1)*(Y3-Y1
 Y1));
                                                               if(t*t < rr+1e-6)
                                                                                                                                                                    return 0;
                                          return 1;
                      return 0;
                                                                                                                                                                                                                     求凸包
 struct node
                      int x, y;
                      friend bool operator < (node a, node b)
                                          if (a. x == b. x) return a. y < b. y;
                                          return a. x < b. x;
 }a[1005], ans[1005];
 int n, m;
 int cross (node a, node b, node c) //向量积
                      return (a. x-c. x)*(b. y-c. y)-(b. x-c. x)*(a. y-c. y);
 int convex(int n)//求凸包上的点
                      sort(a+1, a+n+1);
                      int m = 0;
                      for (int i = 1; i \le n; i++)
```

```
while (m > 1 \&\& cross(ans[m], a[i], ans[m-1]) \le 0) m--;
        ans[++m] = a[i];
    int k = m;
    //求得上凸包
    for (int i = n-1; i >= 1; i--)
        while (m > k \&\& cross (ans[m], a[i], ans[m-1]) \le 0)
        ans[++m] = a[i];
              m--;//起始点重复。
    if(n > 2)
    return m;
                                        求三角形外心
point waixin (point a, point b, point c)
                                         //求三角形外心
    double a1 = b. x-a. x, b1 = b. y-a. y, c1 = (a1*a1+b1*b1)/2;
    double a2 = c. x-a. x, b2 = c. y-a. y, c2 = (a2*a2+b2*b2)/2;
    double d = a1*b2-a2*b1;
    return point (a. x+(c1*b2-c2*b1)/d, a. y+(a1*c2-a2*c1)/d);
                                      计算两圆相交面积
double circlrarea (point al, int rl, point a2, int r2)
                                                          //计算两圆相交面积
    double d = dis(a1, a2);
    if (r1+r2 < d+eps) return 0;
    if(d < fabs(r1-r2)+eps)
        double r = min(r1, r2);
        return PI*r*r;
    double x = (d*d+r1*r1-r2*r2)/(2*d);
    double t1 = a\cos(x/r1), t2 = a\cos((d-x)/r2);
    return r1*r1*t1+r2*r2*t2-d*r1*sin(t1);
                                        平面最近点对
int n, b[100005];
struct xx
    double x, y;
    friend bool operator (xx a, xx b)
        return a. x < b. x;
a[100005];
bool cmp(int x, int y)
    return a[x].y < a[y].y;
double dis(xx a, xx b)
    return sqrt ((a. x-b. x)*(a. x-b. x)+(a. y-b. y)*(a. y-b. y));
double f(int 1, int r)
    if(1 == r) return 0;
    if(r-1 == 1)
                    return dis(a[1], a[r]);
    if(r-1 == 2)
        double minn = 1e18;
```

```
minn = min(minn, dis(a[1], a[1+1]));
        minn = min(minn, dis(a[1], a[r]));
        minn = min(minn, dis(a[1+1], a[r]));
        return minn;
    int mid = (1+r)/2;
    double minn = min(f(1, mid), f(mid+1, r));
    int cnt = 0;
    for (int i = 1; i \le r; i++)
        if(a[i]. x \ge a[mid]. x-minn && a[i]. x \le a[mid]. x+minn) b[++cnt] = i;
    sort(b+1, b+1+cnt, cmp);
    for (int i = 1; i \le cnt; i++)
        for (int j = i+1; j \leftarrow cnt; j++)
             if(a[b[j]]. y-a[b[i]]. y > minn) break;
             minn = min(minn, dis(a[b[i]], a[b[j]]));
    return minn;
int main()
    while (scanf ("%d", &n) && n)
                                       scanf ("%1f%1f", &a[i].x, &a[i].y);
        for (int i = 1; i \le n; i++)
        sort(a+1, a+1+n);
printf("%. 2f\n", f(1, n)/2);
    return 0;
                                            半平面交
int sgn(double x)
    if(fabs(x) < eps)
                        return 0;
    if(x < 0)
               return -1;
    return 1;
struct point
    double x, y;
    point() {};
    point (double a, double b):x(a), y(b) {};
                                                     //向量加法
    friend point operator+(point a, point b)
        return point (a. x+b. x, a. y+b. y);
    friend point operator—(point a, point b)
                                                     //向量减法
        return point (a. x-b. x, a. y-b. y);
    friend double operator*(point a, point b)
                                                     //点积
        return a. x*b. x+a. y*b. y;
                                                     //叉积
    friend double operator (point a, point b)
        return a. x*b. y-a. y*b. x;
```

```
}ans[20005];
struct line
    point s, e;
    double k:
    line() {};
    line (point a, point b):s(a), e(b)
         k = atan2(e. y-s. y, e. x-s. x);
    friend point operator&(line a, line b)
         point ans = a.s;
         double t = ((a. s-b. s) \hat{b. s-b. e})/((a. s-a. e) \hat{b. s-b. e});
         ans. x += (a. e. x-a. s. x)*t;
         ans. y += (a. e. y-a. s. y)*t;
         return ans;
}a[20005], q[20005];
bool HPIcmp(line a, line b)
                                 //直线左边
    \begin{array}{ll} \textbf{if} (fabs (a. k - b. k) > eps) & \textbf{return} \ a. k < b. k; \\ \textbf{return} \ ((a. s - b. s) \hat{\ } (b. e - b. s)) < 0; \end{array}
//返回核的凸包点
void HPI(line *a, int n, point *ans, int &cnt)
    sort(a+1, a+n+1, HPIcmp);
    int tot = 0;
    for (int i = 1; i \le n; i++)
                                                a[++tot] = a[i]:
         if(fabs(a[i].k-a[i-1].k) > eps)
    int head = 0, tail = 1;
    q[0] = a[1];
    q[1] = a[2];
    cnt = 0;
    for (int i = 3; i \le tot; i++)
         if(fabs((q[tail].e-q[tail].s)^(q[tail-1].e-q[tail-1].s)) < eps | |
fabs((q[head]. e-q[head]. s)^(q[head+1]. e-q[head+1]. s)) < eps) return;
         while (head < tail && (((q[tail]&q[tail-1])-a[i].s)^(a[i].e-a[i].s)) > eps)
tail--;
         while (head \langle tail \&\& (((q[head]\&q[head+1])-a[i].s)^(a[i].e-a[i].s)) \rangle eps)
head++;
         q[++tai1] = a[i];
    while (head < tail && (((q[tail]&q[tail-1])-q[head].s)^(q[head].e-q[head].s)) > eps)
    while (head < tail && (((g[head]&g[head-1])-g[tail].s)^(g[tail].e-g[tail].e)) > eps)
head++;
    if(tail <= head+1) return;</pre>
    for (int i = head; i < tail; i++) ans [++cnt] = q[i] & q[i+1];
                            ans[++cnt] = q[head]&q[tai1];
    if (head < tail-1)
```

### 旋转卡壳求最远点对

```
point a[50005], ans[50005];
int n, m;
int cross(point a, point b, point c)//向量积
    return (a. x-c. x)*(b. y-c. y)-(b. x-c. x)*(a. y-c. y);
int dis2(point a, point b)
    return (a-b)*(a-b);
int convex(int n)//求凸包上的点
    sort(a+1, a+n+1);
    int m = 0;
    for (int i = 1; i \le n; i++)
        while (m > 1 \&\& cross(ans[m], a[i], ans[m-1]) \le 0) m--;
        ans[++m] = a[i];
    int k = m;
    //求得上凸包
    for (int i = n-1; i >= 1; i--)
        while (m > k \&\& cross(ans[m], a[i], ans[m-1]) \le 0)
        ans[++m] = a[i];
    if(n > 2)
                 m--;//起始点重复。
    return m;
int rotatingcalipers(point *a, int n)
    int ans = 0, now = 2, ne = 3;
    if (ne > n) ne = 1;
    for (int i = 1; i \le n; i++)
        int j = i+1;
        if (j > n) j = 1;
point t = a[i]-a[j];
        while ((\hat{t}(a[ne]-a[now])) < 0)
             now = ne;
             ne = now+1;
             if (ne > n) ne = 1;
        ans = \max(ans, \max(dis2(a[i], a[now]), dis2(a[j], a[ne])));
    return ans;
int main()
    while ("scanf ("%d", &n))
        for (int i = 1; i \le n; i++)
                                       scanf ("%d%d", &a[i].x, &a[i].y);
        m = convex(n);
        printf("%d\n", rotatingcalipers(ans, m));
```

#### 旋转卡壳球平面点最大三角形面积

```
//已求得凸包
int rotatingcalipers(point *a, int n)
    int ans = 0;
    for (int i = 1; i \le n; i++)
         int j = i+1;
         if(j > n) \quad j = 1;
         int k = j+1;
         if(k > n) k = 1;
         int kk = k+1;
         if(kk > n) \quad kk = 1;
         while (j != i && k != i)
              ans = \max(ans, abs((a[j]-a[i]^(a[k]-a[i]))));
              while (((a \lfloor i \rfloor - a \lfloor j \rfloor)) (a \lfloor kk \rfloor - a \lfloor k \rfloor)) < 0)
                   k = kk;
                   kk++;
                   if(kk > n) kk = 1;
    return ans;
                                        旋转卡壳求两凸包最小距离
```

```
double dis4(point a, point b, point c)
    return dis(a, dis3(a, line(b, c)));
double dis5 (point a, point b, point c, point d)
    double ans1 = min(dis4(a, c, d), dis4(b, c, d));
    double ans2 = min(dis4(c, a, b), dis4(d, a, b));
    return min(ans1, ans2);
double getangel (point a, point b, point c, point d)
    return (b-a)^(d-c);
int rotatingcalipers(point *a, int n, point *b, int m)
    int sa = 1, sb = 1;
    for (int i = 1; i \le n; i++)
         if(sgn(a[i]. y-a[sa]. y) < 0) sa = i;
    for (int i = 1; i \le m; i++)
         if(sgn(b[i].y-b[sb].y) < 0) sb = i;
    double t, ans = dis(a[sa], b[sb]);
    int na = sa+1, nb = sb+1;
    if(na > n)  na = 1;

if(nb > m)  nb = 1;
    for (int i = 1; i \le n; i++)
         while (sgn(t = getangel(a[sa], a[na], b[sb], b[nb])) < 0)
```

```
sb = nb;
            nb++;
            if(nb > m) \quad nb = 1;
        if(sgn(t) == 0) ans = min(ans, dis5(a[sa], a[na], b[sb], b[nb]));
               ans = min(ans, dis4(b[sb], a[sa], a[na]));
        sa = na;
        na++:
        if (na > n) na = 1;
    return ans;
//已求得凸包
double solve (point *a, int n, point *b, int m)
    return min(rotatingcalipers(a, n, b, m), rotatingcalipers(b, m, a, n));
                                           高精度
//输入的两个数要求为正。
int compare (string strl, string str2)
    if (str1. length() > str2. length())
                                          return 1;
    else if (str1. length() < str2. length())
            return strl. compare (str2);
    else
string add(string strl, string str2)//加法
    string str;
    int len1 = strl.length(), len2 = str2.length();
    if(1en1 < 1en2)
        for (int i = 1; i \le 1en2-1en1; i++)
                                               str1 = "0" + str1;
    else
        for (int i = 1; i \le 1en1-1en2; i++) str2 = "0" + str2;
    int cf = 0, temp;
    for (int i = str1. length()-1; i \ge 0; i--)
        temp = str1[i]-'0'+str2[i]-'0'+cf;
        cf = temp/10;
        temp %= 10;
        str = char(temp+'0')+str;
                   str = char(cf+'0')+str;
    if(cf != 0)
    return str;
string sub(string strl, string str2)//减法
    string str;
    int flag = 0;
    if (compare (str1, str2) < 0)
        flag = 1;
        swap(str1, str2);
    int tmp = str1.length()-str2.length(), cf = 0;
    for (int i = str2. length()-1; i \ge 0; i--)
```

```
if(str1[tmp+i] < str2[i]+cf)
            str = char(str1[tmp+i]-str2[i]-cf+'0'+10)+str;
        else
            str = char(str1[tmp+i]-str2[i]-cf+'0')+str;
            cf = 0;
    for (int i = tmp-1; i >= 0; i--)
        if(str1[i]-cf >= '0')
            str = char(str1[i]-cf)+str;
            cf = 0;
        else
            str = char(str1[i]-cf+10)+str;
            cf = 1;
    str.erase(0, str.find_first_not_of('0'));
    if (str. empty()) str = "0";
if (flag) str = "-"+str;
    return str;
string mul(string strl, string str2)//乘法
    string str;
    int len1 = str1.length();
    int 1en2 = str2. 1ength();
    string tempstr;
    for (int i = 1en2-1; i >= 0; i--)
        tempstr = "";
        int temp = str2[i]-'0', t = 0, cf = 0;
        if(temp != 0)
                                                   tempstr += "0";
            for (int j = 1; j \le 1en2-1-i; j++)
            for (int j = 1en1-1; j >= 0; j--)
                 cf = (temp*(str1[j]-'0')+cf);
                 t = cf%10;
                 cf /= 10;
                 tempstr = char(t+'0')+tempstr;
            if(cf != 0) tempstr = char(cf+'0')+tempstr;
            str=add(str, tempstr);
    str.erase(0, str.find_first_not_of('0'));
    if (str. empty()) str = "0";
    return str;
void div(string strl, string str2, string &quotient, string &residue)//除法取余
    quotient = "";
    residue = "":
```

```
if(str2 == "0")
    quotient = "ERROR";
    residue = "ERROR";
    return:
if(str1 == "0")
    quotient = "0";
    residue = "0";
    return:
int res = compare(str1, str2);
if (res < 0)
    quotient = "0";
    residue = str1;
    return;
else
    int len1 = str1.length();
    int 1en2 = str2. 1ength();
    string tempstr;
    tempstr.append(str1, 0, 1en2-1);
    for (int i = 1en2-1; i < 1en1; i++)
        tempstr = tempstr+str1[i];
        tempstr.erase(0, tempstr.find_first_not_of('0'));
        if (tempstr. empty())   tempstr = "0";
        for (char ch = '9'; ch >= '0'; ch--)
            string str, tmp;
            str = str+ch;
            tmp = mul(str2, str);
             if (compare(tmp, tempstr) <= 0)</pre>
                 quotient = quotient+ch;
                 tempstr = sub(tempstr, tmp);
                 break;
    residue = tempstr;
quotient.erase(0, quotient.find_first_not_of('0'));
if (quotient.empty()) quotient="0";
```

# 其他

## 

归并排序

```
void merge(int a[], int first, int mid, int last, int temp[])//合并
    int i = first, j = mid+1, k = 0;
    while (i \le mid \&\& j \le last)
        if(a[i] < a[j])  temp[k++] = a[i++];
                temp[k++] = a[j++];
        else
                         temp[k++] = a[i++];
    while (i \le mid)
    while (j \le last)
                          temp[k++] = a[j++];
    for (i = 0; i < k; i++)
                             a[first+i] = temp[i];
void _merge_sort(int arr[], int temp[], int first, int last)//归并
    if(first < last)</pre>
        int mid = (first+last)/2;
        merge sort (arr, temp, first, mid);
        _merge_sort(arr,temp,mid+1,last);
        merge (arr, first, mid, last, temp);
void merge_sort(int arr[], int n)//归并排序接口
    int* temp = (int*)malloc(n*sizeof(int));
    _merge_sort(arr, temp, 0, n-1);
    free (temp);
                                           堆排序
void percdown(int arr[], int p, int n)//调整最大堆
    int Parent, Child, x = arr[p];
    for(Parent = p;Parent*2+1 < n;Parent = Child)</pre>
        Child = Parent*2+1;
        if (Child != n-1 && arr[Child] < arr[Child+1])</pre>
                                                            Child++;
        if(x \ge arr[Child]) break;
                arr[Parent] = arr[Child];
    arr[Parent] = x;
void heap_sort(int arr[], int n)//堆排序接口
    for (int i = n/2-1; i >= 0; i--)
                                      percdown(arr, i, n);
    for (int i = n-1; i > 0; i--)
        swap(arr[0], arr[i]);
        percdown (arr, 0, i);
                                          快速排序
void quick sort(int arr[], int 1, int r)
    if(1 < r)
        int i = 1, j = r, x = arr[1];
        while (i < j)
            while (i < j \&\& arr[j] >= x)
                                             j--;
```

```
if(i < j) arr[i++] = arr[j];
             while (i < j \&\& arr[i] < x)
                          arr[i--] = arr[i];
             if(i < j)
         arr[i] = x;
         \_quick\_sort(arr, 1, i-1);
         _quick_sort(arr, i+1, r);
                                              头文件
#pragma comment (linker, "/STACK:102400000, 102400000")
#include<cstdio>
#include iostream
#include < iomanip >
#include <algorithm>
#include<cmath>
#include<cstring>
#include<string>
#include<vector>
#include<stack>
#include<queue>
#include(set)
#include < map >
#include<list>
#include < deque >
#include <bitset>
#include<cstdlib>
#define LL long long
#define PI acos(-1)
#define eps 1e-8
#define lowbit(x) (x&-x)
#define MOD 1000000007
#define INF 0x3f3f3f3f
#define MEM(a, x) memset(a, x, sizeof(a))
using namespace std;
                                             输入外挂
void read(int &ans)
    ans = 0;
    int flag = 0, ch;
    if((ch = getchar()) == '-') flag = 1;
else if(ch >= '0' && ch <= '9') ans = ch-'0';
while((ch = getchar()) >= '0' && ch <= '9')
                                                           ans = ans*10+ch-'0';
    if (flag)
                ans *= -1;
                                              随机数
#include<random>
std::random device rd;
rd();
std::mt19937 mt(rd());
mt();
```